

# Parallel Graph Algorithms on the Xeon Phi Coprocessor

Master Thesis presentation

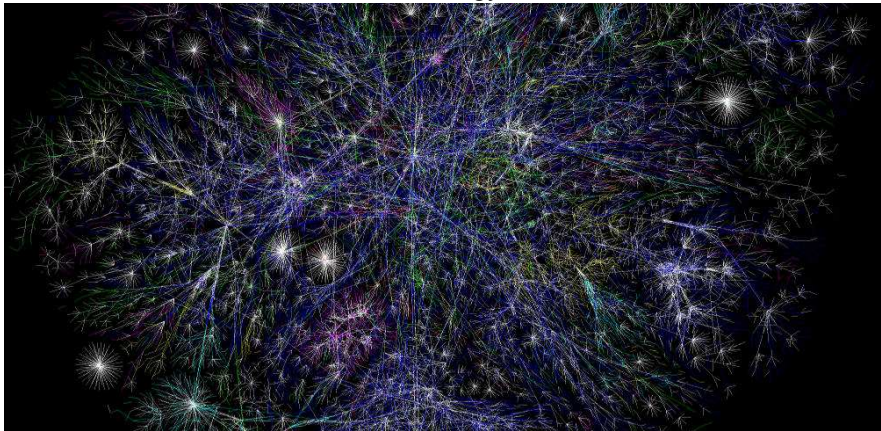
Dennis Felsing | 2015-09-07

INSTITUTE OF THEORETICAL INFORMATICS, RESEARCH GROUP PARALLEL COMPUTING



# Motivation: Graphs

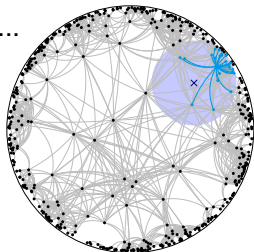
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- Occur in social networks, cell biology, the internet...



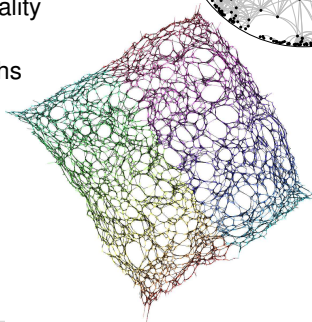
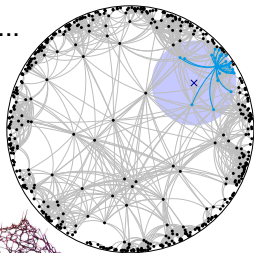
Map of the Internet, <http://www.opte.org/maps/>

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  - Preserves privacy and confidentiality
  - No need to transfer big data
  - Scale to smaller and bigger graphs

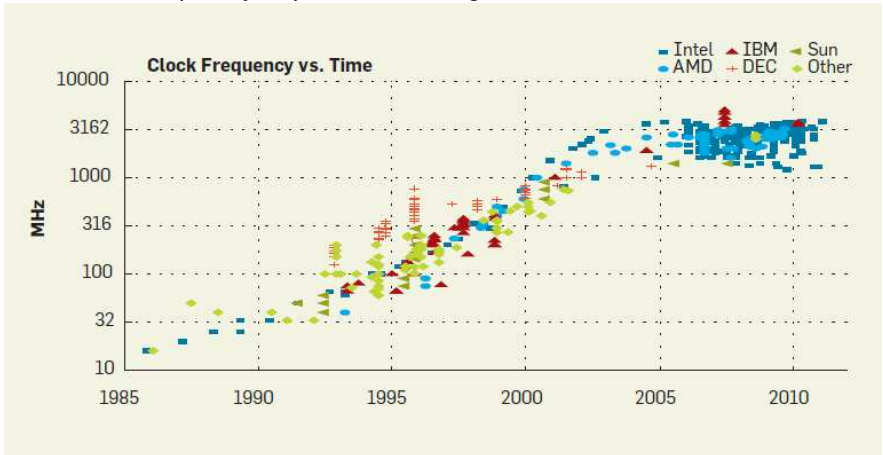


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- **Graph Drawing**
  - Lay out graph visually, mainly for human perception



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⇒ Port two graph algorithms to Xeon Phi  
and evaluate the porting and their performance



- 1 Xeon Phi Coprocessor
  - Hardware Architecture
  - Programming
- 2 Generation of Massive Complex Networks
  - Algorithm
  - Results
- 3 Graph Drawing using Graph Clustering
  - Algorithm
  - Results

