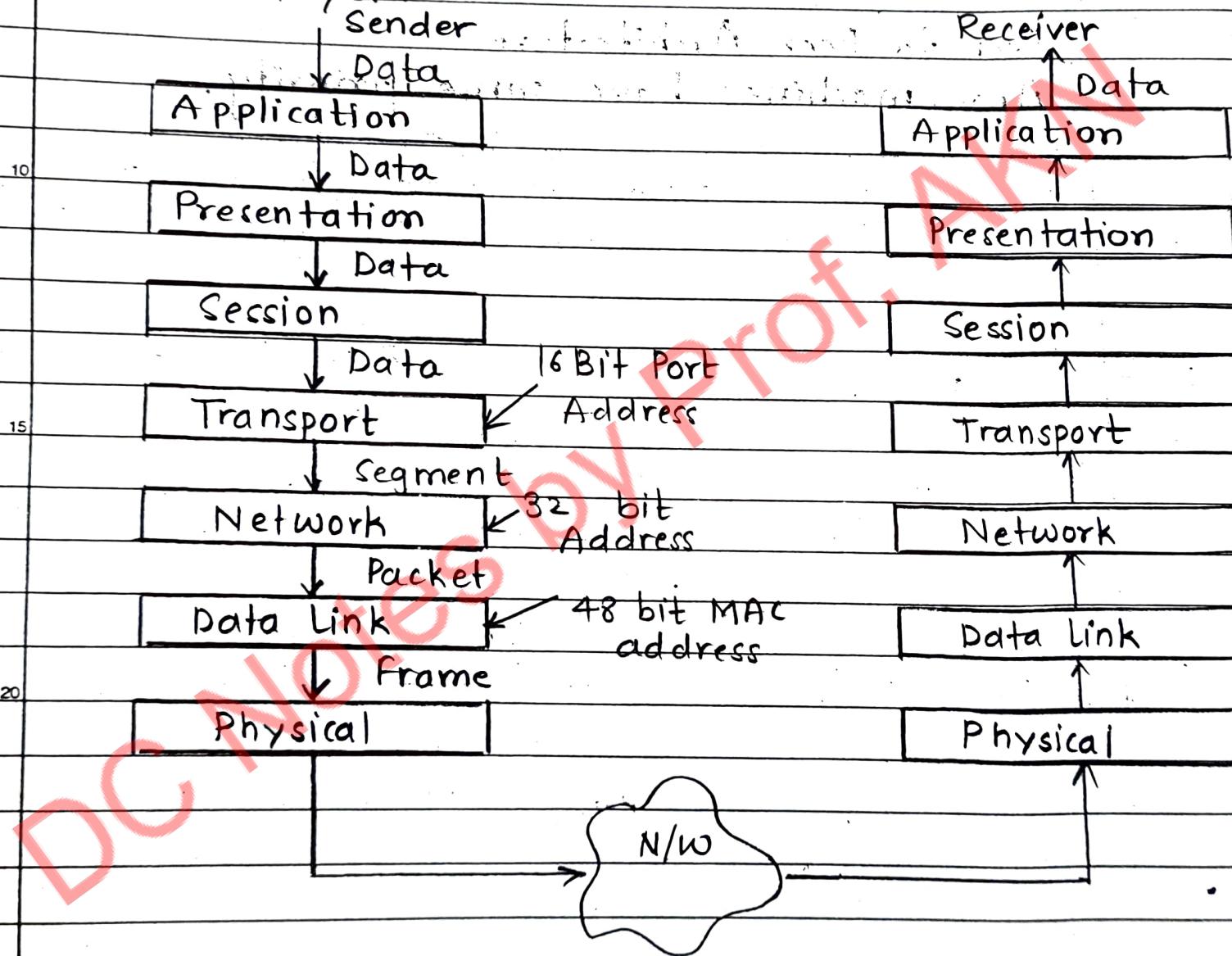


## Open System Interconnection :- (OSI)

It is a 7 layer architecture defined by ISO (International Standard Organisation) defines how exactly the n/w works i.e. functionality of each layer.



## Application Layer :-

- This layer is responsible for interacting with the user to run its application.
- Eg: Imagine a web browser (http) to access the web pages.

