



Your Ultimate Guide To Landing
Top AI roles



DECODE
AiML

2.3.9

Tower of Hanoi

Problem: Tower of Hanoi Puzzle consists of 3 towers namely Source, Auxiliary and Destination. you have n disks arranged in the Source tower in such a way that smaller disk is kept over larger disks.

Goal: you need to move n disks from src tower to dest tower using Aux tower by below rules

- ① Only 1 disks can be moved at a time
- ② No disks can be placed on top of smaller disks.

For $n=3$



```
def TOH(n, src, Target, aux):
```

Base Case

Recursive Case

Base Case



```
def TnH(n, src, Target, Aux)  
    → Move 1 from A to C
```

Recursive Case

```
def TOH(n, src, Target, Aux)
```

→ Move 1 from A to B $\Rightarrow \text{TOH}(1, A, B, C)$
→ Move 2 from A to C
→ Move 1 from B to C $\Rightarrow \text{TOH}(1, B, C, A)$

Recursive Case



```
def TOH(3, src, Target, Aux)
```

- Move 1 from A to C
- Move 2 from A to B
- Move 1 from C to B
- move 3 from A to C
- move 1 from B to A
- move 2 from B to C
- move 1 from A to C

$\text{TOH}(2, \text{A}, \text{B}, \text{C})$

$\text{TOH}(2, \text{B}, \text{C}, \text{A})$

```
def TOH(n, src, Target, aux):
```

```
    if n == 1:
```

```
        print(f"Move 1 disk from {src} to {Target}");
```

Base
case

```
TOH(n-1, src, Aux, Target)
```

```
print(f"Move {n} disk from {src} to {target}")
```

```
TOH(n-1, Aux, Target, src)
```

Recursive
case

Time Complexity

def TOH(n, src, Target, aux): → T(n)

if n == 1:

 print(f"Move 1 disk from {src} to {Target}");

} O(1)

TOH(n-1, src, Aux, Target) → T(n-1)

print(f"Move {n} disk from {src} to {target}"); → O(1)

TOH(n-1, Aux, Target, src) → T(n-1)

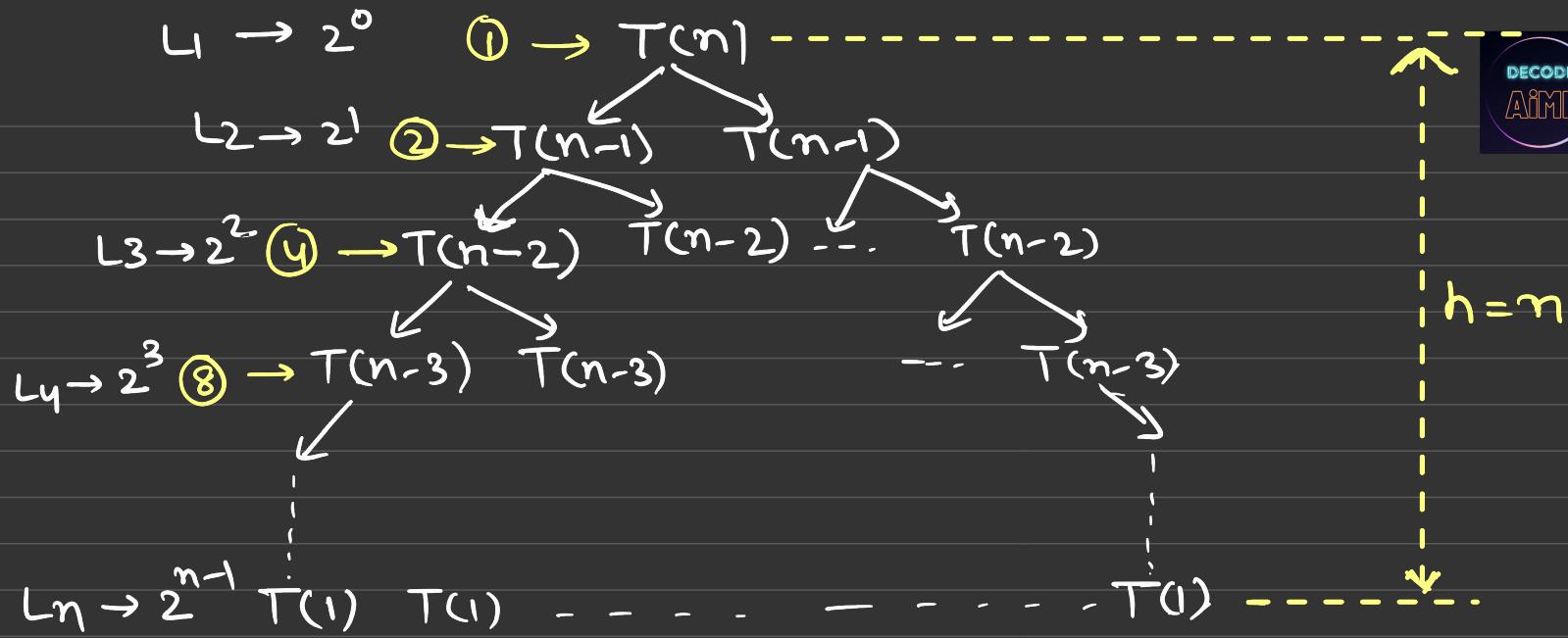
$$T(n) = O(1) + 2 \cdot T(n-1) + O(1)$$

$$T(n) = 2T(n-1) + O(1)$$

$$T(1) = O(1)$$

← Recurrence Relation.

→ Let's Solve this using Recursion Tree



Total # Recursion call = $2^0 + 2^1 + 2^2 + \dots + 2^{n-1}$

$$[GP] \quad S_n = \frac{a(r^n - 1)}{r - 1} = \frac{2 \cdot (2^{n-1} - 1)}{2 - 1} = \boxed{\frac{2^n - 2}{2}}$$

Total time Complexity = # Recursion Calls \times Time at each rec. call

$$= (2^n - 2) \times O(1)$$

$$= (2^n - 2) \times C$$

$$= C \cdot 2^n - 2C$$

$$= O(2^n)$$

$$T(n) = O(2^n)$$



Space Complexity

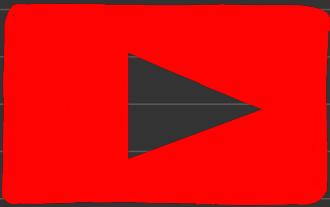
→ Input Space = $O(1)$

Auxiliary Space = Max depth of Recursion
= $O(n)$ [From Recursion Tree]

Space Complexity = $O(n)$

Total Space Complexity = $O(n)$

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