- 1. Proof required.
- 2.

$$\left(\frac{x^{10}}{1-x^{10}}\right)\left(\frac{x^{20}}{1-x^{20}}\right)\left(\frac{x^{50}}{1-x^{50}}\right)$$

- 3. (a)  $\{0\}^*(\{1\}\{1\}^*\{0\}\{0\}^* \setminus \{1\}\{11\}^*\{0\}\{00\}^*)^*\{1\}^*$ 
  - (b)  $a_n = a_{n-1} + a_{n-2}, \quad n \ge 2.$

$$a_0 = 1, a_1 = 2.$$

- 4.  $b_n = (2-n)2^n$ .
- 5. (a) a connected graph with no cycles.
  - (b) Let  $n_i$  be the number of vertices of degree i. Then using sum of degrees = 2q (Handshake theorem), and also q = p 1 for trees, we can see that

$$n_1 - n_3 - 2n_4 - 3n_5 - \dots = -2.$$

This is true for all trees. Substituting  $n_1 = 7$  and  $n_4 = 2$  we get

$$7 - n_3 - 2 \times 4 - 3n_5 - \dots = -2.$$

Thus

$$1 = n_3 + 3n_5 + 4n_6 + \cdots$$

Since the  $n_i$  are nonnegative integers, the only possibility is  $n_3 = 1$  (and  $n_5 = n_6 = \cdots = 0$ ).

- (c) Should be easy.
- 6. (a) The vertices at levels are as follows (if these are correct, you probably have the right tree):

level 0: a

level 1: b e s

- 2: cifort
- 3: g j d p n u
- 4: h k q
- 5: m
- (b) The graph is not bipartite because there is an edge between t wo vertices at the same level (gu). Tracing back up the tree we find the cycle dfguqmhd.
- 7. Use the handshake theorem: sum of vertex degrees is 5p = 2q.

Also the analogue for faces ('face-shaking'): sum of face degrees is 3s = 2q (assuming s denotes the number of faces).

So 
$$p = 2q/5$$
 and  $s = 2q/3$ .

Plug these into Euler's theorem: since the embedding is connected,

$$p - q + s = 2$$

- 8. (a) is planar. (Draw it.)
  - (b) is not planar. It contains an edge subdivision of  $K_{3,3}$ .
- 9. Requires a proof. Similar to proof of 6 colour theorem.
- 10. (a) state Hall's theorem (Bipartite graph with bipartition A, B has matching saturating all vertices in A iff  $|N(D)| \ge |D|$  for all  $D \subseteq A$ .)
  - (b) Let  $D \subseteq A$ . If  $D = \emptyset$  then  $|N(D)| \ge |D|$  immediately. Next, every vertex has at least 10 neighbours. So every nonempty set D of vertices has  $|N(D)| \ge 10$ . Thus  $|N(D)| \ge |D|$  if  $|D| \le 10$ . Now suppose  $|D| \ge 11$  for  $D \subseteq A$ . We show that |N(D)| = 24. If not, then there is a vertex  $b \in B$  such that  $b \notin N(D)$ . Then b is adjacent to no vertex in D. That is, N(b) is disjoint from D. But the graph is bipartite and b has degree at least 10, so at least 10 vertices in A are not in D (i.e. the vertices in N(b)). Since |A| = 20 this means |D| must be at most 10. (Contradiction). So |N(D)| = 24 if  $|D| \ge 10$ . Since  $|D| \le 20$  this implies  $|N(D)| \ge |D|$ , as required.

Finally, since  $|N(D) \ge |D|$  for all  $D \subseteq A$ , the graph has a matching saturating all vertices in A by Hall's theorem.

11. (a) Vertices added in the following order:

To X: 3,5

To Y: 2,4,6,10

To X: 1,9,7,15

To Y: 14, 18, 20

To X: 13, 17,19

To Y: 8, 12, 16

Stop: 12 is unsaturated.

Augmenting path is 5, 10, 15, 14 13, 12.

New matching has edges 1 2, 6 7, 4 9, 5 10, 15 14, 13 12, 8 11, 17 18, 19 20.

(b) Minimum cover 2,4,7.