## Presentation Layer :-

This layer is responsible for taking data from application layer and perform following 3 things

- (1) Converting data from user understandable form to system understandable form
- (2) Compression if required
- (3) Encryption for security

## Session Layer :-

This layer is responsible for dialogue control and token management to establish connection with server.

## Transport Layer :-

This layer performs process to process communication for which it uses 16 bit port addressing.

It also performs:-

(1) Error control

(2) Flow control

(3) Congestion control

## Network Layer :-

This layer performs host to host communication for which it uses 32 bit IP address.

It also performs fragmentation if required.

### ~~Notes by Prot.~~ Data Layer :-

This layer maps a packet and converts it into a frame, this is known as framing. It also performs error control and has a sub layer known as MAC (Medium Access Control) that takes care of collision avoidance and prevention technique.

### Physical layer :-

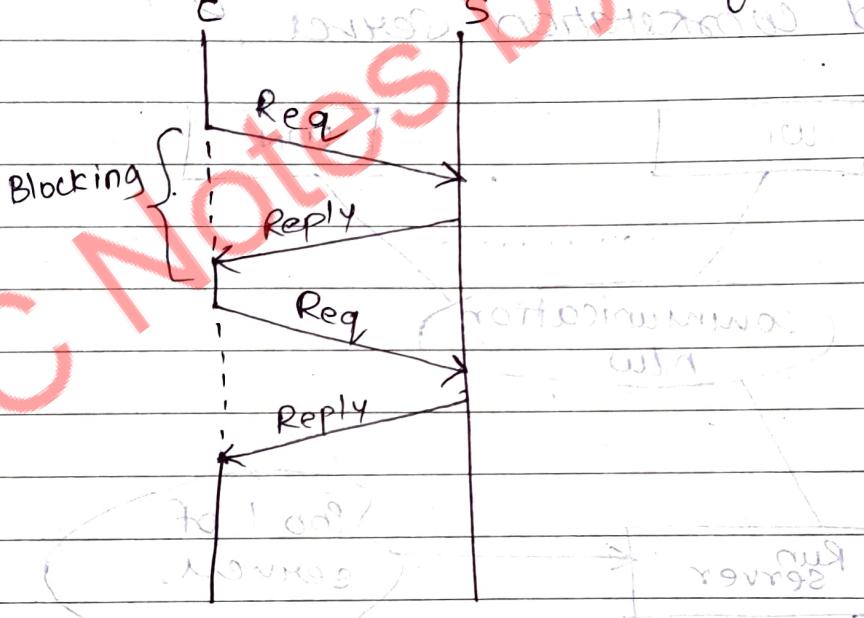
This layer performs transmission of raw bits of 0's and 1's i.e. electrical s/g to receiver's physical layer through network.

## \* Message oriented communication. [Types of communication]

1. Synchronous (blocking)
2. Asynchronous (Non blocking).
3. Persistent.
4. Transient.

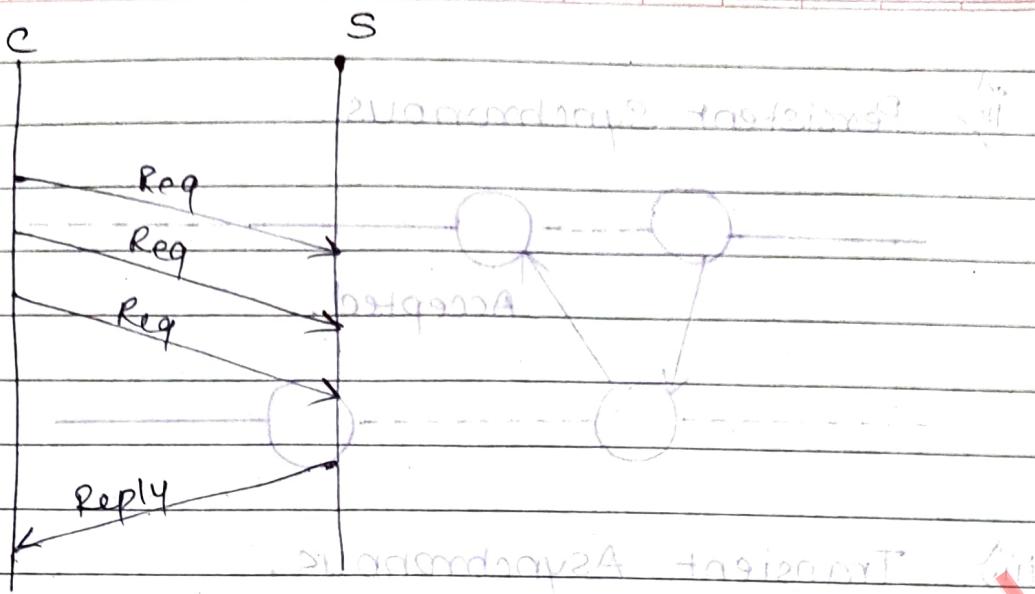
### 1. Synchronous / blocking.

