

**Your Ultimate Guide To Landing
Top AI roles**



2.12

Hash Table Implementation in Python

→ A hash table is a data structure that stores key-value pairs and allows fast access (search) to values based on their keys.

↳ Amortized $O(1)$ time.

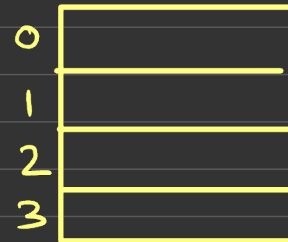
→ In Python, `dict` and `set` are implemented using Hash Table data structure

→ Each element in set is stored as a dict key with a dummy value (usually None)

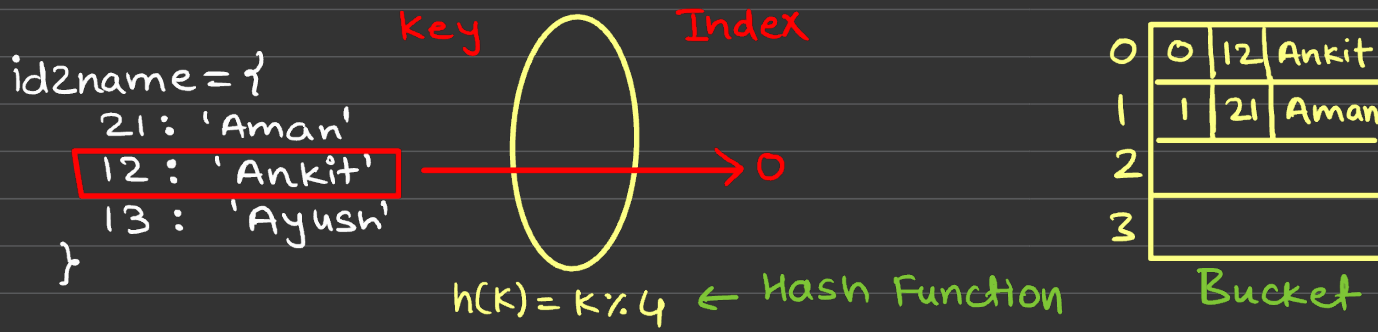
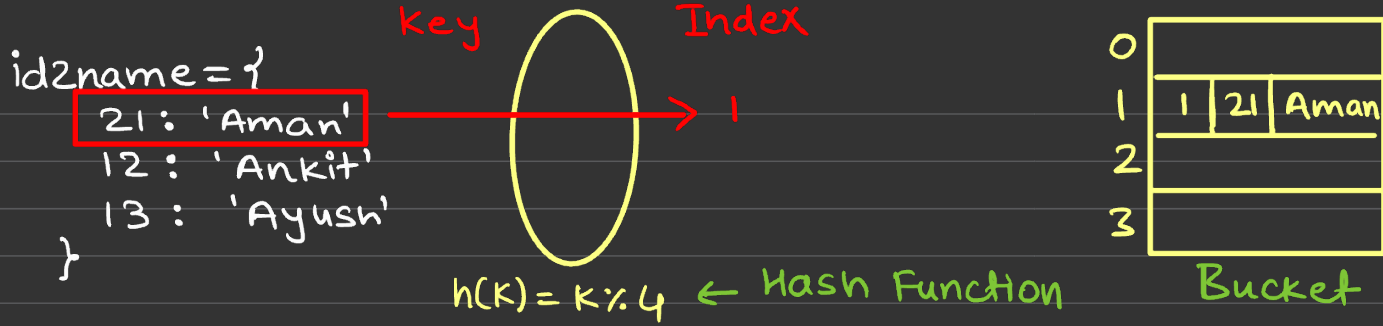
```
id2name = {  
    21: 'Aman'  
    12: 'Ankit'  
    13: 'Ayush'  
}
```

