



# CSS

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VIVA  
User Authentication: checking validity of user

eg. Password, FOTP, Biometrics, smart card.

VIVA  
Entity Authentication:

- Both ~~parties~~  $S^x$  &  $R^x$  has to verify each other.

eg Mutual Authentication

### Authentication Protocols

VIMP (VIVA 2 Interview)

① One way authentication

- A & B communicates
- B authenticates A

eg: 1 One factor

- password

eg: 2 Two factor

- password + OTP

- Card + PIN

② Mutual authentication  
(Also called as 2nd party authentication)

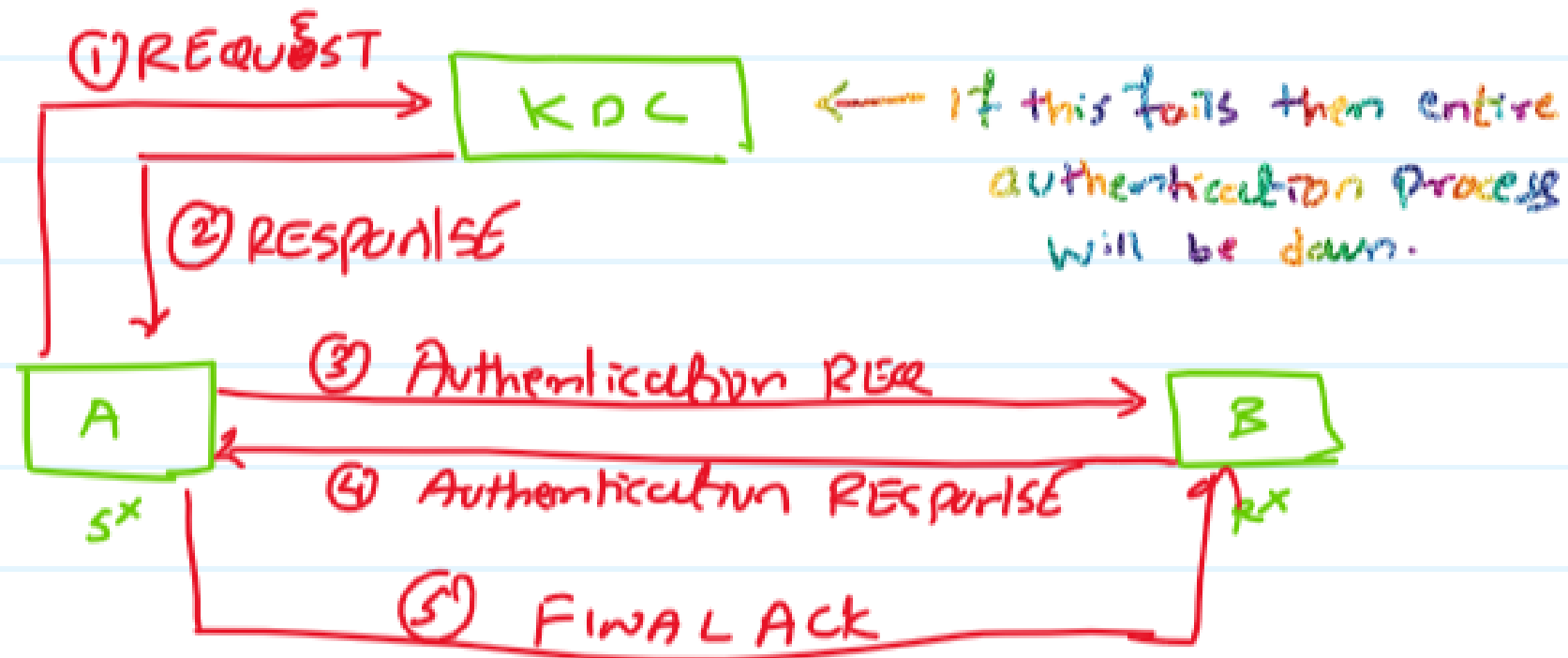
- A & B communicates
- A authenticates B
- B authenticates A

eg: KDC

Needham  
Schroeder  
Kerberos

Key Distribution  
Center

Needham Schroeder : 1<sup>st</sup> mutual authentication algo  
developed in 1978, uses KDC  
ie 3<sup>rd</sup> party authentication.



I A sends Request to KPC telling its wants to communicate with B.



