

SELECTED TOPICS IN ADVANCED NETWORKING FOR SCIENTIFIC APPLICATIONS

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Agenda

- Introduction
- Part I: Application-Aware Networks
 - Dynamic Circuits
 - OGF Network Services Interface
 - Examples (ANSE)
- Part II: Software Defined Networking
 - Introduction to SDN
 - OpenFlow
 - Programmable Networks
 - Use cases
- Part III:
 - Content Centric Networking
- Additional Resources
 - Networks for experimentation



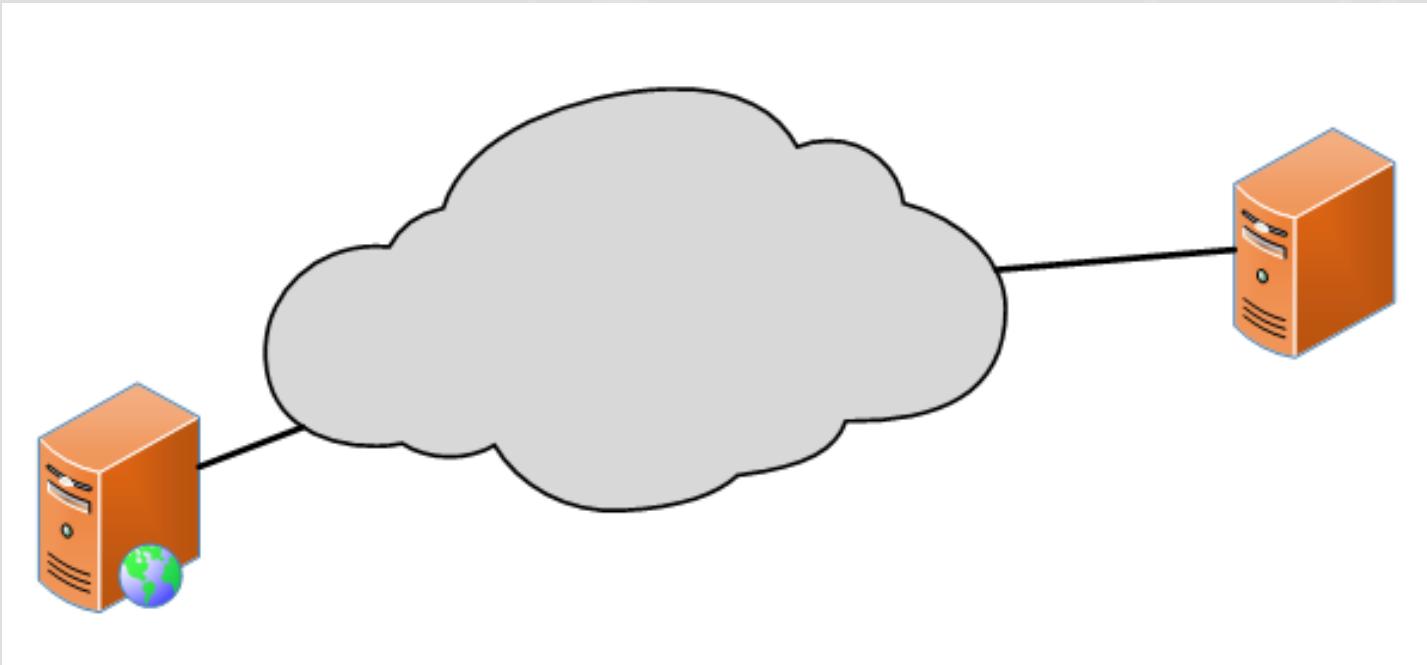
INTRODUCTION



Different Views of a Network

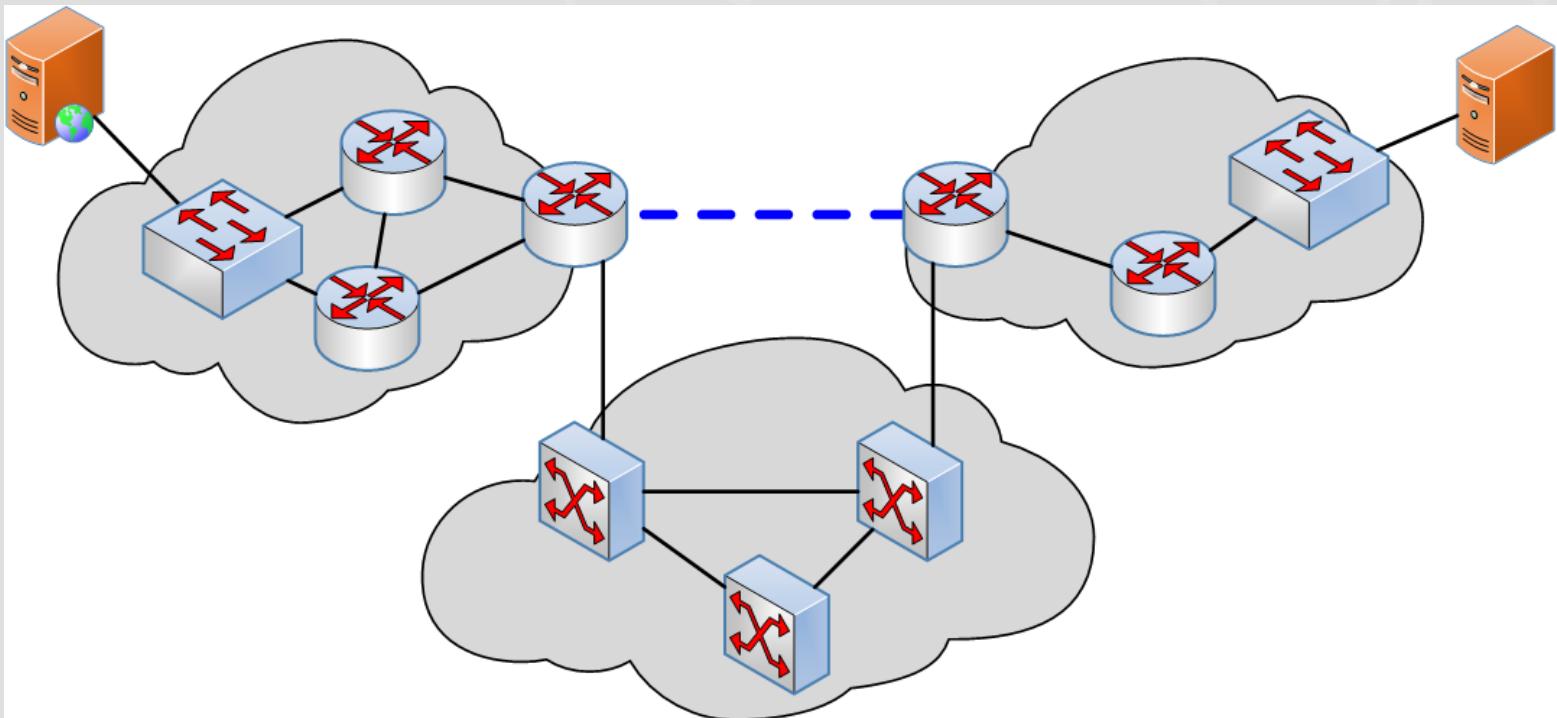


... what a user might see:



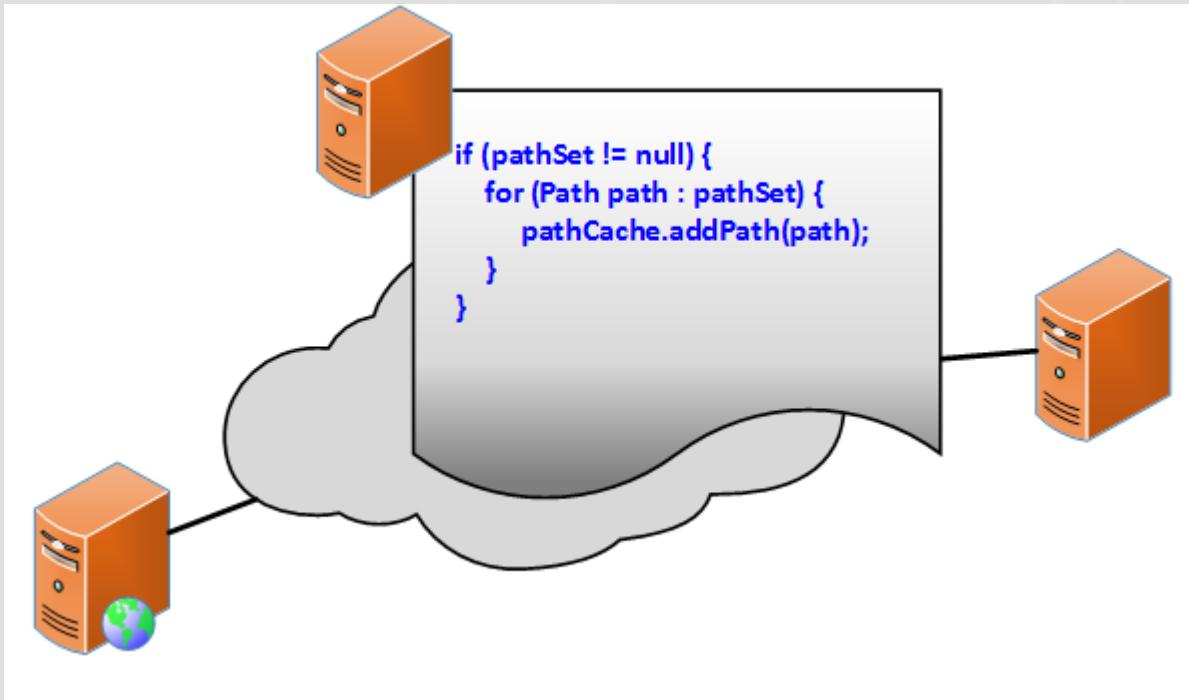
Different views of a Network

... what a network engineer would see:



Different Views of a Network

... what an SDN network engineer would see:



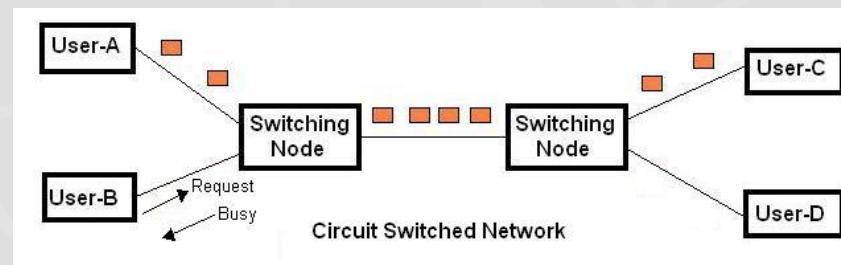
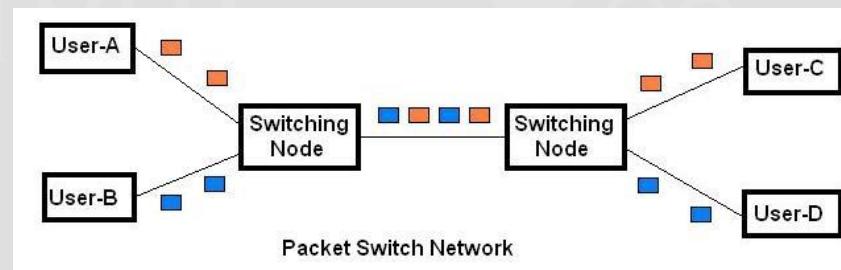
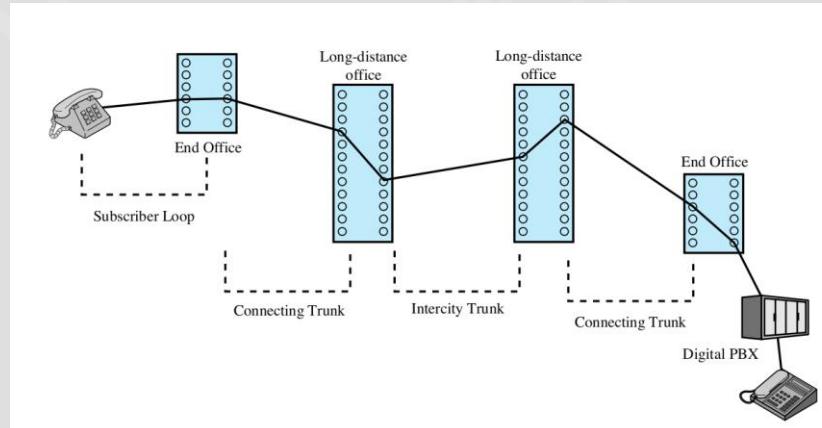
But first...

- A refresher of basics and terms...
(not a formal course on networking, just a collection of terms and definitions needed for the discussion later)



Circuit vs Packet Switched Networks

- Circuit Switching
 - Dedicated communication path between two stations
 - Set up prior to data exchange
 - Usually through several nodes in the network
 - Example: telephone network
- Packet Switching
 - Data sent in packets
 - Each packet's header is inspected at each network node
 - Packets are passed from node to node based on header information and (local) routing database
 - Example: IP network
- Virtual Circuit Switching
 - Emulation of circuit switching on packet switched infrastructure

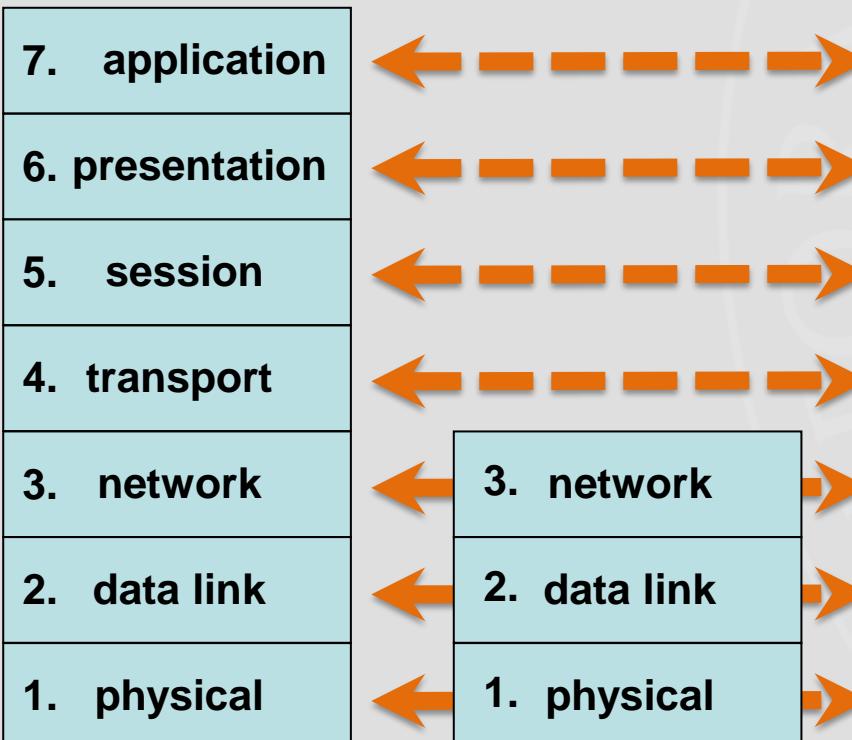


Network Layers

- Communication happens between corresponding entities in a layered structure

OSI Reference Layers

End System



End System



Function

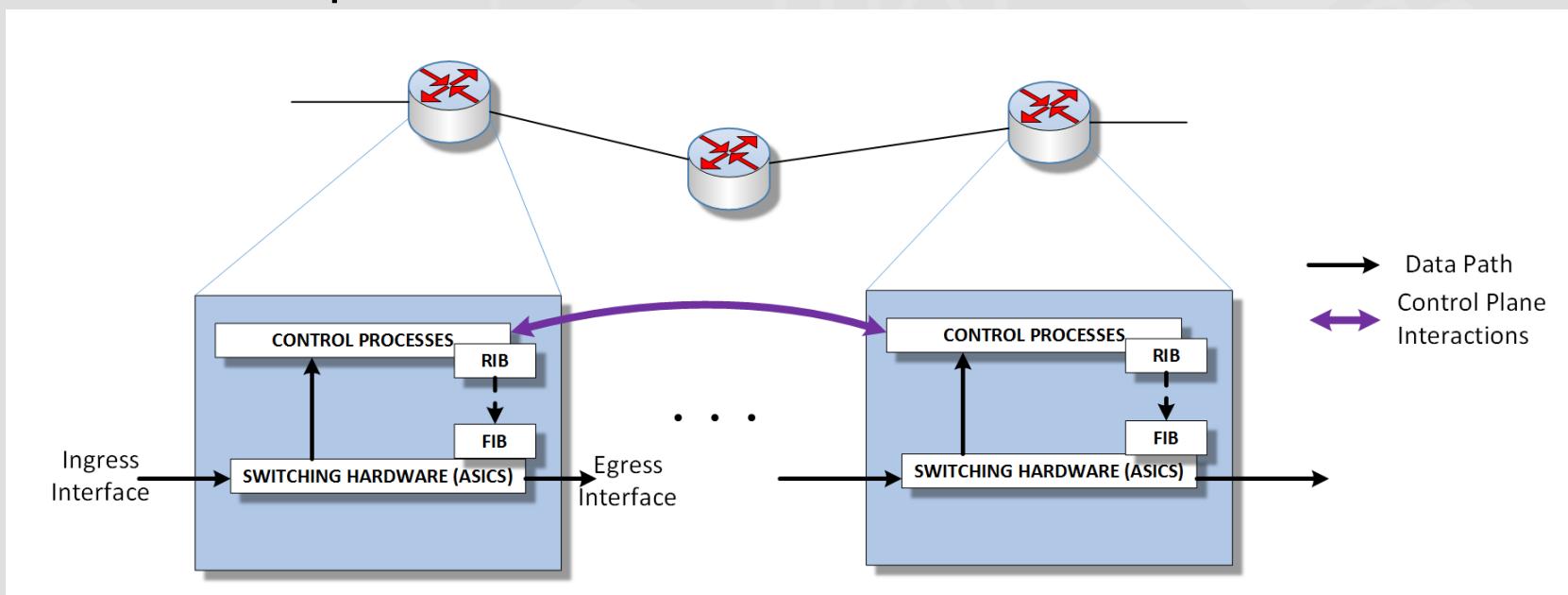
| | |
|--|----------|
| Interface between application and communication sw | http |
| Data formats, encryption | |
| Start/control/stop sessions | |
| Segmentation and reassembly, error recovery | TCP, UDP |
| End-to-end delivery of packets | IP |
| Data delivery across a link or medium | Ethernet |
| Physical characteristics | SONET |

Packet Router



Control & Data Planes

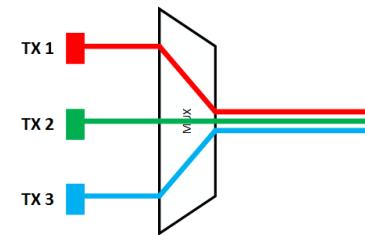
- **Data Plane:** processing of incoming data packets
 - Inspect, forward or drop
- **Control plane:** processes to build topology (RIB) and forwarding tables (FIB)
 - Needed to populate Forwarding Information Base used in the data plane
- In a traditional networks, each node operates processes in both control and data plane



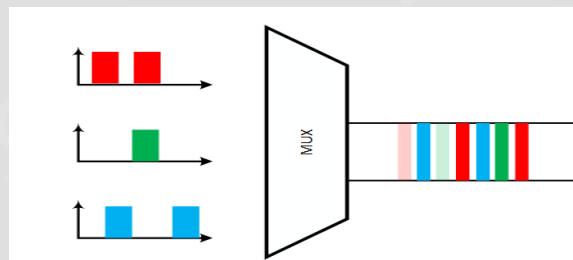
Multiplexing

- Multiplexing is used to enable sharing of transmission medium between multiple devices
- Most common multiplexing schemes:

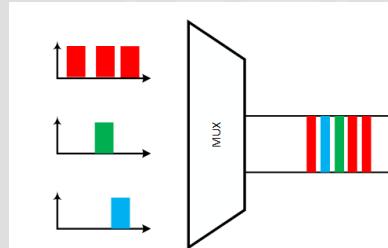
- Wavelength-Division multiplexing



- Time-Division Multiplexing



- Statistical multiplexing



- But also Space-Division multiplexing...

Circuit switching revisited



- Optical circuit switching equipment operate at Layer 1
 - Or even at “Layer 0” like e.g. MEMS switches
- Layer 1 optical equipment can switch based on wavelengths - Called Lightpath or Lambda-switching
- **Virtual circuit** connections above physical layer
 - SONET/SDH: TDM channels with defined capacity
 - MPLS emulates circuit connections using bandwidth profiles
 - TCP: a logical circuit connection between two end-systems
 - As opposed to UDP’s datagrams



Packets, Circuits and Flows



- Packet switches and routers forward based on each individual packet's header information
 - In IP networks, typically only IP Destination address is matched against the Routing Information Base
 - plus QoS fields
 - sometimes also source address (PBR)
- Flow-based forwarding on the other hand...
 - Flow definition based on a set of parameters , such as e.g. {IP_SRC, IP_DST, TCP_PORT}
 - Network device forwards packets based on forwarding database information for that flow – each packet in the flow takes the same path
- Flow-based forwarding is encountered in e.g. Link Aggregation scenarios (LAG, ECMP), as well as being the basis of OpenFlow



Other contents of a network environment

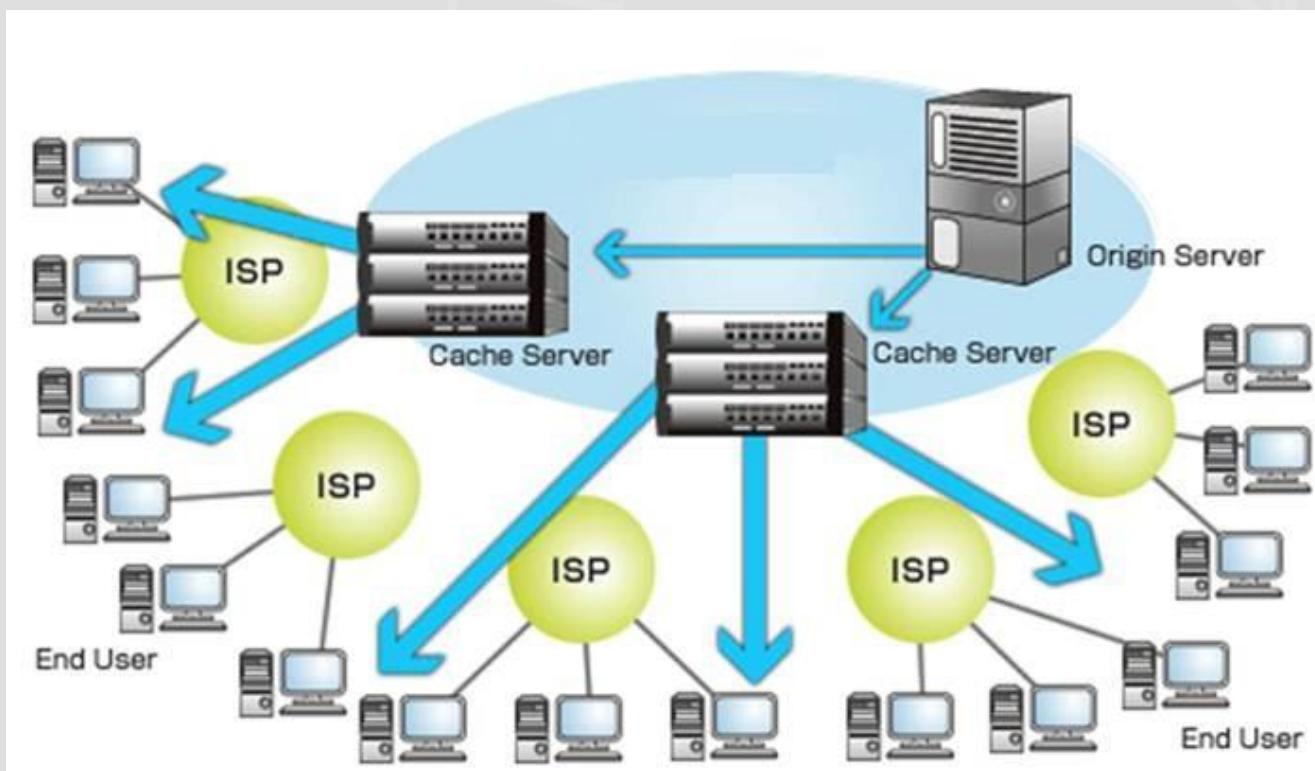


- In practical deployment, networks do not just forward packets
- Other services needed to function:
 - DNS
 - Possibly DHCP
 - AAA
 - NIS, LDAP, Shibboleth, etc.
 - Monitoring
- Networks deliver content...



