

Quantum Computers

Last updated: March 4th 2019, at 9.42am

Contents

1	Limits of Moore's Law	2
1.1	Moore's Law	2
1.2	Parallelism	3
1.2.1	Example	3
1.2.2	Exponential Parallelism	3
2	Implementing Qubits	3
2.1	Ion Traps	3
2.2	Linear Optics	4
2.3	Others	4
3	Quantum Circuits	4
3.1	Reversible gates	4
3.2	Qubits	5
3.3	Quantum gates	5
3.3.1	Single qubit gates	5
3.3.2	Multiple qubit gates	6
3.4	Universal quantum gate sets	8
3.5	Example — Deutsch's Algorithm	8
3.5.1	Problem statement	8
3.5.2	Classical circuits	9
3.5.3	Quantum circuit	9
3.5.4	Deutsch's circuit	10

Why?

Complexity and Intractability

What:

- is 2×7 ?
- are the factors (divisors) of 14?

Quantum Computers

- is 5×11 ?
- are the factors of 55?
- is 13×19 ?
- are the factors of 247?
- is 229×557 ? (you may use pen and paper)
- are the factors of 127, 553? (you may use pen and paper)
- is $573,260,813 \times 879,193,169$? (you may use a calculator)
- are the factors of 504,006,965,615,712,893? (you may use a calculator)

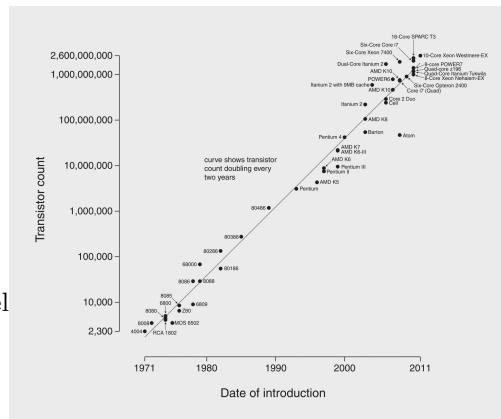
1 Limits of Moore's Law

1.1 Moore's Law

The processing power of chips doubles every...

- ... year (1965)
- ... two years (1975)

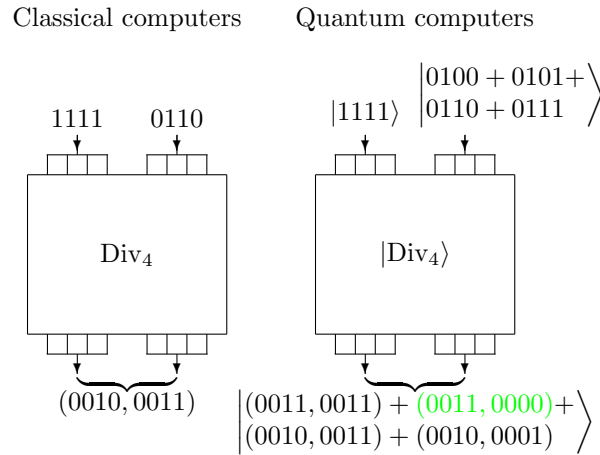
Gordan Moore, Intel



But difficult to make a chip smaller than a hydrogen atom

1.2 Parallelism

1.2.1 Example



1.2.2 Exponential Parallelism

One ...

bit Zero *or* one

qubit Zero *and* one

Two ...

bits Zero *or* one *or* two *or* three

qubits Zero *and* one *and* two *and* three

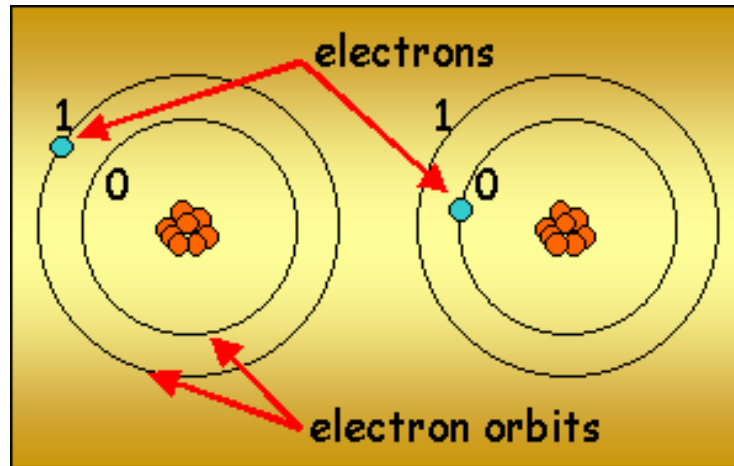
...

