



**CME 213**  
**SPRING 2017**

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## PTHREADS

- `pthread_create`, `pthread_exit`, `pthread_join`
- **Mutex:** locked/unlocked; used to protect access to shared variables (read/write)
- **Condition variables:**
  - used to allow threads to become idle and wake up when a condition becomes true.
  - used in combination with a mutex to protect access to the condition variable (boolean)
  - `cond_wait`
  - `cond_signal`



## PTHREADS/OPENMP

- Pthreads gives you maximum flexibility.
- It's a low level API that allows you to implement pretty much any parallel computation exactly the way you want it.
- However, in many cases, the user only wants to parallelize certain common situations:
  - For loop: partition the loop into chunks and have each thread process one chunk.
  - Hand-off a block of code (computation) to a separate thread
- This is where OpenMP is useful. It simplifies the programming significantly.
- In some cases, adding one line in a C code is sufficient to make it run in parallel.
- As a result, OpenMP is the standard approach in scientific computing for multicore processors.

## OPENMP

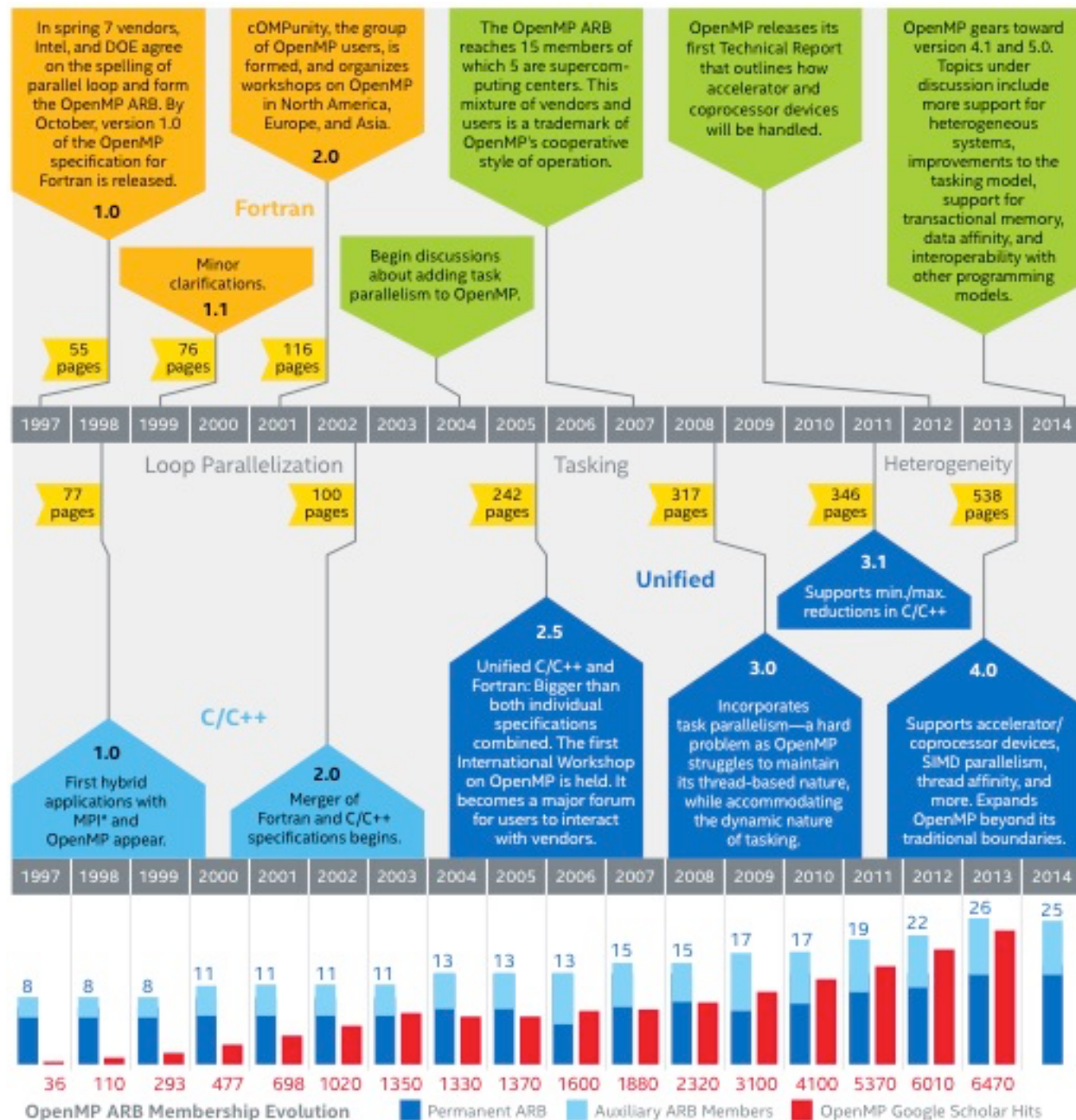
### What is OpenMP?