Content Delivery Networks / CDN

- Goal: reduce WAN latencies for data delivery
- Strategically placed Cache Servers
- Data replicated from the Origin Server(s)
- Application-level technology
- Usually an overlay on top of existing IP network infrastructure



APPLICATION AWARE NETWORKING

And network-aware applications



Network – Application Interface



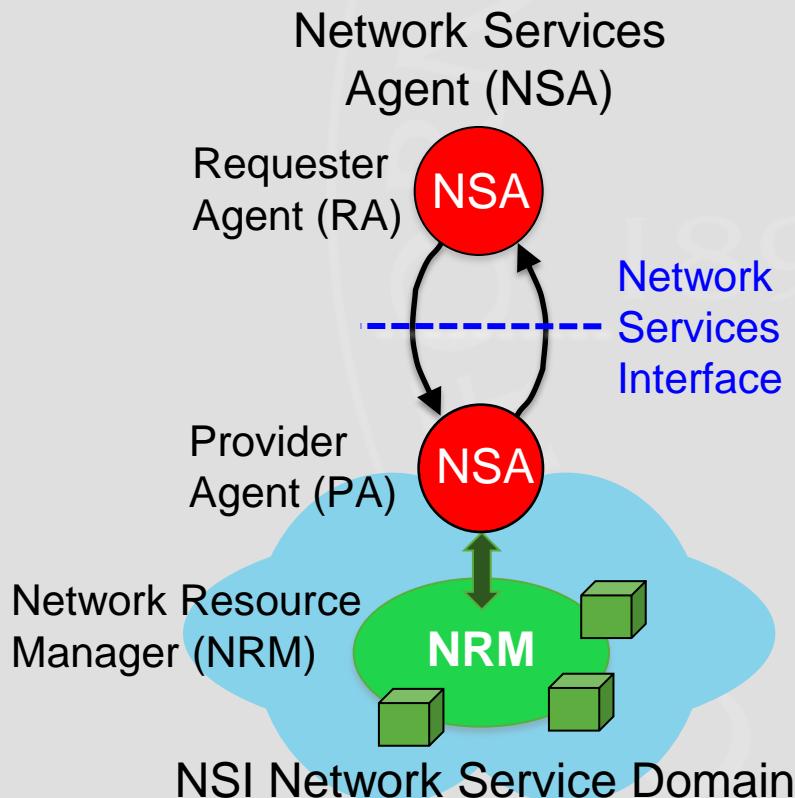
- Any distributed system needs some form of network interaction
- Basic programming interface: Sockets
 - “puts bits on wire”
 - Restricted QoS
- Network Control
 - Reserve capacity
 - usually a NOC procedure, unless BoD system used
 - Prioritize traffic
- Network Monitoring and Analytics
 - To base smart decisions on
 - Reachability
 - Topology
 - Available capacity



OGF Network Services Interface (NSI)



- A standardized service interface between network domains
 - Note: A computing site is also a network domain!
- Open Grid Forum Working Group (NSI-WG)

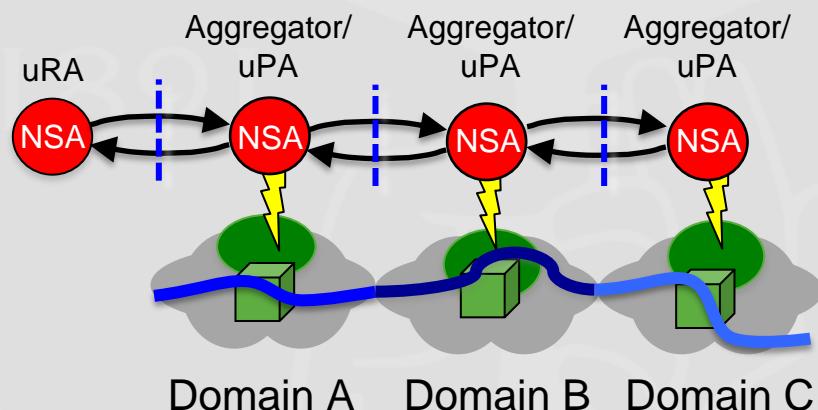
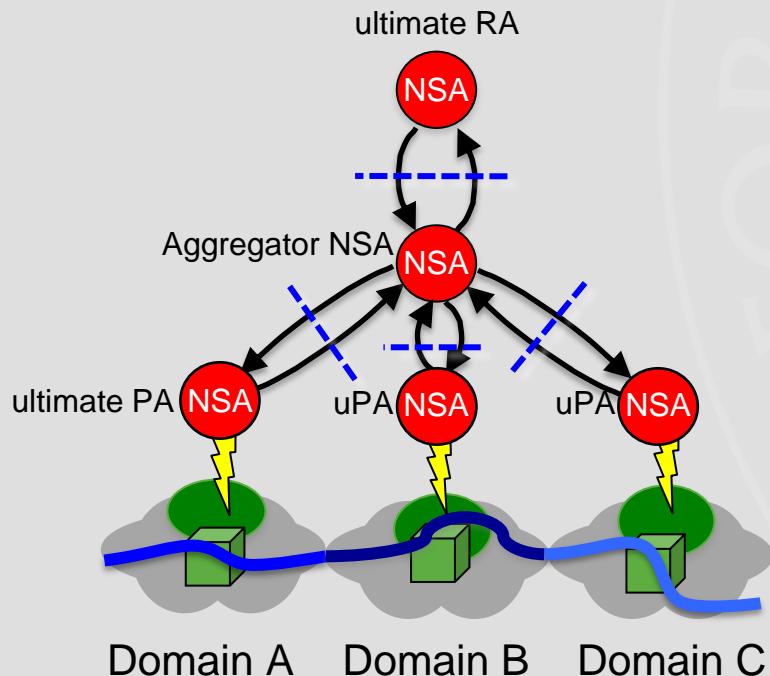


picture C. Guo, ESnet



NSI multi-domain service construction

- Two ways defined for “chaining” services: tree and chain
- Note: “ultimate Requester Agent” can be an end-user app



pictures C. Guok, ESnet

- Currently foreseen services:
 - Connection Service (NSI-CS)
 - Topology Service (NSI-TS)
 - Discovery Service (NSI-DS)
 - Switching Service (NSI-SS)
- Future Services:
 - Monitoring Service
 - Protection Service
 - Verification Service
 - ...



- NSI-CS: first NSI service standardized, currently v2.0
- Advance-reservation protocol
 - Mandatory Reservation parameters:
 - A-point, Z-point
 - Optional parameters:
 - Start time, end time
 - Bandwidth
 - Labels/VLAN IDs
- V2.0 supports optional modification of a reservation
 - Start time, end time and bandwidth



NSI Request Messages



| NSI CS Message (abbreviation) | SM | Synch. /Asynch. | Short Description |
|--|-----|--------------------|--|
| reserve (rsv.rq) | RSM | Asynch | The <i>reserve</i> message allows an RA to send a request to reserve network resources to build a Connection between two STP's. |
| reserveCommit (rsvcommit.rq) | RSM | Asynch | The <i>reserveCommit</i> message allows an RA to request the PA commit a previously allocated Connection reservation or modify an existing Connection reservation. |
| reserveAbort (rsvabort.rq) | RSM | Asynch | The <i>reserveAbort</i> message allows an RA to request the PA to abort a previously requested Connection that was made using the <i>reserve</i> message. |
| provision (prov.rq) | PSM | Asynch | The <i>provision</i> message allows RA to request the PA to transition a previously requested Connection into the Provisioned state. A Connection in Provisioned state will activate associated data plane resources during the scheduled reservation time. |
| release (release.rq) | PSM | Asynch | The <i>release</i> message allows an RA to request the PA to transition a previously provisioned Connection into Released state. A Connection in a Released state will deactivate the associated resources in the data plane. The reservation is not affected. |
| terminate (term.rq) | LSM | Asynch | The <i>terminate</i> message allows an RA to request the PA to transition a previously requested Connection into Terminated state. A Connection in Terminated state will release associated resources and allow the PA to clean up the RSM, PSM and all related data structures. |

Full messages listing in <http://www.ogf.org/documents/GFD.212.pdf>



NSI Service Agent implementations



- **AutoBAHN** : GÉANT (EU)
- **BoD** : SURFnet (NL)
- **DynamicKL** : KISTI (KR)
- **G-LAMBDA-A** : AIST (JP)
- **G-LAMBDA-K** : KDDI Labs (JP)
- **OpenNSA** : NORDUnet (DK, SE, NO, FI, IS)
- **OSCARS** : ESnet (US)



Automated GOLE Fabric

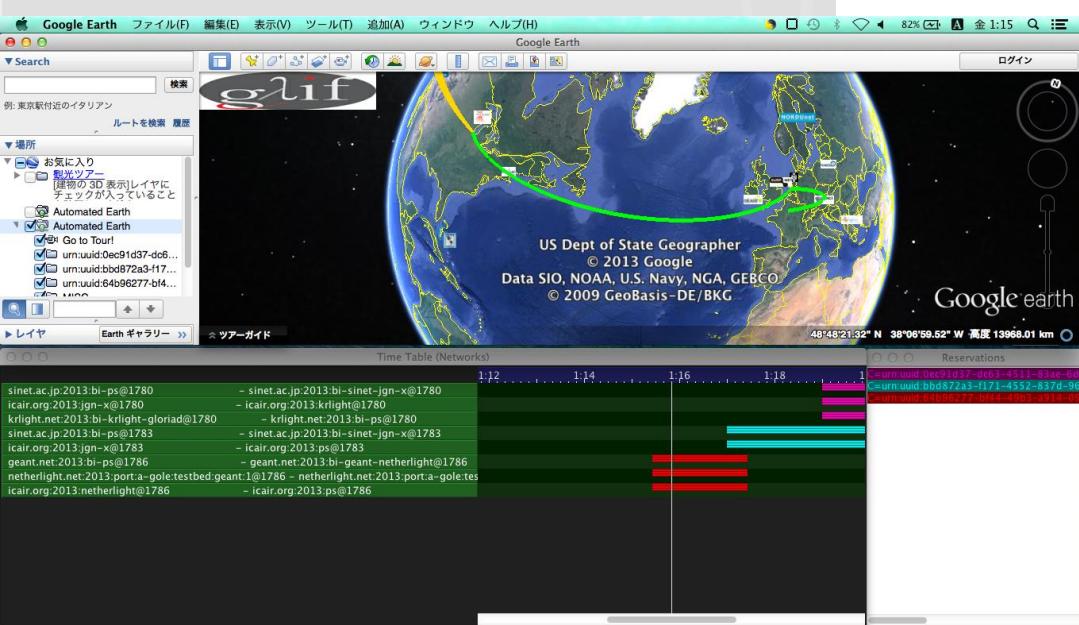


John MacAuley, ESnet

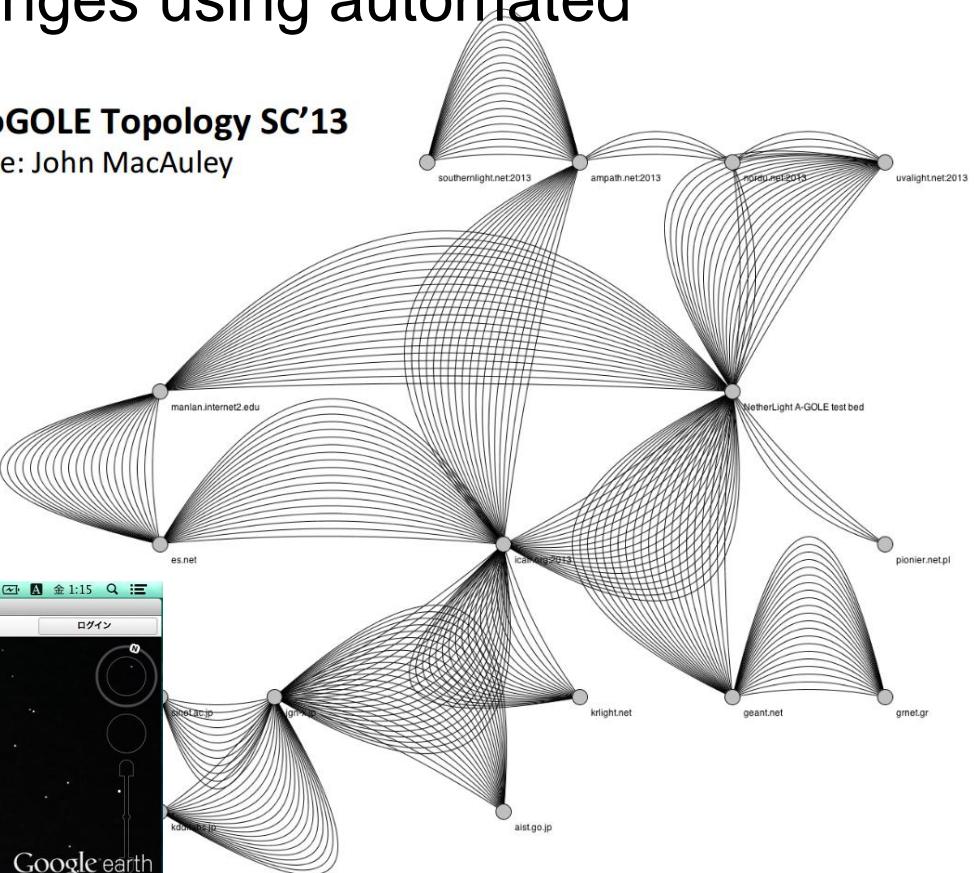


NSI-CS in Action: GLIF AutoGOLEs

- GLIF Open Lightpath Exchanges using automated provisioning
- Currently in R&D
- Demonstrated e.g. at SC'13
- Next demo planned for GLIF and SC'14



AutoGOLE Topology SC'13
Source: John MacAuley



Network of radio telescopes



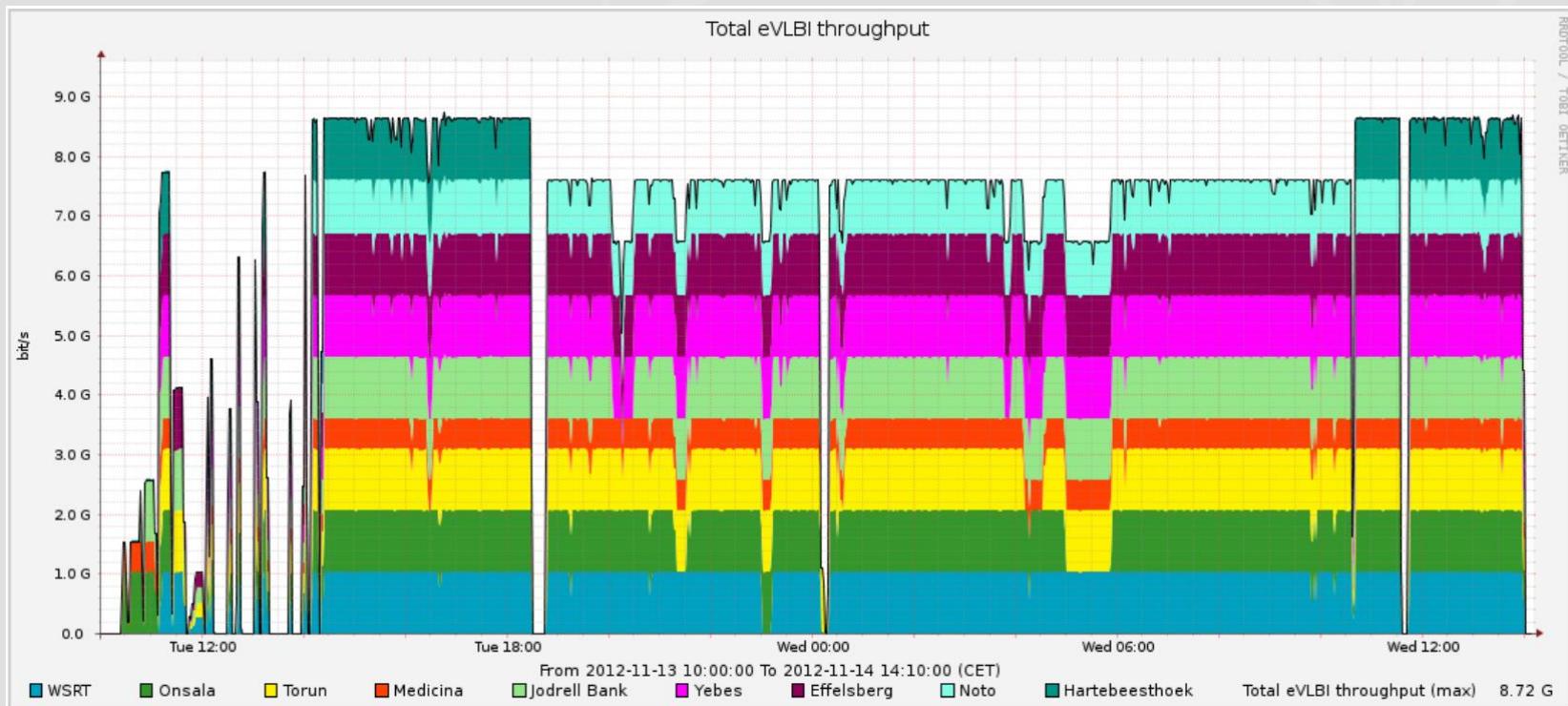
Image by Paul Boven (boven@jive.nl). Satellite image: Blue Marble Next Generation, courtesy of Nasa Visible Earth (visibleearth.nasa.gov).

More info on <http://www.evlbi.org/evlbi/evlbi.html>



A typical eVLBI run

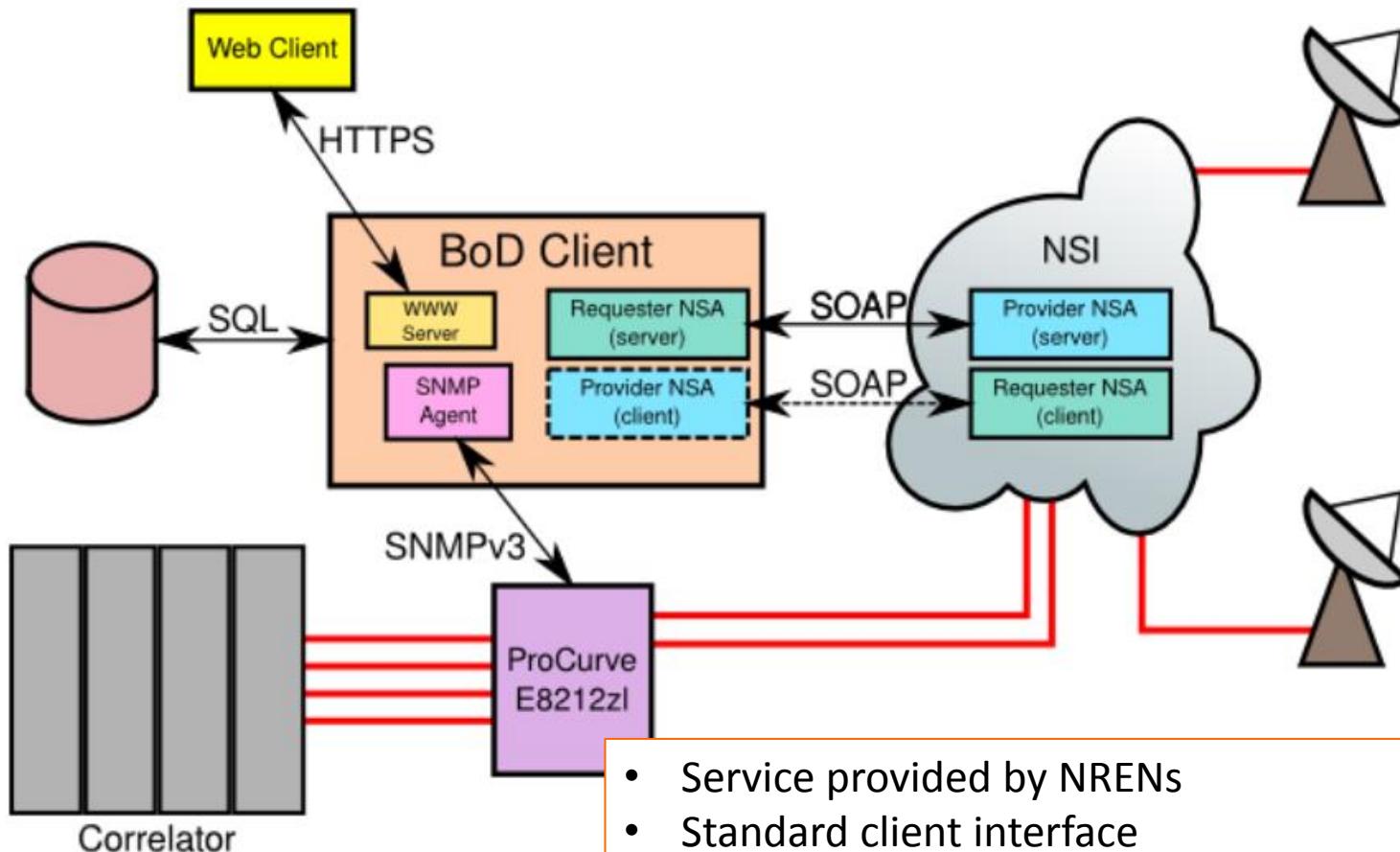
- 8-12 radio telescopes
- 1Gbps per telescope (future: 4Gbps)
 - Steady streams of data from antennas to correlator
 - Low jitter very important
- 8-12 hours
- 30-65 TB



