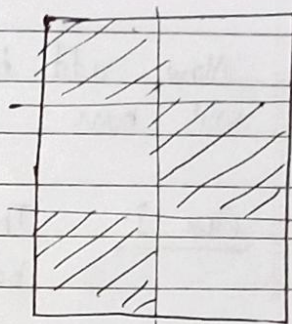


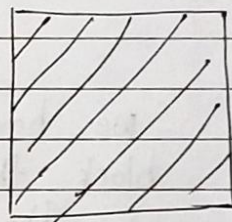
1] Induction Hypothesis:

If we use $n=k$ lines to create polygonal region then it is possible that no two adjacent polygon have same colour and we can colour the entire page using just two colours / max 2 colour are needed



Base Case: For $n=0$

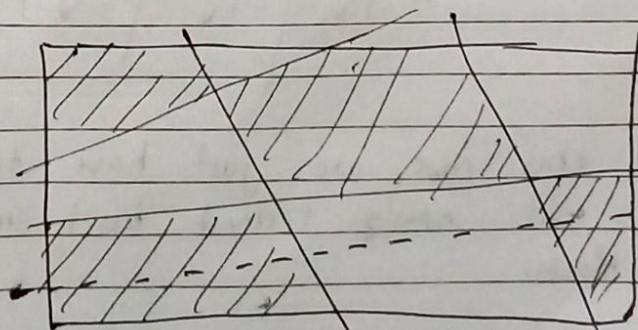
The entire page can be coloured using just one colour, and all other conditions are also satisfied



Induction step

For $n = k+1$

Let there be $k+1$ lines that divide the page into the polygonal region,

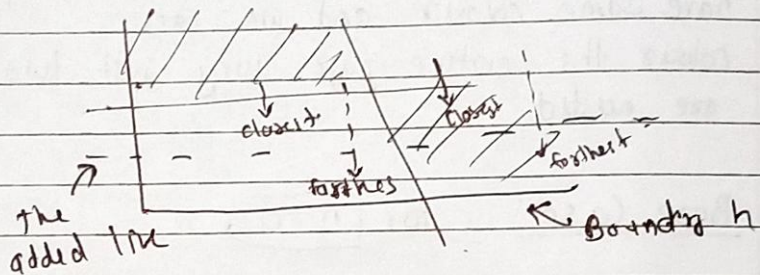


Any arbitrary but particular line

Now, remove any arbitrary, but particular lines. Hence, now we have n lines which satisfy all the conditions from the hypothesis.

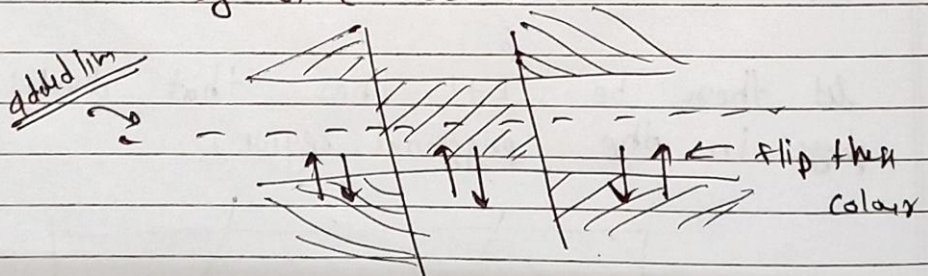
Now add back the removed line. we will have two patterns

Case I: The formed section is near boundary



we have to check. the adjacent of the block that got divided and give the same colour as the adjacent to the one who is most. farthest and the opposite to the the the closest.

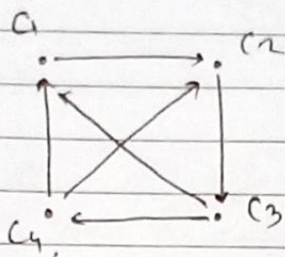
Case II: The formed section is surrounded by other section.



In this case we just have to flip the colours of newly formed block and it's adjacent old blocks.

Hence proved

e]



Induction : Hypothesis.

let the given statement be true for $n = k$ cities.

Base case

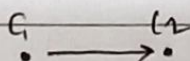
For $n = 0$

It is true as there are no cities to travel

For $n = 1$

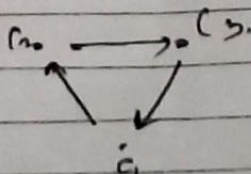
It is true as there is only one city

for $n = 2$,



Here $C2$ is that city.

for $n = 3$



Hence all the cities satisfy the given condition.

