

# LAB ASSIGNMENT 1

## SORTING

1 Given an array nums with n objects colored red, white, or blue, sort them in-place so that objects of the same color are adjacent, with the colors in the order red, white, and blue. We will use the integers: 0 to represent red 1 to represent white 2 to represent blue

### Approach 1 Using Bubble Sort

**How it Works :** Bubble Sort repeatedly compares adjacent elements and swaps them if they are in the wrong order. For an array of 0s, 1s, and 2s, it keeps swapping 2s to the right and 0s to the left in each pass. After multiple passes, all 0s come first, then 1s, and 2s move to the end, sorting the array.

```
#include <bits/stdc++.h>
using namespace std;

void bubbleSort(vector<int>&arr){
    int n = arr.size();
    for(int i=0; i<n-1; i++){
        for(int j=0; j<n-i-1; j++){
            if(arr[j] > arr[j+1]){
                swap(arr[j], arr[j+1]);
            }
        }
    }
}
```

### Approach 2 Selection Sort

**How it Works :** Selection Sort repeatedly finds the **minimum element** from the unsorted part and places it at the beginning. For an array of 0s, 1s, and 2s, it picks the smallest (0 first, then 1) and moves it to the correct position in each pass. After all passes, the array is sorted with 0s first, then 1s, and 2s at the end.

```
void selectionSort(vector<int>&arr){
    int n = arr.size();
    for(int i=0; i<n-1; i++){
        int minIndex = i;
        for(int j=i+1; j<n; j++){
            if(arr[j] < arr[minIndex]) minIndex = j;
        }
        swap(arr[i], arr[minIndex]);
    }
}
```

### Approach 3 Insertion Sort

**How it Works:**

Insertion Sort builds the sorted array one element at a time by taking each element and placing it in its correct position among the already sorted elements. For an array of 0s, 1s, and 2s, it moves each 0 or 1 to the left of larger elements in each pass.

After all passes, the array is sorted with 0s first, then 1s, and 2s at the end.

```
void insertionSort(vector<int>&arr){
    int n = arr.size();
    for(int i=1; i<n; i++){
        int min = arr[i];
        int prev = i - 1;
        while(prev >= 0 && arr[prev] > min){
            arr[prev + 1] = arr[prev];
            prev--;
        }
        arr[prev + 1] = min;
    }
}
```

### Output Using any od Bubble Sort Selection Sort And Insertion Sort

```
Enter the number of objects: 3
Enter the objects color (Red, White, Blue) in any case combination:
Enter color of object 1 out of 3 : Red
Enter color of object 2 out of 3 : White
Enter color of object 3 out of 3 : Blue
Your sorted array of object is :
Red White Blue
```

### Approach 4 Using Merge Sort

**How it Works :** Merge Sort is a **divide and conquer** algorithm that splits the array into two halves, recursively sorts each half, and then merges them into a sorted array. For an array of 0s, 1s, and 2s, it keeps dividing and merging, ensuring that 0s come first, then 1s, and 2s at the end.

### Better Complexity:

- Merge Sort has a time complexity of  $O(n \log n)$ , which is much faster than Bubble, Selection, and Insertion Sort ( $O(n^2)$ ) for larger arrays.
- Space complexity is  $O(n)$  due to temporary arrays used during merging.

```
vector<int>mergeTwoSortedArray(vector<int>&arr1,vector<int>&arr2){
    int i = 0;
    int j = 0;
    int k = 0;
    int size1 = arr1.size();
    int size2 = arr2.size();
    vector<int>arr3(size1+size2);
    while(i<size1 && j<size2){
        if(arr1[i] < arr2[j]) arr3[k++] = arr1[i++];
        else arr3[k++] = arr2[j++];
    }
    while(i<size1) arr3[k++] = arr1[i++];
```

```

        while(j<size2) arr3[k++] = arr2[j++];
        return arr3;
    }

vector<int> mergeSort(vector<int>&arr,int start,int end){
    int mid = (start+end)/2;
    if(start<=end) return {arr[start]};
    vector<int>arr1 = mergeSort(arr,start,mid);
    vector<int>arr2 = mergeSort(arr,mid+1,end);
    return mergeTwoSortedArray(arr1,arr2);
}

```

### Approach 5 : Counting Sort

**How it Works :** Count the number of 0s, 1s, and 2s in the array first. Then, overwrite the array by putting all 0s first, followed by 1s, and then 2s.

