

(10 20)

Delete

(10)

40  
30  
20  
10

Delete

()

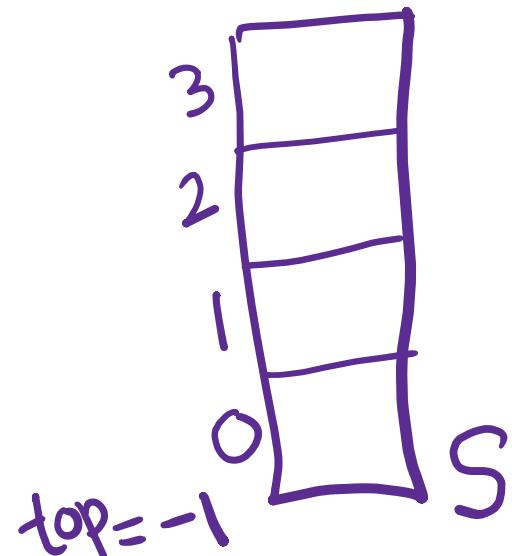
Last In First Out (LIFO)

Insert a new element:— push  
Delete :— POP

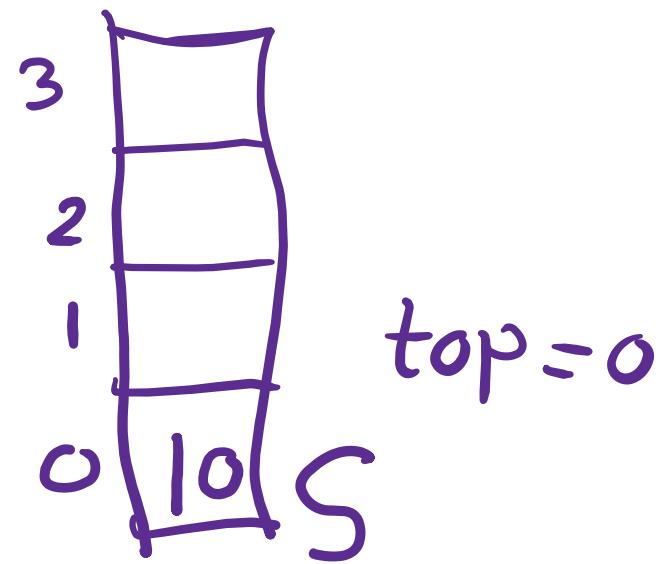
# Implementation of stack using Array:

- Elements will be Kept in an array
- We need one variable to Keep track of the 'Top' of the stack

✓ Push 10



$\text{top} = \text{top} + 1$   
 $\text{top} = 0$   
 $S[\text{top}] = 10$

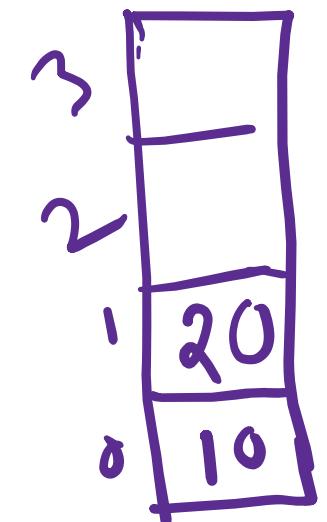


push 20

top = top + 1

top = 1

S[top] = 20



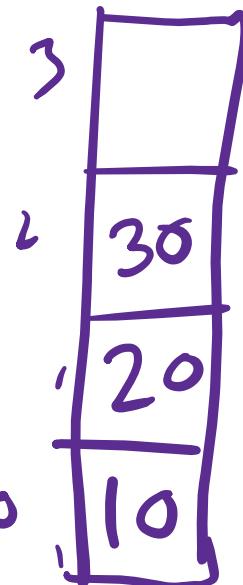
top = 1

push 30

top = top + 1

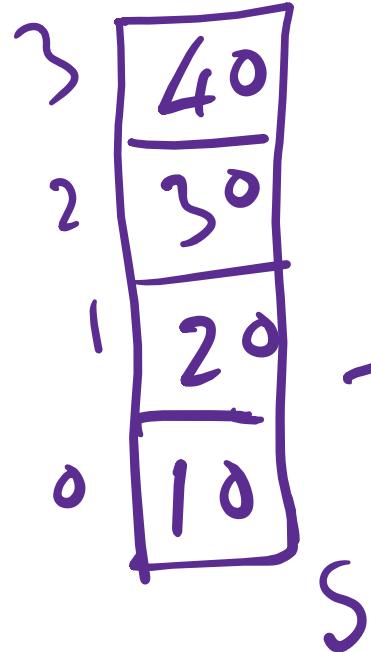
top = 2

S[top] = 30



top = 2

push 40



top = 3

push 50  
stack is full

#define MAXSIZE 4  
int S[MAXSIZE]

# Linear list

Stack

10

20

30

40

( )

(10)

(10 20)

(10 20 30)

(10 20 30 40)  
↑  
top

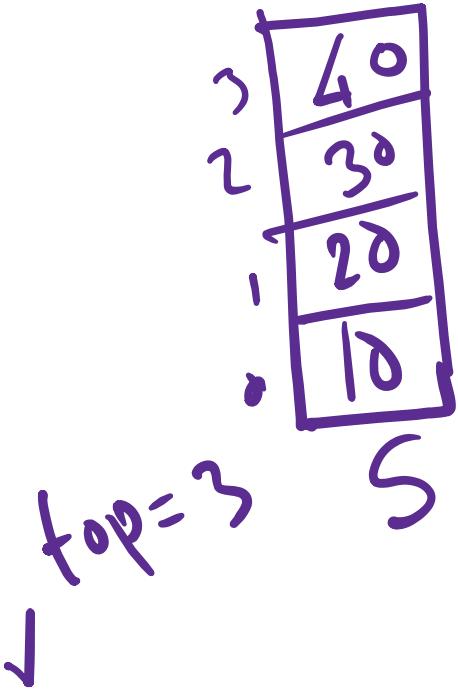
Delete

Element deleted is 40

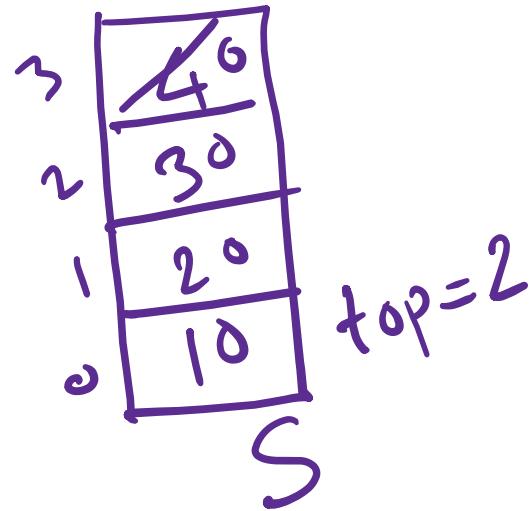
(10 20 30)

↑  
top  
Delete  
(10 20)

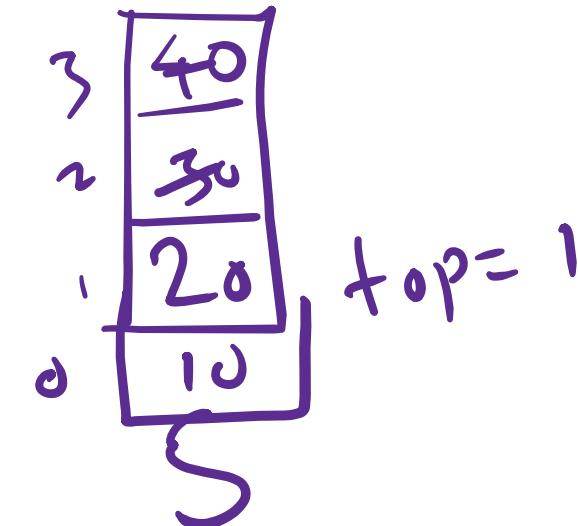
# Stack Overflow



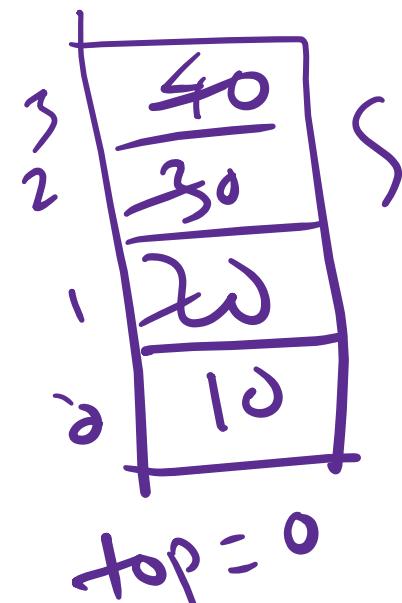
Pop

$$\text{top} = \text{top} - 1$$
$$\text{top} = 2$$


Pop

$$\text{top} = \text{top} - 1$$
$$\text{top} = 1$$


Pop

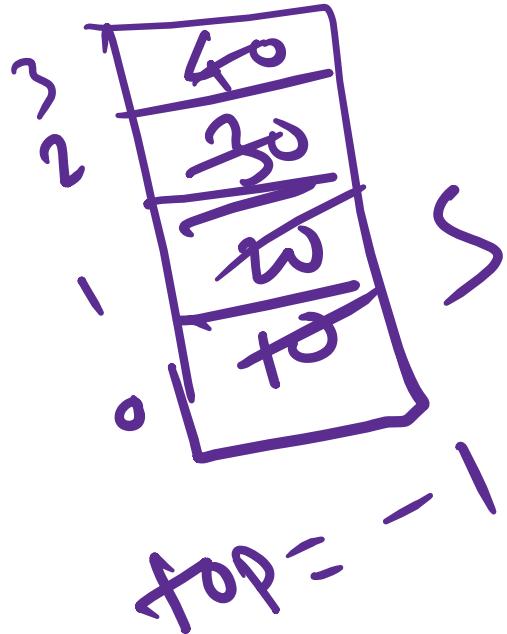
$$\text{top} = \text{top} - 1$$
$$\text{top} = 0$$


Stack from: 0 to top

Pop

$$\text{top} = \text{top} - 1$$

$$\text{top} = -1$$



Stack is empty

Pop

Stack underflow

If ( $\text{top} == -1$ )

Then 'Underflow'

Procedure push(S, MAXSIZE, top, element)

1. If ( $\text{top} == \text{MAXSIZE} - 1$ ) Then ↵  
2.     Write('Stack Overflow') ↵  
3.     Return ↵  
4. EndIf  
5. .  $\text{top} = \text{top} + 1$  ↵  
6. .  $S[\text{top}] = \text{element}$  ↵  
7. Return ↵

constant time  
 $O(1)$

procedure pop(s,top) O(1)

1. If ( $\text{top} == -1$ ) Then
2.     Write('Stack underflow')
3.     Return
4. EndIf
5.     Write('Deleted Element:', s[top])
6.      $\text{top} = \text{top} - 1$
7.     Return

procedure display(S, top) { }

1. If ( $\text{top} == -1$ ) Then  $N$   
    Write('Stack Underflow')  
    Return  $N \text{ elements}$
- 2.
- 3.
4. EndIf  $O \text{ to } N-1$
5. Repeat step 6 for  $I=0, 1, 2, \dots, \text{top}$   $O(N)$   
    Write('The Element: ', S[I])  $N/2$
- 6.
7. Return

5.  $I = top$

6. Repeat Step 7 and 8 while  $I \geq 0$

7.  $Write(S[I])$

8.  $I = I - 1$

9.

s : array

push( )

function: elements of s

pop( )

Where to declare s?

display( )

int s[20];

- globally? — outside all functions

- inside main: local to main

- pass as arguments

```
#include <cs50.h>
#include <stdio.h>

int main()
{
    int s[20];
    push(s);
}

void push()
```

The C Programming  
Language

— Kernighan &  
Ritchie

```
#define <stdio.h>
#define maxsize 20
int s[maxsize], top;
int main(){
    top = -1;
    push();
    push();
    push();
    pop();
    push();
    display();
    return;
```

```
void push(){
    int element;
    printf("Element: ");
    scanf("%d", &element);
    if (top == maxsize - 1)
    {
        printf("Stack Overflow");
        return;
    }
    s[++top] = element;
}
return;
```

```
void pop() {  
    if (top == -1)  
    {  
        printf("Stack underflow");  
        return;  
    }  
    printf("Popped Element is: %d\n", s[top-1]);  
    return;  
}
```

```
Void display( ) {  
    int i;  
    for (i=0; i<=top; i++)  
        printf("Element is: %d\n", s[i]);  
    return;  
}  
  
for (i=top; i>=0; i--)  
    printf.....
```

`push(s&top, element);`

```
int main() {
    int s[maxsize];
    int top, element;
    .
    .
    push(s, &top, element);
```

array is  
implicitly  
passed by  
reference

Queue: A queue is a linear list where all the insertions are done at one end of the list (known as rear) and all the deletions are done at the other end of the list (known as front). A queue exhibits FIFO property.

Insert : 10 20 30

( )

(10)

(10 20)

(10 20 30)

Delete: Deleted element is 10

( 20 30 )

Delete:  $(\underline{20} \quad 30)$

Deleted element is 20

$(30)$

FIFO

Delete: Deleted element is 30

$()$

Insertion Order: 10 20 30

Deletion Order: 10 -20 -30

## Implementation of Queue using Array

We use an array to keep the elements of the queue

We need two more variables to keep track of the 'rear' end and 'front' end.

We call these: REAR and FRONT



0	1	2	3
10	20	30	40

Rear = 3, Front = 0

1) Insert 10 :

$$\begin{aligned} \text{Rear} &= \text{Rear} + 1 \\ Q[\text{Rear}] &= 10 \end{aligned}$$

$$\text{Front} = \text{Front} + 1$$

2) Insert 20 :

$$\begin{aligned} \text{Rear} &= \text{Rear} + 1 \\ Q[\text{Rear}] &= 20; \end{aligned}$$

Insert 30:

$$\begin{aligned} 3) \text{ Rear} &= \text{Rear} + 1 \\ Q[\text{Rear}] &= 30 \end{aligned}$$

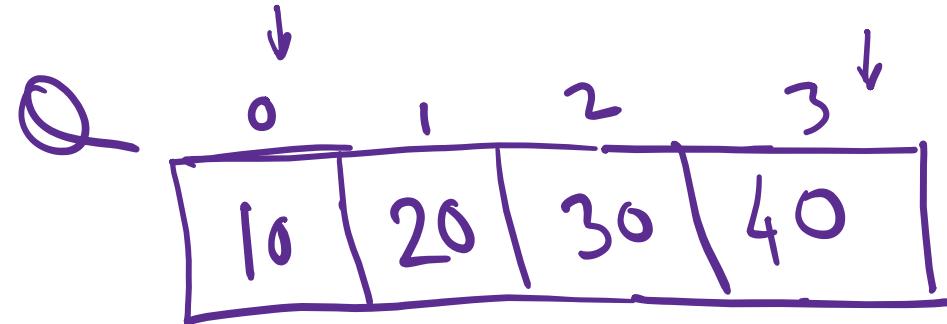
4) Insert 40 :

$$\begin{aligned} \text{Rear} &= \text{Rear} + 1 \\ Q[\text{Rear}] &= 40 \end{aligned}$$

5) Insert 50

Queue Overflow or

Queue Full Cond<sup>n</sup>.



rear=3, front=0

If (rear == maxsize - 1)

# define maxsize 4

int  $Q[\text{maxsize}]$ ;

Deletions: i) Delete

Front = Front + 1

0	1	2	3
10	20	30	40

Rear = 3, Front = 3

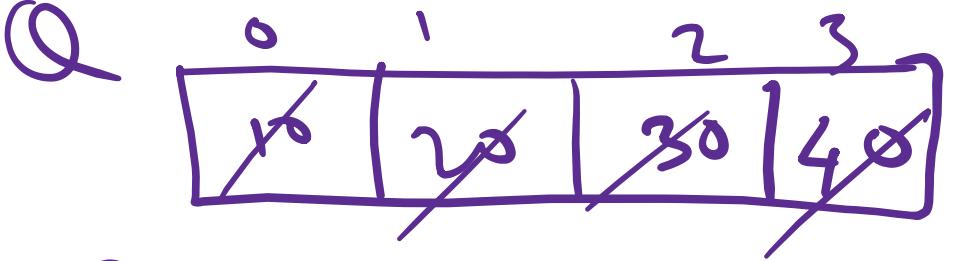
Queue is always  
from index 'Front'  
to index 'Rear'

2) Delete : Front = Front + 1

20 is deleted

3) Delete : Front = Front + 1

30 is deleted



Rear=3, Front=3

40 is deleted

Rear= -1      Front= -1

Deletion of  
only element in  
queue

5) Delete<sup>x</sup> : Queue Underflow or  
Queue Empty Condition

# Circular Queue :

maxsize=4

initially,  
Rear=-1  
Front=-1

Q	0	1	2	3
	50	60	30	40

Rear=1, Front=2

Delete: Q[Front] or q[0]

