

**DUE: NOON Friday 21 October 2011** in the drop boxes opposite the Math Tutorial Centre MC 4067 or next to the St. Jerome's library for the St. Jerome's section.

1. For each of the following sets  $A$ , determine whether  $A^*$  is uniquely created. Prove your assertion in each case. (Hint: to prove a set  $A^*$  is uniquely created, try using induction on the length.)

(a)  $A = \{1, 0010, 000110, 011100\}$

(b)  $A = \{0, 0110, 001100, 0001100\}$ .

**SOLUTION.** In the first case, we show that the elements of  $A^*$  are uniquely created. To see this, let  $\sigma$  be a string in  $A^*$ . We proceed by induction on the length  $\ell$  of  $\sigma$ , the base case  $\ell = 0$  being easy, since the only way to get the empty string is to use none of the elements of  $A$ .

For the induction step,  $\ell > 0$ . There is some way to express  $\sigma = a_1 a_2 \cdots a_k$ , with each  $a_i \in A$ . Consider now the first entry in  $\sigma$ . If it is 1, then  $a_1$  must be 1. Since  $a_2 \cdots a_k \in A^*$  has length  $< \ell$ , the inductive assumption implies it is uniquely created and, therefore,  $\sigma$  is uniquely created.

On the other hand, the first entry in  $\sigma$  might be 0. It follows that, in this case,  $a_1$  is one of 0010, 000110, and 011100. In all cases,  $a_1$  is determined by the number of leading 0's in  $\sigma$  (there can only be one, two, or three leading 0's in  $\sigma$  and each determines which of 0010, 000110, and 011100 is equal to  $a_1$ ). Again,  $a_2 \cdots a_k$  is uniquely created by induction, so we see that  $\sigma$  is uniquely created, as claimed.

For the second case, the elements of  $A^*$  are not uniquely created. For example, 001100001100 can be expressed as  $(0)(0110)(0001100)$  and as  $(001100)(001100)$ .

2. For each of the following sets of binary strings  $S$ , write a decomposition that precisely describes  $S$ , in which the elements are uniquely created. Justify why each has the uniquely created property.

- (a)  $S$  is the set of binary strings in which each occurrence of 1 must be immediately followed by a string of at least two 0's.

**SOLUTION.** In the 1-decomposition, we have  $\{0\}^* (\{1\} \{0\}^*)^*$ . In the strings under consideration for this question, the next two elements after a 1 must both be 0's. Thus, the following is the desired decomposition:  $\{0\}^* (\{1\} \{00\} \{0\}^*)^*$ .

We notice that any decomposition of a string  $\sigma$  into  $\{0\}^* (\{1\} \{00\} \{0\}^*)^*$  is also a decomposition into the 1-decomposition  $\{0\}^* (\{1\} \{0\}^*)^*$  (anything in  $\{00\} \{0\}^*$  is in  $\{0\}^*$ ). Since the 1-decomposition uniquely creates strings, it follows that so does the new decomposition.

- (b)  $S$  is the set of binary strings in which each block of 1's has even length and is followed by a block of 0's whose length is precisely one half that of the length of the block of 1's.

**SOLUTION.** If the block of 1's has length  $a$ , then the succeeding block of 0's must have length  $a/2$ . The, the possibilities are: 110, 111100, 111111000, .... The standard block decomposition is  $\{0\}^*(\{1\}\{1\}^*\{0\}\{0\}^*)^*\{1\}^*$ . A block of 1's at the end of a string cannot occur in this context, so all our strings are contained in  $\{0\}^*(\{1\}\{1\}^*\{0\}\{0\}^*)^*$ . However, we only want some of these. Instead of all the possibilities  $\{1\}\{1\}^*\{0\}\{0\}^*$ , we only want  $\{110, 111100, 111111000, \dots\}$ . Thus, the desired decomposition is  $\{0\}^*(\{110, 111100, 111111000, \dots\})^*$ .

Every way of expressing a string using  $\{0\}^*(\{110, 111100, 111111000, \dots\})^*$  is a way of expressing the string using the standard block decomposition. Since the block decomposition uniquely creates all strings, so does our decomposition.

