

# ITP20001/ECE 20010 Data Structures

# **Data Structures**

# **Chapter 3**

- stacks & queues applications
  - infix to postfix

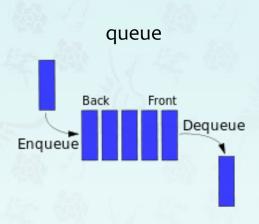




## 3.3 Queues

**Queue:** An ordered list in which **enqueues** (insertion or add) at the **rear** and **dequeues** (deletion or remove) take place at different end or **front**. It is also known as a Fist-in-first-out(FIFO) list.



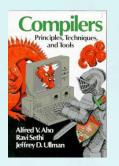


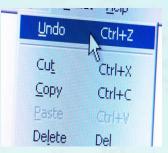
Items can only be added at the rear of the queue and the only item that can be removed is the one at the front of the queue.





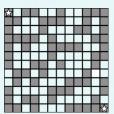
## 3.4 Stack and queue applications

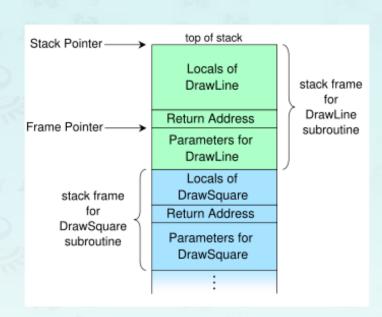






- Parsing in a compiler. (p.127)
- Undo in a word processor.
- Back button in a Web browser.
- PostScript language for printers.
- Backtracking as in a maze (p.121)
- Implementing function calls in a compiler. (p.108)
- ...









## 3.4 Stack and queue applications

In a computer OS: Requests for services come in unpredictable order and timing, sometimes faster than they can be serviced.

- print a file
- need a file from the disk system
- send an email
- job scheduling



## 3.6 Arithmetic expression evaluation

Goal: Convert an infix expression to a postfix expression using a stack.



Stack: ( Output:

Stack: ( Output: 1

Stack: ( + Output: 1

Stack: ( + Output: 1 2

Stack:

Output: 12+

Stack: \*

Output: 12+

Stack: \*

Output: 12 + 3

Stack:

Output: 1 2 + 3 \*

- Operands are output immediately

- Stack operators until right parens

Unstack until left parens
 Delete left parens

- In general, higher precedence operator must be output before lower one.)

postfix



## 3.6 Arithmetic expression evaluation

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Stack: ( Output:

Stack: ( Output: 1

Stack: ( + Output: 1

Stack: ( + Output: 1 2

Stack:

Output: 12+

Stack: \*

Output: 1 2 +

Stack: \*

Output: 12 + 3

Stack:

Output: 1 2 + 3 \*

- Operands are output immediately
- Stack operators until right parens
- Unstack until left parens
   Delete left parens
- In general, higher precedence operator must be output before lower one.)



## 3.6 Arithmetic expression evaluation

Goal: Evaluate infix expressions.

put parenthesis wherever possible

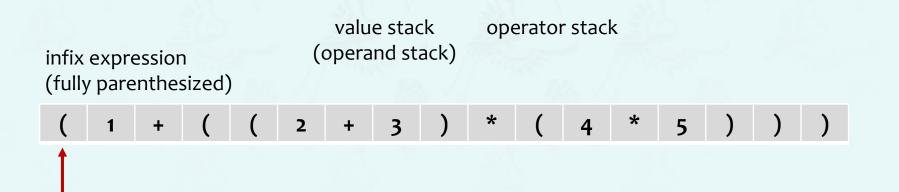
# Two-stack algorithm. [E. W. Dijkstra]

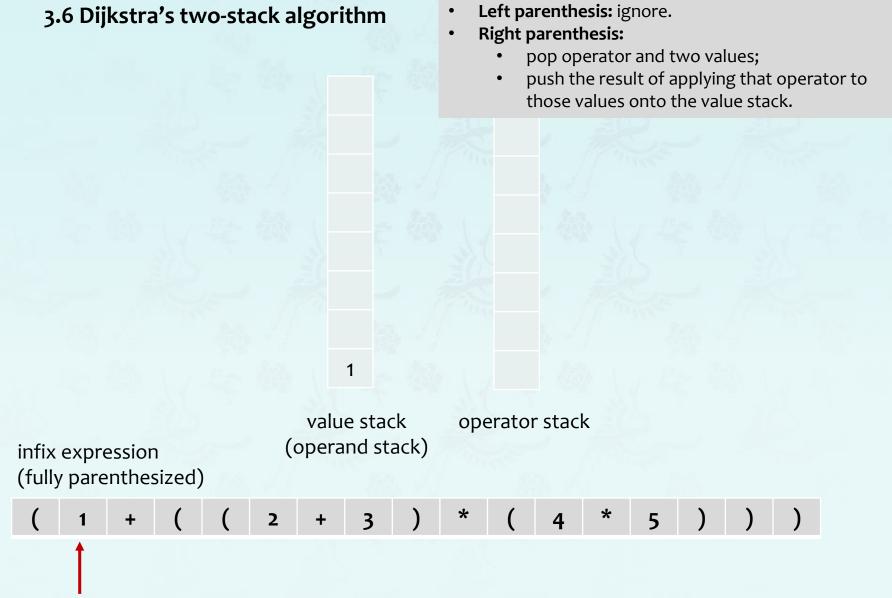
- Value: push onto the value stack.
- Operator: push onto the operator stack.
- Left parenthesis: ignore.
- Right parenthesis:
  - pop operator and two values;
  - push the result of applying that operator to those values onto the value stack.

