(Byzantine) Vertical Paxos

Dahlia Malkhi EPFL Summer School on Blockchains, June 2017



How a typical application becomes distributed





How a typical application becomes distributed,







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Impossibility of distributed consensus with one faulty process

MJ Fischer, NA Lynch, MS Paterson - Journal of the ACM (JACM), 1985 - dl.acm.org
Abstract The **consensus** problem involves an asynchronous system of processes, some of
which may be unreliable. The problem is for the reliable processes to agree on a binary
value. In this paper, it is shown that every protocol for this problem has the possibility of
Cited by 4162 Related articles. All 98 versions. Cite. Save

[BOOK] Distributed consensus in multi-vehicle cooperative control

W Ren, RW Beard - 2008 - Springer

Recent advances in miniaturizing of computing, communication, sensing, and actuation have made it feasible to envision large numbers of autonomous vehicles (air, ground, and water) working cooperatively to accomplish an objective. Cooperative control of multiple Cited by 1715 Related articles All 14 versions Cite Save

On the minimal synchronism needed for distributed consensus

<u>D Doley, C Dwork, L Stockmeyer - Journal of the ACM (JACM), 1987 - di.acm.org</u>
Abstract Reaching agreement is a primitive of **distributed** computing. Whereas this poses no problem in an ideal, failure-free environment, it imposes certain constraints on the capabilities of an actual system: A system is viable only if it permits the existence of Cited by 745 Related articles All 27 versions Cite Save

Stability of continuous-time distributed consensus algorithms

L Moreau - Decision and Control, 2004. CDC. 43rd IEEE ..., 2004 - ieeexplore.ieee.org
Abstract: We study the stability properties of linear time-varying systems in continuous time
whose system matrix is Metzler with zero row sums. This class of systems arises naturally in
the context of **distributed** decision problems, coordination and rendezvous tasks and
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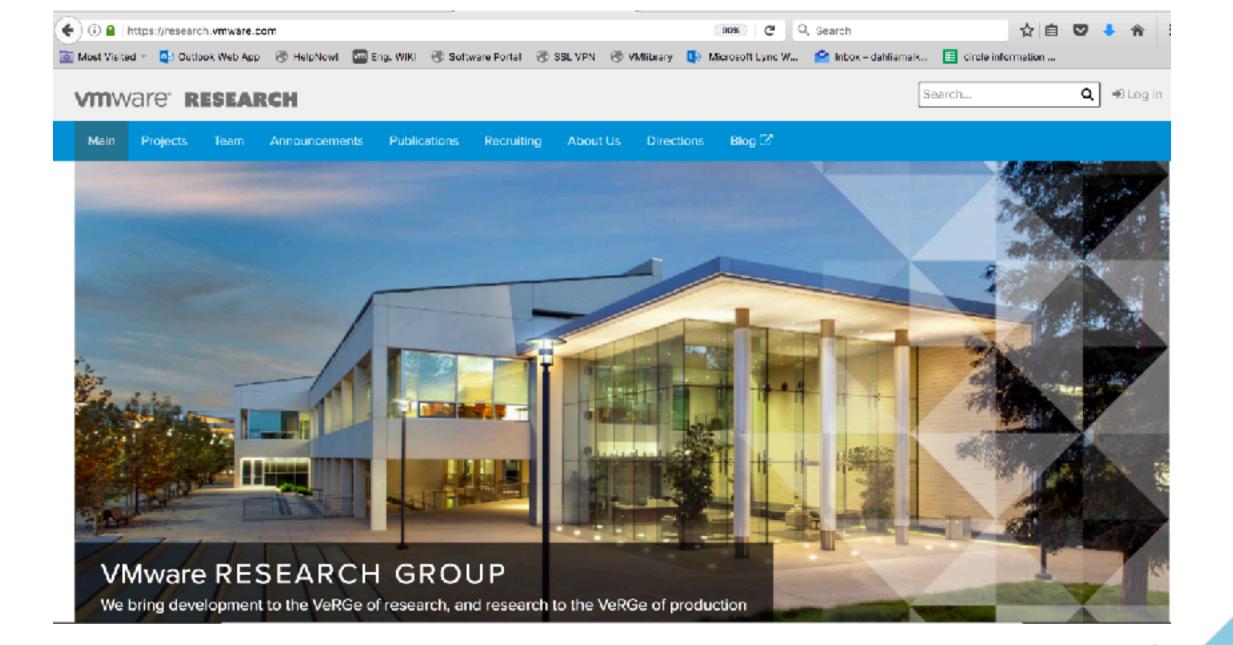
[PDF] arxiv.org



