Chapter 8

Tree Algorithms - Keyword Query



Acknowledgements

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Chapter Outline

- Tree data
- Tree encoding
- Keyword query semantics
- Keyword query processing
- Distributed algorithms

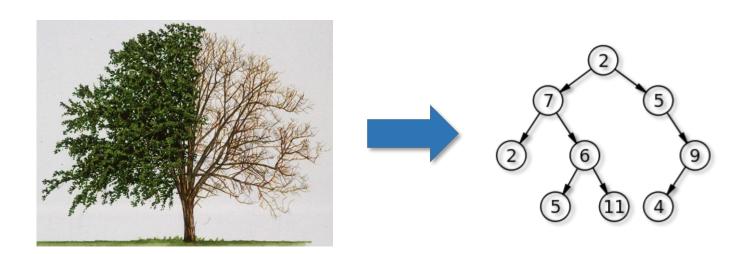
Tree Data





Tree (Data Structure)

- In computer science, a tree is a widely used abstract data type (ADT)—or data structure implementing this ADT
- which simulates a hierarchical tree structure, with a root value and subtrees of children with a parent node, represented as a set of linked nodes





Tree Applications

- Data types
 - XML
 - HTML
 - JSON
- Applications
 - Catalogue
 - Operating System paths
 - Organizations
 - Playoffs (Reverse)















Terminology

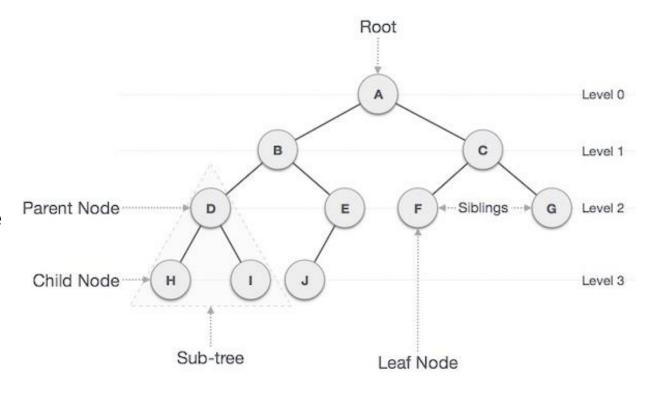
- Root: The top node in a tree.
- Child: A node directly connected to another node when moving away from the Root.
- Parent: The converse notion of a child.
- **Siblings**: A group of nodes with the same parent.
- **Descendant**: A node reachable by repeated proceeding from parent to child.
- Ancestor: A node reachable by repeated proceeding from child to parent.
- Leaf: (less commonly called External node) A node with no children.
- Branch / Internal node: A node with at least one child.
- Degree: The number of subtrees of a node.
- Edge: The connection between one node and another.
- Path: A sequence of nodes and edges connecting a node with a descendant.
- Level: Defined by 1 + (the number of connections between the node and the root).
- Height of node: The number of edges on the longest path between that node and a leaf.
- Height of tree: The height of a tree is the height of its root node.
- **Depth**: The depth of a node is the number of edges from the tree's root node to the node.
- Forest: A forest is a set of $n \ge 0$ disjoint trees.





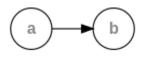
Terminology

- Root
- Child
- **Parent**
- Siblings
- **Descendant**
- **Ancestor**
- Leaf
- **Branch / Internal node**
- **Degree**
- Edge
- **Path**
- Level
- **Height of node**
- **Height of tree**
- **Depth**
- **Forest**

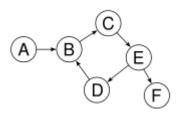


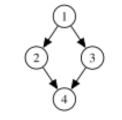


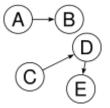
Situations











- Each linear list is trivially a tree
- NOT a tree: cycle A→A. A is the root but it also has a parent
- NOT a tree: cycle $B \rightarrow C \rightarrow E \rightarrow D \rightarrow B$. B has more than one parent (inbound edge)
- NOT a tree: undirected cycle 1-2-4-3. 4 has more than one parent (inbound edge).
- NOT a tree: two non-connected parts,
 A→B and C→D→E. There is more than one root.



