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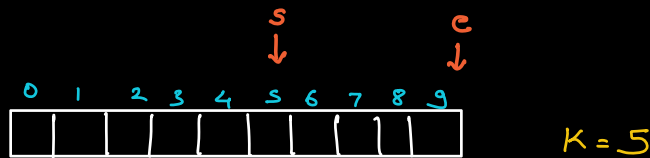
Q Given an array & a no. K .

Find sum of all sub-arrays of length K .

A: $-3, 4, -2, 5, 3, -2, 8, 2, -1, 4$

$K = 5$

| s | e | sum |
|---|---|-----|
| 0 | 4 | 7 |
| 1 | 5 | 8 |
| 2 | 6 | 12 |
| 3 | 7 | 16 |
| 4 | 8 | 10 |
| 5 | 9 | 11 |



$s = 0;$

$[a, b] \rightarrow b - a + 1$

$e = K - 1; \quad // \quad e: [K-1, N-1]$

while ($e < N$) { \rightarrow # iterations : $N - K + 1$

// iterate from s to e & calculate the sum

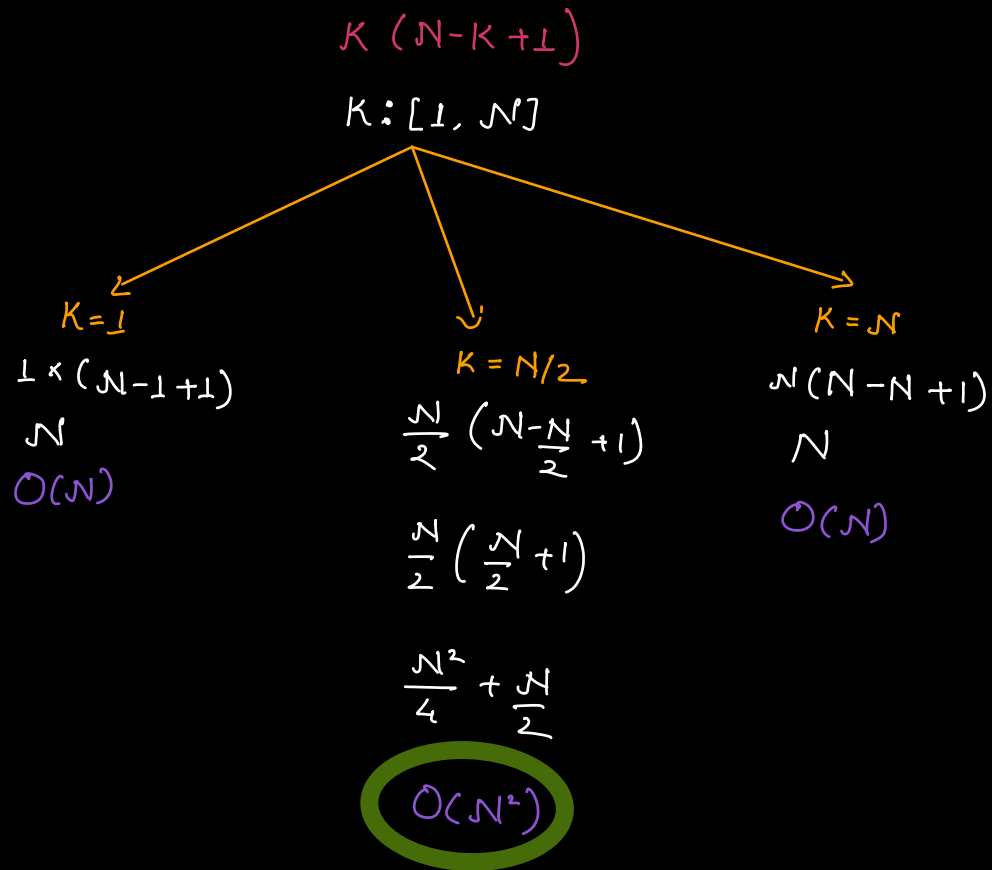
Sum = 0;