The process where the request is sent and then blocked until reply is received is known as synchronous / blocking.



### 2. Asynchronous / Non-blocking -

The process where the request is sent continuously without getting blocked the reply is received is known as asynchronous non-blocking.



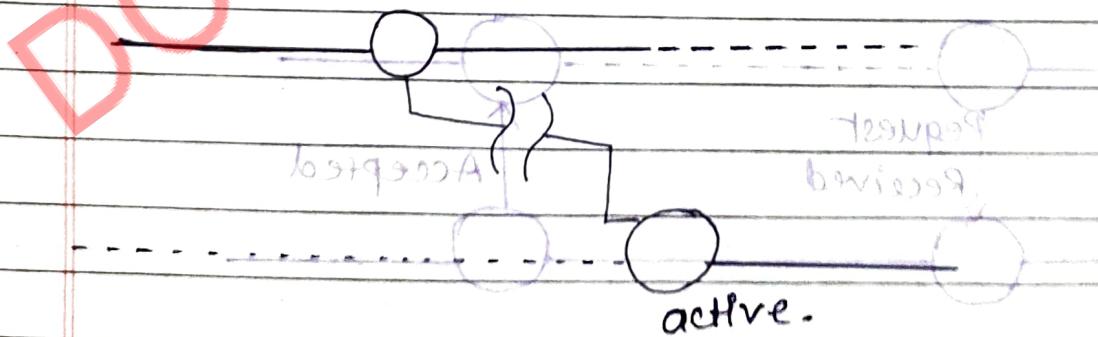
### 3. Persistent

The fact where the client and server both may or may not be online.

### 4. Transient

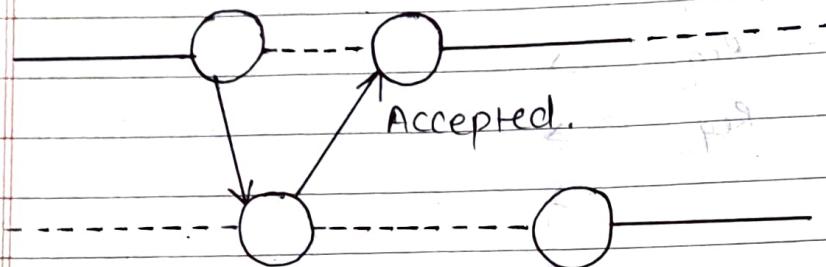
The fact where the client and server both has to be online.

\* POSSIBLE COMBINATION OF SYNCHRONOUS, ASYNCHRONOUS, PERSISTENT, AND TRANSIENT.

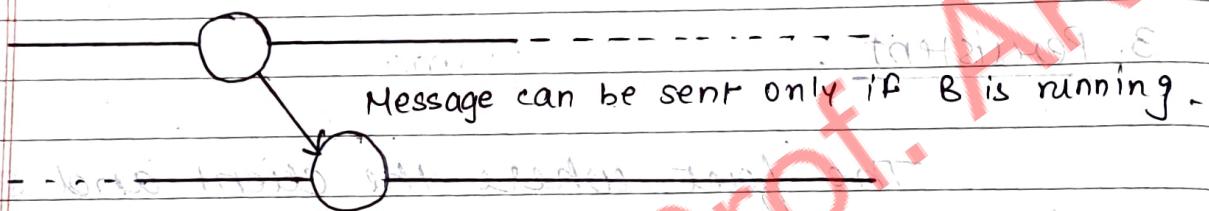


↳ Persistent Asynchronous.

