

# NoSQL

## HISTORY



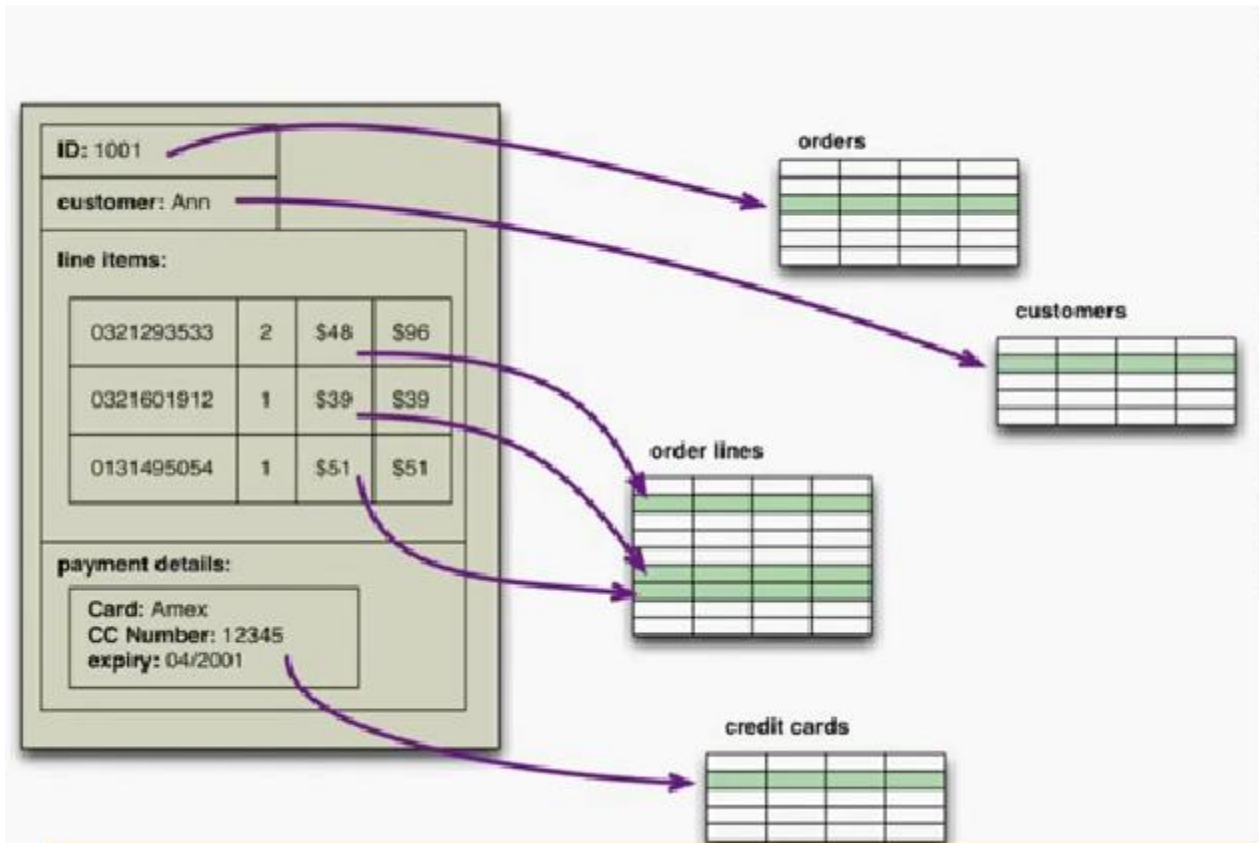
## Benefits of RDBMS

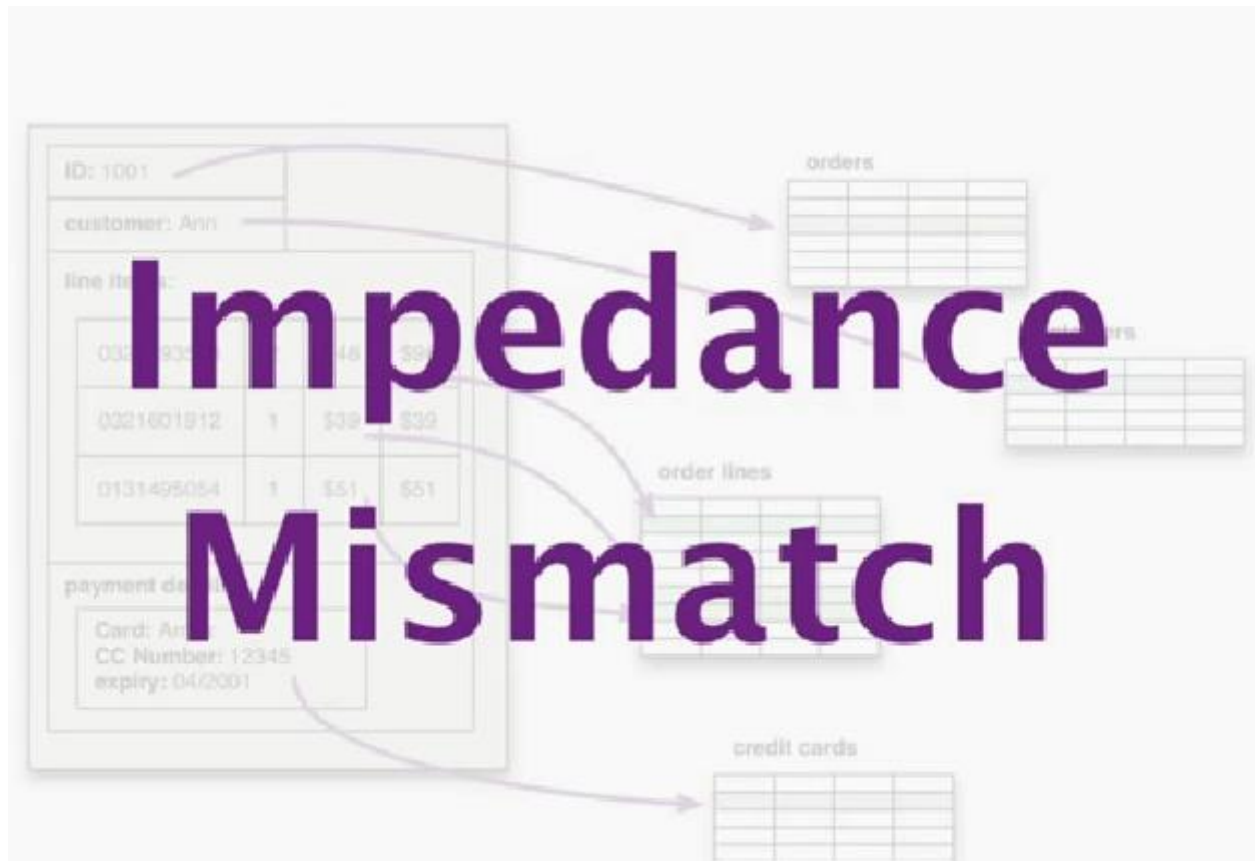
**Persistence**  
**Integration**

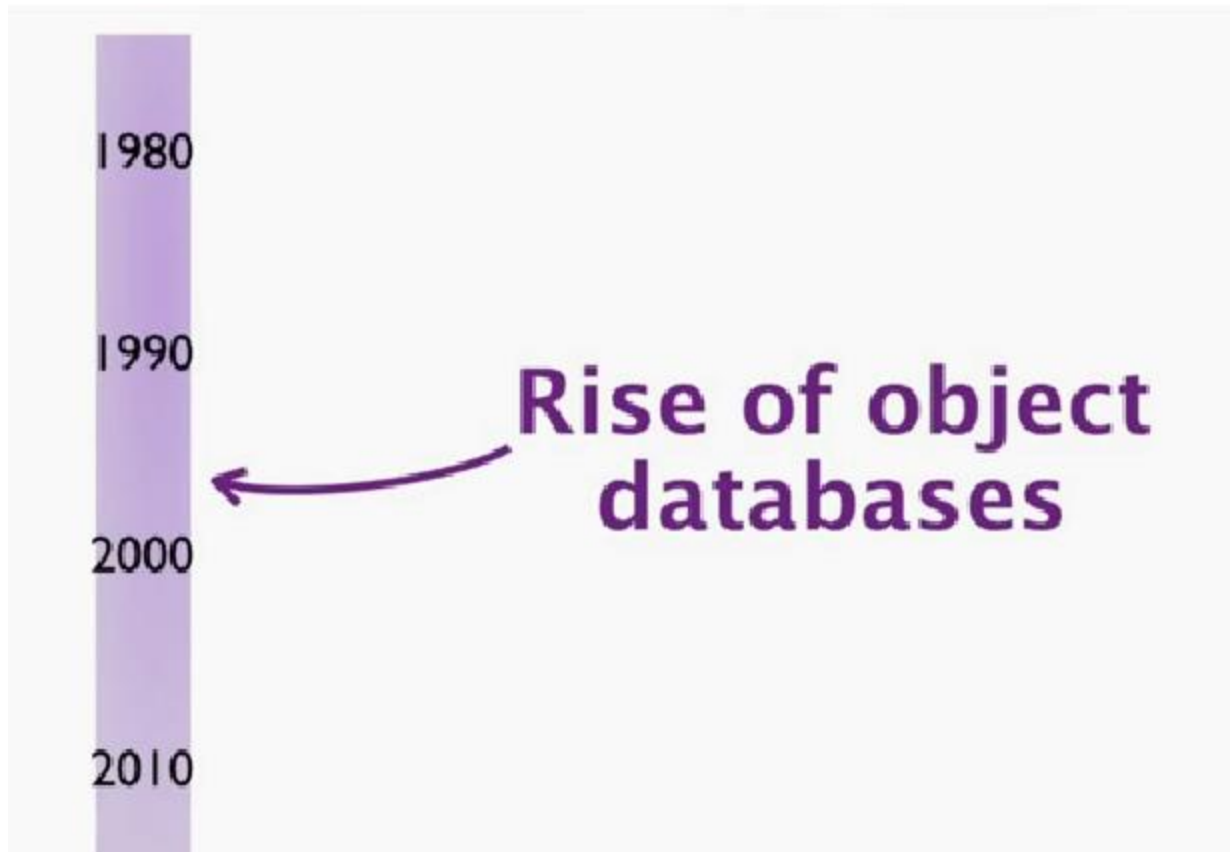
**SQL Transactions**

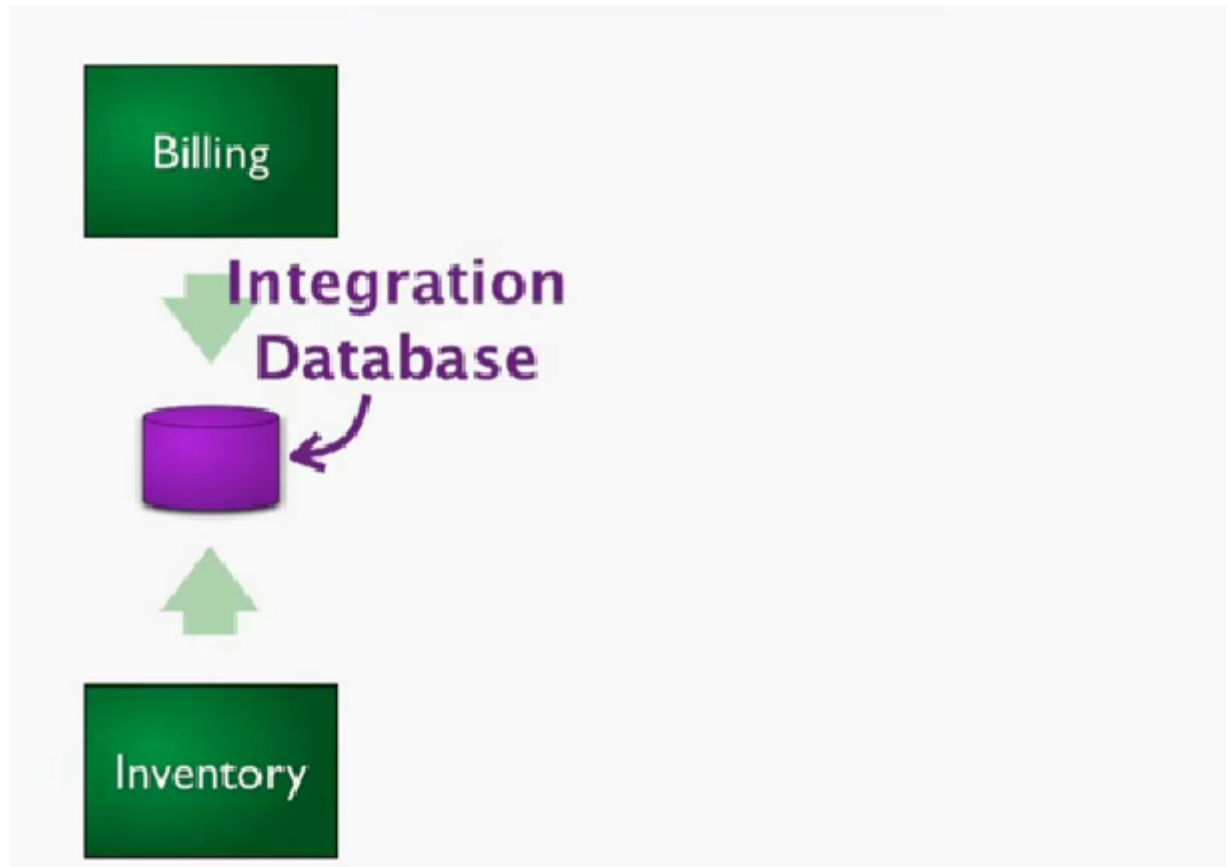
**Reporting**

## Problems with RDBMS model



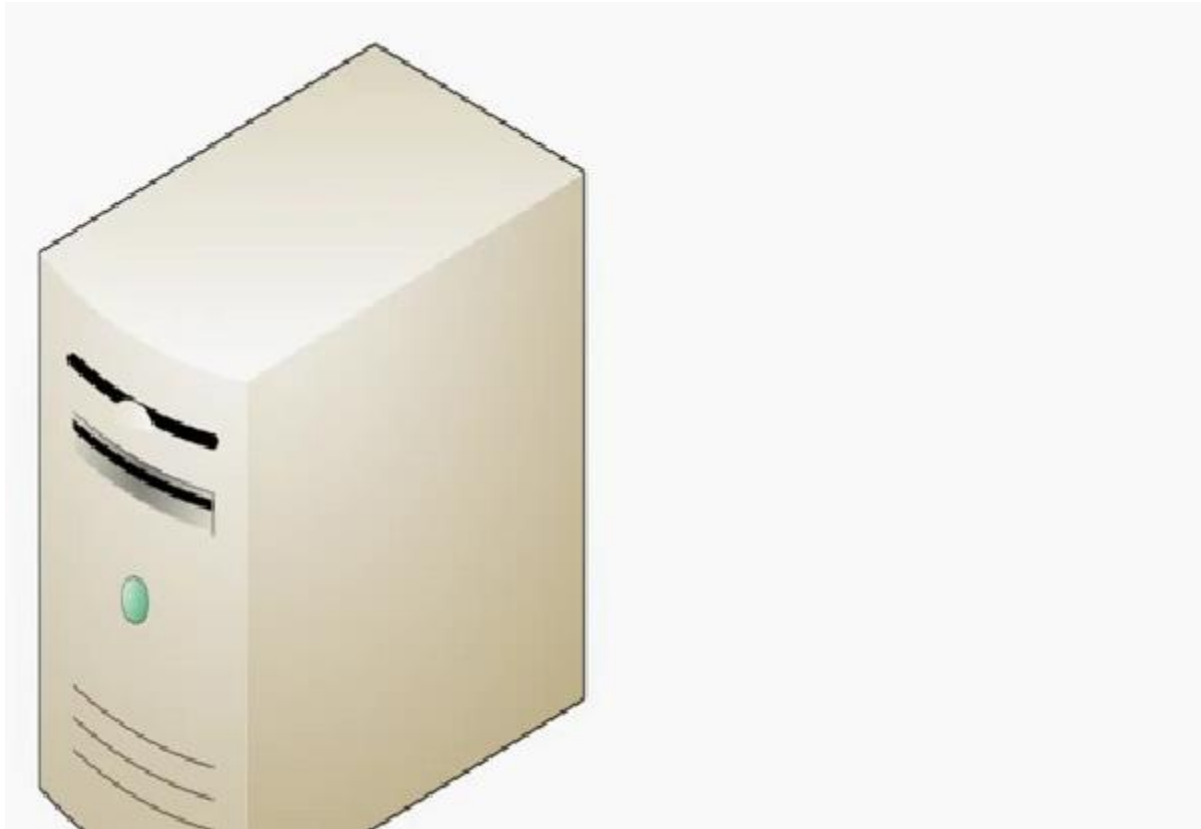
















Google™

amazon.com™

Google



**Bigtable**

amazon.com



**Dynamo**

# "NoSQL"



**Johan Oskarsson**

(from Hadoop)

Twitter hashtag

#nosql



← This event has ended

[View current events hosted by Last.fm](#)

## NOSQL meetup

Thursday, June 11, 2009 from 10:00 AM to 5:00 PM (PT)  
San Francisco, CA

### Ticket Information

TYPE	REMAINING	END	PRICE	QUANTITY
