

Regression Verification: Project Proposal

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Introduction

Context

- Formal Methods of Software Development
- Project Group at KIT
- Student research project

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Main Question

How to prevent regressions in software development?

BILD Devolution

Introduction

Formal Verification

Formally prove correctness of software
⇒ Requires formal specification

Regression Testing

Discover new bugs by testing for them
⇒ Requires test cases

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Formal Verification

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Regression Testing

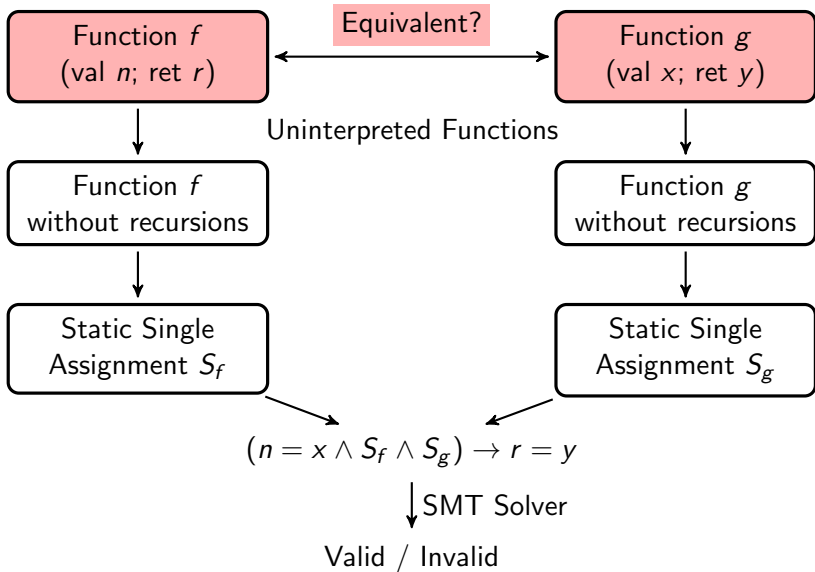
Discover new bugs by testing for them
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Regression Verification

Formally prove there are no new bugs

Regression Verification

Overview



Regression Verification

Formally prove there are no new bugs

- Goal: Proving the equivalence of two **closely related** programs
- No formal specification or test cases required
- Instead use old program version
- Make use of similarity between programs

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Regression Verification

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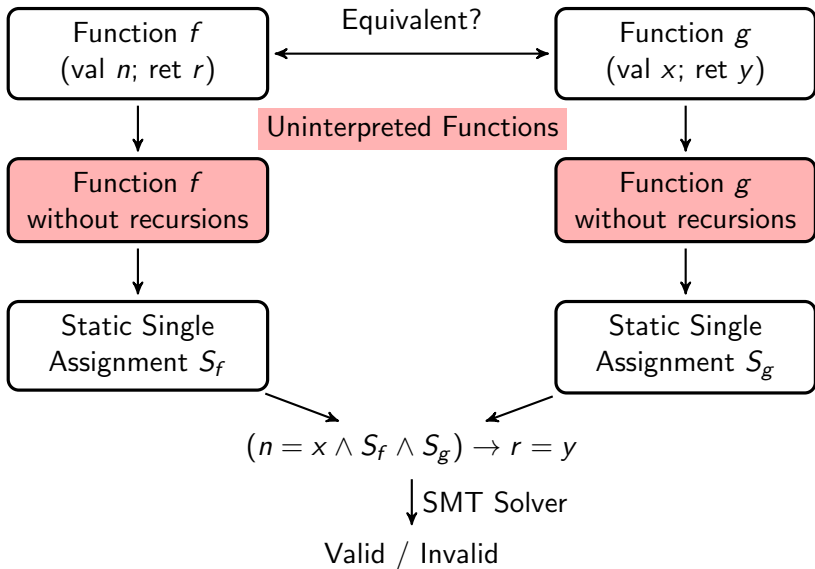
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Uninterpreted Functions

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Uninterpreted Functions

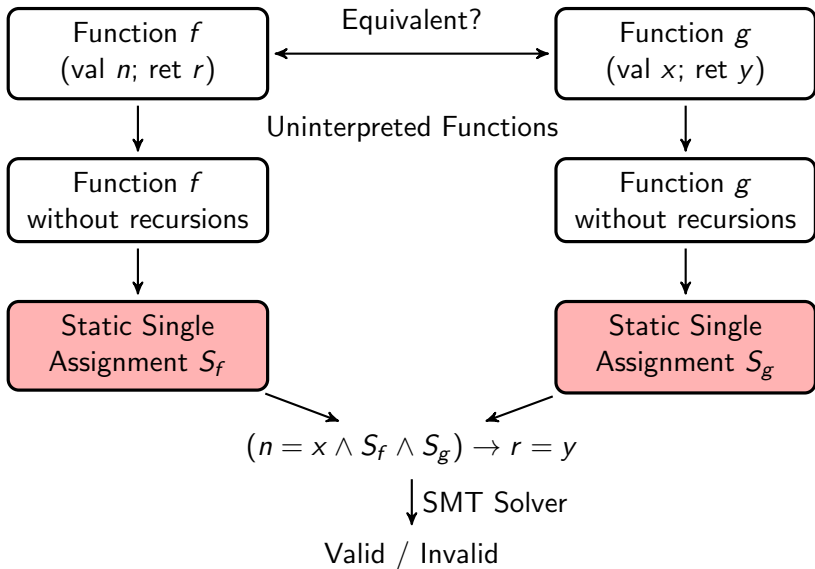
- Given the same inputs an **Uninterpreted Function** always returns the same outputs.
- Motivation: Proof by Induction, to prove $f(n) = g(n)$ assume $f(n-1) = g(n-1)$

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Static Single Assignment

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Static Single Assignment

- Translate program functions to formulas
- Recursions: Abstraction by Uninterpreted Function
- In assignments $x = exp$ replace x with a new variable x_1
- Represents the states of the program

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$$S_f = \begin{pmatrix} r_0 = 0 \\ n \leq 0 \rightarrow r_1 = n \end{pmatrix}$$

\wedge

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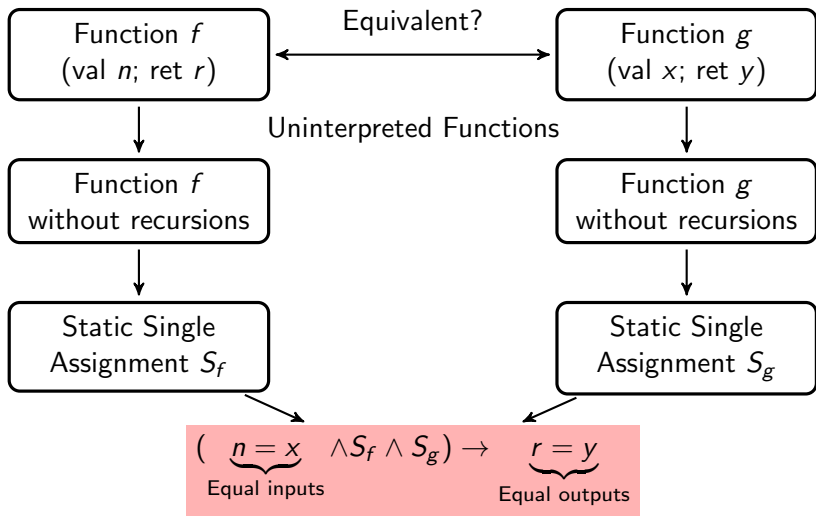
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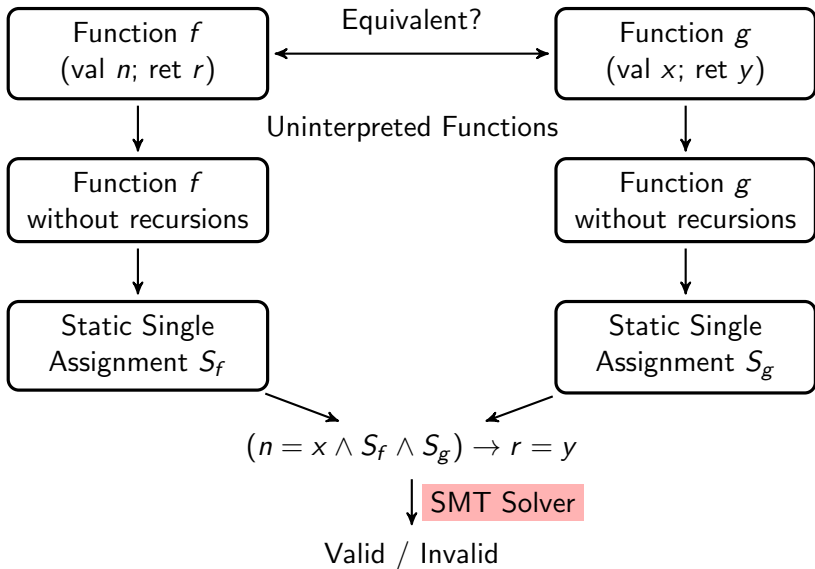
Formula

