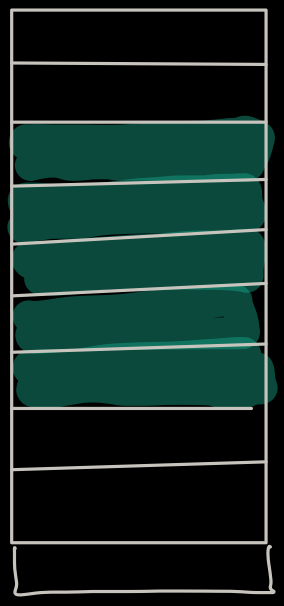
Advantage of Array

↳ array / arrayList

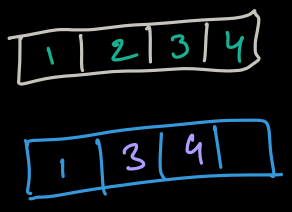
$A[i] \rightarrow O(1)$ — O(1) random access

```
int [] a = new int [5];
int [] b = new int [5] x
```

Don't have enough memory



	At start	At end	At random (k^{th} pos)
Insertion	O(n)	O(1)	O(n)
Deletion	O(n)	O(1)	O(n)
Update	O(1)	O(1)	O(1)



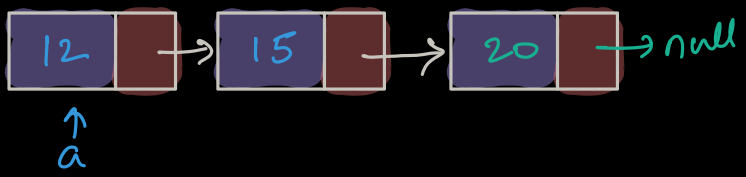
inverted index

word → page No.

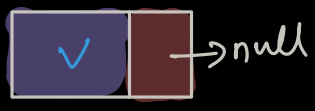
- milk → 47
- phone → [_ _]
- iphone → [. . .]
- iph → [. . . .]

Linked List

```
class Node {
    int val;
    Node next;
    Node(int v) {
        this.val = v
        this.next = null
    }
}
```




```
Node a = new Node(12)
a.next = new Node(15)
a.next.next = new Node(20)
```



* this: reference variable that refers to current object.

```
class pair {
    int x, y;
    pair(int x, int y) {
        x = x
        y = y
    }
}
```

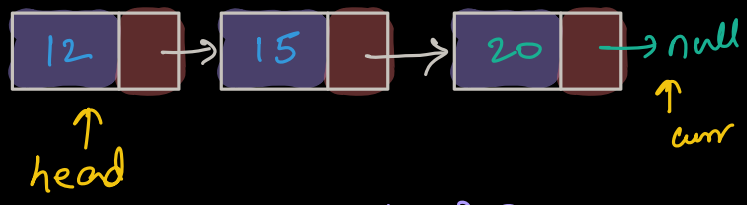
x, y = 0, 0

```
class pair {
    int x, y
    pair(int x, int y) {
        this.x = x
        this.y = y
    }
}
```

```
pair(int a, int b) {
    this.x = a
    this.y = 5
}
```

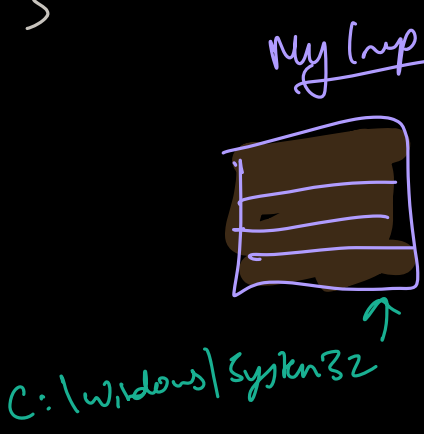
Print the Linked List

```
void printLL(Node head) {
    Node curr = head
    while(curr != null) {
        print(curr.val)
        curr = curr.next
    }
}
```



12 15 20

TC: O(n)



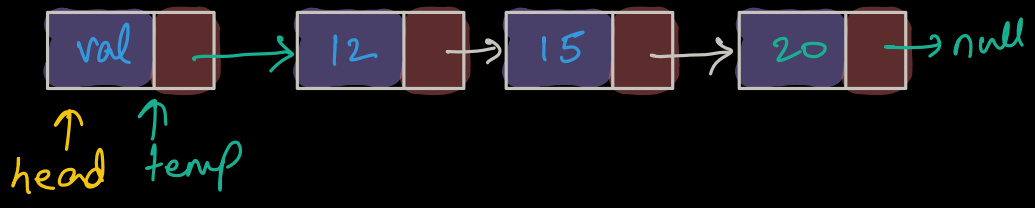
10 GB

cut.