- OpenMP is an Application Programming Interface (API), jointly defined by a group of major computer hardware and software vendors.
- OpenMP provides a portable, scalable model for developers of shared memory parallel applications.
- The API supports C/C++ and Fortran on a wide variety of architectures.

Hence, it is more portable and general than Pthreads.

- OpenMP website: [openmp.org](http://openmp.org)
- Wikipedia: [en.wikipedia.org/wiki/OpenMP](http://en.wikipedia.org/wiki/OpenMP)

## ARB: Architecture Review Board



## COMPILING YOUR CODE

### First things first

- Header file:

```
#include <omp.h>
```

- This is only needed if you explicitly use the OpenMP API.
- Compiler flags:

Compiler	Flag
icc icpc ifort	-openmp
gcc g++ g77 gfortran	-fopenmp

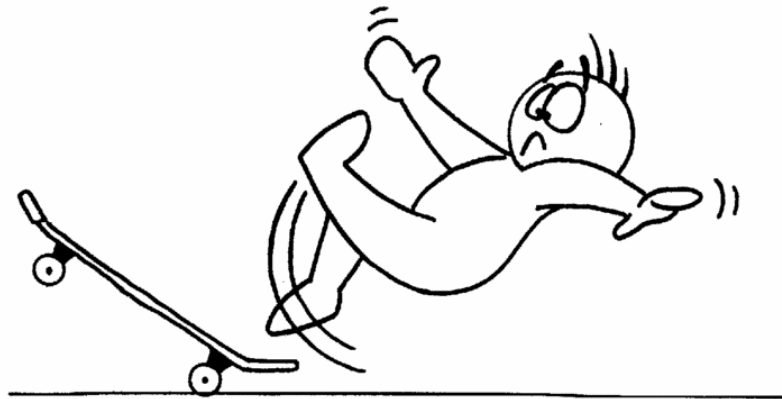
A large, light gray watermark of the Stanford University seal is centered in the background. The seal is circular with a diamond-patterned border. Inside the border, the text "LELAND STANFORD JUNIOR UNIVERSITY" is written in a circular path. Below this, the German phrase "DIE LUFT DER FREIHEIT WEHT" is written. In the center of the seal is a detailed illustration of a redwood tree standing on a rocky outcrop. At the bottom of the seal, the year "1891" is inscribed.

## **PARALLEL REGIONS**



## DIRECTIVES

- OpenMP is based on directives.
- Powerful because of simple syntax.
- Dangerous because you may not understand exactly what the compiler is doing.

