TWINX

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1.* Introduction.   This short program compiles a master index for a set of programs that have been
processed by CTWILL. To use it, you say, e.g., twinx *.tex >index.tex. The individual programs should
define their names with a line of the form '\def\title{NAME}'.

#**include** <stdio.h>
#**include** <stdlib.h>
#**include** <string.h>
#**include** <ctype.h>
  ⟨ Type definitions 4* ⟩
  ⟨ Global variables 2 ⟩
  ⟨ Procedures 5* ⟩
  int *main*(**int** *argc*, **char** *∗argv*[ ])
  {
      ⟨ Local variables 9 ⟩;
      ⟨ Initialize the data structures 8 ⟩;
      while (−−*argc*) {
        if ((*f* ← *fopen*(*∗++argv*, "r")) ≡ Λ)
            fprintf(*stderr*, "twinx:␣Couldn't␣open␣file␣%s␣for␣reading!\n", *∗argv*);
        else {
            ⟨ Scan file *f* until coming to the title 3* ⟩;
            fclose(*f*); *strncpy*(*∗argv* + *strlen*(*∗argv*) − 3, "idx", 3);
            if ((*f* ← *fopen*(*∗argv*, "r")) ≡ Λ)
                fprintf(*stderr*, "twinx:␣Couldn't␣open␣file␣%s␣for␣reading!\n", *∗argv*);
            else {
                ⟨ Copy the index file *f* into the data structures 10 ⟩;
                fclose(*f*);
            }
        }
      }
      ⟨ Output the data structures to make a master index 13 ⟩;
      return 0;
  }

3.*   For your convenience, TWINX grabs the first "word" in \title and turns it into uppercase form.

⟨ Scan file *f* until coming to the title 3* ⟩ ≡
  while (1) {
      if (*fgets*(*buf*, *buf_size*, *f*) ≡ Λ) {
          fprintf(*stderr*, "twinx:␣(no␣title␣found␣in␣file␣%s)\n", *∗argv*); *title*[0] ← '\0'; **break**;
      }
      if (*strncmp*(*buf*, "\\def\\title\{", 11) ≡ 0) { **register char** *∗p*, *∗q*;
          for (*p* ← *buf* + 11, *q* ← *title*; *∗p* ∧ *∗p* ≠ '␣' ∧ *∗p* ≠ '}'; *p*++) *∗q*++ ← *toupper*(*∗p*);
          ∗q ← '\0'; **break**;
      }
  }
This code is used in section 1*.

4*. Data structures.   Our main task is to collate a bunch of texts associated with keys that have already
been sorted. It seems easiest to do this by repeatedly merging the new data into the old, even though this
means we'll be passing over some of the same keys 30 times or more; the computer is fast, and this program
won't be run often.

Further examination shows that a merging strategy isn't so easy after all, because the sorting done by
`CTWILL` (and by `CWEAVE`) is weird in certain cases. When two index entries agree except for their "ilk," the
order in which they appear in the index depends on the order in which they appear in the program. Thus,
they might well appear in different order in two of the indexes we are merging. (There's also another glitch,
although not quite as devasting: When two index entries have the same letters and the same ilk, but differ
with respect to uppercase versus lowercase, the order in which they appear depends on the hash code used
in `CWEB`'s `common.w` code!)

So we'll use Plan B: All index entries will first be copied into a long list. The list will almost always consist
of many sorted sublists, but we will not assume anything about its order. After all the copying has been
done, we will use a list-merge sort to finish the job.

The data structure is built from nodes that each contain three pointers. The first pointer is to an *id* string;
the third pointer is to the *next* node; and the second pointer is either *data.s*, a pointer to a string of text,
or *data.n*, a pointer to a node. In the main list, the *id* fields are the keys of the index, and the *data.n* fields
point to lists of associated texts. In the latter lists, the *id* fields are the individual program titles, while the
data.s fields are the texts.

⟨ Type definitions 4* ⟩ ≡
  typedef union {
    char ∗*s*;
    struct node_struct ∗*n*;
  } **mixed**;
  typedef struct node_struct {
    const char ∗*id*;
    mixed *data*;
    struct node_struct ∗*next*;
  } **node**;
This code is used in section 1*.

5* We copy strings into blocks of storage that are allocated as needed. Here's a routine that stashes away a given string. It makes no attempt to handle extremely long strings, because such strings will arise only if the input is all screwed up.

