

# Error Handling

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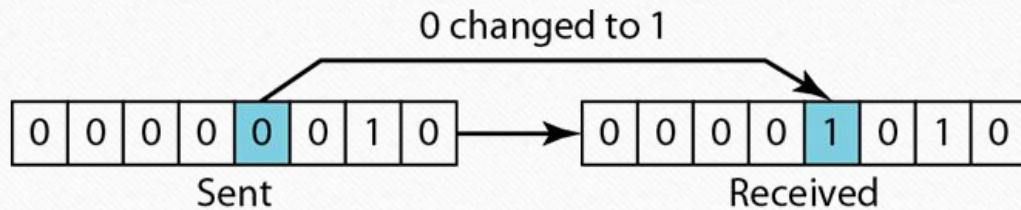
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Assistant Professor

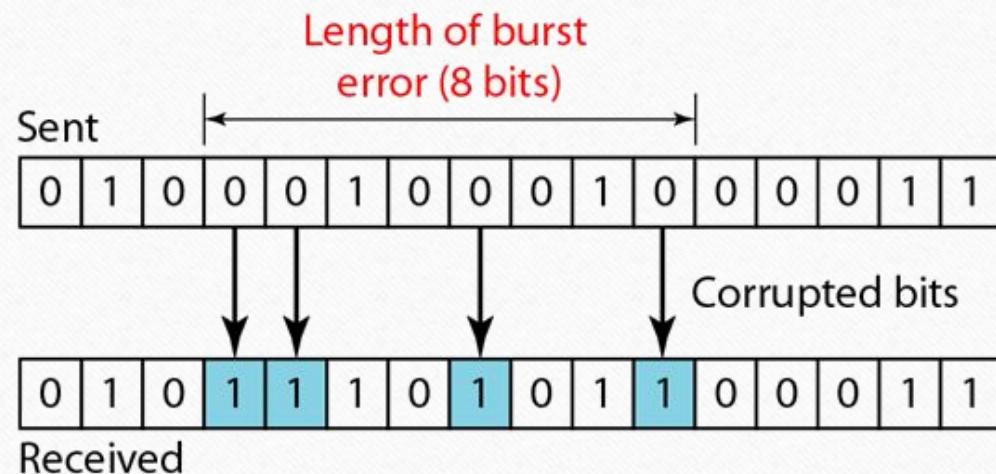
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## Error



**Single-Bit Error:** The term **single-bit error** means that only 1 bit of a given data unit (such as a byte, character, or packet) is changed from 1 to 0 or from 0 to 1.



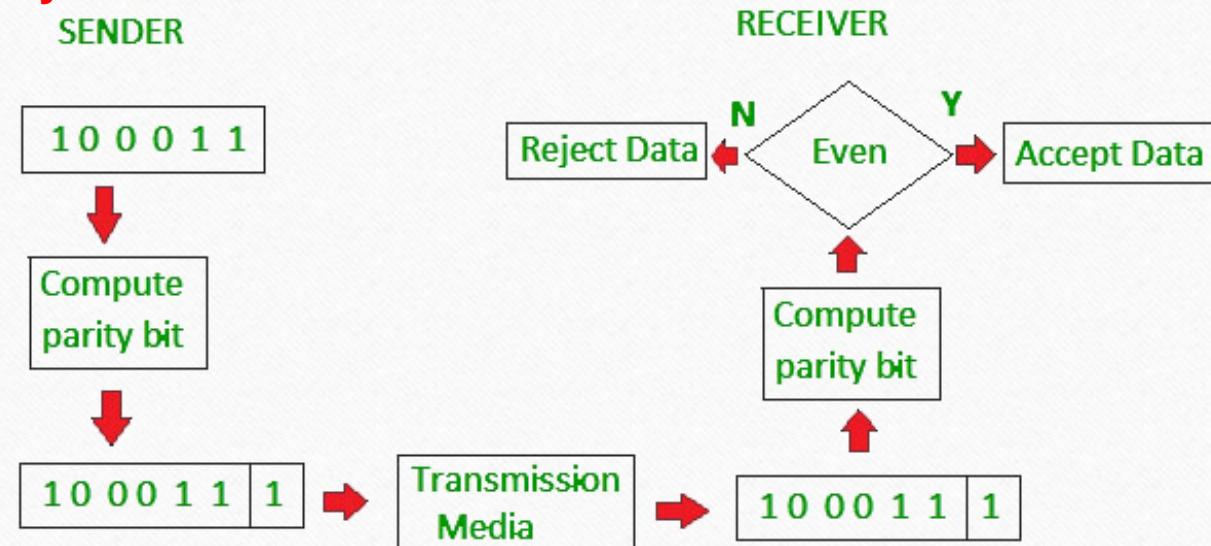
**Burst Error:** The term **burst error** means that 2 or more bits in the data unit have changed from 1 to 0 or from 0 to 1.

