Closest Pair of Points using Divide and Conquer algorithm

We are given an array of n points in the plane, and the problem is to find out the closest pair of points in the array. This problem arises in a number of applications. For example, in air-traffic control, you may want to monitor planes that come too close together, since this may indicate a possible collision. Recall the following formula for distance between two points p and q.

The Brute force solution is $O(n^2)$, compute the distance between each pair and return the smallest. We can calculate the smallest distance in O(nLogn) time using Divide and Conquer strategy. In this post, a $O(n \times (Logn)^2)$ approach is discussed. We will be discussing a O(nLogn) approach in a separate post.

Algorithm

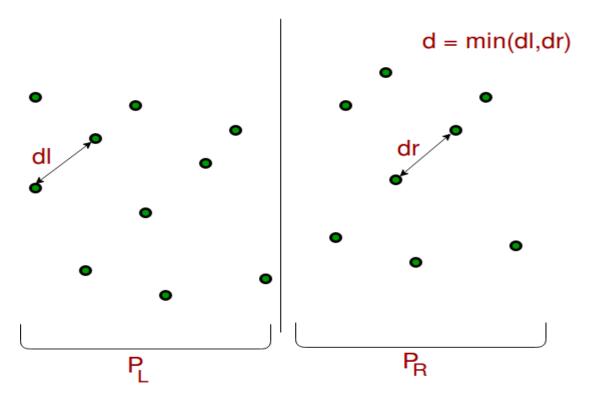
Following are the detailed steps of a O(n (Logn)^2) algortihm.

Input: An array of n points *P[]*

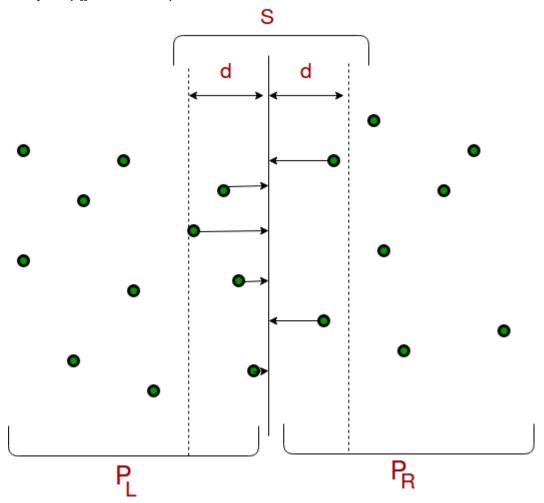
Output: The smallest distance between two points in the given array.

As a pre-processing step, the input array is sorted according to x coordinates.

- 1) Find the middle point in the sorted array, we can take P[n/2] as middle point.
- 2) Divide the given array in two halves. The first subarray contains points from P[0] to P[n/2]. The second subarray contains points from P[n/2+1] to P[n-1].
- 3) Recursively find the smallest distances in both subarrays. Let the distances be dl and dr. Find the minimum of dl and dr. Let the minimum be d.



4) From the above 3 steps, we have an upper bound d of minimum distance. Now we need to consider the pairs such that one point in pair is from the left half and the other is from the right half. Consider the vertical line passing through P[n/2] and find all points whose x coordinate is closer than d to the middle vertical line. Build an array strip[] of all such points.



- **5)** Sort the array strip[] according to y coordinates. This step is O(nLogn). It can be optimized to O(n) by recursively sorting and merging.
- **6)** Find the smallest distance in strip[]. This is tricky. From the first look, it seems to be a O(n^2) step, but it is actually O(n). It can be proved geometrically that for every point in the strip, we only need to check at most 7 points after it (note that strip is sorted according to Y coordinate). See <u>this</u> for more analysis.
- 7) Finally return the minimum of d and distance calculated in the above step (step 6)

```
import java.util.ArrayList;
import java.util.Arrays;
import java.util.Comparator;
import java.util.List;
public class ClosestPairOfPoint {
        public static void mindistance(List<Point> list) throws
IllegalArgumentException{
IllegalArgumentException("We need atleast 2 points");
                if(list.get(i)==null)
                    throw new IllegalArgumentException("Point is not
            Arrays.sort(pointsbyX, new Comparator<Point>() {
                @Override
                public int compare(Point o1, Point o2) {
                    if(o1.x!=o2.x)
                if(pointsbyX[i]==pointsbyX[i+1]){
                    break;
            for (int i = 0; i < n; i++)
```

```
both in right subarray
            double delta1 = closest(pointsByX, pointsByY, aux, lo, mid);
v-coordinate
than delta
                // a geometric packing argument shows that this loop iterates
at most 7 times
                for (int j = i+1; (j < m) && (aux[j].y - aux[i].<math>y < delta);
                            min = delta;
        private static void merge(Point[] a, Point[] aux, int lo, int mid, int
                aux[k] = a[k]:
```

```
else if (less(aux[j], aux[i])) a[k] = aux[j++];
                else
        private static boolean less(Point v, Point w) {
            return v.x<w.x;
            return Math.sqrt(x*x+y*y);
        public static void main(String[] args) {
            Point p6 = new Point(3,4);
            List<Point> list = new ArrayList<>();
// BruteForce Approach
public class ClosestPairOfPoint {
    public static class Point {
        public Point(int x, int y) {
            this.y = y;
```

```
public List<Point> mindistance(List<Point> list){
        if(list==null || list.size()==0) return new ArrayList<>();
        List<Point> result = new ArrayList<>();
        int n = list.size();
        int min = Integer.MAX_VALUE;
        for(int i=0;i<n;i++){</pre>
            for(int j=i+1;j<n;j++){
                int distance = getDistance(list.get(i), list.get(j));
                if(min>distance | result.size()==0){
                     result.clear();
                     result.add(list.get(i));
                     result.add(list.get(j));
                    min = distance;
        return result;
    public int getDistance(Point a, Point b){
        int y = a.y - b.y;
        return x*x + y*y;
    public void display(List<Point> list){
        Point a = list.get(0);
        Point b = list.get(1);
        System.out.print("{"+a.x+" "+a.y+"}");
System.out.print("{"+b.x+" "+b.y+"}");
        System.out.println(" ");
    public static void main(String[] args) {
        ClosestPairOfPoint m = new ClosestPairOfPoint();
        Point p1 = new Point(2,3);
        Point p2 = new Point(12,30);
        Point p3 = new Point(40,50);
        Point p4 = new Point(5,1);
        Point p5 = new Point(12,10);
        Point p6 = new Point(3,4);
        List<Point> list = new ArrayList<>();
        list.add(p1); list.add(p2); list.add(p3); list.add(p4); list.add(p5);
list.add(p6);
        System.out.println("The closest pair of points are ");
        List<Point> result = m.mindistance(list);
        m.display(result);
        System.out.println("The distance between the closest points is
'+Math.sqrt(m.getDistance(result.get(0), result.get(1))));
```