# Xeon Phi: Hardware Architecture

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- Simple cores based on original Pentium design from 1994
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Name	2× Xeon E5-2680	Xeon Phi 5110P
Release Date	2012	2012
Clock Frequency	2.7 GHz	1.05 GHz
Cores	2 × 8 (32 threads)	60 (240 threads)
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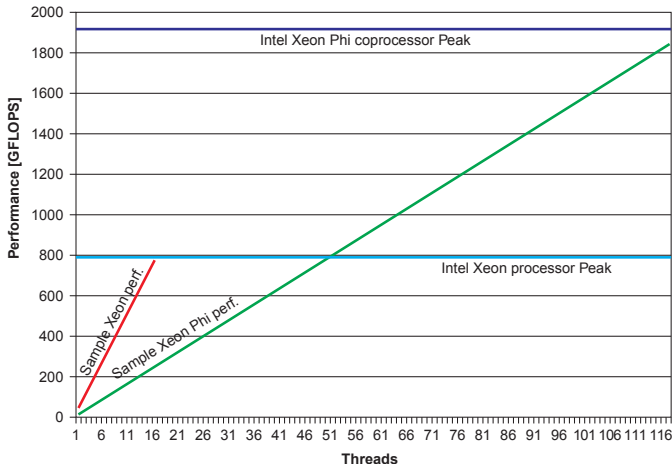
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⇒ Parallelization and vectorization necessary to reach high performance

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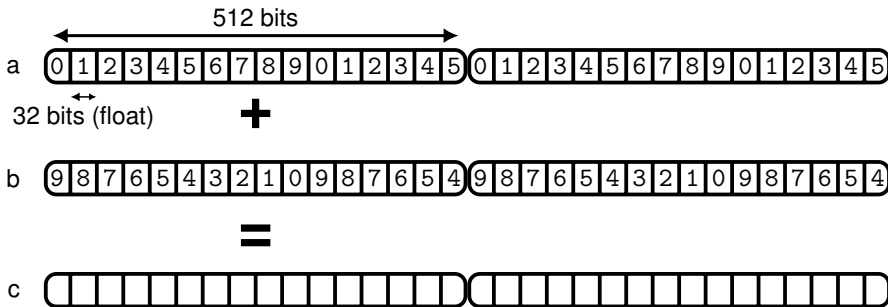
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#pragma offload_wait target(mic:0) wait(c)
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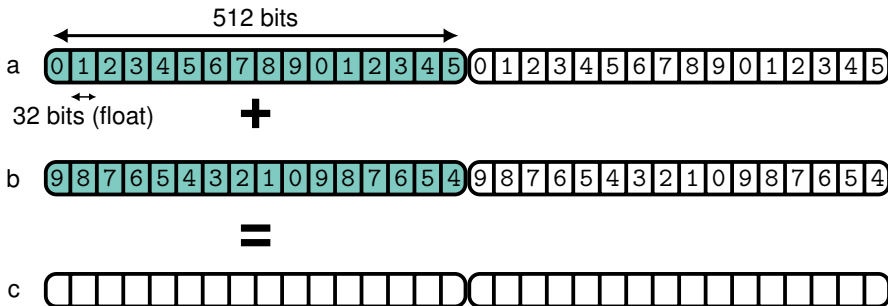
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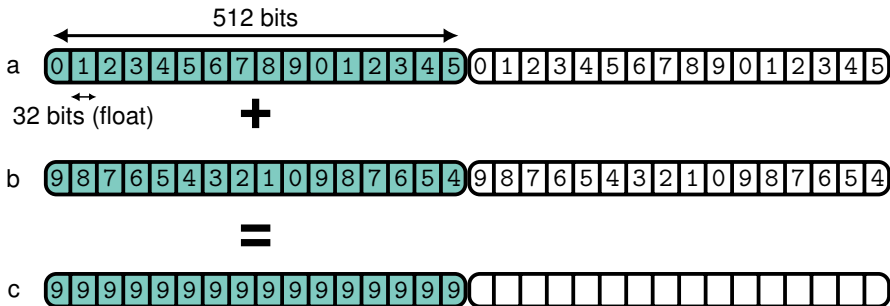
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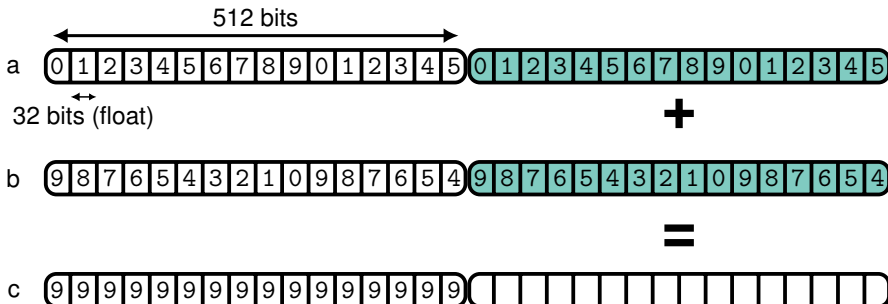
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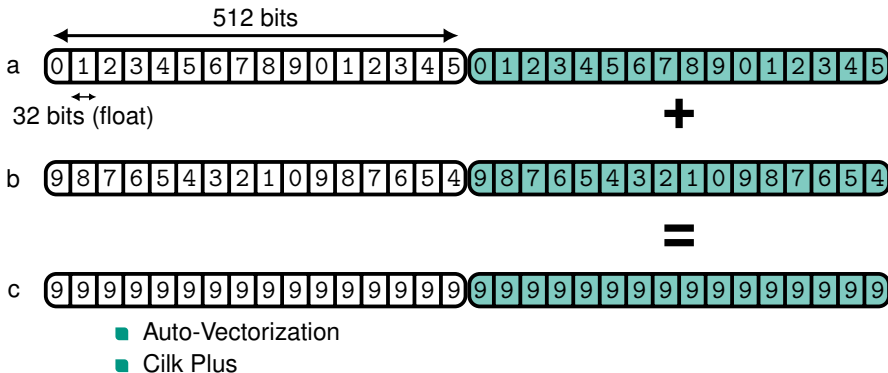