3] Base case

Every positive integer n can be expressed as distinct sum of power of 2.

For $n = 1$

$$\boxed{1 = 2^0} \quad \text{Hence true}$$

For $n = 2$

$$\boxed{2 = 2^1} \quad \text{Hence true}$$

Induction Hypothesis

The the statement stand true for $n = km$
 st $1 \leq m \leq k$

Induction step

let $n = k+1$.

Case I : $(k+1)$ is even
 $\therefore \left(\frac{k+1}{2}\right)$ is a integer

$$\therefore 1 \leq \frac{k+1}{2} \leq k.$$

$\therefore \frac{k+1}{2}$ is ~~true~~ satisfies the condition

$2 \left(\frac{k+1}{2}\right)$ also satis-fies the

condition and has the ~~*~~ distinct power of two as 2 is added to even power

Induction step

For $n = k+1$

Now let's ~~remove~~ ^{remove} one city

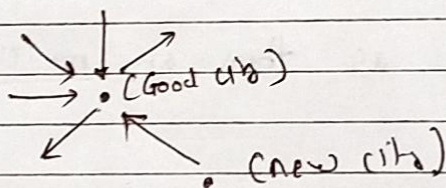
\therefore Number of cities = k

By Induction Hypothesis

For $n = k$

there exist a city which satisfies all the given condition.

Let that city be "Good city"



Now build a new city and build a road which goes from 'new city' to 'good city',

$\therefore n = k+1$

and 'Good city' satisfies all conditions

Hence proved

Set which was used to represent $\frac{(k+1)}{2}$.

Case II: $(k+1)$ is odd

$\therefore k$ is even

$\therefore k$ does not have 2^0 and is represented using distinct power of two

$\therefore k+1 = \text{power set of } k + 2^0$

Hence proved

Q1] We know, for the graph with n vertices
max No. edge in simple graph

$$= \frac{n(n-1)}{2}$$

$$\therefore \text{Max value of } m = \frac{n(n-1)}{2}$$

But it is possible that all the edges may not be present.

$$\therefore m \leq \frac{n(n-1)}{2}$$

$$\therefore \frac{2m \leq n(n-1)}{2m \leq n^2 - n}$$

Q5] Let us assume that every one in the group has different No. of friends.

So for n people in group:

Following set represent the No. of friend each person have

$$0, 1, 2, 3, \dots, n-1$$

But according to this assumption one person is friends with every other person but at the same time there is a person with No friends.

Hence our assumption is false

\therefore The given statement is true.

Q8) $r \rightarrow$ Degree of each vertex
 $n \rightarrow$ No. of vertex
 $m \rightarrow$ No. of edge

As r is Degree of each vertex
and n is No. of vertex

$r \times n \Rightarrow$ summation of degree's of all vertices

We know


If we consider one edge then it add
2 to the summation of degrees

$$\therefore 2m = r \times n$$

$$\therefore m = \frac{r \times n}{2}$$

Q7] Base case:

$$n = 2.$$

i.e. 

$$\delta(G) = 1$$

$$\frac{2}{2} = 1.$$

$$\therefore \delta(G) \geq \frac{1}{2}.$$

Induction Hypothesis

Let us assume that the statement stands true for $n = k$.

$$\therefore \delta(G) \geq \frac{k}{2} \text{ and it is connected}$$

Induction Step

$$\text{Let } n = k+1$$

ign ~~lets~~ ~~ignore~~ the re.

To create this graph let us use the graph G' with $n = k$.

$$\therefore \delta(G') \geq \frac{k}{2}$$

Now add a vertex and draw any number of edges to it, to that ^{new} vertex from the remaining vertex of G' .

\therefore minimum value of
 $\delta(G) = \delta(G')$

Case I: If k is even.

$$\text{the } \frac{k+1}{2} = \frac{k}{2}$$

as they both are integers

$$\therefore \delta(G) \geq \frac{k+1}{2}$$

Case II: If k is odd.

then \Rightarrow

$$\max \text{ of } \frac{k+1}{2} = \delta(G)$$

Hence proved

Q2] Let us assume there is a point P with odd degree

Let this be a part of simple connected graph.

By, by previous Lemma we know that the No. of even odd ordered vertices is always even

\therefore there must be one more vertex with odd order,

And as this graph is connect, then there always exist a path from P to every other vertex including the vertex with odd order

Examples:

