1. (a) [3 marks] Write each of the following rational functions f(x) in form of a formal power series $f(x) = \sum_{i \geq 0} a_i x^i$.

$$\frac{1}{(1-2x)^3}, \qquad \frac{x^3}{(1+4x^2)^5}$$

Solution.

$$\frac{1}{(1-2x)^3} = \sum_{n\geq 0} \binom{n+2}{n} (2x)^n = \sum_{n\geq 0} \binom{n+2}{n} 2^n x^n = \sum_{n\geq 0} \binom{n+2}{2} 2^n x^n.$$

$$\frac{x^3}{(1+4x^2)^5} = x^3 \sum_{n\geq 0} \binom{n+4}{n} (-4x^2)^n = \sum_{n\geq 0} \binom{n+4}{n} (-4)^n x^{2n+3} = \sum_{n\geq 0} \binom{n+4}{4} (-4)^n x^{2n+3}.$$

(b) [1 mark] Determine the coefficient $[x^5] \frac{x^2}{(1+x)^7}$. Solution.

$$[x^5] \frac{x^2}{(1+x)^7} = [x^5] x^2 \sum_{i>0} \binom{i+6}{i} (-x)^i = [x^3] \sum_{i>0} \binom{i+6}{i} (-1)^i x^i = -\binom{9}{3}.$$

(c) [3 marks] Write the following formal power series in closed form:

$$\sum_{i\geq 0} \left(-\frac{1}{2}\right)^i x^{3i}, \qquad \sum_{i\geq 0} \binom{i+2}{i} 5^i x^i.$$

Solution.

$$\sum_{i \ge 0} \left(-\frac{1}{2} \right)^i x^{3i} = \sum_{i \ge 0} \left(\frac{-x^3}{2} \right)^i = \frac{1}{1 + \frac{x^3}{2}} = \frac{2}{2 + x^3}.$$

$$\sum_{i>0} \binom{i+2}{i} 5^i x^i = \sum_{i>0} \binom{i+2}{i} (5x)^i = \frac{1}{(1-5x)^3}.$$

2. (a) [2 marks] Use the Binomial Theorem to prove that $\sum_{i=0}^{n} {n \choose i} = 2^n$. Solution.

By the Binomial Theorem,

$$(1+x)^n = \sum_{i=0}^n \binom{n}{i} x^i.$$

Letting x = 1, we get

$$2^n = \sum_{i=0}^n \binom{n}{i}.$$

(b) [3 marks] Prove that

$$\sum_{r=0}^{n} \sum_{s=0}^{r} \binom{n}{r} \binom{r}{s} = 3^{n}.$$

(Hint: Use the Binomial Theorem to expand $(1+2)^n$.)

Solution.

By part (a) and the Binomial Theorem,

$$3^{n} = (1+2)^{n} = \sum_{r=0}^{n} {n \choose r} 2^{r} = \sum_{r=0}^{n} {n \choose r} \left(\sum_{s=0}^{r} {r \choose s} \right) = \sum_{r=0}^{n} \sum_{s=0}^{r} {n \choose r} {r \choose s}$$

as required.

- 3. An even composition of n is a composition of n into an even number of parts such that each part is a positive even number. For example, the even compositions of 8 are (2,6), (4,4), (6,2), (2,2,2,2).
 - (a) [4 marks] Determine the generating function for the number of even compositions of n having 2k parts. Use this generating function to determine the number of 2k-part even compositions of n.

Solution. Let $E = \{2, 4, 6, \ldots\}$. By definition,

$$\Phi_E(x) = x^2 + x^4 + x^6 + \dots = \frac{x^2}{1 - x^2}.$$

The set of all even compositions with 2k parts is $S = E^{2k}$ (the Cartesian product of E with itself 2k times). By the Product Lemma,

$$\Phi_S(x) = (\Phi_E(x))^{2k} = \left(\frac{x^2}{1-x^2}\right)^{2k} = \frac{x^{4k}}{(1-x^2)^{2k}}$$

The number of such even compositions of n is

$$[x^n]\Phi_S(x) = [x^n] \frac{x^{4k}}{(1-x^2)^{2k}} = [x^{n-4k}] \sum_{i>0} {2k+i-1 \choose i} x^{2i} = {n \choose \frac{n}{2}-1 \choose \frac{n}{2}-2k}.$$

(Remark: this binomial coefficient is zero unless n is even and $n \geq 4k$.)

(b) [3 marks] Determine the generating function for the the number of all even compositions of n.