# Error Detection

- 1. Parity Check**
- 2. Checksum**
- 3. Cyclic Redundancy Check (CRC).**

## Parity Check

The parity check is done by adding an extra bit, called parity bit to the data to make a number of 1s either even in case of even parity or odd in case of odd parity.



## 2D Parity Check

Original Data

|          |          |          |          |
|----------|----------|----------|----------|
| 10011001 | 11100010 | 00100100 | 10000100 |
|----------|----------|----------|----------|

Parity check bits are calculated for each row and column.

Row parities

|                 |   |
|-----------------|---|
| 1 0 0 1 1 0 0 1 | 0 |
| 1 1 1 0 0 0 1 0 | 0 |
| 0 0 1 0 0 1 0 0 | 0 |
| 1 0 0 0 0 1 0 0 | 0 |
| 1 1 0 1 1 0 1 1 | 0 |

Column parities →

|           |           |           |           |           |
|-----------|-----------|-----------|-----------|-----------|
| 100110010 | 111000100 | 001001000 | 100001000 | 110110110 |
|-----------|-----------|-----------|-----------|-----------|

Data to be sent

# Checksum

Original Data

|          |          |          |          |
|----------|----------|----------|----------|
| 10011001 | 11100010 | 00100100 | 10000100 |
|----------|----------|----------|----------|

1

2

3

4

k=4, m=8

Sender

1 10011001

2 11100010

101111011  
1

01111100

3 00100100

10100000

4 10000100

100100100  
1

Sum: 00100101

CheckSum: 11011010

Reciever

1 10011001

2 11100010

101111011  
1

01111100

3 00100100

10100000

4 10000100

100100100  
1

00100101

11011010

Sum: 11111111

Complement: 00000000

Conclusion: Accept Data



## Error Detection and Correction

Error correction techniques find out the exact number of bits that have been corrupted and as well as their locations.

### Hamming Code

Hamming code is a set of error-correction codes that can be used to detect and correct the errors that can occur when the data is moved or stored from the sender to the receiver. It is technique developed by R.W. Hamming for error correction

$$2^r \geq m + r + 1$$

where, r = redundant bit, m = data bit

Suppose the number of data bits is 7, then the number of redundant bits can be calculated using:

$$= 2^4 \geq 7 + 4 + 1$$

Thus, the number of redundant bits= 4

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# Hamming Example

| Position | R8 | R4 | R2 | R1 |
|----------|----|----|----|----|
| 0        | 0  | 0  | 0  | 0  |
| 1        | 0  | 0  | 0  | 1  |
| 2        | 0  | 0  | 1  | 0  |
| 3        | 0  | 0  | 1  | 1  |
| 4        | 0  | 1  | 0  | 0  |
| 5        | 0  | 1  | 0  | 1  |
| 6        | 0  | 1  | 1  | 0  |
| 7        | 0  | 1  | 1  | 1  |
| 8        | 1  | 0  | 0  | 0  |
| 9        | 1  | 0  | 0  | 1  |
| 10       | 1  | 0  | 1  | 0  |
| 11       | 1  | 0  | 1  | 1  |

Determining the position of redundant bits –

These redundancy bits are placed at the positions which correspond to the power of 2.

As in the above example:

1. The number of data bits = 7
2. The number of redundant bits = 4
3. The total number of bits = 11
4. The redundant bits are placed at positions corresponding to power of 2- 1, 2, 4, and 8

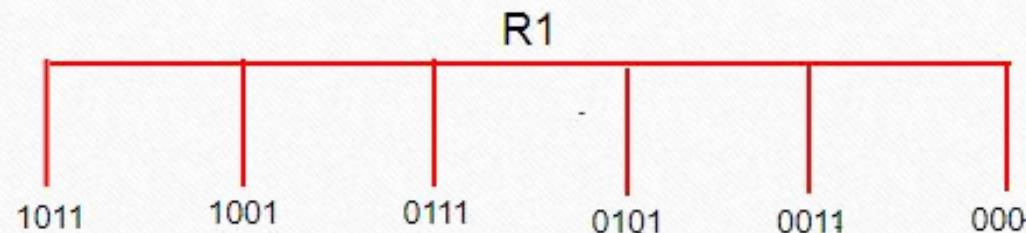


## Hamming Example

Suppose the data to be transmitted is 1011001, the bits will be placed as follows:

| 11 | 10 | 9 | 8  | 7 | 6 | 5 | 4  | 3 | 2  | 1  |
|----|----|---|----|---|---|---|----|---|----|----|
| 1  | 0  | 1 | R8 | 1 | 0 | 0 | R4 | 1 | R2 | R1 |