$(ID_A, ID_B, N_i)$

$ID_A$ : Identity of A ( $CS^*$ )     $N_i$ : Nonce (Random no.)  
 $ID_B$ : Identity of B ( $CP^*$ )

II KPC will now create RESPONSE packet, ~~with~~ for A,  
which also contains Response for B.



FORMAT:  $E(\text{key}, [\text{msg}])$

msg is Encrypted using key

$\rightarrow E(K_a, \underbrace{[K_s, ID_A, ID_B, N_i]}_A \parallel \underbrace{E(K_b, [K_s, ID_A])}_B)$

$K_a$ : A's secret key  $\leftarrow$  A's Ke password PE hash applied

$K_b$ : B's secret key  $\leftarrow$  B's Ke password PE hash applied

$K_s$ : Session key  $\leftarrow$  Randomly created



III Now A will decrypt this message using  $K_A$  & will get  $(K_S, ID_A, ID_B, M_1)$  & encrypted message of B. A will forward the encrypted message of B to B.



$$E(K_B, [K_S, ID_A])$$

IV B decrypts message using  $K_B$  & understands A wants to communicate ( $ID_A$ ) for which  $K_S$  is session key.

Now B creates a random nonce  $N_2$  & encrypts using  $K_S$  & sends to A.



$$E(K_S, [N_2])$$

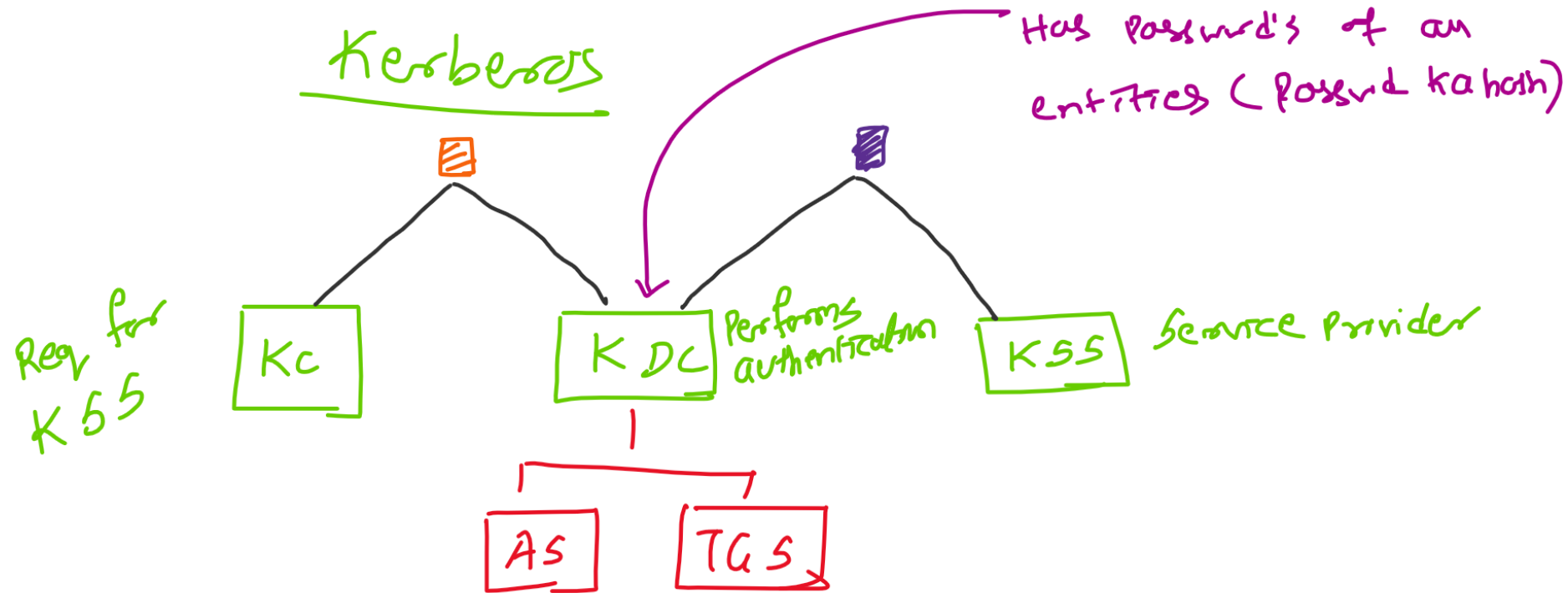
V A decrypts message using  $K_S$  & gets  $N_2$ .

