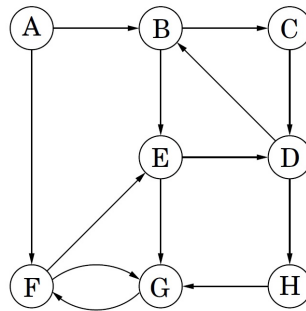


### 3. WERKCOLLEGE 3

You are allowed to answer in Dutch.

3.1. Perform depth-first search on the following graph; whenever there's a choice of vertices, pick the one that is alphabetically first. Classify each edge as a tree edge, forward edge, back edge, or cross edge.



3.2. Suppose a CS curriculum consists of  $n$  courses, all of them mandatory. The prerequisite graph  $G$  has a node for each course, and an edge from course  $v$  to course  $w$  if and only if  $v$  is a prerequisite for  $w$ . Find an algorithm that works directly with this graph representation, and computes the minimum number of semesters necessary to complete the curriculum (assume that a student can take any number of courses in one semester). The running time of your algorithm should be linear.

3.3. You are given a rooted binary tree  $T = (V, E)$ , with a designated root node  $r \in V$ . We say that  $u$  is an ancestor of  $v$  in the rooted tree if the unique path from  $r$  to  $v$  in  $T$  passes through  $u$ . You wish to preprocess the tree so that queries of the form is  $u$  an ancestor of  $v$ ? can be answered in constant time. The preprocessing itself should take linear time. How can this be done?

3.4. Most graph algorithms that take an adjacency-matrix representation as input require time  $\mathcal{O}(|V|^2)$ , but there are some exceptions. Show how to determine whether a directed graph  $G = (V, E)$  contains a *universal sink*, i.e., a vertex with in-degree  $n-1$  and out-degree 0, in time  $\mathcal{O}(|V|)$  given an adjacency matrix for  $G$ .

3.5. Give an efficient algorithm which takes as input a directed graph  $G = (V, E)$ , and determines whether or not there is a vertex  $s \in V$  from which all other vertices are reachable.

3.6. Another way to perform topological sorting on a directed acyclic graph  $G = (V, E)$  is to repeatedly find a vertex of in-degree 0, output it, and remove it and all of its outgoing edges from the graph. Explain how to implement this idea so that it runs in time  $\mathcal{O}(|V| + |E|)$ .