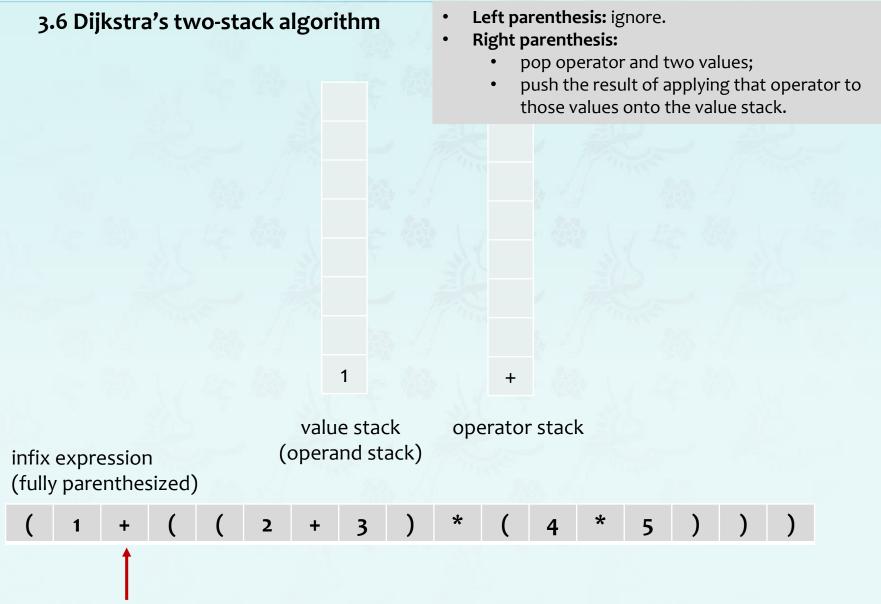


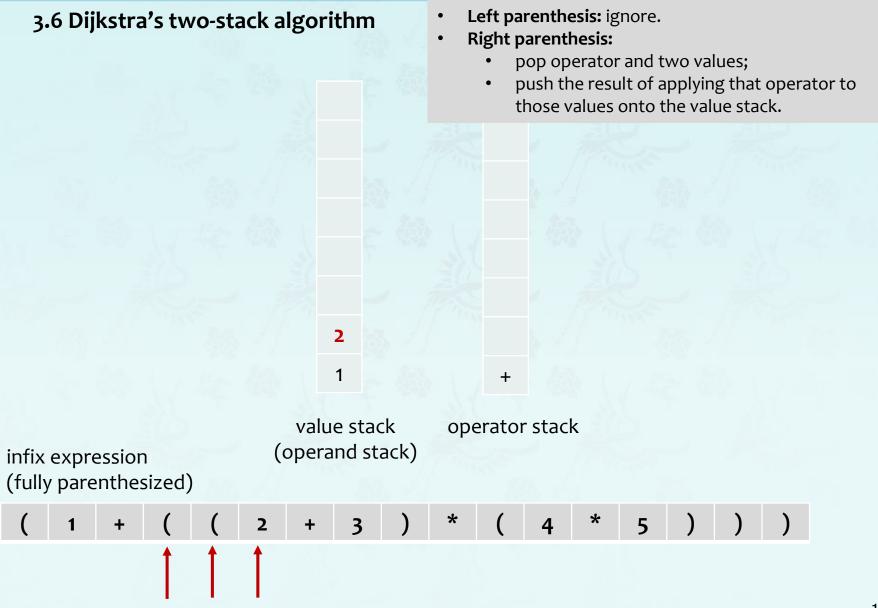


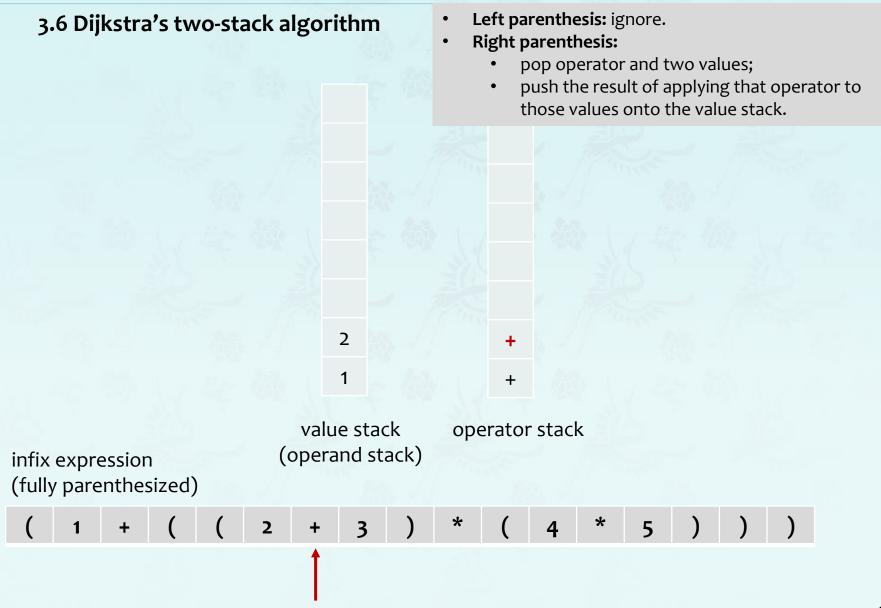
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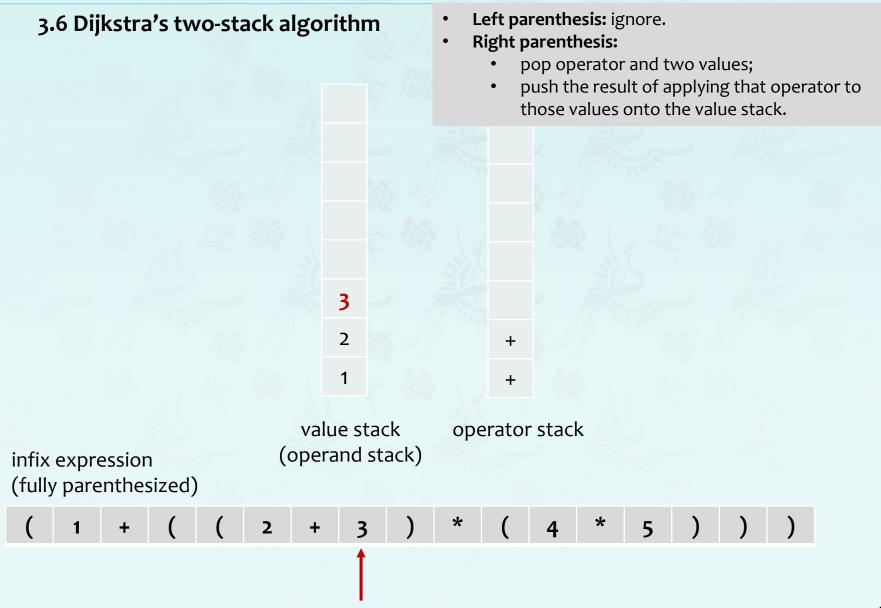