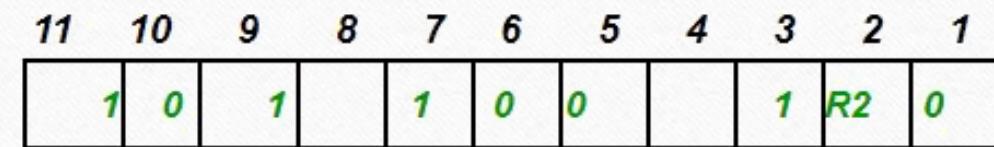
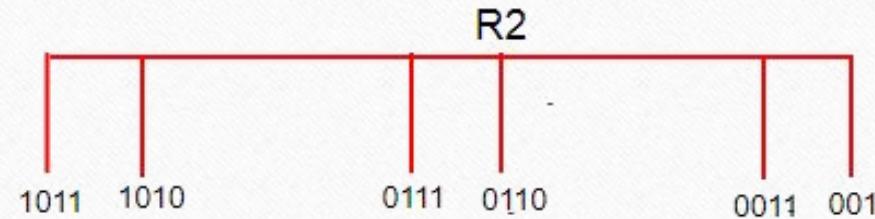
R1: bits 1, 3, 5, 7, 9, 11



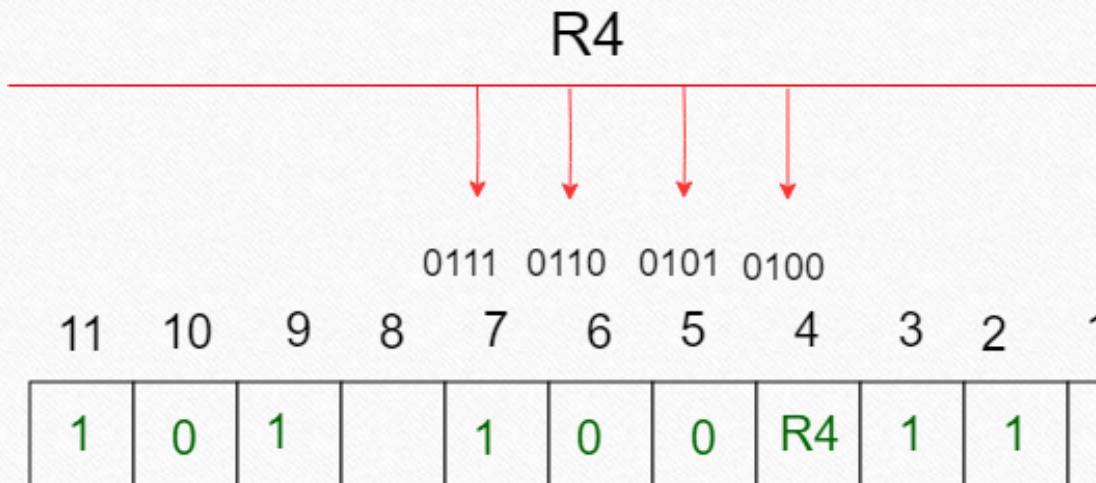
| 11 | 10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1  |
|----|----|---|---|---|---|---|---|---|---|----|
| 1  | 0  | 1 |   | 1 | 0 | 0 |   | 1 |   | R1 |