#### Better Complexity:

- Time Complexity: **O(n)** (single pass to count + single pass to overwrite)
- Space Complexity: **O(1)** if only three counters are used, much better than Bubble, Selection, or Insertion Sort.

```

// sort array of 0s 1s 2s
// basic approach count no of 0s 1s 2s basically and add in arr
// good approach 0n is DNF 3 pointer

void countAndArrangeInArray(vector<int>& arr) {
    int cntZero = 0, cntOne = 0, cntTwo = 0;
    for (int i : arr) {
        if (i == 0) cntZero++;
        else if (i == 1) cntOne++;
        else if (i == 2) cntTwo++;
    }
    int i = 0;
    while (cntZero-->0) arr[i++] = 0;
    while (cntOne-->0) arr[i++] = 1;
    while (cntTwo-->0) arr[i++] = 2;
}

```

### Approach 6 : Dutch National Flag (DNF) or Three Pointers Approach

#### How it Works:

DNF Algorithm sorts 0s, 1s, and 2s in a **single pass** using three pointers: low for 0s, mid for current element, and high for 2s.

- If arr[mid] is 0 → swap with arr[low] and move low and mid forward.

- If `arr[mid]` is 1 → just move `mid` forward.
- If `arr[mid]` is 2 → swap with `arr[high]` and move `high` backward.  
After the pass, all 0s, 1s, and 2s are in order.

Complexity:

- Time Complexity: **O(n)** (single pass)
- Space Complexity: **O(1)** (no extra space) — the most efficient approach for this problem.

```
void printArr(vector<int>arr){
    for(auto i : arr) cout<<i <<" ";
    cout<<endl;
}
// 3 pointers
void dnfAlgo(vector<int>&arr){
    int start = 0;
    int mid = 0;
    int end = arr.size()-1;
    while(mid<=end){
        if(arr[mid] == 0){
            swap(arr[start],arr[mid]);
            printArr(arr);
            start++;
            mid++;
        }
        else if(arr[mid] == 1){
            printArr(arr);
            mid++;
        }
        else{
            swap(arr[mid],arr[end]);
            printArr(arr);
            end--;
        }
    }
}
```

**Output Window for DNF approach after every iteration we print arr for better understanding**

```
Enter the number of objects: 6
Enter the objects color (Red, White, Blue) in any case combination:
Enter color of object 1 out of 6 : Blue
Enter color of object 2 out of 6 : White
Enter color of object 3 out of 6 : Red
Enter color of object 4 out of 6 : White
Enter color of object 5 out of 6 : Red
Enter color of object 6 out of 6 : Blue
2 1 0 1 0 2
0 1 0 1 2 2
```

```
0 1 0 1 2 2
0 1 0 1 2 2
0 0 1 1 2 2
0 0 1 1 2 2
Your sorted array of object is :
Red Red White White Blue Blue
```

Main Functions to call all other functions

```
int colorToNumber(const string &color) {
    string temp = color;
    // Convert to lowercase for uniformity
    for(auto &c : temp) c = tolower(c);
    if(temp == "red") return 0;
    else if(temp == "white") return 1;
    else if(temp == "blue") return 2;
    else return -1; // invalid color
}

int main(){
    int n;
    cout << "Enter the number of objects: ";
    cin >> n;
    vector<int> arr(n);
    cin.ignore(); // flush newline after reading n

    cout << "Enter the objects color (Red, White, Blue) in any case
combination:\n";

    for(int i = 0; i < n; i++){
        string color;
        cout << "Enter color of object " << i+1 << " out of " << n << " : ";
        getline(cin, color);

        int val = colorToNumber(color);
        if(val == -1){
            cout << "Invalid input! Please enter Red, White, or Blue
only.\n";
            return 1;
        }
        arr[i] = val;
    }
    bubbleSort(arr);
    // selectionSort(arr);
    // insertionSort(arr);
    // countAndArrangeInArray(arr);
    // dnfAlgo(arr);
    // vector<int>mergeSortAns = mergeSort(arr,0,arr.size()-1);
    cout<<"Your sorted array of object is : \n";
    for(auto i : arr){
        if(i == 0) cout << "Red ";
        else if(i == 1) cout << "White ";
    }
}
```

```
    else if(i == 2) cout << "Blue ";  
  }  
}
```

### **Conclusion : Which algorithm is best and why**

In this assignment, we explored different approaches to sort an array containing only 0s, 1s, and 2s.

