Todays Content -s Recursion

- -> How to write recursive code / trace
- -> TC/SC of recursive code in Next class.

Why Recursion?

- -) Merge Sort / Quick Sort
- -> Binary Tree / BST/ Bolanced BST / Segment Trees / Tries
- -> Dyramic Programming
- -> Back tracking
- -> Graphs

Recursion ? => Function calling itself

6 Solving a problem using smaller instance of same problem

$$Sum(N) = [+2+3+...+N-(+N)]$$

 $Sum(N) = Sum(N-1)+N$

$$Sum(u) = \frac{Sum(3) + 4}{6}$$

$$Sub - problem$$

On! How to write remosive code?

Assumption: Fix what your function should do

Main Logic! Solve the assumption using the problem

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Base Condition: Inputs for which you want to stop the recursion.

Problems

int Sum (N) Σ Assumption: Given N, calc & return the sum of first N natural numbers.

if (N = =1) Σ return 1 Σ

return SUM(N-1)+N

Main Logic:

$$Sum(N) = 1+2+3+.-+N-1+N$$

$$Sum(N) = Sum(N-1) + N$$

$$Sum(1) = Sum(0) + 1$$

$$= 1.$$

On: fact(N)? $(N \ge 1)$ fact(3) = 3x 2x1 = 6, fact(u) = 4x 3x 2x1 = 24.

fact(3) = 3x 2x1 = 6, fact(u) = 4x 3x 2x1 = 24.

int fact(N) & Assumption: Griven N, find & return N.

int fact(N) & Section 1 & Main Logic:

fact(N) = Nx N - 1x N - 1x

Function Call Tracing

int add (N, m)?

return N+M

mul (N, y)?

return x*y

sub (a, b)?

returns a-b

returns a-b

main () $\frac{2}{2}$ x = 10; y = 20 x = 10; y = 1

Dutput: 825

11 Data Structure

add(x,y): returns (30) [remove]

mul (add(x,y), 30): returns (900) [remove]

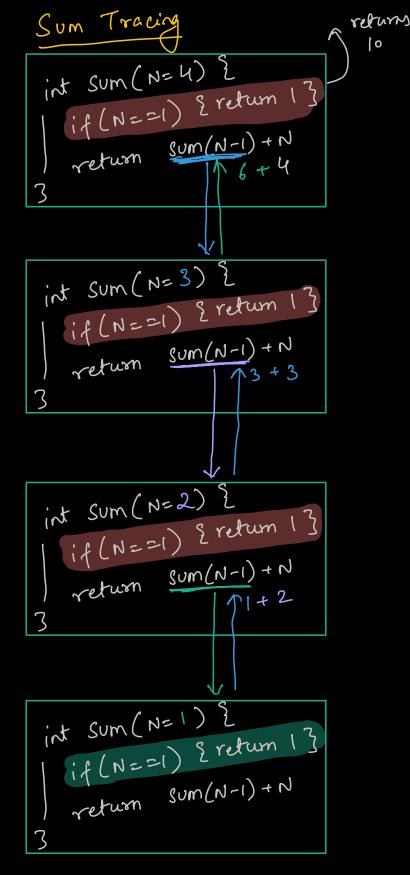
sub (mul (add(x,y), 30), 75): returns
825

