

## CHAPTER - 1

# *Introduction & Basics*

Neso Academy

## Introduction to OPERATING SYSTEM



Windows



Linux



Ubuntu

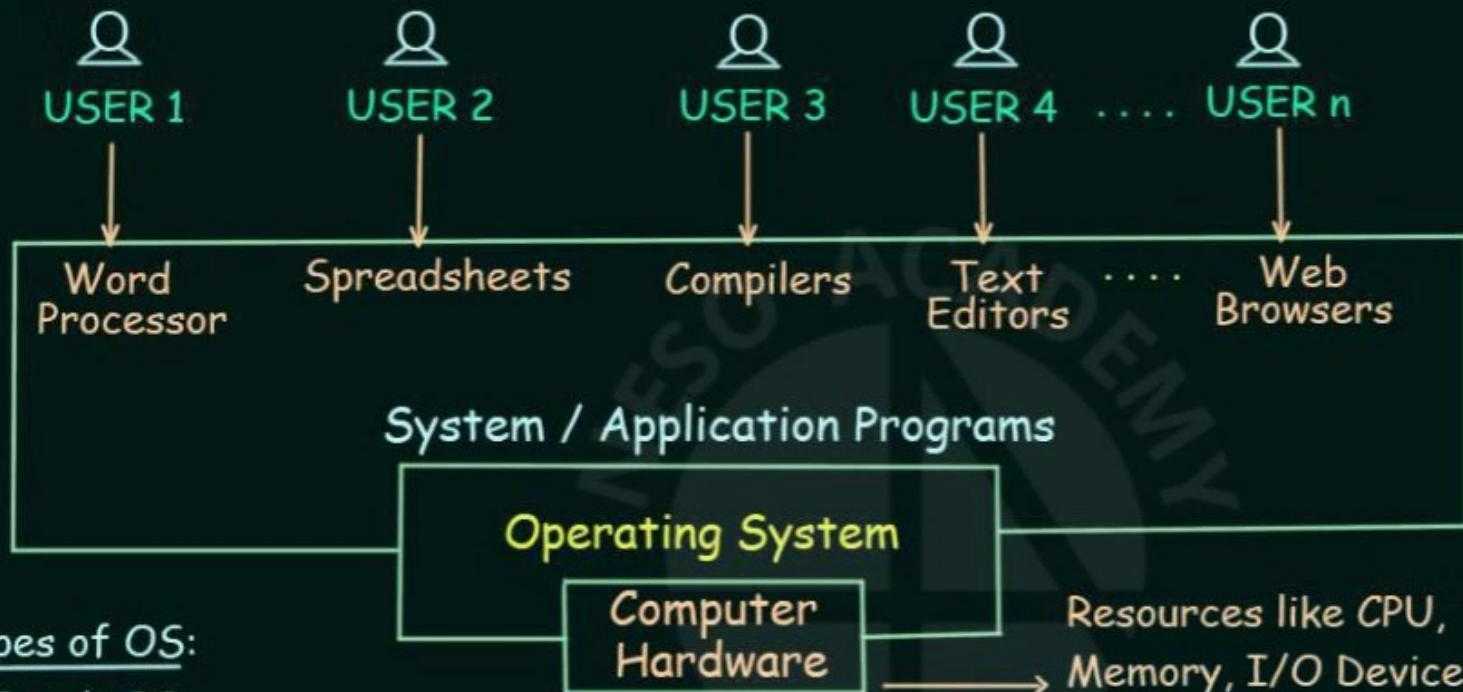


Mac OS X  
iOS



Android

- An Operating System (OS) is a program that manages the computer hardware.
- It also provides a basis for Application Programs and acts as an intermediary between computer User and computer Hardware.



#### Types of OS:

- Batch OS
- Time sharing OS
- Distributed OS
- Network OS
- Real Time OS
- Multi Programming/  
Processing/ Tasking OS