Node Represenation

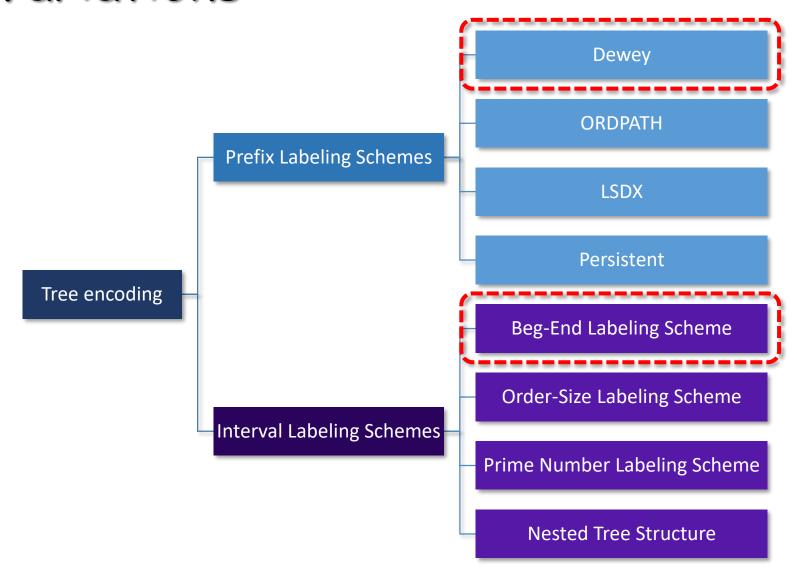
• E.g., Binary tree

```
struct node {
  int data;
  struct node *leftChild;
  struct node *rightChild;
};
```

Tree Encoding



Variations



Tree Encoding

Interval Labeling Schemes Prefix Labeling Schemes

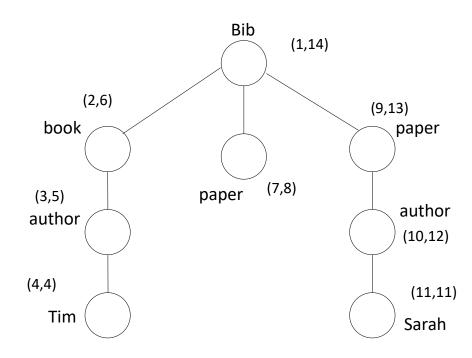


Beg-End Labeling Scheme

- A pair of numbers is assigned to each node in an XML document according to its sequential traversal order.
 - Starting from the root element, each node is given a "Beg" number.
 - If the end of an attribute, an attribute value, or an ending tag element is reached, the "End" number is assigned. The "End" number is equal to the next sequential number.
 - If the value of the element is a leaf the "Beg" number = "End" number



Example



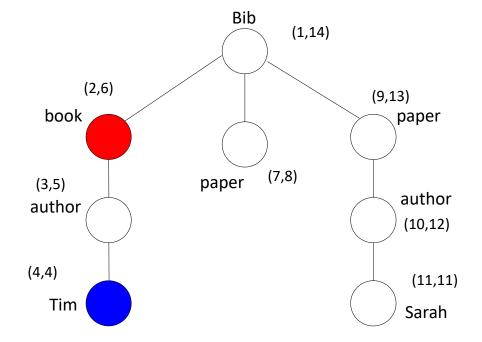


Properties 1

• A "Level" is added to the (Beg, End) label to form a nodetriplet identification label (Beg, End, Level) for each node in the tree.

Ancestor-descendant relationship:

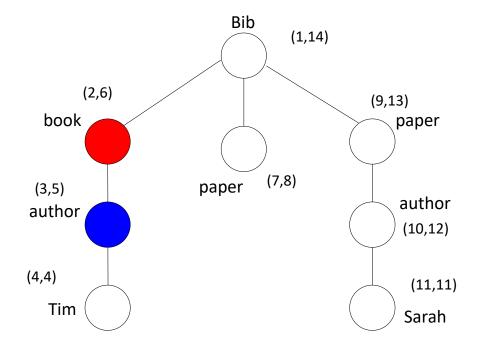
 In a given data-tree, node "x" is an ancestor of node "y" iff $(x.Beg < y.Beg \le y.End < x.End)$ (preorder property).





