

Today's class

- Intro
- Flipping
- Sort ch[]
- Reverse string
- Longest Palindromic Substring.

String :
 → array of characters
 → bunch of characters
 → sequence of characters

$\{acb\}$
 $\{abc\}$] Order matters

Characters

How is a char stored? \Rightarrow ASCII value

ASCII val $\xrightarrow{+32}$ 'a' - 97
 'A' \rightarrow 65 $\xleftarrow{-32}$ 'b' - 98
 'B' \rightarrow 66
 'C' \rightarrow 67
 'c' \rightarrow 99
 'Z' \rightarrow 90
 'z' \rightarrow 122

'0' - 48
 '1' - 49
 '2' - 50
 ...
 '9' - 57
 '10' -

\hookrightarrow Not a single char

Size of char = 1 byte \Rightarrow 8 bits

7 6 5 4 3 2 1 0
 0 0 1 1 1 0 0 1

\Rightarrow 57

ch = '9' (57)
 ch = ch + 8 (57+8)
 print(ch) (65)
 \hookrightarrow (A)

Strings: array of characters

```
String s = "abcd"
print(s[0])
    ↓
    'a'
```

```
char ch[] = "abcd"
print(ch[0])
    ↳ 'a'
```

Q1: Given a char[], Toggle everything -

↳ Small \Leftrightarrow Capital

Note: char[] contains only
lower case
or
uppercase alphabets.

Eg: ch[] = AnaConDa

Toggle: aNAcONdA

```
Toggle(char[] s) {
```

```
    int n = s.length
```

```
    for (i=0; i < n; i++) {
```

```
        if (s[i] >= 65 & s[i] <= 90) {
```

```
            s[i] = s[i] + 32
```

```
        } else {
```

```
            s[i] = s[i] - 32
```

```
        }
```

```
    }
    return s
```

Tc: $O(n)$
Sc: $O(1)$

Solve w/o

if-else

$s[i] = s[i] \wedge 32$

or
 $s[i] = s[i] \wedge (1 \ll 5)$

}

④ HINT: 32 is a special number

$2^7 \ 2^6 \ 2^5 \ 2^4 \ 2^3 \ 2^2 \ 2^1 \ 2^0$

'A' (65): $\underline{0} \ \underline{1} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{1}$

'B' (66): $\underline{0} \ \underline{1} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{1} \ \underline{0}$

'C' (67): $\underline{0} \ \underline{1} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{1} \ \underline{1}$

⋮
'Z' (90): $\underline{0} \ \underline{1} \ \underline{0} \ \underline{1} \ \underline{1} \ \underline{0} \ \underline{1} \ \underline{0}$

'a': $\overset{2^7}{\underline{0}} \ \overset{2^6}{\underline{1}} \ \overset{2^5}{\underline{1}} \ \overset{2^4}{\underline{0}} \ \overset{2^3}{\underline{0}} \ \overset{2^2}{\underline{0}} \ \overset{2^1}{\underline{0}} \ \overset{2^0}{\underline{1}}$

'b': $\underline{0} \ \underline{1} \ \underline{1} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{1} \ \underline{0}$

'c': $\underline{0} \ \underline{1} \ \underline{1} \ \underline{0} \ \underline{0} \ \underline{0} \ \underline{1} \ \underline{1}$

⋮
'z': $\underline{0} \ \underline{1} \ \underline{1} \ \underline{1} \ \underline{1} \ \underline{0} \ \underline{1} \ \underline{0}$

└────────────────────────────────┘

Convert

Toggle the 5th bit

Upper case
5th bit is
unset

Small case
5th bit
is set

Q2: Given a `char[]`, which contains only lower case alphabets.
Sort the given `char[]` in alphabetical order

Eg: `s = d a b a c d b`
 \downarrow Sort
`a a b b c d d`

Constraint:

$$1 \leq N \leq 10^5$$

$$'a' \leq \text{ch}[i] \leq 'z'$$

ideas: ① Sort `ch[]` using bubble sort

TC: $O(N^2)$

$\rightarrow 10^{10} \rightarrow$ TLE

② Use in-built library + custom comparator (if required):

TC: $O(N \log N)$ ✓

#idea:

`s = d a b a c d b`

`'a' - 2`
`'b' - 2`
`'c' - 1`
`'d' - 2`

\Rightarrow ans \Rightarrow `a a b b c d d`

\downarrow
`s =` ~~`a`~~ ~~`a`~~ ~~`b`~~ ~~`b`~~ ~~`c`~~ ~~`c`~~ ~~`c`~~ ~~`e`~~ ~~`e`~~
 \uparrow

`= a a b b b c c c e e`

`'a' - 2`
`'b' - 3`
`'c' - 3`
`'d' - 0`
`'e' - 2`

ans = `a a b b b c c c e e`

$\begin{matrix} & 0 & 1 & 2 & 3 & 4 & 5 & & 25 \end{matrix}$
`cnt[26] =`

2	3	3	0	2	-	-	-	1	0
---	---	---	---	---	---	---	---	---	---

$$\text{idx} = s[i] - 97$$

$$\text{cnt}[\text{idx}] = \text{cnt}[\text{idx}] + 1$$

How many distinct chars: 26

int cnt[26] = {0}

cnt[0] → freq of 'a' (97)

cnt[1] → freq of 'b' (98)

cnt[2] → freq of 'c' (99)

⋮
cnt[25] → freq of 'z' (122)

c - 97

or

c - 'a'

or

c = c % 97

Sort String(char s[]) {

n = s.length

int cnt[26] = {0}

for(i=0; i<n; i++) { // iterate string

idx = s[i] - 97

cnt[idx]++

}

// modify original string.

k=0

for(i=0; i<26; i++) { // iterate cnt array.