A 16 qubit word represents 65,536 values simultaneously, a 32 qubit word 4,294,967,296 values, and a 64 qubit word 18,446,744,073,709,551,616 values.

2 Implementing Qubits

2.1 Ion Traps

Use electron orbits to represent bits



- Ion trapped by electromagnetic field
- Use lasers to set and measure states

Long coherence time, reliable, but slow, and difficult to scale.

2.2 Linear Optics

- Uses polarisation of photons
- Difficult to entangle

2.3 Others

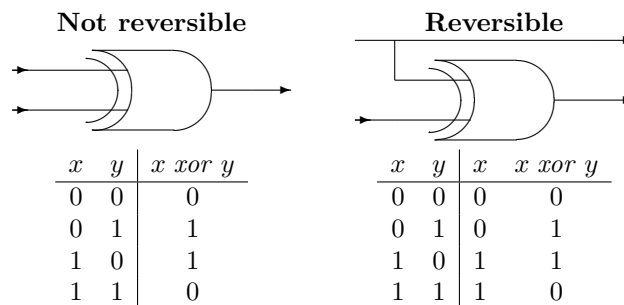
NMR Qubit = spin state of many molecules in a fluid

SQP¹ Qubit = frequency of oscillations in superfluids

3 Quantum Circuits

3.1 Reversible gates

All quantum gates must be *reversible*. E.g.



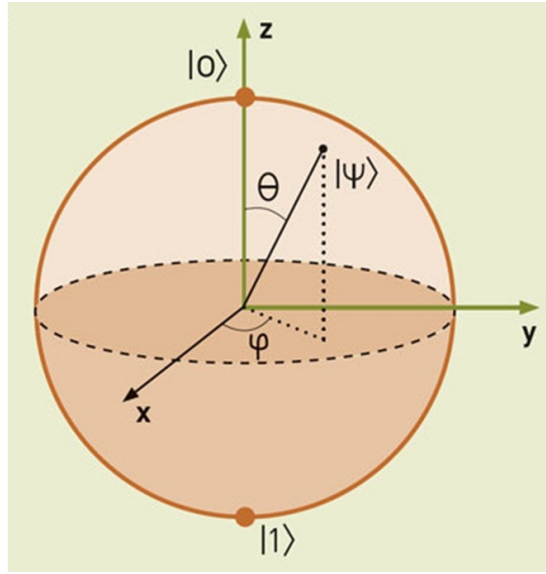
3.2 Qubits

A qubit is a matrix with complex numbers:

$$\begin{bmatrix} c_0 \\ c_1 \end{bmatrix}$$

with $|c_0|^2 + |c_1|^2 = 1$ and $|c_n|^2$ (with $n \in \{0, 1\}$) the probability the qubit is in state $|n\rangle$.

A single qubit can be represented as a point on a *Bloch sphere*.