#### **1. Naive Approaches ( $O(n^2)$ complexity):**

- a. **Bubble Sort:** Repeatedly swaps adjacent elements if they are in the wrong order.
- b. **Selection Sort:** Selects the minimum element from the unsorted portion and places it in the correct position.
- c. **Insertion Sort:** Builds the sorted array one element at a time by inserting elements into their correct position.

These approaches are easy to understand and implement but have **time complexity  $O(n^2)$**  in the worst case. For large arrays, these algorithms are inefficient.

#### **2. Optimized Counting Approach (Two-Pass, $O(n)$ complexity):**

- a. Count the number of 0s, 1s, and 2s.
- b. Overwrite the original array with the correct number of 0s, 1s, and 2s.
- c. Time Complexity:  $O(n)$ , Space Complexity:  $O(1)$ .
- d. This is better than  $O(n^2)$ , but it requires **two passes** over the array.

#### **3. Dutch National Flag (DNF) Algorithm (Single-Pass, $O(n)$ complexity):**

- a. Uses **three pointers**: low, mid, and high.
- b. Traverses the array **once**, swapping elements to place 0s at the beginning, 2s at the end, and 1s in the middle.
- c. **Time Complexity:**  $O(n)$
- d. **Space Complexity:**  $O(1)$
- e. **Advantage:** Optimal for this specific problem because it **sorts in a single pass without extra space**.

#### **Final Observation:**

While traditional sorting algorithms (Bubble, Selection, Insertion) work for any array, they are **not efficient** for arrays with only three distinct elements. The **DNF algorithm** is the most suitable solution for this specific problem due to its **single-pass, in-place sorting** property, making it the **best choice** in terms of both time and space efficiency.

## LAB ASSIGNMENT 2

### ARRAYS

Q1. Take an input from the user in form of 1D array. Convert the given array in nxn form. If the items are lesser than nxn then insert '0' in the last left elements. Now based on the user input element and base address, determine the row major and column major address.

```
#include<bits/stdc++.h>
using namespace std;

// row major
int addUsingRowMajor(int i, int j, int noOfCols, int baseAddress, int rowLB
= 0, int colLB = 0) {
    int sizeOfElement = 4;
    return baseAddress + ((i - rowLB) * noOfCols + (j - colLB)) *
sizeOfElement;
}

int addUsingColMajor(int i, int j, int noOfRows, int baseAddress, int rowLB
= 0, int colLB = 0) {
    int sizeOfElement = 4;
    return baseAddress + ((j - colLB) * noOfRows + (i - rowLB)) *
sizeOfElement;
}

void printMatrix(vector<int>& arr, int n) {
    for(int i = 0; i < n; i++) {
        for(int j = 0; j < n; j++) {
            cout<<arr[i * n + j]<<" ";
        }
        cout<<endl;
    }
}

pair<int,int>findElement(vector<int>arr,int elem){
    int n = sqrt(arr.size());
    for(int i = 0; i < n; i++) {
        for(int j = 0; j < n; j++) {
            if(arr[i * n + j] == elem) {
                return {i, j};
            }
        }
    }
}
```

```

    return {-1, -1};
}

int main1(){
    int size,baseAddress;
    cout<<"Enter Number of Elements in the array : ";
    cin>>size;
    cout<<"Enter the Base Address : ";
    cin>>baseAddress;
    vector<int>arr(size);
    cout<<"Enter the Elements of the array : ";
    for(int i=0; i<size; i++) cin>>arr[i];
    // your n*n array means needed to append zero check current size of the
    array and nearest n*n
    int n = ceil(sqrt(size));
    arr.resize(n*n, 0);

    printMatrix(arr, n);
    // if user also enters number to get that particular address (but but
    this fails when he enters zeros !!!)
    // then simple using n^2 find element and then store its i and j

    int elemI = 0, elemJ = 0;
    int elem;
    cout<<"Enter Element to find its Address : ";
    cin>>elem;
    pair<int,int> pos = findElement(arr, elem);
    elemI = pos.first;
    elemJ = pos.second;

    // add logic to take lower bound ask user if wants to enter then yes
    else by defaultly steed as 0
    int rowLB = 0, colLB = 0;
    cout<<"Enter Lower Bound of Row (default 0) : ";
    cin>>rowLB;
    cout<<"Enter Lower Bound of Column (default 0) : ";
    cin>>colLB;

    // to print address to complete matrix
    // for(int i = 0; i < n; i++){
    //     for(int j = 0; j < n; j++){
    //         int addressRowMajor = addUsingRowMajor(i, j, n, baseAddress,
    rowLB, colLB);
    //         int addressColMajor = addUsingColMajor(i, j, n, baseAddress,
    rowLB, colLB);
    //         cout<<"Address of element at ("<<i<<","<<j<<") in Row Major
    Order : "<<addressRowMajor<<endl;
    //         cout<<"Address of element at ("<<i<<","<<j<<") in Column
    Major Order : "<<addressColMajor<<endl;
    //         cout<<endl;
    //     }
    // }