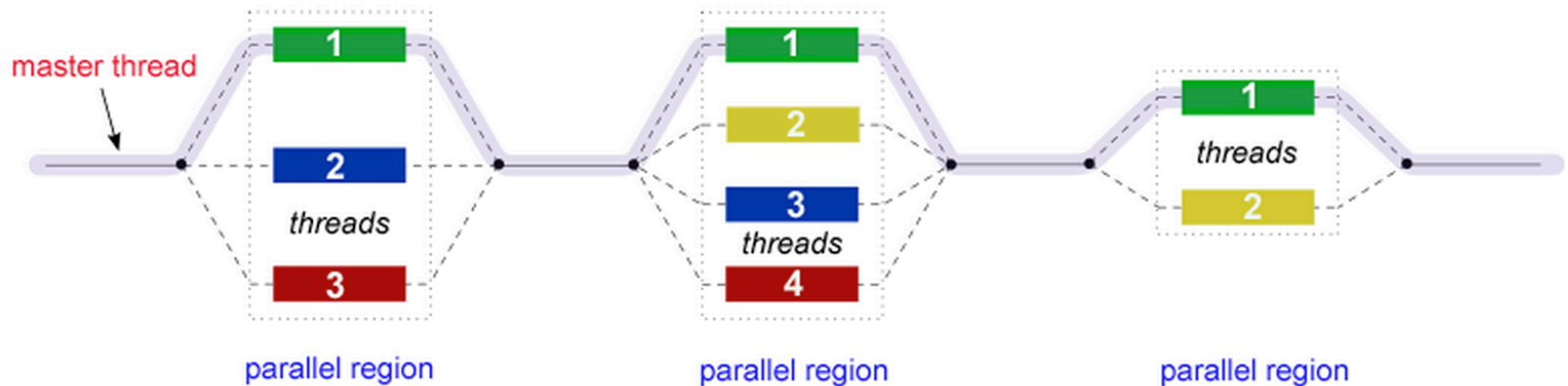


## PARALLEL REGION

The most basic directive is

```
#pragma omp parallel  
{ // structured block ... }
```

This starts a new parallel region. OpenMP follows a fork-join model:



Upon entering a region, if there are no further directives, a team of threads is created and all threads execute the code in the parallel region.

## BASIC EXAMPLE

`hello_world_openmp.cc`

Compilation:

```
g++ -o hello_world hello_world_openmp.cc -fopenmp
```

3. 1415926535897932384626433832795028841971693993751058209749445  
923078164062862089986280348253421170679821480865132823066470938  
446095505822317**253594081284811174502841027019385211055**59644  
62294895493**038196442881097566593344612847564823378678**31652  
712019091**45648566923460348610454326648213393607260249**14127  
37245870**066063155881748815209209628292540917153643678**92590  
3600113**3053054882046652138414695194151160943305727036**57595  
919530**9218611738193261179310511854807446237**996274956735188575  
27248**91227938183011949129833673362440656643**086021394946395224  
73719**07021798609437027705**39217176293**1767523**8467481846766940513  
2000**568127145263560827785**77134275778**9609173**6371787214684409012  
24953430146549585371**105079**22796892589**2354201**9956112129021960864  
03441815981362977477**13099**60518707211**3499999**9837297804995105973  
1732816096318595024**459455**34690830264**2522308**2533446850352619311  
8817101000313783875**28865**875332083814**2061717**7669147303598253490  
4287554687311595628**63882**35378759375**1957781**85778053217122680661  
300192787661119590**9216420**1989380952**5720106**54858632788659361533  
818279682303019520**353018**52968995773**6225994**13891249721775283479  
13151557485724245**4150695**9508295331**16861727**85588907509838175463  
7464939319255060**40092770**16711390098**4882401**28583616035637076601  
047101819429555**9619894**6767837449448**2553797**74726847104047534646  
20804668425906**94912933**136770289891**52104752**1620569660240580381  
5019351125338**243003558**764024749647**32639141**992726042699**227**9678  
23547816360**09341721641**219924586315**030286182**9745557067**498**38505  
4945885869**26995690927**2107975093029**553211653**449872027**559**602364  
806654991**19881834797**75356636980742**6542527862551818417**5746728  
90977772**793800081647**060016145249192**17321721477235014**14419735  
68548161**36115735255**21334757418494684**385233239073941**433345477  
624168625**189835694**8556209921922218427**2550254256887**67179049460  
165346680498**8627**2327917860857843838279679**7668145**41009538837863  
609506800642251252051173929848960841284886269456042419652850222  
106611863067442786220391949450471237137869609563643719172874677

## PI ALGORITHM

In our code, Pi is computed using:

$$\frac{\pi}{2} = 1 + \frac{1}{3} \left( 1 + \frac{2}{5} \left( 1 + \frac{3}{7} \left( 1 + \frac{4}{9} \left( 1 + \dots \right) \right) \right) \right)$$

Using this expansion, can you show that the code computes the digits of Pi, **4 at a time**, assuming that:

$$\text{carry} + \text{sum} / \text{SCALE} < 10,000$$

- Is the previous algorithm parallel?
- Is this a good multicore implementation?
- How would you improve it?

