



**Your Ultimate Guide To Landing
Top AI roles**



**DECODE
AiML**

2.12

Hash Table Implementation in Python

- A hash table is a data structure that stores key-value pairs and allows fast access (search) to values based on their keys.
 - ↳ Amortized $O(1)$ time.
- In Python, `dict` and `set` are implemented using Hash Table data structure.
- Each element in set is stored as a dict key with a dummy value (usually None)

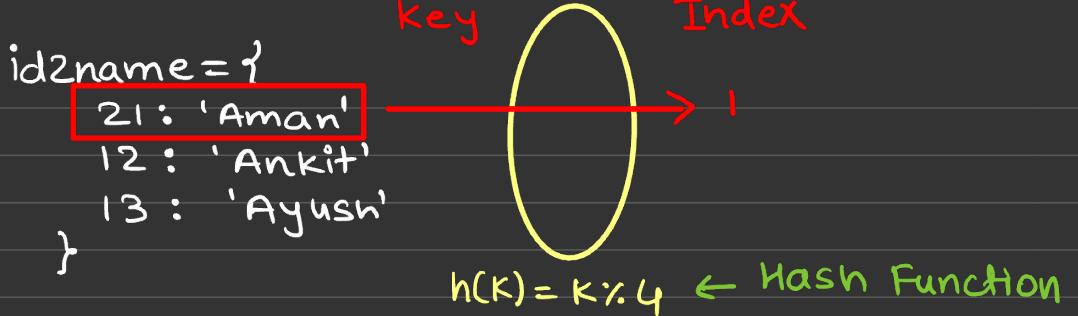
```
id2name = {  
    21: 'Aman'  
    12: 'Ankit'  
    13: 'Ayush'  
}
```



$h(K) = K \% 4 \leftarrow \text{Hash Function}$

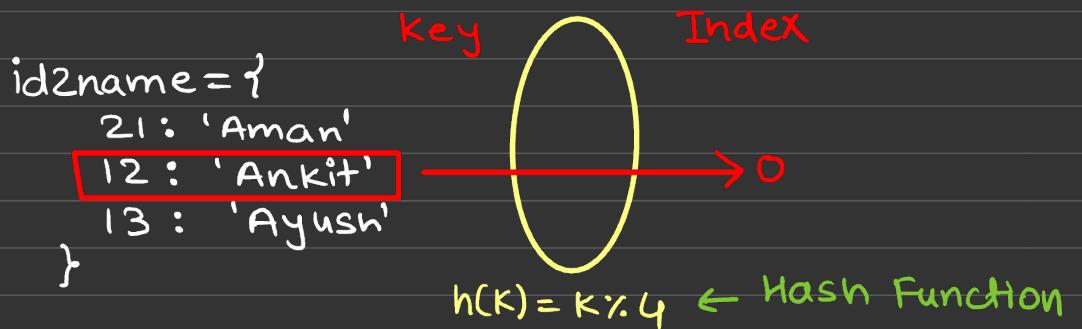
0	
1	
2	
3	

Bucket



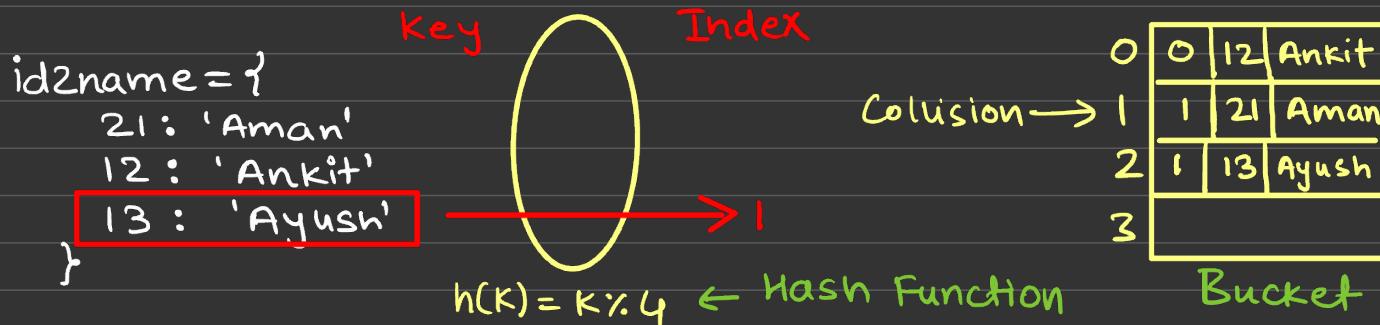
0			
1	1	21	Aman
2			
3			

Bucket



0	0	12	Ankit
1	1	21	Aman
2			
3			

Bucket



→ Collision happens when we try to insert value with key 13

→ So we use Quadratic probing in double hashing.

$$h'(K, i) = (K + c_1 \cdot i + c_2 \cdot i^2) \% 4 \quad \text{if } c_1=0 \text{ and } c_2=1$$

$$h'(13, 1) = (13 + 1^2) \% 4 = 2$$

→ Python's Hash Table Implementation details

- ① uses Open Addressing (Quadratic Probing) to resolve Collision
- ② Bucket stores entry directly in the array.
- ③ A bucket stores
 - ① dict → <hash, key, value>
 - ② set → <hash, key>
- ④ uses dynamic resizing when load factor gets too high ($\sim 2/3$)
 - ↳ during insertion.
- ⑤ Hashing, Probing and Resizing logic of set is same as a dict.

Operations in Hash Table

Expected no of probes = $\frac{1}{\alpha}$



$$T(n) = O\left(\frac{1}{\alpha}\right)$$

Operation	Amortized T(n)	Worst T(n)
Insertion	O(1)	O(n)
Search	O(1)	O(n)
Delete	O(1)	O(n)
Update	O(1)	O(n)

$$\alpha = \frac{n}{m}$$

If α is constant (~ 0.66)

$$\hookrightarrow T(n) = O(1)$$

→ In hash table, the expensive operation is resizing.

↳ When the load factor (α) pass a threshold, double size table.

↳ Resizing requires rehashing all keys $\rightarrow O(n)$

↳ But, resizing is a rare operation.

↳ Amortized cost = O(1)

Hash Table using Dict

- Refer Lecture 1.8.3 (Dictionary in one shot)
- for a dict item, the values can be any datatype, but key should be Hashable. and immutable

→ Initialization $\leftarrow O(1)$

```
my_dict = {}
```

→ Insertion $\leftarrow O(1)$ Amortized

```
my_dict["name"] = "Sanjeev"  
my_dict["age"] = 26
```

→ Search ← O(1) Amortized

my-dict["name"]

When the load factor stays below a threshold ($\approx 2/3$), Colision are rare

→ Updation ← O(1) Amortized

my-dict["age"] = 26

→ Deletion ← O(1) Amortized

del my-dict["name"]

Hash Table using Set

- Refer Lecture 1.8.2 (Sets in one shot)
- Each element in set is stored as a dict key with a dummy value (usually None)
- for a set datatype, the key should be hashable and mutable.

→ Initialization $\leftarrow O(1)$

```
my_set = set()
```

→ Insertion $\leftarrow O(1)$ Amortized

```
my_set.add("Aman")
my_set.add("Ankit")
```

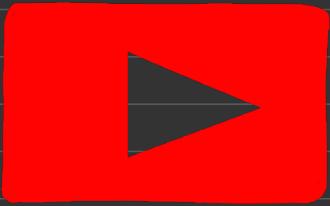
→ Search ← O(1) Amortized

"Aman" in my_set

→ Deletion ← O(1) Amortized

my_set.remove ("Aman")

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