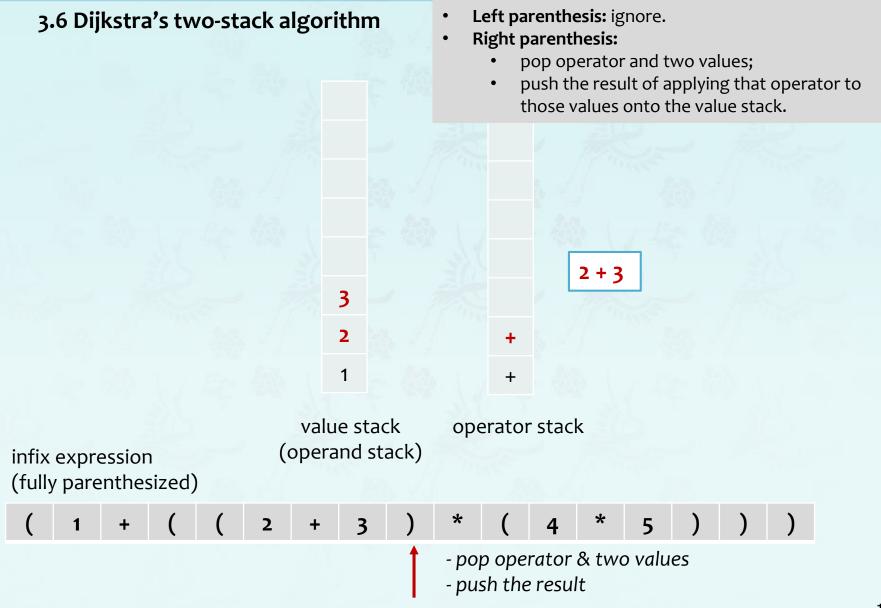


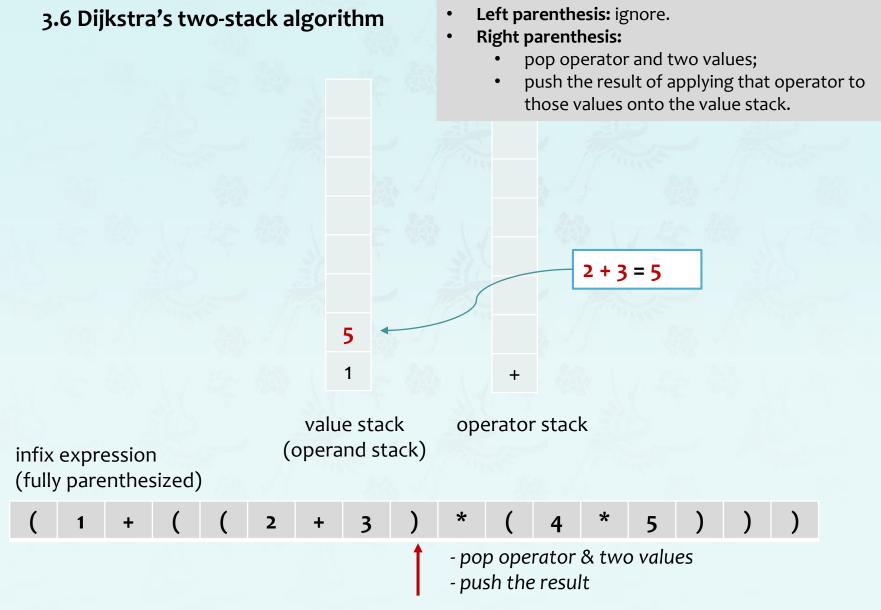


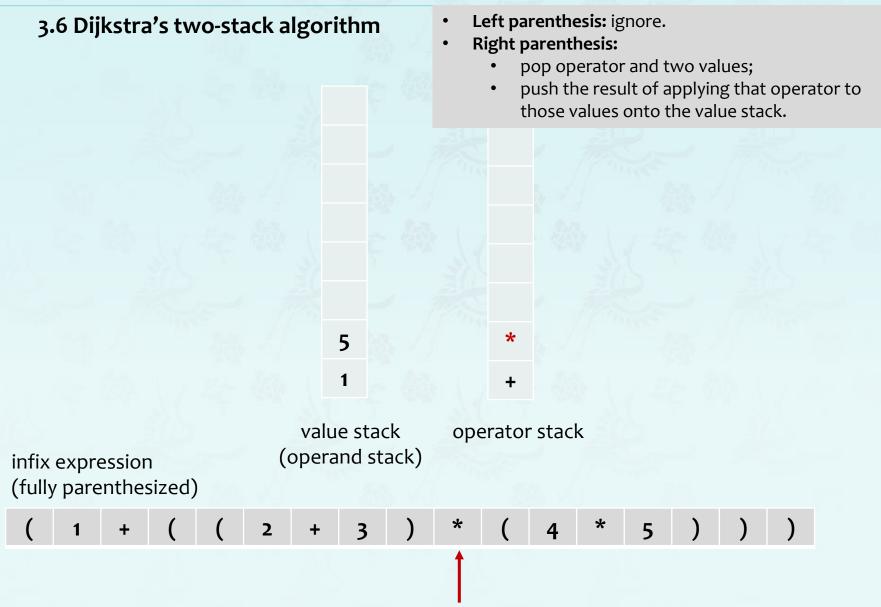




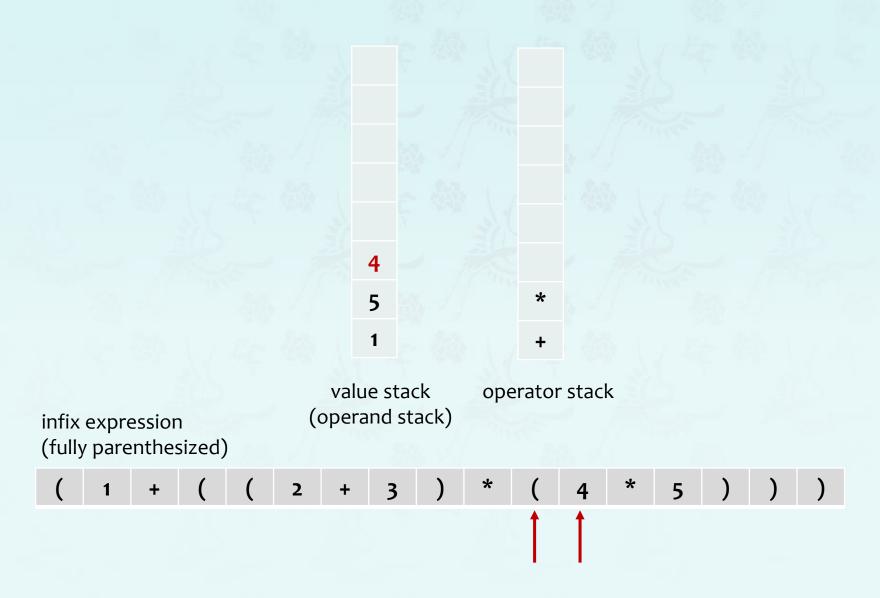




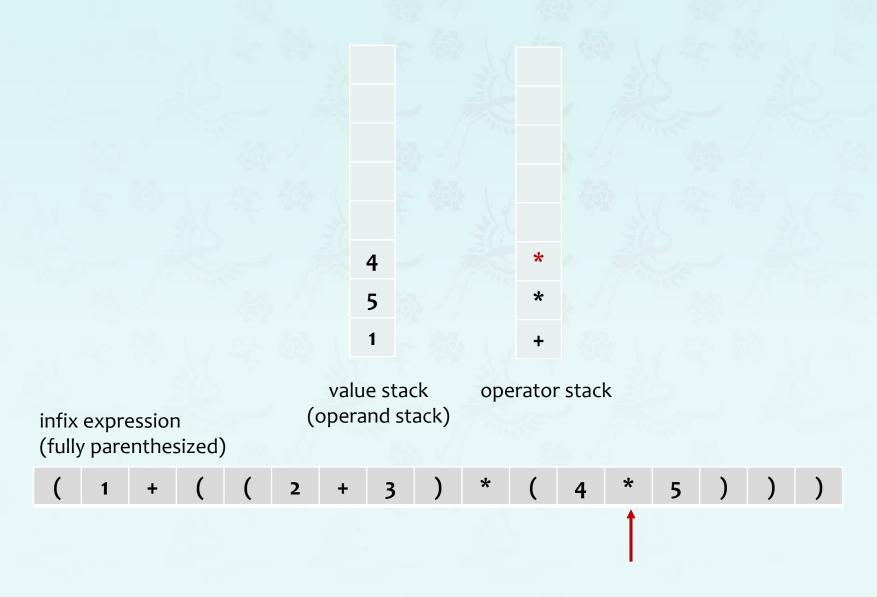




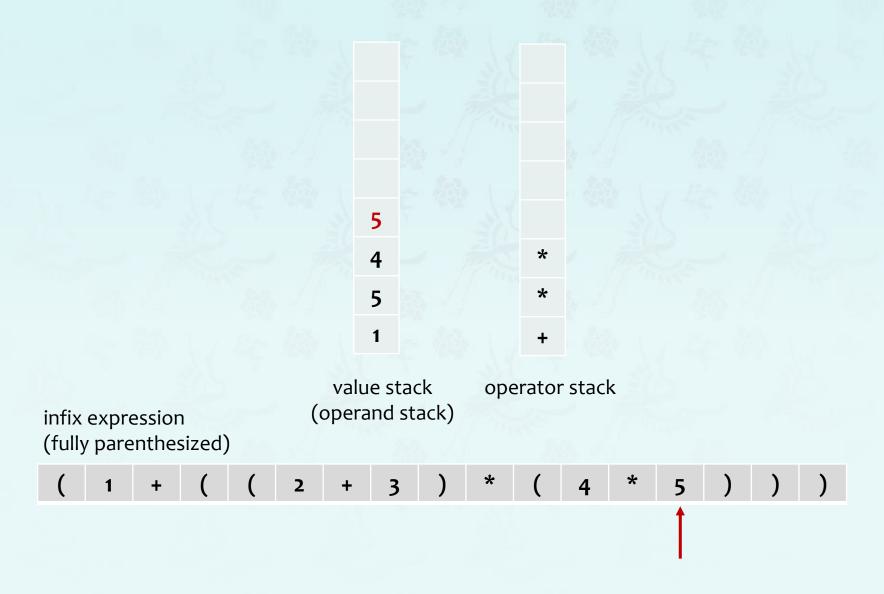




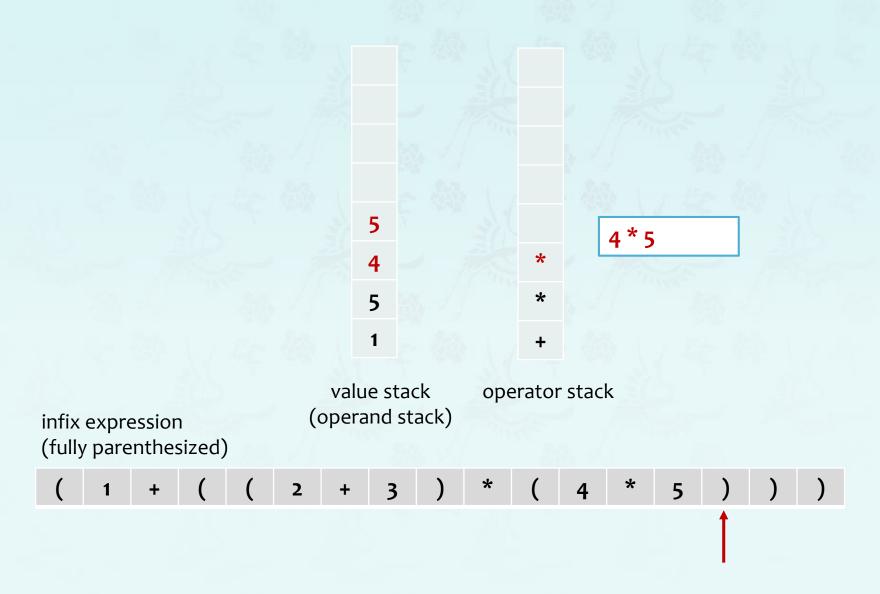


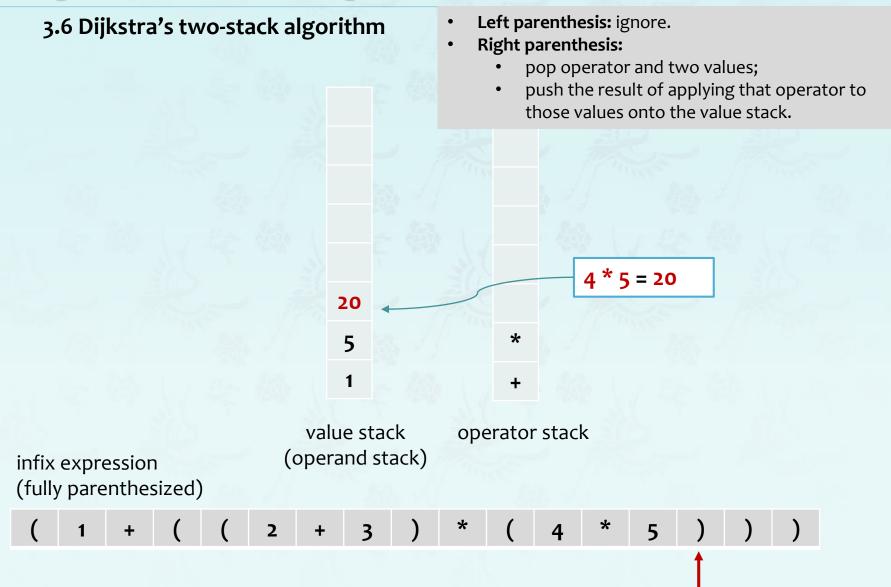


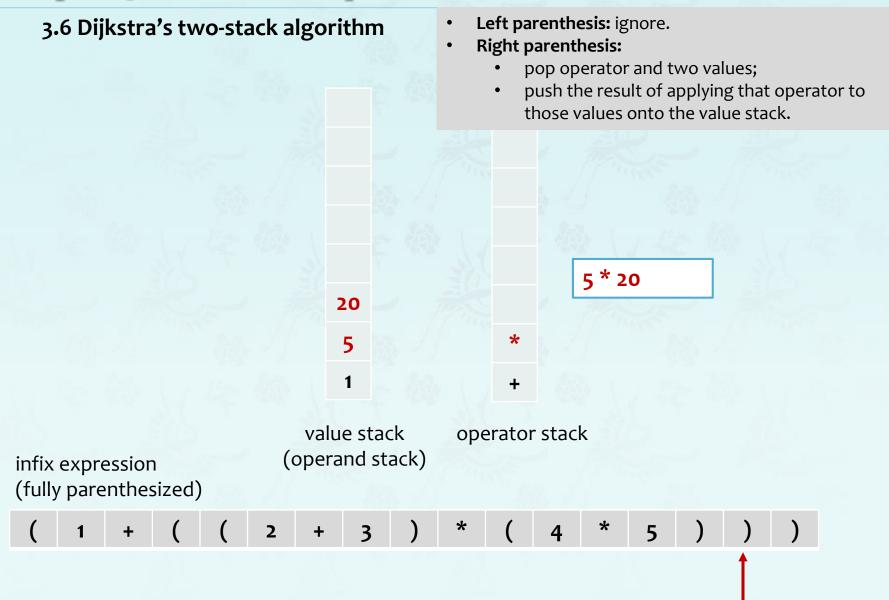


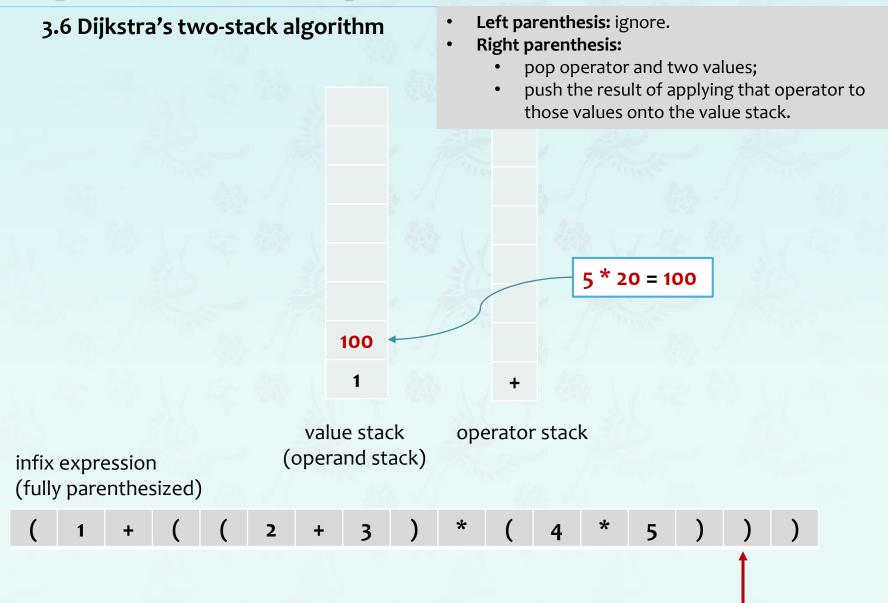


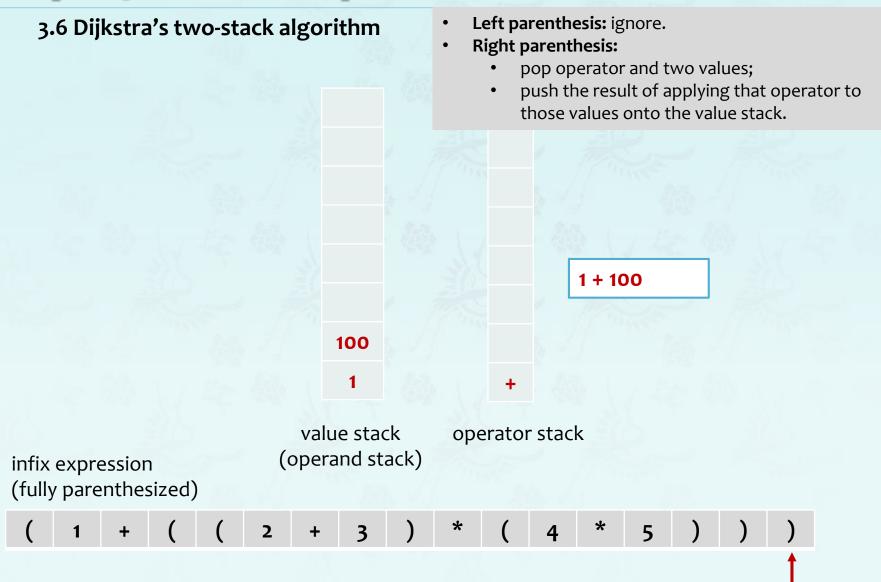




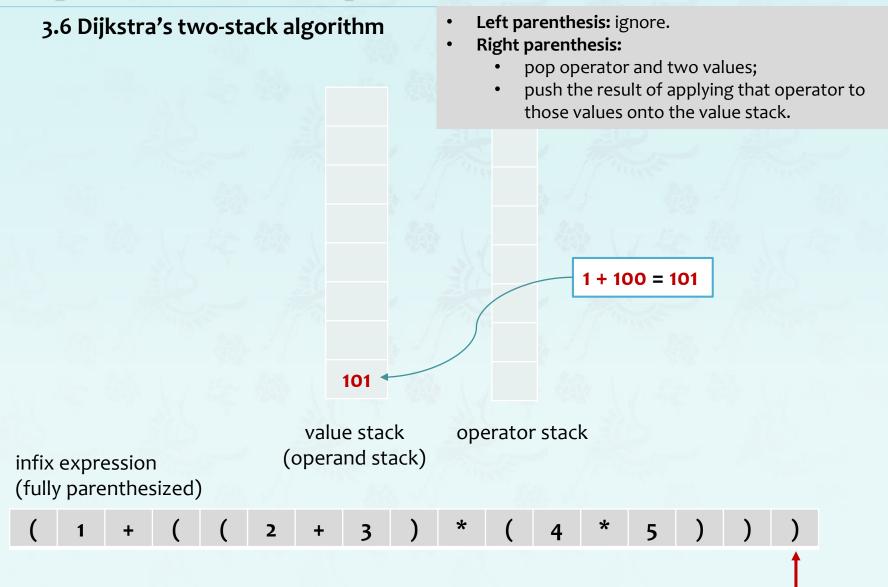




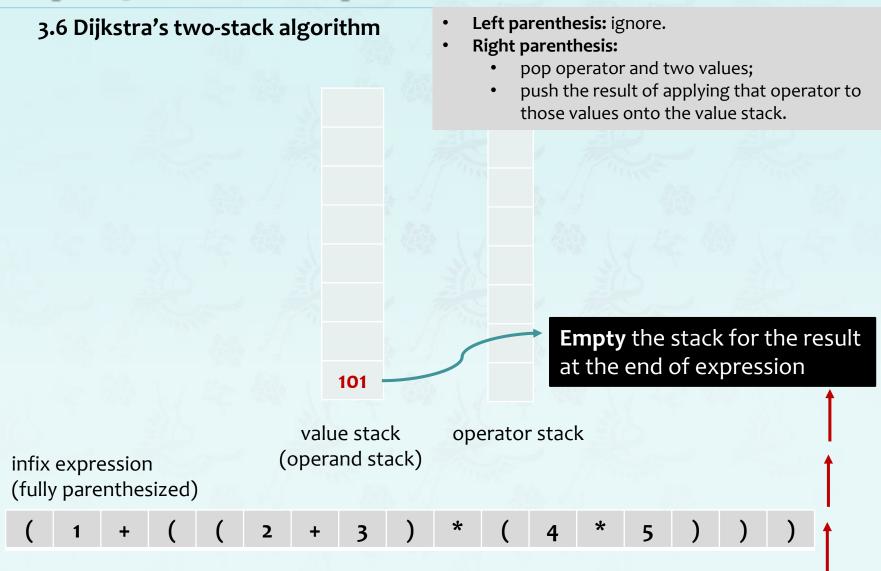














```
public class ArithmeticExpression {
public static void main(String[] args) {
 Stack<Character> ops = new Stack<Character>();
 Stack<Double> vals = new Stack<Double>();
 String e = JOptionPane.showInputDialog(null,
    "Enter an expression", "Stack application", JOptionPane.QUESTION MESSAGE);