Free ticket	Sold Out	Ended	Free	Sold Out

Share this! [Share](#) [Tweet](#) [Like](#) [Be the first of your friends to like this.](#)

### Event Details

#### Introduction

This meetup is about "open source, distributed, non relational databases".

Have you run into limitations with traditional relational databases? Don't mind trading a query language for scalability? Or perhaps you just like shiny new things to try out? Either way this meetup is for you.

Join us in figuring out why these newfangled Dynamo clones and BigTables have become so popular lately. We have gathered presenters from the most interesting projects around to give us all an introduction to the field.

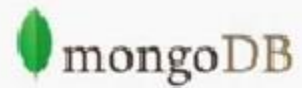
#### Preliminary schedule

### When & Where



**CBS Interactive, Magma room**  
235 Second Street  
San Francisco, CA 94105

Thursday, June 11, 2009 from 10:00 AM to 5:00 PM (PT)



Dynomite

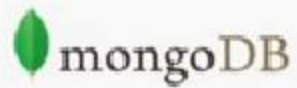
APACHE  
HBASE

# Definition of NoSQL

# Characteristics of NoSQL

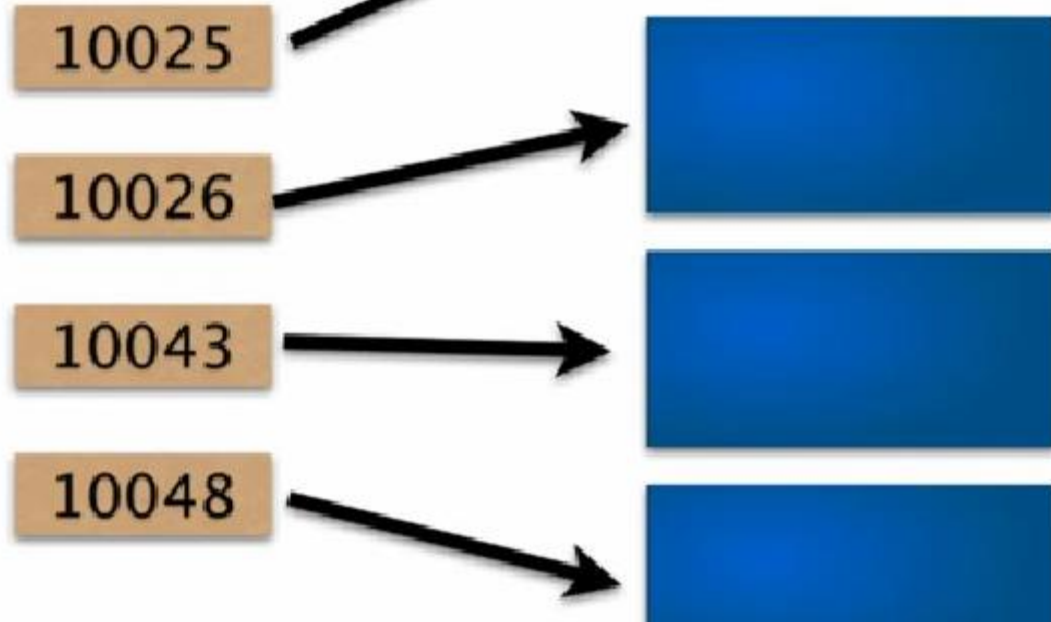
**non-  
relational open-source  
cluster-friendly  
21st Century  
Web  
schema-less**

# Data Model





# Key-value





# Document

```
{ "id": 1001,  
  "customer_id": 7231,  
  "line-items": [  
    { "product_id": 4555, "quantity": 8 },  
    { "product_id": 7655, "quantity": 4 }, { "product_id": 8755,
```

```
    { "product_id": 2155, "quantity": 4 } ],  
  "discount-code": "Y" }
```

