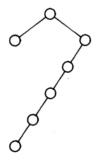
### Drawing Rooted Trees . . .

Aditya Siram (@deech)

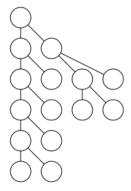
February 15, 2016

## Outline

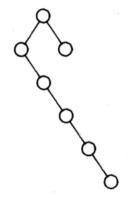
# Ugly Tree



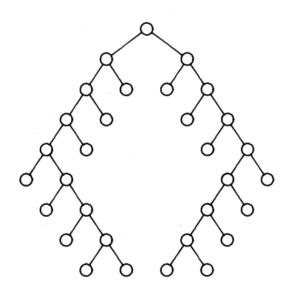
# Ugly Tree



### Nice Tree



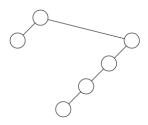
### Nice Tree



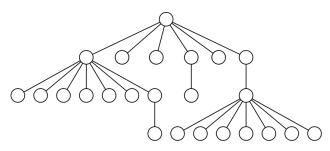


### Knuth

• Knuth (1971)



• 'nuf said.

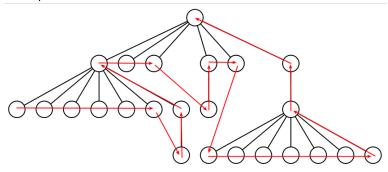


- Quadratic is easy
  - Update every node in case of conflicts
  - Undergraduate exercise
- Linear is hard
  - 25 years (1981 2006)
  - 33 years with Knuth (1971 2006)

- Two major players
- Walker (1989)
- Buccheim (2006)
  - Assisted by Reingold (1981)

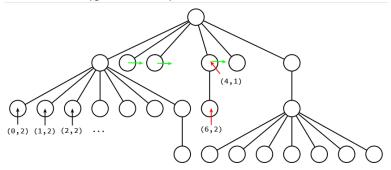
- Two passes
- First pass sets X coordinates naively, records conflicts.
- Second pass updates X coordinates accounting for conflicts.

Do a post order traversal

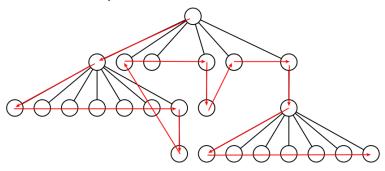


• This is an ugly tree

- First pass
- Save each node's X
- Record shifts (green arrows) but don't execute



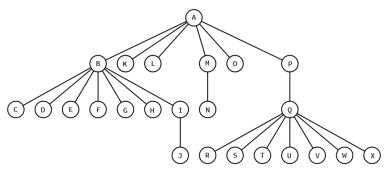
• Pre-order second pass



Add shifts to X-coordinate

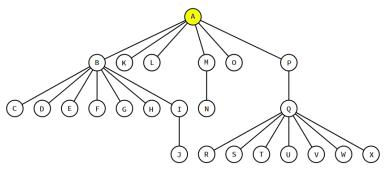
• Left neighbors

[]

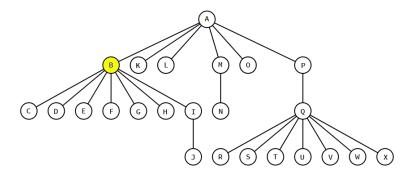


Left neighbors

[A]

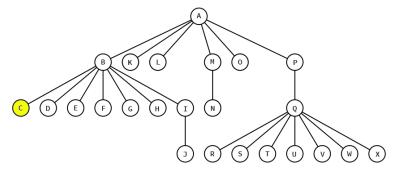


Left neighbors[A,B]



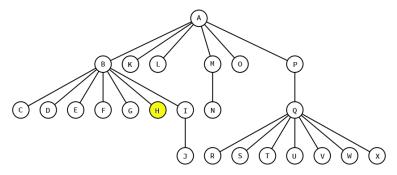
• Left neighbors

[A,B,C]

