



Worksheet 2.1

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Subject Name: Advance Programming Lab **Subject Code**: 20CSP-334

1. Aim/Overview of the practical:

a) From a given vertex in a weighted connected graph, find shortest paths to other vertices using Dijkstra's algorithm.

b) Compute the transitive closure of a given directed graph using Warshall's algorithm.

2. Task to be done:

- a) From a given vertex in a weighted connected graph, find shortest paths to other vertices using Dijkstra's algorithm.
- b) Compute the transitive closure of a given directed graph using Warshall's algorithm.

3. Algorithm/Flowchart (For programming-based labs):

a) Dijkstra's algorithm

- Create a set **sptSet** (shortest path tree set) that keeps track of vertices included in the shortest-path tree, i.e., whose minimum distance from the source is calculated and finalized. Initially, this set is empty.
- Assign a distance value to all vertices in the input graph. Initialize all distance values as **INFINITE**. Assign the distance value as 0 for the source vertex so that it is picked first.
- While **sptSet** doesn't include all vertices
 - Pick a vertex **u** which is not there in **sptSet** and has a minimum distance value.
 - Include u to **sptSet**.
 - Then update distance value of all adjacent vertices of u.
 - To update the distance values, iterate through all adjacent vertices.
 - For every adjacent vertex v, if the sum of the distance value of u (from source) and weight of edge u-v, is less than the distance value of v, then update the distance value of v.

b) Floyd Warshall:

- For the first step, the solution matrix is initialized with the input adjacent matrix of the graph. Let's name it as reach.
- Next we need to iterate over the number of nodes from {0,1,....n} one by one by considering them strating vertex. Similarly, another iteration is performed over the nodes {1,2,....,n} by considering ending vertex one by one.







- For the shortest path, we need to form another iteration which ranges from {1,2,...,k-1}, where vertex k has been picked up as an intermediate vertex.
- For every pair (i, j) of the starting and ending vertices respectively, there are two possible cases.
- if k is an intermediate vertex in the shortest path from i to j, then we check the condition reach[i][j] > reach[i][k] + reach[k][j] and update reach[i][j] accordingly.
- Otherwise, if k is not an intermediate vertex, we don't update anything and continue the loop.

Transitive Closure condition:

Only one difference of the condition to be checked when there is an intermediate vertex k exits between the starting vertex and the ending vertex. We need to check two conditions and check if any of them is true,

- Is there a direct edge between the starting vertex and the ending vertex? If yes, then update the transitive closure matrix value as 1.
- For k, any intermediate vertex, is there any edge between the (starting vertex & k) and (k & ending vertex)? If yes, then update the transitive closure matrix value as 1.

4. CODE:

```
a) Dijkstra
// Dijkstra's single source shortest path using adjacency matrix representation of the graph
#include <bits/stdc++.h>
using namespace std;
// Number of vertices in the graph
#define V 9
// function to find the vertex with minimum
// distance value, from the set of vertices not yet included
// in shortest path tree
int minDistance(int dist[], bool sptSet[])
        // Initialize min value
        int min = INT MAX, min index;
        for (int v = 0; v < V; v++)
                if (\operatorname{sptSet}[v] == \operatorname{false \&\& dist}[v] <= \min)
                        min = dist[v], min\_index = v;
        return min_index;
}
void printSolution(int dist[])
        cout << "Vertex \t\t Distance from Source" << endl;</pre>
```







```
for (int i = 0; i < V; i++)
               cout \ll i \ll " \t t \ll " \dist[i] \ll endl;
}
void dijkstra(int graph[V][V], int src)
       int dist[V]; // The output array. dist[i] will hold the shortest
       // distance from src to i
       bool sptSet[V]; // sptSet[i] will be true if vertex i is included in shortest path tree
                        // or shortest distance from src to i is finalized
       // Initialize all distances as INFINITE and stpSet[] as false
       for (int i = 0; i < V; i++)
               dist[i] = INT_MAX, sptSet[i] = false;
       // Distance of source vertex from itself is always 0
       dist[src] = 0;
       // Find shortest path for all vertices
       for (int count = 0; count < V - 1; count++) {
               // Pick the minimum distance vertex from the set of vertices not yet processed.
               //u is always equal to src in the first iteration.
               int u = minDistance(dist, sptSet);
               // Mark the picked vertex as processed
               sptSet[u] = true;
               // Update dist value of the adjacent vertices of the picked vertex.
               for (int v = 0; v < V; v++)
                       // Update dist[v] only if is not in sptSet,
                       // there is an edge from u to v, and total
                       // weight of path from src to v through u is
                       // smaller than current value of dist[v]
                       if (!sptSet[v] && graph[u][v]
                               && dist[u]!= INT MAX
                               && dist[u] + graph[u][v] < dist[v]
                               dist[v] = dist[u] + graph[u][v];
        }
       // print the constructed distance array
       printSolution(dist);
}
```







```
int main()
        int graph[V][V] = \{ \{ 0, 4, 0, 0, 0, 0, 0, 8, 0 \}, \}
                \{4, 0, 8, 0, 0, 0, 0, 11, 0\},\
                \{0, 8, 0, 7, 0, 4, 0, 0, 2\},\
                \{0, 0, 7, 0, 9, 14, 0, 0, 0\},\
                \{0, 0, 0, 9, 0, 10, 0, 0, 0\},\
                \{0, 0, 4, 14, 10, 0, 2, 0, 0\},\
                \{0, 0, 0, 0, 0, 0, 2, 0, 1, 6\},\
                \{8, 11, 0, 0, 0, 0, 1, 0, 7\},\
                \{0, 0, 2, 0, 0, 0, 6, 7, 0\}\};
        dijkstra(graph, 0);
        return 0;
}
b) Transitive Closure using Warshall
#include<bits/stdc++.h>
using namespace std;
#define V 4
void printSolution(int reach[][V])
{
  for (int i = 0; i < V; i++)
     for (int j = 0; j < V; j++)
         if(i == j)
           printf("1");
         else
           printf("%d", reach[i][j]);
     cout<<"\n";
   }
void transitiveClosure(int graph[][V])
  int reach[V][V], i, j, k;
  for (i = 0; i < V; i++)
     for (j = 0; j < V; j++)
        reach[i][j] = graph[i][j];
  for (k = 0; k < V; k++)
```







```
for (i = 0; i < V; i++)
        for (j = 0; j < V; j++)
          reach[i][j] = reach[i][j] ||
           (reach[i][k] && reach[k][j]);
        }
     }
  cout<<"Following matrix is transitive closure of the given graph\n";
  printSolution(reach);
}
int main()
  int graph[V][V] = \{ \{1, 1, 0, 1\}, \}
                \{0, 1, 1, 0\},\
                \{0, 0, 1, 1\},\
                \{0, 0, 0, 1\}
               };
  cout<<"Given graph: \n";</pre>
  printSolution(graph);
  transitiveClosure(graph);
  return 0;
  /*
         10
    (0)---->(3)
    5
                 | 1
    \|/
    (1)---->(2)
}
```