no

```
anOrder["price"] * anOrder["quantity"]
```

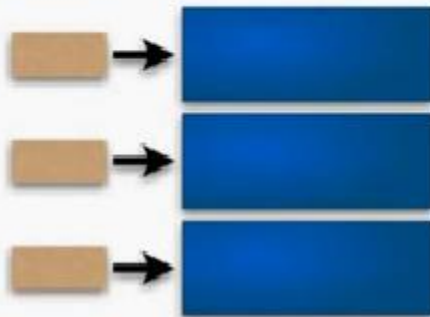


`anOrder["price"] * anOrder["quantity"]`

implicit  
schema



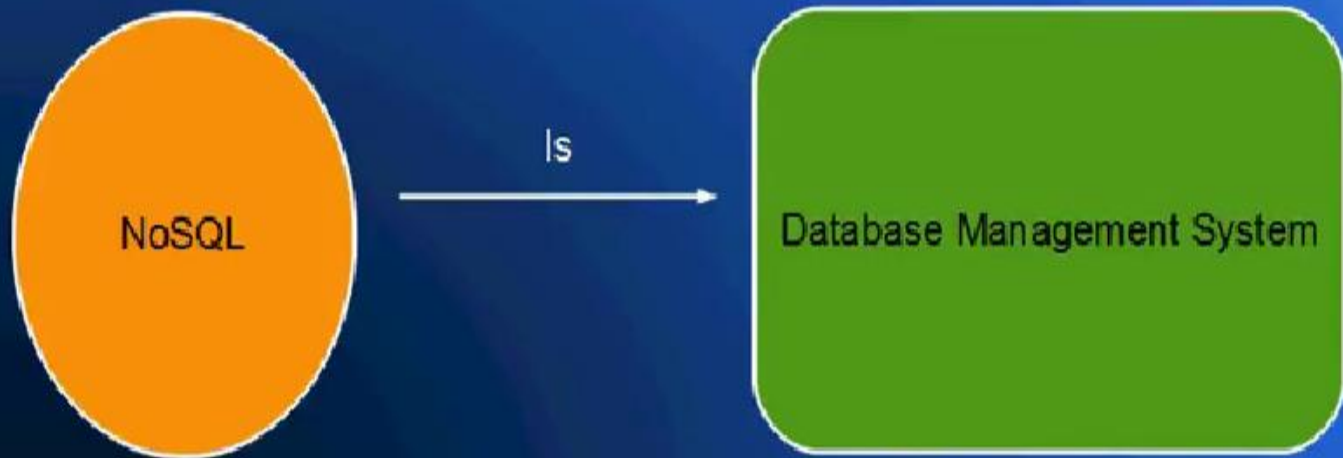
## Key-Value



## Document

```
{ "id": 1001,  
  { "id": 1002,  
    "customer_id": 7231,  
    "line-items": [  
      { "product_id": 4555, "quantity": 8},  
      { "product_id": 7655, "quantity": 4},  
      { "product_id": 8755, "quantity": 3}]  
    }  
  }
```

# What is ?



# Database Management System

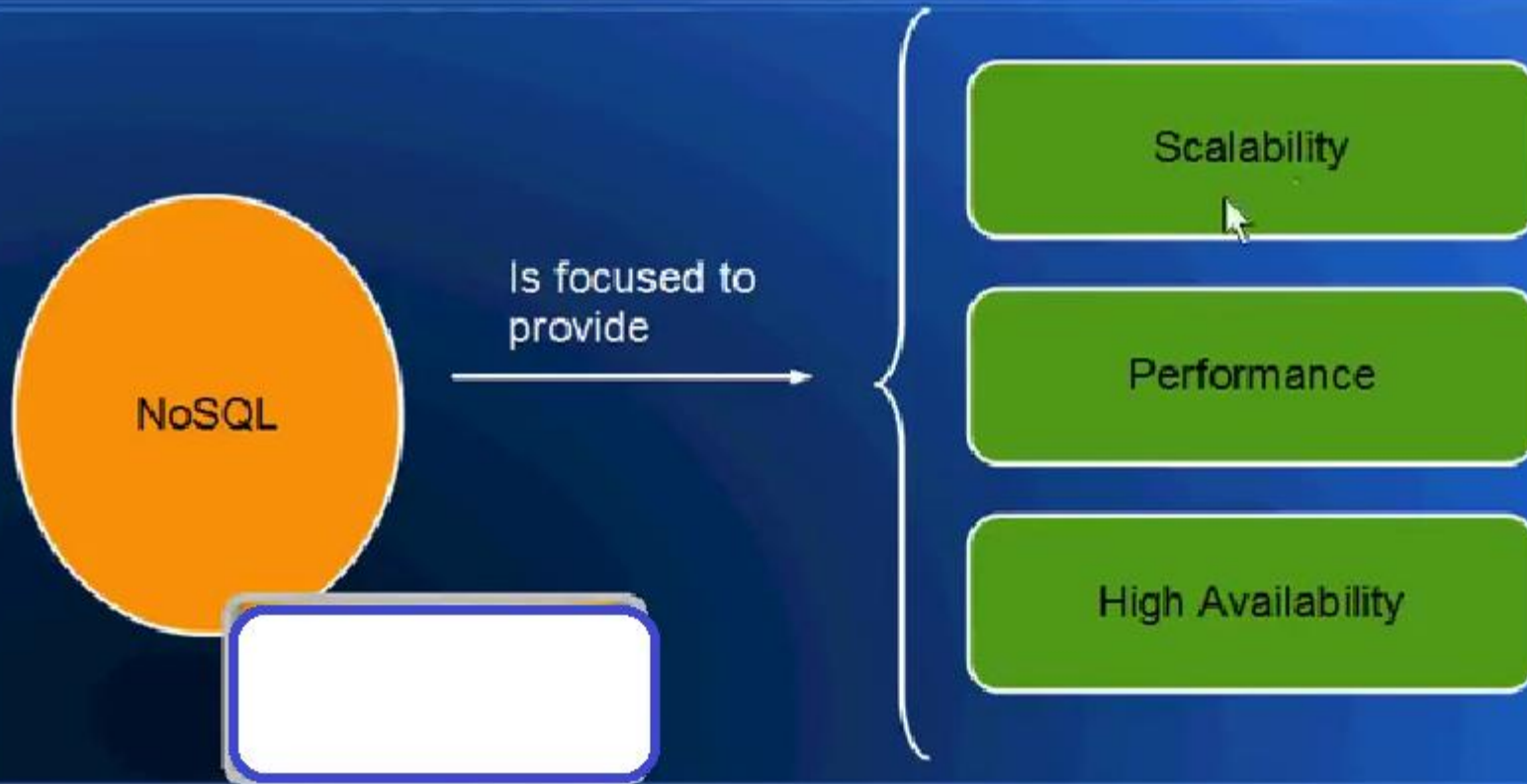
Database Management Systems

RDBMS  
(Relational)

OLAP

NoSQL

# Objective



# Performance

More functionality  
Less Performance

Less functionality  
More Performance

Database Management Systems

RDBMS  
(Relational)

OLAP

NoSQL



# Unstructured

Structured Data

Structured Data (or)  
Unstructured Data

Database Management Systems

RDBMS  
(Relational)