#**define** *string_block_size* 8192   ▷ number of bytes per string block ◁

⟨ Procedures 5* ⟩ ≡

  char *save_string(**char** *s)

  {

    register char *p, *q;

    register int *l*;

    for $(p \leftarrow s; *p; p{+}{+})$;

    $l \leftarrow p - s + 1$;

    if $(l > string\_block\_size)$ {

      fprintf(*stderr*, "twinx:␣Huge␣string␣'%.20s...'␣will␣be␣truncated!\n", *s*);

      $l \leftarrow string\_block\_size$; $s[l-1] \leftarrow$ '\0';

    }

    if $(next\_string + l \geq bad\_string)$ {

      $next\_string \leftarrow$ (**char** *) *malloc*(*string_block_size*);

      if $(next\_string \equiv \Lambda)$ {

        fprintf(*stderr*, "twinx:␣Not␣enough␣room␣for␣strings!\n"); $exit(-1)$;

      }

      $bad\_string \leftarrow next\_string + string\_block\_size$;

    }

    for $(p \leftarrow s, q \leftarrow next\_string; *p; p{+}{+})$ $*q{+}{+} \leftarrow *p$;

    $*q \leftarrow$ '\0'; $next\_string \leftarrow q + 1$; **return** $next\_string - l$;

  }

See also sections 6*, 17*, and 20*.

This code is used in section 1*.

6* Nodes are allocated with a similar but simpler mechanism.

#**define** *nodes_per_block* 340

⟨ Procedures 5* ⟩ +≡

  node *new_node(**void**)

  {

    if $(next\_node \equiv bad\_node)$ {

      $next\_node \leftarrow$ (**node** *) *calloc*(*nodes_per_block*, **sizeof**(**node**));

      if $(next\_node \equiv \Lambda)$ {

        fprintf(*stderr*, "twinx:␣Not␣enough␣room␣for␣nodes!\n"); $exit(-2)$;

      }

      $bad\_node \leftarrow next\_node + nodes\_per\_block$;

    }

    $next\_node{+}{+}$; **return** $next\_node - 1$;

  }

11* ⟨Copy a new index entry into *cur_name* and *cur_node* 11*⟩ ≡

  if (*buf*[4] ≠ '{') {

    fprintf(*stderr*, "twinx:␣missing␣brace␣in␣file␣%s:␣'%.20s...'\n", ∗*argv*, *buf*); **break**;

  }

  { **register char** ∗*p*, ∗*q*; **register int** *bal* ← 1;

    cur_name[0] ← *buf*[2]; *cur_name*[1] ← *buf*[3]; *cur_name*[2] ← '{';

    for (*p* ← *buf* + 5, *q* ← *cur_name* + 3; ∗*p* ∧ (*bal* ∨ ∗*p* ≡ '{'); *p*++) {

      switch (∗*p*) {

      case '\\': ∗*q*++ ← ∗*p*++; **break**;

      case '{': *bal*++; **break**;

      case '}': *bal*−−; **break**;

      }

      ∗*q*++ ← ∗*p*;

    }

    if (*bal*) {

      fprintf(*stderr*, "twinx:␣unbalanced␣entry␣in␣file␣%s:␣'%.20s...'\n", ∗*argv*, *buf*); **break**;

    }

    if (∗*p*++ ≠ ',') {

      fprintf(*stderr*, "twinx:␣missing␣comma␣in␣file␣%s:␣'%.20s...'\n", ∗*argv*, *buf*); **break**;

    }

    if (∗*p*++ ≠ '␣') {

      fprintf(*stderr*, "twinx:␣missing␣space␣in␣file␣%s:␣'%.20s...'\n", ∗*argv*, *buf*); **break**;

    }

    ∗*q* ← '\0'; ⟨Copy the text part of the index entry into *cur_node* 12⟩;

  }

This code is used in section 10.

17* Comparison is a three-stage process in general. First we compare the keys without regarding case or format type. If they are equal with respect to that criterion, we try again, with case significant. If they are still equal, we look at the format characters (the first two characters of the *id* field).