Solution. Using part (a), the set of all even compositions is

$$T = E^0 \cup E^2 \cup E^4 \cup E^6 \cup \cdots$$

By the Sum and Product Lemmas,

$$\Phi_T(x) = (\Phi_E(x))^0 \cup (\Phi_E(x))^2 \cup (\Phi_E(x))^4 \cup (\Phi_E(x))^6 \cup \cdots
= \left(\frac{x^2}{1-x^2}\right)^0 + \left(\frac{x^2}{1-x^2}\right)^2 + \left(\frac{x^2}{1-x^2}\right)^4 + \left(\frac{x^2}{1-x^2}\right)^6 + \cdots
= \frac{1}{1-\left(\frac{x^2}{1-x^2}\right)^2} = \frac{1-2x^2+x^4}{1-2x^2} = 1 + \frac{x^4}{1-2x^2}.$$

(c) [3 marks] Use your answer in part (b) to determine a_n , the number of even compositions of n.

Solution.

$$a_{n} = [x^{n}]\Phi_{T}(x)$$

$$= [x^{n}]1 + x^{4} \sum_{i \geq 0} 2^{i}x^{2i}$$

$$= [x^{n}]1 + \sum_{i \geq 0} 2^{i}x^{2i+4}$$

$$= \begin{cases} 1 & \text{if } n = 0, \\ 2^{\frac{n-4}{2}} & \text{if } n \text{ is even and } n > 2, \\ 0 & \text{otherwise.} \end{cases}$$

- 4. For each of the following sets, write down a decomposition that uniquely creates the elements of that set.
 - (a) [2 marks] The $\{0,1\}$ -strings that have no blocks of 1s with length 3, and no substrings of 0s of length 2.

Solution.

The decomposition for these strings is

$$(\{1\}^* \setminus \{111\}) \quad (\{0\}(\{1\}\{1\}^* \setminus \{111\}))^* \quad \{\varepsilon, 0\}.$$

(b) [2 marks] The set of $\{0,1\}$ -strings with no occurrence of the substring 010. Solution.

The decomposition for this set of strings is

$$\{1\}^* \quad (\{0\}\{0\}^*\{11\}\{1\}^*)^* \quad (\{0\}^* \cup \{0\}\{0\}^*\{1\}).$$

5. (a) [3 marks] Find the generating function $\Phi_{S_1}(x)$ for the following set of strings:

$$S_1 = \{0\}^* (\{1\}\{111\}^*\{00\}\{000\}^*)^* \{1\}^*.$$

Solution.

By the Product Lemma, the generating function for S_1 is

$$\Phi_{S_1}(x) = \left(\frac{1}{1-x}\right) \left(\frac{1}{1-\left(\frac{x}{1-x^3}\right)\left(\frac{x^2}{1-x^3}\right)}\right) \left(\frac{1}{1-x}\right) \\
= \frac{(1+x+x^2)^2}{1-3x^3+x^6}.$$

(b) [4 marks] Find the generating function $\Phi_{S_2}(x)$ for the following set of strings:

$$S_2 = \bigcup_{k=0}^{\infty} \{0\}^k (\{1\}\{111\}^*\{00\}\{000\}^*)^* \{1\}^k.$$

Solution.

Using the Product and Sum Lemmas, we get that

$$\Phi_{S_2}(x) = \sum_{k=0}^{\infty} x^k \left(\frac{1}{1 - \left(\frac{x}{1 - x^3}\right) \left(\frac{x^2}{1 - x^3}\right)} \right) x^k
= \frac{(1 - x^3)^2}{1 - 3x^3 + x^6} \sum_{k=0}^{\infty} x^{2k}
= \frac{(1 - x^3)^2}{1 - 3x^3 + x^6} \left(\frac{1}{1 - x^2} \right) = \frac{1 + x + x^2 - x^3 - x^4 - x^5}{1 + x - 3x^3 - 3x^4 + x^6 + x^7}.$$

(c) [1 mark] Find a string that is in S_1 but not in S_2 .

Solution.

The string 1 is in S_1 but is not in S_2 .

(d) [1 mark] Let $a_n = [x^n]\Phi_{S_1}(x)$ and $b_n = [x^n]\Phi_{S_2}(x)$. Without evaluating a_n or b_n , show that $a_n \geq b_n$ for all $n \geq 0$.

Solution.

Clearly every string of S_2 is also in S_1 . Therefore $a_n \geq b_n$.

