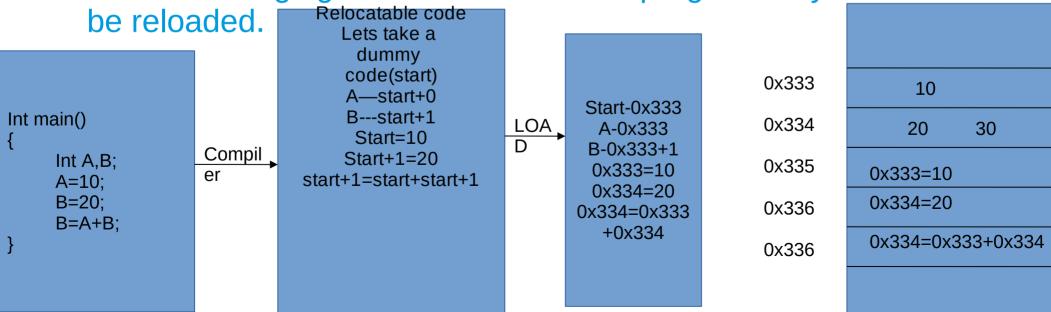
Binding at load time

Here the compiler produces only relocatable code

Binding to absolute address is done during load time.

In case of changing in fence address the program only needs to



Any Disadvantages

The fence address can not change when program is executing.

Binding at execution time dynamic binding

The fence register is now called base register and every address generated by a user process at the time it is sent to memory.

The user never sees the real physical addresses.

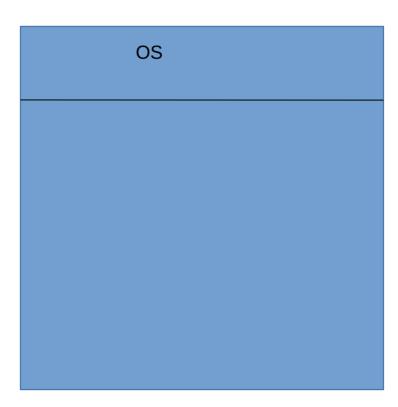
The user program deals with logical addresses. The memory management unit converts logical addresses into physical addresses.

The user generates only logical addresses.

Allocation of memory partition

Ready queue

P1 P2 p3



Fixed partition (MFT)

Also known as multiprogramming with a fixed number of tasks.

Divide the memory into a number of fixed size partitions.

Each partition may contain exactly one process.

There fore the degree of multi programming is bound by the number of partitions.

When a partition is free a process is selected from the ready queue and is loaded into the free partition.

When a process terminates the partition becomes available for another process.

P2

P5

p6

Degree of multiprogramming is 3

OS
p1
p3
p4

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P2 P5

p6

Degree of multiprogramming is 3

OS	
p1	
p3	
p4	

P2 P5

Degree of multiprogramming is 3

OS
p1
p3
p6

32 kB memory can be divided into regions of the following sizes:

Os----10KB

Very large programs----12 KB

Average programs-----6 KB

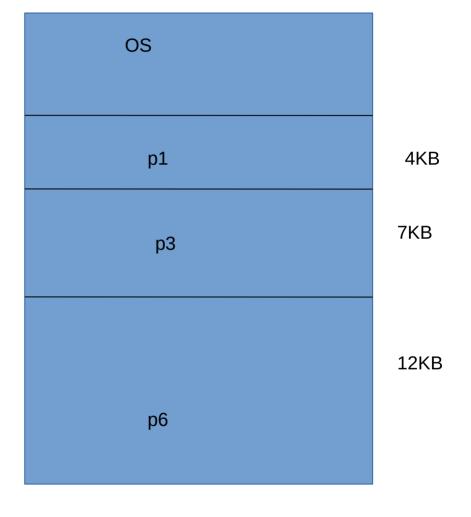
Small programs-----4 KB

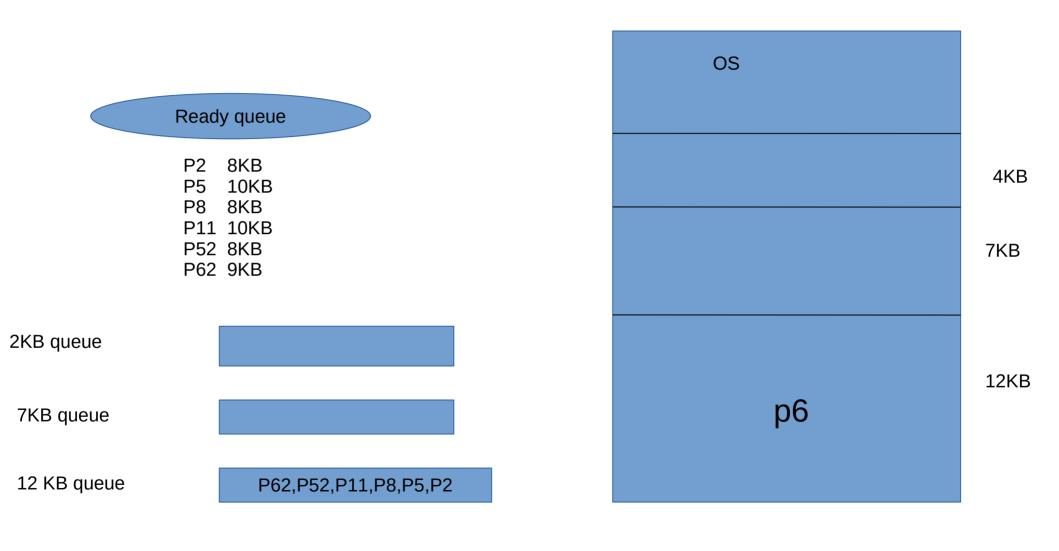
One variation is have multiple queues, a queue for each region.

The user specifies the maximum amount of memory required or the OS can attempt to determine the memory requirements automatically.

Accordingly a process could be assigned to queue.

Re	eady queue
P P P	22 2KB 25 10KB 28 8KB 211 5KB 252 3KB 262 6KB
2KB queue	P52,P2
7KB queue	P62,P11
12 KB queue	P8,P5





Disadvantage

Internal and external fragmentation may occur.

Dynamic Allocation Variable partition Allocation (MVT)

The OS keeps a table indicating which part of memory are available (called holes) are available and which are occupied.

Initially all memory is available for user processes.

When a process arrives we select a hole which is large enough to hold this process.

We allocate as much memory is required for the process and the rest is kept as a hole which can be used for later requests.

If a hole large enough for this process cannot be found, this process waits until some other process(es) finishes and a large enough hole is available.

If a new hole is adjacent to other holes OS merge the adjacent holes to form a large hole.