```



```

    // to print address to particular matrix
    if(elemI != -1 && elemJ != -1) {
        int addressRowMajor = addUsingRowMajor(elemI, elemJ, n, baseAddress,
rowLB, colLB);
        int addressColMajor = addUsingColMajor(elemI, elemJ, n, baseAddress,
rowLB, colLB);
        cout<<"Address of element at ("<<elemI<<","<<elemJ<<") in Row Major
Order : "<<addressRowMajor<<endl;
        cout<<"Address of element at ("<<elemI<<","<<elemJ<<") in Column
Major Order : "<<addressColMajor<<endl;
    } else {
        cout<<"Element not found in the array.";
    }
}

```

### Output

```

Enter Number of Elements in the array : 5
Enter the Base Address : 1000
Enter the Elements of the array : 1
2
3
4
5
1 2 3
4 5 0
0 0 0
Enter Element to find its Address : 3
Enter Lower Bound of Row (default 0) : 0
Enter Lower Bound of Column (default 0) : 0
Address of element at (0,2) in Row Major Order : 1008
Address of element at (0,2) in Column Major Order : 1024

```

### Question 3 :

A company has some interns under a project. In his coding class, who are practicing problems. Given the difficulty of the problems that the students have solved in order, help the Chef identify if they are solving them in non-decreasing order of difficulty. Non-decreasing means that the values in an array are either increasing or remaining the same, but not decreasing. That is, the students should not solve a problem with difficulty  $d_1$ , and then later a problem with difficulty  $d_2$ , where  $d_1 > d_2$ . Output "Yes" if the problems are attempted in non-decreasing order of difficulty rating and "No" if not.

```

bool isOrderCorrect(vector<int>order){
    int i = 0;
    while(i < order.size() - 1){
        if(order[i] > order[i + 1]) return false;
        i++;
    }
    return true;
}

int main2(){
    // basically take input of solving levels by student order by which
problems in solved
    // ausmme as numbers from 1 to INT_MAX;
    int questions;

```

```

    cout<<"Enter Number of Questions solved by Student : ";
    cin>>questions;
    vector<int>order(questions);
    cout<<"Enter Difficulty level of questions from 1 to INT_MAX that you
solved : \n";
    for(int i=0; i<questions; i++){
        cout<<"Enter difficulty level of "<<i+1 <<" question : ";
        cin>>order[i];
    }
    if(isOrderCorrect(order)) cout<<"Student Solved question in correct
order.";
    else cout<<"Student Solved Question in Wrong order.";
}

```

Output :

```

Enter Number of Questions solved by Student : 6
Enter Difficulty level of questions from 1 to INT_MAX that you solved :
Enter difficulty level of 1 question : 34
Enter difficulty level of 2 question : 23
Enter difficulty level of 3 question : 12
Enter difficulty level of 4 question : 56
Enter difficulty level of 5 question : 67
Enter difficulty level of 6 question : 1001
Student Solved Question in Wrong order.

```

## LAB ASSIGNMENT 3

### Complexity Analysis

Question1: Solve the following problem and also analyze its complexity.

A **left rotation** operation on a circular array shifts each of the array's elements 1 unit to the left. The elements that fall off the left end reappear at the right end. Given an integer  $d$ , rotate the array that many steps to the left and return the result.