#### Goals of OS:

- i) Convenience
- ii) Efficiency
- iii) Both

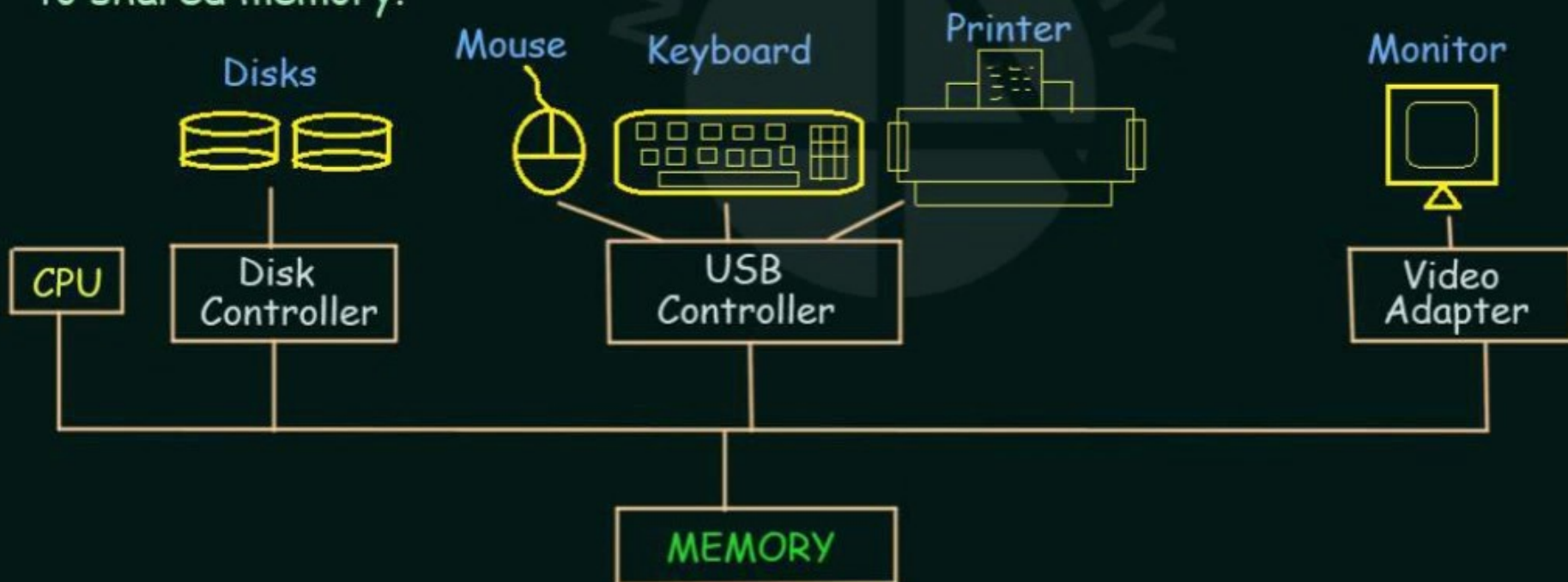
#### Functions of OS:

- It is an interface between User & Hardware
- Allocation of Resources
- Management of Memory, Security, etc.

## Basics of Operating System (Computer System Operation)

Some basic knowledge of the structure of Computer System is required to understand how Operating Systems work.

→ A modern general-purpose computer system consists of one or more CPUs and a number of device controllers connected through a common bus that provides access to shared memory.





- Each device controller is in charge of a specific type of device
- The CPU and the device controllers can execute concurrently, competing for memory cycles
- To ensure orderly access to the shared memory, a memory controller is provided whose function is to synchronize access to the memory

#### Some important terms:

- 1) **Bootstrap Program:** → The initial program that runs when a computer is powered up or rebooted.
  - It is stored in the ROM.
  - It must know how to load the OS and start executing that system.
  - It must locate and load into memory the OS Kernel.
- 2) **Interrupt:** → The occurrence of an event is usually signalled by an Interrupt from Hardware or Software.
  - Hardware may trigger an interrupt at any time by sending a signal to the CPU, usually by the way of the system bus.

3) System Call (Monitor call): → Software may trigger an interrupt by executing a special operation called System Call.