- Computing pi in parallel is difficult.
- Many algorithms use sequential calculations using **high-precisions arithmetic**, that is you compute using numbers with a lot of digits.
- This leads to the natural question:

Is it possible to compute the  
n-th digit of pi independently  
from the others?

# BAILEY–BORWEIN–PLOUFFE FORMULA

```

3. 1415926535897932384626433832795028841971693993751058209749445
923078164062862089986280348253421170679821480865132823066470938
44609550582231725359408128481117450284102701938521105559644
6229489549303819644288109756659334461284756482337867831652
7120190914564856692346034861045432664821339360726024914127
3724587006606315588174881520920962829254091715364367892590
3600113305305488204665213841469519415116094330572703657595
9195309218611738193261179310511854807446237996274956735188575
2724891227938183011949129833673362440656643086021394946395224
73719070217986094370277053921717629317675238467481846766940513
20005681271452635608277857713427577896091736371787214684409012
24953430146549585371050792279689258923542019956112129021960864
03441815981362977477130996051870721134999999837297804995105973
17328160963185950244594553469083026425223082533446850352619311
88171010003137838752886587533208381420617177669147303598253490
42875546873115956286388235378759375195778185778053217122680661
30019278766111959092164201989380952572010654858632788659361533
81827968230301952035301852968995773622599413891249721775283479
13151557485724245415069595082953311686172785588907509838175463
74649393192550604009277016711390098488240128583616035637076601
04710181942955596198946767837449448255379774726847104047534646
2080466842590694912933136770289891521047521620569660240580381
5019351125338243003558764024749647326391419927260426992279678
2354781636009341721641219924586315030286182974555706749838505
4945885869269956909272107975093029553211653449872027559602364
806654991198818347977535663698074265425278625518184175746728
90977772793800081647060016145249192138523323907394114419735
685481613611573525521334757418494661385233239073941133345477
62416862518983569485562099219222184135562312566871179049460
1653466804986272327917860857843838279679766814541009538837863
609506800642251252051173929848960841284886269456042419652850222
106611863067442786220391949450471237137869609563643719172874677
    
```

$$\pi = \sum_{k=0}^{\infty} \frac{1}{16^k} \left( \frac{4}{8k+1} - \frac{2}{8k+4} - \frac{1}{8k+5} - \frac{1}{8k+6} \right)$$



## COMPUTING THE N-TH BIT

This problem can now be reformulated as:

Can we compute the fractional part of

$$16^n \pi$$

Take:

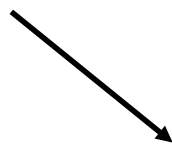
**Only a few terms  
are needed**

$$\sum_{k=0}^{\infty} \frac{16^{n-k}}{8k+1} = \sum_{k=0}^n \frac{16^{n-k}}{8k+1} + \sum_{k=n+1}^{\infty} \frac{16^{n-k}}{8k+1}$$

**whole numbers  
can be removed**

## COMPUTING THE FIRST SUM

$$\sum_{k=0}^n \frac{16^{n-k}}{8k+1}$$



$$\sum_{k=0}^n \frac{16^{n-k} \bmod (8k+1)}{8k+1}$$

**This can be easily  
computed**

## CLAUSE

- This is one of the tricky points of OpenMP.
- Recall in Pthreads that:
  - Variables passed as argument to a thread are **shared** (they might be pointers in a `struct` for example)
  - **Variables** inside the function that a thread is executing are **private** to the thread.
- OpenMP needs a similar mechanism: some variables are going to be shared (all threads can read and write), others need to be private.
- There are (complicated) rules to figure out whether a variable is private or shared.

## SHARED/PRIVATE

- See `shared_private_omp.cc`
- In a `parallel` construct, variables defined outside are shared by default.
- You can declare explicitly whether a variable is shared or private using

```
private(variable_name)  
shared(variable_name)
```

the value of a private variable like `should_be` after the program exits the parallel region?

