

A DETAILED DESCRIPTION OF CODE GENERATION (C-F07)

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Initial program fed to compiler:

```

int g;                // a global array

int dog(int x, y)     // two arguments
{
    int z;            // local variable

    z = x + y;        // expression using local vars

    return g*z;       // expression using local and global vars
}

int h[10];            // global array

void cat(int x[], y)  // two arguments one of which is an array
{
    int z[10];        // a local array

    z[8] = z[9];      // regular local array stuff
    h[8] = h[9];      // regular global array stuff
    z[7] = x[y];      // parm array on right hand side
    x[7] = z[y];      // parm array on left hand side
}                    // use default return

void main()
{
    int a[10];        // local array
    int b;            // local var

    dog(b, 999);      // call dog
    cat(a, 6);        // call cat
    b = dog(777,888); // get return value from call to dog
}

```

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Semantic Analysis

Besides doing the type checking of the tree during traversal this will annotate the nodes in the tree with the data necessary for code generation phase. FOFFSET is used to hold the local offset GOFFSET is used to hold the global offset.

```

// init goffset=0

int g;                // decl node is annotated with
                    //   location: 0, size: 1, referenceType: global
                    // update goffset=1

```

```

// enter scope of dog in symbol table
// foffset=-2

int dog(int x, y)
// decl node for x is annotated with
//   location: -2, size: 1, referenceType: parameter
// update local offset
// decl node for y is annotated with
//   location: -3, size: 1, referenceType: parameter
// update local offset
{
    int z;
// decl node for z is annotated with
//   location: -4, size: 1, referenceType: local
// update local offset
    z = x + y;
// annotate the id nodes with location/size/referenceType
via
    return g*z;
// lookup in symbol table
// annotate the id nodes with location/size/referenceType
via
// lookup in symbol table
}
// save size of frame (5) for dog in global symbol dog.
//   This is in foffset.
// exit scope of dog in symbol table

int h[10];
// decl node is for global variable annotated with
//   location: -2, size: 10+1 for array size,
referenceType: global
//   allocate 1 for size, then 10 for the array it self.
Location
//   is array itself
// update goffset=-12

// enter scope of cat in symbol table
// foffset=-2

void cat(int x[], y)
// decl node for x is annotated with
//   location: -2, size: 1, referenceType: parameter
//   (arrays are passed by reference and so size is 1.
//   see code generation section.)
// update local offset foffset=-3
// decl node for y is annotated with
//   location: -3, size: 1, referenceType: parameter
// update local offset foffset=-4
{
    int z[10];
// decl node for z is annotated with
//   location: -5, size: 10+1 for size, referenceType:
local
// update local offset foffset=-15

    z[8] = z[9];
// annotate the id nodes with location/size/referenceType
via
    h[8] = h[9];
// lookup in symbol table
// annotate the id nodes with location/size/referenceType
via
    z[7] = x[y];
// lookup in symbol table
// annotate the id nodes with location/size/referenceType
via

```

```

// lookup in symbol table
via    x[7] = z[y];    // annotate the id nodes with location/size/referenceType
// lookup in symbol table
}

// save size of frame (15) for cat in global symbol cat.
// This is the value of the foffset.
// exit scope of cat in symbol table

// enter scope of main in symbol table
// foffset=-2
void main()
{
    int a[10];    // decl node for a is annotated with
                  // location: -3, size: 10+1, referenceType: local
    int b;        // update local offset foffset=-13
                  // decl node for b is annotated with
                  // location: -13, size: 1, referenceType: local
                  // update local offset
via    dog(b, 999);    // annotate the id node with location/size/referenceType
// lookup in symbol table
// annotate the func node with size via symbol table
via    cat(a, 6);    // annotate the id node with location/size/referenceType
// lookup in symbol table
// annotate the func node with size via symbol table
    b = dog(777,888); // annotate id nodes with location/size/referenceType via
// lookup in symbol table
// annotate the func node with size via symbol table
}

// save size of frame (14) for main in global
// symbol dog. This is foffset. exit scope
// of cat in symbol table

```

===== Code Generation

Now we traverse the tree again using the information at each node to tell us what instructions to generate. The information we need will be:

1. the size of the global space so you know where to put the frame for main. Get this from goffset.
2. the symbol table, which at the end of the semantic phase contains all the global symbols. This is used to look up a function location at call time.
3. the annotated tree from the semantic phase. This has all of the sizes and locations of the data elements attached to each ID

You will need to keep track of a compile-time variable called toffset

which is the temporary offset during compilation.

