## 1 Sec -> 108-109 iterations

$$N \leq 10_e$$
 O(N)  $\sim$ 

Q Amazon MS Adobe

Given an array of 180. We can replace one of the with a 1. Return the count of mare consecutive L's in the array.

Ex: 1 1 0 1 1 0 1 1

0 1 2 3 4 5 6 7 8 1 1 0 1 1 1 1 1 1 1

Ex:  $0 \stackrel{1}{1} \stackrel{2}{1} \stackrel{3}{1} \stackrel{4}{1} \stackrel{5}{0} \stackrel{6}{1} \stackrel{7}{1} \stackrel{8}{1} \stackrel{3}{1} \stackrel{1}{0} \stackrel{1}{1} \stackrel{1}{1} \stackrel{1}{0} \stackrel{1}{1} \stackrel{1}{1$ 

for every 0 in the array:

- . Count number of consecutive 1's on left  $\rightarrow$  l
- · Count number of consecutive 1's on right -> 27
- · if ( l+2+1 > ans ) { ans = l+2+1}

Edge Care: 97 all are 1 (no 0's) 

Every element is getting accessed at man 3 times

i. # ideration - 3N
TC: O(N)

Amazon Q Given an array of I & O. We can surap one of the Direct - i with a 1. Return the count of max consecutive I's in the array. Aclobe Ex; 1 1 0 1 1 0 1 1 1

0 1 2 3 4 5 6 7 8 0 1 0 1 1 1 1 1 1

- · Calculate total count of 1:3
- · for every 0 in the array:
  - Count number of consecutive 1's on left  $\rightarrow$  l
  - · Count number of consecutive 1's on right -> 2

Cerunt =  $\begin{cases} 1+2x & \text{if } (1+x) = 1 \\ 1+2x+1 & \text{if } (1+2x) < 1 \end{cases}$ Cerunt =  $\begin{cases} 1+2x & \text{if } (1+2x) < 1 \\ 1+2x+1 & \text{if } (1+2x) < 1 \end{cases}$ 

· if ( Count > ans) { ans = Count }

1 1 0 <u>1 1</u> total Count of 1's -> 4 (1+7) = = 4 ? 0111

Amazon Q No of triplets

Given an array. Count the number of triplets i, j 4 k such that

i < j < k & A[i] < A[j] < A[k]

j 7 K (i)A AGj A[K] 0 T 2 2 6 g 1 0 2 6 70 3 4 0 2 4 70 0 2 4 2 9 T O T 4 6 9 10

icleq: iterate over all possible striplets

cerunt = 0;

for ( i=0; i < N; i++) {

For ( j=i+1; j < N; j++) {

for ( k=j+1; k < N; k++) {

if ( A[i] < A[j] as A[i] < A[k])

}

count ++;

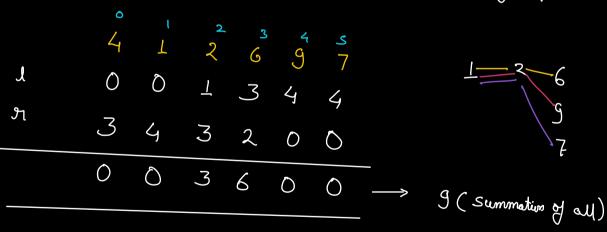
}

Hint: In how many triplets index 3 is in middle

No of triplets in which any A[x] will be in middle:

No of elements smaller than Airi on left of x

No of elements greater than AINI on right of n



for every element A[i]; -> N iterations

(N-1)

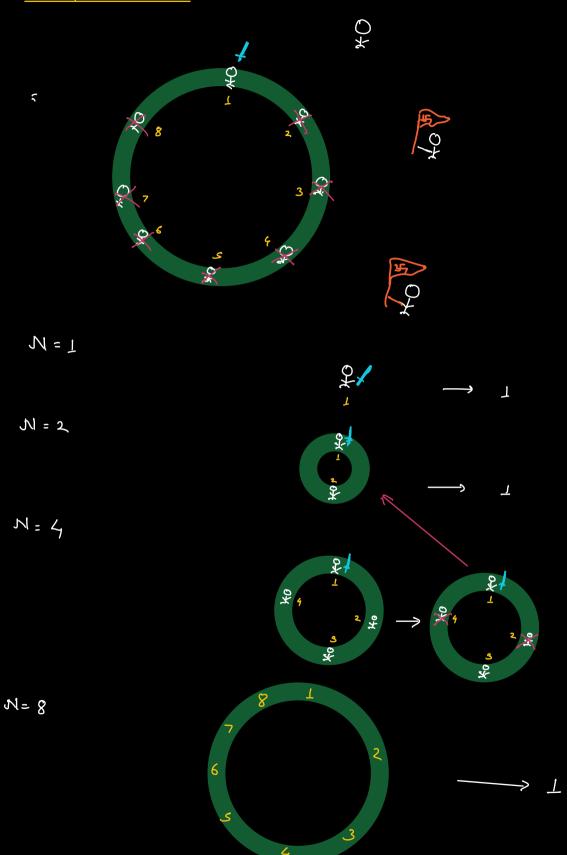
Iterates from inden 0 to i-1 & Count the elements smaller than A[i] -> l

uterates from inden i+1 to N-1 & Count the elements greater than A[i] -> n

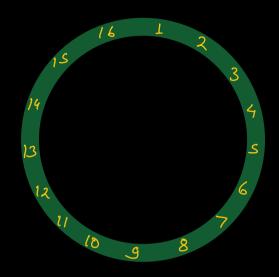
· ans = ans + 1x9

TC: O(N2)

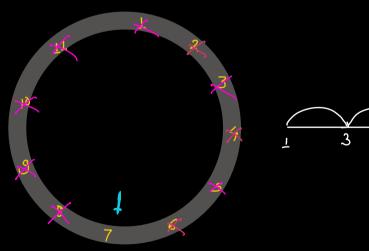
## Josephus Problem

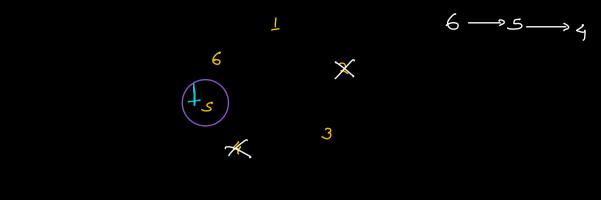


N = 16



If N is a fearer of 2 — 1 always win in that case (Who ever starts Killing)





$$N=100 \xrightarrow{36 \text{ kill}} \rightarrow 64 \qquad 1 + 36x2$$

$$\Rightarrow 73$$

$$N=11 \xrightarrow{3 \text{ kill}} \rightarrow 8$$

$$1 + 2x3$$

$$\Rightarrow 7$$