The value it had during the last assignment on line 36

0

Undefined

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## What is the problem with read\_write?

All threads are  
trying to write

All threads are  
trying to read

All threads are

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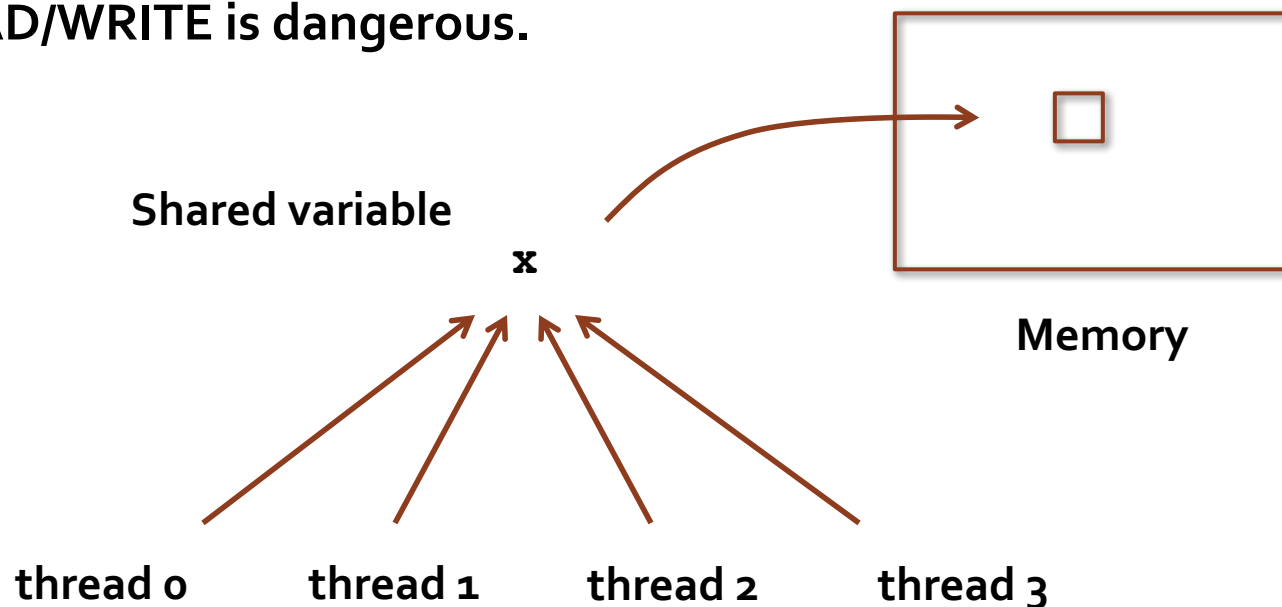
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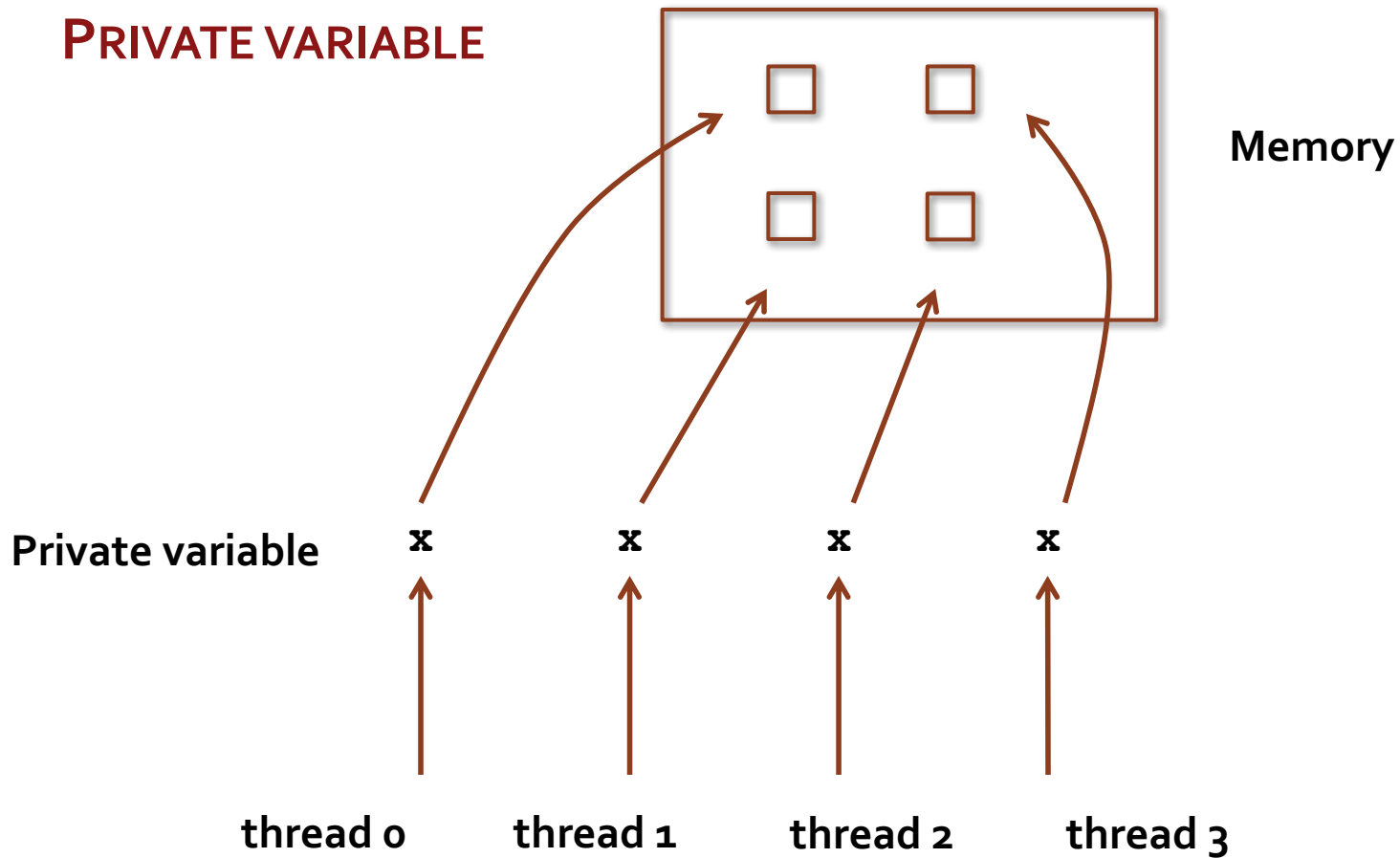
## SHARED VARIABLE