First let's look at the memory layout by looking at the example code below and the configuration of memory using tm. We will stop the code at four spots and look at memory. This can be done using the breakpoint command in tm.

Consider this code:

```
int g;

int dog(int x, y)
{
    int z;
    z = 300;
                                <- stop point D at SECOND call to dog inside expression
                                <- stop point B in first call to dog

    z = x + y;
    return g*z;
}

int h[10];

void cat(int x[], y)
{
    int z[10];

    z[0] = 500;
    z[9] = 509;
                                <- stop point C

    z[8] = z[9];
    h[8] = h[9];
    z[7] = x[y];
    x[7] = z[y];
}

void main()
{
    int a[10];
    int b;

    a[0] = 100;
    a[9] = 109;
    b = 200;
    g = 300;
    h[0] = 400;
    h[9] = 409;
                                <- stop point A

    dog(b, 999);
    cat(a, 6);
    b = 555 * (666 + dog(777,888));
}
```

At stop point A:

Enter command: r

```
r[0]: 9999   r[1]: 9987   r[2]: 0       r[3]: 409
r[4]: 9      r[5]: 9988   r[6]: 0       r[7]: 148
```

Enter command: d 9999 -40

addr:	value	instr	that last assigned this loc
9999:	300	133	<- g <- reg 0 (globals) <
9998:	10	190	<- *h
9997:	400	140	<- h[0]
9996:	0	unused	
9995:	0	unused	
9994:	0	unused	
9993:	0	unused	
9992:	0	unused	
9991:	0	unused	
9990:	0	unused	
9989:	0	unused	
9988:	409	147	<- h[9]
9987:	9987	192	<- old frame pointer <- reg 1 <
9986:	195	113	<- return address
9985:	10	115	<- *a
9984:	100	122	<- a[0]
9983:	0	unused	
9982:	0	unused	
9981:	0	unused	
9980:	0	unused	
9979:	0	unused	
9978:	0	unused	
9977:	0	unused	
9976:	0	unused	
9975:	109	129	<- a[9]
9974:	200	131	<- b
9973:	9	142	<- was temp used in computing subscript 9 in previous
expression			
9972:	0	unused	
9971:	0	unused	
9970:	0	unused	
9969:	0	unused	
9968:	0	unused	
9967:	0	unused	
9966:	0	unused	
9965:	0	unused	
9964:	0	unused	

At stop point B:

Enter command: d 9999 -30

addr:	value	instr	that last assigned this loc
9999:	300	133	<- g <- reg 0 (globals) <
9998:	10	190	<- *h
9997:	400	140	<- h[0]
9996:	0	unused	
9995:	0	unused	
9994:	0	unused	
9993:	0	unused	

9992:	0	unused		
9991:	0	unused		
9990:	0	unused		
9989:	0	unused		
9988:	409	147	<- h[9]	<
9987:	9987	192	<- old frame pointer	<
9986:	195	113	<- return address	
9985:	10	115	<- *a	
9984:	100	122	<- a[0]	
9983:	0	unused		
9982:	0	unused		
9981:	0	unused		
9980:	0	unused		
9979:	0	unused		
9978:	0	unused		
9977:	0	unused		
9976:	0	unused		
9975:	109	129	<- a[9]	
9974:	200	131	<- b	<
9973:	9987	148	<- old frame pointer	<
9972:	156	30	<- return address	
9971:	200	150	<- parm 1	
9970:	999	152	<- parm 2	
9969:	300	32	<- z	<
9968:	0	unused		
9967:	0	unused		
9966:	0	unused		
9965:	0	unused		
9964:	0	unused		
9963:	0	unused		

At stop point C:

Enter command: r

```
r[0]: 9999    r[1]: 9973    r[2]: 359700    r[3]: 509
r[4]: 9       r[5]: 9959    r[6]: 0       r[7]: 69
```

```
Enter command: d 9999 -50
```

addr:	value	instr	that last assigned this loc
9999:	300	133	<- g
9998:	10	190	<- *h
9997:	400	140	<- h[0]
9996:	0	unused	
9995:	0	unused	
9994:	0	unused	
9993:	0	unused	
9992:	0	unused	
9991:	0	unused	
9990:	0	unused	
9989:	0	unused	
9988:	409	147	<- h[9]
9987:	9987	192	<- old frame pointer
9986:	195	113	<- return address
9985:	10	115	<- *a
9984:	100	122	<- a[0]
9983:	0	unused	