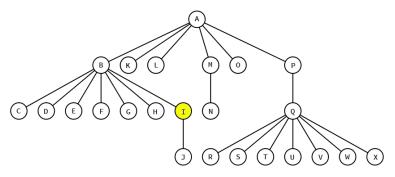


• Left neighbors

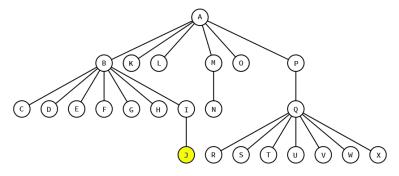
[A,B,H]



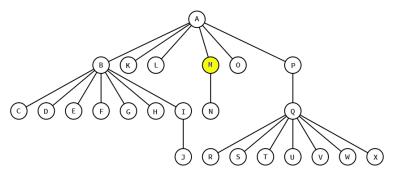
• Left neighbors



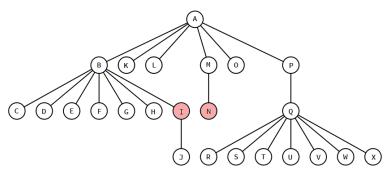
• Left neighbors



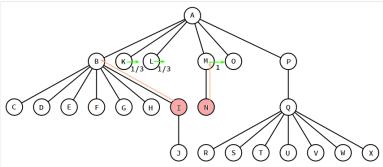
Left neighbors



• Left neighbors

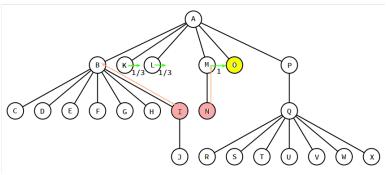


- APPORTION
- Track back to greatest common ancestors
- Record shifts in siblings between conflicting roots



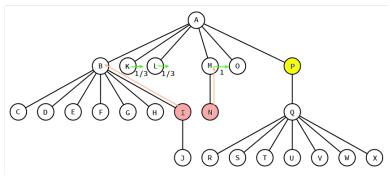
Left neighbors

[A,O,N,J]

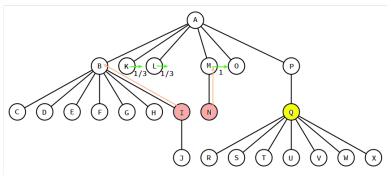


• Left neighbors

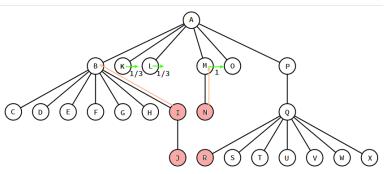
[A,P,N,J]



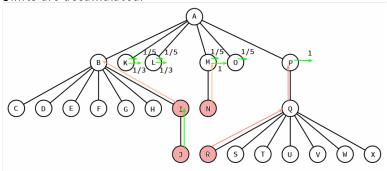
• Left neighbors



• Left neighbors



- APPORTION
- Track back to greatest common ancestors and shift
- Shifts are accumulated!

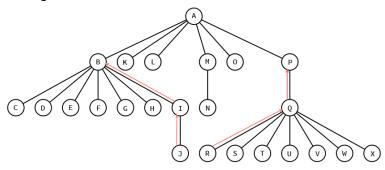


- Second pass, pre-order
- Shifts propogate to subtrees
- Pretty tree

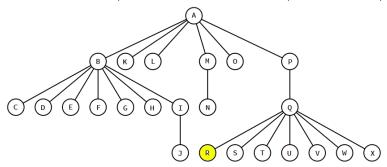
- Not linear!
- Discovered (and fixed) by Buccheim in 2006!
  - 16 years
- All the problems are in APPORTION.

### Buccheim

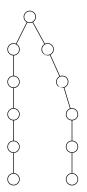
• Finding common ancestors is not linear!



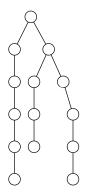
- Finding the leftmost node is also not linear
- Not as obvious
- Buccheim shows a pattern of trees that cause quadratic blowup.



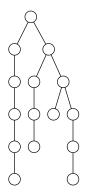
- Pick a k, say, 3
  - Hang two chains of (2 \* k) from root.