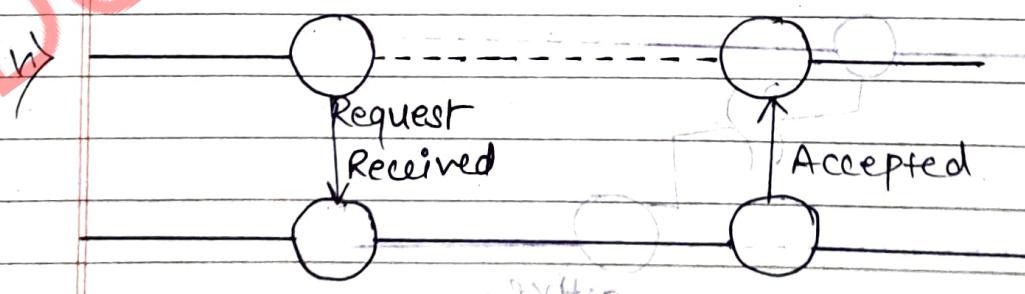
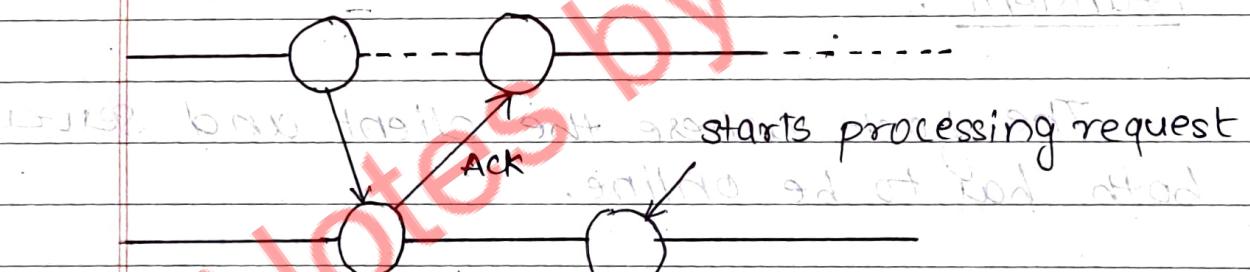
i) Persistent Synchronous.



ii) Transient Asynchronous.

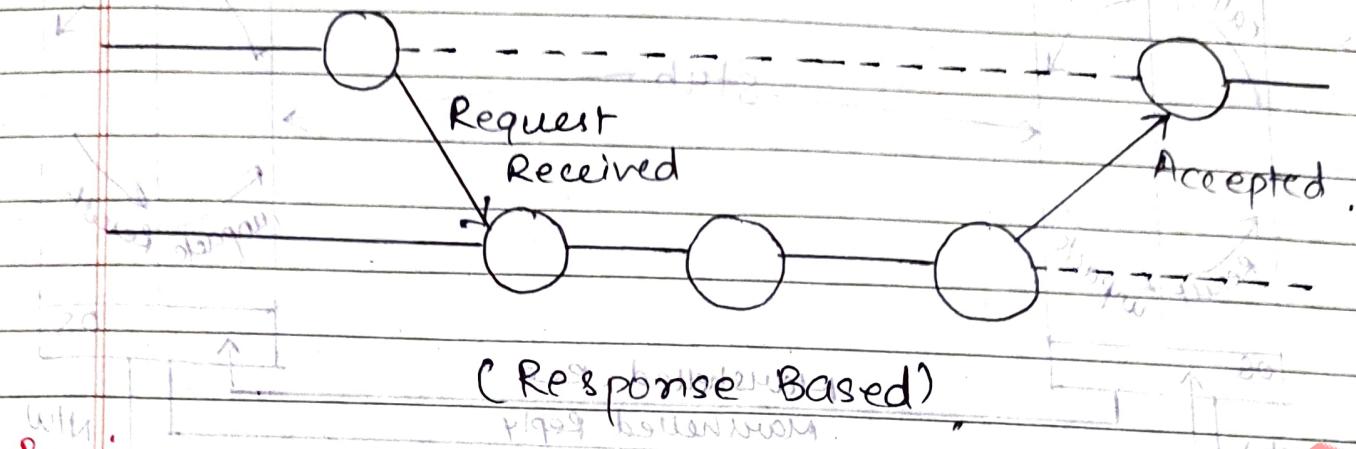


iv) Transient Synchronous.