9982:	0	unused		
9981:	0	unused		
9980:	0	unused		
9979:	0	unused		frame for
9978:	0	unused		main
9977:	0	unused		
9976:	0	unused		
9975:	109	129	<- a[9]	
9974:	200	131	<- b	<
9973:	9987	157	<- old frame pointer <- reg 1	<
9972:	165	52	<- return address	
9971:	9984	159	<- parm 1 ptr to array x	
9970:	6	161	<- parm 2 y	
9969:	10	54	<- *z	
9968:	500	61	<- z[0]	
9967:	0	unused		
9966:	0	unused		frame for
9965:	0	unused		cat
9964:	0	unused		
9963:	0	unused		
9962:	0	unused		
9961:	0	unused		
9960:	0	unused		
9959:	509	68	<- z[9]	<
9958:	9	63	<- temp	
9957:	0	unused		

At stop point D

Enter command: r

r[0]: 9999	r[1]: 9971	r[2]: 0	r[3]: 300
r[4]: 7	r[5]: 9977	r[6]: 0	r[7]: 33

Enter command: d 9999 -60

addr:	value	instr that	last assigned this loc	
9999:	300	133	<- g <- reg 0 (globals)	<
9998:	10	190	<- *h	
9997:	400	140	<- h[0]	
9996:	0	unused		
9995:	0	unused		
9994:	0	unused		
9993:	0	unused		globals
9992:	0	unused		
9991:	0	unused		
9990:	0	unused		
9989:	409	88		
9988:	409	147		<
9987:	9987	192	<- old frame pointer	<
9986:	195	113	<- return address	
9985:	10	115	<- *a	
9984:	100	122	<- a[0]	
9983:	0	unused		
9982:	0	unused		
9981:	0	unused		
9980:	0	unused		frame for
9979:	0	unused		main

```

9978:    0    unused
9977:    0    108
9976:    0    unused
9975:  109   129    <- a[9]
9974:  200   131    <- b
9973:  555   167    <- temp!
9972:  666   169    <- temp!
9971: 9987   170    <- old frame pointer <- reg 1 <
9970:  178    30    <- return address (see below) |
9969:  777   172    <- parm 1                      dog frame
9968:  888   174    <- parm 2                      |
9967:  300    32    <- z                          <
9966:    0    unused    <- temp!
9965:    0    unused
9964:    0    unused
9963:    0    unused
9962:    0    unused
9961:    0    98      <- junk left from previous call to cat
9960:  509    78      <- junk left from previous call to cat
9959:  509    68      <- junk left from previous call to cat
9958:    7   100     <- junk left from previous call to cat
9957:    0    unused
9956:    0    unused

```

IMPORTANT: notice that pending temporaries for the expression
 $555 * (666 + \text{dog}(777, 888))$
are stored at 9972 and 9973! This is what the toffset is for.

the temporary at 9966 was generated in the expression
 $z = x + y$
in dog where x was 777 and y was 888.

As you go through the tree generating code you need to keep track of toffset.

1. It is set to the end of the current frame when you enter a function.
2. As expressions push and pop stuff off the temporary stack after the current frame the toffset gets incremented when you save a value there and decremented when that value is no longer needed. Remember the stack you are managing is a RUNTIME stack not a compile time stack. You are using toffset to predict where the temporaries go! First study how a simple expression is computed and then just use toffset in your emit commands as an offset from register 1. Here is the code for $x = y + z$ from my compiler:

```

38:    LD  3,-2(1)    Load variable x          <- LHS code
39:    ST  3,-5(1)    Save left side in temporary <- after LHS code
40:    LD  3,-3(1)    Load variable y          <- RHS code
41:    LD  4,-5(1)    Load left into acl from temporary <- after RHS code
42:    ADD 3,4,3      Op +                      <- do op
43:    ST  3,-4(1)    Store variable z          <- assignment

```

the size of the dog frame was 5. So -5(1) means the first spot after

the frame (because 0(1) is the first spot in the frame).

3. When you enter a function, as in the case of entering dog for a second time, as above, note that we leave some temps between the main frame and the dog frame. This means new frames are always added as if they are on the temp stack of the previous frame! That is, frames are added after the current temp on the stack. This position is toffset. So you should use toffset while creating the ghost frame.