- (c)  $S$  is the set of binary strings in which each block of 1's has even length and is followed by a block of 0's whose length is at least one half that of the length of the block of 1's.

**SOLUTION.** Let  $S_h$  denote the set  $\{110, 111100, 111111000, \dots\}$  of strings from the solution to (b). The set of strings consisting of a block of 1's of even length followed by a block of 0's of length at least half that of the block of 1's is  $S_h\{0\}^*$ . Thus, the desired decomposition is  $\{0\}^*(S_h\{0\}^*)^*$ .

As  $S_h\{0\}^*$  is a subset of  $\{1\}\{1\}^*\{0\}\{0\}^*$ , our decomposition is again a special case of the standard block decomposition. Therefore the elements of  $\{0\}^*(S_h\{0\}^*)^*$  are uniquely created.

3. The set  $S_e$  of binary strings in which every block of 1's has even length has decomposition  $\{11\}^*(\{0\}\{0\}^*\{11\}\{11\}^*)^*\{0\}^*$ . (You may assume that this is a correct decomposition and that strings are uniquely created by this decomposition.)

- (a) Find the generating function for  $S_e$  with weight function being the length of the string.

**SOLUTION.** Because length adds over the parts of a decomposition, the Product Lemma implies that the generating function for  $S_e$  is the product of the generating functions for  $\{11\}^*$ ,  $(\{0\}\{0\}^*\{11\}\{11\}^*)^*$ , and  $\{0\}^*$ .

By the \*-Lemma,

$$\Phi_{\{11\}^*}(x) = \frac{1}{1 - \Phi_{\{11\}}} = \frac{1}{1 - x^2}.$$

The same analysis implies

$$\Phi_{\{0\}^*}(x) = \frac{1}{1 - x}.$$

Likewise

$$\begin{aligned} \Phi_{(\{0\}\{0\}^*\{11\}\{11\}^*)^*}(x) &= \frac{1}{1 - \Phi_{\{0\}\{0\}^*\{11\}\{11\}^*}} \\ &= \frac{1}{1 - x \left(\frac{1}{1-x}\right) x^2 \left(\frac{1}{1-x^2}\right)} \end{aligned}$$

Therefore,

$$\begin{aligned}
\Phi_{S_e}(x) &= \left( \frac{1}{1-x^2} \right) \left( \frac{1}{1-x \left( \frac{1}{1-x} \right) x^2 \left( \frac{1}{1-x^2} \right)} \right) \left( \frac{1}{1-x} \right) \\
&= \frac{1}{(1-x)(1-x^2) - x^3} \\
&= \frac{1}{1-x-x^2}.
\end{aligned}$$

- (b) Let  $S_{\geq 1 \text{ odd}}$  denote the set of the binary strings in which at least one block of 1's has odd length. Find the generating function for  $S_{\geq 1 \text{ odd}}$ .

**SOLUTION.** Let  $S_{all}$  denote the set of all binary strings. Then  $S_{all} = S_{\geq 1 \text{ odd}} \cup S_e$ . By the Sum Lemma,  $\Phi_{S_{all}}(x) = \Phi_{S_{\geq 1 \text{ odd}}}(x) + \Phi_{S_e}(x)$ . Since we know that

$$\Phi_{S_{all}}(x) = \frac{1}{1-2x},$$

and (from (a))

$$\Phi_{S_e}(x) = \frac{1}{1-x-x^2},$$

we deduce that

$$\begin{aligned}
\Phi_{S_{\geq 1 \text{ odd}}}(x) &= \Phi_{S_{all}}(x) - \Phi_{S_e}(x) \\
&= \frac{1}{1-2x} - \frac{1}{1-x-x^2} \\
&= \frac{(1-x-x^2) - (1-2x)}{(1-2x)(1-x-x^2)} \\
&= \frac{x-x^2}{1-3x+x^2+2x^3}.
\end{aligned}$$

4. The generating function for some set  $S$  with weight function  $w$  satisfies the equation

$$\Phi_S(x) = \frac{x-x^2}{1-3x+x^2+2x^3}.$$

Find a simple closed form expression (not a recurrence relation) for the number of elements of  $S$  that have weight  $n$ .