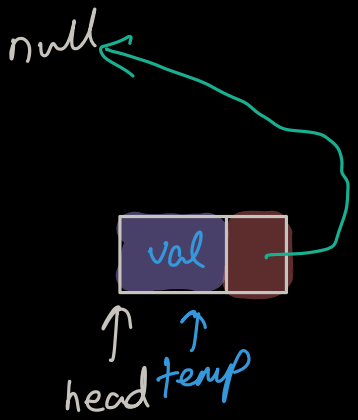
instant!

Insertion at start in LL

```
Node insertAtStart(head, val) {
    Node temp = new Node(val)
    temp.next = head
    head = temp
    return head
}
```

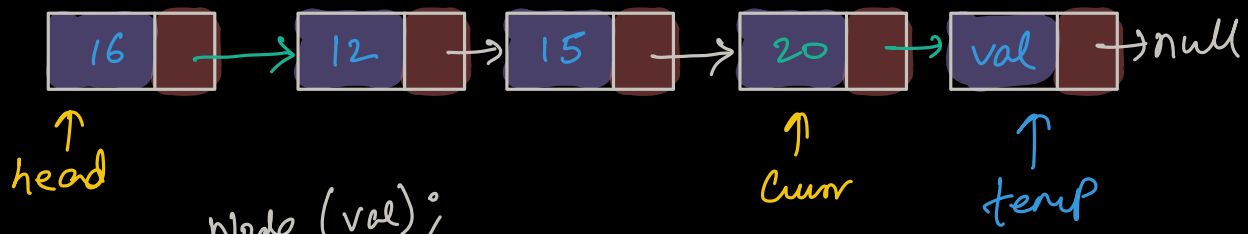


TC: O(1)

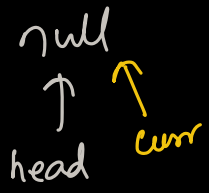


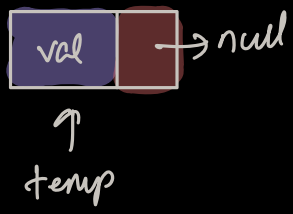
Insert at the end

```
Node insertAtEnd(head, val) {
```




```
    Node temp = new Node(val);
    if(head == null) {
        head = temp
        return head
    }
    else {
        Node Curr = head
        while(Curr.next != null) {
            Curr = Curr.next
        }
        Curr.next = temp
        return head
    }
}
```

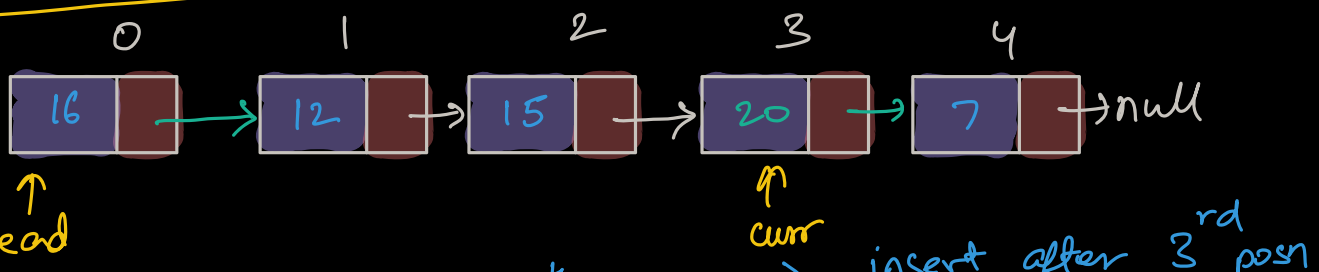




TC: O(n)

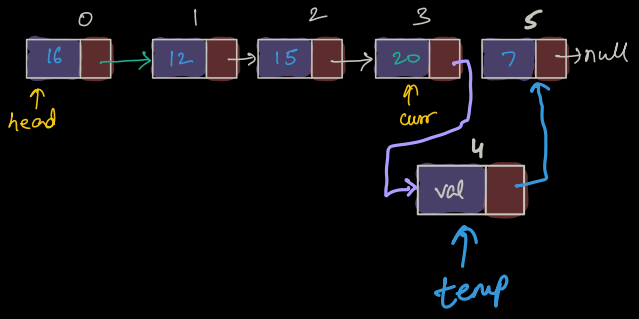
Break: till 8:43 am.