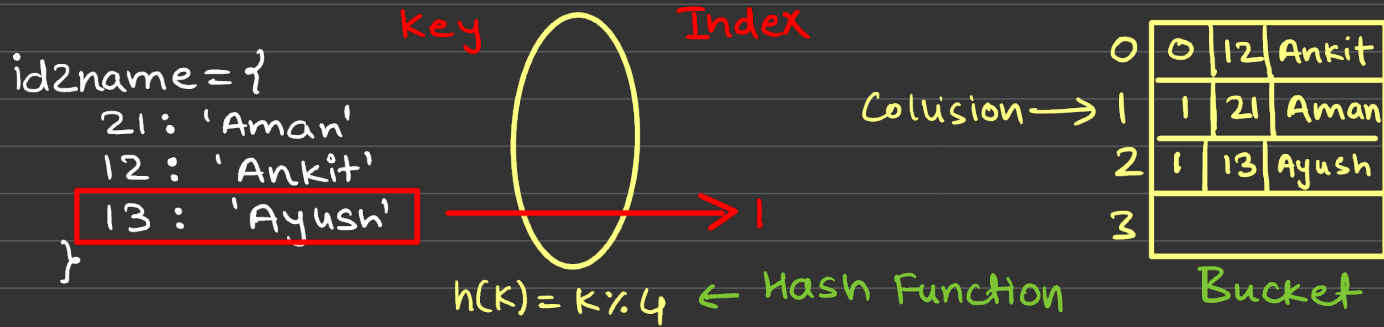


$h(K) = K \% 4$ ← Hash Function



Bucket





→ Collision happens when we try to insert value with key 13

→ So we use Quadratic probing in double hashing.

$$h'(k, i) = (k + c_1 \cdot i + c_2 \cdot i^2) \% 4$$

if $c_1 = 0$ and $c_2 = 1$

$$h'(13, 1) = (13 + 1^2) \% 4 = 2$$

→ Python's Hash Table Implementation details

- ① Uses Open Addressing (Quadratic Probing) to resolve Collision
- ② Bucket stores entry directly in the array.
- ③ A bucket stores
 - ① dict $\rightarrow \langle \text{hash}, \text{key}, \text{value} \rangle$
 - ② set $\rightarrow \langle \text{hash}, \text{key} \rangle$
- ④ Uses dynamic resizing when load factor gets too high ($\sim 2/3$)
 ↳ during Insertion.
- ⑤ Hashing, Probing and Resizing logic of set is same as a dict.

Operations in Hash Table

Expected no of probes $= \frac{1}{1-\alpha}$

$$T(n) = O\left(\frac{1}{1-\alpha}\right)$$



Operation	Amortized $T(n)$	Worst $T(n)$
Insertion	$O(1)$	$O(n)$
Search	$O(1)$	$O(n)$
Delete	$O(1)$	$O(n)$
Update	$O(1)$	$O(n)$

$$\alpha = \frac{n}{m}$$

If α is constant (~ 0.66)

$$\hookrightarrow T(n) = O(1)$$

→ In hash table, the expensive operation is resizing.

↳ When the load factor (α) pass a threshold, double size table.

↳ Resizing requires rehashing all keys $\rightarrow O(n)$

↳ But, resizing is a rare operation.

↳ Amortized cost = $O(1)$

Hash Table using Dict

→ Refer Lecture 1.8.3 (Dictionary in one shot)

→ for a dict item, the values can be any datatype, but key should be Hashable, and immutable

→ Initialization $\leftarrow O(1)$

```
my_dict = {}
```

→ Insertion $\leftarrow O(1)$ Amortized

```
my_dict["name"] = "Sanjeev"  
my_dict["age"] = 26
```

→ Search $\leftarrow O(1)$ Amortized

```
my_dict["name"]
```

When the load factor stays below a threshold ($\approx 2/3$), Collision are rare

→ Updation $\leftarrow O(1)$ Amortized

```
my_dict["age"] = 26
```

→ Deletion $\leftarrow O(1)$ Amortized

```
del my_dict["name"]
```


Hash Table using Set

- Refer Lecture 1.8.2 (Sets in one shot)
- Each element in set is stored as a dict key with a dummy value (usually None)
- for a set datatype, the key should be hashable and mutable.
- Initialization $\leftarrow O(1)$

```
my_set = Set()
```

- Insertion $\leftarrow O(1)$ Amortized

```
my_set.add("Aman")  
my_set.add("Ankit")
```

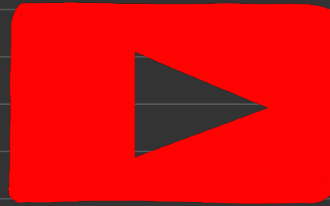
→ Search $\leftarrow O(1)$ Amortized

```
"Aman" in my_set
```

→ Deletion $\leftarrow O(1)$ Amortized

```
my-set.remove("Aman")
```

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