**SOLUTION.** The roots of  $x^3 - 3x^2 + x + 2$  are 2 and  $(1 \pm \sqrt{5})/2$ . Therefore,

$$[x^n]\Phi_S(x) = A2^n + B \left( \frac{1+\sqrt{5}}{2} \right)^n + C \left( \frac{1-\sqrt{5}}{2} \right)^n.$$

We use the facts that  $\Phi_S(x) = \frac{x-x^2}{1-3x+x^2+2x^3}$  and  $\Phi_S(x) = \sum_{n \geq 0} a_n x^n$  to determine the first few  $a_n$ . We note they satisfy the recurrence  $a_n - 3a_{n-1} + a_{n-2} + 2a_{n-3} = 0$  for  $n \geq 3$ .

We compute  $a_0$ ,  $a_1$ , and  $a_2$  from the equation

$$x - x^2 = (1 - 3x + x^2 + 2x^3) \left( \sum_{n \geq 0} a_n x^n \right).$$

- Comparing the coefficients of  $x^0$  on both sides, we see  $0 = a_0$ .
- Comparing the coefficients of  $x^1$  on both sides, we see  $1 = a_1 - 3a_0$ . Since  $a_0 = 0$ , we conclude  $a_1 = 1$ .
- Comparing the coefficients of  $x^2$  on both sides, we see  $-1 = a_2 - 3a_1 + a_0$ . Since  $a_0 = 0$  and  $a_1 = 1$ , we conclude that  $a_2 = 2$ .

These three values  $a_0 = 0$ ,  $a_1 = 1$ , and  $a_2 = 2$  allow us to compute  $A$ ,  $B$ , and  $C$ .

- $n = 0$ :  $0 = A + B + C$ .
- $n = 1$ :  $1 = 2A + \frac{1+\sqrt{5}}{2}B + \frac{1-\sqrt{5}}{2}C$ .
- $n = 2$ :  $2 = 4A + \left(\frac{1+\sqrt{5}}{2}\right)^2 B + \left(\frac{1-\sqrt{5}}{2}\right)^2 C$ , so  $2 = 4A + \frac{(3+\sqrt{5})}{2}B + \frac{(3-\sqrt{5})}{2}C$ .

Rewriting the second and third equations, we get:

$$\begin{aligned} 1 &= 2A + \frac{1}{2}(B + C) + \frac{\sqrt{5}}{2}(B - C) \quad \text{and} \\ 2 &= 4A + \frac{3}{2}(B + C) + \frac{\sqrt{5}}{2}(B - C). \end{aligned}$$

Using  $\frac{1}{2}(A + B + C) = 0$  and  $\frac{3}{2}(A + B + C) = 0$ , these become

$$\begin{aligned} 1 &= \frac{3}{2}A + \frac{\sqrt{5}}{2}(B - C) \quad \text{and} \\ 2 &= \frac{5}{2}A + \frac{\sqrt{5}}{2}(B - C). \end{aligned}$$

Subtracting the first from the second implies  $A = 1$ . The equations  $A + B + C = 0$  and  $1 = (3/2)A + (\sqrt{5}/2)(B - C)$  become

$$\begin{aligned} -1 &= B + C \quad \text{and} \\ -\frac{1}{2} &= \frac{\sqrt{5}}{2}(B - C). \end{aligned}$$

The second is the same as  $B - C = -\sqrt{5}/5$ . Adding this to  $B + C = -1$ , we conclude  $2B = (-5 - \sqrt{5})/5$ , or  $B = (-5 - \sqrt{5})/10$ . Thus,  $C = (-5 + \sqrt{5})/10$ .

It follows that

$$a_n = 2^n - \left( \frac{5 + \sqrt{5}}{10} \right) \left( \frac{1 + \sqrt{5}}{2} \right)^n + \left( \frac{-5 + \sqrt{5}}{10} \right) \left( \frac{1 - \sqrt{5}}{2} \right)^n.$$