# Hamming Example

**R2: bits 2,3,6,7,10,11**

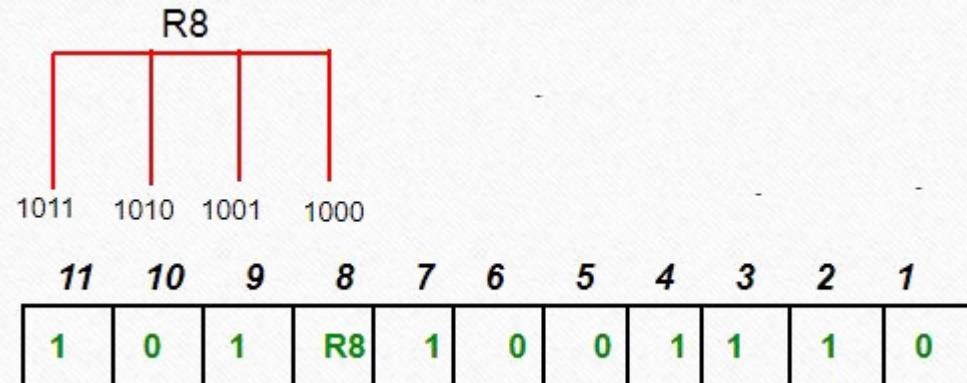


**R4: bits 4, 5, 6, 7**

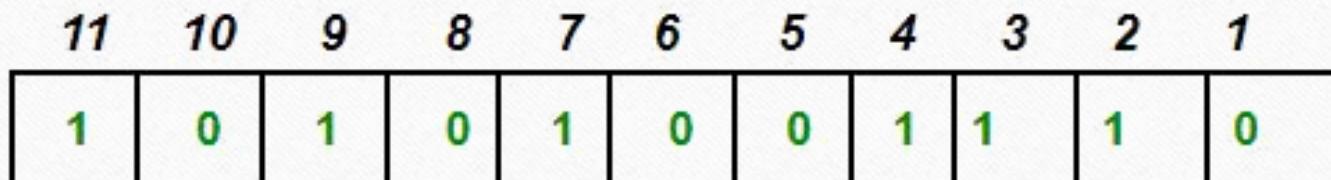


# Hamming Example

**R8: bit 8,9,10,11**

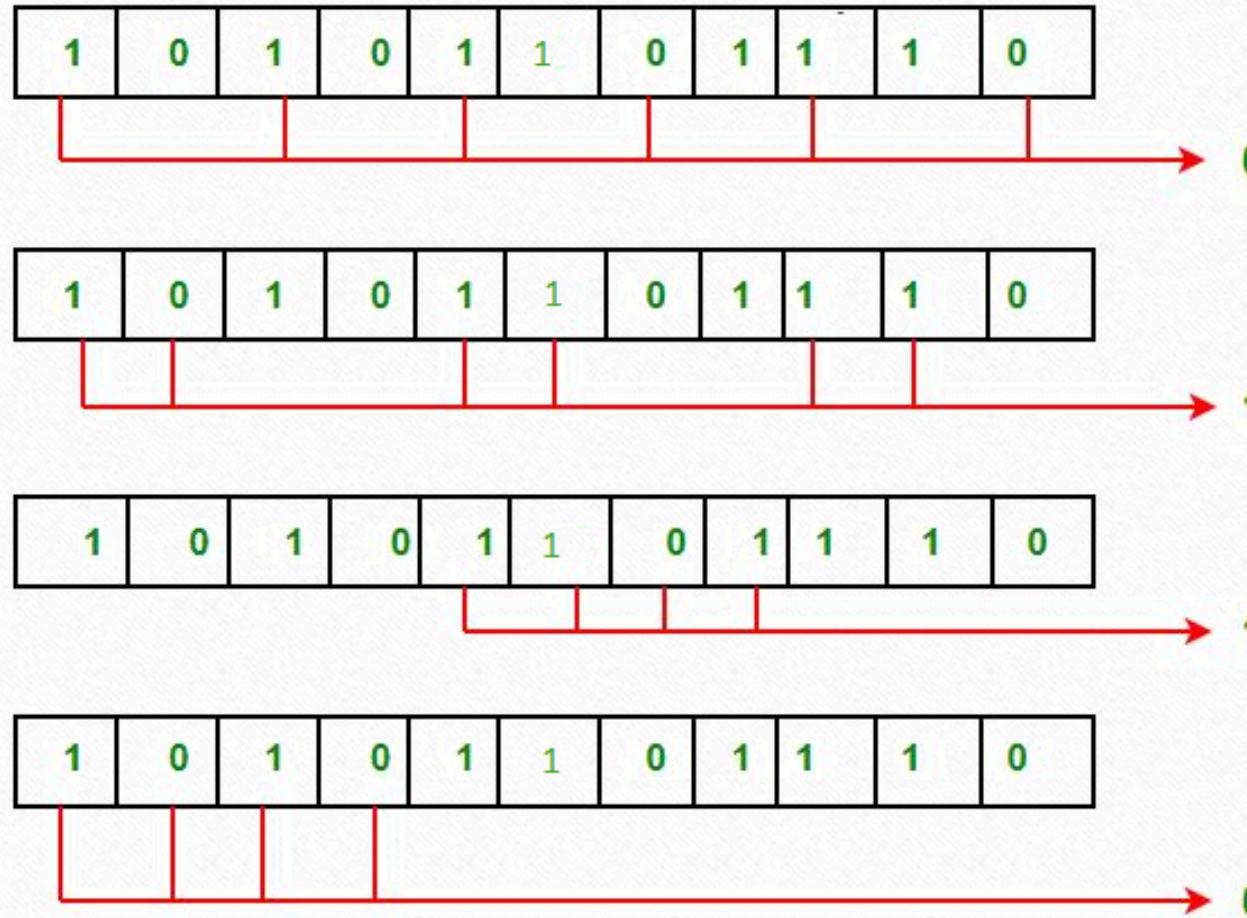


**Thus, the data transferred is:**



# Hamming Example

## Error detection and correction –



**PROF. AMIT K. NERURKAR**



# Thank You

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