Final Ack:  $N_2 + 1$  & encrypt using  $K_S$  & send to B



$$E(K_S, [N_2 + 1])$$

B now decrypts this using  $K_S$ .

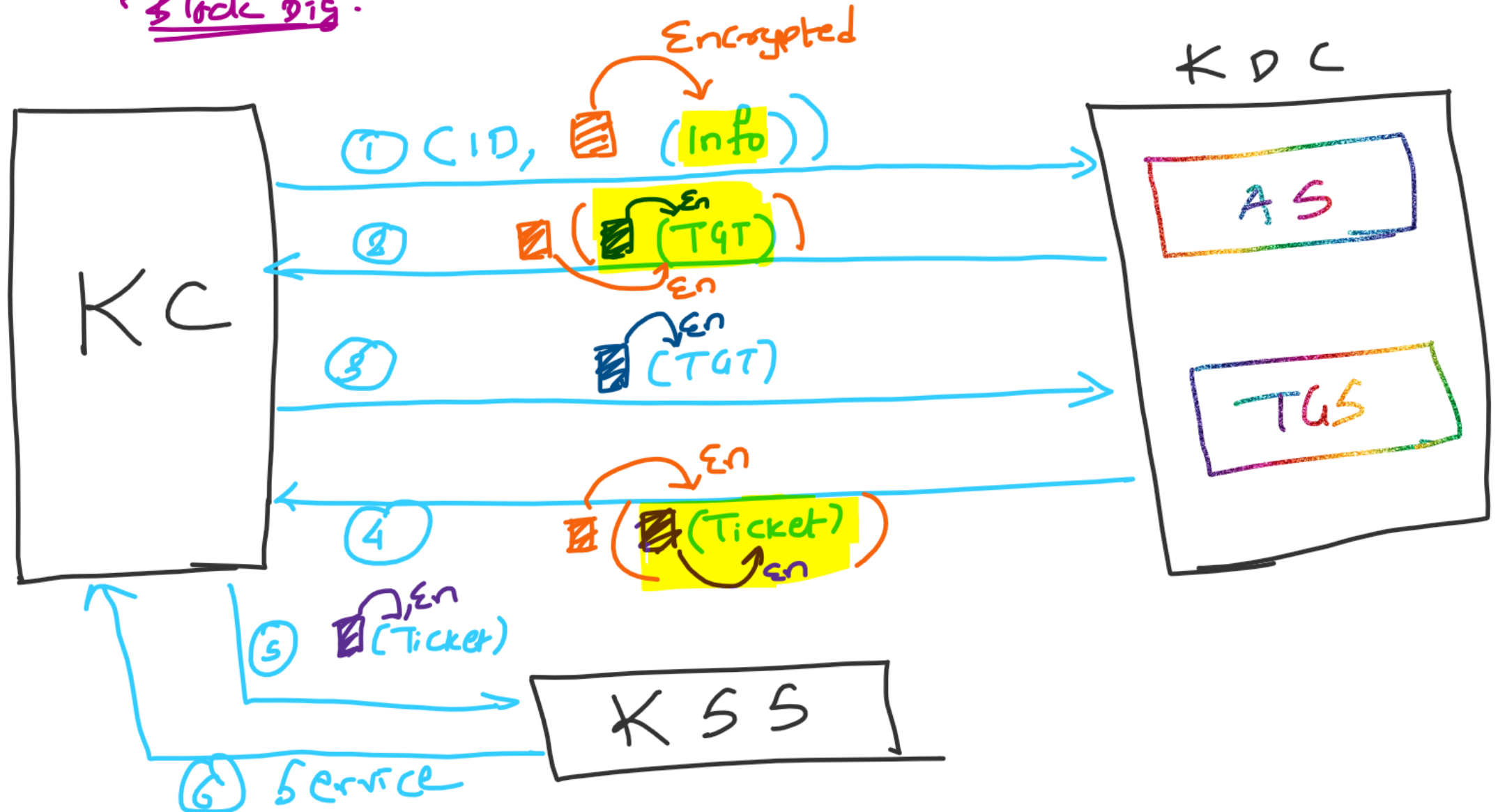


■ : Client's secret key [Client ke password pe hash apply]

■ : Session key [TGS ke password pe hash apply]

■ : Server's secret key [Server ke password pe hash apply]

Block dig:



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# Thank You

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