 if (e == null) return;  // Check "Cancel"
 for (int i = 0; i < e.length(); i++) {
   Character c = e.charAt(i);
    if (c.equals(' ') || c.equals('('));
   else if (c == '+') ops.push(c);
   else if (c == '*') ops.push(c);
   else if (c == ')') {
       Character op = ops.pop();
              (op.equals('+')) vals.push(vals.pop() + vals.pop());
       if
       else if(op.equals('*')) vals.push(vals.pop() * vals.pop());
   else {
       String s = "" + c;
       vals.push(Double.parseDouble(s));
  JOptionPane.showMessageDialog(null, e + " = " + vals.pop());
```

## 3.6 Arithmetic expression evaluation

Q: How does it work?

A: When algorithm encounters an operator surrounded by two values within parentheses, it leaves the result on the value stack.

as if the original input were:

Repeating the argument:

Extensions: More ops, precedence order, associativity.



## 3.6 Arithmetic expression evaluation

**Observation 1.** Dijkstra's two-stack algorithm computes the same value if the operator occurs **after** the two values.

**Observation 2.** All of the parentheses are redundant!

Bottom line: Postfix or "reverse Polish" notation.

Applications: Postscript, calculators, JVM, ....



# 3.6 Arithmetic expression evaluation

infix	postfix
2 + 3 * 4	2 3 4 * +
a * b + 5	a b * 5 +
(1 + 2) * 7	
a * b / c	
(a/(b-c+d))*(e-a)*c	
a / b – c + d * e – a * c	

Figure 3.13 Infix and postfix notation



## 3.6 Arithmetic expression evaluation

Goal: Evaluate postfix expressions.



## 3.6 Arithmetic expression evaluation

Goal: Evaluate postfix expressions.



value stack (operand stack)

a

a b c - d + / e a - \* c \*

push the operands until an operator comes up.



## 3.6 Arithmetic expression evaluation

Goal: Evaluate postfix expressions.