4. By carefully incrementing and decrementing toffset you should find that you only need to set toffset to the position right after the frame you just created when you enter a function (as you are generating the code). Expressions should always end with toffset back at the space right after the current frame. But the important idea is that this will happen. So make sure you always pair increment with decrement.

5. toffset is a compile time variable and changes as you pass through the code lexically.

6. Remember as you generate code (compile time) you are traversing the tree which is essentially reading the program from the beginning to the end (lexical order). When you run the code at run time, then you move through the code in execution order.

=====

TM code for this example

```
* C- compiler version C-F07
* Author: Robert B. Heckendorn
* Backend adapted from work by Jorge Williams (2001)
* File compiled: z.c-
* Nov 29, 2007
* BEGIN function input
1:      ST  3,-1(1)      Store return address
2:      IN  2,2,2        Grab int input
3:      LD  3,-1(1)      Load return address
4:      LD  1,0(1)       Adjust fp
5:      LDA 7,0(3)       Return
* END of function input
* BEGIN function output
6:      ST  3,-1(1)      Store return address
7:      LD  3,-2(1)      Load parameter
8:      OUT 3,3,3        Output integer
9:      LDC 2,0(6)       Set return to 0
10:     LD  3,-1(1)      Load return address
11:     LD  1,0(1)       Adjust fp
12:     LDA 7,0(3)       Return
* END of function output
* BEGIN function inputb
13:     ST  3,-1(1)      Store return address
14:     INB 2,2,2        Grab bool input
15:     LD  3,-1(1)      Load return address
16:     LD  1,0(1)       Adjust fp
17:     LDA 7,0(3)       Return
* END of function inputb
```

```
* BEGIN function outputb
18:    ST  3,-1(1)    Store return address
19:    LD  3,-2(1)    Load parameter
20:    OUTB 3,3,3      Output bool
21:    LDC 2,0(6)     Set return to 0
22:    LD  3,-1(1)    Load return address
23:    LD  1,0(1)     Adjust fp
24:    LDA 7,0(3)     Return
* END of function outputb
* BEGIN function outnl
25:    ST  3,-1(1)    Store return address
26:    OUTNL 3,3,3     Output a newline
27:    LD  3,-1(1)    Load return address
28:    LD  1,0(1)     Adjust fp
29:    LDA 7,0(3)     Return
* END of function outnl
* BEGIN function dog
30:    ST  3,-1(1)    Store return address.
* BEGIN compound statement
* EXPRESSION STMT
31:    LD  3,-2(1)    Load variable x
32:    ST  3,-5(1)    Save left side
33:    LD  3,-3(1)    Load variable y
34:    LD  4,-5(1)    Load left into acl
35:    ADD 3,4,3      Op +
36:    ST  3,-4(1)    Store variable z
* RETURN
37:    LD  3,0(0)     Load variable g
38:    ST  3,-5(1)    Save left side
39:    LD  3,-4(1)    Load variable z
40:    LD  4,-5(1)    Load left into acl
41:    MUL 3,4,3      Op *
42:    LDA 2,0(3)     Copy result to rt register
43:    LD  3,-1(1)    Load return address
44:    LD  1,0(1)     Adjust fp
45:    LDA 7,0(3)     Return
* END compound statement
* Add standard closing in case there is no return statement
46:    LDC 2,0(6)     Set return value to 0
47:    LD  3,-1(1)    Load return address
48:    LD  1,0(1)     Adjust fp
49:    LDA 7,0(3)     Return
* END of function dog
* BEGIN function cat
50:    ST  3,-1(1)    Store return address.
* BEGIN compound statement
51:    LDC 3,10(6)     load size of array z
52:    ST  3,-4(1)     save size of array z
* EXPRESSION STMT
53:    LDC 3,8(6)      Load constant
54:    ST  3,-15(1)    Save index
55:    LDC 3,9(6)      Load constant
56:    LDA 4,-5(1)     Load address of base of array z
57:    SUB 3,4,3        Compute offset of value
58:    LD  3,0(3)      Load the value
59:    LD  4,-15(1)     Restore index
60:    LDA 5,-5(1)     Load address of base of array z
61:    SUB 5,5,4        Compute offset of value
```

```

62:      ST  3,0(5)      Store variable z
* EXPRESSION STMT
63:      LDC  3,8(6)      Load constant
64:      ST  3,-15(1)     Save index
65:      LDC  3,9(6)      Load constant
66:      LDA  4,-2(0)     Load address of base of array h
67:      SUB  3,4,3       Compute offset of value
68:      LD   3,0(3)      Load the value
69:      LD   4,-15(1)     Restore index
70:      LDA  5,-2(0)     Load address of base of array h
71:      SUB  5,5,4       Compute offset of value
72:      ST  3,0(5)      Store variable h
* EXPRESSION STMT
73:      LDC  3,7(6)      Load constant
74:      ST  3,-15(1)     Save index
75:      LD   3,-3(1)     Load variable y
76:      LD   4,-2(1)     Load address of base of array x
77:      SUB  3,4,3       Compute offset of value
78:      LD   3,0(3)      Load the value
79:      LD   4,-15(1)     Restore index
80:      LDA  5,-5(1)     Load address of base of array z
81:      SUB  5,5,4       Compute offset of value
82:      ST  3,0(5)      Store variable z
* EXPRESSION STMT
83:      LDC  3,7(6)      Load constant
84:      ST  3,-15(1)     Save index
85:      LD   3,-3(1)     Load variable y
86:      LDA  4,-5(1)     Load address of base of array z
87:      SUB  3,4,3       Compute offset of value
88:      LD   3,0(3)      Load the value
89:      LD   4,-15(1)     Restore index
90:      LD   5,-2(1)     Load address of base of array x
91:      SUB  5,5,4       Compute offset of value
92:      ST  3,0(5)      Store variable x
* END compound statement
* Add standard closing in case there is no return statement
93:      LDC  2,0(6)      Set return value to 0
94:      LD   3,-1(1)     Load return address
95:      LD   1,0(1)      Adjust fp
96:      LDA  7,0(3)      Return
* END of function cat
* BEGIN function main
97:      ST  3,-1(1)      Store return address.
* BEGIN compound statement
98:      LDC  3,10(6)     load size of array a
99:      ST  3,-2(1)      save size of array a
* EXPRESSION STMT
100:     ST  1,-14(1)      Store old fp in ghost frame
101:     LD   3,-13(1)     Load variable b
102:     ST  3,-16(1)      Store parameter
103:     LDC  3,999(6)     Load constant
104:     ST  3,-17(1)      Store parameter
105:     LDA  1,-14(1)     Load address of new frame
106:     LDA  3,1(7)       Return address in ac
107:     LDA  7,-78(7)     call dog
108:     LDA  3,0(2)       Save the result in ac
* EXPRESSION STMT
109:     ST  1,-14(1)      Store old fp in ghost frame

