# Why Linear Types Are The Future Of Systems Programming

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# Outline

- ATS programming language
  - ML
  - linear types
  - refinement types
  - dependant types
  - As fast as C! ("blazing fast")
- Lots of typelevel madness
  - No optimizations
- Hongwei Xi
  - Boston University

- Very hard!
  - Research language
  - hbox overfull with ideas
  - Tons of accidental complexity
  - Keywords everywhere . . .
  - Zero docs
- And that's OK!
  - Our job to make usable things

- Goals
  - Not evangelism!
  - Not adoption!
  - Be dissatisfied
  - Inspire your next language

- Very difficult to present
  - Linear/dependant/refinement types, ML, C all converge
- Concrete motivating examples
  - High level handwaving
- Assuming comfort with ML like langs and basic C
- Start by taste of the ML & C side

• First from the ML side

• A linear Option (explanations come later ...)

```
datavtype Option_vt (a:vt@ype, bool) =
   | Some_vt(a, true) of (a)
   | None_vt(a, false)
```

• probably more familiar (\_vt for viewtype)

```
datavtype Option_vt =
    | Some_vt of (a)
    | None_vt
```

• Indexed on a type-level bool, dependent types!

```
datavtype Option_vt =
   | Some_vt(a, true) of (a)
   | None_vt(a, false)
```

• Parameterized on a view type, linear types!

```
datavtype Option_vt (a:vt@ype, bool) =
   | Some_vt(a, true) of (a)
   | None_vt(a, false)
```

All ADTs in ATS are GADTs

```
datavtype Option_vt (a:vt@ype, bool) =
   | Some_vt(a, true) of (a)
   | None_vt(a, false)
```

• A linear C array

```
absvtype arrayptr (a:vt@ype, 1:addr, n:int) = ptr(1)
vtypedef arrayptr (a:vt@ype, n:int) =
[1:addr] arrayptr(a, 1, n)
```

• Just a pointer to some address, that's it

```
vtypedef arrayptr
...
1:addr = ptr(1)

^^^^^
```

• Parameterized on a linear viewtype & size (should be size\_t)

```
vtypedef arrayptr (a:vt@ype, n:int) =
...
```

• Returns an arrayptr to an existential (unknown) address type

```
l:addr = ptr(1)
vtypedef arrayptr =
[1:addr]
```

- Don't worry if this isn't clear
- Just a taste ...
- Tons type level concepts to learn!
- we'll get to some later . . .

• Now from the C side!

What resources are leaked?

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int));
  *i = 10;
  FILE* fp = fopen("test.txt","r");
  return 0;
}
```

Memory!

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int)); // <--- LEAK!!
  *i = 10;
  FILE* fp = fopen("test.txt","r");
  return 0;
}</pre>
```

File descriptor

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int)); // <--- LEAK!!
  *i = 10;
  FILE* fp = fopen("test.txt","r"); // <-- LEAK!!
  return 0;
}</pre>
```

• Equivalent ATS program

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
  free(pfset | i);
  fclose(pfFile | fp);
end
```

• "Client-facing" code, analogous, safe, this is why ATS is "fast"

• malloc produces a linear proof pf, consumed by ptr\_set

• ptr\_set produces a proof pfset

• fopen produces a proof of the file descriptor pfFile

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
  free(pfset | i);
  fclose(pfFile | fp);
end
```

• What happens when free and fopen are deleted?

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
```

end

• pfset is left unconsumed

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset <---
  val (pfFile | fp) = fopen("test.txt", "r")
in</pre>
```

end

• pfFile is left unconsumed

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset <---
  val (pfFile <---
in</pre>
```

end

- Free to write your all your code this way!
  - safe from buffer overflows & pointer bugs
  - ... there's sugar for implicitly passing proofs around
- Reuse decades of design sensibilities (safely!)
- But you're not benefitting from Functional Programming<sup>™</sup>...

- First "big" example
  - Read a number from the user between 1 and 10
  - Allocate an array of that length
  - Fill it.
  - Print it to console
  - Exit
- Doesn't seem like it but it's a LOT

```
fun read_input():Option_vt(a) = ...
fun make_array (len:int n): arrayptr = ...
implement main0() = begin
    println! ("Length of array? (1-10):");
    case+ read_input<int>() of
    | ~None_vt() => println! ("Not a number!")
    | ~Some_vt(len) =>
        if (len >= 1) * (len <= 10) then
          make_array(len)
        else println! ("Bad number!")
```

- Overall structure, types simplified
- Not too far from a functional program

• Simplified make\_array type signature

```
fun make_array (len:int n): arrayptr = ...
...
...
```

• Real make\_array type signature

```
fun make_array
  {n:int| n >= 1; n <= 10}
  (len:int n): [l:addr] arrayptr(int,1,n) =
   ...</pre>
```

• len is indexed with a refined int sort, n.

```
fun make_array
  {n:int| n >= 1; n <= 10} <-- refines it
  (len:int n): [1:addr] arrayptr(int,1,n) =</pre>
```

• Array pointer at *some* address

```
fun make_array
  {n:int| n >= 1; n <= 10}
  (len:int n): [l:addr] arrayptr(int,1,n) =</pre>
```

• Length between 1 & 10!

```
fun make_array
  {n:int| n >= 1; n <= 10}
  (len:int n): [l:addr] arrayptr(int,1,n) =</pre>
```

... being called here

how does it know {n:int| n >= 1; n <= 10}?!!</li>

• It statically understands runtime checks!

• Runtime checks translate to type constraints at compile time.

 Now anything in make\_array's call graph inherits the refinement