



**Your Ultimate Guide To Landing
Top AI roles**



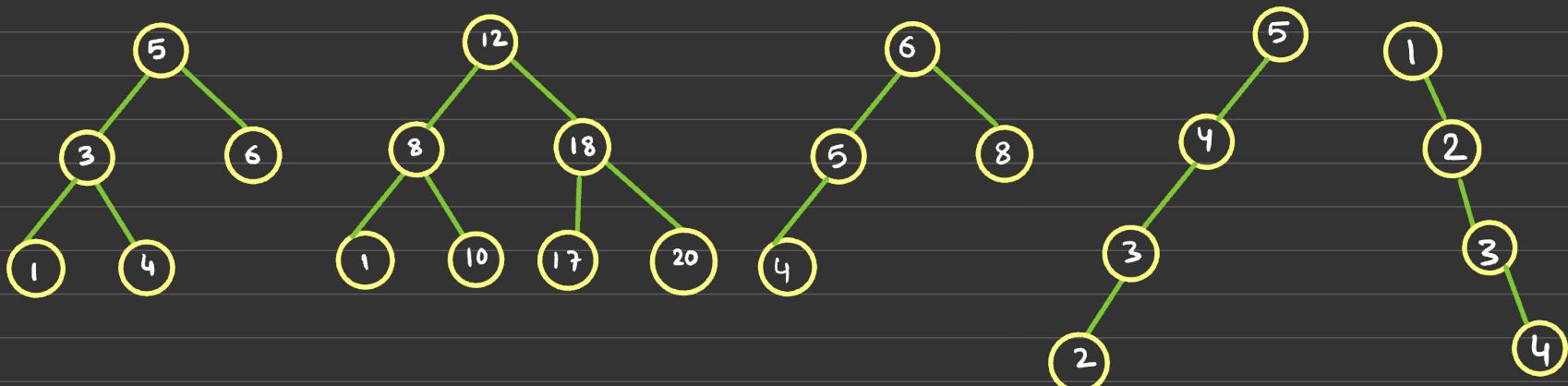
**DECODE
AiML**

2.16.1

Binary Search Tree



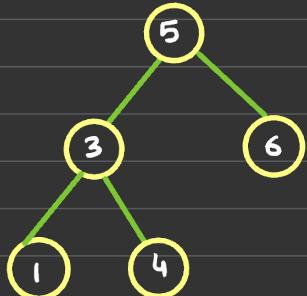
- A binary Search Tree (BST) is a binary tree such that
 - ① all left subtree elements should be less than root data
 - ② all right subtree elements should be greater than root data.
- This property should be satisfied at every node in the Tree.



Properties of BST

- If we perform Inorder Traversal of a BST, it will get all the elements in increasing order
- In a BST, leftmost element will be the least and rightmost element will be the greatest element.
- No of BST possible with n distinct keys

$$\# \text{BST} = \frac{2^n C_n}{n+1}$$



Inorder: 1, 3, 4, 5, 6

BST Implementation in Python

Class TreeNode:

```
def __init__(self, value):
    self.value = value
    self.left = None
    self.right = None
```

Value		
Left R	right R	

Node structure

Key Operations in a BST

① Insertion

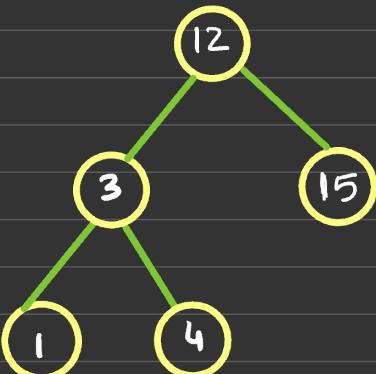
→ Insert a key in a BST

② Search

→ Search for a key in a BST

③ Deletion

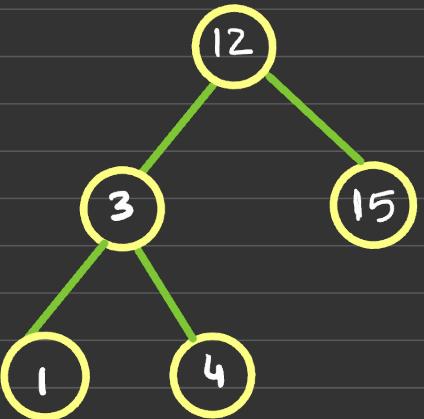
→ Delete a key in a BST



Search in a BST

```
def search(root, key):  
    if root is None:  
        return root  
    if root.key == key:  
        return root  
    if key > root.key:  
        return search(root.right, key)  
    return search(root.left, key)
```

