

CSCI-665 Homework 3

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Problem 1

Problem 1

CODE 1

A B C D E F
~~35%~~ 15% 15% 8% 7%
 35% 20%

(i) $EKF \rightarrow \alpha_1$

A B C D α_1
~~35%~~ 20% 15% 15% 15%
 35%

(ii) $CKD \rightarrow \alpha_2$

A B α_1 α_2
~~35%~~ 20% 15% 30%

(iii) $BK\alpha_1 \rightarrow \alpha_3$

(iv) A α_2 α_3
~~35%~~ 30% 35%

AK $\alpha_2 \rightarrow \alpha_4$

α_3 α_4
 35% ~~65%~~
 65%

(v) $\alpha_3 K \alpha_4 \rightarrow \star$

A \rightarrow 11
 B \rightarrow 00
 C \rightarrow 100
 D \rightarrow 101
 E \rightarrow 010
 F \rightarrow 011

BPS = $.35 \times 2 + .2 \times 2 + .15 \times 3 + .15 \times 3 + .08 \times 3 + .07 \times 3$
 $= 2.45$

CSCI-665 Homework 3

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CODE 2

A	B	C	D	E	F
35	20	15	15	8	7

(i) $EKF \rightarrow \alpha_1$

(ii) $DK\alpha_1 \rightarrow \alpha_2$

A	B	C	D	α_1
35	20	15	15	15

(iii) $BKC \rightarrow \alpha_3$

A	α_2	α_3
35	30	35

(iv) $\alpha_2 \& \alpha_3 \rightarrow \alpha_4$

A	α_4
35	65

(v) $AK\alpha_4 \rightarrow \star$

$A \rightarrow 0$
 $B \rightarrow 110$
 $C \rightarrow 111$
 $D \rightarrow 100$
 $E \rightarrow 1010$
 $F \rightarrow 1011$

$BPS = .35 \times 1 + .20 \times 3 + .15 \times 3 + .15 \times 3$
 $+ .08 \times 4 + .07 \times 4$
 $= 2.45$

Problem 2

min. gaps
earliest finishing interval

same as provided in example

(1) (a) $(1, 4), (4, 5), (1, 4), (2, 7)$

(a.2)

(1) Earliest finish interval

(a) Intervals S & F :-
Same as provided in the question :-
 $(1, 4), (2, 7), (4, 5), (6, 10), (8, 9)$
 $S = 1, F = 10$

(b) If Earliest finish intervals are considered the sorting would look like :-

Intervals sorted by Finish time

Intervals selected
 $= \{(1, 4), (4, 5), (8, 9)\}$
 Gap $= |(10-9)| + |8-5| = 4$

(c) As we know the optimal intervals & corresponding gap are :-
 $(1, 4), (4, 5), (6, 10) \rightarrow \text{gap} = 2$

(2) Earliest start time

(a) Intervals & S & F :-
 $(1, 4), (2, 7), (4, 5), (6, 10)$
 $S = 1, F = 10$

(b) $\{(1, 4), (6, 10)\}$
 Gap $= 4$

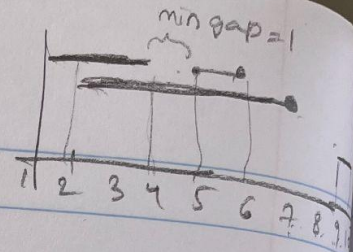
(c) optimal: $\{(2, 7), (8, 10)\}$
 Corresponding gap $= |2-1| + |8-7| = 2$

3.

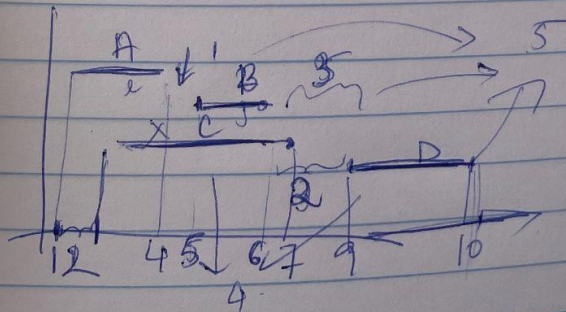
(a) $(1, 4), (2, 7), (5, 6), (9, 10)$
 $s = 1, f = 10.$

(b) $(1, 4), (5, 6), (9, 10)$
 $\text{Gap} = 4$

(c) optimum: $(2, 7), (9, 10)$
 $\text{Gap} = 3$



Reason:-



A & B have
 least gap.
 $\therefore C$ is discarded

$j \rightarrow B$
 $i \rightarrow A$

Problem 3

Given are n distinct points in 2-dimensional space. Design an $O(n^2 \log n)$ algorithm that determines whether any three of these points are evenly spaced, co-linear points. A $\Theta(n^3)$ algorithm can earn at most 8 out of 18 marks.

BONUS: only can be earned after you successfully complete this problem in $O(n^2 \log n)$ complexity. Additionally, design a solution that runs in expected $O(n^2)$ time. You may use a hash-table data structure for this part of the problem, and only for this part of the problem. Document the performance of this approach compared to your $O(n^2 \log n)$ solution.

Design an $O(n)$ algorithm that determines the number of youngest kids that end up sitting on the left side while balancing the boat as much as possible.

Hint: start out by asking yourself how you would solve this problem if the children were already sorted by increasing age. You can't afford to sort the data, however, as none of the $O(n)$ sorts applies (you can't make any assumptions about the distribution of the age values).

$O(n^2 \log n)$ solution :

The algorithm is as follows:

We use the concept of computing midpoints since it helps discover the colinearity and equidistance w.r.t. the other two points

Breaking down the steps:

We sort the 2-d array by y and then x values ($O(n \log n)$)

We take in two point combinations and compute mid points. ($O(n^2)$)

We check if these mid-points are present in the data array using binary search. The binary search on the 2 d array is done in such a way that it checks if x is already the mid-value. If yes, then it checks the y value (1st index) and decides to go to the left or right half. If x itself has not been found, depending on x value of point, we go to the left or right sub-array midpoints. ($O(\log n)$ for each midpoint)). The loop breaks if one mid-point is found.

This makes the algorithm $O(n^2(\log n))$ since midpoint is searched for all n^2 combinations

CSCI-665 Homework 3

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$O(n \log n)$:

Here a hashset is used to store points instead of binary search which easily finds if midpoint computed is present in the hashset. Adding and finding points in hashset is $O(1)$ steps.

This makes it a $O(n^2)$ algorithm.

The code is commented out in the uploaded data

For test cases

Binary search	vs	hashset
T1:		
0.0002572080120444298		0.00012516698916442692
T2:		
0.00023329199757426977		6.008398486301303e-05
T3:		
0.035936458996729925		0.003771957999560982
T4:		
0.0655184590141289		0.0016313329979311675
T5:		
14.235483291995479		2.750484917021822

As we can see there is a vast difference between the two approaches

Problem 4

We use median value to compute the number of kids and use a recursive function to shift kids from left and right buckets to get optimum minimum variance w.r.t. the balance weight ($\text{total}/2$) for the left and right buckets. We consider both teachers on both sides and then compute the least difference provided combination w.r.t. Balance weight for the left side weight.

Algorithm:

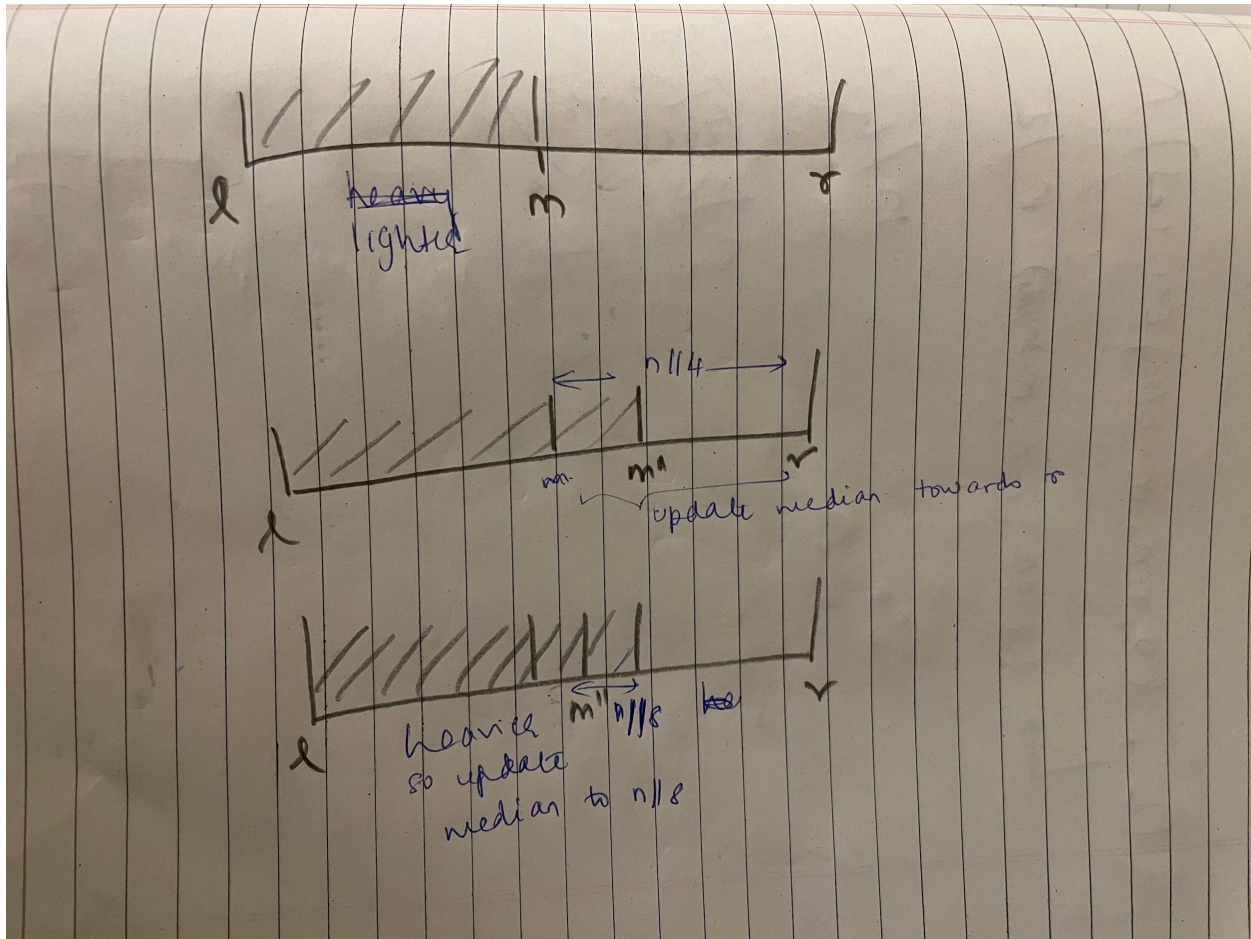
We use the k select algorithm for n datapoints(kids). The expected linear time for this algorithm is $O(n)$. If we use k select recursively then we run it for $n/2, n/4, \dots$ which leads to geometric progression that leads to 2 as the constant for n ($1/1-\alpha$ formula). This confirms that this

CSCI-665 Homework 3

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doesn't lead to a more complex algorithm is greater than $O(n)$.



Problem 5

Suppose that instead of encoding with a prefix code such as determined in Problem 1, the symbols are encoded using the codewords given in the table below:

Symbol Codeword

A 0

B 1

C 10

CSCI-665 Homework 3

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D 01

E 111

F 011

Given an input encoding of size n bits, design an $O(n)$ algorithm that determines how many different decodings of those n bits are possible using the table above.