Paul Boven, JIVE



NEXPREs NSI client

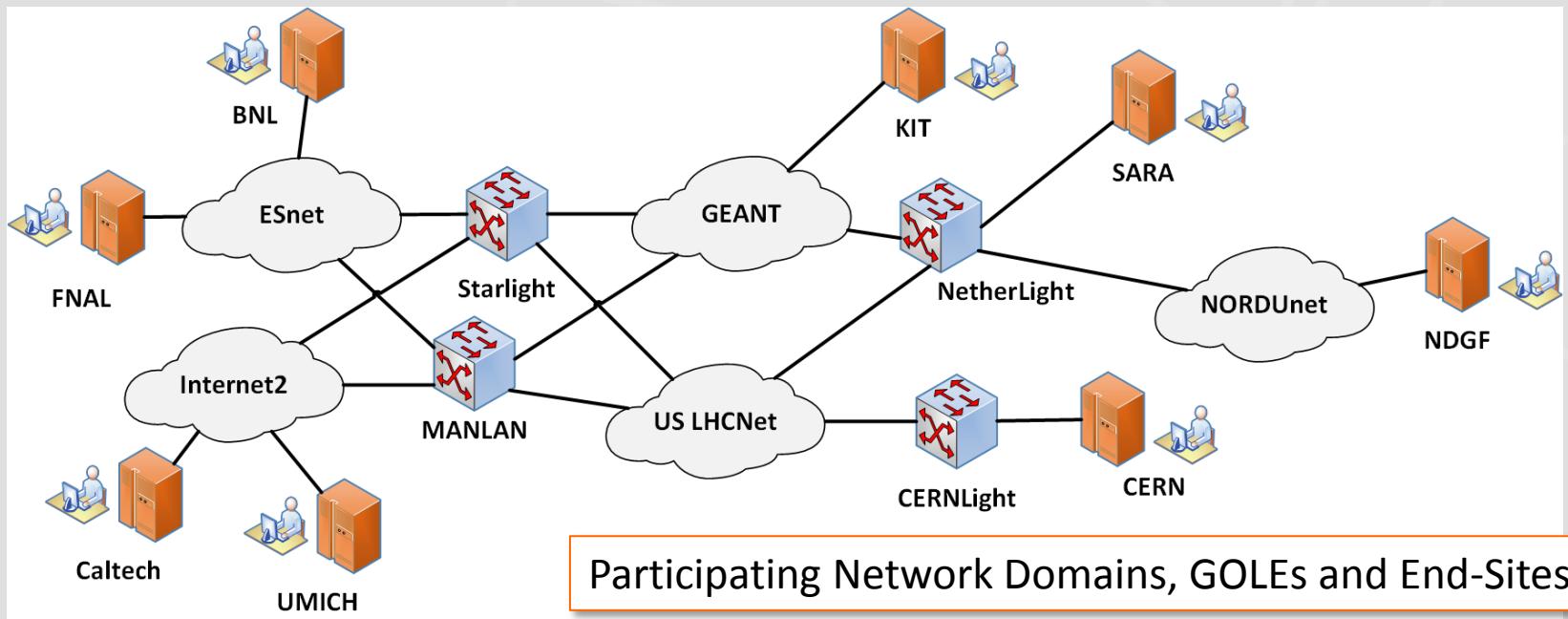


- Service provided by NRENs
- Standard client interface
- Client software developed by eVLBI collaboration

NSI for LHC: LHCONE Point-to-Point Service



- LHC Open Network Environment
 - VRF for current multipoint production use
 - Experiment/demonstration: Bandwidth-on-Demand
- Target: demonstrate multi-domain bandwidth reservation capability



- Status: under construction
 - Multi-domain service created first, then connect end-sites

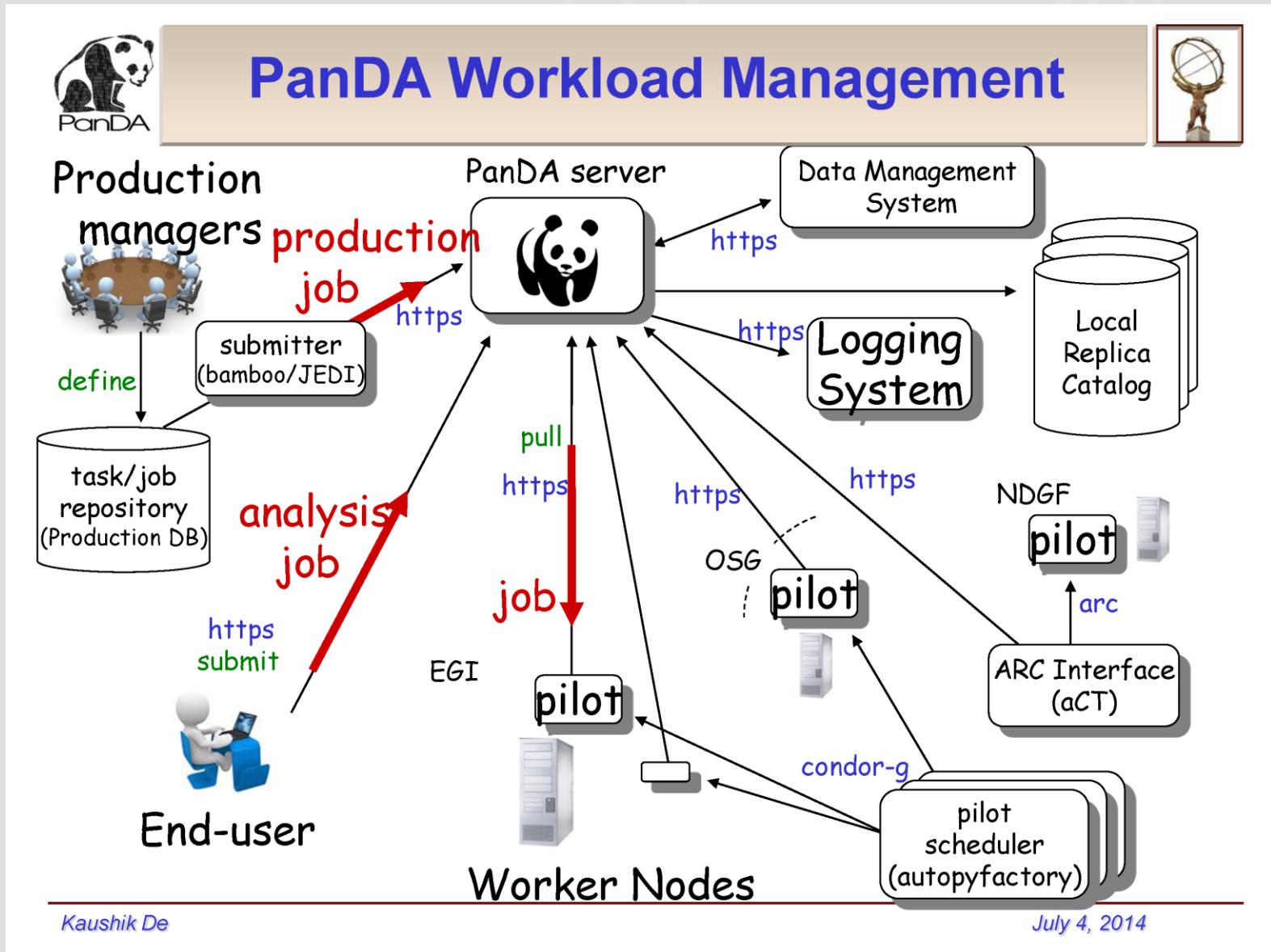


Building Network Awareness in LHC



- General: network monitoring in OSG and WLCG
- Based on information from monitoring systems such as MonALISA and PerfSONAR
- Specific project in CMS and ATLAS Experiments:
Advanced Network Services for Experiments
 - Network Integration into
 - Workflow management (PanDA)
 - Data movement management (PhEDEx)
 - Measurement: PerfSONAR and MonALISA integration
 - Control: interface to provisioning systems
(DYNES/OSCARS, NSI)





Kaushik De

July 4, 2014

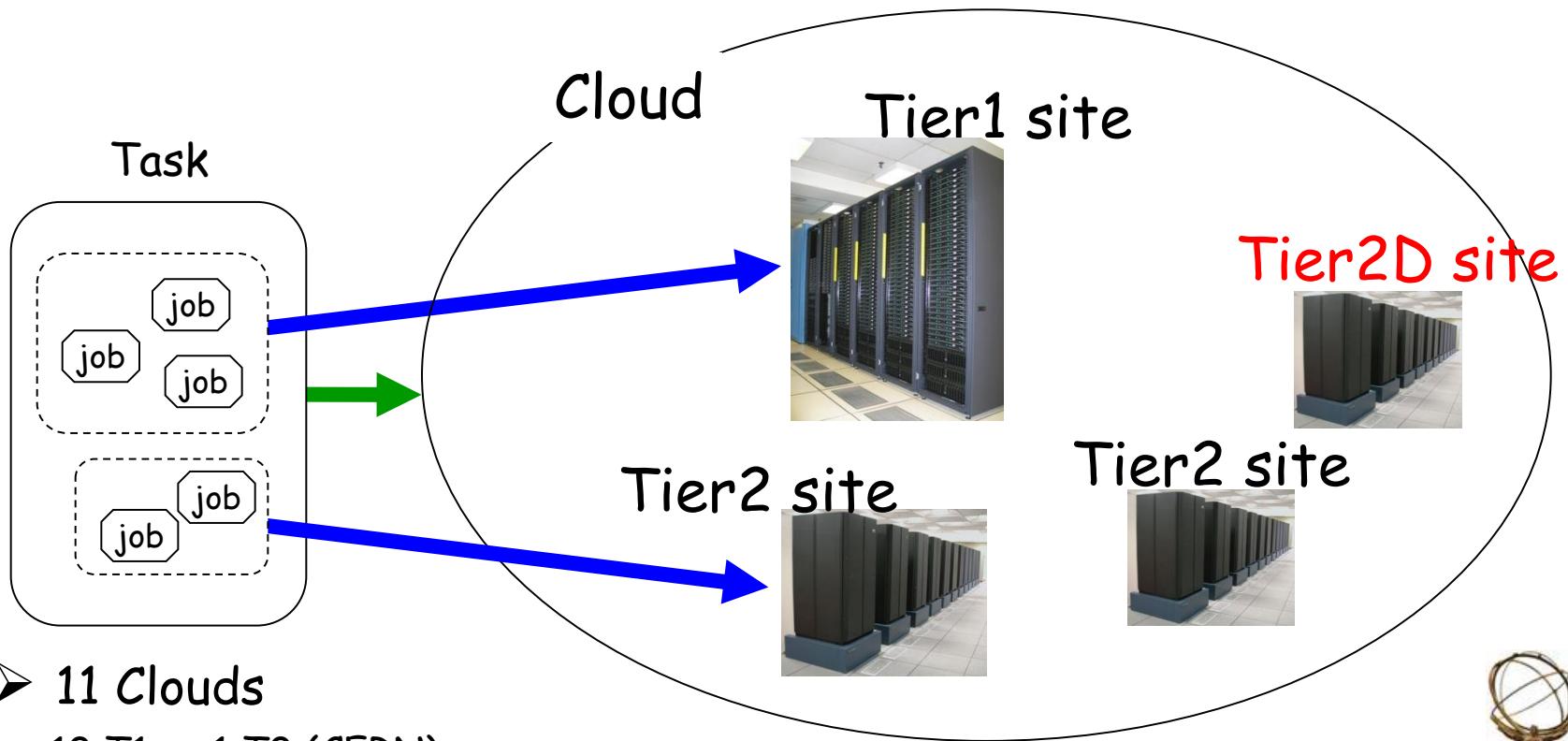
Network Integration in PanDA



- Concept: utilize network as a resource like other resources such as CPU, disk storage
 - Use network information for FAX brokerage
 - Brokerage should use concept of nearby sites
 - Jobs are sent to site with best weight, not necessarily the site with local data or available CPUs
 - Use network information for cloud selection
 - Best T2D site should be selected based on throughput from T1 to T2D measurements
- Network measurements are available at SSB (Site Status Board, Network view)
 - FAX xrdcp rate metric used for FAX brokerage
 - DDM Sonar metrics used for cloud selection



ATLAS Computing Model



➤ 11 Clouds

10 T1s + 1 T0 (CERN)

Cloud = T1 + T2s + T2Ds (except CERN)

T2D = multi-cloud T2 sites

➤ 2-16 T2s in each Cloud

Task → Cloud
Task brokerage
Jobs → Sites
Job brokerage



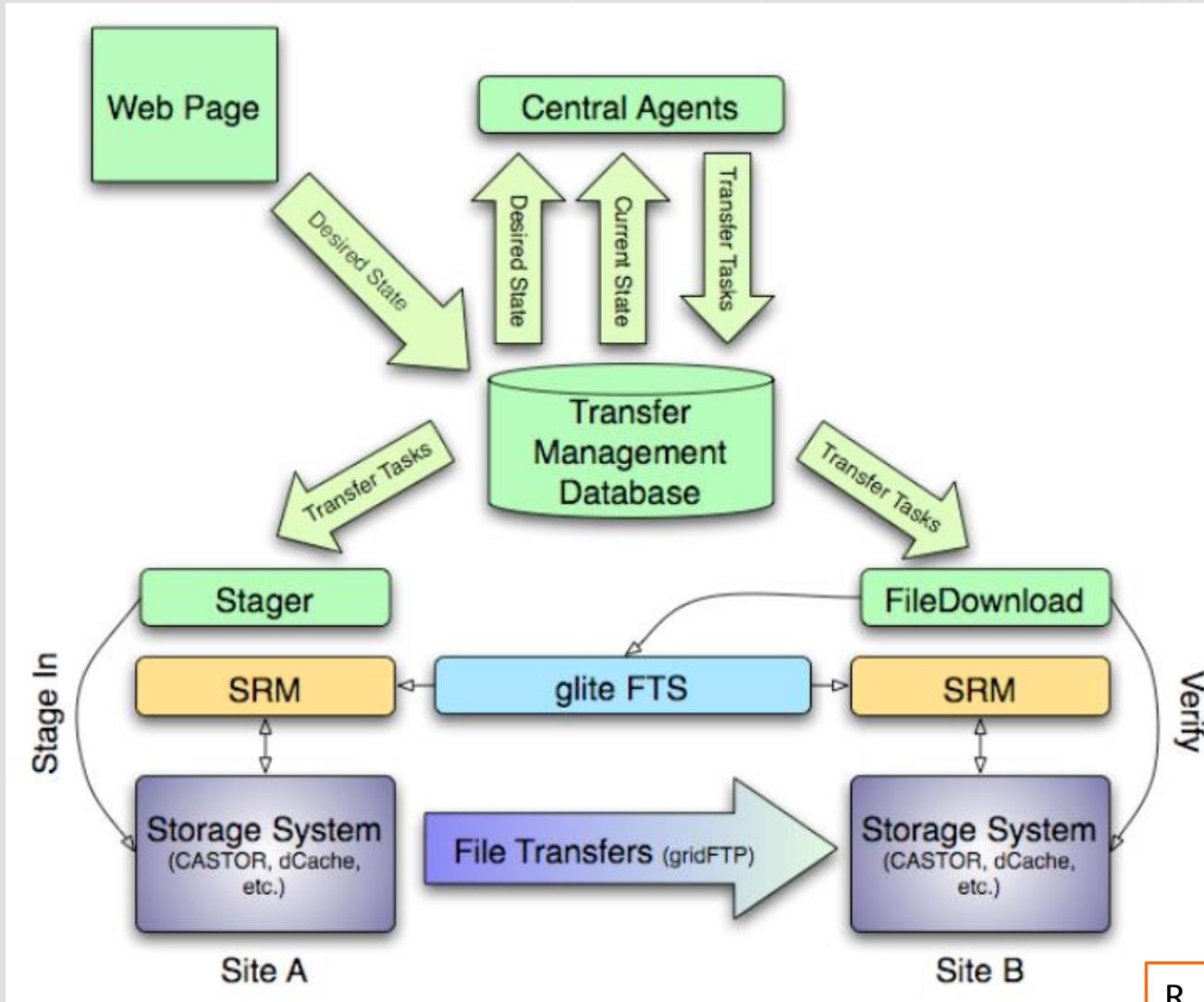
Kaushik De, UTA



- Based on
 - Important to PanDA users
 - Enhance workload management through use of network
 - Should provide clear metrics for success/failure
- 1. Improve User Analysis Workflow
 - Include network information for routing of jobs to T1/T2 sites
- 2. Cloud Selection:
 - Optimize choice of T1-T2 pairings
 - Automate using network information
- Both use cases are development and testing phase



- PhEDEx is the CMS data management toolkit

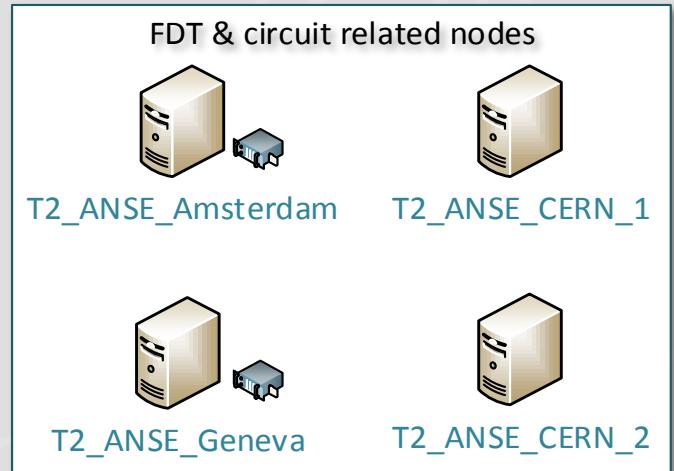


R. Egeland, 2008

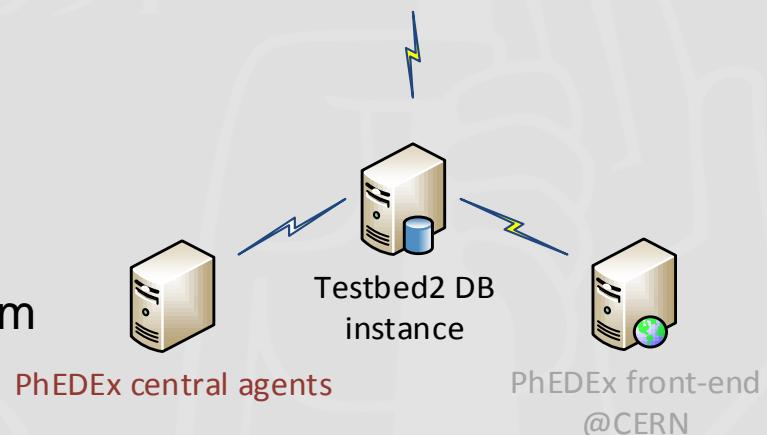
ANSE/PhEDEx – Dynamic Circuits



- Several points of circuit integration into PhEDEx
 - At the transfer-job level
 - At the link level (FileDownload agent)
 - At the instance level (FileRouter agent)



- Currently using a distributed testbed
 - T2_ANSE_CERN_1 & T2_ANSE_CERN_2
 - Both PhEDEx and storage nodes
 - T2_ANSE_Geneva & T2_ANSE_Amsterdam
 - PhEDEx and storage nodes separate
 - High speed link between AMS & GVA
 - 4x4 SSD software RAID 0 arrays

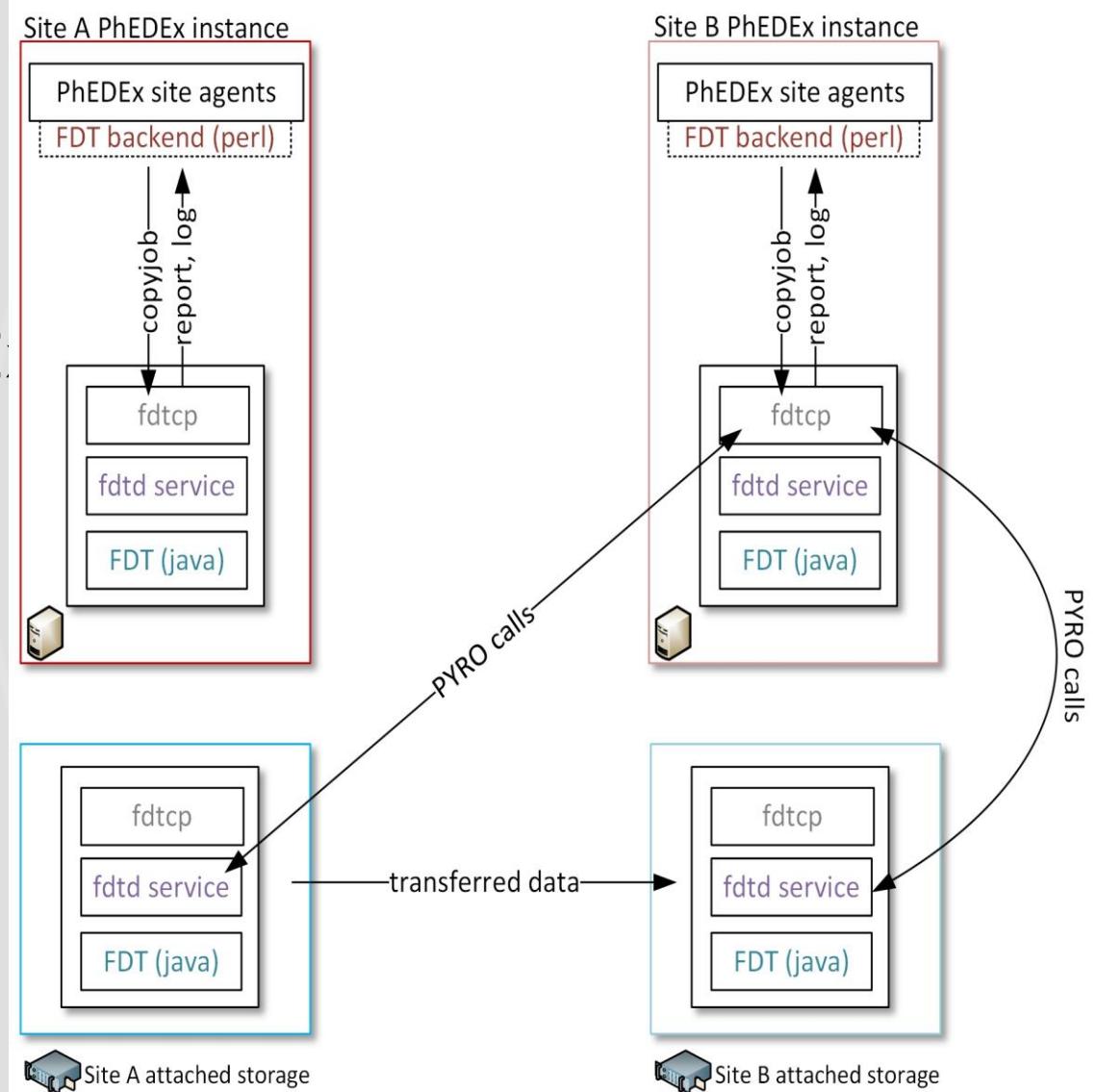


V. Lapadatescu, T. Wildish



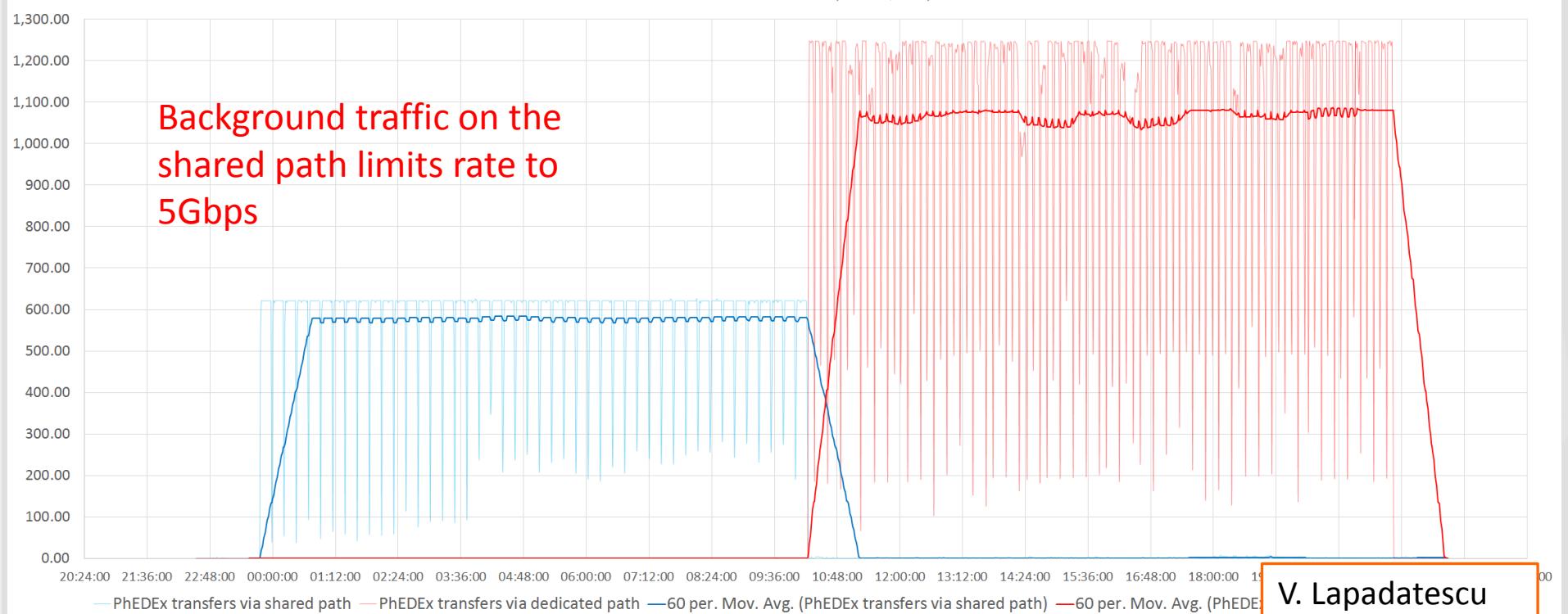
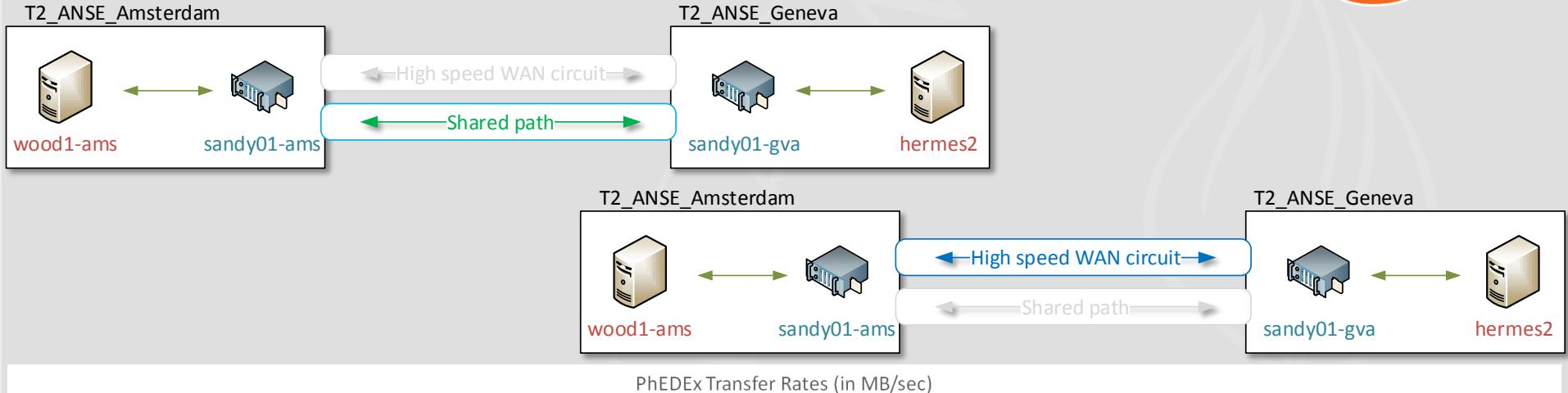
Circuits in PhEDEx at transfer level