Properties 2

- Parent-child relationship:
 - In a given data-tree, node "x" is a parent of node "y" iff
 (x.Beg < y.Beg ≤ y.End < x.End) and (y.Level = x.Level + 1).
- There is **no way** to locate the siblings of a given node, using only the knowledge of its index numbers.



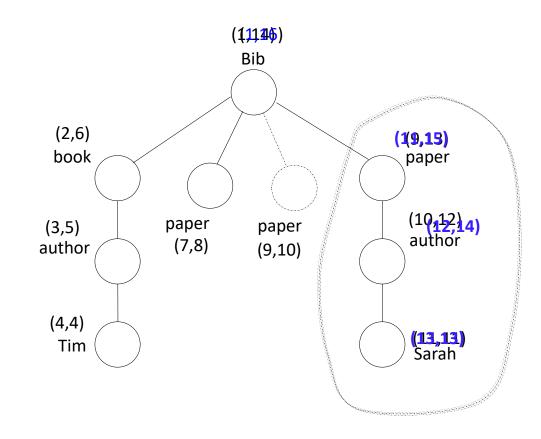


Are updates possible?

- Updating the labeling (numbering) scheme of Beg-End is costly.
- When a new node is inserted into the tree, then all the nodes in the tree, except the left sibling subtrees of the inserted node, have to be updated.
- On the other hand when a node is deleted no relabeling is needed.



Update Example



Tree Encoding

Interval Labeling Schemes Prefix Labeling Schemes



Dewey Encoding

- In a prefix labeling scheme, the label of a node in the XML tree often consists of:
 - A **prefix**, which often represents the label of all the ancestors of the node.
 - A delimiter, which in most cases is the fullstop "."
 - A positional identifier, which indicates the position of the node relative to its siblings.



Comparison

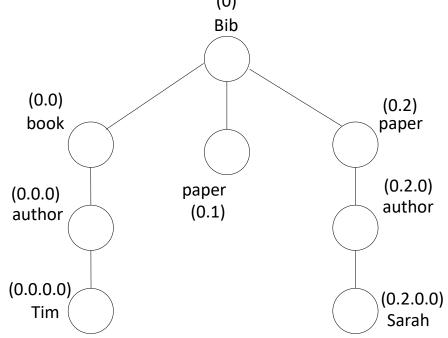
- Prefix Labeling Schemes:
 - Can handle updates easier and more efficient than Interval Labeling Schemes
 - Support sibling relationship
- However:
 - Extra space required to store paths
 - Its storage size increases quickly as the depth and the breath of the tree increases
 - Infer a bit more costly ancestor/descendant relationship



Dewey - Structure

- Each node is assigned a label that represents the path from the document's root to the node.
- Each component of the label represents the local order of an ancestor node.

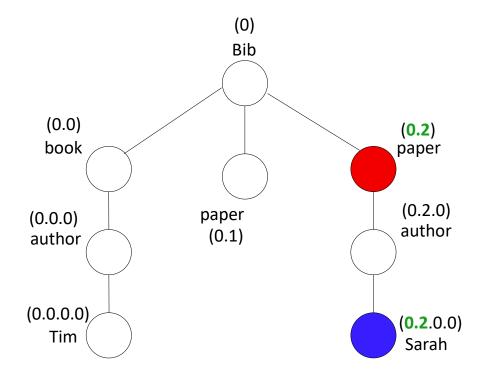
• Nodes with the same number of delimiters (".") in their label are in the same level. (0)





Dewey - Supported Queries

- Ancestors / Descendants
 - Node "X" is an ancestor of node "Y" if the label of node "X" is a substring of the label of node "Y".

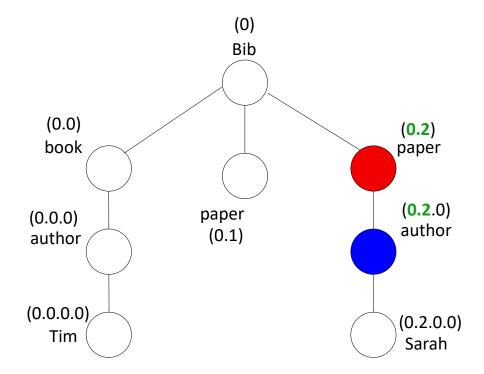




Dewey - Supported Queries

Parent / Child

- Node "X" is parent of node "Y" if:
 - The label of node "X" is a substring of the label of node "Y" and
 - frags(X) = frags(Y) 1, where frags(X) is the number of delimiters of the label of node X and frags(Y) is the number of delimiters of label of node Y.