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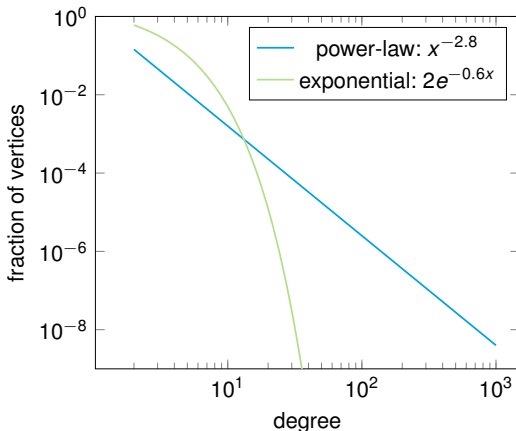
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Desired properties of a complex network:

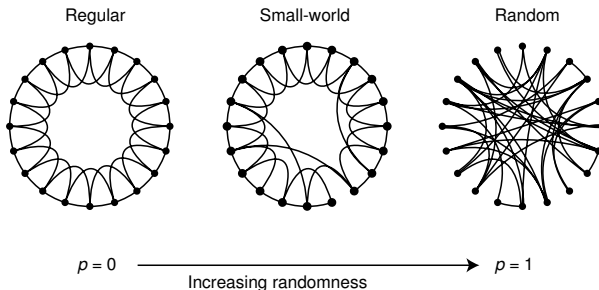
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# Network Generation: Algorithm

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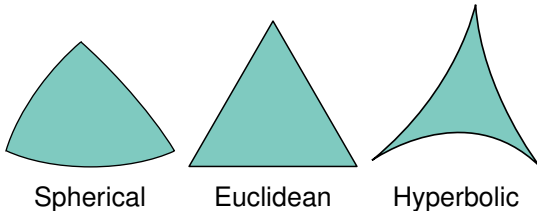


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⇒ Generator using **hyperbolic geometry** performs well in both properties



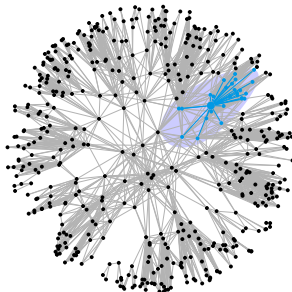
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  - ⇒ Natural embedding of graphs with tree-like structure
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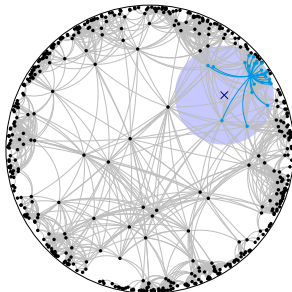
M.C. Escher: Circle Limit IV

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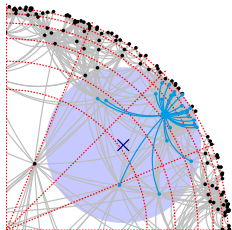