- FDT transfer tool integrates IDCP (OSCARS) calls
- Integrating FDT as transfer tool in PhEDEx naturally includes BoD capability
- Work ongoing on integration at FileDownload agent level



V. Lapadatescu

PhEDEx BoD Trials – FileDownload agent



V. Lapadatescu



Network-Aware Applications Summary



- Network awareness can improve overall system performance
 - through acting on precise, real-time data on network state
 - through creating application-specific topologies such as point-to-point virtual circuits
- Network Services Interface (NSI) standard released, several implementations in development at many NRENs
- More examples of network-application interaction in the SDN part later



SDN SOFTWARE DEFINED NETWORKING

Where we encounter OpenFlow and intelligent networks



Network business the traditional way

- Proprietary hardware, proprietary software
 - IPR
 - provide business edge
 - vendor lock-in
- Effects:
 - closed software
 - innovation slow, driven by vendors only
 - difficult to develop and evaluate new ideas



Drivers behind SDN



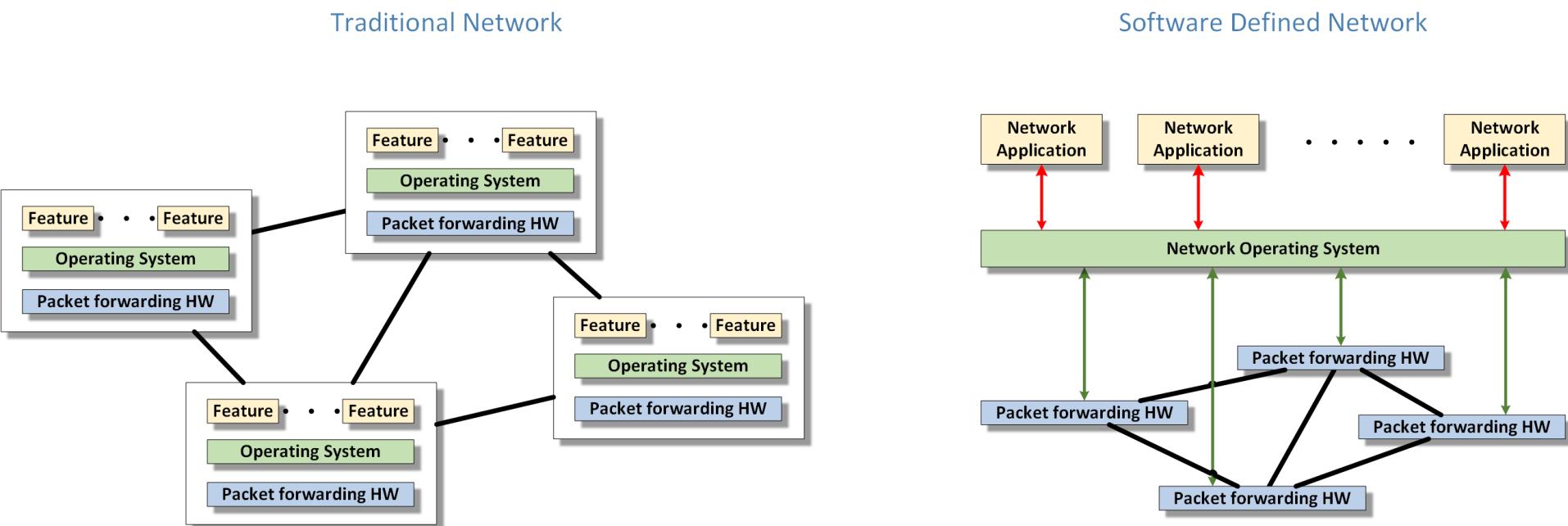
- Change in traffic patterns
 - away from single client - single server
 - local as well as wide area
- Appearance of cloud services
 - need for security, flexibility, scalability
- Manage complexity on large scales
 - Appearance of huge data centers
 - Multi-tenant facilities
 - Often global connectivity requirements

More in ONF SDN whitepaper at <http://www.opennetworking.org>



Software Defined Networking

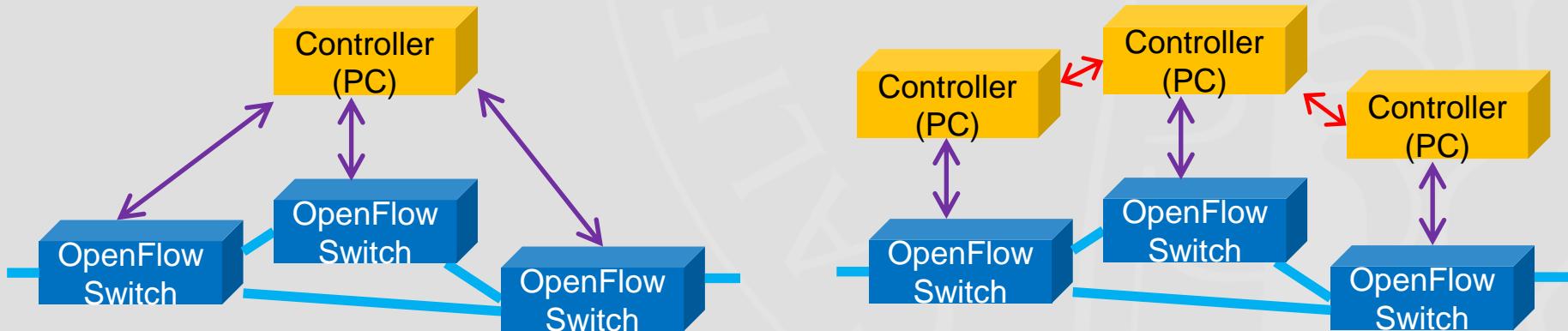
- Basic SDN Paradigm:
Separation of Network control plane from the data plane



- Enables network control by applications; provides an API to **programmatically** define network functionality

Central control

- Logically centralized control:
 - simplified operation
 - cost reduction
 - faster reconfiguration -> increased efficiency
- Physically distributed infrastructure:
 - scalability
 - redundancy



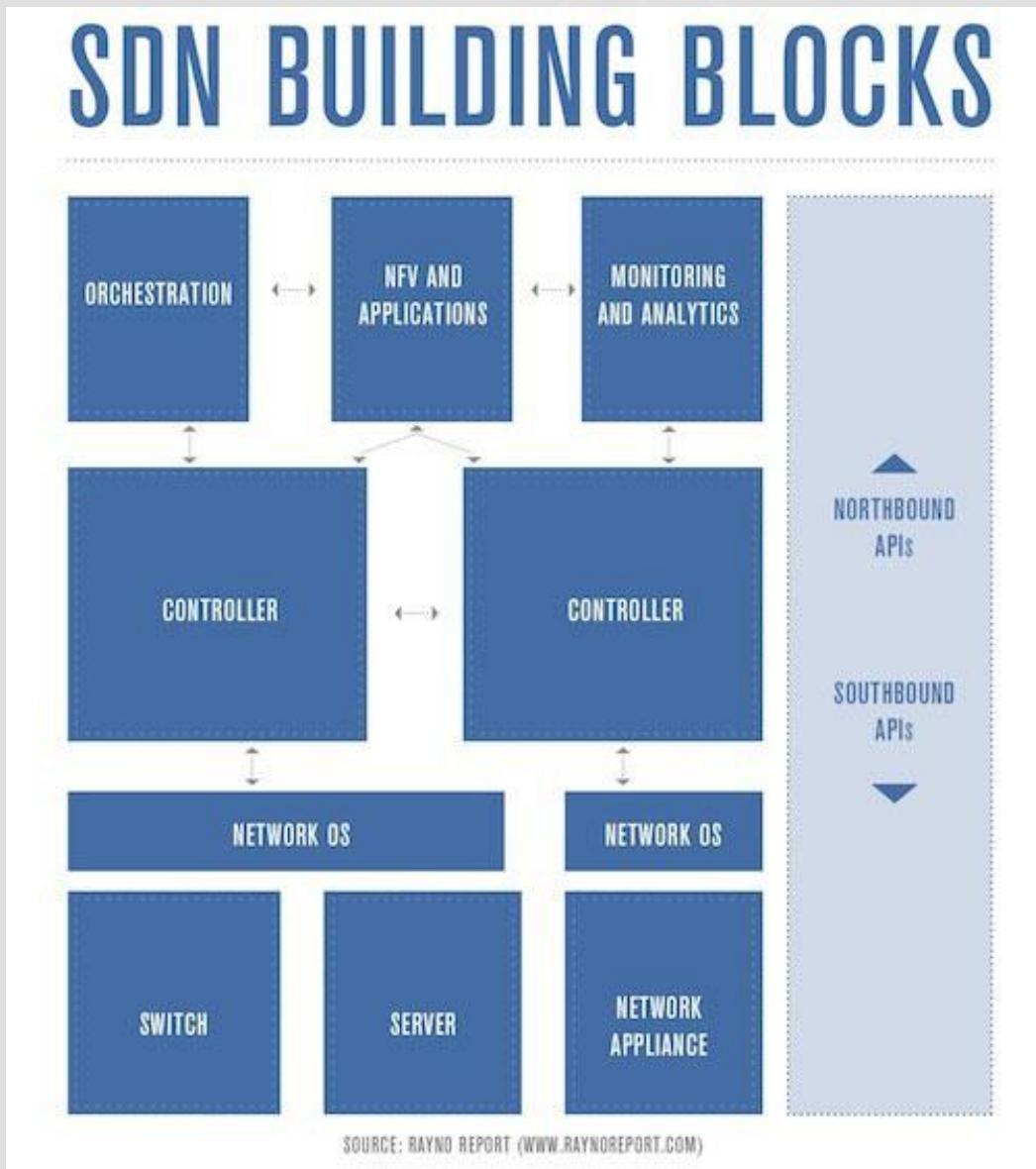
Network Programmability



- Network devices expose interface to third-party applications
- Applications provide the value
 - Networks applications:
 - Routing
 - Traffic Engineering
 - Flow Management
 - Network load balancing
 - End-user and service provider applications:
 - Access control and filtering
 - Computing resource load balancing
- Standards provide uniformity across vendor platforms



The Building Blocks



What is SDN good for

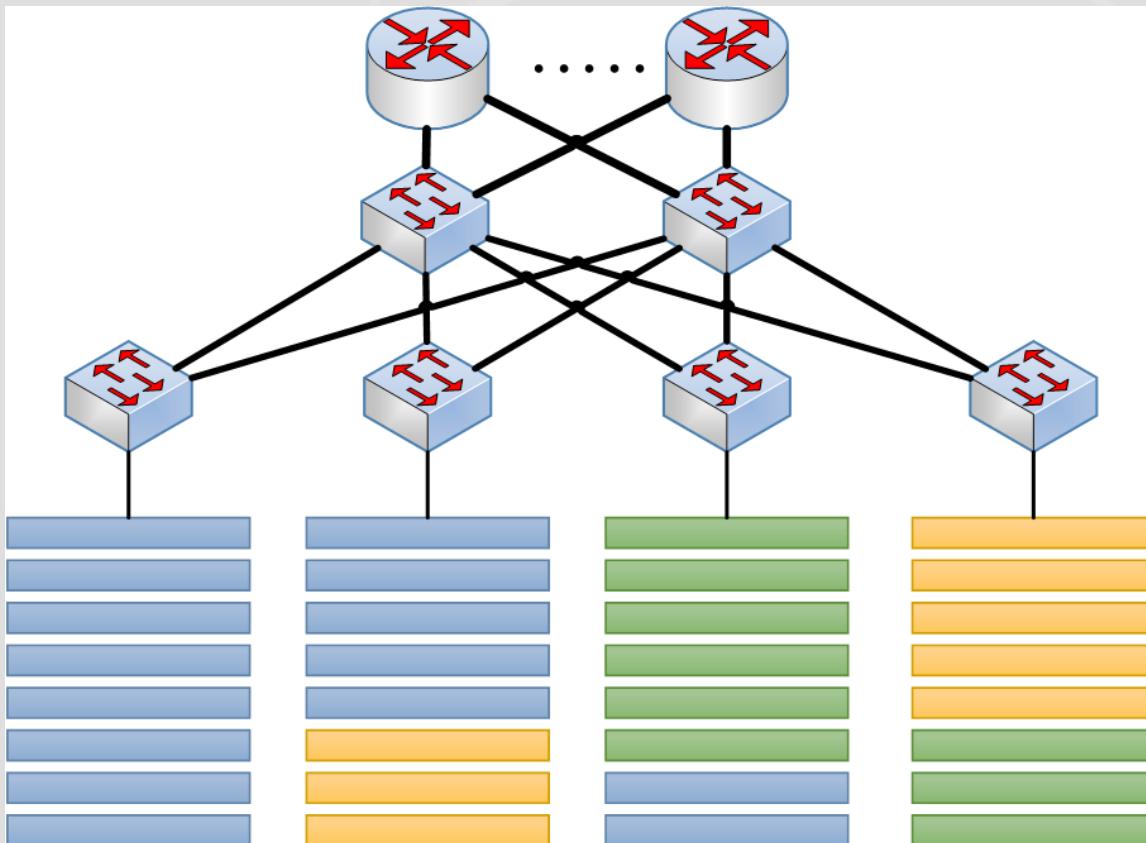


- SDNs are used for
 - Network virtualization
 - Scalability
 - Robustness
 - Security
 - Logical separation (multi-tenant environments)
 - Centralization of management
 - Simplify operational aspects and workload
 - R&D
 - Fast development and deployment of new or non-IP protocols
- SDNs are/can/will be used in
 - Data center networks
 - Cloud systems (intra-/inter-site)
 - WANs
 - Transport networks



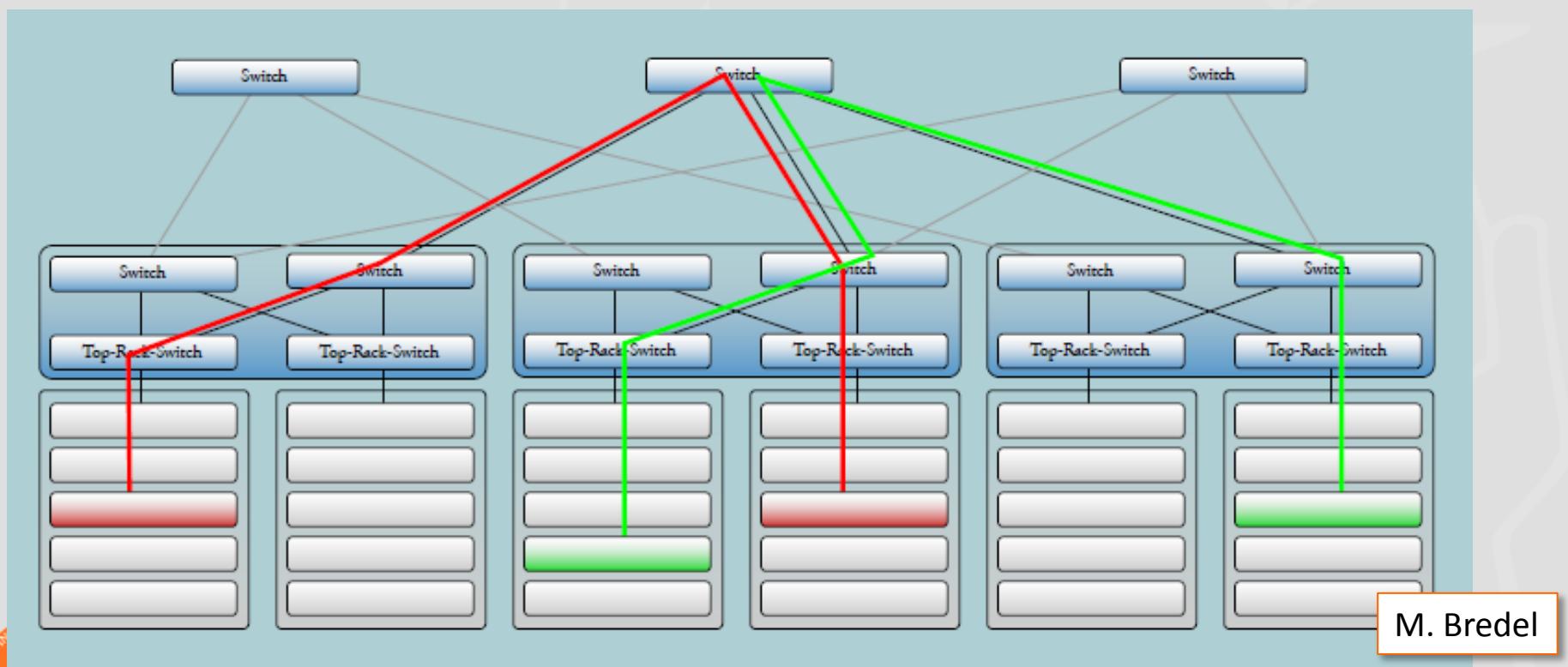
Example: Multi-Tenant Datacenter

- Some challenges in multi-tenant large data centers are
 - scalability
 - change management in large/complex deployments
 - elasticity, fast
 - ...



Data Center Example

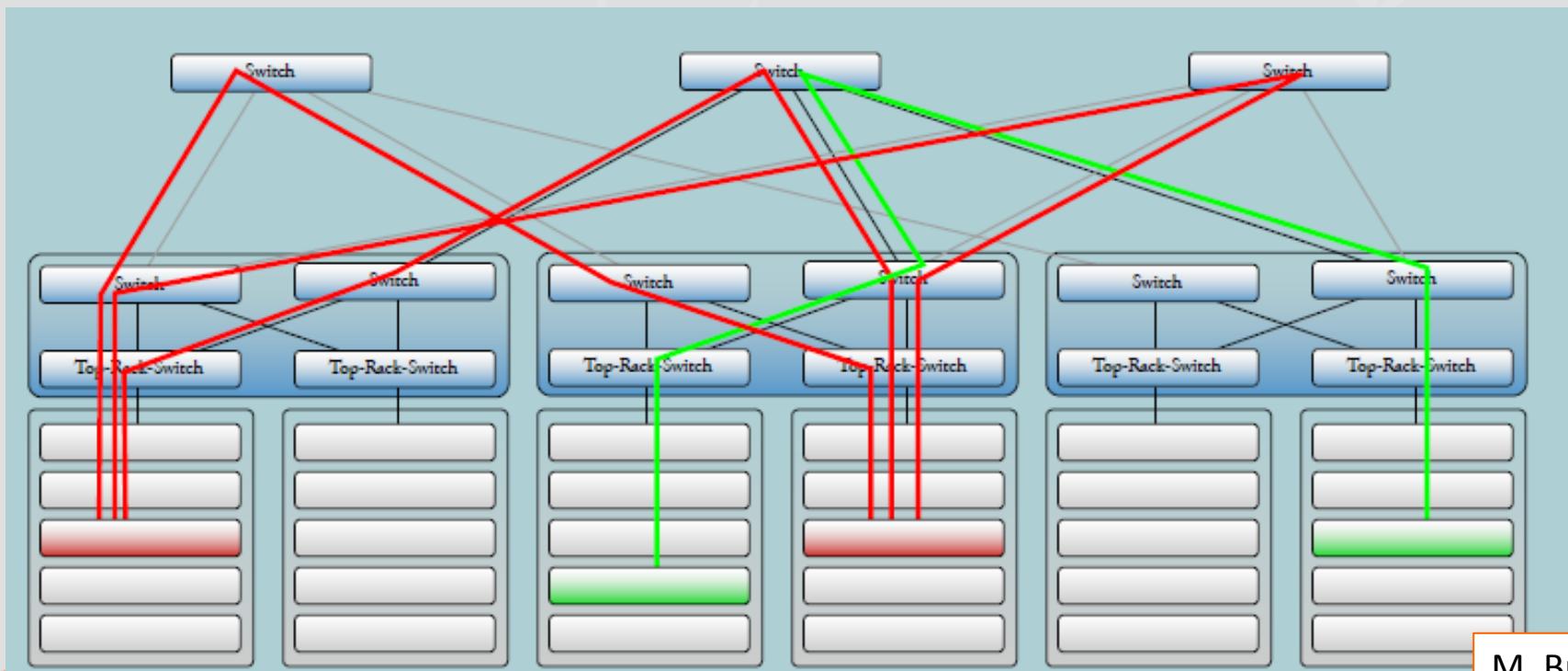
- Current techniques are limiting performance:
 - Spanning Tree for loop avoidance
 - LAGs are link-local
 - scaling up involves much configuration work on each involved device



M. Bredel

Multipath in Data Center

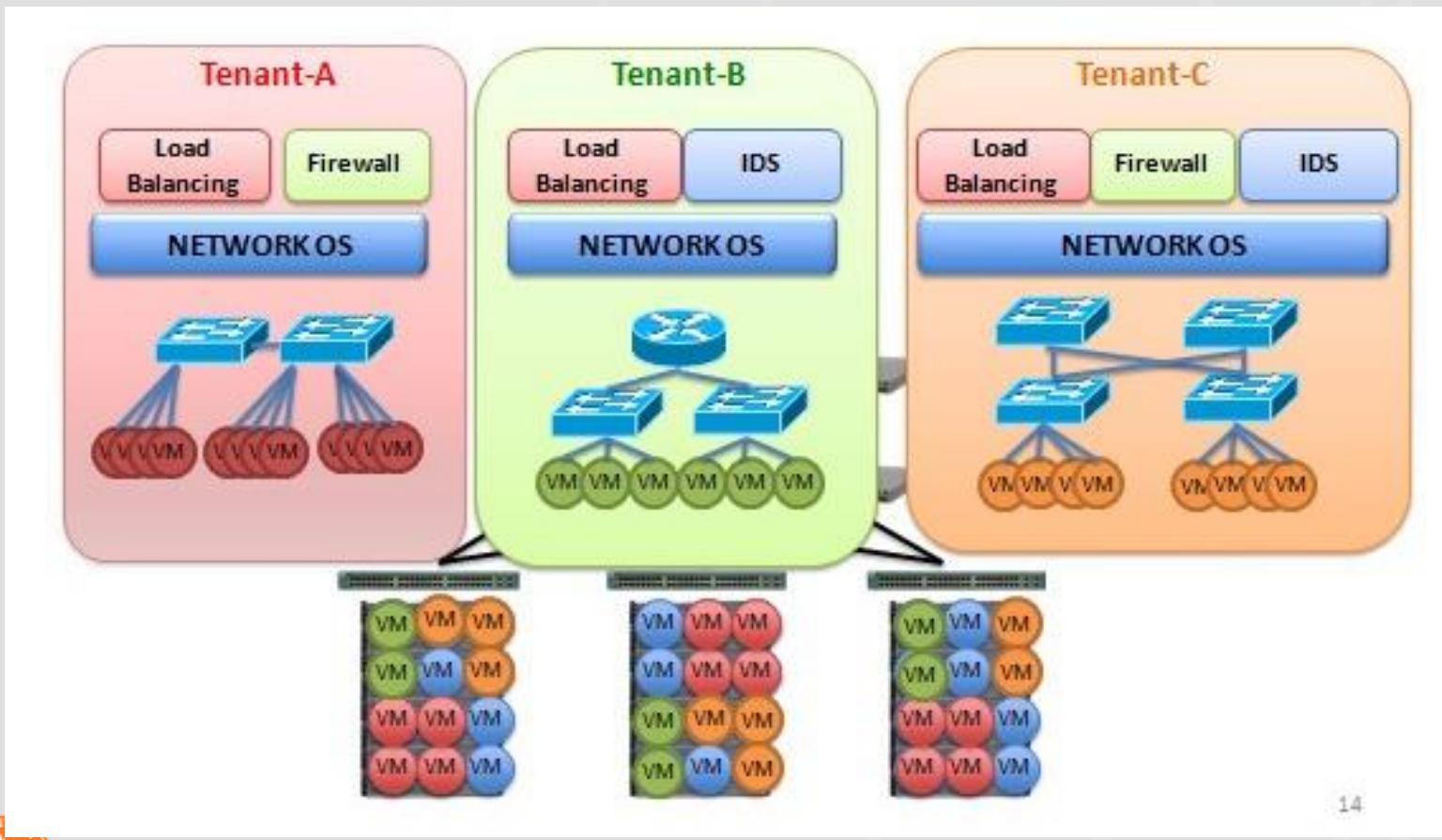
- Multipath can be achieved in several ways, e.g.
 - Multipath-TCP (IETF RFC 6824)
 - TRILL (IETF RFC 6325)
 - SPB (IEEE 802.1aq)
 - **And/Or Load Balancing algorithms in SDN!**



M. Bredel

Example: Multi-Tenant Datacenter

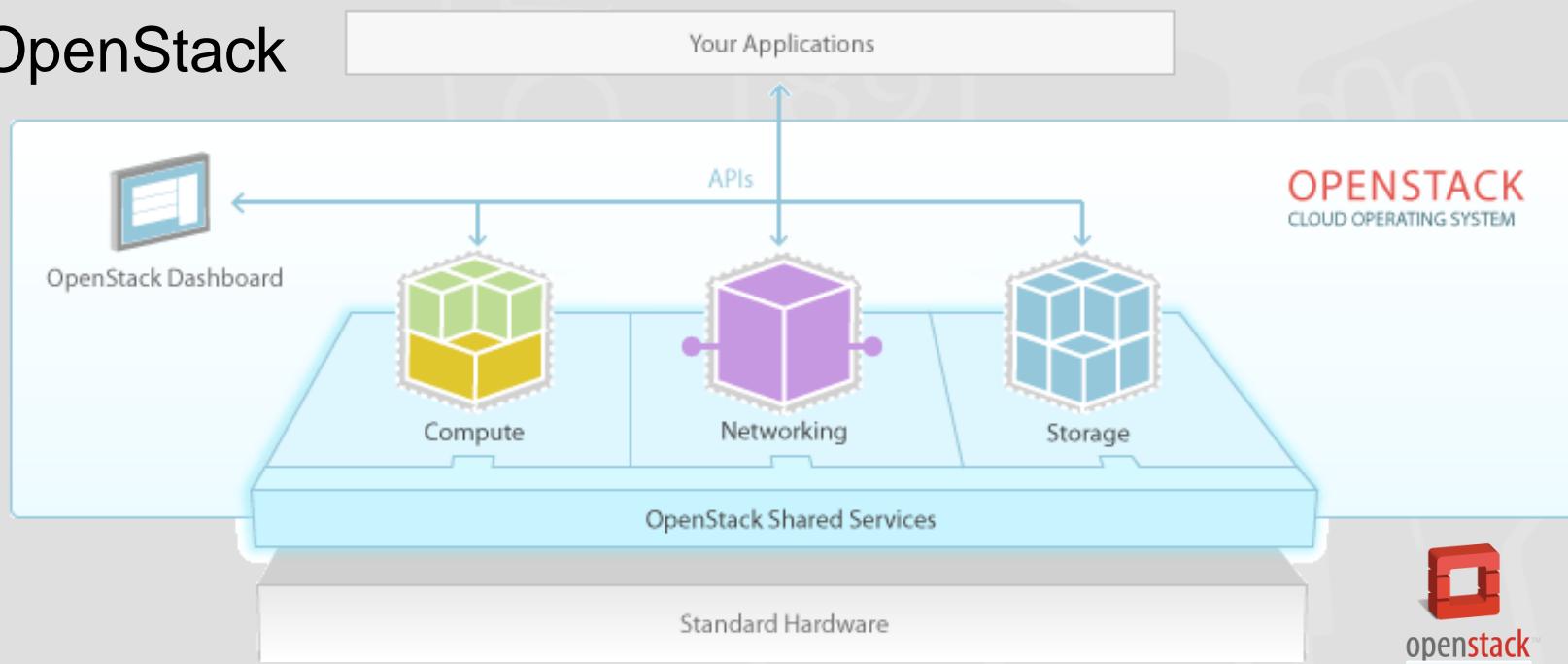
- In addition, virtualization enables
 - host sharing
 - client-specific topologies



14

Example: Orchestration

- (Wikipedia: “...automated arrangement, coordination, and management of complex computer systems, middleware, and services”)
- For full service deployment need to orchestrate Storage, Compute and Network resources
- E.g. OpenStack



OpenFlow – SDN's favourite protocol

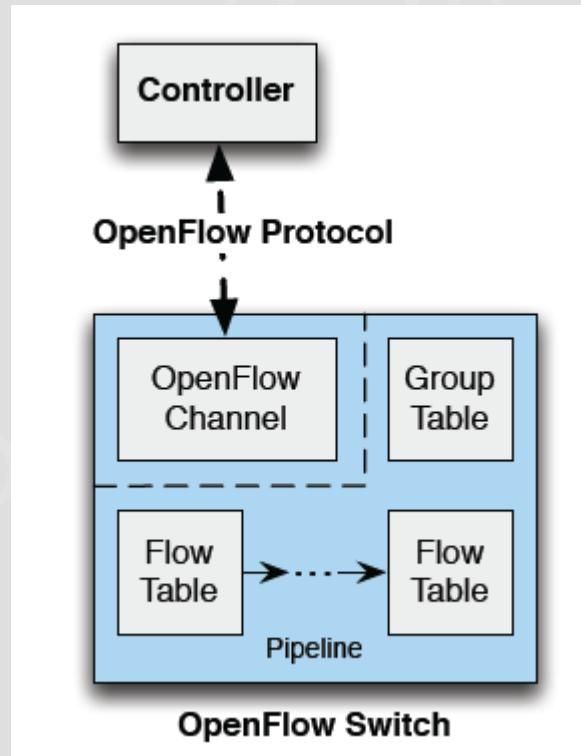


- OpenFlow ≠ SDN
 - SDN is (technically, and with limitations) possible with SNMP, CLI, etc.
- OpenFlow = open standard for
 - Protocol for controller – device communication
 - Definition of packet processing in the switch
- Standardized by the Open Networking Foundation



OpenFlow switch components

- For packet look-up and forwarding
 - Flow Tables
 - Group Tables
- Control Channel
 - add, update, remove flow table entries
- OpenFlow Switch Ports:
 - Physical
 - Logical
 - e.g. LAG, tunnel, etc.
 - Reserved
 - ALL, CONTROLLER, TABLE, etc.



ONF <http://www.opennetworking.org>

Flow Tables

| Match fields | Priority | Counters | Instructions | Timeouts | Cookie | Flags |
|--------------|----------|----------|--------------|-----------|------------|-------|
| MAC src | MAC dst | IP src | IP dst | TCP dport | ... | Count |
| * | 50:25:.. | * | * | * | | 531 |
| * | * | * | 1.2.3.* | 80 | | 77 |
| * | * | * | * | * | * | 2755 |
| | | | | | Controller | * |

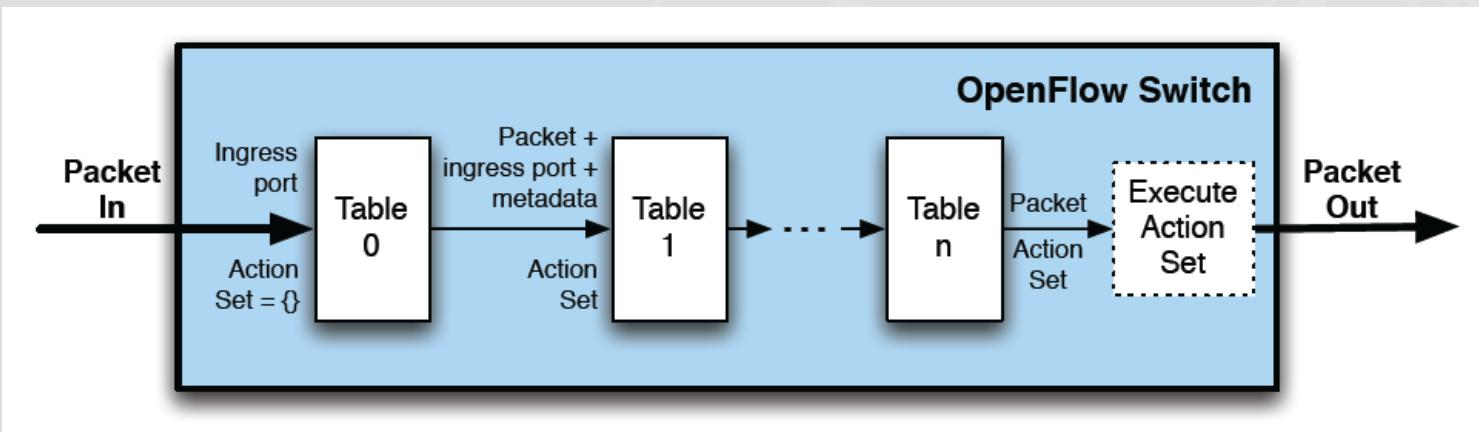
- Matching fields at Layer 2, 2.5, 3 and 4
- Wildcards allowed
- Table miss entry - default action:
 - forward to controller, port or drop (default)

It is what the controller writes into the flow tables that determines the network behaviour



Table Pipeline

- Tables are processed in a pipeline



- For each table:
 - Find highest priority matching flow entry
 - Apply Instructions
 - apply actions
 - update action set
 - update metadata
 - Send match data and action set to next table

Instructions and Actions



- (Some) Instructions:
 - Apply actions <actions>
 - Clear actions <actions>
 - Write actions <actions>
 - Meter <actions>
 - Goto <table>
- (Some) Actions:
 - Output <port>
 - Drop
 - Push tag
 - Pop tag
- Tags specified (v1.3) can be
 - VLAN
 - MPLS
 - PBB

ONF OpenFlow Standard v1.3

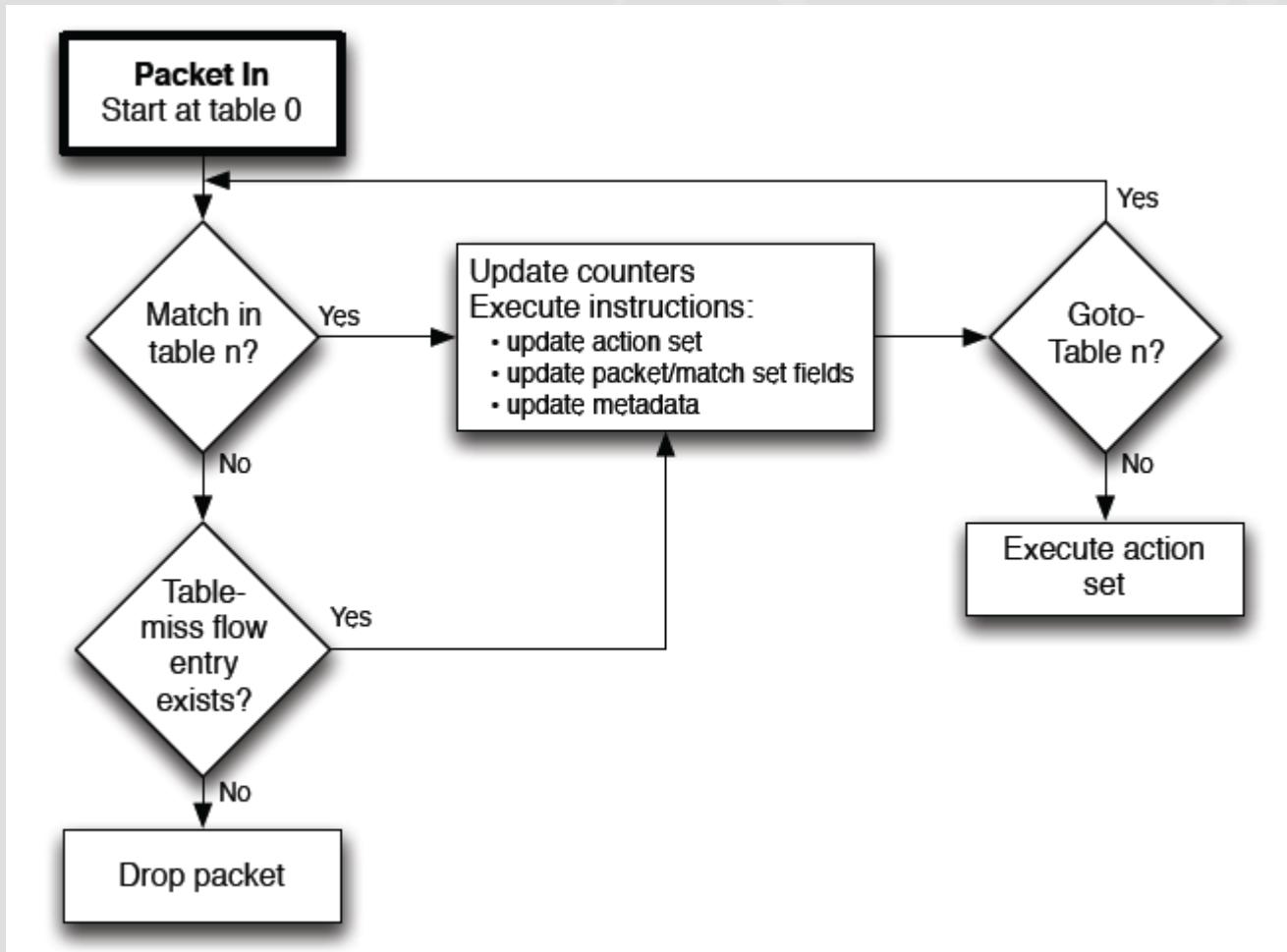
For full document, see <http://www.opennetworking.org>



Packet Processing in OpenFlow Switch

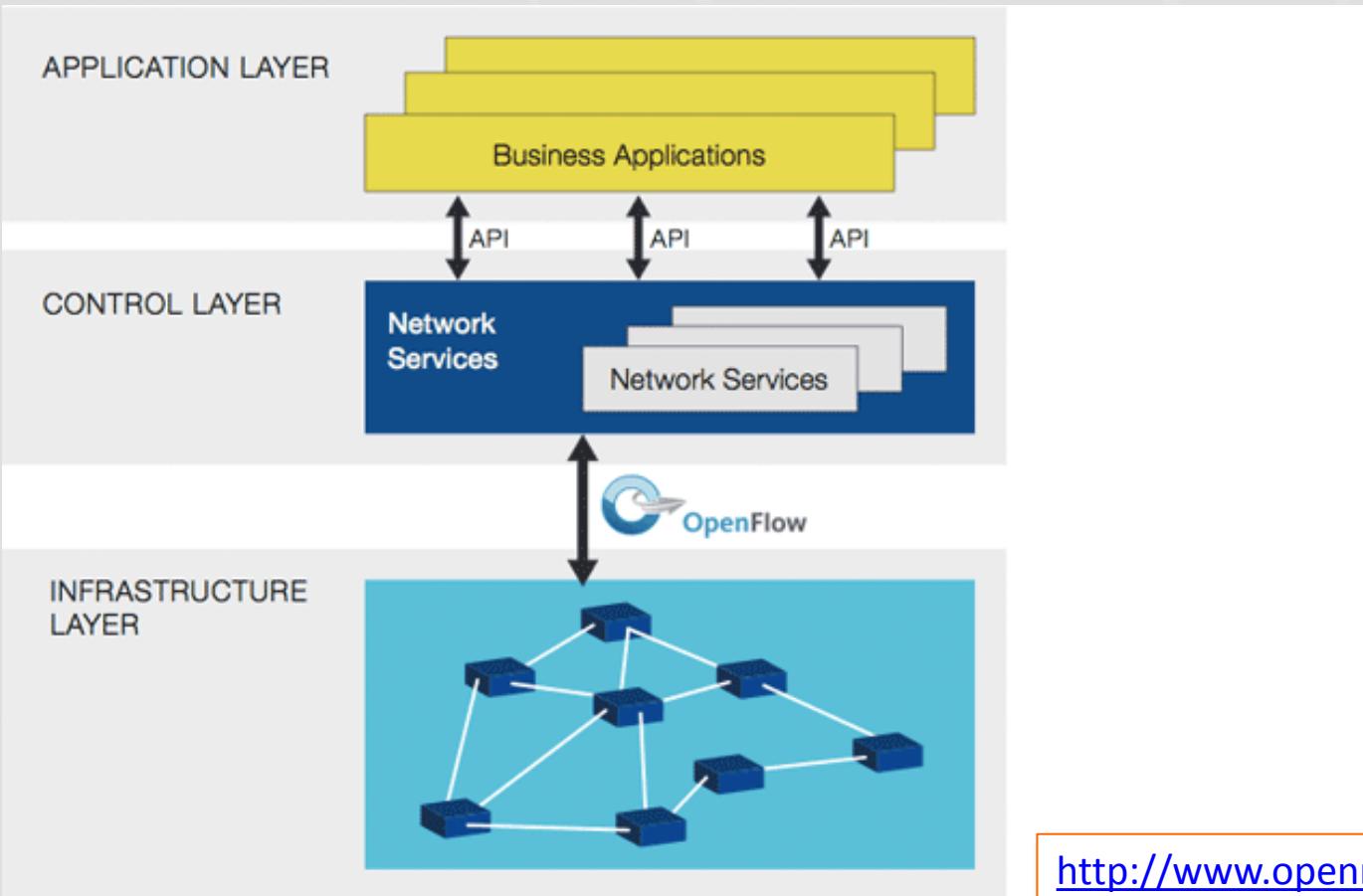
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US LHCNET



OpenFlow - The Controller

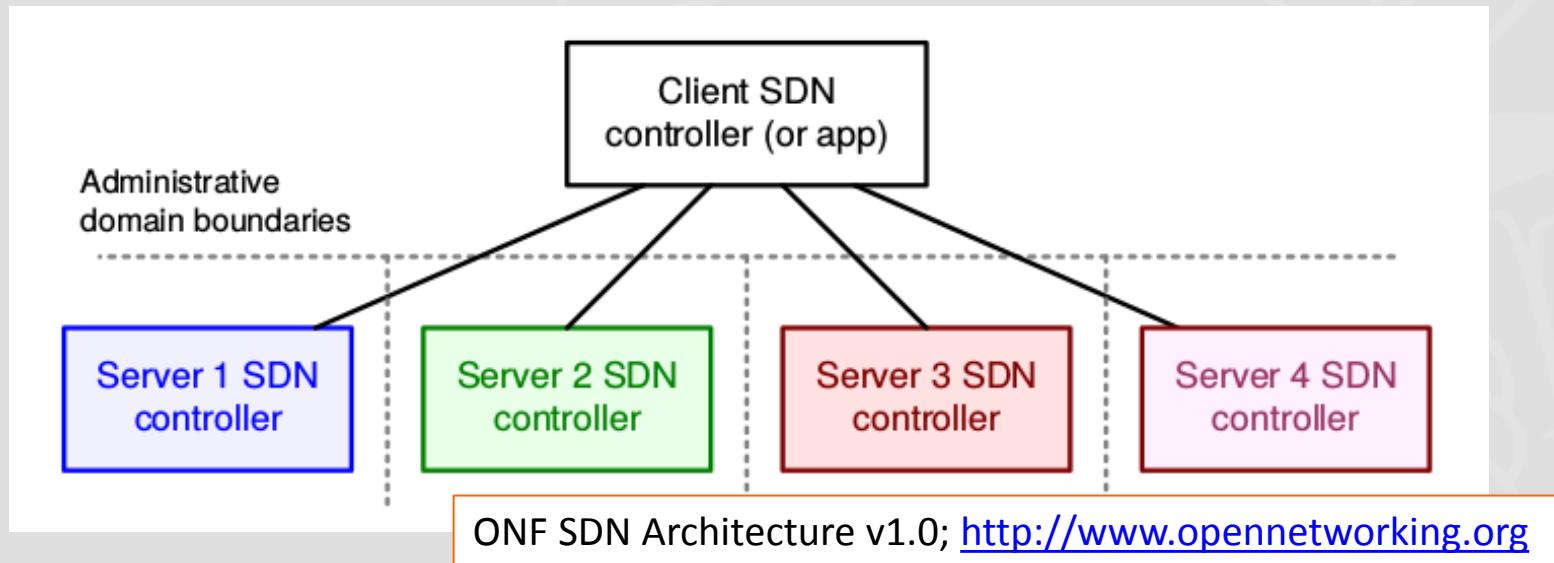
- Software typically running on commodity hardware
- Provides the API to user applications
 - Aka Northbound interface



<http://www.opennetworking.org>

Path to multi-domain SDN

- Not to forget: interactions between administrative domains



Popular Controller Examples

- NOX (C++)
 - <http://www.noxrepo.org/>
- POX (Python)
 - <http://www.noxrepo.org/>
- Ryu (Python)
 - <http://osrg.github.io/ryu/>
- Floodlight (Java)
 - <http://www.projectfloodlight.org/floodlight/>
- OpenDaylight (Java)
 - <http://www.opendaylight.org/>