P2 8KB

P5 10KB

P8 8KB

P11 10KB

P52 8KB

P62 9KB

Free---24KB

OS	
p2	
	32KB

P5 10KB P8 8KB P11 10KB P52 8KB P62 9KB

Free---24KB

OS	
p2	
P5	
	32KB

P8 8KB P11 10KB P52 8KB P62 9KB

Free---14KB

OS	
p2	
P5	
	32KB

P8 8KB P11 10KB P52 8KB P62 9KB

Free---14KB

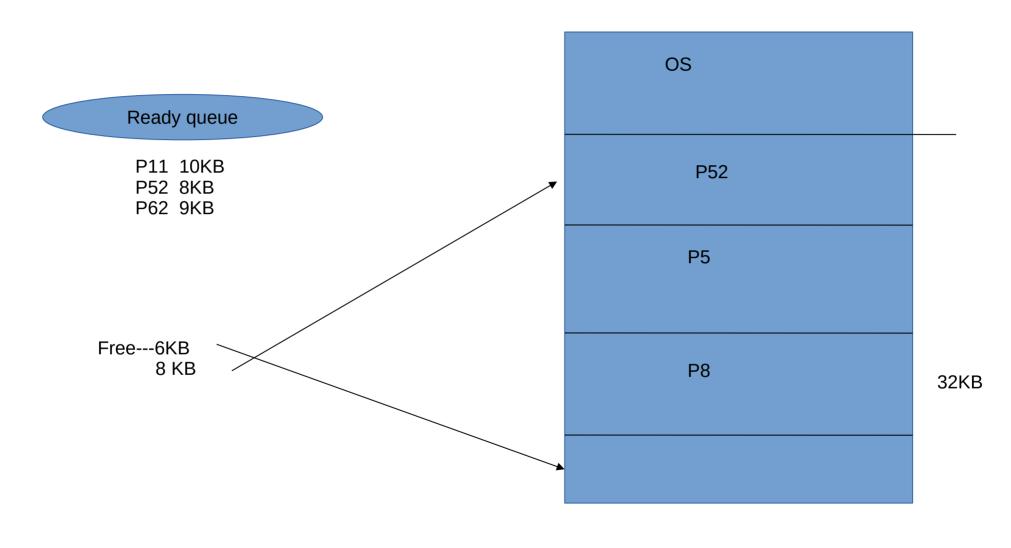
OS	
p2	
P5	
P8	32KB

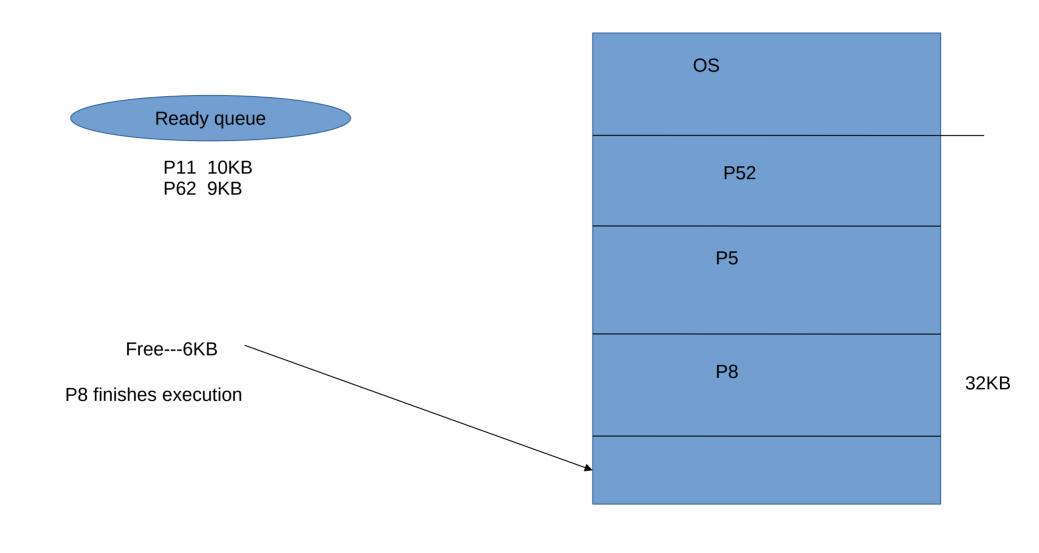
P11 10KB P52 8KB P62 9KB

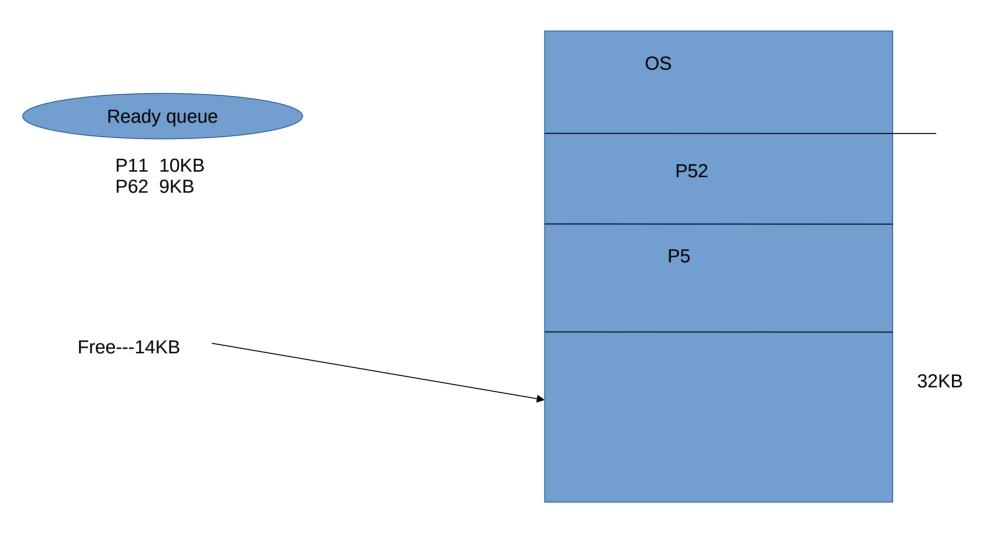
Free---6KB

P2 finishes execution

OS	
p2	
P5	
P8	32KB







	OS	
Ready queue		
P62 9KB	P52	
	P5	
Free4KB	P11	32KB