- Those are typically READ ONLY.
- Using a shared variable w/ READ/WRITE is dangerous.



Variable refers to the same memory location for all threads.

## PRIVATE VARIABLE



Variable refers to a different memory location for each thread. Those variables are typically READ/WRITE.



# WORKSHARING CONSTRUCTS

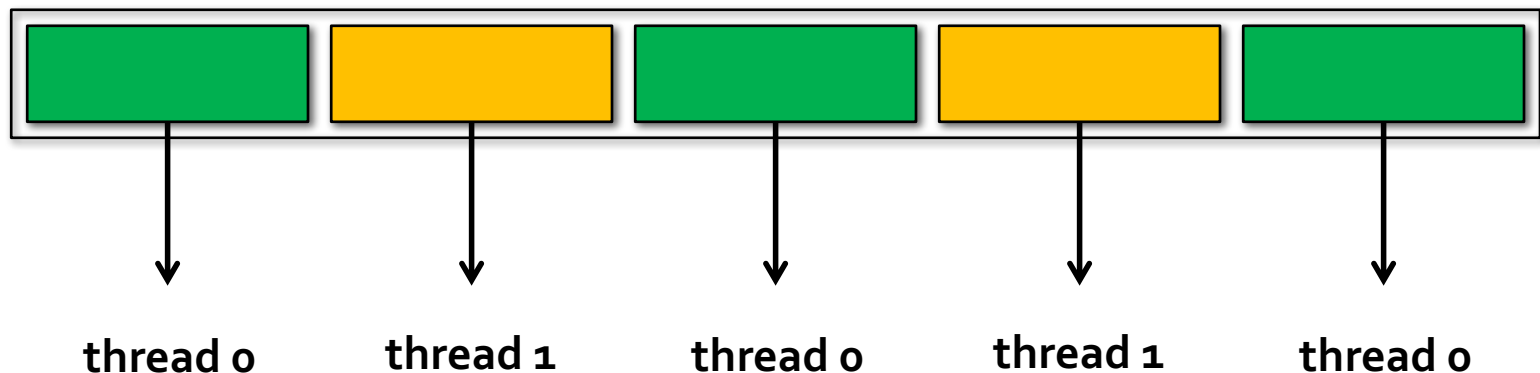


## PARALLEL FOR LOOP

The most common approach to parallelize a computation on a multicore processor is to parallelize a `for` loop.

OpenMP has some special constructs to do that.

```
#pragma omp for [clause [clause] ... ]  
for (i = lower bound; i op upper bound; incr expr) {  
    ...  
}
```



## EXAMPLE

- Let's consider again the matrix-matrix example we used for Pthreads.
- See `matrix_prod_openmp.cc`
- One line of code is sufficient to parallelize the calculation! This is the power of OpenMP.

```
$ ./matrix_prod_openmp -n 4000 -p 32
```

```
$ top
```

Is it possible to move #pragma omp parallel for to line 122?

Yes, it makes no difference

Yes but this is not recommended

No

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## SCHEDULING FOR LOOPS

- How are the iterates in a for loop split among threads? This is important to fine-tune the optimization of your code.
  - This is a problem of load-balancing: how should we distribute the work so that we minimize the total execution time?
1. `schedule(static, block_size)`: iterations are divided into pieces of size `block_size` and then statically assigned to threads. This is the best **default** option.
  2. `schedule(dynamic, block_size)`: iterations are divided into pieces of size `block_size`, and dynamically scheduled among the threads; when a thread finishes one chunk, it is dynamically assigned another. This is useful when the **work per iteration is irregular**.
  3. `schedule(guided, block_size)`: specifies a dynamic scheduling of blocks but with decreasing size. It is appropriate for the case in which the **threads arrive at varying times** at a `for` construct (with each iteration requiring about the same amount of work).

## OTHER WORKSHARING CONSTRUCTS: SECTIONS



- There are situations where two independent pieces of work can be executed concurrently. For example, we may need to update two vectors independently.
- In that case, we would like to assign one thread to do each operation in parallel.
- This can be done using sections.
- The compiler is allowed to schedule the execution of the code inside each section concurrently.
- See `section.cc`

How many threads have work to do in section.c

Default  
number of  
threads (24)

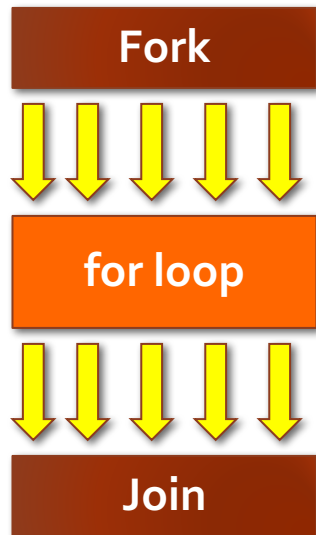
Two  
threads

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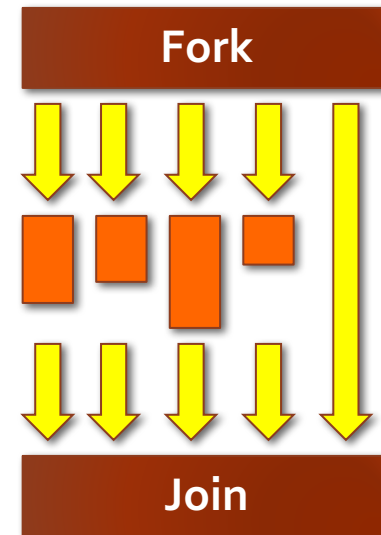
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## SUMMARY



Large number of  
iterates.  
Parallel for loop.



Small and fixed number  
of independent tasks.  
Parallel sections.