// cnt[i] = freq of ('a'+i)

char ch = 'a'+i

for(j=1; j<=cnt[i]; j++) {

s[k] = ch

k++

}

}
return s

TC: O(N)

SC: O(26)
= O(1)

N x 26

TC: O(N)

SC: O(1)

TC: O(N)

SC: O(1)

Table:

i	j	#iter
0	[1 cnt[0]]	cnt[0]
1	[1 cnt[1]]	cnt[1]
2	[1 cnt[2]]	cnt[2]
⋮		⋮
25	[1 cnt[25]]	cnt[25]

→ this sum = No of iter.

freq of 'a' + freq of 'b' + freq of 'c'
+ + freq of 'z'

Length of string = N

Break :

8:30 am

Substring : concept is same as subarray

- ① Continuous part of string
- ② Full string is a substring
- ③ A single char is also a substring.

Qn: Check if a given substring is palindrome or not.

Eg: man madam
 dad radar
 civic tenet

malayalam

ch[11]: 0 1 2 3 4 5 6 7 8 9 10
 a n a m a d a m s p e

[3 7] : madam ✓

[8 10] : spe ✗

[0 2] : ana ✓

bool isPalindrome(char c[], int s, int e) {
 // Check if substring [s e] is a palindrome.

 while(s < e) {
 if(c[s] != c[e]) {
 return false
 }

 s++, e--

 }
 return true

Tc: O(n)
Sc: O(1)

Qn: Given a string, calc length of longest palindromic substring.

Eg1: a ba cab
len = 5

Eg: a b c d e
len = 1

Idea: For every substring, check if its a palindrome or not.
Get max length.

of substrings \times is Palindrome check
 $\frac{N(N+1)}{2} \times N = \underline{O(N^3)}$

$1 \leq N \leq 3 \times 10^2$

27×10^9 X

```
int long Palindrome(char s[]) {  
    int n = s.length  
    int ans = 0
```

```
    for(i = 0; i < n; i++) { // start
```

```
        for(j = i; j < n; j++) { // end
```

```
            // substring [i, j]
```

```
            if(isPalindrome(s, i, j)) { // len = j - i + 1
```

```
                ans = max(ans, j - i + 1)
```

```
            }
```

```
        }
```

```
    }
```

```
    return ans
```

TC: $O(N^3)$
SC: $O(1)$

→ won't work
→ TLE

Eg: x b d y z z y d b d y z y d x
 ↑ ↑ ↑ ↓ ↑ ↓ ↓
 p₁ c₁ c₂ p₁ p₂ c₁ p₂

length: $p_2 - p_1 - 1$

even length:

* Take every adjacent char as centre & expand.

Adj Centre: $n-1 \times \text{Expansion}$
TC: $O(n) \times O(n)$
 $= O(n^2)$

odd length

length excluding p₁p₂

$[p_1, p_2] : p_2 - p_1 + 1$

$(p_1, p_2) = p_2 - p_1 - 1$

* Take every char as centre & expand left & right around it, get max palindromic substring
 $p_2 - p_1 - 1$

TC: #No of centre * expansion
 $n \times n$
 $= O(n^2)$

```

int expand(char s[], int p1, int p2) {
    while(p1 >= 0 && p2 < N && s[p1] == s[p2]) {
        p1--;
        p2++;
    }
    return p2 - p1 - 1;
}

```

$p1 \geq 0$
 $p2 < N$
 $s[-1]$
 $TC: O(N)$

```

int long Palindrome(char s[]) {
    int n = s.length;
    int ans = 0;
    for(i = 0; i < n; i++) { // odd length
        // Centre: s[i]
        p1, p2 = i
        ans = max(ans, expand(s, p1, p2))
    }
    for(i = 0; i < n-1; i++) { // even length
        // Centre: s[i], s[i+1]
        p1 = i, p2 = i+1
        ans = max(ans, expand(s, p1, p2))
    }
    return ans;
}

```

n^2
 n^2

$i < n$
 $i < n-1$
 both works

$TC: O(N^2)$
 $SC: O(1)$

Longest Palindromic Substring

- ✓ ① BF: $O(N^3)$
- ✓ ② Expand around centres $O(N^2)$
- ✓ ③ $O(N^2)$: DP
↳ adv module

- ✗ ④ $O(N \log N)$: Binary Search + Rabin Karp
- ✗ ⑤ $O(N)$: Manacher algo
→ optimal class for adv

Qn: Given a string, reverse the words.

Eg:- "i hate love dsa"

Ans:- dsa love hate i

SC: O(1)

① Reverse each word

i etah evol asd

② Reverse whole string

dsa love hate i

Todo

reverse array

```
for (i=0; i<n/2; i++)  
    swap(a[i], a[n-i-1])
```

or

```
s=0, e=n-1  
while (s<e) {  
    swap(a[s], a[e])  
    s++, e--  
}
```