- (k (current level))-th odd number chains on right
  - current level = 1, k = 3, (3 1)th odd number = 3

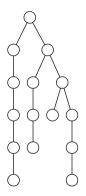


- (k (current level))-th odd number chains on right
  - current level = 2, k = 3, (3 2)th odd number = 1

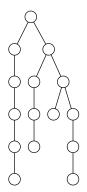


Total nodes = sum(k even numbers) + k

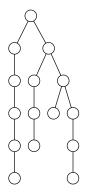
$$\bullet$$
 6 + 4 + 2 + 3 = 15



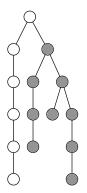
• Total nodes = k \* (k + 1) + k• 3 \* (3 + 1) + 3 = 15



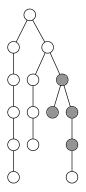
• Total nodes =  $k^2 + k + k = \theta(k^2)$ • 3 \* (3 + 1) + 3 = 15



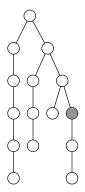
- Number of comparisons
  - level 1 = 9



- Number of comparisons
  - level 2 = 4



- Number of comparisons
  - level 3 = 1

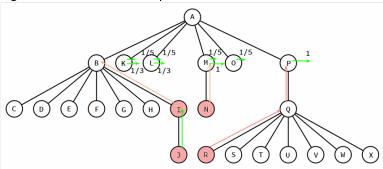


- Total # of calls
  - 1 + 4 + 9
- sum of k squares

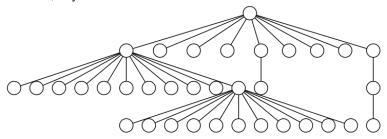
$$k^3/3 + k^2/2 + k/6$$

- $\theta(k^3)$
- nodes, k<sup>2</sup>
- calls, k<sup>3</sup>
- k, nodes<sup>(1/2)</sup>

- Shifting is also not linear!
- Again, Buccheim shows a pattern of trees

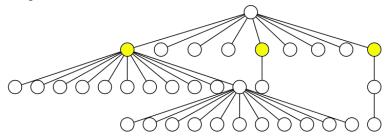


• Pick a k, say 3

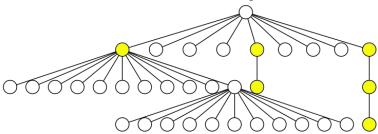


• Will cause quadratic shifting

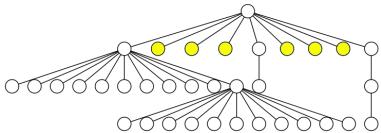
• Hang k nodes from root



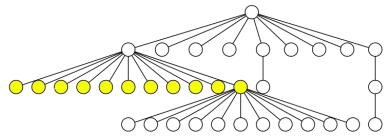
• Each of the k nodes is a chain k nodes high



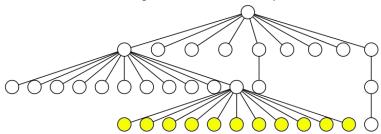
• Add k nodes in between those nodes



• Add 2k + 5 children to leftmost node



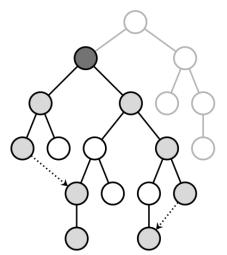
• Add 2k + 5 nodes to rightmost of leftmost upto k - 1



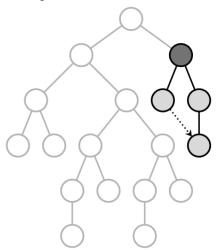
- Again shows k<sup>2</sup> nodes and k<sup>3</sup> runtime
- Similar procedure as before
- Omitted for time!

- Ancestor and leftmost node problem fixed by contours
- Introduced by Reingold (1981)
- Maintains the shape of the tree outlines, not left/rightmost nodes

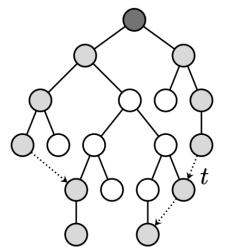
- Every inner/outer leaf stores a link to left/rightmost child
- Just re-uses the same space for left/right nodes



• Adding new subtree is added, it's inner links are now useless.



- Adds a link the rightmost child of the left tree
- If the right tree were bigger the left outer leaf adds a link

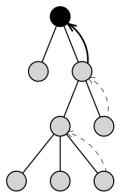


- At most 1 link added per subtree combination linear!
- To get the leftmost, just follow the left link linear!

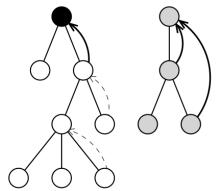
- Finding the common ancestors in linear time
- The ancestor is the right subtree is just the root

- Each node points to it's parent
- Also a pointer called defaultAncestor
- If parent isn't the common ancestor, default to defaultAncestor

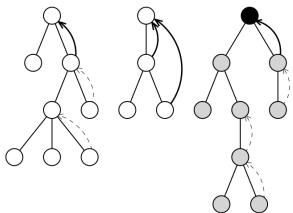
- Leftmost subtree set defaultAncestor to root.
- Solid arrows are correct links, dashed are expired links



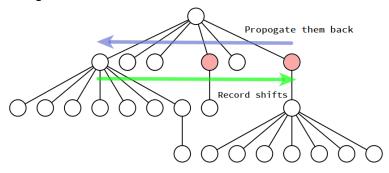
- When right is smaller, right contour is updated
- No expired links, always point to the correct ancestor



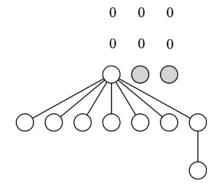
- When right is bigger, defaultAncestor now points to it's root
- The right shadows everything before it!



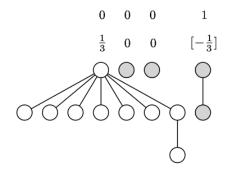
- Don't update every intermediate sibling for each conflict!
- Records shift, and change and aggregate them backwards once last sibling has been reached



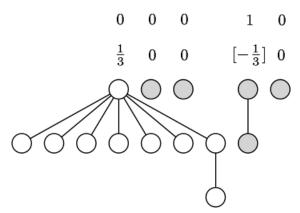
- Maintain pointer to leftmost sibling.
- Top number is *shift*, bottom is *change*.
- No conflicts so far



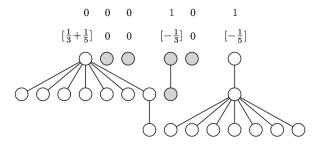
- Conflict!
- Shift is 1
- Change is 1/(number of intermediate subtrees + 1) = 1/3
- Add -1/3 on conflict node, add +1/3 to leftmost node
- This is confusing!



No clashes on the next one.

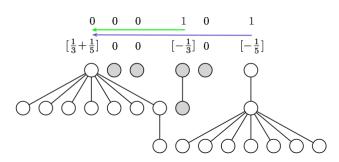


- Conflict!
- Shift is 1
- Change is 1/(number of intermediate subtrees + 1) = 1/5
- Add -1/5 on conflict node, accumulate +1/5 to leftmost node
- This is still confusing!



# Shifting (taking stock)

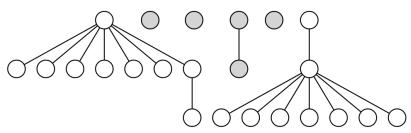
- shift of 1 from 6th sibling needs to be distributed back by fifths (blue line)
- shift of 1 from 4th sibling needs to be distributed back by thirds (green line)



- Top two numbers are the same as before
- Bottom two are the new shift and propagated change.
- Rightmost node does not move, shift is 0