⟨Procedures 5*⟩ +≡

  int *compare*(**node** ∗*p*, **node** ∗*q*)

  { **register unsigned char** ∗*pp*, ∗*qq*;

    for (*pp* ← (**unsigned char** ∗) *p*⃗*id* + 3, *qq* ← (**unsigned char** ∗) *q*⃗*id* + 3; ∗*pp* ∧ *ord*[∗*pp*] ≡ *ord*[∗*qq*];

      pp++, *qq*++) ;

    if (∗*pp* ∨ ∗*qq*) **return** *ord*[∗*pp*] − *ord*[∗*qq*];

    for (*pp* ← (**unsigned char** ∗) *p*⃗*id* + 3, *qq* ← (**unsigned char** ∗) *q*⃗*id* + 3; ∗*pp* ∧ ∗*pp* ≡ ∗*qq*;

      pp++, *qq*++) ;

    if (∗*pp* ∨ ∗*qq*) **return** (**int**) ∗*pp* − (**int**) ∗*qq*;

    if (*p*⃗*id*[0] ≠ *q*⃗*id*[0]) **return** *p*⃗*id*[0] − *q*⃗*id*[0];

    return *p*⃗*id*[1] − *q*⃗*id*[1];

  }

19* The right brace is placed lowest in collating order, because each key is actually followed by a right brace when we are sorting.

Apology: I haven't had time to update this part of the program to allow 8-bit characters. At present the data is assumed to be 7-bit ASCII, as it was in the early versions of CWEAVE.

⟨Initialize the data structures 8⟩ +≡
  $collate[0] \leftarrow 0$; $strcpy(collate + 1,$
      "}␣\1\2\3\4\5\6\7\10\11\12\13\14\15\16\17\20\21\22\23\24\25\26\27\30\31\32\33\34\
      \35\36\37!\42#$%&'()*+,-./:;<=>?@[\\]^'{|~_abcdefghijklmnopqrstuvwxyz0123456789");
  { **register int** j;

     for $(j \leftarrow 1;\ collate[j];\ j{+}{+})\ ord[(\textbf{int})\ collate[j]] \leftarrow j$;
     $ord[128] \leftarrow j$;   ▷ this affects the ordering of $sentinel.id$ ◁
     for $(j \leftarrow \text{'A'};\ j \leq \text{'Z'};\ j{+}{+})\ ord[j] \leftarrow ord[tolower(j)]$;
  }

20* When two lists are combined, we put the data from the second node before the data from the first node, because we are going to reverse the order when printing. After this procedure has acted, the field $q{\rightarrow}data.n$ should not be considered an active pointer.

⟨Procedures 5*⟩ +≡
  void $collapse(\textbf{node} *p, \textbf{node} *q)$
  { **register node** $*x$;

     for $(x \leftarrow q{\rightarrow}data.n;\ x{\rightarrow}next;\ x \leftarrow x{\rightarrow}next)$;
     $x{\rightarrow}next \leftarrow p{\rightarrow}data.n$; $p{\rightarrow}data.n \leftarrow q{\rightarrow}data.n$;
  }

22*   ⟨Output $x{\rightarrow}id$ in suitable TEX format 22*⟩ ≡
  { **register const char** $*p \leftarrow x{\rightarrow}id$;

     if $(*p \equiv \text{'}\text{␣}\text{'})$ {
        if $(*(p + 1) \neq \text{'}\text{␣}\text{'})$ **goto** $unknown$;
        goto $known$;
     }
     if $(*p \neq \text{'}\backslash\backslash\text{'})$ **goto** $unknown$;
     switch $(*(p + 1))$ {
     case '\\': **case** '|': **case** '.': **case** '&': **case** '9': $printf(\texttt{"\\\\\%c"}, *(p + 1))$; **goto** $known$;
     case '$': $printf(\texttt{"\$\\\\"})$;
        for $(p\ {+}{=}\ 3;\ *p \neq \text{'}\}\text{'};\ p{+}{+})$
           if $(*p \equiv \text{'}\_\text{'})\ putchar(\text{'x'})$;
           else $putchar(*p)$;
        $putchar(\text{'\$'})$; **goto** $done$;
     default: **goto** $unknown$;
     }
  $unknown$: $fprintf(stderr, \texttt{"twinx:␣'\%s'␣has␣unknown␣format!\textbackslash n"}, p)$;
  $known$:
     for $(p\ {+}{=}\ 2;\ *p;\ p{+}{+})$ {
        if $(*p \equiv \text{'}\_\text{'} \wedge *(p - 1) \neq \text{'}\backslash\backslash\text{'})\ putchar(\text{'}\backslash\backslash\text{'})$;
        $putchar(*p)$;
     }
  $done$: ;
  }
This code is used in section 21.

24*. Index.

The following sections were changed by the change file: 1, 3, 4, 5, 6, 11, 17, 19, 20, 22, 24.