OpenDaylight – Industry Driven

BROCADE

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Your Network is Information

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VERSA
NETWORKS

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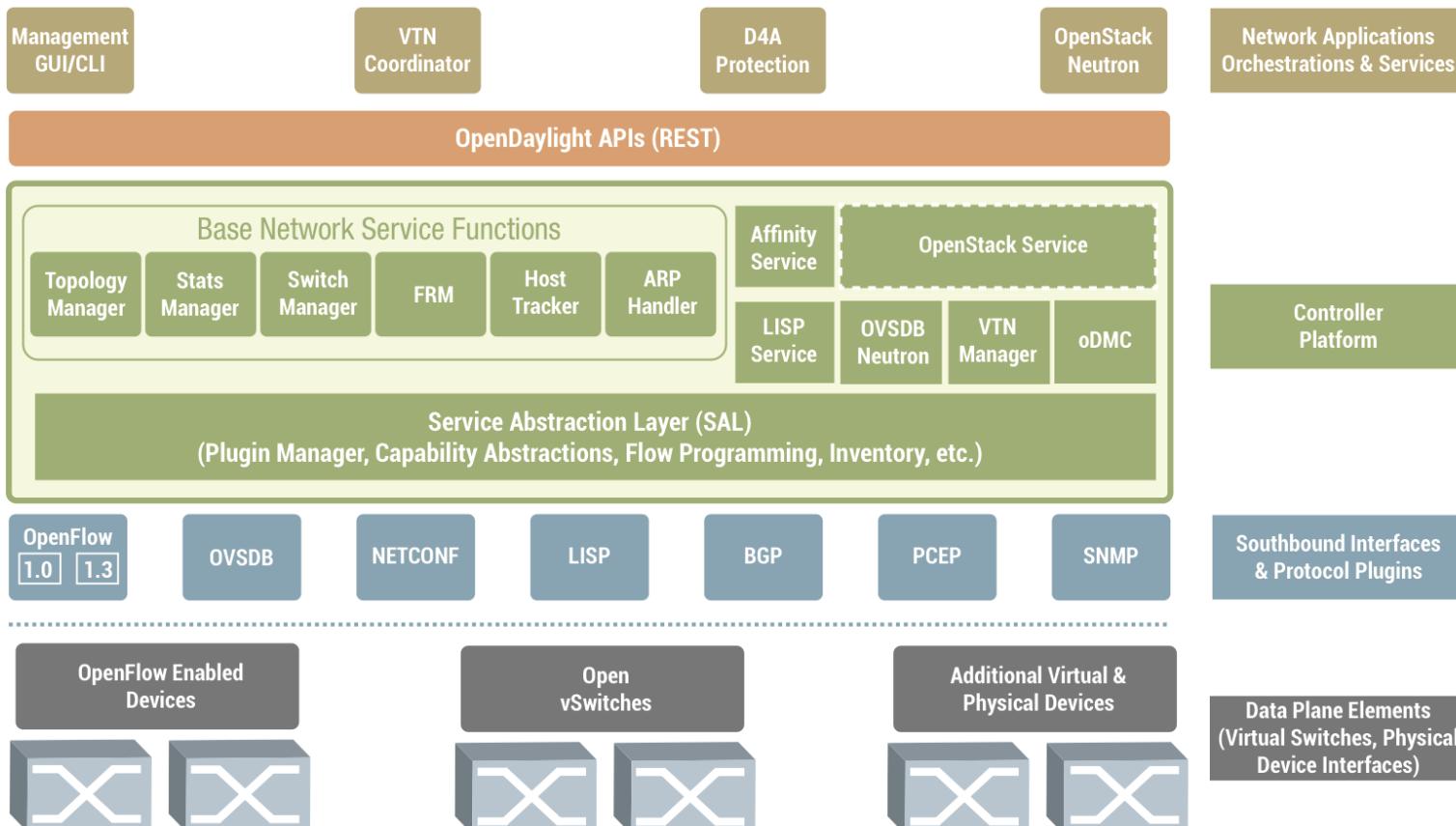
vmware®



OPEN DAYLIGHT

“HYDROGEN”

VTN: Virtual Tenant Network
oDMC: Open Dove Management Console
D4A: Defense4All Protection
LISP: Locator/Identifier Separation Protocol
OVSDB: Open vSwitch DataBase Protocol
BGP: Border Gateway Protocol
PCEP: Path Computation Element Communication Protocol
SNMP: Simple Network Management Protocol
FRM: Forwarding Rules Manager
ARP: Address Resolution Protocol



Some Important Components

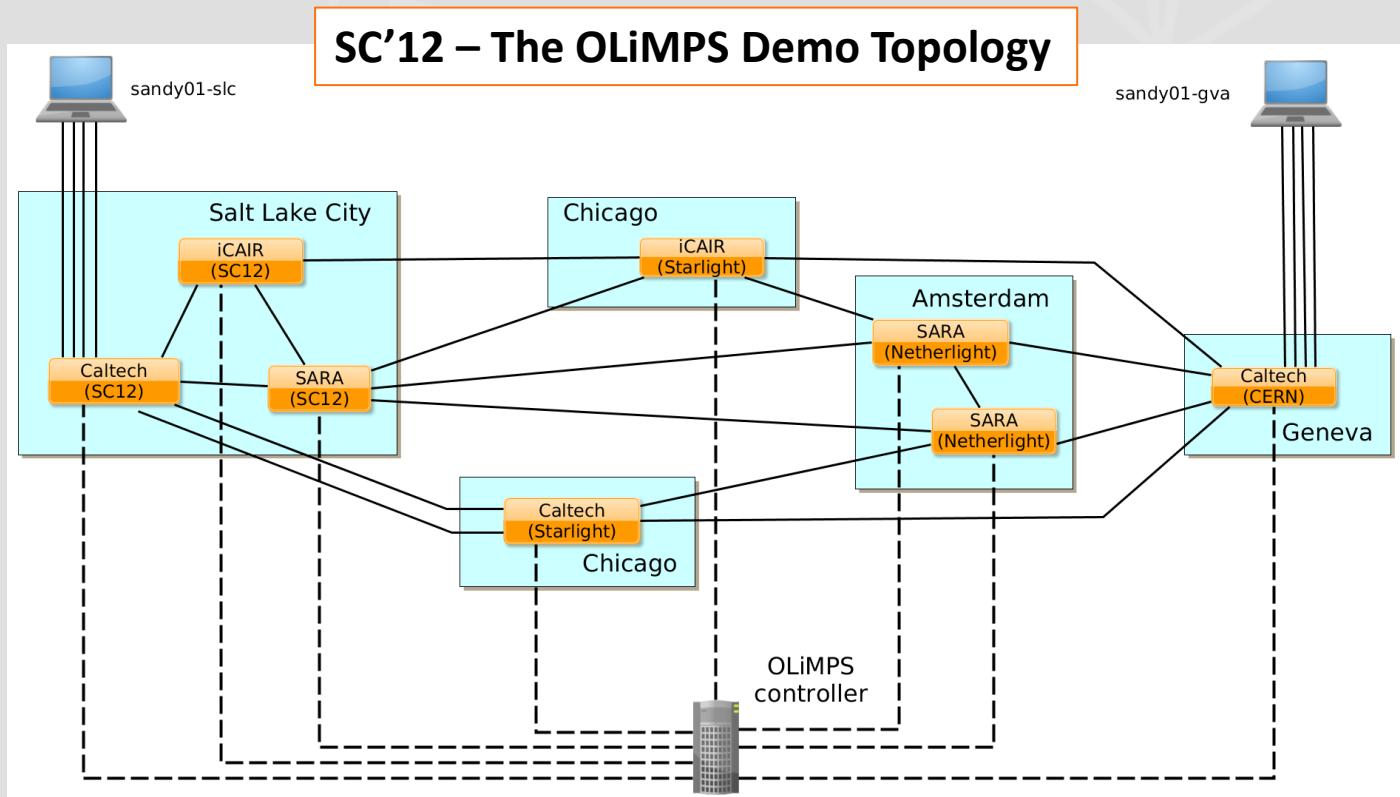


- Northbound interface: REST and OSGi
- BGP-LS (BGP-Link State)
- PCEP (Path Computation Engine Protocol)
- Southbound interface supporting OpenFlow and non-OpenFlow devices



Application Example: Multipath Controller

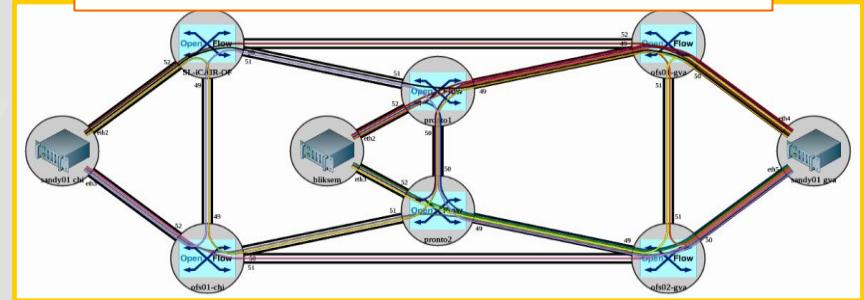
- OpenFlow Link-layer Multi-Path Switching, OLiMPS
- DOE funded project
- Extending capabilities of the Floodlight controller
- Load-balancing traffic over multiple possible end-to-end paths



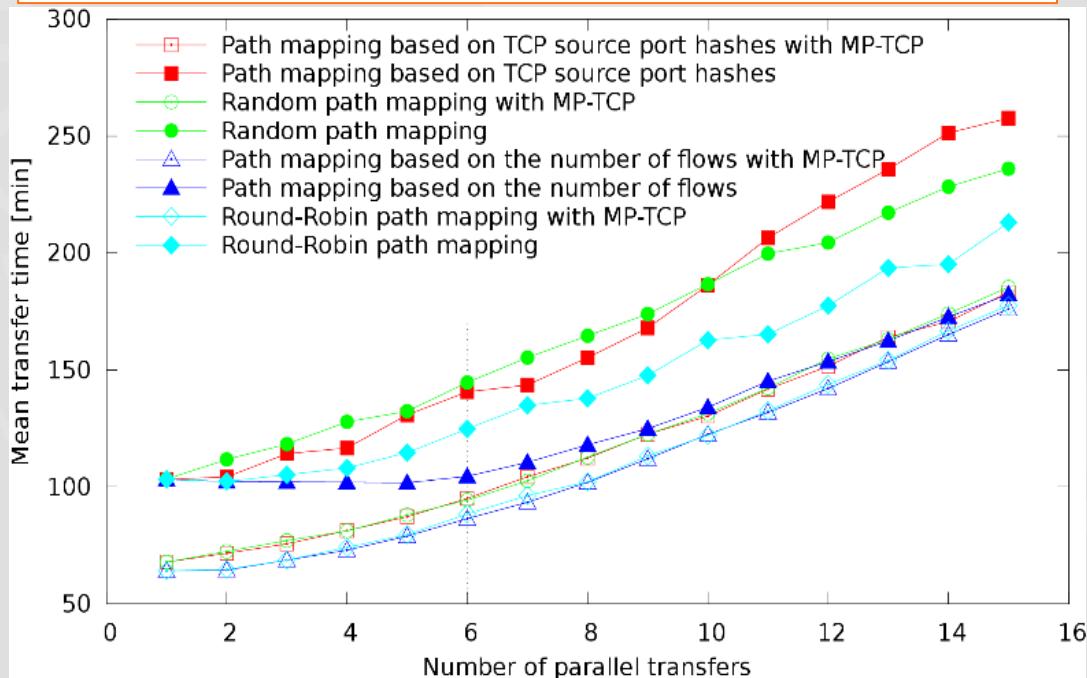
Meshed Networks, Multipath

- Data connections (TCP) are point-to-point
- Classical IP routing constrains flows to a single path
- Reality: Networks are meshed, many paths possible, only one used
- Multipath forwarding helps increasing network efficiency
- Application “telling” the network controller its intentions increases efficiency even further
- Implemented using Floodlight controller
- Paper to be presented at HotSDN 2104

Example of WAN demo topology



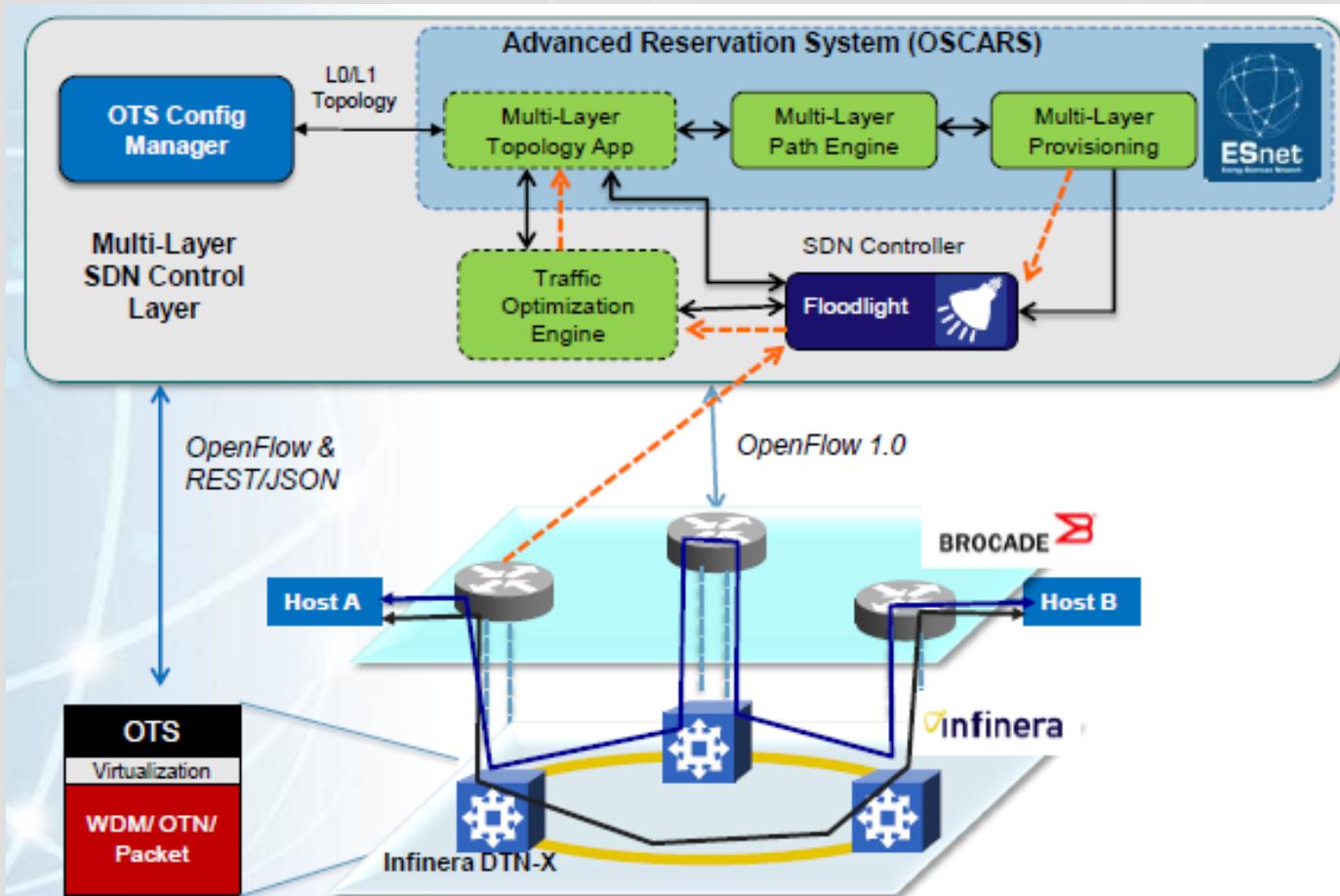
From OLiMPS project – multipath with openflow



Local testbed with Application-Network interface

SDN + Dynamic Circuits I

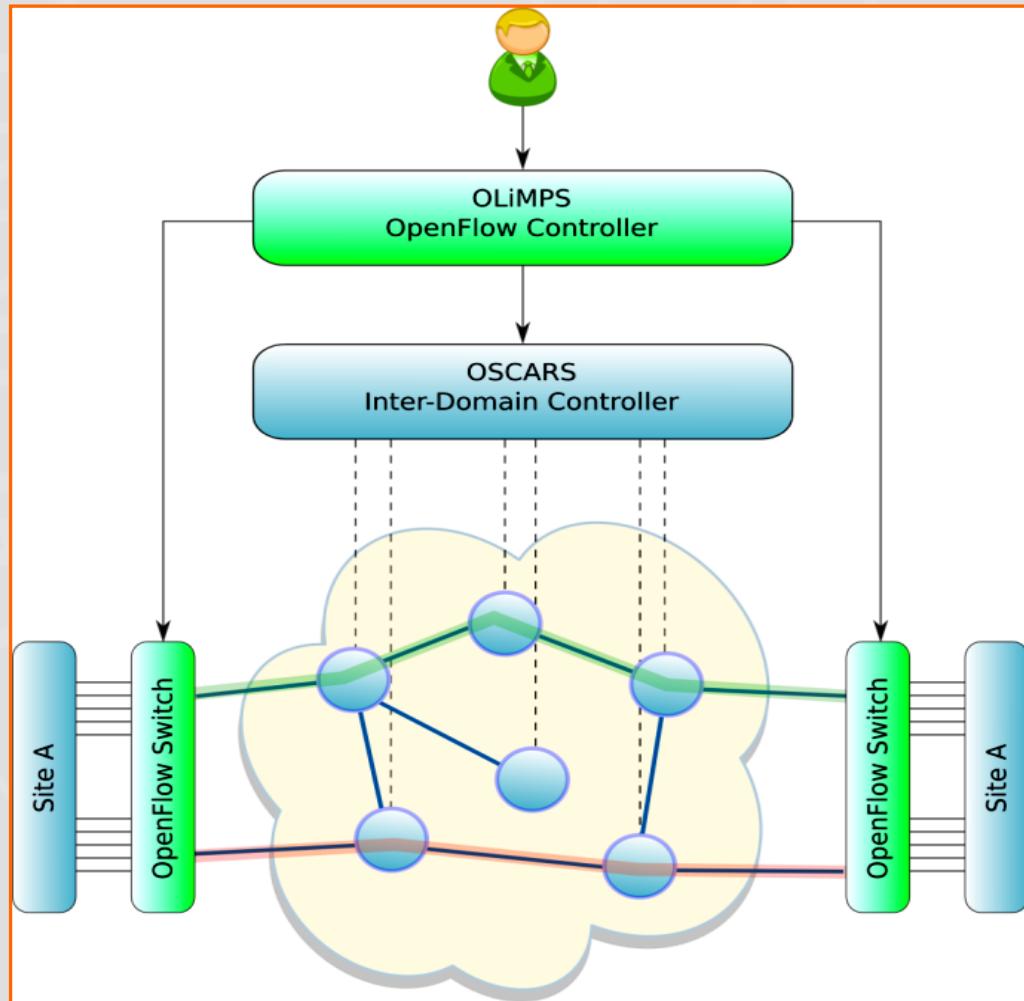
- ESnet's OSCARS management system using OpenFlow controller for traffic optimization



Demo: <http://www.sdncentral.com/events/brocade-infinera-esnet-sdn-demo/>

SDN + Dynamic Circuits II

- Caltech's OLiMPS project created an interface between Floodlight OpenFlow controller and the OSCARS dynamic circuit system
- Additional capability of the controller:
Create additional paths between OpenFlow devices
- I.e. create a topology optimized to the load distribution in the network
- Fits OpenDaylight architecture



SDN Summary



- SDN provides a new possibility for programmatic network interaction
- HEP computing should be involved in defining services provided by the networks, built on SDN



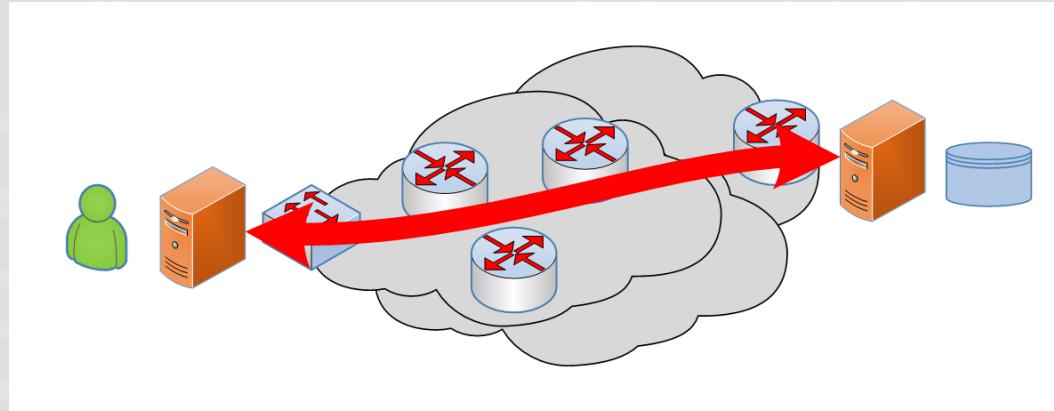
CONTENT CENTRIC NETWORKING

Where we meet CCN, NDN and friends

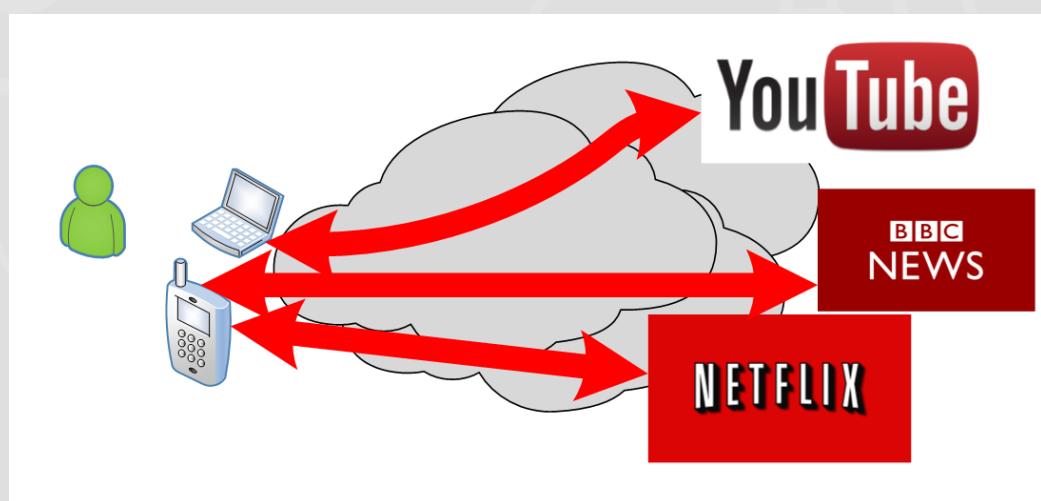


CCN Background

- When Internet was invented, it was connecting resources
 - TCP/IP: point-to-point connections between two entities
 - IP: delivering packets to destination **hosts**



- Today's applications, ours included, care about **content**



CCN Background, cont.



- Applications deal with “what”, while the network deals with “where”
 - Lots of middleware needed to match these
 - Web services, CDNs, P2P, ...
- Complexity arises when dealing with failover, security, etc.
 - E.g. if the server at A.B.C.D does not respond, it’s the application to react and possibly find a backup source for the data
 - E.g. you trust the server, but it’s the content that’s potentially dangerous
- Lot of the work in CDNs, redirection, caching deals with this mismatch
- Can we do a better design instead?
- Identify data rather than hosts?