- 6. Let A be the set of binary strings a such that in each substring of three consecutive positions in a, there is at least one 1 and at least one 0. Let a_n be the number of strings of length n in A.
 - (a) [2 marks] For n = 1, 2, 3, 4, write down all the strings of length n in A. Solution.

$$\begin{array}{ll} n=1 & 0,1 \\ n=2 & 00,01,10,11 \\ n=3 & 001,010,100,011,101,110 \\ n=4 & 0010,0011,0100,0101,0110,1001,1010,1011,1100,1101 \end{array}$$

(b) [3 marks] Prove that the generating function for a_n is given by

$$\Phi_A(x) = \frac{1 + 2x + 3x^2 + 2x^3 + x^4}{1 - x^2 - 2x^3 - x^4}.$$

Solution.

A decomposition for the strings of A is

$$A = \{\varepsilon, 1, 11\}\{01, 011, 001, 0011\}^* \{\varepsilon, 0, 00\}.$$

By the Sum and Product Lemmas, the generating function for A is

$$\Phi_A(x) = (1+x+x^2) \left(\frac{1}{1-x^2-2x^3-x^4}\right) (1+x+x^2)
= \frac{1+2x+3x^2+2x^3+x^4}{1-x^2-2x^3-x^4}$$

as required.

(c) [2 marks] From part (b), find a recurrence relation for the numbers a_n , together with an appropriate set of initial conditions.

Solution.

By Theorem 1.26, a linear homogeneous recurrence relation for a_n 's is

$$a_n - a_{n-2} - 2a_{n-3} - a_{n-4} = 0 \quad \forall n \ge 5.$$

From part (a), the initial conditions that uniquely determine the sequence $\{a_n\}$ are

$$a_0 = 1$$
, $a_1 = 2$, $a_2 = 4$, $a_3 = 6$, $a_4 = 10$.

7. (a) [4 marks] Solve the linear, homogeneous recurrence equation

$$a_n + a_{n-1} - 8a_{n-2} - 12a_{n-3} = 0 \quad \forall \quad n \ge 3$$

with initial conditions $a_0 = 4$, $a_1 = 8$, $a_2 = 2$.

Solution.

The characteristic polynomial of this recurrence equation is

$$x^{3} + x^{2} - 8x - 12 = (x+2)^{2}(x-3).$$

By Theorem 1.42, the general solution of this recurrence equation is

$$a_n = (A + Bn)(-2)^n + C3^n$$
.

Setting n equal to 0, 1, 2, we get the following 3 linear equations in A, B, C:

$$A + C = 4,$$

 $-2A - 2B + 3C = 8,$
 $4A + 8B + 9C = 2.$

The solution to this system of equations is (A, B, C) = (2, -3, 2). Therefore, the solution for a_n is

$$a_n = (2 - 3n)(-2)^n + 2(3^n) \quad \forall n \ge 0.$$

(b) [1 marks] What is the asymptotic value of a_n ?

Solution.

From part (a) we get that the asymptotic value of a_n is $a_n \sim 2(3^n)$.

(c) [3 marks] Find a particular solution of the nonhomogeneous linear recurrence equation

$$a_n - a_{n-1} - 5a_{n-2} - 3a_{n-3} = \frac{15}{8}(-2)^n \quad \forall \quad n \ge 3.$$

Solution.

Substitute $a_n = \alpha(-2)^n$ in this recurrence equation and solve for α .

$$\frac{15}{8}(-2)^n = \alpha(-2)^n - \alpha(-2)^{n-1} - 5\alpha(-2)^{n-2} - 3\alpha(-2)^{n-3}$$
$$= \alpha(-2)^{n-3}(-8 - 4 + 10 - 3)$$
$$= -5\alpha(-2)^{n-3} = \frac{5}{8}\alpha(-2)^n.$$

Therefore, $\alpha = 3$, and we obtain the particular solution $a_n = 3(-2)^n$.

8. (a) [2 marks] Let G be a 3-regular graph with 10 vertices. (This means every vertex has degree 3.) How many edges does G have? Justify your answer.

Solution.

Let the vertices of G be v_1, v_2, \ldots, v_{10} and let the number of edges be q. By the Handshake Theorem,

$$2q = \sum_{i=1}^{10} \deg(v_i) = \sum_{i=1}^{10} 3 = 30.$$

Therefore, the number of edges is q = 15.

(b) [3 marks] Let G be a connected graph with 27 edges and exactly one cycle. How many vertices does G have? Justify your answer.