Overview



SMT Solver

Overview



Extensions

- SMT solver still complains:

$$f(n) = \begin{cases} -1 & \text{if } n = 0 \\ g(n) & \text{otherwise} \end{cases}$$

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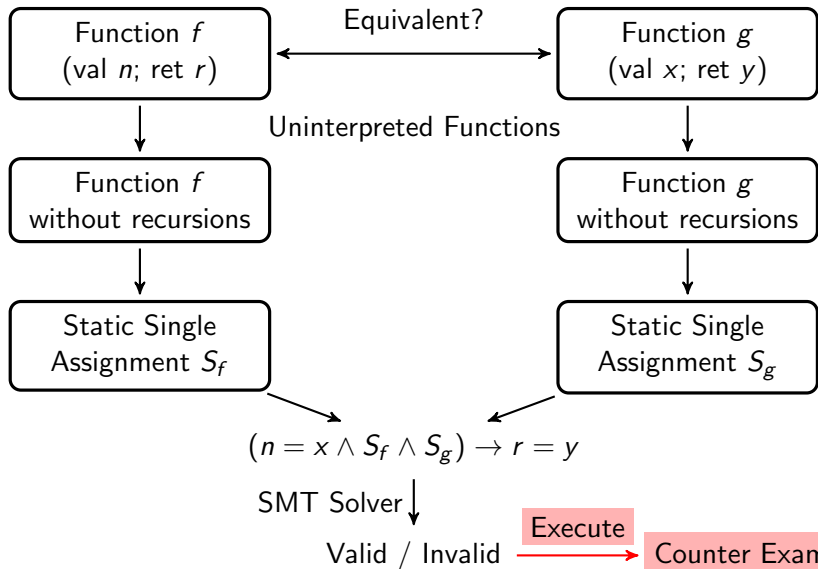
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- But we can fix it:

$$f(0) = 0$$

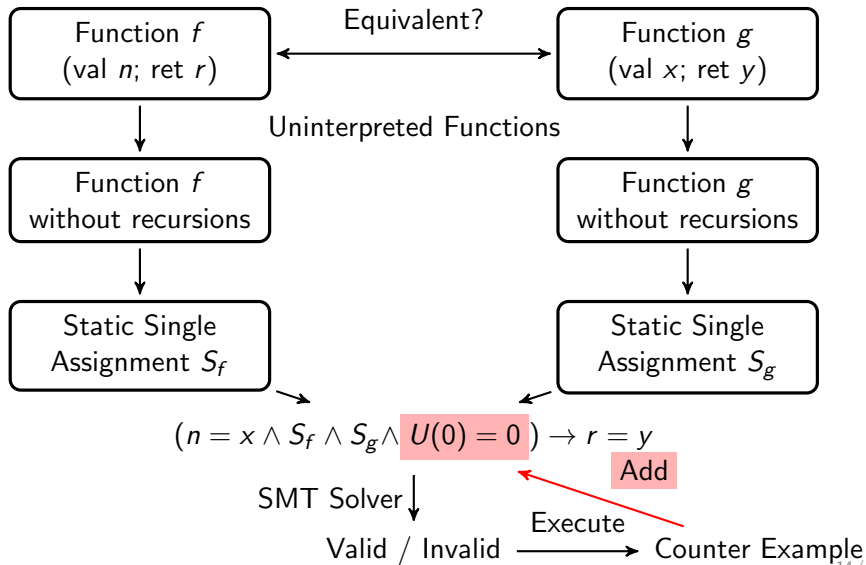
Extensions

Finding Counter Examples



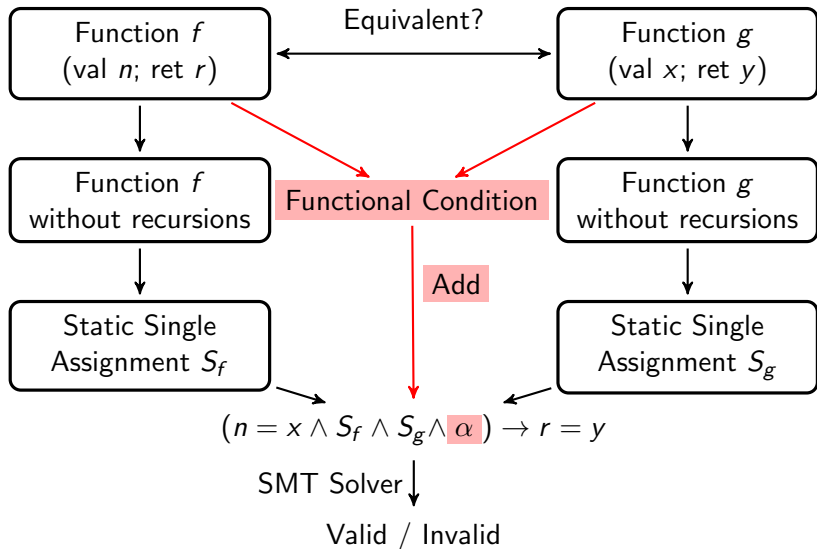
Extensions

Determining Corner Cases



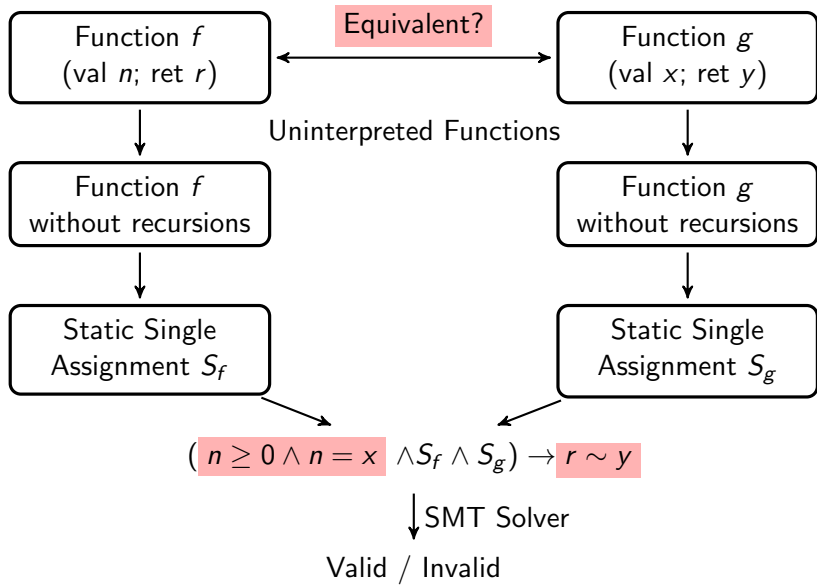
Extensions

Functional Condition Extraction



Extensions

Relational Equivalence



Example Catalog

- Collect examples: Papers, Refactoring Rules, ...
- 51 program pairs so far
- Test how well approach and extensions work

Conclusion

Regression Verification

- Better chance of being adopted than Formal Verification
- More powerful than Regression Testing

My Contributions

- Implementation for simple language
- Various extensions to cover more cases
- Example catalog for evaluation