- Latitude — probability of $|0\rangle$, $|1\rangle$
- Longitude — evolution

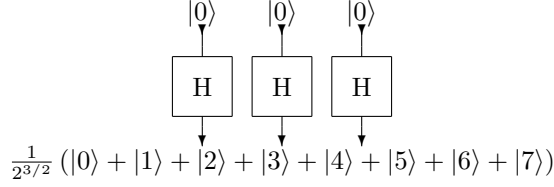
3.3 Quantum gates

3.3.1 Single qubit gates

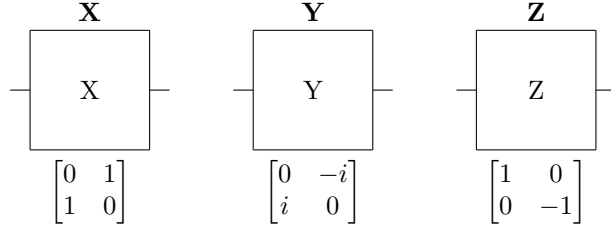
Hadamard gate

$$\boxed{\text{H}} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$H|0\rangle = \frac{|0\rangle + |1\rangle}{\sqrt{2}}, H|1\rangle = \frac{|0\rangle - |1\rangle}{\sqrt{2}}$$



Pauli gates



Rotate the Bloch sphere through 180° around the x, y, z axes

Square root of NOT

$$\sqrt{\text{NOT}} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & -1 \\ 1 & 1 \end{bmatrix}$$

Rotational gates Let $\vec{v} = (x, y, z)$ be a unit vector in the Bloch sphere, then

$$R_{\vec{v}}(\theta) = \cos \frac{\theta}{2} I - i \sin \frac{\theta}{2} (xX + yY + zZ)$$

rotates the Bloch sphere round \vec{v} by θ .

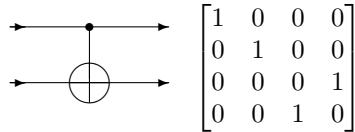
Special cases:

$$\begin{aligned}
 R_{\vec{x}}(\theta) &= \cos \frac{\theta}{2} I - i \sin \frac{\theta}{2} X = \begin{bmatrix} \cos \frac{\theta}{2} & -i \sin \frac{\theta}{2} \\ -i \sin \frac{\theta}{2} & \cos \frac{\theta}{2} \end{bmatrix} \\
 R_{\vec{y}}(\theta) &= \cos \frac{\theta}{2} I - i \sin \frac{\theta}{2} Y = \begin{bmatrix} \cos \frac{\theta}{2} & -\sin \frac{\theta}{2} \\ \sin \frac{\theta}{2} & \cos \frac{\theta}{2} \end{bmatrix} \\
 R_{\vec{z}}(\theta) &= \cos \frac{\theta}{2} I - i \sin \frac{\theta}{2} Z = \begin{bmatrix} e^{-i\theta/2} & 0 \\ 0 & e^{i\theta/2} \end{bmatrix}
 \end{aligned}$$

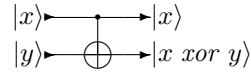
$e^{i\theta/2} R_{\vec{z}}(\theta) = \begin{bmatrix} 1 & 0 \\ 0 & e^{\theta} \end{bmatrix}$ is known as a *phase shift* gate

3.3.2 Multiple qubit gates

Controlled not



Quantum Computers

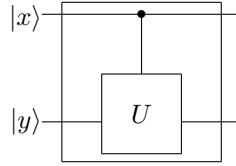


Controlled U If U is a single qubit gate then a controlled U gate is

```

if (x == 0) {
    0,y;
} else {
    1,U(y);
}

```



If $U = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$ then

$${}^C U = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & a & b \\ 0 & 0 & c & d \end{bmatrix}$$

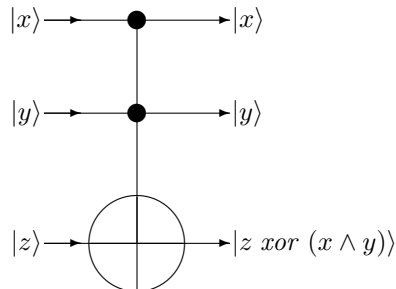
This generalises to n -qubit gates. E.g. if U_2 is

$$\begin{bmatrix} u_{0,0} & u_{0,1} & u_{0,2} & u_{0,3} \\ u_{1,0} & u_{1,1} & u_{1,2} & u_{1,3} \\ u_{2,0} & u_{2,1} & u_{2,2} & u_{2,3} \\ u_{3,0} & u_{3,1} & u_{3,2} & u_{3,3} \end{bmatrix}$$