Transient Synchronous (Delivery Based).

## vi) Transient Synchronous



\* Remote Procedure Call (RPC) [concept]:

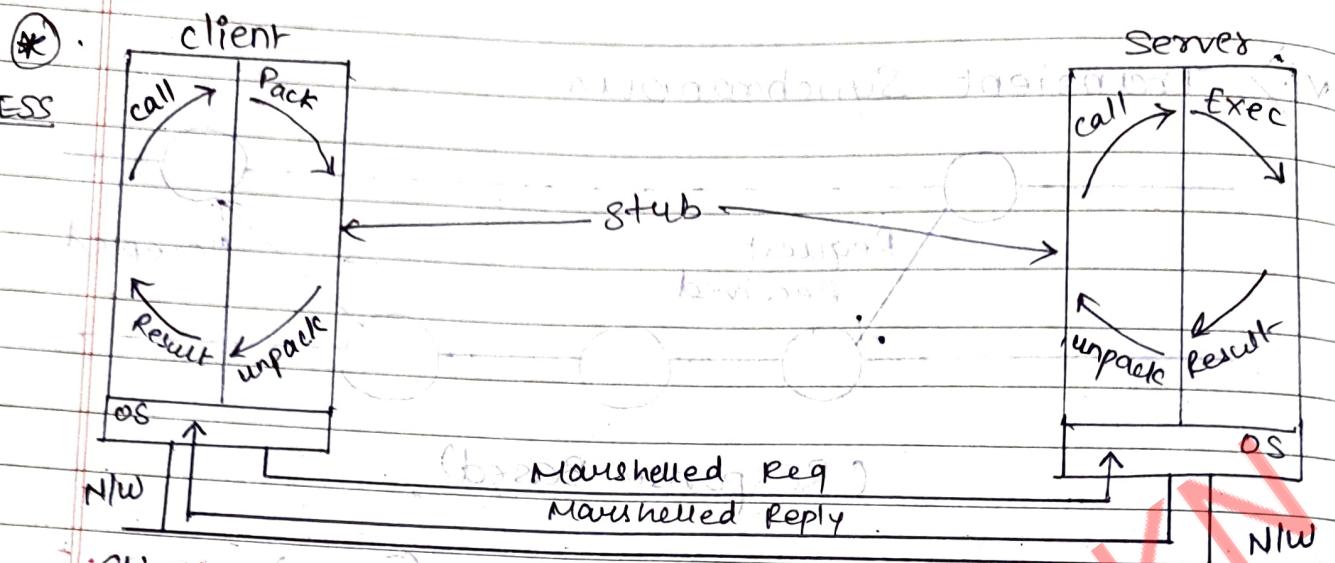
The process of calling a function which is not on the same machine, but present on some other machine in the network, this is known as remote procedure call.

Eg: Given source code of function  $f()$  in file  $P_1$  to  $P_8$  respectively.

list, arguments, etc.  $\rightarrow$  Exec  $f()$

reset functionality to original

Parameter: Parametric input data having  $f()$  function, unique to array.



1. Client process makes the call and realises that the function to be executed is not present on client, hence this request is handed over to server an interface known as "Stub".
2. Client stub locates the server having required function and same version of stub.
3. Client stub packs all the parameters, this is known as "Marshalling".
4. Marshalled request is forwarded to server, server stub unpacks this request, which is known as unmarshalling.
5. Server stub calls the required function, server process executes the function and returns the function to server stub.
6. Server stub packs the result and gives marshalled reply to client.

8. Client stub unpacks the packet, and returns the result to the calling process.

Transparencies to be maintained by RPC.

1. Location (ind. where soft. runs)

2. Migration (w.r.t program).

3. Relocation (w.r.t function).

Parameter passing semantics.