```

```
110:    LDA    3,-3(1)    Load address of base of array a
111:    ST     3,-16(1)   Store parameter
112:    LDC     3,6(6)     Load constant
113:    ST     3,-17(1)   Store parameter
114:    LDA     1,-14(1)  Load address of new frame
115:    LDA     3,1(7)    Return address in ac
116:    LDA     7,-67(7)  call cat
117:    LDA     3,0(2)    Save the result in ac
* EXPRESSION STMT
118:    ST     1,-14(1)  Store old fp in ghost frame
119:    LDC     3,777(6)  Load constant
120:    ST     3,-16(1)  Store parameter
121:    LDC     3,888(6)  Load constant
122:    ST     3,-17(1)  Store parameter
123:    LDA     1,-14(1)  Load address of new frame
124:    LDA     3,1(7)    Return address in ac
125:    LDA     7,-96(7)  call dog
126:    LDA     3,0(2)    Save the result in ac
127:    ST     3,-13(1)  Store variable b
* END compound statement
* Add standard closing in case there is no return statement
128:    LDC     2,0(6)    Set return value to 0
129:    LD      3,-1(1)   Load return address
130:    LD      1,0(1)    Adjust fp
131:    LDA     7,0(3)    Return
* END of function main
0:    LDA     7,131(7)  Jump to init
* BEGIN Init
132:    LD      0,0(0)    Set the global pointer
* BEGIN init of global array sizes
133:    LDC     3,10(6)   load size of array h
134:    ST     3,-1(0)    save size of array h
* END init of global array sizes
135:    LDA     1,-12(0)  set first frame at end of globals
136:    ST     1,0(1)     store old fp (point to self)
137:    LDA     3,1(7)    Return address in ac
138:    LDA     7,-42(7)  Jump to main
139:    HALT    0,0,0     DONE!
* END Init
```