then a ${}^C U_2$ gate is

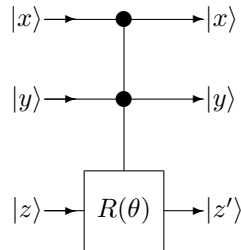
$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & u_{0,0} & u_{0,1} & u_{0,2} & u_{0,3} \\ 0 & 0 & 0 & 0 & u_{1,0} & u_{1,1} & u_{1,2} & u_{1,3} \\ 0 & 0 & 0 & 0 & u_{2,0} & u_{2,1} & u_{2,2} & u_{2,3} \\ 0 & 0 & 0 & 0 & u_{3,0} & u_{3,1} & u_{3,2} & u_{3,3} \end{bmatrix}$$

Toffoli gate



The Toffoli gate is ${}^C ({}^C \text{NOT})$

Deutsch gates



```

if (|x>==|1> && |y>==|1>) {
    |z'> = R(θ) |z>
} else {
    |z'> = |z>
}

```

3.4 Universal quantum gate sets

- $\{H, {}^C\text{NOT}, R(\cos^{-1} \frac{3}{5})\}$
- $\{D(\theta)\}$, for some θ for which $\frac{\pi}{\theta}$ is irrational

are both universal quantum gate sets

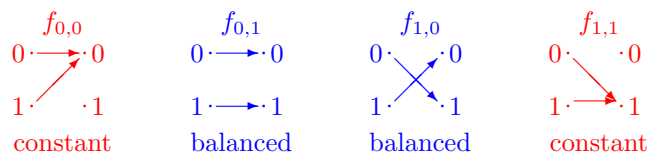
3.5 Example — Deutsch's Algorithm

Implementing algorithms

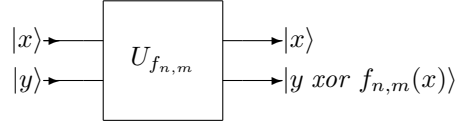
- Start in a classical state
- Move to a superposition of states
- Act on the superposition
- Measure qubits

3.5.1 Problem statement

Considers functions from $\{0, 1\}$ to $\{0, 1\}$.



Given a $U_{f_{n,m}}$ “black box”



decide if $f_{n,m}$ is constant or balanced

3.5.2 Classical circuits

Table for, e.g., $U_{f_{1,0}}$

x	y	$f_{1,0}(x)$	$y \text{ xor } f_{1,0}(x)$	$U_{f_{1,0}}(x, y)$
0	0	1	1	01
0	1	1	0	00
1	0	0	0	10
1	1	0	1	11

All functions

x	y	$U_{f_{0,0}}$	$U_{f_{0,1}}$	$U_{f_{1,0}}$	$U_{f_{1,1}}$
0	0	00	00	01	01
0	1	01	01	00	00
1	0	10	11	10	11
1	1	11	10	11	10

On each line:

- boxed outputs are identical
- unboxed outputs are identical
- one boxed output is from a **constant** function, one from a **balanced**
- one unboxed output is from a **constant** function, one from a **balanced**

so cannot find input that will discriminate

3.5.3 Quantum circuit

Constructing matrices Construct matrix for, e.g. $U_{f_{1,0}}$

x	0	0	1	1
y	0	1	0	1
$f_{1,0}(x)$	1	1	0	0
$y \text{ xor } f_{1,0}(x)$	1	0	0	1
$ x, y \text{ xor } f_{1,0}(x)\rangle$	$ 01\rangle$	$ 00\rangle$	$ 10\rangle$	$ 11\rangle$
	0	1	0	0
	1	0	0	0
	0	0	1	0
	0	0	0	1

Quantum Computers

Matrices:

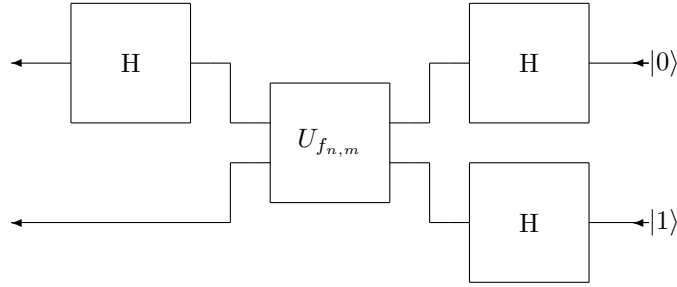
$$\begin{array}{c|cc}
 & m=0 & m=1 \\
 \hline
 n=0 & \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} & \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix} \\
 \hline
 n=1 & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} & \begin{bmatrix} 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}
 \end{array}$$

In general

$$U_{f_{n,m}} = \begin{array}{c} \begin{matrix} 00 & 01 & 10 & 11 \\ \begin{bmatrix} \bar{n} & n & 0 & 0 \\ n & \bar{n} & 0 & 0 \\ 0 & 0 & \bar{m} & m \\ 0 & 0 & m & \bar{m} \end{bmatrix} \end{matrix} \end{array}$$

The top row gives the $|xy\rangle$ input, the column is the matrix for the output. E.g., the output for input 11 is $[0 \ 0 \ m \ \bar{m}]^T$ (which, e.g., for $f_{1,1}$ is $[0 \ 0 \ 1 \ 0]^T = |10\rangle$)

3.5.4 Deutsch's circuit



The matrix representation for this is

$$(H \otimes I) * U_{f_{n,m}} * (H \otimes H) * |01\rangle$$

where H is the matrix for the Hadamard gate, I is a two-by-two identity matrix, $U_{f_{n,m}}$ is the matrix for our mystery controlled function gate (Note: n and m will have concrete values — we just don't know which ones), and $|01\rangle$ is our input. The reasoning behind these values is (from right to left):

- $|01\rangle$ — the input. This is given.
- $H \otimes H$ — the parallel composition of two Hadamard gates

Quantum Computers

- $U_{f_{n,m}}$ — the mystery gate
- $H \times I$ — the parallel composition of a Hadamard gate with an “identity gate” (i.e. the wire that just passes the lower qubit straight through).

The circuit is a sequential composition of these components, hence the use of ordinary matrix multiplication ($*$) to put them together.

We now need the matrix representations of these values:

- $|01\rangle$

This is the tensor product of $|0\rangle$ and $|1\rangle$.

$$|0\rangle \otimes |1\rangle = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \otimes \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix}$$

- $H \times H$

The matrix for a Hadamard gate is

$$\frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

so this parallel composition is

$$H \otimes H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix}$$

- $U_{f_{n,m}}$

This is given.

$$U_{f_{n,m}} = \begin{bmatrix} \bar{n} & n & 0 & 0 \\ n & \bar{n} & 0 & 0 \\ 0 & 0 & \bar{m} & m \\ 0 & 0 & m & \bar{m} \end{bmatrix}$$

- $H \otimes I$

The two-by-two identity matrix is

$$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

so our parallel composition is

$$H \otimes I = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \otimes \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & -1 & 0 \\ 0 & 1 & 0 & -1 \end{bmatrix}$$