Solution.

Let e be any edge of the cycle in G. Then G - e is connected (because e is not a bridge) and has no cycle (removing edge e destroys the unique cycle in G). Therefore, G - e is a tree with 26 edges. Therefore, by Theorem 2.29, G - e has 27 vertices, as does G.

(c) [3 marks] Is there a bipartite graph with 9 vertices and 23 edges? Explain.

Solution.

Let G be a bipartite graph with bipartition A, B, where A and B have a and b vertices respectively. If every vertex of A is adjacent to every vertex of B, then G has ab edges. Clearly, no bipartite graph with bipartition A, B can have more than ab edges.

If G is a bipartite graph with 9 vertices, then (a, b) must be one of (0, 9), (1, 8), (2, 7), (3, 6) or (4, 5). In each case, $ab \le 20$. Therefore, no bipartite graph on 9 vertices has 23 edges.

- 9. Let Q_n denote the *n*-cube graph: the vertex set of Q_n consists of all binary sequences (i.e. binary strings) of length n and the edge set of Q_n consists of all pairs of binary sequences which differ in exactly one position.
 - (a) [3 marks] Determine the number of vertices and edges in Q_n . Solution.

Since there are 2^n binary strings of length n, Q_n has 2^n vertices.

Since each binary string has n positions, it is adjacent to n binary strings; that is, each vertex has degree n. By the Handshake Theorem,

$$2q = \sum_{v \in V(Q_n)} \deg(v) = \sum_{v \in V(Q_n)} n = n(2^n).$$

Therefore, the number of edges in Q_n is $q = n(2^{n-1})$.

- (b) [2 marks] Show that Q_n is bipartite.
 - Solution.

Let E be the set of all binary strings that have an even number of 1's, and let O be the set that have an odd number of ones. Let $S = s_1 \dots s_n$ and $T = t_1 \dots t_n$ be adjacent binary strings. They must be identical in n-1 positions and different in 1 position, say position i. Without loss of generality, we may assume the $s_i = 1$ and $t_i = 0$. Since S has one more 1 than T has, one of S and T is in E and the other is in O. Since every edge joins a vertex in E to a vertex in O, O, O is a bipartite graph with bipartition O.

(c) [3 marks] Prove that Q_n has no bridges.

Solution.

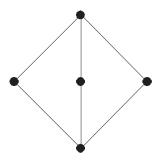
For n = 1, Q_n is the complete graph on 2 vertices, which has precisely one edge. This one edge is a bridge. We now prove that for $n \ge 2$, Q_n has no bridges.

Let S and T be any pair of adjacent vertices as described in part (b). Let $U = u_1 \dots u_n$ be another vertex adjacent to T. Assume U differs from T in position $j \neq i$. Then $u_j = 1 - t_j = 1 - s_j$ and $u_i = t_i = 0$. Let $V = v_1 \dots v_n$ be the string that has

$$v_i = s_i = 1, \quad v_j = u_j = 1 - t_j = 1 - s_j$$

and is the same as S, T and U in all other positions. Vertex V is adjacent to both U and S. Hence S, T, U, V, S is a cycle of length 4 in Q_n . Since edge $\{S, T\}$ is in a cycle of length 4, it is not a bridge.

(d) [4 marks] (Continued from page 10.) Prove that that the graph Q_n does not contain a subgraph isomorphic to the graph drawn below:



Solution.

Let the top vertex be $S=s_1\ldots s_n$, let the middle 3 vertices be $A=a_1\ldots a_n,\ B=b-1\ldots b_n$ and $C=c_1\ldots c_n$, respectively, and let the bottom vertex be $W=w_1\ldots w_n$. Let S differ from vertex A in position i and be the same as A in all other positions. Similarly, let S differ from vertices B and C in positions j and k respectively. Then $(a_i,a_j,a_k)=(1-s_i,s_j,s_k),\ (b_i,b_j,b_k)=(s_i,1-s_j,s_k)$ and $(c_i,c_j,c_k)=(s_i,s_j,1-s_k)$. Since vertex W is adjacent to both A and B, and $W\neq S$, $(w_i,w_j,w_k)=(1-s_i,1-s_j,s_k)$ and is identical to vertices S,A,B,C in all other positions. (Note that this solution for W is uniquely determined by S,A and B.) Now W differs from C in each of positions i,j,k. Hence, W is not adjacent to vertex C. Therefore, Q_n does not contain a subgraph isomorphic to the above graph.