1. Call by value.

2. Call by Reference.

~~Info~~ RPC Failures (Process resilience)

i) Client cannot locate the server.

→ Client function is not available.

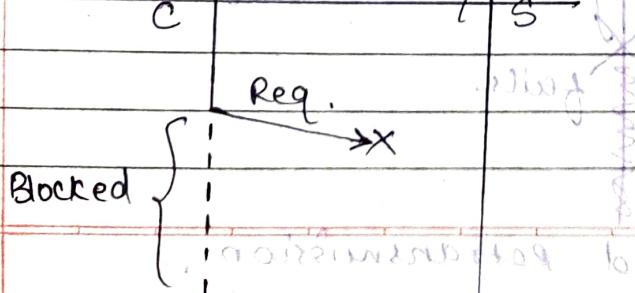
→ Server has crashed or gone down.

→ Stub mismatch.

→ Client new failure.

→ Server new failure.

ii) Client lost request.

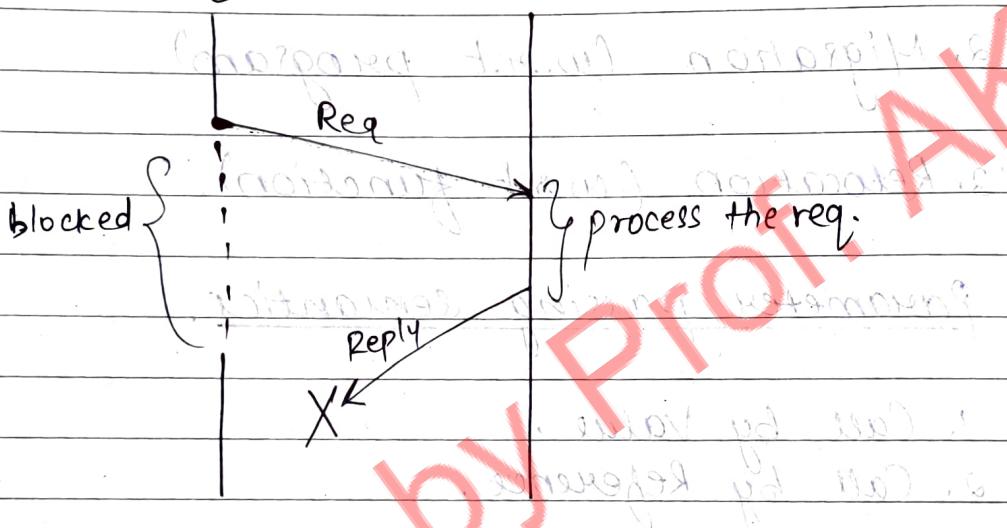


The request is being sent to server by Client but the request is lost.

Sol<sup>n</sup> : Timeout and retransmission.

3) Server lost reply.

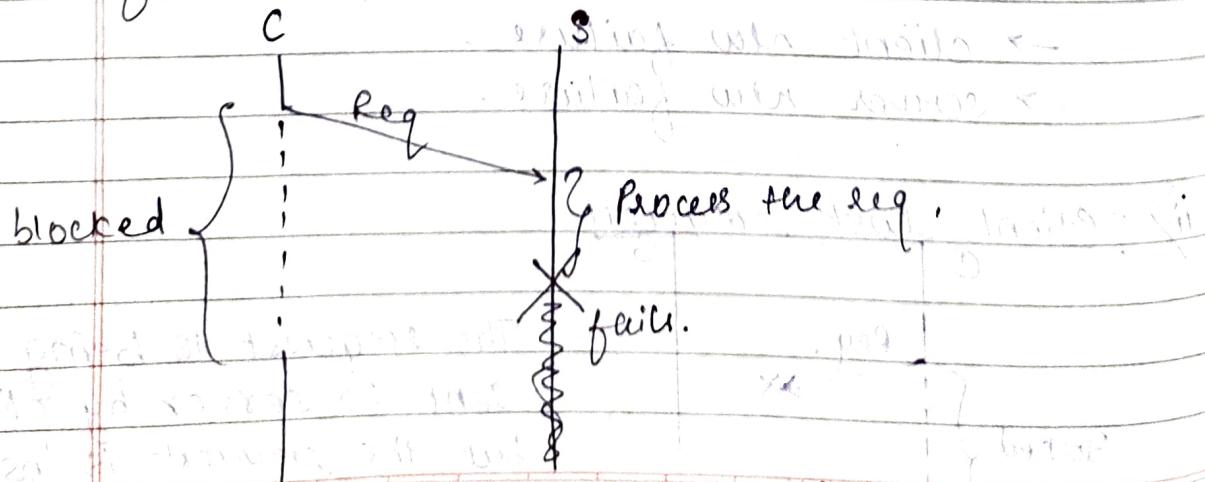
client sends the request and gets blocked,  
server sends the reply, but it gets lost.



