Why Linear Types Are The Future Of Systems Programming

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Outline

Introduction

Introduction

Standard Hello World

```
implement main0 = println! ("hello world")
```

• First from the ML side

• a linear list type

```
datatype list_vt (a:vt@ype, int) =
   | list_nil(a, 0) of ()
   | {n:int | n >= 0}
    list_cons(a, n+1) of (a, list(a, n))
```

• probably more familiar (_vt for viewtype)

• indexed on numbers, dependant types, just like (S(S(..)))

• numbers can be constrained, refinement types!

• parameterized on a view, linear types!

```
datatype list_vt (a:vt@ype, int) =
   | list_nil(a, 0) of ()
   | {n:int | n >= 0}
    list_cons(a, n+1) of (a, list(a, n))
```

• all ADT's are GADT's in ATS

```
datatype list_vt (a:vt@ype, int) =
   | list_nil(a, 0) of ()
   | {n:int | n >= 0}
    list_cons(a, n+1) of (a, list(a, n))
```

- tons type level concepts to learn!
- we'll get to some later . . .

• Now from the C side!

What resources are leaked?

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int));
  *i = 10;
  FILE* fp = fopen("test.txt","r");
  return 0;
}
```

Memory!

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int)); // <--- LEAK!!
  *i = 10;
  FILE* fp = fopen("test.txt","r");
  return 0;
}</pre>
```

File descriptor

```
int main(int argc, char** argv) {
  int* i = (int*)malloc(sizeof(int)); // <--- LEAK!!
  *i = 10;
  FILE* fp = fopen("test.txt","r"); // <-- LEAK!!
  return 0;
}</pre>
```

• Equivalent ATS program

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
  free(pfset | i);
  fclose(pfFile | fp);
end
```

• "Client-facing" code, analogous, safe, this is why ATS is "fast"

malloc produces a linear proof pf, consumed by ptr_set

• ptr_set produces a proof pfset

• fopen produces a proof of the file descriptor pfFile

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
  free(pfset | i);
  fclose(pfFile | fp);
end
```

• What happens when free and fopen are deleted?

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset | ()) = ptr_set(pf | i, 10)
  val (pfFile | fp) = fopen("test.txt", "r")
in
```

end

• pfset is left unconsumed

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset <---
  val (pfFile | fp) = fopen("test.txt", "r")
in</pre>
```

end

• pfFile is left unconsumed

```
implement main0 () = let
  val (pf | i) = malloc (sizeof<int>)
  val (pfset <---
  val (pfFile <---
in</pre>
```

end

- Free to write your all your code this way!
 - safe from buffer overflows & pointer bugs
 - ... there's sugar for implicitly passing proofs around
- Reuse decades of design sensibilities (safely!)
- But you're not benefitting from Functional Programming[™]...