Tables

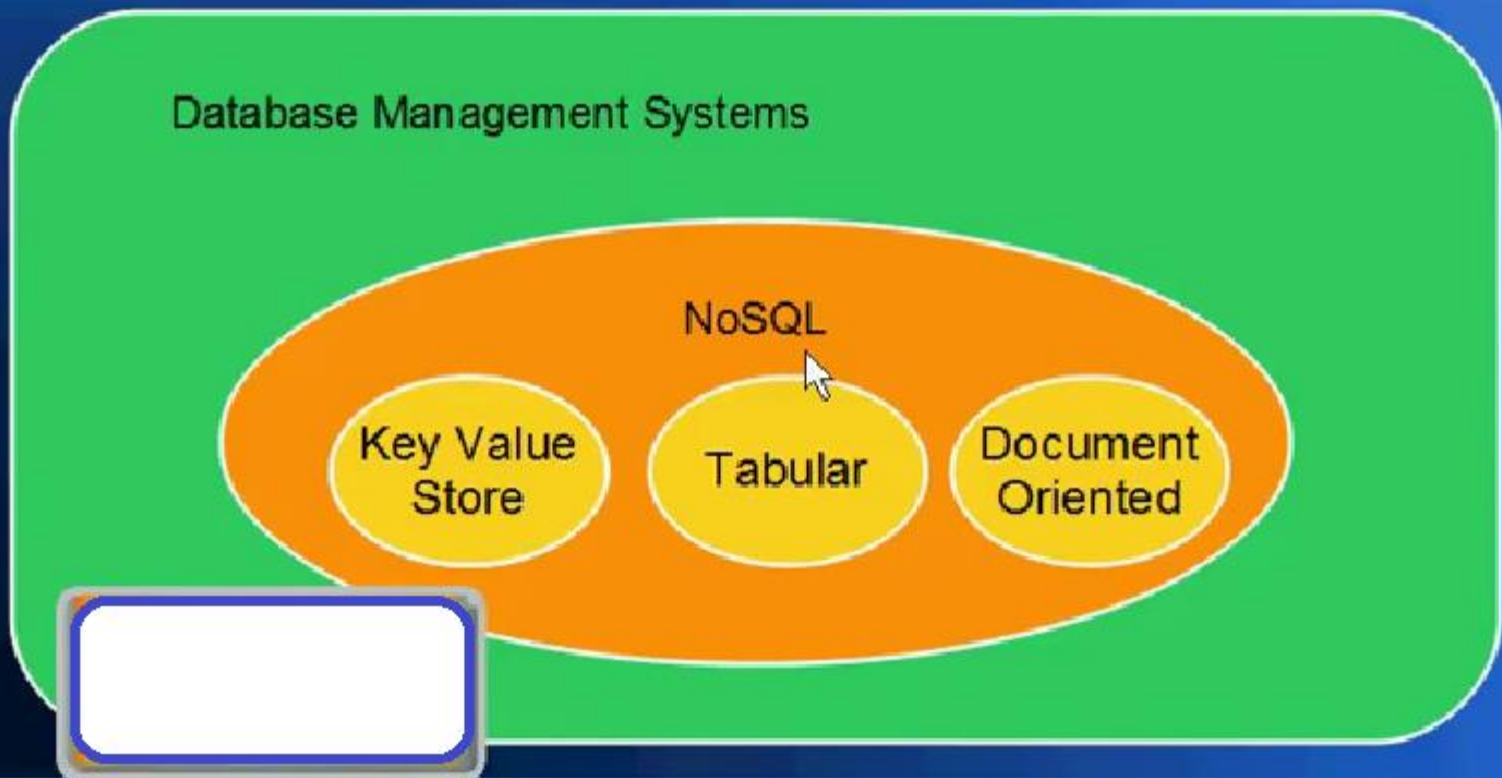
OLAP

Cubes

NoSQL

Collections

# Types of NoSQL Databases



Key Value  
Store

Examples:  
Memcached  
Coherence  
Redis



Tabular

Examples:  
BigTable  
Hbase  
Accumulo

Document  
Oriented

Examples:  
MongoDB  
Couch DB  
Cloudant

## NoSQL – What is Missing



# NoSQL – What is Available?

"Not Only SQL"

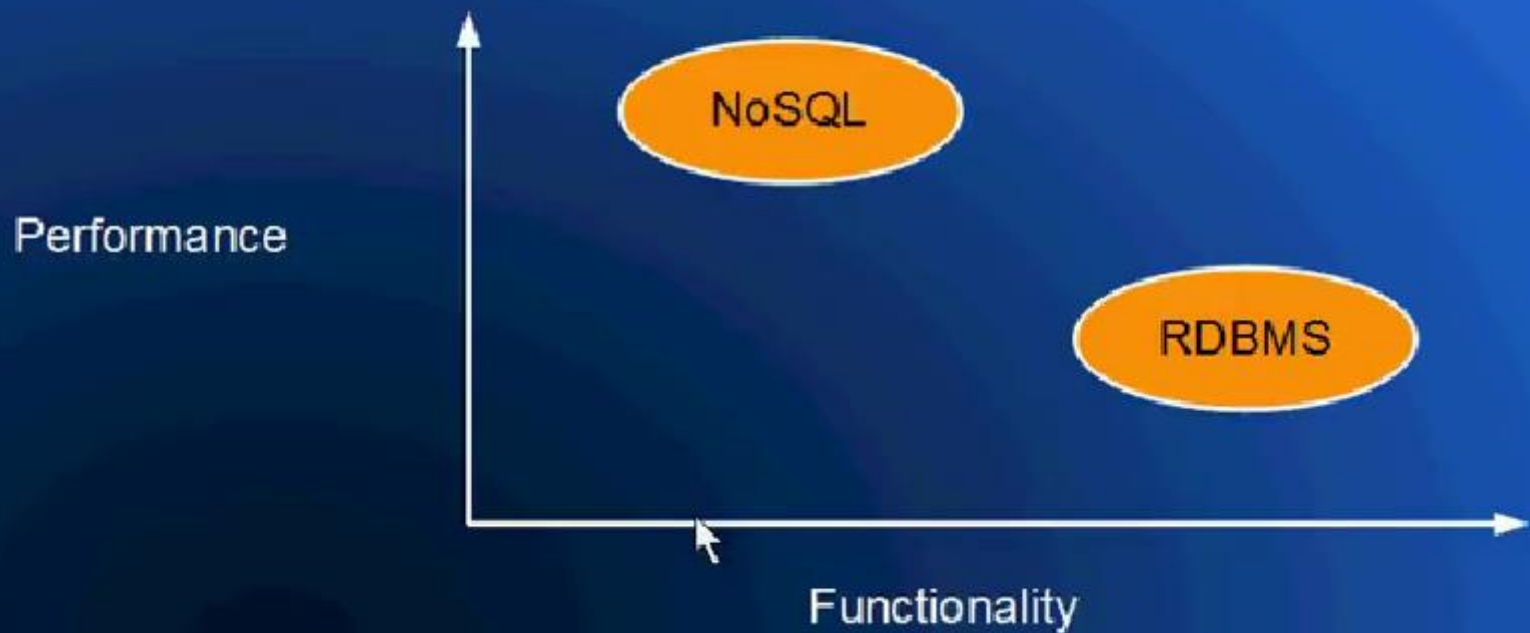
Query Language (Other than SQL)



Fast Performance

Horizontal Scalability

# NoSQL Attributes





## When to use NoSQL

The ability to store and retrieve great quantities of data is important

Storing relationships between the elements is not important

Dealing with growing lists of elements:  
Twitter posts, Internet server logs, Blogs

The data is not structured or the structure is changing with time

Prototypes or fast applications need to be developed

Constraints and validations logic is not required to be implemented in database

## When not to use NoSQL

**Complex Transactions need to be handled**

**Joins must be handled by databases**

**Validations must be handled by databases**



## Column-Oriented Storage

Each column is stored in a separate file

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M
3	Daffy	Duck	NY	10013	(212)227-1810	35	M
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M
5	Witch	Hazel	MA	01970	(978)744-0991	57	F

Each column for a given row is at the same offset  
(auto-indexing)

e.g. InfiniDB (e.g. find average age of all Male customers)