Range sum
s to e
 $O(1)$ using
PSI

{ for ($i = s; i \leq e; i++$) {
Sum = Sum + A[i],
} } # iterations : K

Print (Sum);
s++;
e++;
}

Total iterations = $K(N - K + 1)$



Using PS

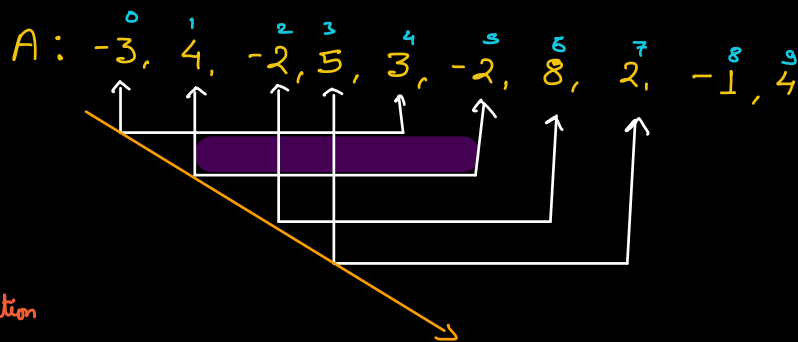
iterations = N
 Build PS

+ $N-K+1$ $\rightarrow 2N-K+1$
 iterate over subarrays
 of size K

Tc : $O(N)$

Sc : $O(N)$

Sliding Window



| s | e | sum | |
|---|---|-----|--|
| 0 | 4 | 7 | // iteration |
| 1 | 5 | 8 | $\Sigma[1, 5] = \Sigma[0, 4] - A[0] + A[5] \rightarrow 8$ |
| 2 | 6 | 12 | $\Sigma[2, 6] = \Sigma[1, 5] - A[1] + A[6] \rightarrow 12$ |
| 3 | 7 | 16 | $\Sigma[3, 7] = \Sigma[2, 6] - A[2] + A[7] \rightarrow 16$ |
| 4 | 8 | 10 | $\Sigma[4, 8] = \Sigma[3, 7] - A[3] + A[8] \rightarrow 10$ |
| 5 | 9 | 11 | $\Sigma[5, 9] = \Sigma[4, 8] - A[4] + A[9] \rightarrow 11$ |

$$\Sigma[i, j] = \Sigma[i-1, j-1] - A[i-1] + A[j]$$

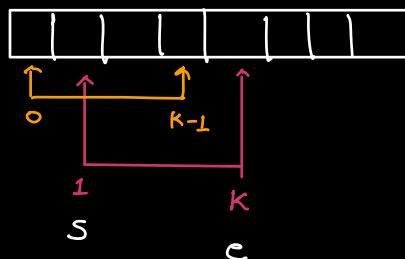
// first window of length K

```

sum = 0;
for (i = 0; i < K; i++) {
    sum = sum + A[i];
}
Print(sum);

```

K iteration



```

s = 1;
e = K; // e: [K, N-1]

```

```

while (e < N) {

```

```

    sum = sum - A[s-1] + A[e];

```

```

    Print(sum);

```

```

    e++;

```

```

    s++;

```

```

}

```

(N-K) iterations

Total iterations = K + N - K \rightarrow N

TC: $O(N)$

SC: $O(1)$

Q Given an array & a no. k .
Find max sum of all sub-arrays of length k .

Q Given an array & a no. k .
Find min sum of all sub-arrays of length k .

Q Given an array & a no. k .
Find average of every window of size k . (moving average)

Q Given an array of size N & a number B .
Return the minimum no. of swaps required to bring all elements less than or equals to B together.

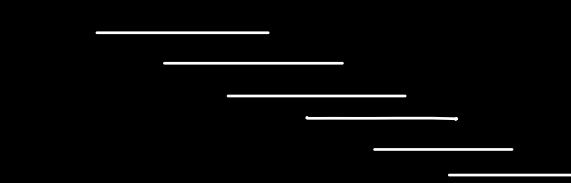
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A : 1, 12, 10, 3, 14, 10, 5 B : 8
14 12, 10, 3, 1, 5, 10 \rightarrow 2

A : 5, 17, 100, 11 B : 20
5, 17, 11, 100

A : 7, 12, 19, 11, 6, 5, 3, 14 B : 8

A : 1, 12, 10, 3, 14, 10, 5, 12, 9, 14 B = 8

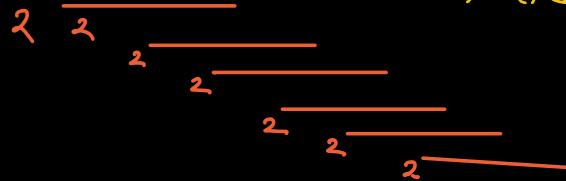


At the end all the numbers $\leq B$ will end up in a window of size same as count of these numbers (assume k)

In how min swaps can we bring all k number in a given window? \Rightarrow No of elements greater than B in that window.