SMR Reconfiguration

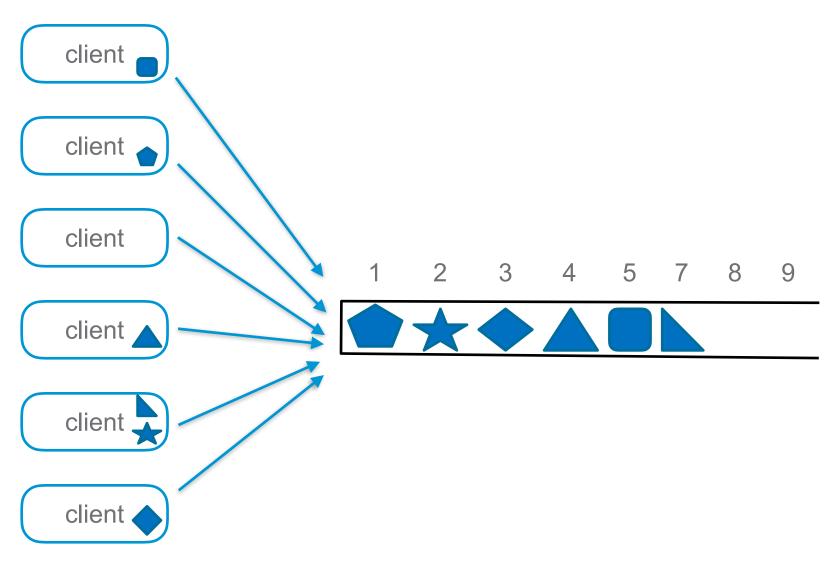
- Overview SMR
- (Vanilla) Paxos reconfiguration
- Contiguous log reconfiguration: VR, ZK, Raft
- Explicit reconfiguration: Cheap Paxos, Vertical Paxos, Virtually Synchronous Paxos, 700 BFTs
- Byzantine VP









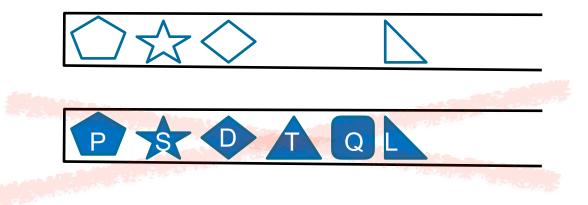






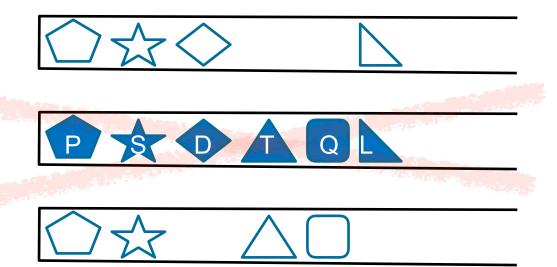


F-resilience





F-resilience

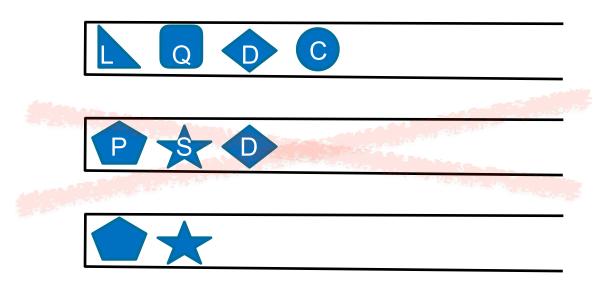




Glossary

- replica(s)
 - leader/coordinator
- command sequence
- clients
- learn

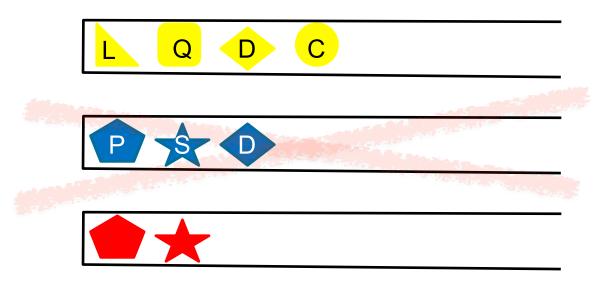
A replicated system after a stormy night





A replicated system after a stormy night

with leader ranks

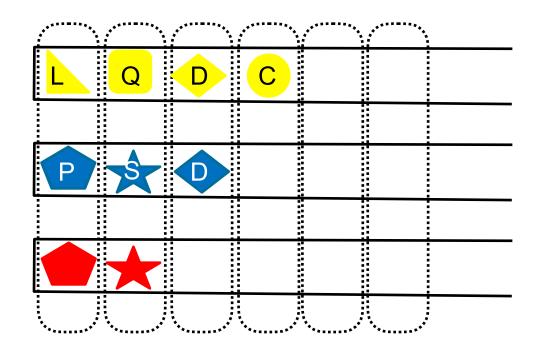




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Paxos: Two-Phase SMR Protocol