LIFO Structure

Last-In

First Out

DS: Stack bandles the implementation



Stack

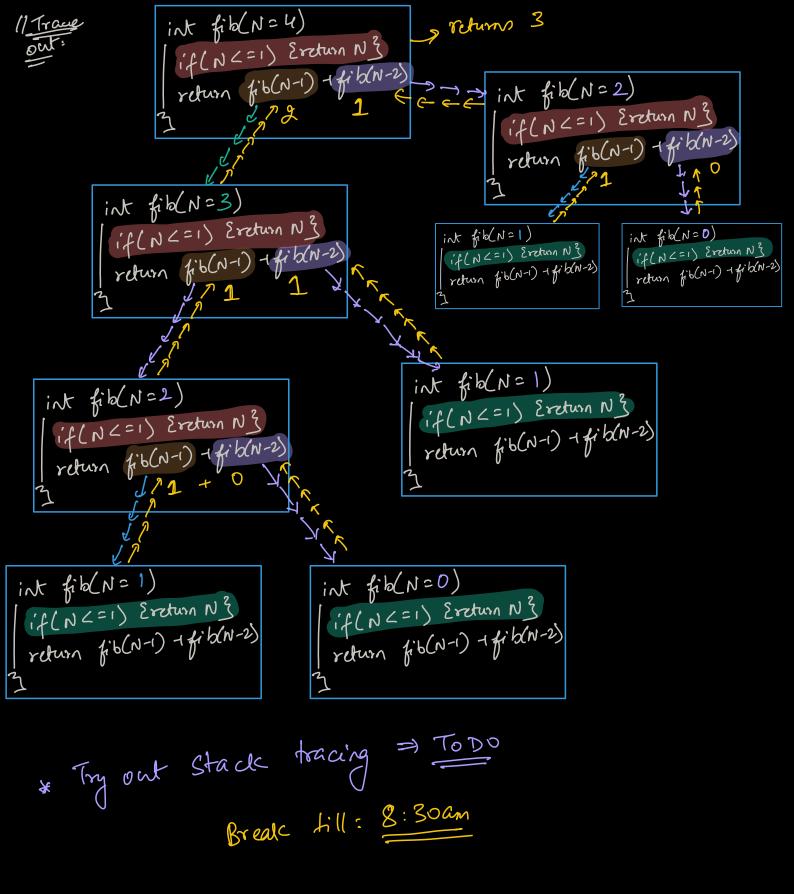
Sum (1): return 1 sum(1)+2 Sum(2): return (um (3): return sum(2)+3 Sum(4): returns sum(3)+4= 10

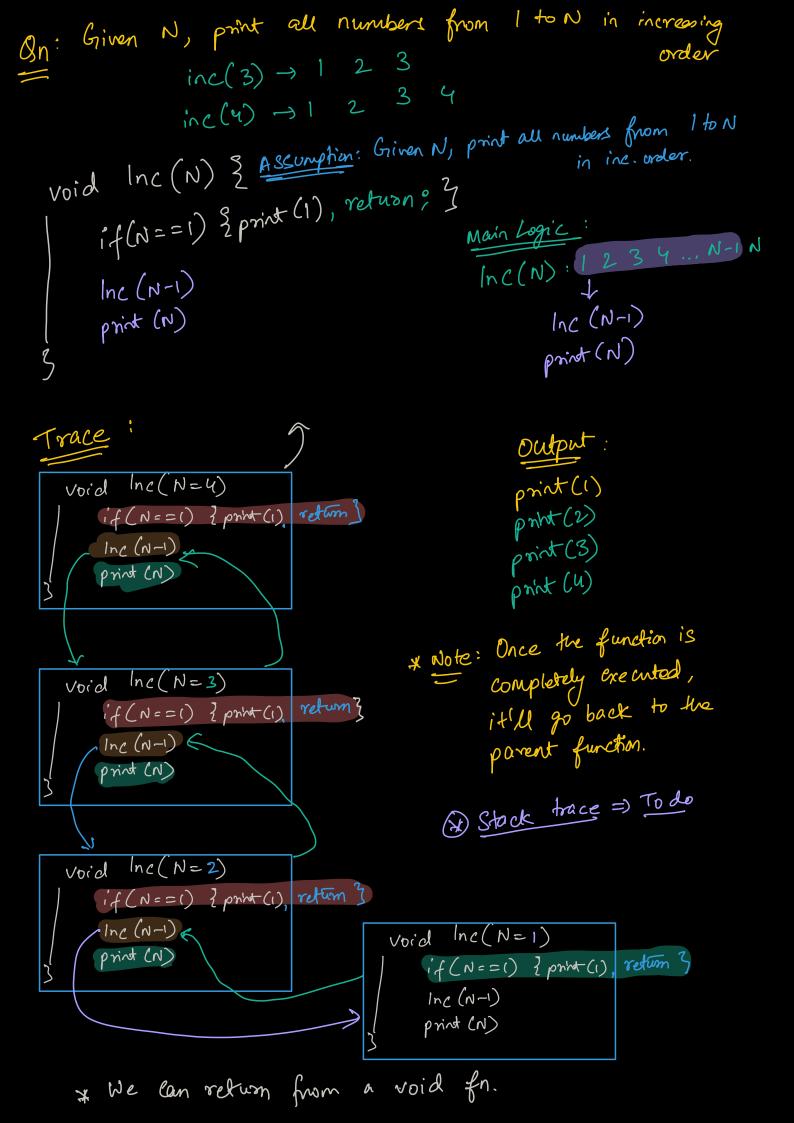
What happens if base cords not there? b) MLE V a) TLE