## Read Columns, Not Rows

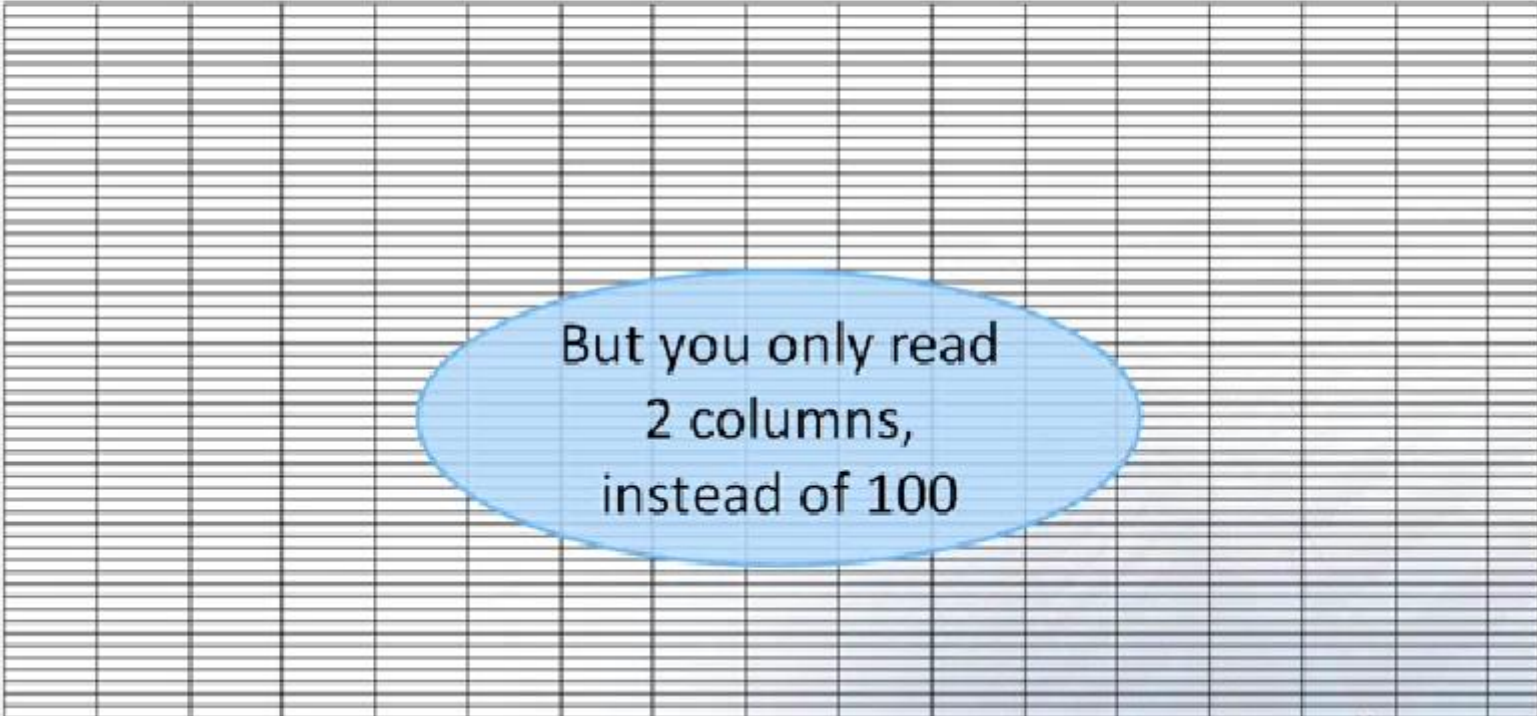
Only read the files you need

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718) 938-3235	34	M
2	Yosemite	Sam	CA	95389	(209) 375-6572	52	M
3	Daffy	Duck	NY	10013	(212) 227-1810	35	M
4	Elmer	Fudd	ME	04578	(207) 882-7323	43	M
5	Witch	Hazel	MA	01970	(978) 744-0991	57	F

Also get improved compression because all data in one file is the same data type.

## I/O Reduction

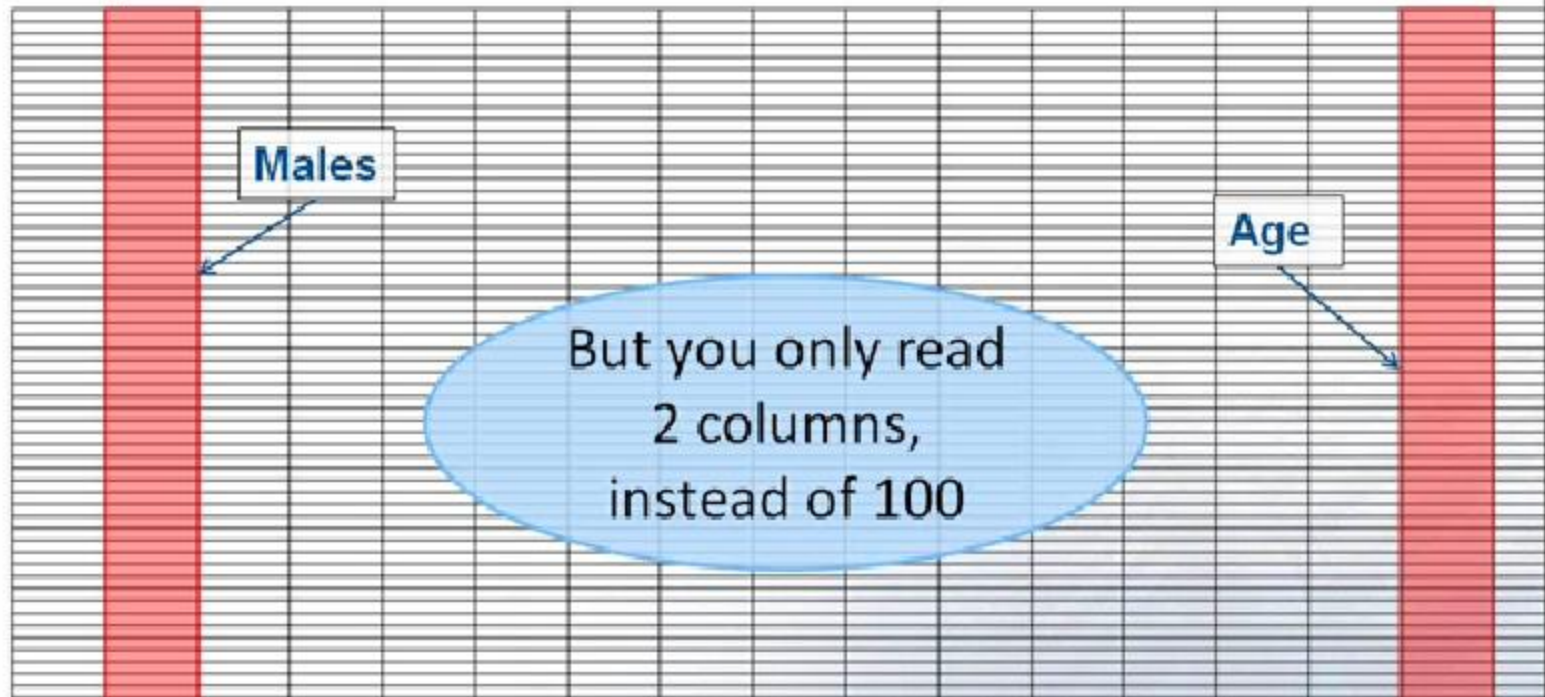
So you still have 100 million rows, with 100 columns...



But you only read  
2 columns,  
instead of 100

## I/O Reduction

So you still have 100 million rows, with 100 columns...