Delete q[1]

Insert 50-

Insert 60

Circularly Updated modulus (%) operator

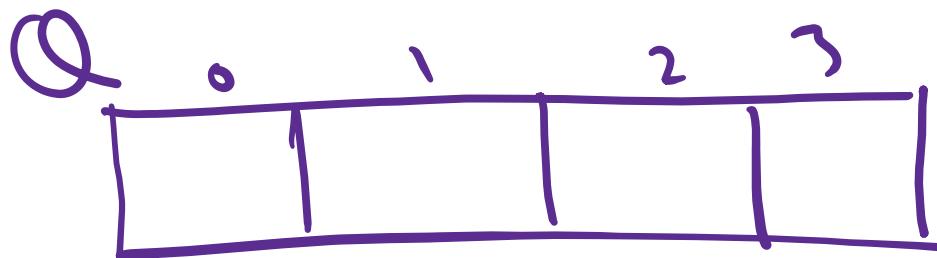
$$\begin{aligned}\text{Rear} &= (\text{Rear} + 1) \bmod \text{maxsize} \\ &= (3 + 1) \% 4 \\ &= 4 \% 4 \\ &= 0\end{aligned}$$

$$\begin{aligned}\text{Rear} &= (\text{Rear} + 1) \% \text{maxsize} \\ &= (0 + 1) \% 4 \\ &= 1 \% 4 \\ &= 1\end{aligned}$$

# Circular Queue Insertion and Deletion

Insert 10, 20, 30, 40

1)

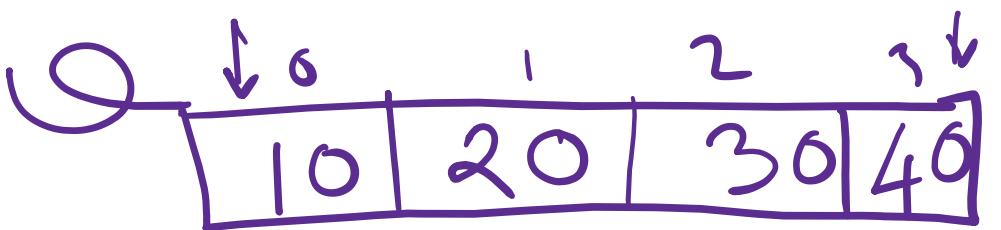


Rear = -1, Front = -1

2)

Insert 10

$$\begin{aligned}\text{Rear} &= (\text{Rear} + 1) \% 4 \\ &= (-1 + 1) \% 4 \\ &= 0\end{aligned}$$



Rear = 3, Front = 0

3) Insert 20

$$\text{Rear} = (0+1)\%4 \\ = 1$$

4) Insert 30

$$\text{Rear} = (1+1)\%4 \\ = 2$$

5) Insert 40

$$\text{Rear} = (2+1)\%4 \\ = 3$$

If (Front == 0)

or (Rear == maxsize)

6) Delete :  $\text{Q} \begin{array}{|c|c|c|c|} \hline 0 & 1 & 2 & 3 \\ \hline 10 & 20 & 30 & 40 \\ \hline \end{array}$

$\text{Rear} = 3, \text{Front} = 1$

$$\begin{aligned}\text{Front} &= (\text{Front} + 1) \% \text{maxSize} \\ &= (0 + 1) \% 4 \\ &= 1 \% 4 \\ &= 1\end{aligned}$$

7) Delete :  $\text{Front} = (1 + 1) \% 4$   
 $= 2$

$\text{Q} \begin{array}{|c|c|c|c|} \hline 0 & 1 & 2 & 3 \\ \hline \cancel{10} & \cancel{20} & 30 & 40 \\ \hline \end{array}$   
 $\text{Rear} = 3, \text{Front} = 2$

7) Insert + 50

Q	0	1	2	3
	50	20	30	40

Rear = 0, Front = 2

$$\begin{aligned}\text{Rear} &= (3+1)\gamma \cdot 4 \\ &= 0\end{aligned}$$

$$Q[\text{Rear}] = 50$$

8)