A : 1, 12, 10, 3, 14, 10, 5, 12, 9, 14 B = 8



1, 2, 10, 3, 14, 10

B = 8

Count = 1
min = 1

4, 6, 5, 9, 18, 10, 11, 7, 1 $B = 8$

13, 14, 15, 16, 17, 7, 6, 9 .. $B = 8$

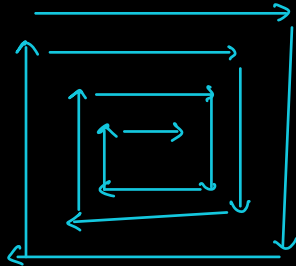
Q.

Given a matrix (N x N). Print it in spiral form.

Amazon
ms
DP world

| | | | | |
|----|----|----|----|----|
| 1 | 2 | 3 | 4 | 5 |
| 6 | 7 | 8 | 9 | 10 |
| 11 | 12 | 13 | 14 | 15 |
| 16 | 17 | 18 | 19 | 20 |
| 21 | 22 | 23 | 24 | 25 |

1, 2, 3, 4, 5, 10, 15, 20, 25, 24, 23, 22, 21,
16, 11, 6, 7, 8, 9, 14, 19, 18, 17, 12, 13



$B = 6$

3, 10, 15, 7, 12, 16, 9, 6, 5

~~3~~, ~~10~~, 15, 7, 12, 16 $C \rightarrow 3 \rightarrow 4 \rightarrow \text{X} \rightarrow 4$

Print Boundary Elements

row = 0;

col = 0;

```
for (K=1; K<N; K++) {
    Print (A[row][col]),
    col ++,
}
```

```
for (K=1; K<N; K++) {
    Print (A[row][col]),
    row ++,
}
```

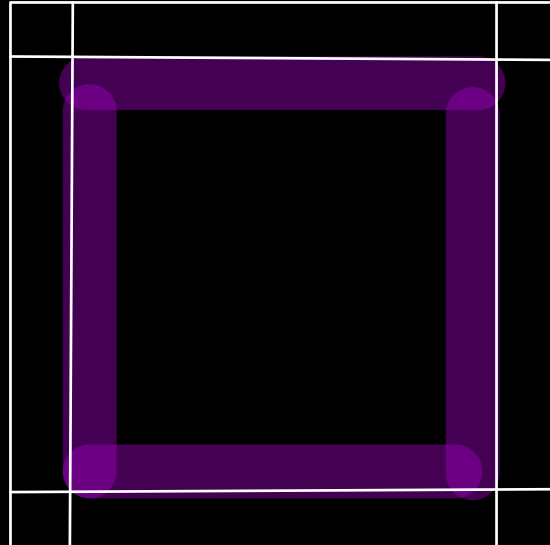
```
for (K=1; K<N; K++) {
    Print (A[row][col]),
    col --,
}
```

```
for (K=1; K<N; K++) {
    Print (A[row][col]),
    row --,
}
```

// row → 0, col → 0

| | 0 | 1 | 2 | 3 | 4 | ↓ |
|---|----|----|----|----|----|---|
| 0 | 1 | 2 | 3 | 4 | 5 | |
| 1 | 6 | 7 | 8 | 9 | 10 | |
| 2 | 11 | 12 | 13 | 14 | 15 | |
| 3 | 16 | 17 | 18 | 19 | 20 | |
| 4 | 21 | 22 | 23 | 24 | 25 | ← |

| K | A[row][col] | col |
|---|-----------------|-----|
| 1 | A[0][0] | 1 |
| 2 | A[0][1] | 2 |
| 3 | A[0][2] | 3 |
| 4 | A[0][3] | 4 |
| 5 | → <u>Breaks</u> | |



```

row = 0;
col = 0;
While (N - -) {
    for (K=1; K<N; K++) {
        Print (A[row][col]);
        col ++;
    }

```

$N > 0$
 $N > 1$
 $N > 2$
 $row \leq N/2$

?

TC: $O(N^2)$

SC: $O(1)$

Enter

```

    for (K=1; K<N; K++) {
        Print (A[row][col]);
        row ++;
    }

```

```

    for (K=1; K<N; K++) {
        Print (A[row][col]);
        col --;
    }

```

```

    for (K=1; K<N; K++) {
        Print (A[row][col]);
        row --;
    }

```

row ++;

col ++

N = N - 1

Edge case

}

$$J4 = J4 - 2;$$

H.W : \rightarrow ① While condition

② Edge case (The test case for which current algo will not work)

③ Solve same problem for $N \times M$ matrix