# Vertical Partitioning

Columnar databases produce automatic vertical partitioning

1	Bugs	Bunny	Brooklyn	NY	11217	(718) 938-3235
2	Yosemite	Sam	Wawona	CA	95389	(209) 375-6572
3	Daffy	Duck	New York	NY	10013	(212) 227-1810
4	Elmer	Fudd	Wiscasset	ME	04578	(207) 882-7323
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮	⋮	⋮	⋮
8m	Snoopy	Brown	Springfield	MA	01105	(413) 781-6500



# Horizontal Partitioning

InfiniDB also automatically creates horizontal partitions of 8 million rows (default)

1	Bugs	Bunny	Brooklyn	NY	11217	(718) 938-3235
2	Yosemite	Sam	Wawona	CA	95389	(209) 375-6572
3	Daffy	Duck	New York	NY	10013	(212) 227-1810
4	Elmer	Fudd	Wiscasset	ME	04578	(207) 882-7323
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
8m	Snoopy	Brown	Springfield	MA	01105	(413) 781-6500
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...
...	...	...	...	...	...	...

Knowing what values are in each partition allows for partition elimination at query time

## Bonus: Easy to Add a New Column

Row-oriented: Usually requires rebuilding table

Key	Fname	Lname	State	Zip	Phone	Age	Sex	Golf
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M	Y
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M	N
3	Daffy	Duck	NY	10013	(212)227-1810	35	M	Y
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M	Y
5	Witch	Hazel	MA	01970	(978)744-0991	57	F	N

Addition of  
column shifts  
every row

Column-oriented: Just create another file

Key	Fname	Lname	State	Zip	Phone	Age	Sex	Golf
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M	Y
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M	N
3	Daffy	Duck	NY	10013	(212)227-1810	35	M	Y
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M	Y
5	Witch	Hazel	MA	01970	(978)744-0991	57	F	N

## Single-Row Operations

Because of the nature of columnar storage, single-row operations can underperform.

***Do not attempt OLTP-style transactions on a columnar database.***

More details on individual DML statements follow...



## Single-Row Operations: Insert

Row-oriented: new rows appended to the end

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M
3	Daffy	Duck	NY	10013	(212)227-1810	35	M
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M
5	Witch	Hazel	MA	01970	(978)744-0991	57	F
6	Marvin	Martian	CA	91602	(818)761-9964	26	M

Columnar: new value must be added to each file

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M
3	Daffy	Duck	NY	10013	(212)227-1810	35	M
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M
5	Witch	Hazel	MA	01970	(978)744-0991	57	F
6	Marvin	Martian	CA	91602	(818)761-9964	26	M

## Insert: Solution

Do batch inserts and use `cpimport`, the bulk loader, instead.

***CPIMPORT is your friend.***

# Single-Row Operations: Delete

Row-oriented: row is deleted

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M
5	Witch	Hazel	MA	01970	(978)744-0991	57	F

Columnar: each column must be deleted from its file

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718)938-3235	34	M
2	Yosemite	Sam	CA	95389	(209)375-6572	52	M
4	Elmer	Fudd	ME	04578	(207)882-7323	43	M
5	Witch	Hazel	MA	01970	(978)744-0991	57	F

## Delete: Solutions

Do batch deletes.

Any extents that contain only data that is to be deleted can be dropped.

Otherwise, consider copying desired rows to a new table using the bulk loader and dropping the old table.

# Single-Row Operations: Update

Row-oriented: value replaced

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718) 852-2352	34	M
2	Yosemite	Sam	CA	95389	(209) 375-6572	52	M
3	Daffy	Duck	NY	10013	(212) 227-1810	35	M
4	Elmer	Fudd	ME	04578	(207) 882-7323	43	M
5	Witch	Hazel	MA	01970	(978) 744-0991	57	F

Column-oriented: value replaced

Key	Fname	Lname	State	Zip	Phone	Age	Sex
1	Bugs	Bunny	NY	11217	(718) 852-2352	34	M
2	Yosemite	Sam	CA	95389	(209) 375-6572	52	M
3	Daffy	Duck	NY	10013	(212) 227-1810	35	M
4	Elmer	Fudd	ME	04578	(207) 882-7323	43	M
5	Witch	Hazel	MA	01970	(978) 744-0991	57	F

*Yeah, this one just works.*



# Document Databases

- Each record in the database is a self-describing document
- Each document has an independent structure
- Documents can be complex
- All databases require a unique key
- Documents are stored using JSON or XML or their derivatives
- Content can be indexed and queried
- Offer auto-sharding for scaling and replication for high-availability



```
{
  "UID": "21f7f844-8051-5889-85",
  "Time": "2011-04-01T13:01:02.42",
  "Server": "A2223E",
  "Calling_Server": "A2223W",
  "Type": "E300",
  "Initiating_User": "dsallings@pspy.net",
  "Details": {
    "IP": "10.1.1.22",
    "API": "InsertDVDQueueItem",
    "Trace": "cdmsnd",
    "Tags": {
      "SERVER": "US/West",
      "API": "API"
    }
  }
}
```

# Relational vs Document data model

R1C1	R1C2	R1C3	R1C4
R2C1	R2C2	R2C3	R2C4
R3C1	R3C2	R3C3	R3C4
R4C1	R4C2	R4C3	R4C4

**Relational data model**

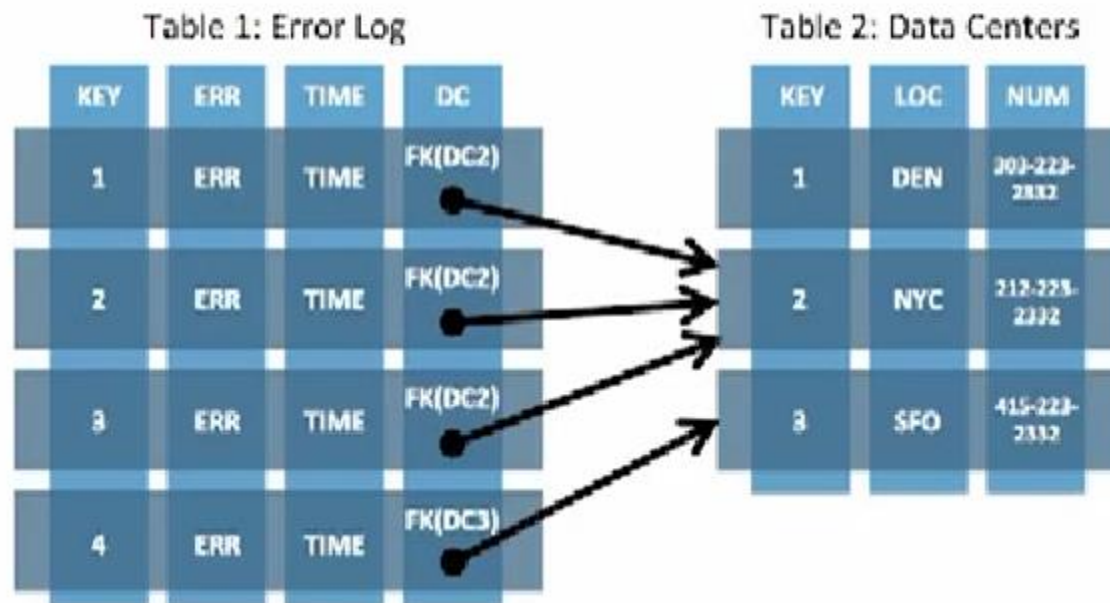
Highly-structured table organization with rigidly-defined data formats and record structure.



**Document data model**

Collection of complex documents with arbitrary, nested data formats and varying "record" format.