0

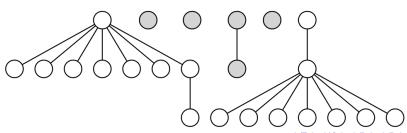
 $\left[-\frac{1}{5}\right]$ 



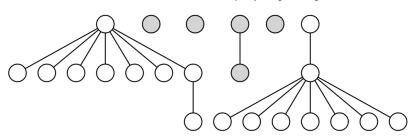
- Back one node, shift for this node is 1 1/5 = 4/5
- 4/5 propagates back

$$[\frac{1}{3} + \frac{1}{5}] \qquad 0 \qquad \qquad 0 \qquad [-\frac{1}{3}] \qquad 0 \qquad [-\frac{1}{5}]$$
 
$$\frac{4}{5} \qquad 0$$

$$\left[-\frac{1}{5}\right] \quad \left[-\frac{1}{5}\right]$$

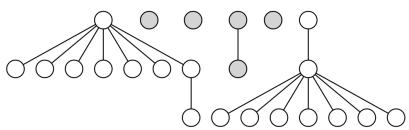


- Back another node, shift for this node is 4/5 1/5 = 3/5
- We're picking up another shift of 1 & a change of -1/3
- shift of 1 3/5 & change -1/5 1/3 propagates back!

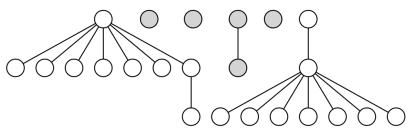


- No conflicts!
- (13/5) + (-1/5 1/3) = 2/3 + 2/5!
- Shift 2/3 + 2/5 and change (-1/5 1/3) propagates back!

0 0 0 1 0 1 
$$[\frac{1}{3} + \frac{1}{5}]$$
 0 0  $[-\frac{1}{3}]$  0  $[-\frac{1}{5}]$  
$$[\frac{2}{5} + \frac{2}{3}]$$
 
$$\frac{3}{5}$$
 
$$\frac{4}{5}$$
 0 
$$[-\frac{1}{5} - \frac{1}{3}][-\frac{1}{5} - \frac{1}{3}][-\frac{1}{5}]$$
 
$$[-\frac{1}{5}]$$



- No conflicts!
- (2/3 + 2/5) + (-1/5 1/3) = 1/5 + 1/3
- Shift (1/5 + 1/3), change (-1/5 1/3) propagates back!

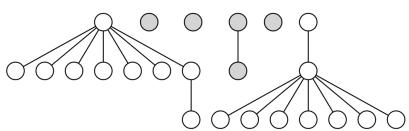


- No conflicts!
- The changes accumulated going left -> right was a stopper!
- (1/5 + 1/3) + (-1/5 1/3) = 0!

0 0 0 1 0 1 
$$[\frac{1}{3} + \frac{1}{5}] \quad 0 \quad 0 \quad [-\frac{1}{3}] \quad 0 \quad [-\frac{1}{5}]$$

$$0 \qquad \left[\frac{1}{5} + \frac{1}{3}\right] \quad \left[\frac{2}{5} + \frac{2}{3}\right] \qquad \frac{3}{5} \qquad \frac{4}{5} \qquad 0$$

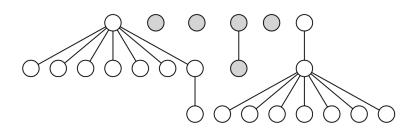
$$0 \qquad [-\frac{1}{5} - \frac{1}{3}][-\frac{1}{5} - \frac{1}{3}][-\frac{1}{5} - \frac{1}{3}][-\frac{1}{5}] \quad [-\frac{1}{5}]$$



## Second pass

- Pre-order, just like Walker
- We're done!

$$0 \qquad \left[\frac{1}{5} + \frac{1}{3}\right] \quad \left[\frac{2}{5} + \frac{2}{3}\right] \qquad \frac{3}{5} \qquad \frac{4}{5} \qquad 0$$



#### Resources

- Bill Mill's page .. http://billmill.org/pymag-trees/
- Reingold's paper ... http://reingold.co/tidier-drawings.pdf
- Walker's paper ... http://www.cs.unc.edu/techreports/89-034.pdf
- Buccheim's paper ...
  https://github.com/tristanpenman/n-puzzle/blob/master/
  doc/reference/Drawing%20Rooted%20Trees%20in%20Linear%
  20Time.pdf