Algorithm:

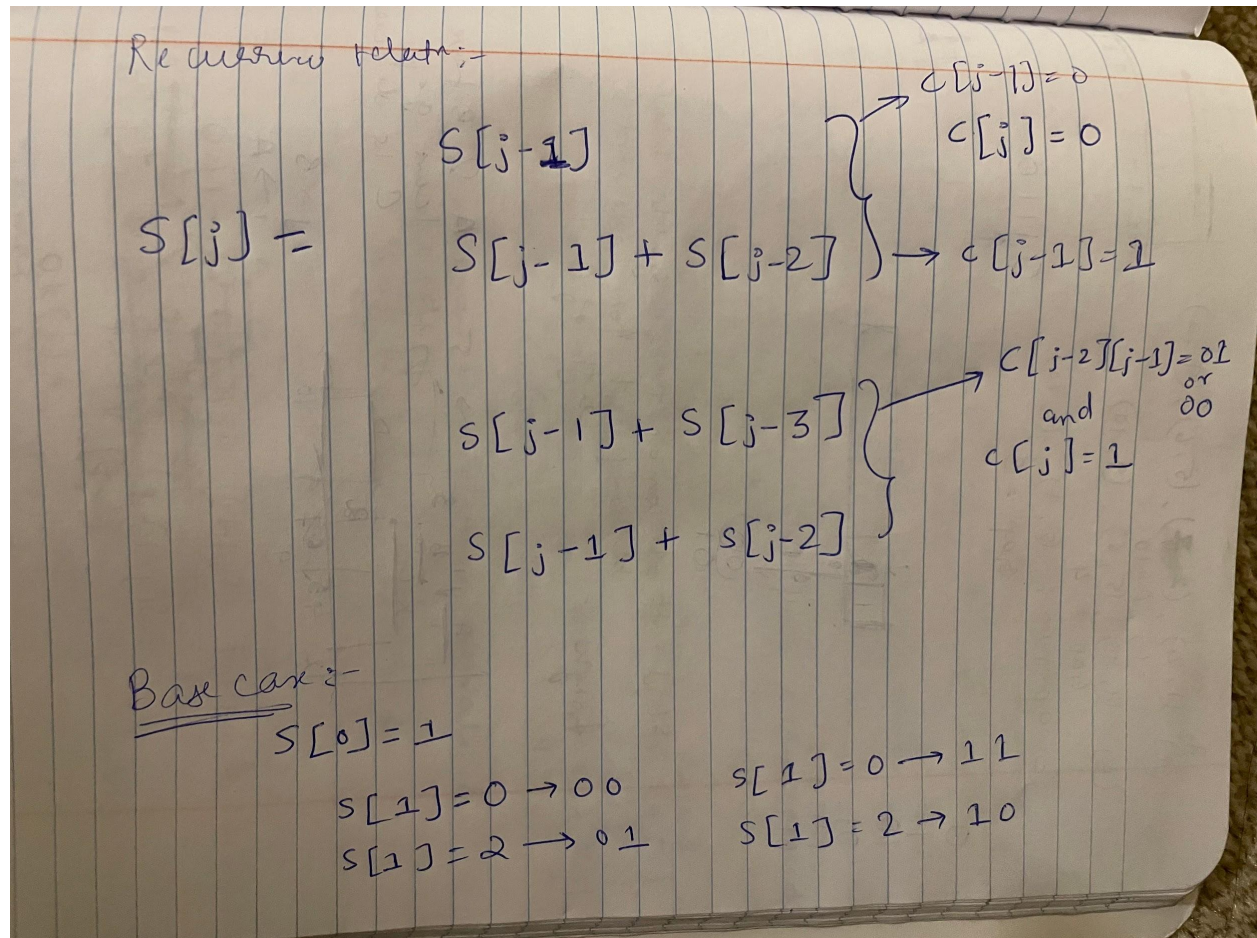
It goes through the bitstream linearly and goes back using index keeping constant time

Heart of algorithm:

Problem description:

$S[j]$ = number of combinations of symbols that are computed (decoded) for bitstream of length from $0 \dots j$

Recurrence equation:



Final output:

$S[n]$ = provides the combinations of codes used in bitstream of size n