

Dynamic Programming

Algorithms: Design and Analysis, Part II

Sequence Alignment
Optimal Substructure

Problem Definition

Recall: Sequence alignment. [Needleman-Wunsch score = Similarity measure between strings]

Input: Strings $X = x_1 \dots x_m$, $Y = y_1 \dots y_n$ over some alphabet Σ (like $\{A,C,G,T\}$)

- Penalty $\alpha_{\rm gap}$ for inserting a gap, α_{ab} for matching a & b [presumably $\alpha_{ab} = 0$ of a = b]

Feasible solutions: Alignments - i.e., insert gaps to equalize lengths of the string

Goal: Alignment with minimum possible total penalty

A Dynamic Programming Approach

Key step: Identify subproblems. As usual, will look at structure of an optimal solution for clues.

[i.e., develop a recurrence + then reverse engineer the subproblems]

Structure of optimal solution: Consider an optimal alignment of X, Y and its final position:



Question: How many <u>relevant</u> possibilities are there for the contents of the final position?

- A) 2 C) 4 B) 3 D) *mn*
- Case 1: x_m , y_n matched, case 2: x_m matched with a gap, case 3: y_n matched with a gap [Pointless to have 2 gaps]

Optimal Substructure

Point: Narrow optimal solution down to 3 candidates.

Optimal substructure: Let
$$X' = X - x_m$$
, $Y' = Y - y_n$.

If case (1) holds, then induced alignment of X' & Y' is optimal. If case (2) holds, then induced alignment of X' & Y is optimal. If case (3) holds, then induced alignment of X & Y' is optimal.

Optimal Substructure (Proof)

Proof: [of Case 1, other cases are similar]

By contradiction. Suppose induced alignment of X', Y' has penalty P while some other one has penalty $P^* < P$.

 $\Rightarrow \text{Appending} \quad \frac{x_m}{y_n} \quad \text{to the latter, get an alignment of X and Y}$ with penalty $P^* / + \alpha_{x_m y_n} < P + \alpha_{x_m y_n}$

Contents of final position Penalty of original alignment

 \Rightarrow Contradicts optimality of original alignment of X & Y. QED!