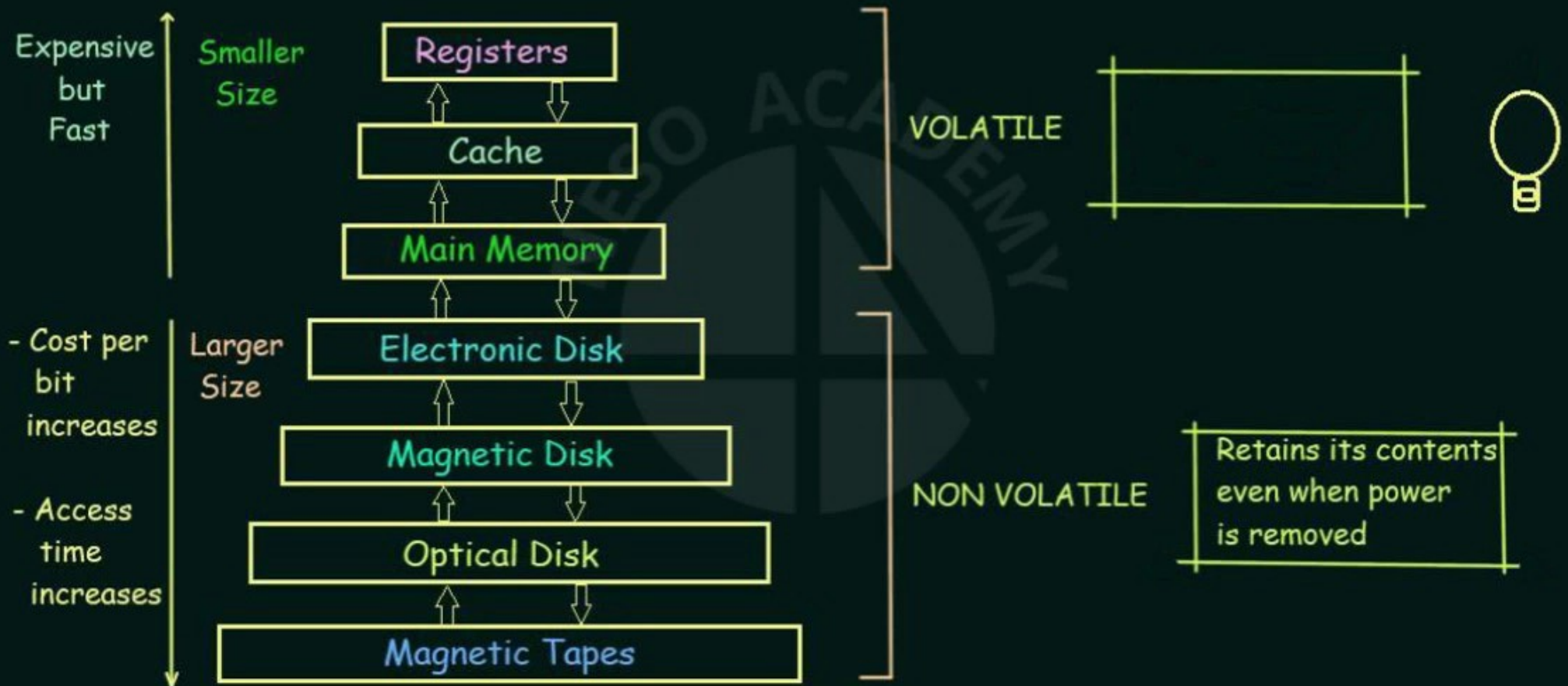
When the CPU is interrupted, it stops what it is doing and immediately transfers execution to a fixed location.

→ The fixed location usually contains the starting address where the Service Routine of the interrupt is located.

The Interrupt Service Routine executes.

On completion, the CPU resumes the interrupted computation.

## Basics of Operating System (Storage Structure)





## Basics of Operating System ( I/O Structure)

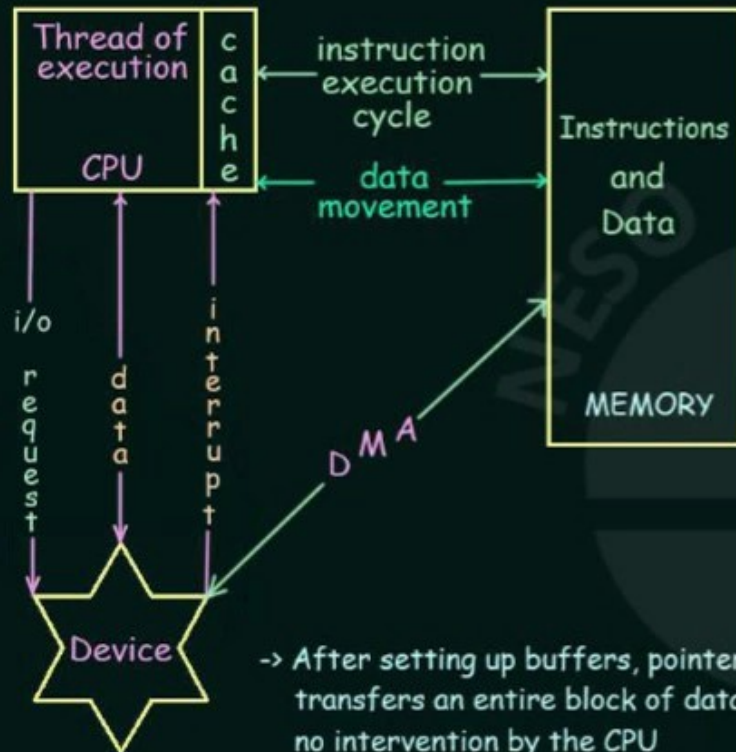
- > Storage is only one of many types of I/O devices within a computer
- > A large portion of operating system code is dedicated to managing I/O, both because of its importance to the reliability and performance of a system and because of the varying nature of the devices
- > A general-purpose computer system consists of CPUs and multiple device controllers that are connected through a common bus
- > Each device controller is in charge of a specific type of device