- First phase: leader replacement, amortized
- Second phase: replication

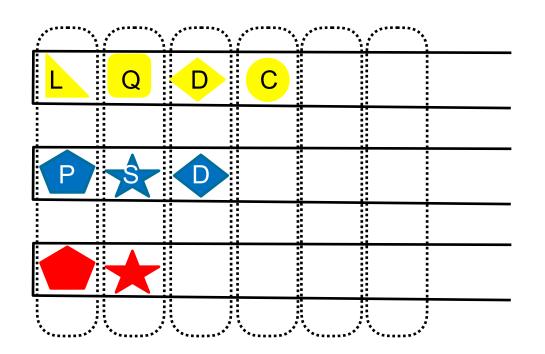






Contiguous log

most recent log wins







Vertical Paxos

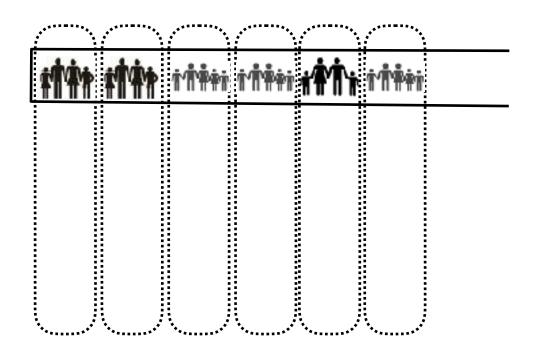
reconfiguration == leader replacement





reconfiguration

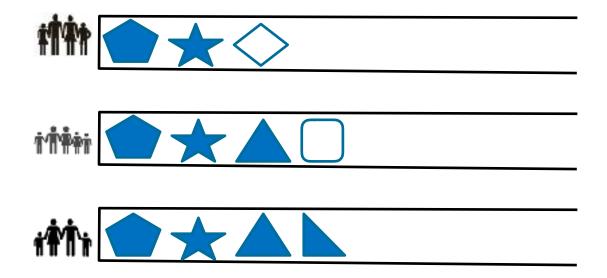
• intra-sequence





reconfiguration

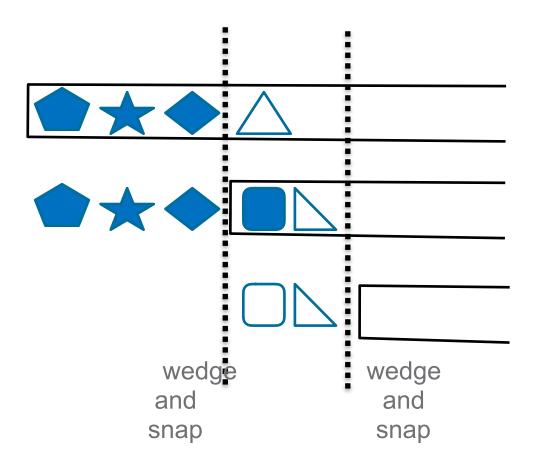
• intra-sequence





Virtually Synchronous Paxos

- inter-sequence
- separate steady state from reconfiguration
- two consensus decisions: state and next





Vertical Paxos

- steady state is F+1
- aux reconfiguration master
- Primary backup: special case, implicit wedge-and-snap
- Cheap Paxos: special case (F+1)-of-(2F+1), F standby's

why is it vertical?





rsection revisited

or Spiegelman!

Each of the phases of Paxos may use non-intersecting quorums. Only quorums from different phases are required to intersect. Majority quorums are not necessary as intersection is required only across

phases.



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Flexible Paxos: Use-Cases





BVP: Byzantine Vertical Paxos

- non-repudiation: interest in one correct replica
- wedge and snap: can "prove" decision



BVP in the synch model

- 3-message-delay
- steady state 2F+1
 - client-leader
 - leader-all (signed)
 - all-client (signed)
 - proceed with 2F+1 signed echoes
- closing state for reconfig
 - synchronously probe 2F+1
- liveness
 - (2F+1)-of-(3F+1)





BVP in the synch model

- 4-message-delay
- steady state F+1
 - client-leader
 - leader-all (signed)
 - all-all (signed)
 - all-client (signed)
 - proceed with F+1 signed echoes
- closing state for reconfig
 - XFT: synchronously probe F+1 by F+1
- liveness
 - (F+1)-of-(2F+1)







BVP in the asynch model

- w/TPM
 - similar to sync
 - steady state F+1
 - proof by HW attestation

A dynamic fault model

- interplay between adversary and system
- begin handover / end handover
- begin handover **to** speculative new: |new| F correct
- end handover from current: all of current can be faulty



Take-aways

- going beyond Paxos
- reconfiguration-based approach



overflow

ask me questions here; otherwise I'll go on.



Flexible Paxos: Use-Cases





Flexible Paxos: Use-Cases

