

Math 239 - Tutorial 6

Feb 26, 2013

1. The **odd graph** O_n is the graph whose vertices are the n -subsets of a $(2n + 1)$ -set, two such subsets being adjacent if and only if they are disjoint.
 - (a) Determine how many vertices O_n has.
 - (b) Prove that O_n is k -regular, and determine k in terms of n .
 - (c) Determine how many edges O_n has.

Solution:

- (a) There are $\binom{2n+1}{n}$ n -subsets of a $(2n + 1)$ -set. So O_n has $\binom{2n+1}{n}$ vertices.
- (b) Let T be a particular n -subset of our $(2n + 1)$ -set. Then there are

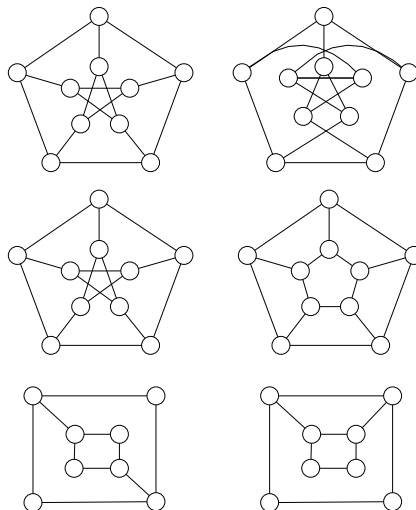
$$\binom{2n+1-n}{n} = \binom{n+1}{n} = n+1$$

n -subsets that do not contain any element of T . So each vertex of O_n has degree $n + 1$, i.e. O_n is $(n + 1)$ -regular.

- (c) By Theorem 4.3.1,

$$\begin{aligned} |E(O_n)| &= \frac{1}{2} \sum_{v \in V(O_n)} \deg(v) \\ &= \frac{1}{2} \binom{2n+1}{n} (n+1) \text{ (since } O_n \text{ is } (n+1)\text{-regular).} \end{aligned}$$

2. Determine which of the following pairs of graphs are isomorphic.



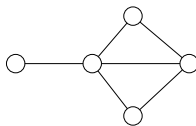
Solution: Pair 1: Isomorphic (just flip horizontally the inside pentagon). Pair 2: Non isomorphic. The second graph has a 4 cycle while the first graph has no 4 cycles. Pair 3: Non isomorphic. The second graph has two adjacent vertices of degree 2, while in the first graph the pairs of vertices of degree 2 are not adjacent.

3. Determine whether or not the following sequences are degree sequences of a graph.

- a) 2,2,2,4,5,5
- b) 3,3,3,4,4,5,5
- c) 1,2,2,3,4

Solution:

- a) There is no graph with this degree sequence. Suppose for the sake of contradiction that there is a graph with degree sequence 2,2,2,4,5,5. Let H be the induced subgraph consisting of the vertices of degree at least 4 in G . H has at most 3 edges, these edges contribute in at most 2 to the sum of the degrees in H , therefore there are at least $4 + 5 + 5 - 3(2) = 8$ edges joining vertices in H and vertices in $V(G) \setminus V(H)$. Let H' be the graph induced by the vertices in $V(G) \setminus V(H)$, as any vertex in H' has degree 2 in G , there are at most $2 + 2 + 2 = 6$ edges joining vertices in H and vertices in $V(G) \setminus V(H)$. Contradiction.
- b) This sequence has an odd number of odd degree vertices. By Theorem 4.3.1 this is impossible.
- c) .



4. A subset of vertices S of G is an **independent set** if no two vertices in S are adjacent in G . Prove that the size of a largest independent set in G is equal to the size of a largest complete subgraph of \bar{G} (the complement of G). (note that by "size", we mean the number of vertices)

Solution:

Recall that

- \bar{G} is the graph with $V(\bar{G}) = V(G)$, and $\{u, v\} \in E(\bar{G})$ if and only if $\{u, v\} \notin E(G)$
- a complete graph is one in which all pairs of distinct vertices are adjacent.

Then (using the definitions) an independent set in G will be a complete subgraph of \bar{G} . So the size of a largest independent set in G must be the same as the size of a largest complete subgraph of \bar{G} .

5. A graph G is called self-complimentary if G is isomorphic to \bar{G} . Show that if a graph G is self-complimentary, then the number of vertices n is congruent to 0 or 1 modulo 4.

Solution: Note that G and \bar{G} have the same number of edges m , and their union is K_n . K_n has $\frac{1}{2}n(n-1)$ edges, and as G has half as many edges as K_n , G has $\frac{1}{4}n(n-1)$ edges. As the number of edges is integer, this means that $n(n-1)$ must be divisible by 4. As either n or $n-1$ is an odd number, this means that the other factor has to be divisible by 4 on its own, meaning that either n or $n-1$ is divisible by 4. But this is equivalent to n being congruent to either 0 or 1 modulo 4.