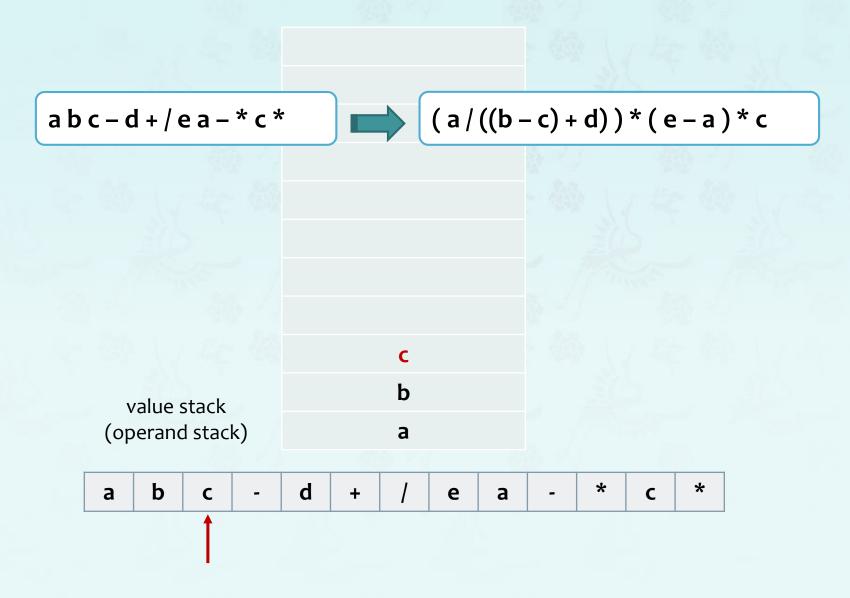
value stack (operand stack) b

a

a b c - d + / e a - \* c \*



## 3.6 Arithmetic expression evaluation





## 3.6 Arithmetic expression evaluation

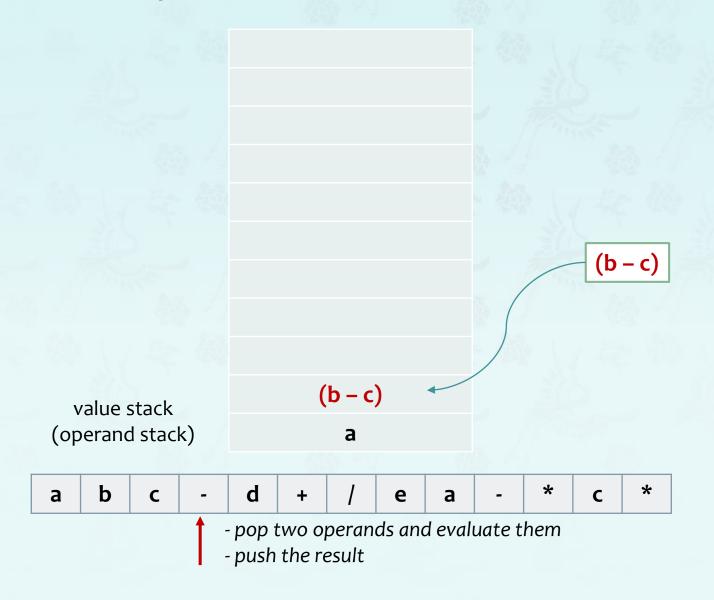
Goal: Evaluate postfix expressions.