- > Typically, operating systems have a device driver for each device controller
- > This device driver understands the device controller and presents a uniform interface to the device to the rest of the operating system



### Working of an I/O Operation:



- > To start an I/O operation, the device driver loads the appropriate registers within the device controller
- > The device controller, in turn, examines the contents of these registers to determine what action to take
- > The controller starts the transfer of data from the device to its local buffer
- > Once the transfer of data is complete, the device controller informs the device driver via an interrupt that it has finished its operation
- > The device driver then returns control to the operating system

This form of interrupt-driven I/O is fine for moving small amounts of data but can produce high overhead when used for bulk data movement

To solve this problem, Direct Memory Access (DMA) is used

- > After setting up buffers, pointers, and counters for the I/O device, the device controller transfers an entire block of data directly to or from its own buffer storage to memory, with no intervention by the CPU
- > Only one interrupt is generated per block, to tell the device driver that the operation has completed
- > While the device controller is performing these operations, the CPU is available to accomplish other works

## Computer System Architecture

Types of Computer Systems based on number of General Purpose Processors:

1. Single Processor Systems



2. Multiprocessor Systems



3. Clustered Systems



## 1. Single Processor Systems



- Has two or more processors in close communication, sharing the computer bus and sometimes the clock, memory, and peripheral devices

### Advantages:



Increased throughput



Economy of scale

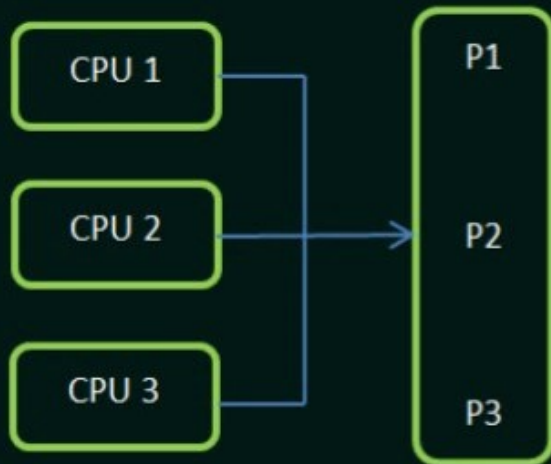


Increased reliability

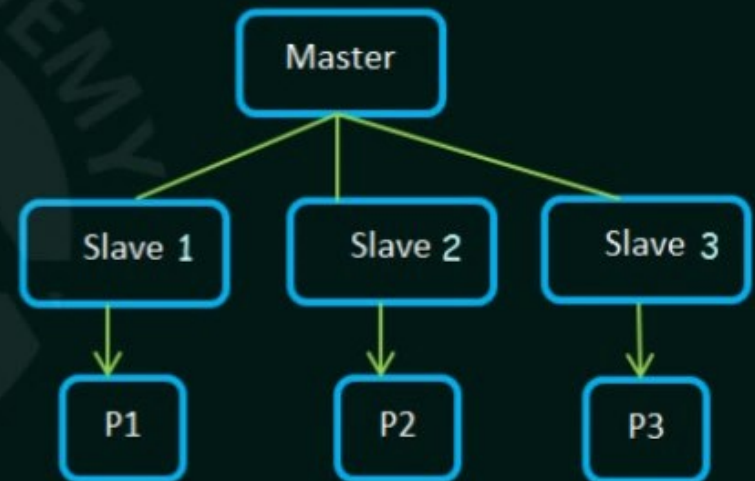


## Types of Multiprocessor Systems:

### Symmetric Multiprocessing



### Asymmetric Multiprocessing



### 3. Clustered Systems



- Like multiprocessor systems, clustered systems gather together multiple CPUs to accomplish computational work.
- They are composed of two or more individual systems coupled together.
- Provides high availability
- Can be structured **asymmetrically** or **symmetrically**



- One machine in Hot-Standby mode
- Others run applications



- Two or more hosts run applications
- Monitors each other