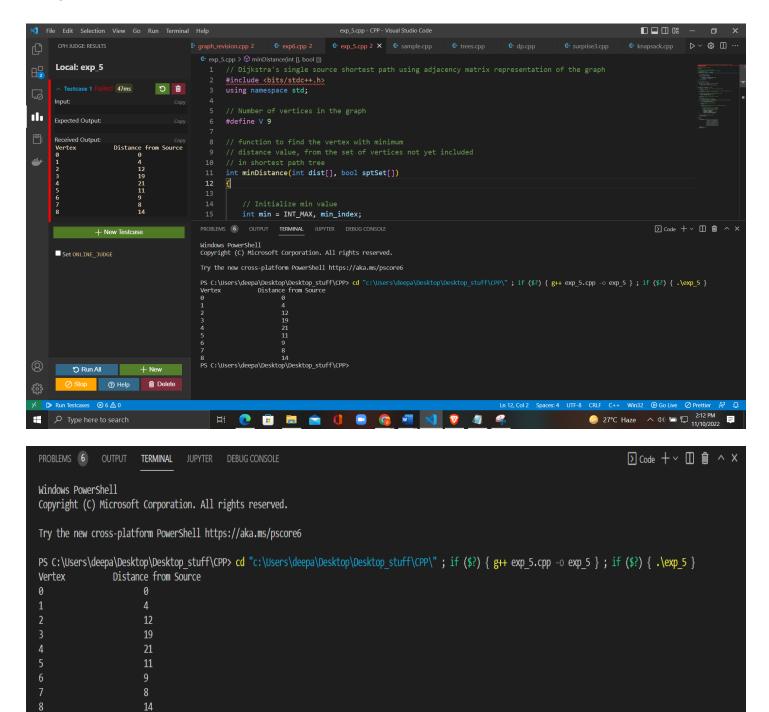






5. Result/Output:

a) Dijkstra



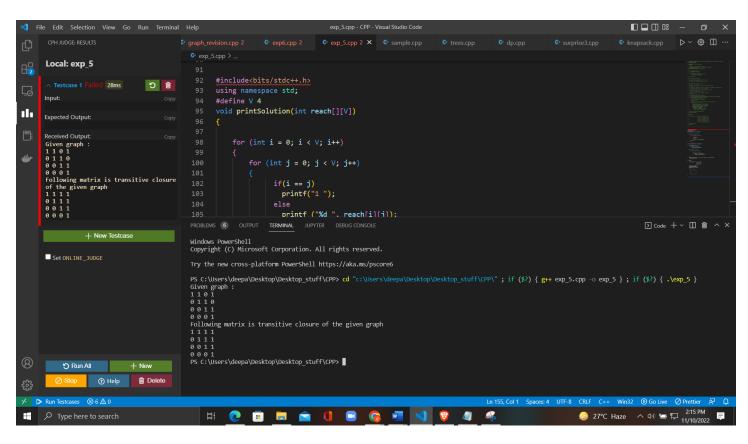


PS C:\Users\deepa\Desktop\Desktop stuff\CPP>





b) Transitive Closure using Warshall:



```
PROBLEMS 6
                                                                   OUTPUT
                                                                                                              TERMINAL
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 Try the new cross-platform PowerShell https://aka.ms/pscore6
 PS C:\Users\deepa\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desktop\Desk
 Given graph:
 1101
 0 1 1 0
 0011
 0001
 Following matrix is transitive closure of the given graph
 1111
 0111
0 0 1 1
 0001
 PS C:\Users\deepa\Desktop\Desktop stuff\CPP>
```







Learning Outcomes:

- Learn shortest paths to other vertices using Dijkstra's algorithm.
- Learn transitive closure of a given directed graph using Warshall's algorithm.

Evaluation Grid (To be created as per the SOP and Assessment guidelines by the faculty):

Maximum Marks