Quantum Computers

So the whole circuit is

$$\begin{aligned}
(H \otimes I) * U_{f_{n,m}} * (H \otimes H) * |01\rangle &= (H \otimes I) * U_{f_{n,m}} * \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix} * \begin{bmatrix} 0 \\ 1 \\ 0 \\ 0 \end{bmatrix} \\
&= (H \otimes I) * U_{f_{n,m}} * \begin{bmatrix} +\frac{1}{2} \\ -\frac{1}{2} \\ +\frac{1}{2} \\ -\frac{1}{2} \end{bmatrix} \\
&= (H \otimes I) * \begin{bmatrix} \bar{n} & n & 0 & 0 \\ n & \bar{n} & 0 & 0 \\ 0 & 0 & \bar{m} & m \\ 0 & 0 & m & \bar{m} \end{bmatrix} * \begin{bmatrix} +\frac{1}{2} \\ -\frac{1}{2} \\ +\frac{1}{2} \\ -\frac{1}{2} \end{bmatrix} \\
&= (H \otimes I) * \begin{bmatrix} \frac{\bar{n}-n}{2} \\ \frac{n-\bar{n}}{2} \\ \frac{\bar{m}-m}{2} \\ \frac{m-\bar{m}}{2} \end{bmatrix} \\
&= \frac{1}{2} \begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & -1 & 1 & -1 \\ 1 & 1 & -1 & -1 \\ 1 & -1 & -1 & 1 \end{bmatrix} * \begin{bmatrix} \frac{\bar{n}-n}{2} \\ \frac{n-\bar{n}}{2} \\ \frac{\bar{m}-m}{2} \\ \frac{m-\bar{m}}{2} \end{bmatrix} \\
&= \frac{1}{\sqrt{2}} \begin{bmatrix} \frac{(\bar{n}+\bar{m})-(n+m)}{2} \\ \frac{(n+m)-(\bar{n}+\bar{m})}{2} \\ \frac{(\bar{n}+m)-(\bar{n}+\bar{m})}{2} \\ \frac{(n+\bar{m})-(\bar{n}+m)}{2} \end{bmatrix}
\end{aligned}$$

The values of the entries in the:

$$\frac{1}{\sqrt{2}} \begin{bmatrix} \frac{(\bar{n}+\bar{m})-(n+m)}{2} \\ \frac{(n+m)-(\bar{n}+\bar{m})}{2} \\ \frac{(\bar{n}+m)-(\bar{n}+\bar{m})}{2} \\ \frac{(n+\bar{m})-(\bar{n}+m)}{2} \end{bmatrix}$$

matrix, for the possible values of n and m are:

n	0	0	1	1
m	0	1	0	1
$\frac{(\bar{n}+\bar{m})-(n+m)}{2\sqrt{2}}$	$+\frac{1}{\sqrt{2}}$	0	0	$-\frac{1}{\sqrt{2}}$
$\frac{(n+m)-(\bar{n}+\bar{m})}{2\sqrt{2}}$	$-\frac{1}{\sqrt{2}}$	0	0	$+\frac{1}{\sqrt{2}}$
$\frac{(\bar{n}+m)-(\bar{n}+\bar{m})}{2\sqrt{2}}$	0	$+\frac{1}{\sqrt{2}}$	$-\frac{1}{\sqrt{2}}$	0
$\frac{(n+\bar{m})-(\bar{n}+m)}{2\sqrt{2}}$	0	$-\frac{1}{\sqrt{2}}$	$+\frac{1}{\sqrt{2}}$	0

Quantum Computers

E.g., if $n = 0$ and $m = 0$ the output of this circuit is

$$\begin{bmatrix} +\frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} & 0 & 0 \end{bmatrix}^T.$$

This is the superposition of two qubits, q_0 and q_1 , say, which can be written as

$$+\frac{|00\rangle - |01\rangle}{\sqrt{2}}.$$

Here q_0 is $|0\rangle$, and q_1 is $\frac{|0\rangle - |1\rangle}{\sqrt{2}}$.

For the other possible values of n and m :

	n	m		q_0	q_1
constant	0	0	$+\frac{ 00\rangle - 01\rangle}{\sqrt{2}}$	$+ 0\rangle$	$\frac{ 0-1\rangle}{\sqrt{2}}$
balanced	0	1	$+\frac{ 10\rangle - 11\rangle}{\sqrt{2}}$	$+ 1\rangle$	$\frac{ 0-1\rangle}{\sqrt{2}}$
balanced	1	0	$-\frac{ 10\rangle - 11\rangle}{\sqrt{2}}$	$- 1\rangle$	$\frac{ 0-1\rangle}{\sqrt{2}}$
constant	1	1	$-\frac{ 00\rangle - 01\rangle}{\sqrt{2}}$	$- 0\rangle$	$\frac{ 0-1\rangle}{\sqrt{2}}$

So measure q_0 .

- If $q_0 = |0\rangle$ function is **constant**.
- If $q_0 = |1\rangle$ function is **balanced**.

The sign also disappears on measurement.