The Future Internet



- CCN is one of the **Future Internet Architectures** being developed and studied
- Specific projects include
 - Content-Centric Networking – CCN
 - Project at PARC
 - Code base developed: CCNx
 - Named Data Networking – NDN
 - NSF funded project since 2010, recently extended
 - Collaboration including PARC
 - and several other similar projects
- I will focus on the Named Data Networking (NDN) project in the following



Named Data Networking

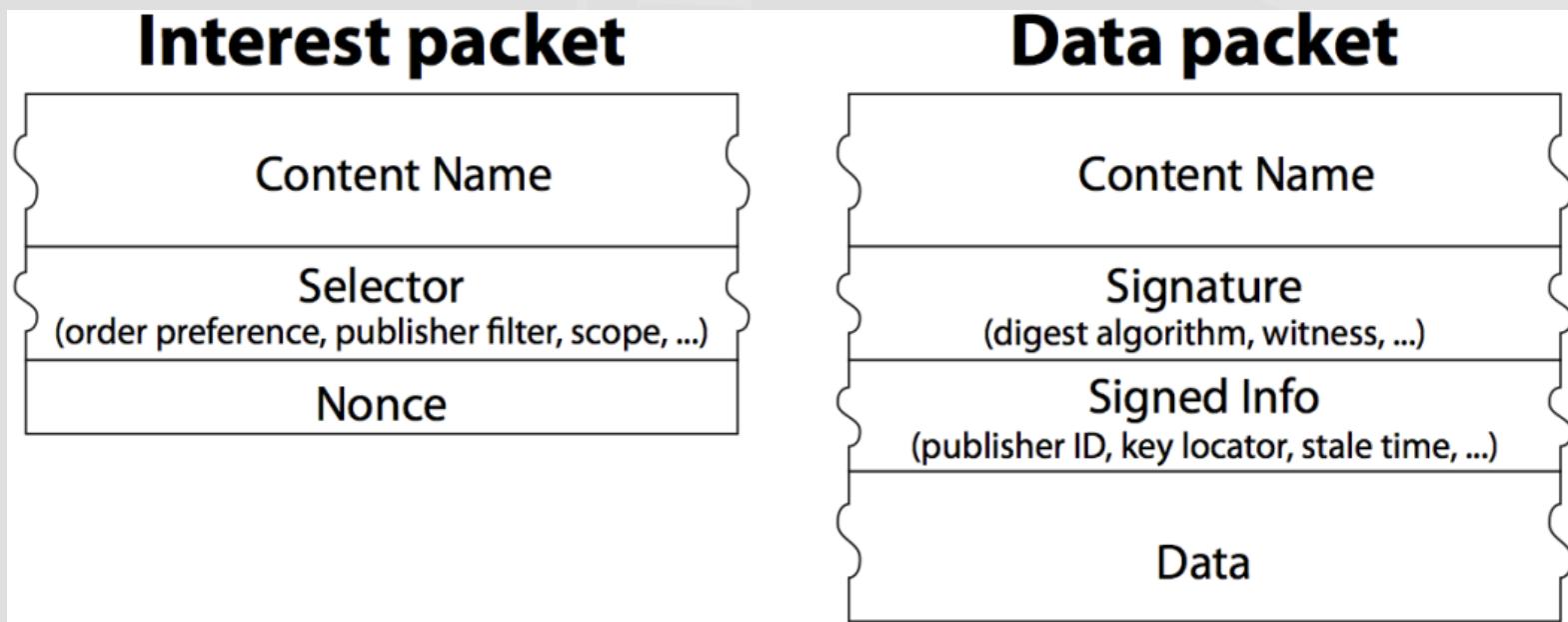


- Basic Principle: Name Data instead of naming end-hosts
- Today's Internet delivers packets to a destination address
- NDN delivers content identified by a given name to the client
- This is a basic change in semantics of the network service



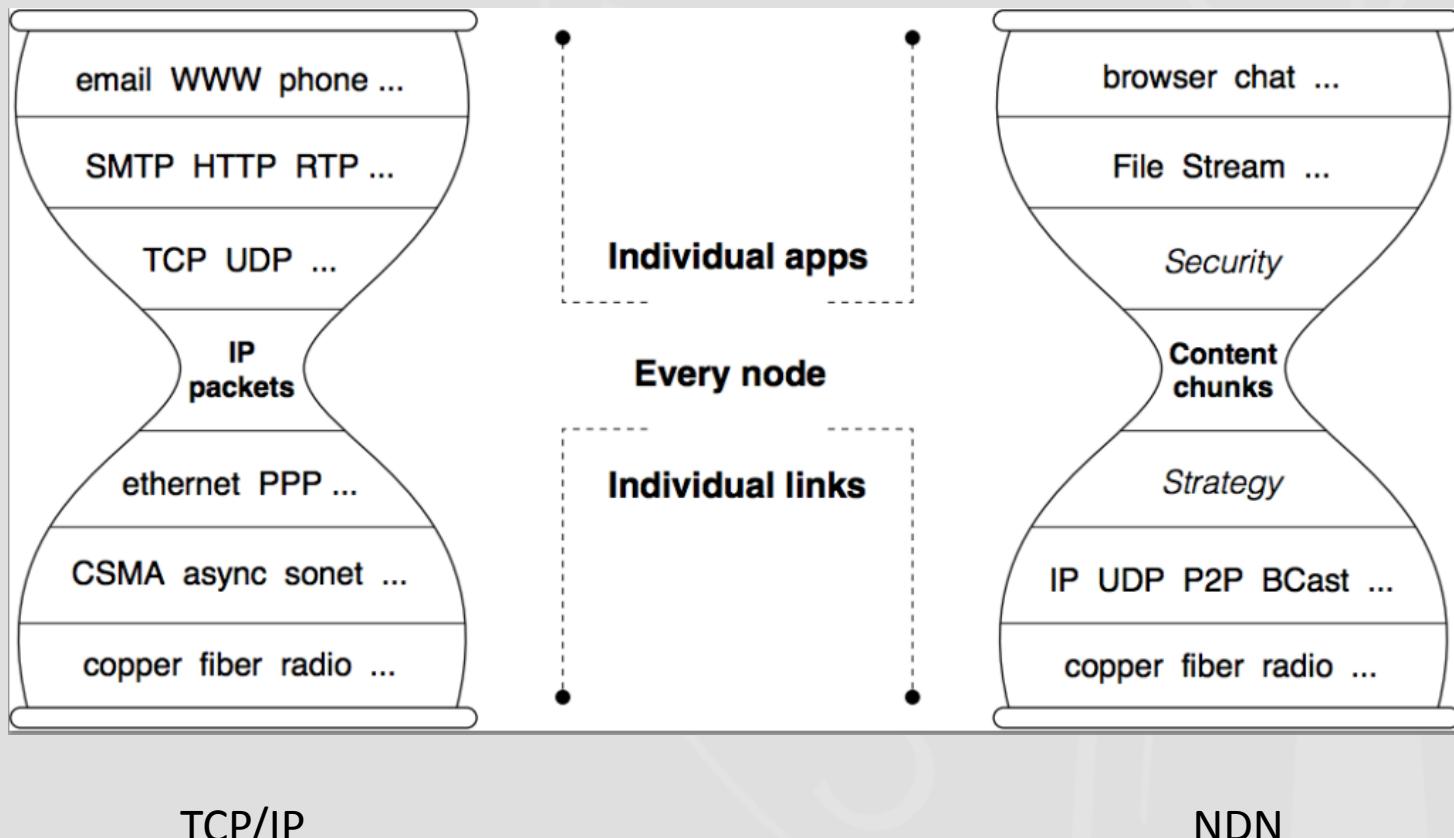
Two NDN packet types

- Communication is driven by the receiving end
- **Interest Packets:**
 - Sent out by the data consumer, identifies desired data
- **Data Packets:**
 - Sent back by the node which has the desired content



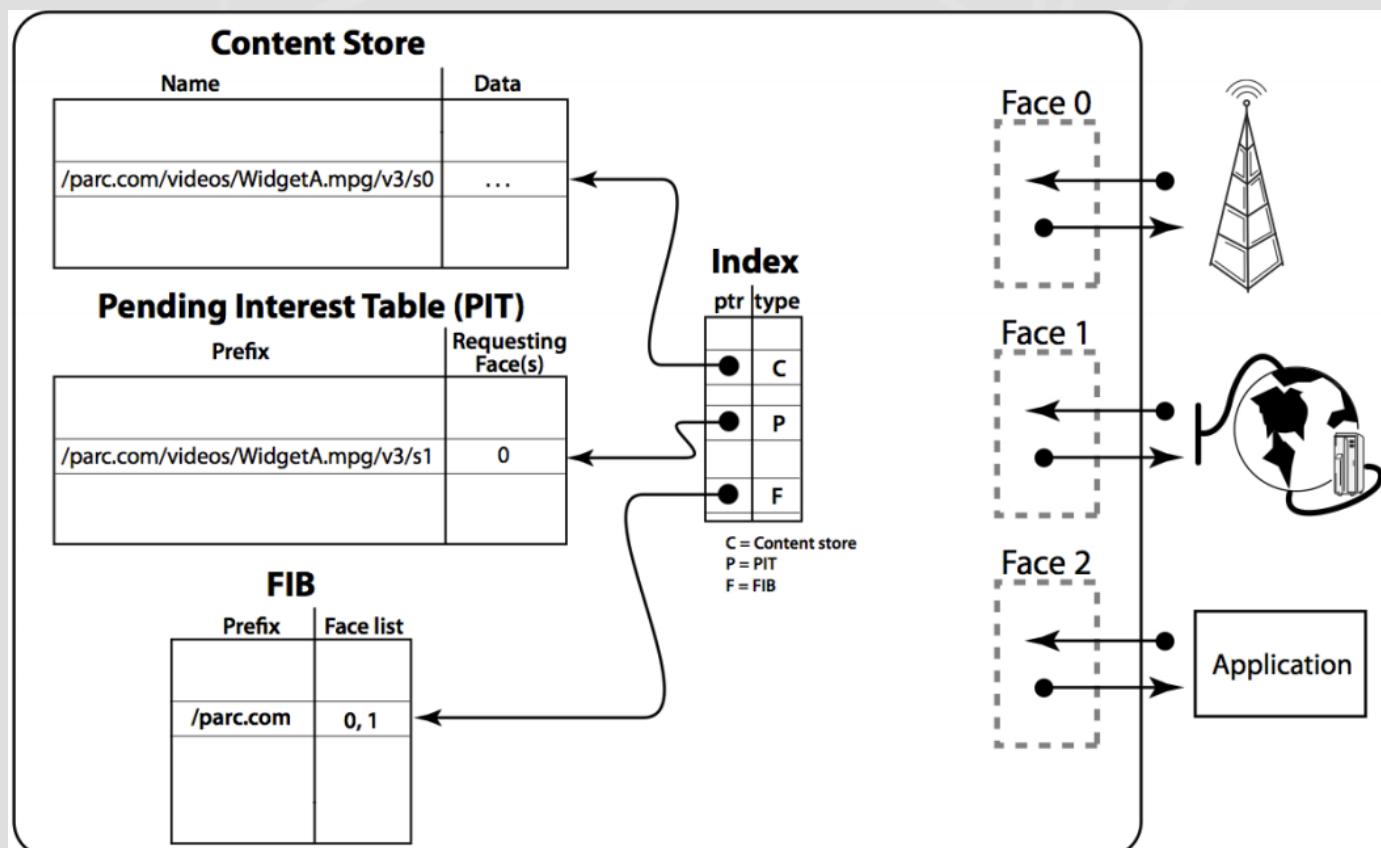
Protocols

- Basis for NDN:
 - Named data replaces named end-points
 - Keeping the thin waist approach



Network node operation

- NDN routers remember the interface the request came in the PIT
- Forwards the Interest Packet looking up the name in the FIB
 - Populated by routing protocol
- Once the Interest Packet reaches a node that has the content, a Data Packet is sent back following the reverse path (as stored in the PIT)



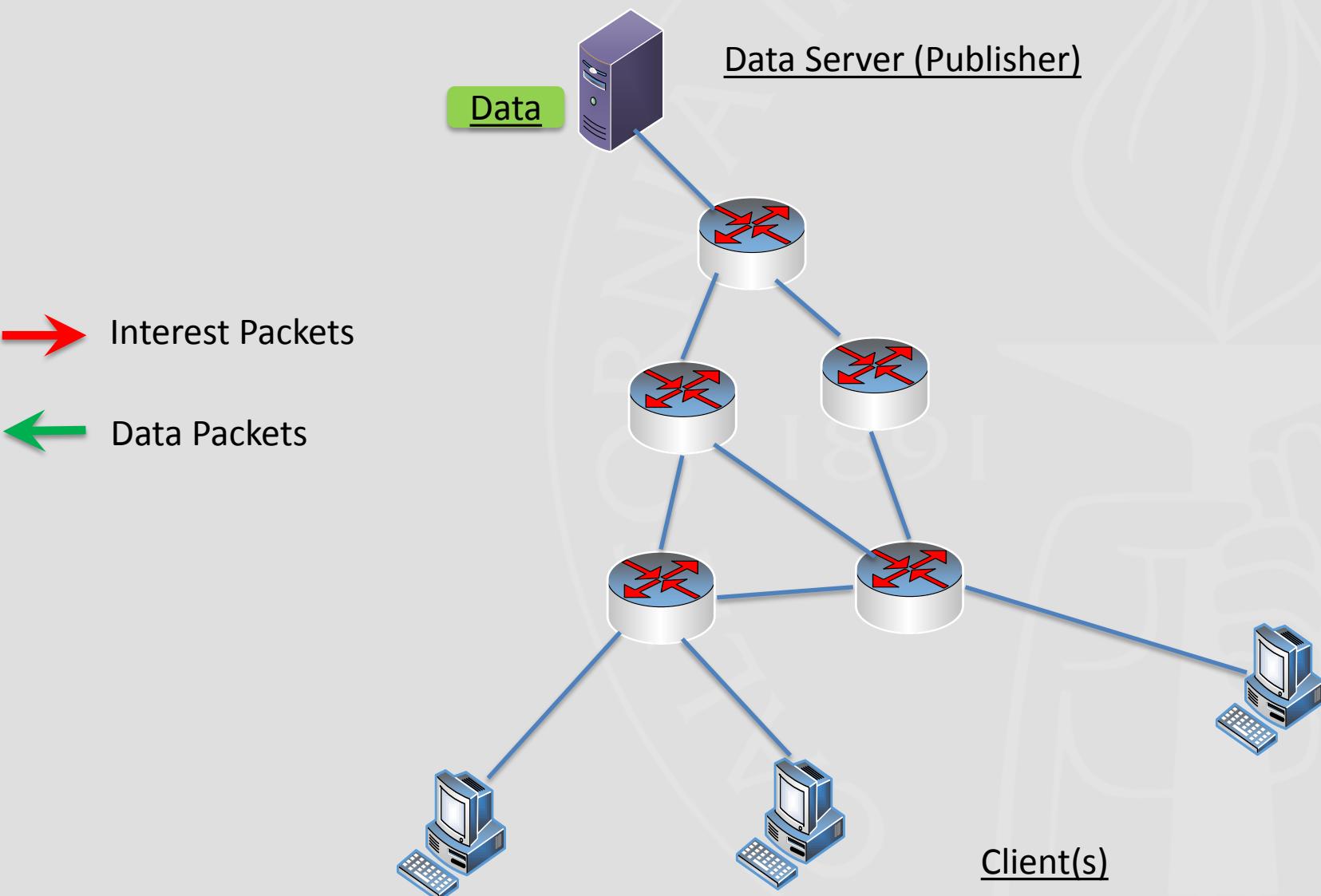
Better-than Multicast



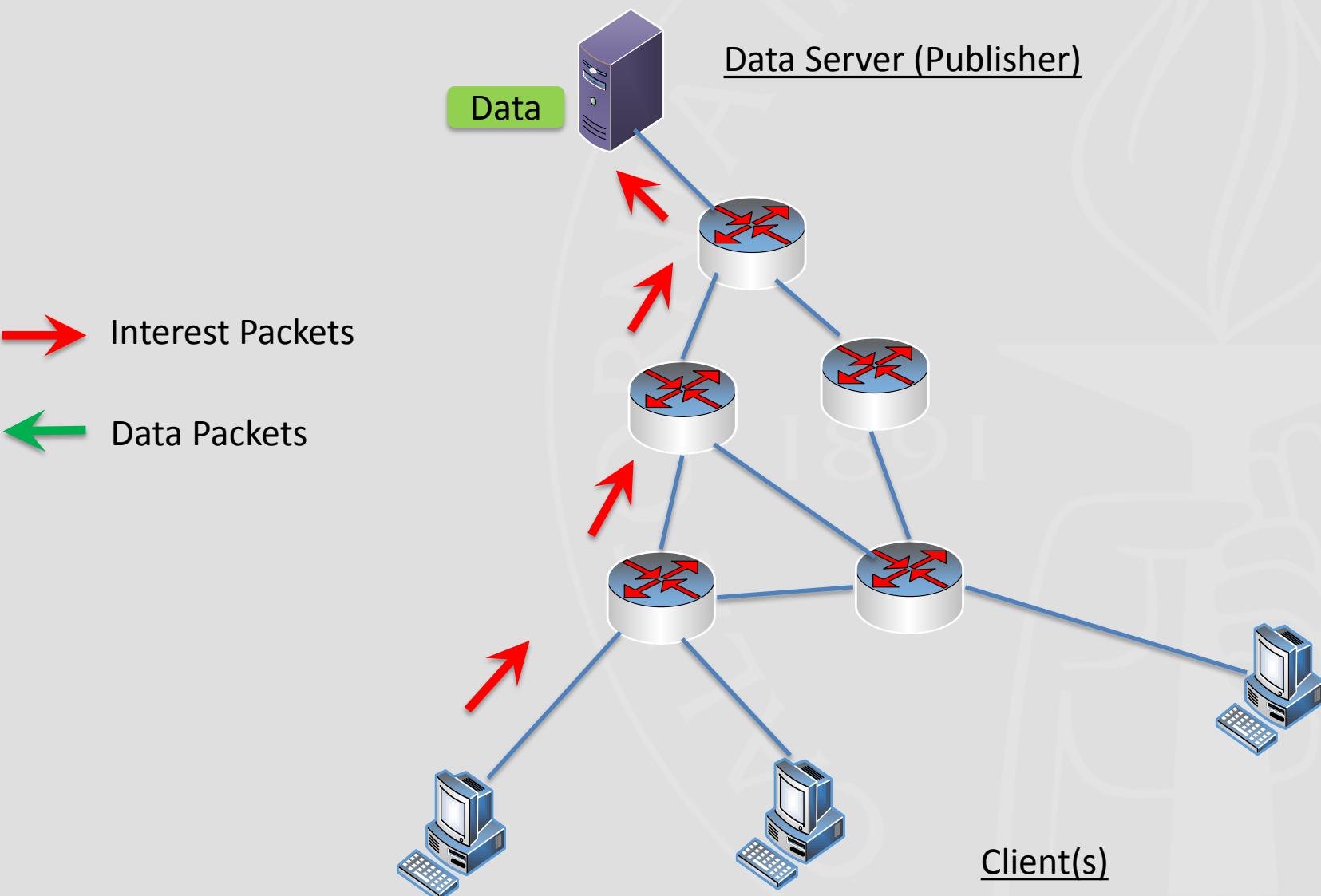
- Data Packets are cached in the routers' Content Store
- Data Packets are then forwarded to all interfaces with registered interest in the router's PIT
 - i.e. if multiple IPs received for same content – multicast!
- When next interest packet arrives for a named data in the content store (cache), a Data Packet is sent from the router, rather than forwarding the IP to the data source
- This provides for additional multicast-like operation
- With one big difference: no multicast request or protocol is necessary
- Added benefit: because of caching, destinations do not have to be synchronised - fits a pull model as opposed to push



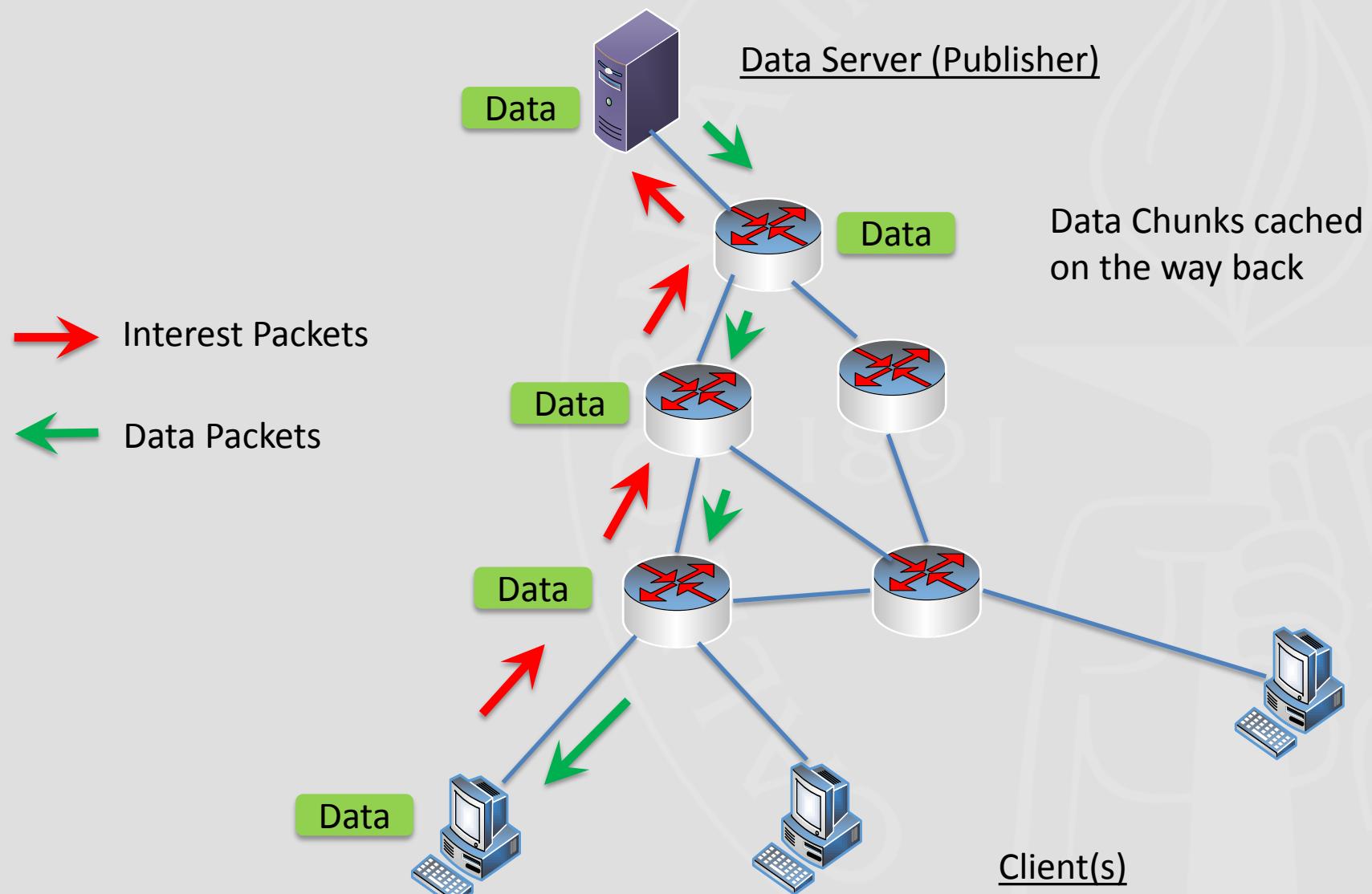
NDN Operation



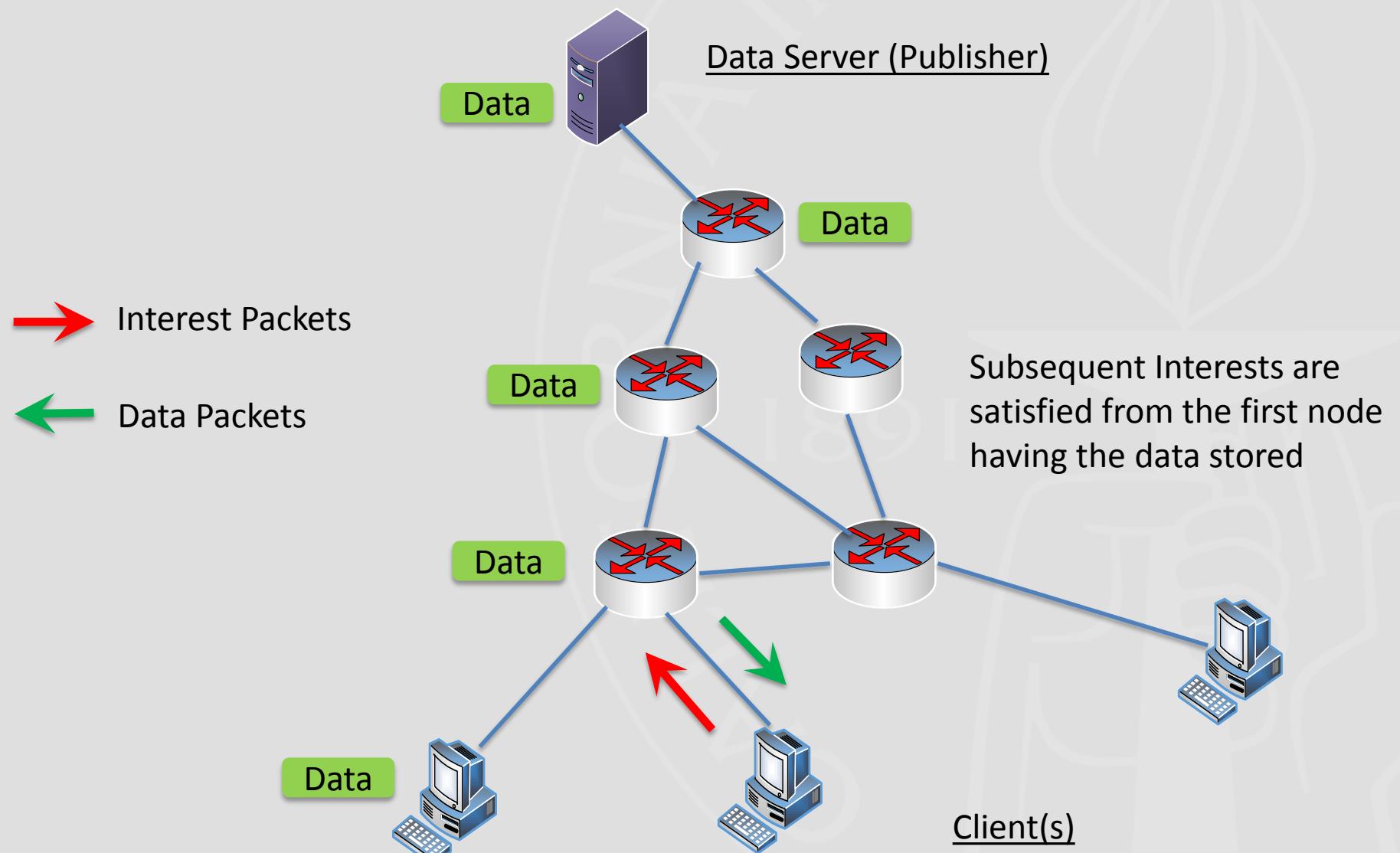
NDN Operation – New request



NDN Operation – New request



NDN Operation – New request



Why should we investigate use of NDN?