Insert 60

Q

0	1	2	3
50	60	30	40

✓

Rear = -1, Front = 1

$$\begin{aligned} \text{Rear} &= (0 + 1) \% 4 \\ &= 1 \end{aligned}$$

Q [Rear] = 60

Delete

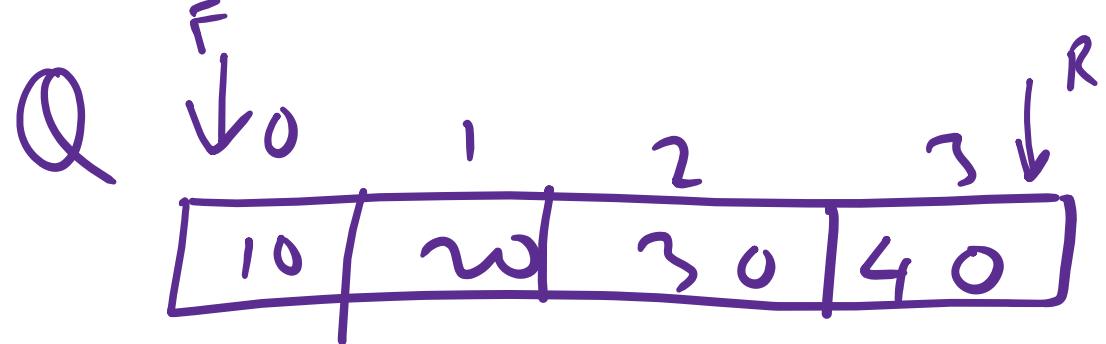
Delete

Delete

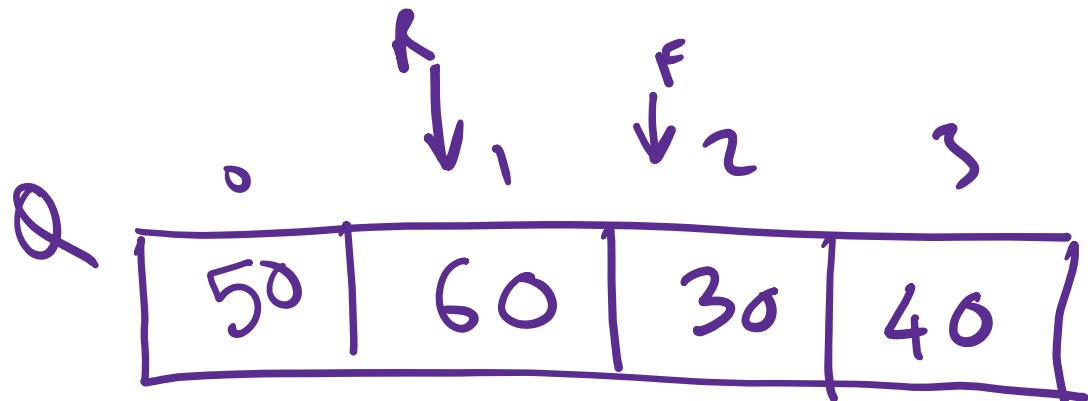
Delete

Front = -1

Rear = -1



Rear = 3, Front = 0



Rear = 1, Front = 2

Circular Queue  
 $(4 \times 4)$   
 $((\text{Rear} + 1) \% \text{maxsize})$   
 == Front  
 $0 == 0$  True

Normal Queue:      Insert: enqueue  
                              Delete: dequeue

procedure Enqueue(Q, Rear, Front, maxsize)

1. If (Real == maxsize - 1) Then
2.     Write ('Queue Overflow')
3.     Return
4. EndIf
5. Rear = Real + 1

6.  $Q[Rear] = \text{element}$
7. If ( $\text{Front} == -1$ ) Then
8.      $\text{Front} = 0$
9. End If
10. Return

Procedure Dequeue(Q, Rear, Front)

1. If ( $\text{Front} == -1$ ) Then
2.     Write('Queue Underflow')
3.     Return
4. Endif
5. Write('Element to be deleted:', Q[Front])
6. If ( $\text{Front} == \text{Rear}$ ) Then
7.      $\text{Front} = \text{Rear} = -1$

8. Else  
9.     Front = Front + 1  
10. EndIf  
11. Return

C Implementation: - Global

- One function for Enqueue
- One " " " Dequeue
- One " " " Display

array  
Real  
Front

- Declare the array, real, front locally within main
- Pass the array, real, front as arguments.

& real  
& front