## Example: Error Logging Use case





## Document design with flexible schema

```
{
  {
    {
      {
        "ID": 1,
        "ERR": "Out of Memory",
        "TIME": "2004-09-16T23:59:58.75",
        "DC": "NYC",
        "NUM": "212-223-2332"
      }
    }
  }
}
```

### SCHEMA CHANGE

```
{
  "ID": 5,
  "ERR": "Out of Memory",
  "TIME": "2004-09-16T23:59:58.75",
  "COMPONENT": "DMS",
  "SEV": "LEVEL1",
  "DC": "NYC",
  "NUM": "212-223-2332"
}
```

MongoDB sponsored by 10gen

# Terminology

database → database

table → collection

row → document

4 MB limit on document size

No limit on nesting depths

# JSON-style Documents

represented as *BSON*

```
{“hello”: “world”}
```



```
\x16\x00\x00\x00\x02hello  
\x00\x06\x00\x00\x00world  
\x00\x00
```

JavaScript Object Notation

[Json.org](http://Json.org)/[Bsonspec.org](http://Bsonspec.org)

# Flexible “Schemas”

```
{“author”: “mike”,  
  “text”: “...”}  
  
{“author”: “eliot”,  
  “text”: “...”,  
  “tags”: [“mongodb”]}
```

## New Post

```
post = {author: "mike",  
        date: new Date(),  
        text: "my blog post...",  
        tags: ["mongodb", "intro"]}  
  
db.posts.save(post)
```

Db.posts is a collection (if it does not exist it will create it)  
It is going to save that post into that collection

# \_id

if not specified drivers will add default:

```
ObjectId("4bface1a2231316e04f3c434")  
      timestamp  
              machine id  
                      process id  
                              counter
```

\_id is lightweight occupying 12 bytes of storage  
Generated on client side to reduce load on database server

# Dynamic Queries



# Posts by Author

```
db.posts.find({author: "mike"})
```

## Last 10 Posts

```
db.posts.find()  
    .sort({date: -1})  
    .limit(10)
```

Find is a method

## Posts Since August 1<sup>st</sup>

```
aug_1 = new Date(2010, 7, 1)  
db.posts.find({date: {$gt: aug_1}})
```

\$gt operator stands for greater than  
Javascript month starts with 0

# Atomic Update Modifiers

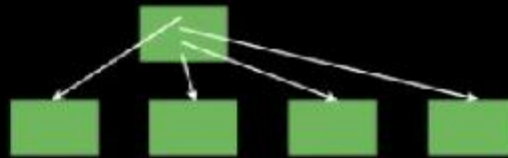
```
c = {author: "eliot",  
     date: new Date(),  
     text: "great post!"}  
  
db.posts.update({_id: post._id},  
                {$push: {comments: c}})
```

\$ sign update modifier

New comments will keep on appending so no worry about locking

# Indexing

B-Tree indexes



```
db.posts.ensureIndex({tags: 1})
```

# Indexing / Querying on Embedded Docs

(dot notation)

```
db.posts.ensureIndex({"comments.author": 1})
```

```
db.posts.find({"comments.author": "eliot"})
```

# Many Supported Platforms / Languages

Drivers for PHP, Perl, C++, C#, .Net, Python, Ruby, etc.  
All OS also 32/64 bit

# Focus on Performance

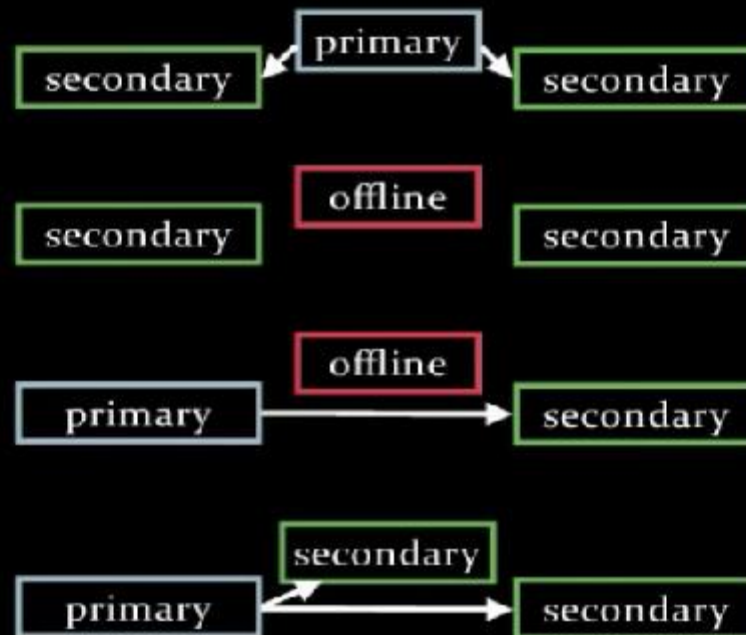
Does not use REST (http protocol)

Binary tcp wired protocol to communicate between client and server

Does not wait for Success Response to operations, proceeds to next one



# Replica Sets

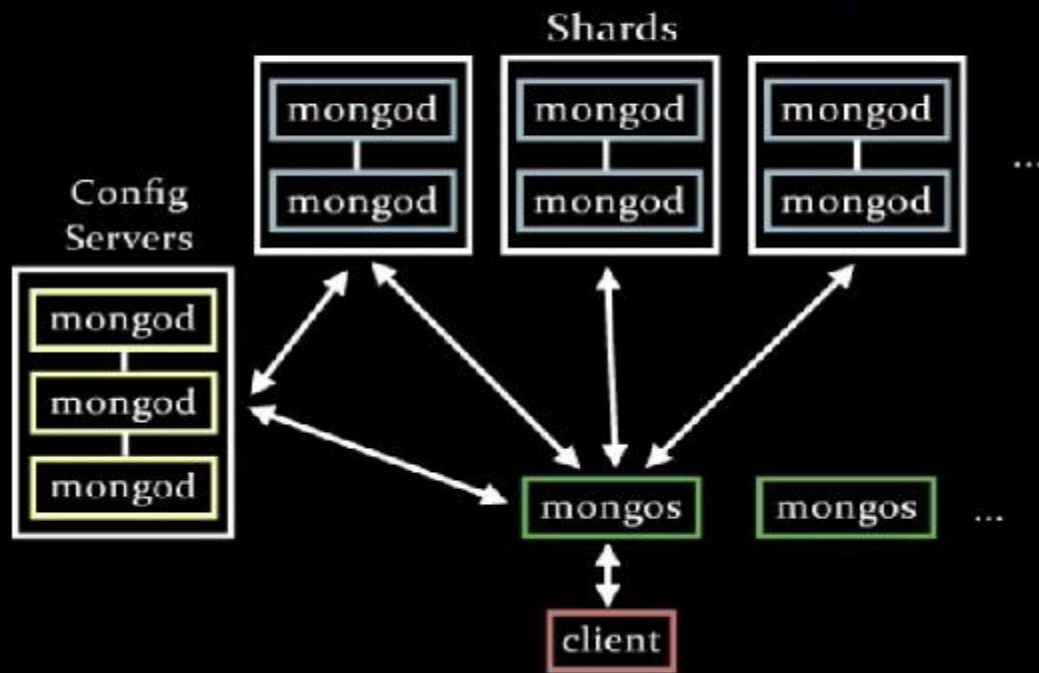


Master/Slave

Primary is not fixed (Automatic failover)

If Primary goes offline then other nodes elect a new Primary

# Auto-sharding



Split up the data, each shard has a subset of the data

Mongos is the process that distributes queries and writes to the shards

Scale out when needed not necessarily in advance

# Other Cool Stuff

aggregation and map/reduce

capped collections

unique indexes

GridFS

geo

