1. Create a BST and delete an element

InsertNode(root, key):

1. If root is NULL:

Create a new node with key and return it

2. If key < root->key:

root->left = InsertNode(root->left, key)

3. Else:

root->right = InsertNode(root->right, key)

4. Return root

DeleteNode(root, key):

1. If root is NULL:

Return root

2. If key < root->key:

root->left = DeleteNode(root->left, key)

3. Else if key > root->key:

root->right = DeleteNode(root->right, key)

4. Else:

If root has one or no child:

Return the non-NULL child or NULL

If root has two children:

Find inorder successor (smallest in right subtree)

Replace root->key with successor's key

Delete successor

5. Return root

2. Breadth-First Search (BFS)

BFS(graph, start):

- 1. Create an empty queue and enqueue start node
- 2. Mark start as visited
- 3. While queue is not empty:

Dequeue a node, say current

For each neighbor of current:

If neighbor is not visited:

Mark as visited and enqueue it

3. Depth-First Search (DFS)

DFS(graph, node, visited):

- 1. Mark node as visited
- 2. For each neighbor of node:

If neighbor is not visited:

Call DFS(graph, neighbor, visited)

4. Matrix Chain Multiplication

```
function MatrixChainMultiplication(p[]):

n = length \ of \ p - 1

let m be a 2D array of size n x n

for i = 1 to n:

m[i][i] = 0 // No cost to multiply one matrix

for length = 2 to n:

for i = 1 to n - length + 1:

j = i + length - 1

m[i][j] = infinity

for k = i to j - 1:

q = m[i][k] + m[k+1][j] + p[i-1] * p[k] * p[j]

if q < m[i][j]:

m[i][j] = q

return m[1][n]
```

5. AVL Tree Insertion and Deletion

InsertAVL(root, key):

- 1. Perform standard BST insertion
- 2. Update height of current node
- 3. Check balance factor:

If unbalanced, perform appropriate rotations (LL, LR, RL, RR)

4. Return root

DeleteAVL(root, key):

- 1. Perform standard BST deletion
- 2. Update height of current node
- 3. Check balance factor:

If unbalanced, perform appropriate rotations (LL, LR, RL, RR)

4. Return root

6. Create and Delete from a Red-Black Tree

InsertRBT(root, key):

- 1. Perform standard BST insertion
- 2. Fix violations using rotations and recoloring

DeleteRBT(root, key):

- 1. Perform standard BST deletion
- 2. Fix violations using rotations and recoloring

7. Heap Sort in Ascending and Descending Order

HeapSort(A):

- 1. Build a max-heap from the array
- 2. For i = n-1 to 1:

Swap A[0] and A[i] Heapify(A, 0, i)

ReverseHeapSort(A):

- 1. Build a min-heap from the array
- 2. For i = n-1 to 1:

Swap A[0] and A[i] Heapify(A, 0, i)

8. Prim's Algorithm

Prim(graph):

- 1. Initialize all nodes as unvisited
- 2. Start from any node and add it to the MST
- 3. While MST doesn't include all nodes:

Select the smallest edge connecting visited and unvisited nodes Add the selected edge to the MST Mark the new node as visited

9. Red-Black Tree Creation and Deletion

// algorithm for deletion in Red-Black Tree

```
rb_delete(t, z)
  if z->left = NULL or z->right = NULL
      y \leftarrow z
  else y \leftarrow tree-successor(z)
  if y->left ≠ NULL
      x \leftarrow y->left
  else x \leftarrow y->right
  x->p ← y->p
  if y - p = NULL
      t->root \leftarrow x
  else if y = y-p-left
      y->p->left \leftarrow x
  else y->p->right \leftarrow x
  if y 3≠ z
      z->key ← y->key
  //copy y's satellite data into z
  if y->color = BLACK
      rb_delete_fixup(t, x)
  return y
 rb_delete_fixup(t, x)
  while x \neq t->root and x->color = BLACK
     do if x = x-p-
         w \leftarrow x - p - right
     if w->color = RED
         w->color ← BLACK //Case 1
     x->p->color ← RED //Case 1
     left-rotate(t, x->p) //Case 1
     w ← x->p->right //Case 1
     if w->left->color = BLACK and w->right->color = BLACK
         w->color ← RED //Case 2
     x \leftarrow x - p //Case 2
     else if w->right->color = BLACK
         w->left->color ← BLACK
                                               //Case 3
     w->color ← RED
                              //Case 3
     right-rotate(t, w)
                            //Case 3
```

```
w ← x->p->right //Case 3
w->color ← x->p->color //Case 4
x->p->color ← BLACK //Case 4
w->right->color ← BLACK //Case 4
left-rotate(t, x->p) //Case 4
x ← t->root /*Case 4*/
else (same as then clause with "right" and "left" exchanged)
x->color ← BLACK
```

10. Strassen's Matrix Multiplication

Strassen(A, B):

1. If size of A or B is 1x1:

Return A * B

- 2. Divide A and B into 4 submatrices each
- 3. Compute 7 matrix products using Strassen's formula:

```
P1 = A11 * (B12 - B22)
P2 = (A11 + A12) * B22
```

. . .

4. Combine results to form the resultant matrix

11. Create a Binary Tree from Preorder and Inorder Traversals

BuildTree(preorder, inorder):

1. If inorder is empty:

Return NULL

- 2. Root = preorder[0]
- 3. Find root in inorder, say index
- 4. Recursively build left subtree from inorder[0:index] and preorder
- 5. Recursively build right subtree from inorder[index+1:end] and preorder
- 6. Return root

12. Merge Sort

MergeSort(A, left, right):

1. If left < right:

mid = (left + right) // 2

MergeSort(A, left, mid)

MergeSort(A, mid+1, right)

Merge(A, left, mid, right)

Merge(A, left, mid, right):

- 1. Create temporary arrays for left and right parts
- 2. Merge elements back into A in sorted order

13. BST Inorder Traversal

InorderTraversal(root):

1. If root is not NULL:

Call InorderTraversal(root->left)

Print root->key

Call InorderTraversal(root->right)

14. Dijkstra's Algorithm

Dijkstra(graph, source):

- 1. Initialize distances of all nodes to ∞, distance[source] = 0
- 2. Create a priority queue and add source
- 3. While queue is not empty:

Extract node with minimum distance, say u

For each neighbor v of u:

If distance[u] + weight(u, v) < distance[v]:

Update distance[v]

Add v to the queue