

### 3.a)

If  $f(n)$  is  $O(g(n))$  then  $g(n)$  is  $O(f(n))$

Counterexample to prove this is false:

$$\begin{aligned}\text{Let } f(n) &= 1 \\ g(n) &= n\end{aligned}$$

$$f(n) \text{ is } O(g(n))$$

$$\text{implies } 0 \leq 1 \leq cn$$

for  $c=1, n_0=1$  the above condition holds good

$$g(n) \text{ is } O(f(n)) \text{ implies } 0 \leq n \leq c$$

As we can't find a constant  $c$  in this case which holds for all  $n \geq n_0$   
 $g(n)$  is  $O(f(n))$  does not hold good.

### 3.b)

If  $f(n)$  is  $\Theta(g(n))$  then  $\lg f(n)$  is  $\Theta(\lg g(n))$

Counterexample to prove this is false:

$$\begin{aligned}\text{Let } f(n) &= 2 \\ g(n) &= 1\end{aligned}$$

$$f(n) \text{ is } \Theta(g(n))$$

$$\text{implies } 0 \leq c_1 \leq 2 \leq c_2$$

The above case can be satisfied for some value of  $c_1 = 1$  and  $c_2 = 3$

$$\lg f(n) \text{ is } \Theta(\lg g(n))$$

$$\text{implies } 0 \leq c_1 \lg g(n) \leq \lg f(n) \leq c_2 \lg g(n)$$

$$\Rightarrow 0 \leq 0 \leq 1 \leq 0$$

which cannot be true.

### 3.c)

**If  $f(n)$  is  $\Theta(g(n))$  then  $2^{f(n)}$  is  $\Theta(2^{g(n)})$**

Counterexample to prove this is false:

Let  $f(n)=2n$ ,  $g(n)=n$

$f(n)$  is  $\Theta(g(n))$

implies  $0 \leq c_1 g(n) \leq f(n) \leq c_2 g(n)$

$\Rightarrow 0 \leq c_1 n \leq 2n \leq c_2 n$

The above case holds for  $c_1=1$ ,  $c_2=3$ ,  $n_0=1$

$2^{f(n)}$  is  $\Theta(2^{g(n)})$

implies  $0 \leq c_1 2^{g(n)} \leq 2^{f(n)} \leq c_2 2^{g(n)}$

$\Rightarrow 0 \leq c_1 2^n \leq 2^{2n} \leq c_2 2^n$

$\Rightarrow c_2 \geq 2^n$  for all  $n \geq n_0$  which is not possible.

**3.d)**

**If  $f(n)$  is  $O(g(n))$  then  $g(n)$  is  $\Omega(f(n))$**

This is true.

$f(n)$  is  $O(g(n))$

implies  $0 \leq f(n) \leq cg(n)$

Let  $c'=1/c$

$\Rightarrow 0 \leq c'f(n) \leq g(n)$

$\Rightarrow g(n)$  is  $\Omega(f(n))$

**3.e)**

**$f(n)$  is  $\Theta(f(n/2))$**

Counterexample to prove this is false:

Let  $f(n) = 2^n$

$f(n)$  is  $\Theta(f(n/2))$

implies  $0 < c_1 (f(n/2)) \leq f(n) \leq c_2 (f(n/2))$

$\Rightarrow 2^n \leq c_2 2^{n/2}$

$\Rightarrow c_2 \geq 2^{n/2}$  for all  $n \geq n_0$  which is not possible.

**3.f)**

**If  $g(n)$  is  $o(f(n))$  then  $f(n) + g(n)$  is  $\Theta(f(n))$**

This is true.

$g(n)$  is  $o(f(n))$

$\Rightarrow g(n) < c f(n)$

$f(n) \leq f(n) + g(n) \leq f(n) + c f(n)$

$\Rightarrow f(n) \leq f(n) + g(n) \leq c' f(n)$  ( $c' = c+1$ )

$\Rightarrow f(n) + g(n)$  is  $\Theta(f(n))$

**3.g)**

**$f(n) + g(n)$  is  $\Theta(\max(f(n), g(n)))$**

This is true.

$f(n) + g(n)$  is  $\Theta(\max(f(n), g(n)))$

$\Rightarrow c_1 \max(f(n), g(n)) \leq f(n) + g(n) \leq c_2 \max(f(n), g(n))$

$c_1 \max(f(n), g(n)) \leq f(n) + g(n)$  holds good for  $c_1 = 1$

as the summation of 2 numbers is always greater than or equal to the maximum of the two values.

$f(n)+g(n) \leq c_2 \max(f(n),g(n))$  holds good for  $c_2=2$

$\Rightarrow (f(n)+g(n))/2 \leq \max(f(n),g(n))$

as the average of 2 numbers is always less than or equal to the maximum of the two values.

Hence  $f(n) + g(n)$  is  $\Theta(\max(f(n),g(n)))$  is true.