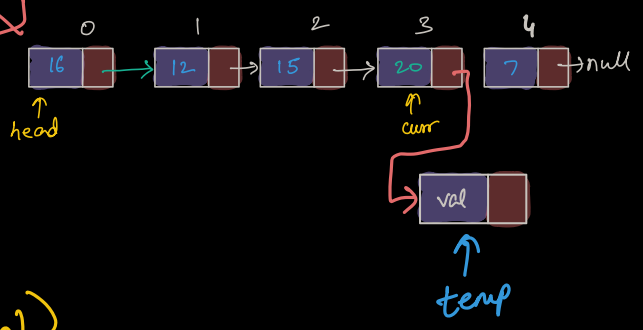
Insert at k^{th} pos



K=4 insert @ 4th pos ⟹ insert after 3rd posn

```
Node insertAtKPos(head, val, K) {   // K=N, code works.
    if(K==0) {
        head = InsertAtStart(head, val)
        return head
    }
    else {
        count=0
        Node temp = new Node(val);
        Node curr = head
        while(count < K-1) {
            curr = curr.next
            count++
        }
        curr.next = temp
        temp.next = curr.next
        curr.next = temp
    }
}
```

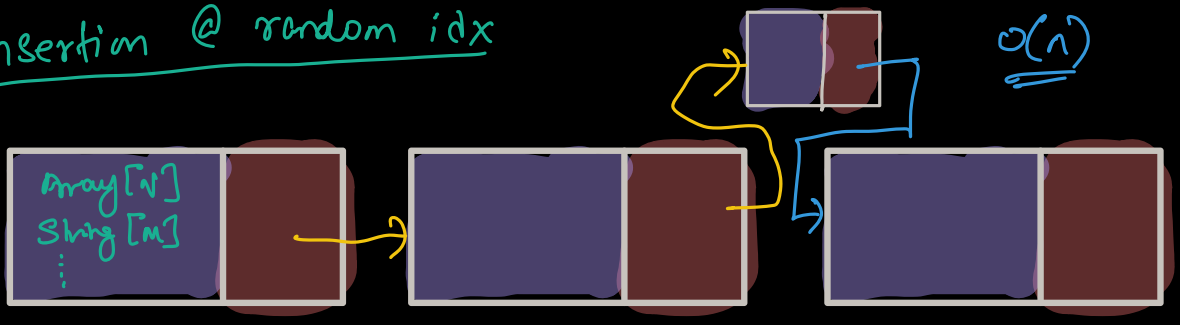




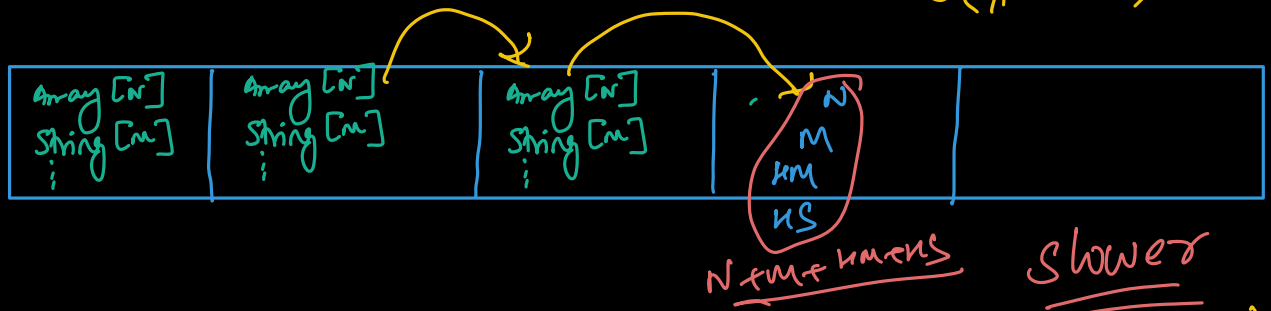
TC: O(N)

Delete: Todo

Insertion @ random idx



O(# items) is for 1 shift



N times

Slower

O (N x (# items))

→ Write code in paper

→ Dry run your code

→ Never do trial & error. Try to figure out the mistake.

→ Check your edge cases

* null
* size 1
* size 2 | 3
* Problem specific cases.