Sol<sup>n</sup> : Timeout and Retransmission.

4) Server fails

Client sends the function, server starts executing and ~~in middle of~~ during execution it fails.

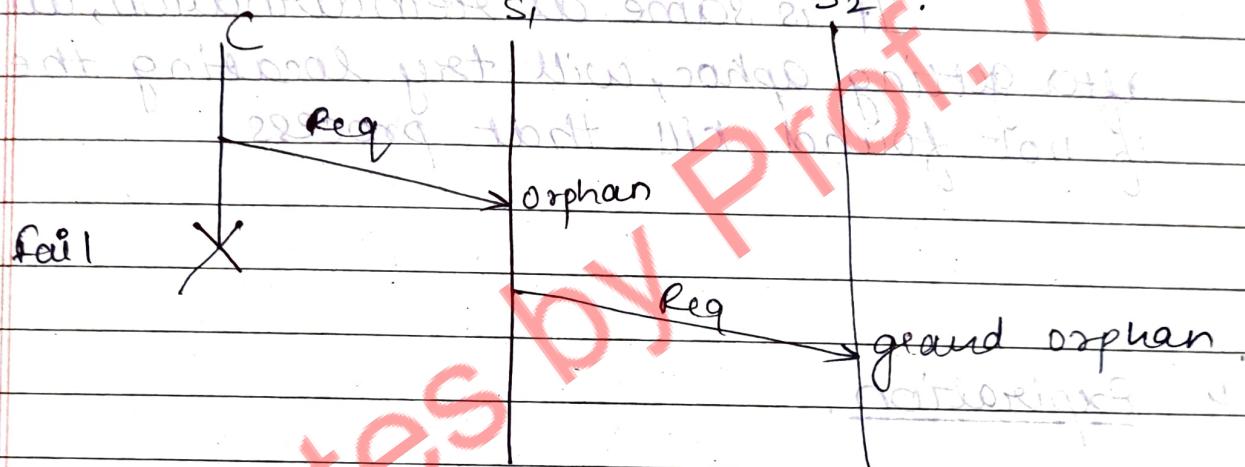


Sol<sup>n</sup> : Timeout and retransmission.

## S) soft client failure

Client sends the request, and crashes  
server processes the request, ~~if their~~ and  
hence that function becomes orphan as  
there is no existence of calling parent.  
and if therefore occurs a situation where  
there is function call within function then  
upcoming all functions would be grand orphans.

class M, Autoremoves S1 space is S2.



Soln: upon login traps main

Ex-termination (i.e. killing)

client will enter details of server (S<sub>1</sub>) in log, when request is made. If client fails and reboot, it will check the log and inform that server (S<sub>1</sub>) to kill all previous transaction of client, this will kill only orphans.

## 2. Reincarnation

When client after failure reboots, it will broadcast epoch (time) and inform all by broadcasting epoch to kill all processes which has started before epoch, even though it will kill orphans & grand orphans it will kill all processes.

## 3. Gentle Reincarnation

It is same as reincarnation, all servers after getting epoch, will try locating the owner, if not found kill that process.

## 4. Expiration

Server keeps a timeout it waits for a client request again during that time period. If client does not request then server checks if client alive. If it is then complete the execution and sends back result.