→ Search Key → 4

* Time & Space Complexity

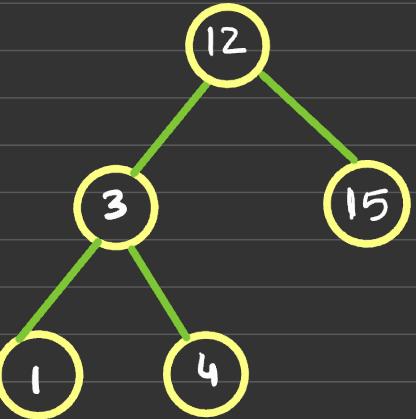
Time Complexity = $O(n)$

Space Complexity = $O(n)$

Insertion in a BST

```
def insert(root, key):
    if root is None:
        return Node(key)
    if key < root.key:
        root.left = insert(root.left, key)
    elif key > root.key:
        root.right = insert(root.right, key)
    return root
```

→ Insert Key → 8



* Time & Space Complexity

Time Complexity = $O(n)$

Space Complexity = $O(n)$

Deletion in a BST

→ When deleting a node, there are 3 Cases

→ Delete key → 18

① Node has no children (leaf node)

→ Just remove it

② Node has 1 child

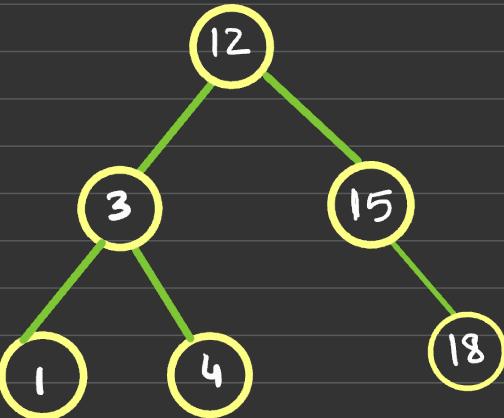
→ Replace the node with its child

③ Node has 2 children

→ Find the inorder Successor (smallest node in right subtree)
or inorder predecessor (largest in left subtree)

→ Swap the value with the target Node.

→ Delete that Successor / Predecessor.



Deletion in a BST

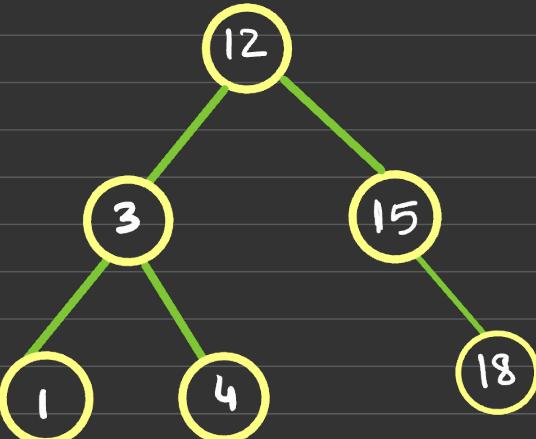
```

def deleteNode(root, key):
    if root is None:
        return root
    if key < root.key:
        root.left = deleteNode(root.left, key)
    elif key > root.key:
        root.right = deleteNode(root.right, key)
    else: #node found
        if root.left == None:
            return root.right
        elif root.right == None:
            return root.left
        temp = inorderPredecessor(root.left)
        root.key = temp.key
        root.left = deleteNode(root.left, temp.key)
    return root

```

Case I or II

Case III



* Time & Space Complexity

Time Complexity = $O(n)$

Space Complexity = $O(n)$

Binary Tree Vs Binary Search Tree

→ If Tree is balanced. means height (h) = $\log n$.

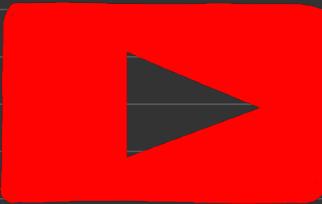
	Binary Tree	BST
Insert	$O(n)$	$\log(n)$
Search	$O(n)$	$\log(n)$
Delete	$O(n)$	$\log(n)$

Next Lecture

→ AVL Tree



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