- A potential candidate technology to solve several issues, but do so at the network layer:
 - Optimal data distribution
 - Data caching
 - Popularity based data placement
 - Latency optimization for remote data access
- NDN could simply be the way the Internet works in the future
- How will this change the way we access and process data?



Some topics for investigation



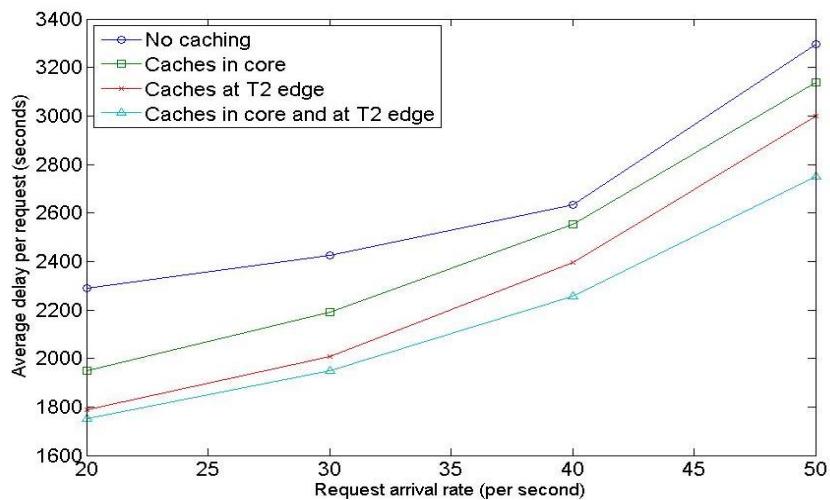
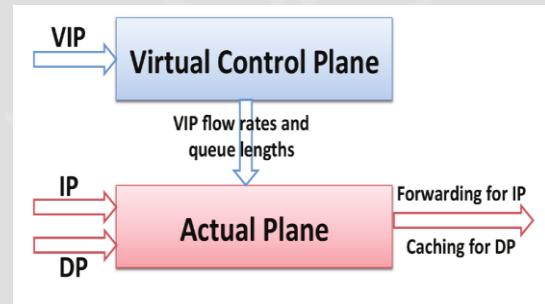
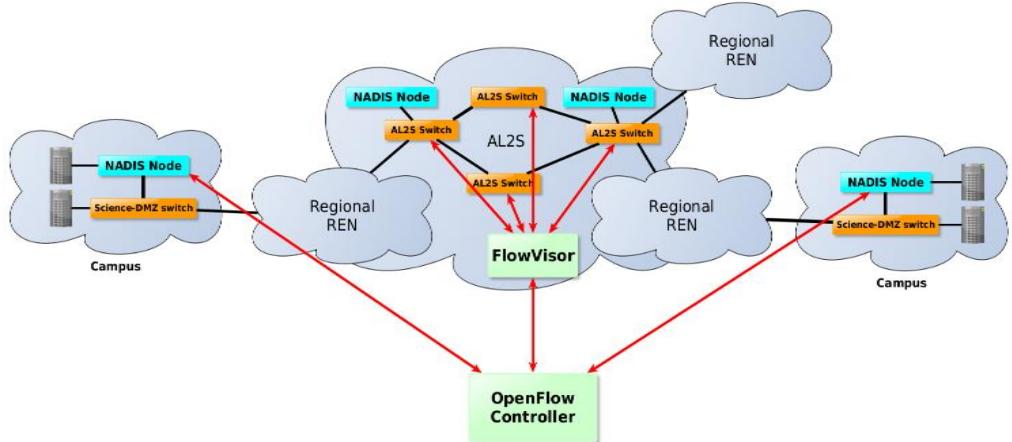
- Caching strategies
 - Could we reduce the storage at end-sites to only permanent copies?
 - Rely on caching in the network?
- What is the correct data chunk?
 - File? Block? Event?
- Bulk data transfer strategies
 - E.g. multipath, multi-source, multicast
- Multipath forwarding
- Network-Application interface
 - Sockets?
 - Calendaring?
- Impact on workflows and job scheduling
 - Reduced latency through caching
 - Rely on remote data access as default?
- QoS and flow prioritization
- ...



Combining SDN and NDN

- For starters, SDN can be an easy way to create a high performance NDN test bed
- In which we want to investigate a possible design suitable for HEP data (and other data intensive science fields)

Caltech and Northeastern proposal to NSF: NDN Architecture for Data Intensive Science (NADIS)



SDN/NDN - Summary



- Software Defined Networking provides a powerful way to interact with the network
 - Needs engagement and collaboration with the network operators
- Named Data Networking is a fresh approach at the design of the Internet of the future
 - Designed with the content rather than end-points as basis for communication
 - Has many features which can benefit LHC data processing
 - Despite it being rather new, basic implementation and a test bed are available
 - **The underlying ideas match very well with distributed data and computing models as in HEP computing**
 - Impact on the LHC data processing models needs careful study



NETWORK TESTBEDS

A non-exclusive list of examples for people interested in practical network innovation



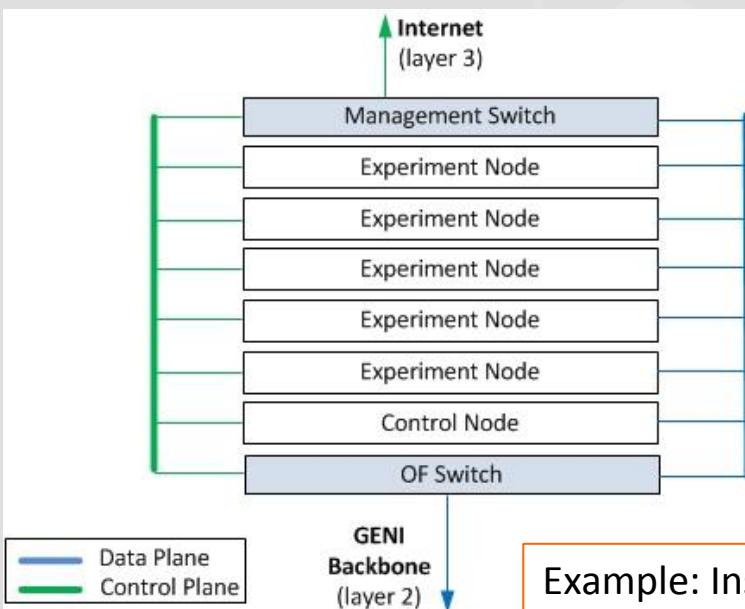
Network Testbeds



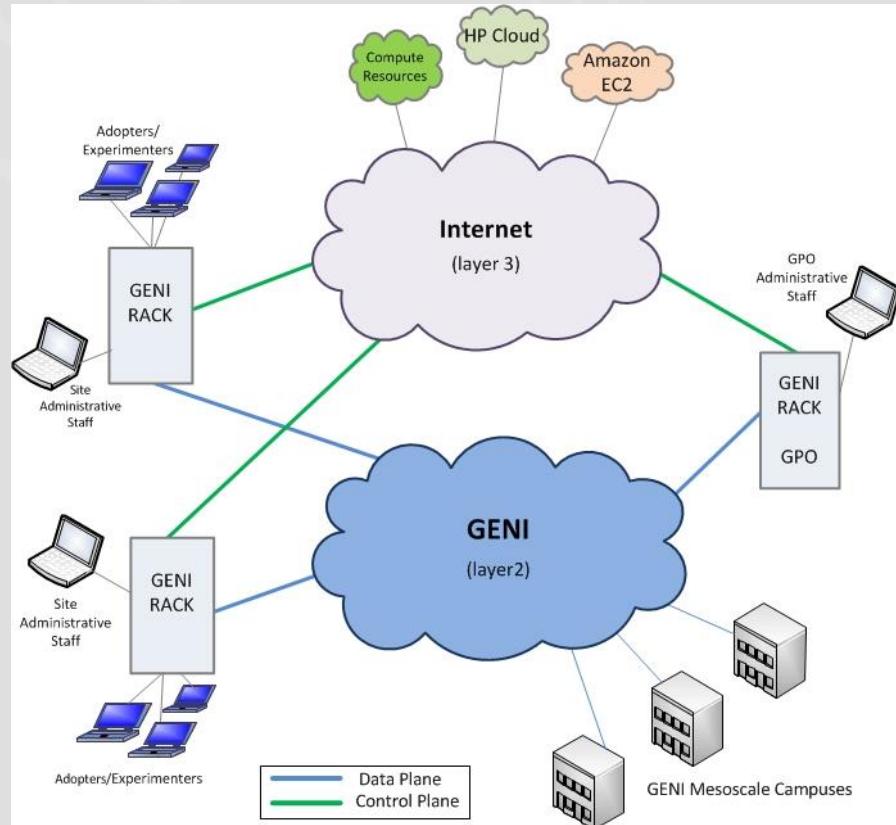
- If you want to test new network ideas, several testbeds might be available:
 - GENI – generic network testbed, mostly US
 - OFELIA – European Openflow testbed
 - GEANT OpenFlow test facility
 - FELIX – EU-JAPAN testbed for FI research
 - ...



- Global Environment for Network Innovations
- “virtual laboratory for networking and distributed systems research and education”
- Mainly US based initiative, but not only
- GENI racks (3 “models”) installed on several campuses



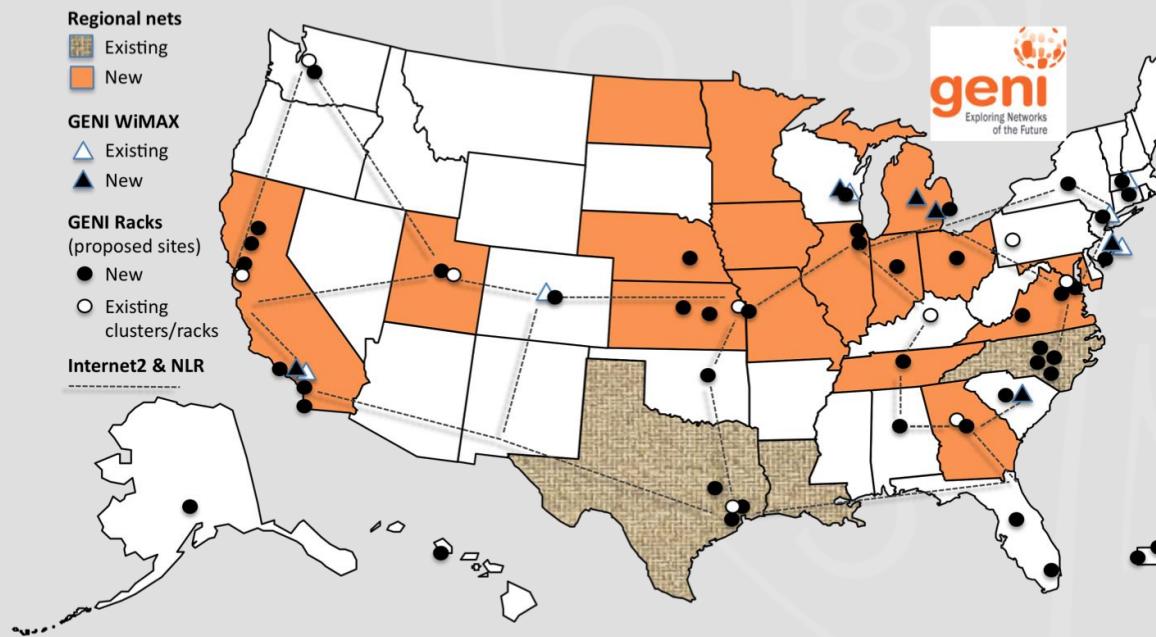
Example: InstaGENI rack layout



<http://www.geni.net>

GENI, cont.

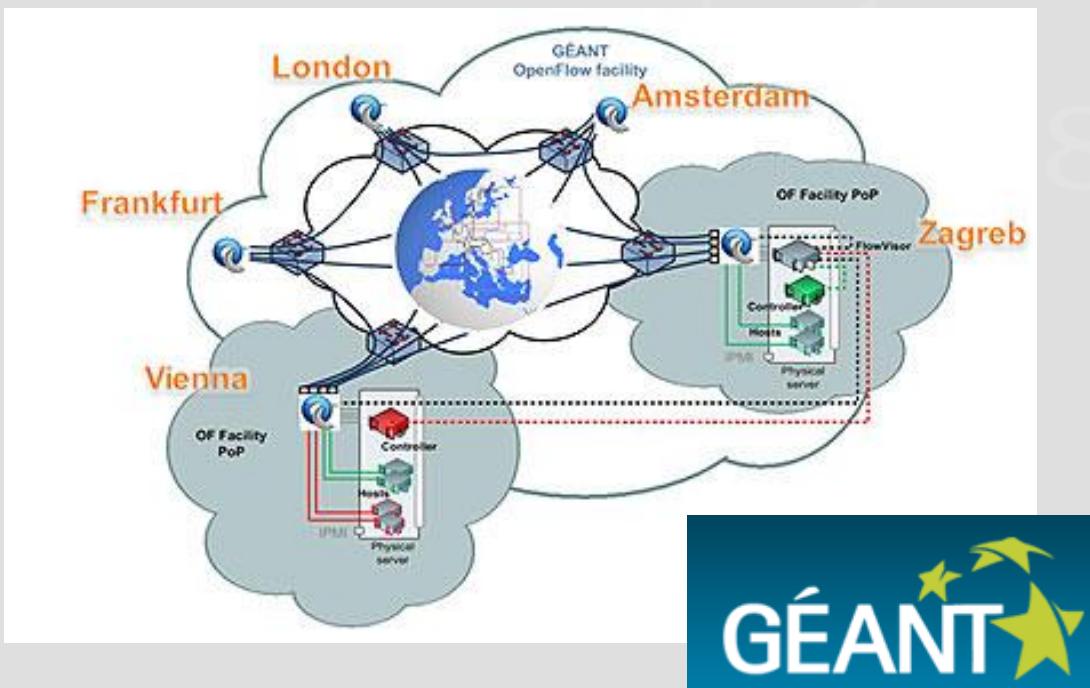
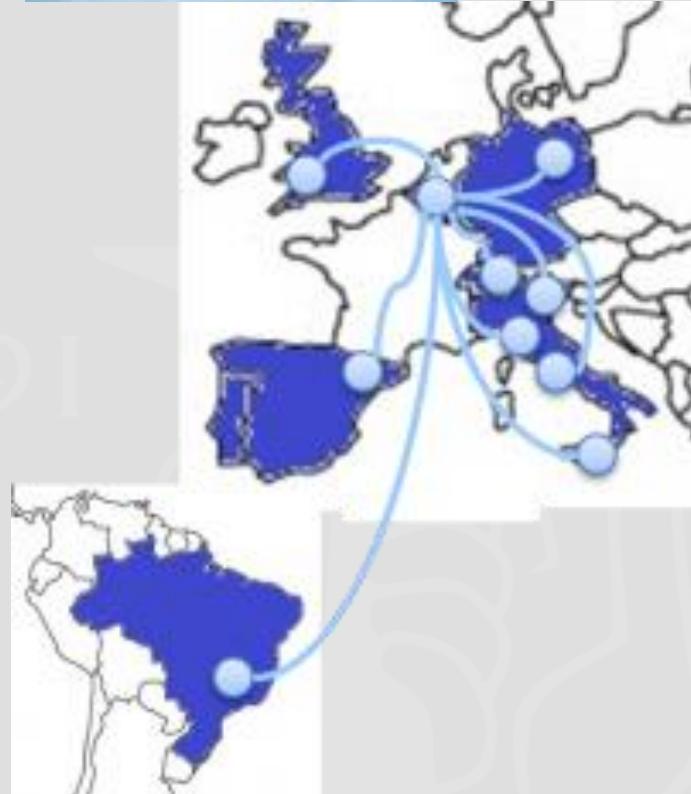
- GENI allows experimenters to:
 - Obtain compute resources from locations around the US
 - Connect compute resources using Layer 2 networks in topologies best suited to their experiments
 - Install custom software or even custom operating systems on these compute resources
 - Control how devices in their experiment handle traffic flows
 - Run their own Layer 3 and above protocols



<http://www.geni.net>



- European FP7 Project
- European OpenFlow Testbed Facility
- Project ended in 2013, but the GEANT Openflow Facility continues

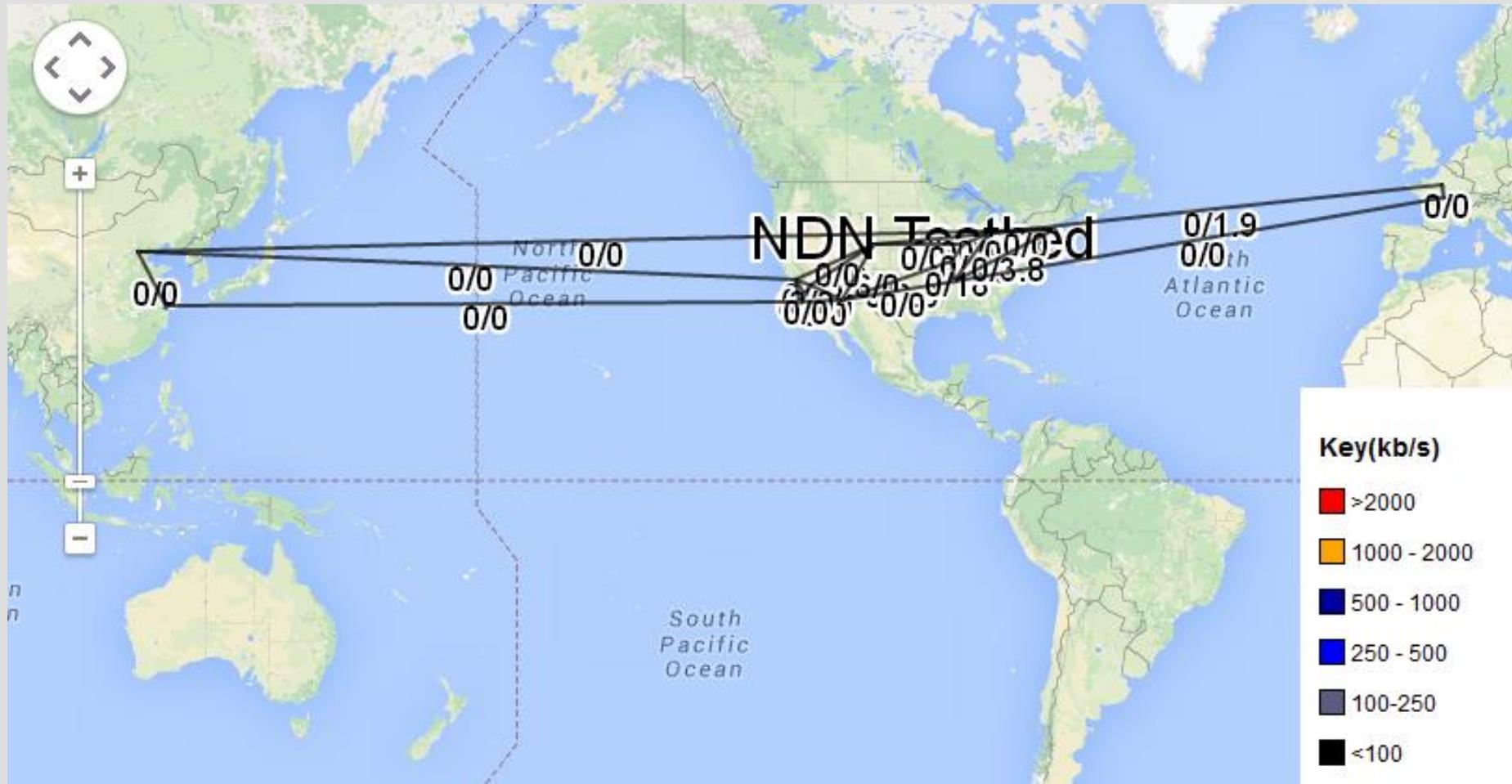


- FEderated Test-beds for Large-scale Infrastructure eXperiments
- EU-JAPAN Project
- SDN-oriented service architecture for federating Future Internet facilities like OFELIA and JGN-X RISE
- Use high-capacity NSI-enabled networks as substrate
 - JGN-X, GEANT, GLIF, NRENs
- On-demand setup of “OpenFlow based network slices”
 - including network, compute and storage



<http://www.ict-felix.eu>

NDN Testbed



Summary and Conclusions



- Networks are not any more providing only transmission of bits between a pair of hosts
- New developments are in areas above providing bandwidth
- In development of distributed computing systems, we should leverage the new capabilities of the network systems
- Engagement with the network service providers (NRENs) is necessary in order to benefit most from it



QUESTIONS & DISCUSSION

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Some Resources



- 1) OGF NSI WG <http://redmine.ogf.org/projects/nsi-wg>
- 2) Open Networking Foundation: <http://www.opennetworking.org>
- 3) Floodlight controller: <http://www.projectfloodlight.org/floodlight/>
- 4) OpenDaylight: <http://www.opendaylight.org/>
- 5) Named Data Networking: <http://named-data.net/>
- 6) CCNx: <http://www.ccnx.org/>