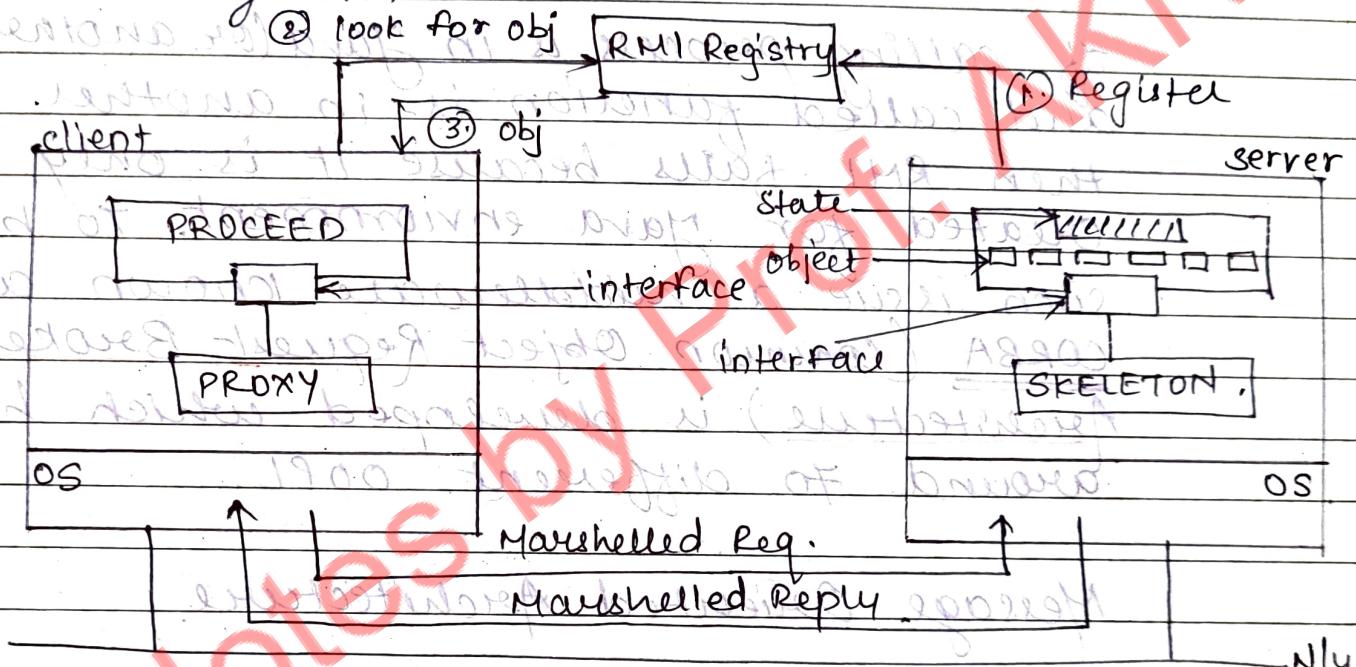
Notes by Prof. AKA

Imp

\* Remote Method Invocation (RMI) / Remote  
Object Invocation. [Implementation of RPC in Java].

RMI is implementation of RPC in Java.  
i.e. both calling and called program should  
be in Java.

\* Block Diagram.



Server initially register interface and object or RMI registry.

Client process makes the call and realises that the required function is not on the same machine and hence forwards this request to proxy. Proxy which is client side stub, looks up for server and object in RMI registry, gets the required object, implements the same interface as that of Server.

Proxy now packs the parameters (marshalling) and sends this marshalled request to server.

Server skeleton, (server stub) unpacks the request (unmarshalling) and with required object calls the function. Server process executes and returns the result to skeleton, which packs this and gives it to client. Client proxy unpacks the result and handover the result to calling process.

If calling process is in Java (or another OOPL) and called function is in another OOPL then RMI fails because it is only created for Java environment. To handle such issue a middleware known as CORBA (Common Object Request Broker Architecture) is developed which handles around 70 different OOPL.

## ~~Message Oriented Architecture~~

- \* Types of Communication, (Same as message oriented communication).
- \* Message Broker Architecture (Refer notes).
- \* Message queuing architecture (Refer notes).
- \* Stream oriented communication. (Refer notes).
- \* Group communication. (Refer notes).