Note: In recursion, if you get MLE, vanify base condition.

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Bn: N >=0 (Sum of prev 2 numbers): Fisonacci # input(N) 0 1 2 3 4 5 6 7 8 9 10 Fib() 0 1 2 3 5 8 13 21 34 55 fib (9) = 34 int fib(N) { Assumption: Given N, Calc & return Nth fibonacci if(n<=1) ? return n? fib(N)= sum of prev 2 terms return fib (N-1) + fib(N-2) fib(N) = N-1 + N-2 fib(N) = fib(N-1) + fib(N-2) Note: Now to figure out base condition. G For which valid inputs, expression is invalid. fib(0) = fib(-1) + fib(-2) x fib(i) = fib(0) + fib(-1) x fib(2) = fib(1) + fib(0)





Qn-	Given N, print all numbers from N to 1 (dec. order) dec (N)
gn:	Given a substing, check if its a palindrome or not. S=4, e=6 => true S=e S=e S=e S=e S=e S=e S=e
60	good dad S=2, e=3 S=2, e=3 S=2, e=3 S=2, e=3 S=2, e=3 S=3 S=2, e=3 S=3, e=3 S
الم	Main Logic: Ch[]: Still e-1 Ch[]: O if (ch[s] == ch[e])
	(i) if (ch[s] == ch[e]) (i) Subshing [s+1, e-i] should be a palindrome is Pal (ch, s+1, e-i) mada m malyala m

```
m a 2 3 4 S=0, a=4 5 return
```

```
bool is Pal (ch[], S=0, e=4)

if (S>e) { return true }

return (ch[S] == ch[e] & 2 is Pal (ch, S+1, e-1))

}
```

```
bool is Pal (ch[], S=1, e=3)

if (S>e) { return true }

return (ch[S] == ch[e] &2 is Pal (ch, S+1, e-1))

}
```

```
bool is Pal (ch[], S=2, e=2)

if (S>e) { return true }

return (ch[S] == ch[e] &2 is Pal (ch, S+1, e-1))
```

```
bool is Pal (ch[], S=3, e=1)

if (S>e) { return true }

return (ch[S] == ch[e] && is Pal (ch, S+1, e-1))
}
```

Input: a n m e t n a S=0, e=6

) return false

```
bool is Pal (ch[], S=0, e=6)

if (S>e) { return true }

return (ch[S] == ch[e] &2 is Pal (ch, S+1, e-1))
```

```
bool is Pal (ch[], S=1, e=5)

if (S>e) { return true }

return (ch[S] == ch[e] & 2 is Pal (ch, S+1, e-1))
```

bool is Pal (ch[], S= 2, e= 4)

if (S>e) { return true }

return (ch[S] == ch[e] &2 is Pal (ch, S+1, e-1))

F