- pop two operands and evaluate them

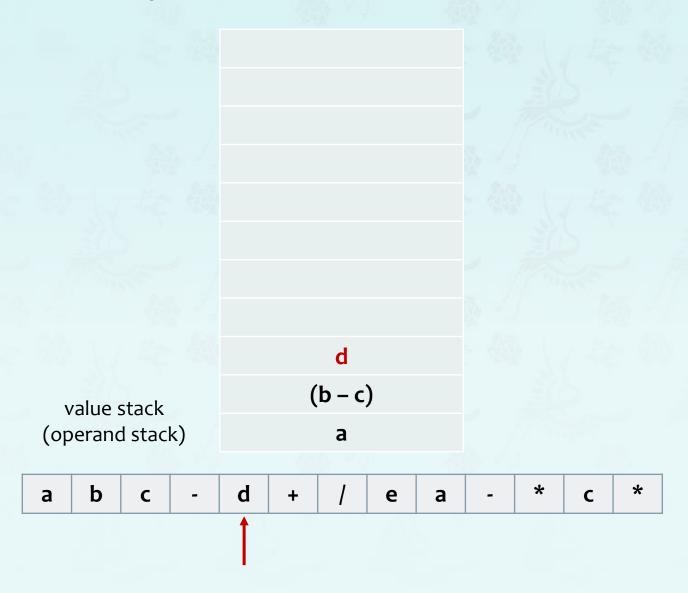
- push the result



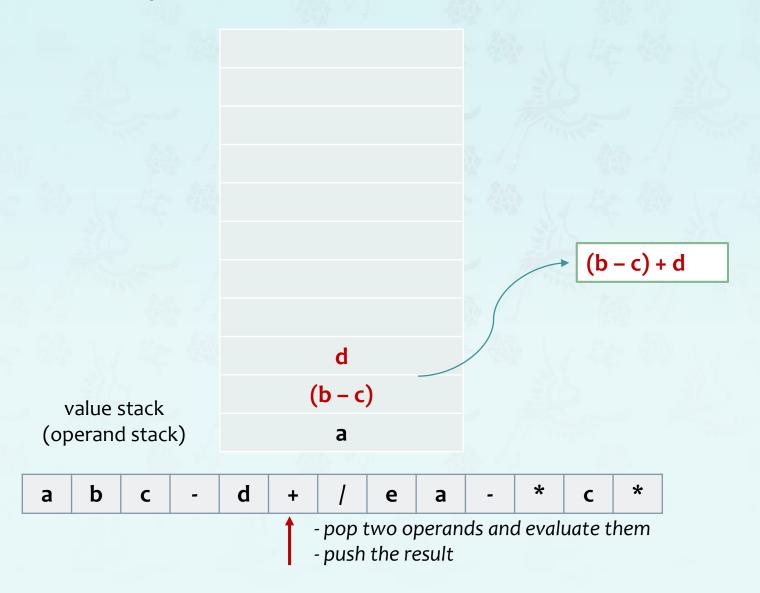
# 3.6 Arithmetic expression evaluation



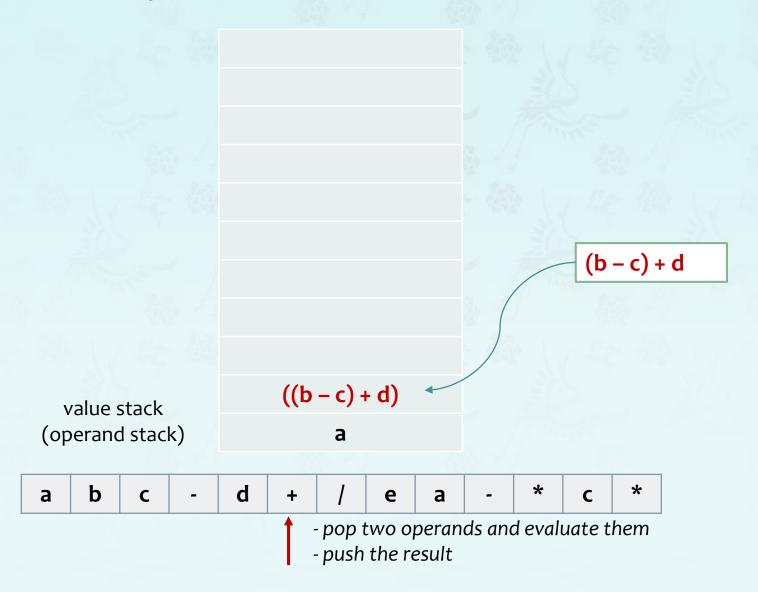




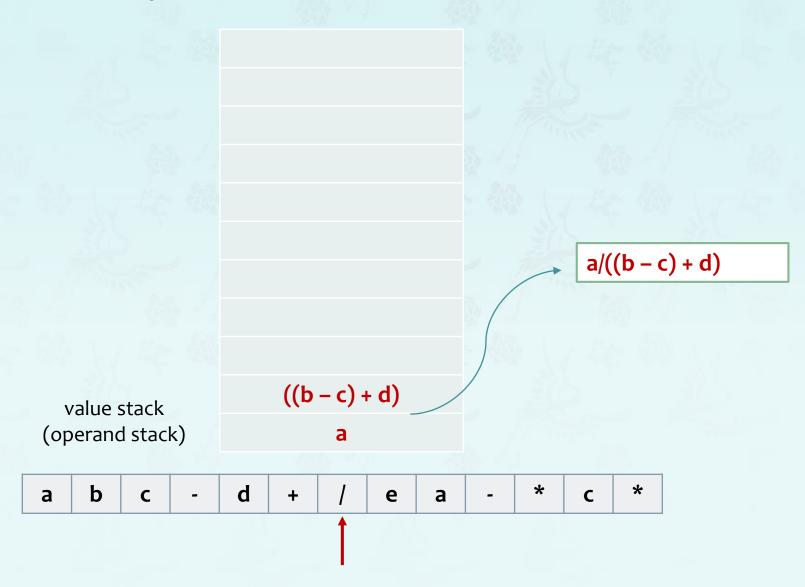




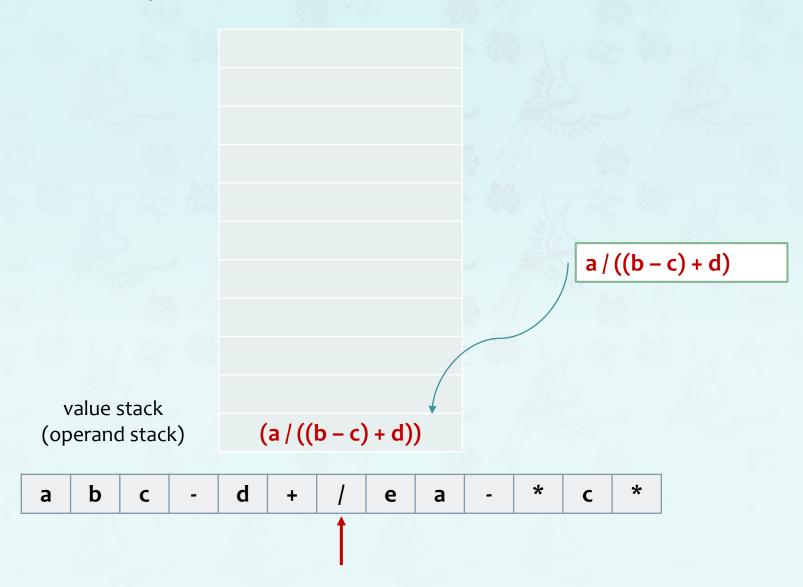




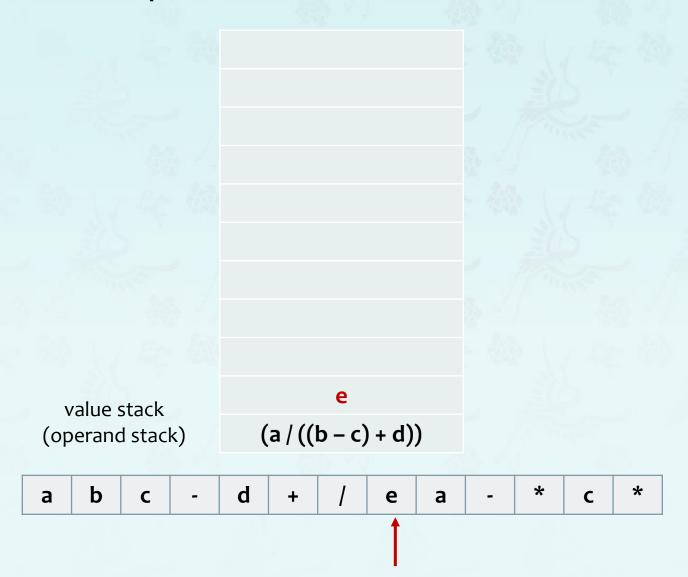




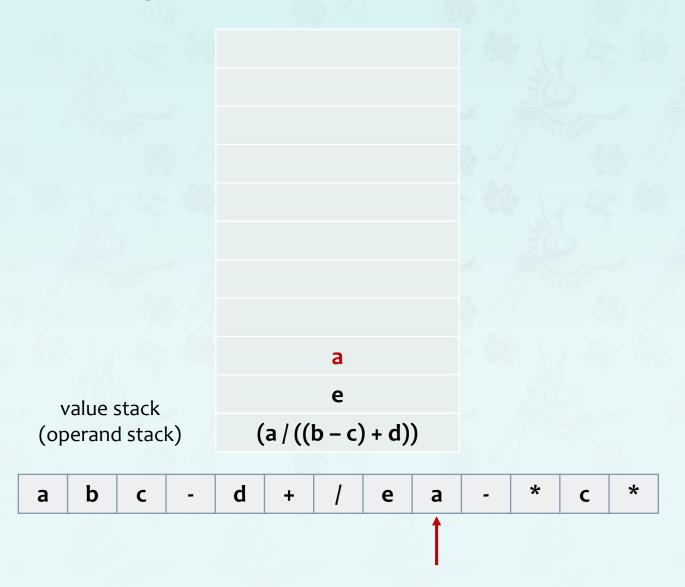




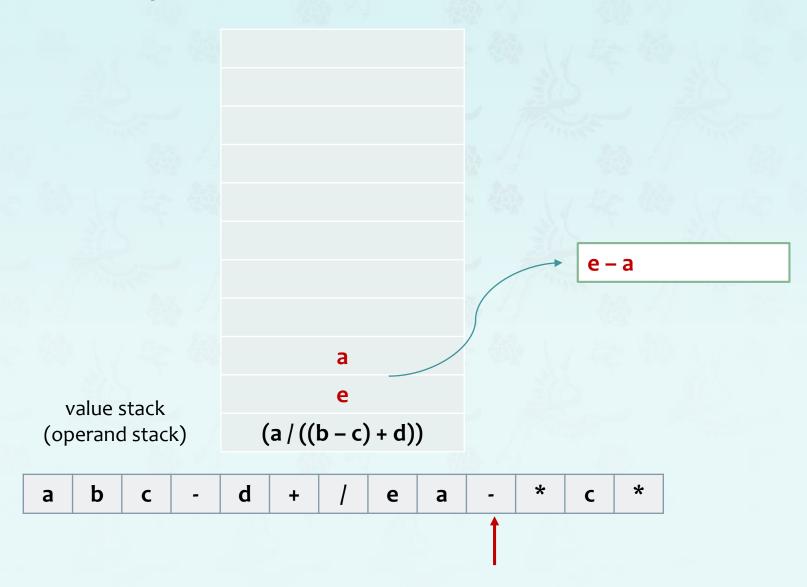




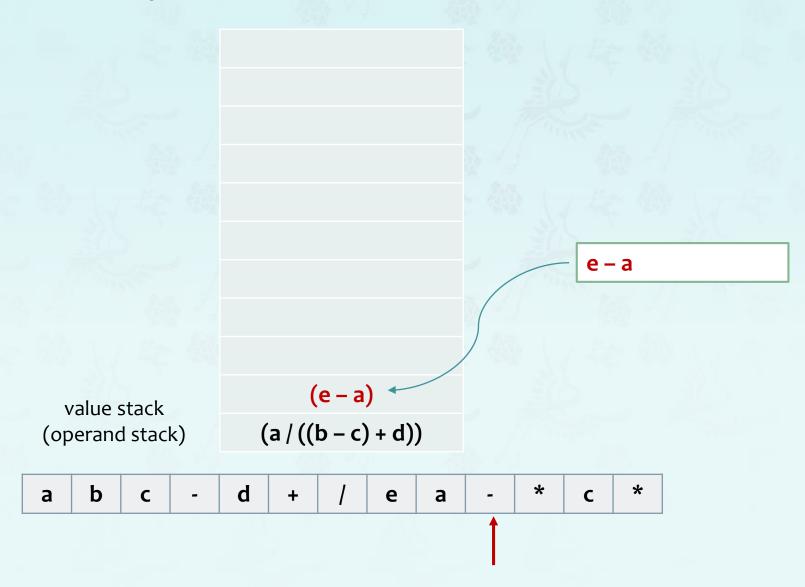




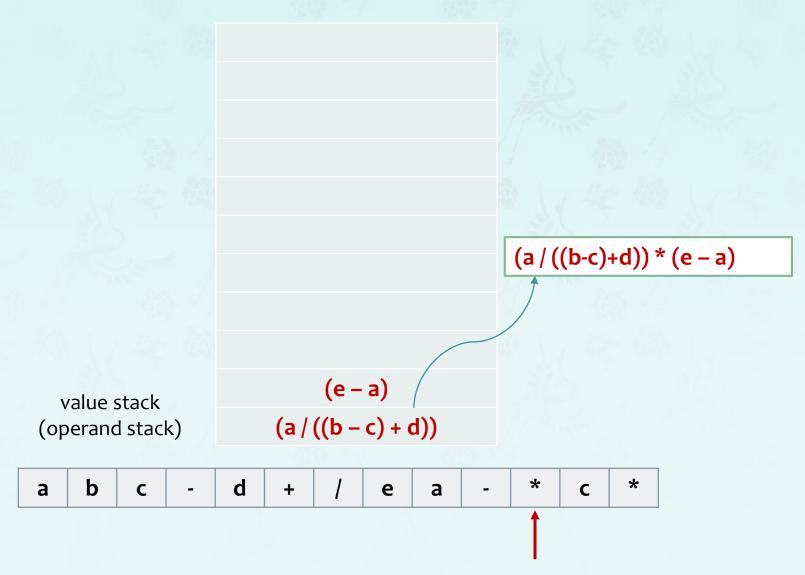




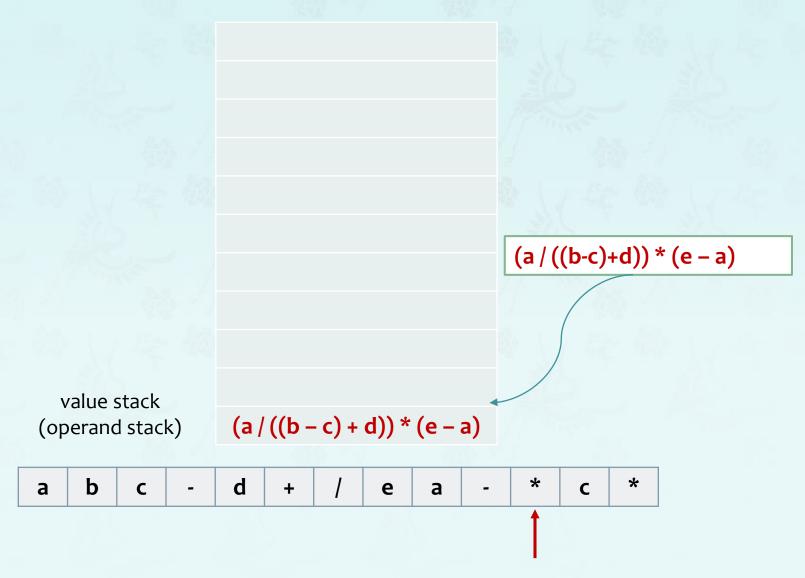




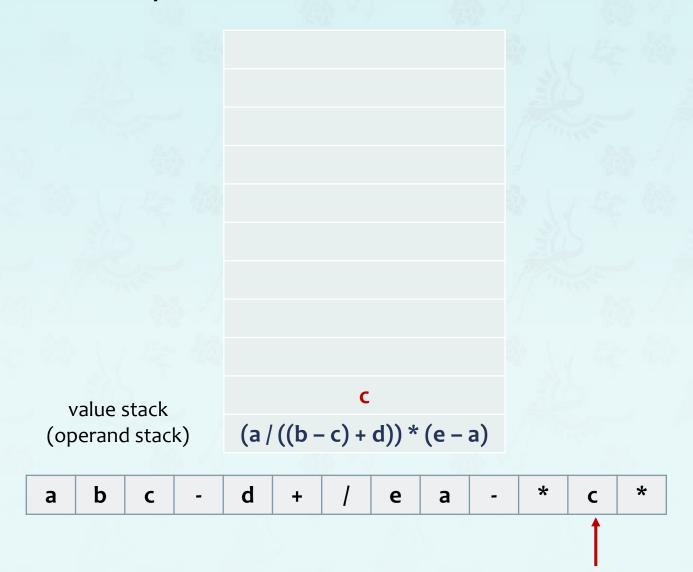














#### 3.6 Arithmetic expression evaluation

Goal: Evaluate postfix expressions.

#### 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion

#### We use a stack.

- When an operand is read, output it.
- When an operator is read,
  - Pop until the top of the stack has an element of lower precedence.
  - Then push it.
- 3. When ) is found, pop until we find the matching (.
- 4. (has the lowest precedence when in the stack but has the highest precedence when in the input.
- 5. When we reach the end of input, pop until the stack is empty.

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 1:

infix: 3+4\*5/6

in	stack(bottom to top)	postfix
3		
+		
4		
*		
5		
/		
6		

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 1:

infix: 3+4\*5/6

in	stack(bottom to top)	postfix
3		3
	+	
4		3 4
*	+ *	
5	5	3 4 5
/	+ /	3 4 5 *
6		3 4 5 * 6
		3 4 5 * 6 / +

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 2:

infix: (1+3)\*(4-2)/(5+7)

in	stack (bottom to top)	postfix	in	stack	postfix
(	(		(	/ (	1 3 + 4 2 - *
1		1	5		1 3 + 4 2 - * 5
+	( +		+	/ ( +	
3		1 3	7		1 3 + 4 2 - * 5 7
)		1 3 +	)		13+42-*57+
*	*				13+42-*57+/
(	* (				
4		1 3 + 4			
-	* ( -				
2		1 3 + 4 2			
)	*	1 3 + 4 2 -			
/	/	1 3 + 4 2 - *			

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 3:

infix: 
$$a - (b + c * d) / e$$

in	stack(bottom to top)	postfix
	stack(bottom to top)	postnik
а		
-		
(		
b		
+		
С		
*		
d		
)		
/		
е		

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 3:

infix: 
$$a - (b + c * d) / e$$

in	stack(bottom to top)	postfix
а		a
-	-	
(	- (	
b		ab
+	- ( +	
С		abc
*	- ( + *	
d		abcd
)	-	abcd*+
/	-/	
е		abcd*+e
		abcd*+e/-

# 3.6 Arithmetic expression evaluation - Infix to Postfix Conversion Example 3:

infix: 
$$A * (B + C * D) + E$$

	in	stack(bottom to top)	postfix
1	Α		
2	*		
3	(		
4	В		
5	+		
6	С		
7	*		
8	D		
9	)		
10	+		
11	Е		
12			



#### 3.6 Arithmetic expression evaluation

Exercise: Infix to postfix, postfix to infix parsing



# **ECE 20010 Data Structures**

#### **Data Structures**

## **Chapter 3**

- abstract data types review
- stacks & queues using dynamic arrays
- some applications