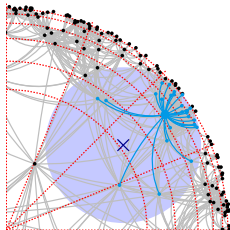
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⇒ Subquadratic running time  $O((n^{3/2} + m) \log n)$



- Implementation part of NetworKit, high level C++11 code
- NetworKit ported to Intel C++ compiler 15.0 and Xeon Phi
- Working around compiler restrictions and bugs with *constexpr*, implicit conversions, null pointers, the standard library, and function traits on lambdas
- Three execution modes implemented and tested:
  - No offloading: Entire code runs on Xeon Phi, not enough memory
  - Full offloading: Offload parts of the calculation, keep results in memory of host system
  - Partial offloading: Offload part of the calculation, other part on the host system

# Network Generation: Results

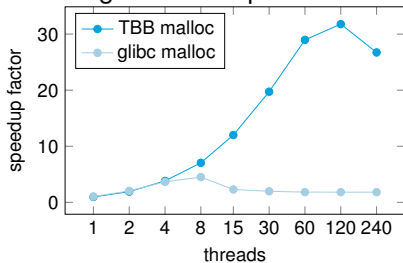
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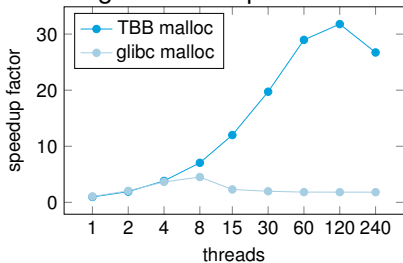
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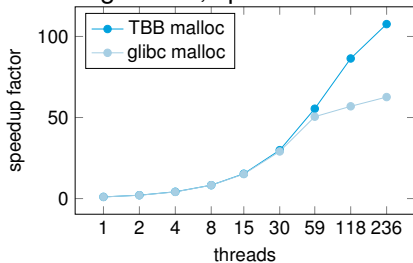
⇒ Use **non-locking allocations** of Intel's **Threading Building Blocks**

⇒ Reduce number of (re)allocations by **reusing memory** and **preallocating expected size**

## ■ Scaling of initial implementation:



## ■ Scaling of final, optimized code:

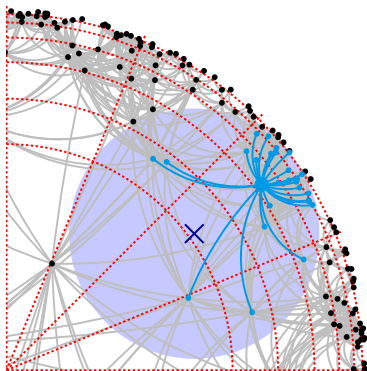


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- Tuning parameters of the algorithm:

**Capacity:** Maximum number of vertices in a leaf cell before split

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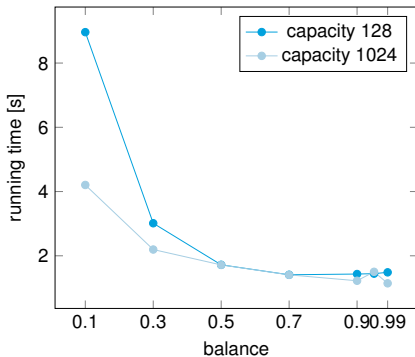
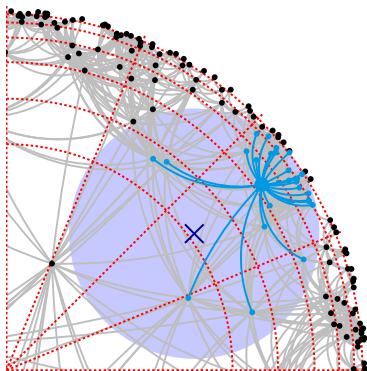


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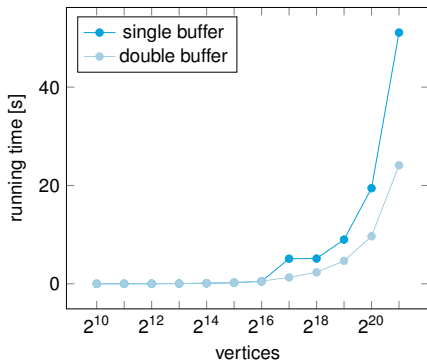


