

MATH 239 Supplementary : Eulerian circuits

(An example is provided to illustrate the ideas of this proof.)

Theorem 1. *Let G be a graph with no isolated vertices. Then G has an Eulerian circuit if and only if G is connected and every vertex of G has even degree.*

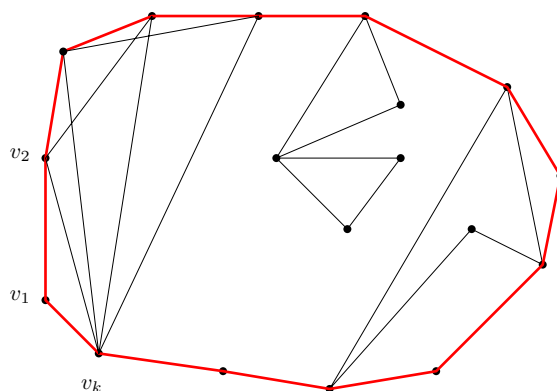
Proof. (\Rightarrow) Exercise.

(\Leftarrow) We will prove by induction on the number of cycles in G .

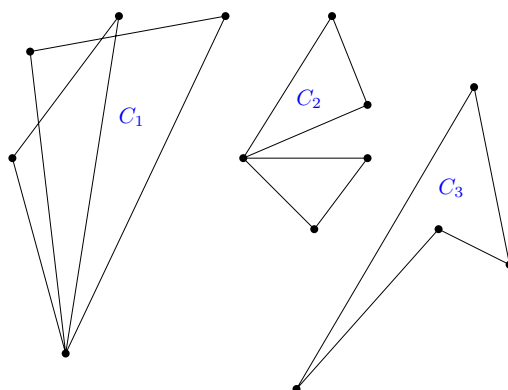
Base case: When G has 1 cycle, it is already an Eulerian circuit.

Induction hypothesis: Assume that any connected graph with fewer cycles than G where every vertex has even degree has an Eulerian circuit.

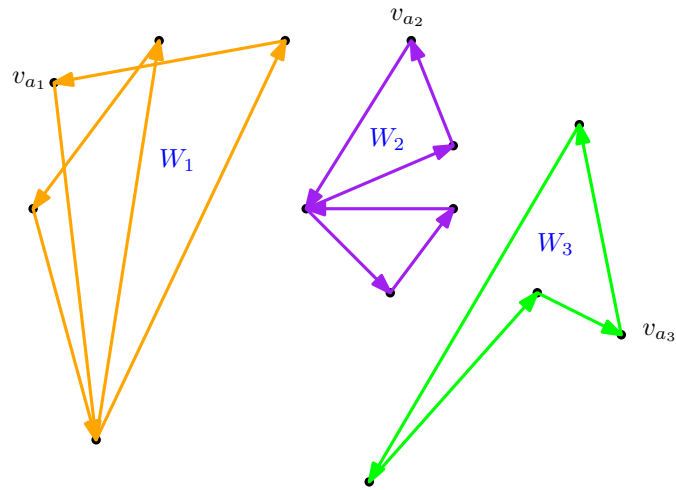
Induction step: Let G be a connected graph where every vertex has even degree. Since there is no isolated vertex, every vertex has degree at least 2. Therefore, there exists a cycle C in G , say the vertices on the cycle in order are $v_1, v_2, \dots, v_k, v_1$.



Remove edges of C from G and remove isolated vertices to obtain G' . Since every vertex is incident with even number of edges in C , every vertex in G' still has even degree. Now G' consists of components C_1, \dots, C_l , each containing fewer cycles than G .



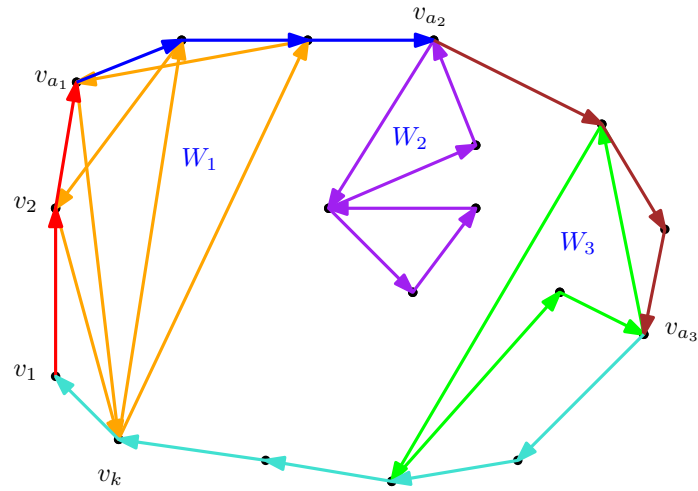
By induction hypothesis, each component C_i has an Eulerian circuit W_i . Moreover, each component must have a common vertex with C , for otherwise G is disconnected.



Let v_{a_i} be one vertex of C_i . Rearrange the components so that $a_1 < a_2 < \dots < a_l$, and let W_i start and end at v_{a_i} .

Then we can construct an Eulerian circuit for G by walking along C and making detours W_i as we hit v_{a_i} :

$$v_1, \dots, v_{a_1} - 1, W_1, v_{a_1} + 1, \dots, v_{a_2} - 1, W_2, v_{a_2} + 1, \dots, v_{a_l} - 1, W_l, v_{a_l} + 1, \dots, v_k, v_1.$$



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