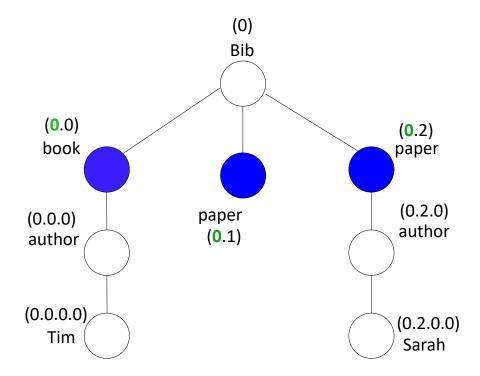




Dewey - Supported Queries

Siblings

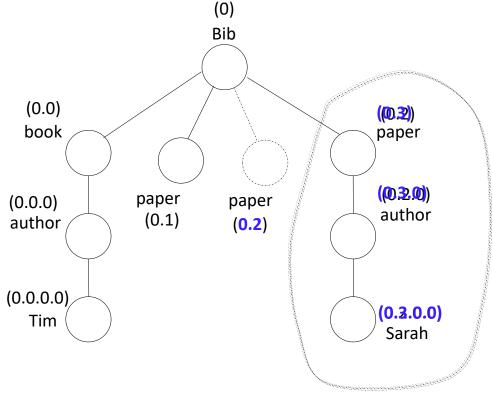
- Nodes "X" and "Y" are siblings if:
 - They have the same number of delimiters in their labels and
 - X.prefix = Y.prefix, where prefix is the label of the node without its positional identifier





Dewey - Updates

- Insertion of new node
 - The label of the nodes in the subtree rooted at the following sibling need to be updated
 - O(n) nodes need relabeling, where n is the number of nodes of the XML file





Dewey: Conclusion

- Not efficient for dynamic XML files with many updates
 - Need to re-label many nodes
- As the depth of the tree increases:
 - Label size of a node increases rapidly
 - Storage size increases rapidly
 - It becomes more costly to infer the supported queries between any two nodes (the string prefix matching becomes longer)
- Overflow problem
 - The original fixed length of bits assigned to store the size of the label is not enough.

Keyword Query

Semantics Query Processing



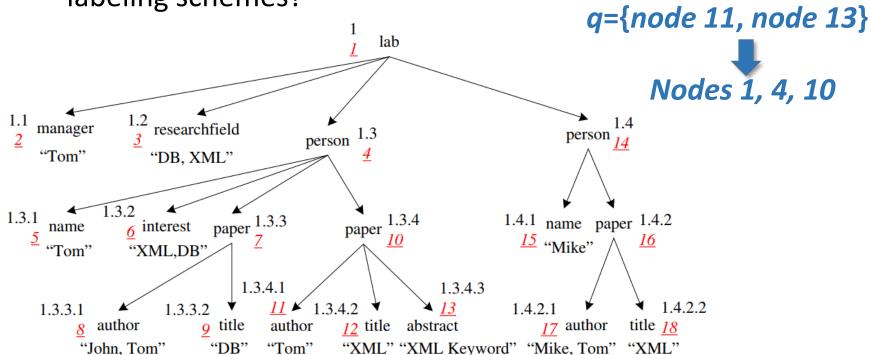
Why do we need keyword search?

- Tree data (i.e., XML, JSON, etc) is becoming a standard in data representation
- Enables users to access information in tree databases
- Inspired by IR (Information Retrieval) style keyword search on the web, that is designed mostly for text databases
- Desirable to support keyword search in tree database without the knowledge of complex query languages
- The extreme success of web search engines makes keyword search the most popular search model for ordinary users



Common Ancestors

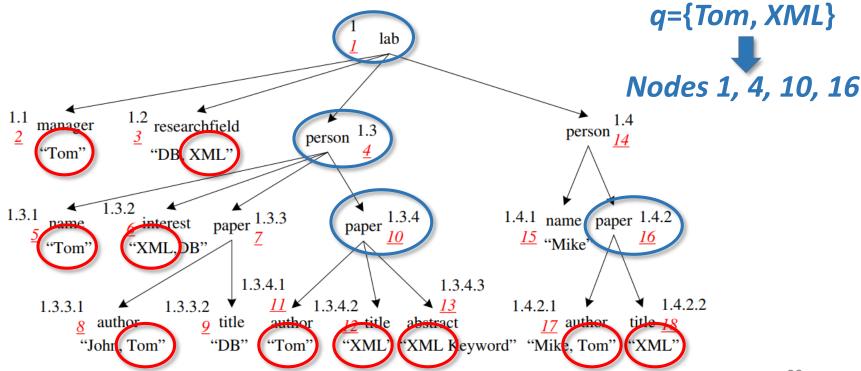
- Given two nodes in a tree
 - What are the common ancestors of these two nodes?
 - How to find them?
 - What's the difference between using Dewey and Interval labeling schemes?





Semantic: LCA

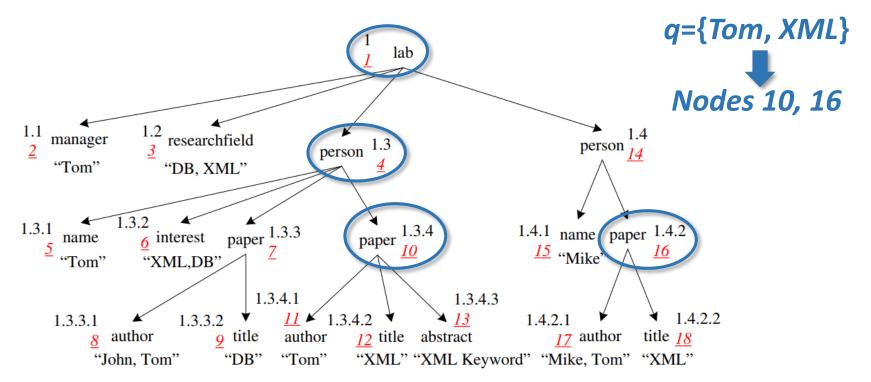
- Lowest Common Ancestor
 - Given an XML document d and a subset of nodes v_1 , v_2 , ..., v_m from d matching a set of query terms t_1 , t_2 , ..., t_n , the LCA of those nodes is a node e that is their common ancestor located farthest from the root of d.





Semantic: SLCA

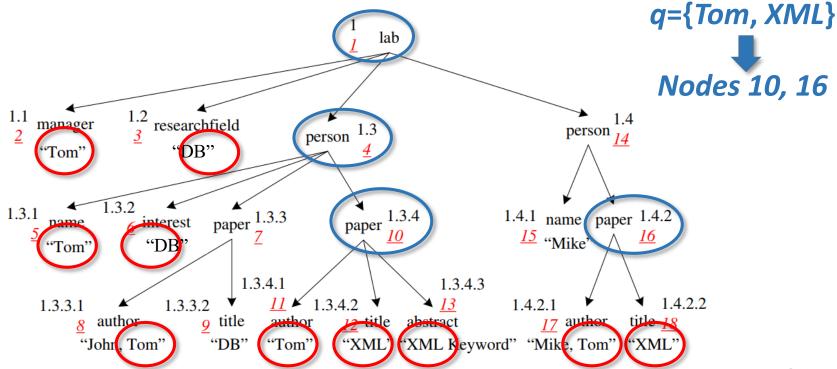
- Smallest Lowest Common Ancestor:
 - Given a set of LCA nodes returned as the result of a query q on an XML document d, the corresponding SLCA nodes are the LCA nodes that contain no other LCA node as descendant.





Semantic: ELCA

- Exclusive Lowest Common Ancestor:
 - Given a set of LCA nodes returned as the result of a query q on an XML document d, the corresponding ELCA nodes are those LCA nodes that contain at least one occurrence of each query term, excluding all their descendent LCA subtrees that also match any of those terms.





Other LCA-based Semaantics

Semantics	Definition	Existing algorithms	Returned nodes in Example 1
LCA	An LCA is a lowest common ancestor of a combination of matching nodes, i.e., each keyword corresponds to at least one matching node in the combination	XRANK Sigmod 2003	{ &o1, &o2, &o3, &o4, &o5 }
ELCA (Exclusive LCA)	*An ELCA is an LCA of a combination of matching nodes *An ELCA has its own witnesses, i.e., it does not share its matching nodes with its descendant ELCA nodes	*Index Stack EDBT 2008 *Hash Count EDBT 2010 *Top-K ICDE 2010 *Set-intersection ICDE 2011	{ &o1, &o4, &o5 }
SLCA (Smallest LCA)	*An SLCA is an LCA of a combination of matching nodes *There is no LCA node as its descendant	*XKSearch Sigmod 2005 *Multiway-SLCA WWW 2007 *Top-K ICDE 2010 *Set-intersection ICDE 2011	{ &o4, &o5 }
VLCA (Valuable LCA)	*A VLCA is an LCA of a combination of matching nodes *For each pair of matching nodes, all nodes in the path connecting them are of different types.	*XSEarch VLDB 2003 *VLCA CIKM 2007 *RLCA ADC 2010	{ &o2, &o3, &o4, &o5 }
MLCA (Meaningful LCA)	*An MLCA is an LCA of a combination of matching nodes *The LCA of matching nodes in the data tree belongs to the LCA of their node types in the schema tree	*MLCA VLDB 2004	{ &o1, &o2, &o3, &o4, &o5 }

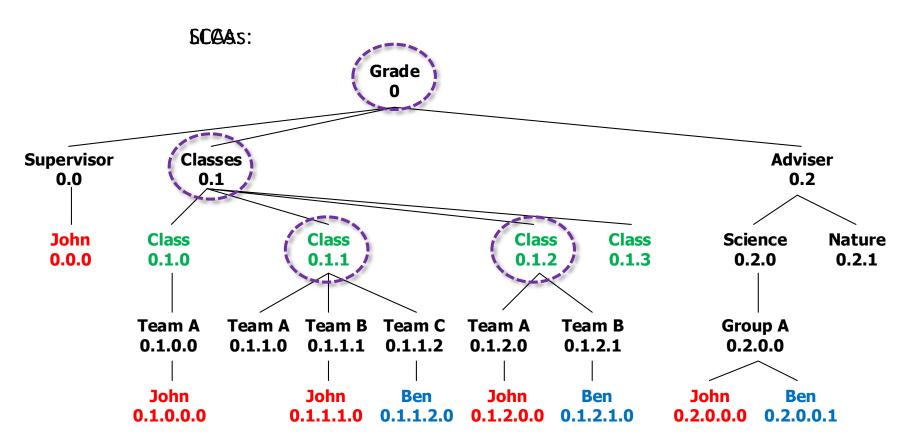
Keyword Query

Semantics Query Processing



Find SLCAs

 $Q = \{John, Ben, Class\}$





Stack Algorithm for SLCA

 $Q = \{John, Ben, Class\}$

 \mathbf{C} J В

0	T	F	F
0	F	F	F
0	F	F	F

(a) node 0.0.0

0	F	F	T
1	F	F	F
0	T	F	F

(b) node 0.1.0

0	T	F	F
0	F	F	F
0	F	F	T
1	F	F	F
0	T	F	F

(c) node 0.1.0.0.0

	(0.0.0,)
	0.1.0.0.0,
John: <	(0.1.1.1.0,)
	0.1.2.0.0,
	(0.2.0.0.0)

1	F	F	T
1	T	F	T
0	T	F	T

(d) node 0.1.1

0	T	F	F
1	F	F	F
1	F	F	T
1	T	F	T
0	T	F	T

(e) node 0.1.1.1.0

2	F	F	T
1	F	F	F
0	F	F	F

(g) node 0.1.2 report 0.1.1 as a SLCA

0	F	T	F
2	F	F	F
1	Т	F	Т
1	T	F	T
0	T	F	T

(f) node 0.1.1.2.0

Ben: -	$\left\{ egin{matrix} 0.1.1.2.0, \ 0.1.2.1.0, \ 0.2.0.0.1 \end{matrix} ight\}$

Class: {

Distributed Algorithms

Reference UNPUBLISHED

End of Chapter 8