**Example:**

$d=2$

$arr=[1,2,3,4,5]$

After 2 rotations,

$arr' = [3,4,5,1,2]$

Solution :

Brute force shifts elements one by one for each rotation  $\rightarrow O(n \cdot k)$  time,  $O(1)$  space.

Temp array directly places elements in new positions using modulo  $\rightarrow O(n)$  time,  $O(n)$  space.

```
// using O(n) and O(n) simple using temp array
void rotateKtimes2(vector<int>&nums,int k,vector<int>&temp){
    for(int i=0; i<nums.size(); i++){ //O(n)
        temp[(i+k) % nums.size()] = nums[i];
    }
}

// Brute force about n square and constant
void rotateKtimes1(vector<int>&nums,int k){
    int n = nums.size();
    for(int i=0; i<k; i++){
        //pickup last elem
        int last = nums[n-1];
        // shift all elems by 1 place
        for(int l=n-1; l>0; l--){
            nums[l] = nums[l-1];
        }
        nums[0] = last;
    }
}

int main(){
    vector<int>nums = {1,2,3,4,5};
    vector<int>temp(nums.size());
    int k = 7;
    k = k%nums.size(); // if 5 times rotate than original array back
    // rotateKtimes1(nums,k);
    rotateKtimes2(nums,k,temp);
    for(auto i : temp) cout<<i <<" ";
}
```

OUTPUT :

3 4 5 1 2

**Question 2: Solve the following problem and also analyze its complexity.**

Alice, an aspiring cryptographer, recently discovered an ancient scroll containing a sequence of mysterious numbers. According to a legend, these numbers hold the key to unlocking a hidden treasure buried centuries ago by an enigmatic mathematician. However, deciphering the scroll requires identifying specific pairs of numbers that follow an ancient numerical pattern.

Alice is given an array of integers `nums`. She wants to determine how many pairs  $(i, j)$  exist such that:

- $0 \leq i < j < \text{nums.length}$
- The **sum of digits of `nums[i]`** is equal to the **sum of digits of `nums[j]`**.

She believes that the correct count of these pairs will reveal a crucial clue needed to decode the next part of the scroll. Your task is to help Alice compute this number so she can continue her quest.

**Task**

Return the **number of special pairs** to assist Alice in uncovering the hidden secret.

**Input format**

- The first line contains a single integer  $N$  (size of the array).
- The second line contains  $N$  space-separated integers representing the array `nums`.

**Example**

Given:

$N = 4$ , `nums` = [51, 71, 17, 42]

Explanation:

- 51 → Sum of digits = 6
- 71 → Sum of digits = 8
- 17 → Sum of digits = 8
- 42 → Sum of digits = 6

Valid pairs:

- (1,4) → 51 & 42 (sum = 6)

- (2,3) → 71 & 17 (sum = 8)

**Total = 2**

**Solution :**

```
// sum of all digits one loop num.size() complexity
int sum(int num){
    int total = 0;
    while(num > 0){
        total += num % 10;
        num /= 10;
    }
    return total;
}

// counting pairs using nested loops matching each pairs basically comlexity
is n square where n is nums.size() and no extra space
int countSpecialPairs1(vector<int>& nums) {
    int n = nums.size();
    int count = 0;
    for(int i = 0; i < n; ++i) {
        int sumI = sum(nums[i]);
        for(int j = i+1; j < n; ++j) {
            int sumJ = sum(nums[j]);
            if(sumI == sumJ) count++;
        }
    }
    return count;
}

// counting pairs with hashmaps first store that paricular sum has that
number of pairs then use combination formula
// 0 n compelxity using 1 for loop and using maps also 0n
int countSpecialPairs2(vector<int>&nums){
    int n = nums.size();
    int ans = 0;
    unordered_map<int,int>freq;
    for(int i=0; i<n; i++){
        int total = sum(nums[i]);
        freq[total]++;
    }
    for(auto f : freq){
        int frequency = f.second;
        if(frequency > 1) ans += frequency*(frequency-1)/2;
    }
    return ans;
}

int main(){
    vector<int>nums = {23,32,14,41,50};
    cout<<countSpecialPairs2(nums);
}
```

## OUTPUT :

10