⇒ **Imbalanced quadtree** with greater space to outer children

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- Transferring parts of graph back to host system is slow  
⇒ Double Buffering on Phi and host

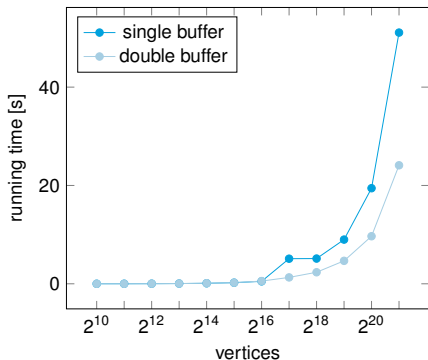
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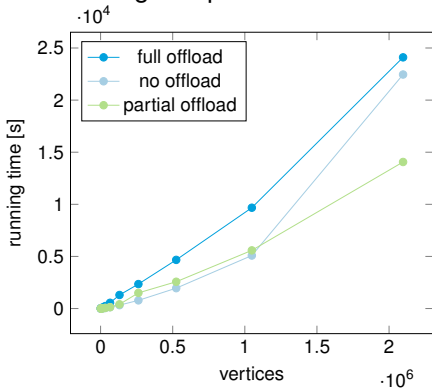
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⇒ With many tricks similar speed as dual Xeon, but not faster

- Offloading comparison:



# Graph Drawing: Algorithm

We assume to have graphs with predefined **target edge lengths**

- **Full Stress Model**: Physical springs connecting all pairs of vertices
- **Maxent-Stress Model**: Minimize stress, maximize entropy:

$$M(x) = \underbrace{\sum_{\{u,v\} \in E} \overbrace{w_{uv}}^{\text{weight factor}} (\underbrace{\|x_u - x_v\|}_{\text{target edge length}} - \overbrace{d_{uv}}^{\text{target edge length}})^2}_{\text{stress for target edge lengths}} - \alpha \underbrace{\sum_{\{u,v\} \notin E} \ln \|\overbrace{x_u - x_v}^{\text{coordinate vector}}\|}_{\text{entropy for rest}}$$

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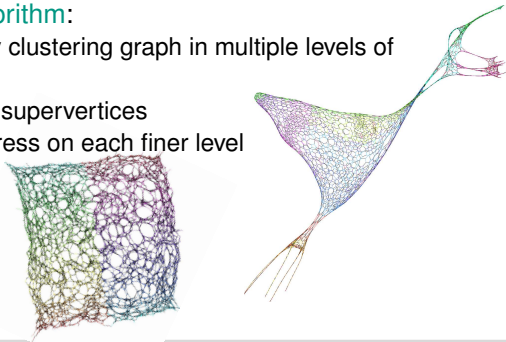
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- **Multilevel Maxent-Stress Algorithm:**

- Minimize maxent-stress by clustering graph in multiple levels of hierarchy
- Contract clusters into new supervertices
- Iteratively solve maxent-stress on each finer level

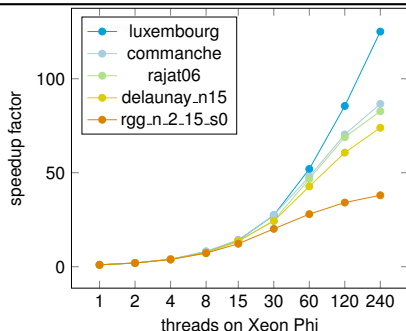
⇒ Parallelizes well



- Preliminary implementation of multilevel maxent-stress graph drawing algorithm
- Parallelized with OpenMP
- Graphs of interest ( $< 10^7$  edges) easily fit into Xeon Phi memory  
⇒ No expensive offloading necessary
- Source code libraries had to be fixed for ICPC
- No major dynamic allocations in algorithm  
⇒ Intel TBB's *malloc* has small effect
- Inner loop vectorizes well when isolated: speedup factor 7.0  
Smaller effect when embedded in real program  
⇒ Other calculations at same time, hyper-threading, memory connection busy

# Graph Drawing: Results

Graph	$n$	$m$	Description	Phi	Host
nyc	264 346	365 050	Road Network	960.0	1845.9
luxembourg	114 599	119 666	Road Network	89.5	166.5
commanche	7920	11 880	Helicopter Mesh	2.6	3.5
rajat06	10 922	18 061	Circuit Simulation	3.5	4.5
delaunay_n15	32 768	98 274	Delaunay Triangulation	7.8	9.0
rgg_n.2.15.s0	32 768	160 240	Random Graph	5.3	3.6

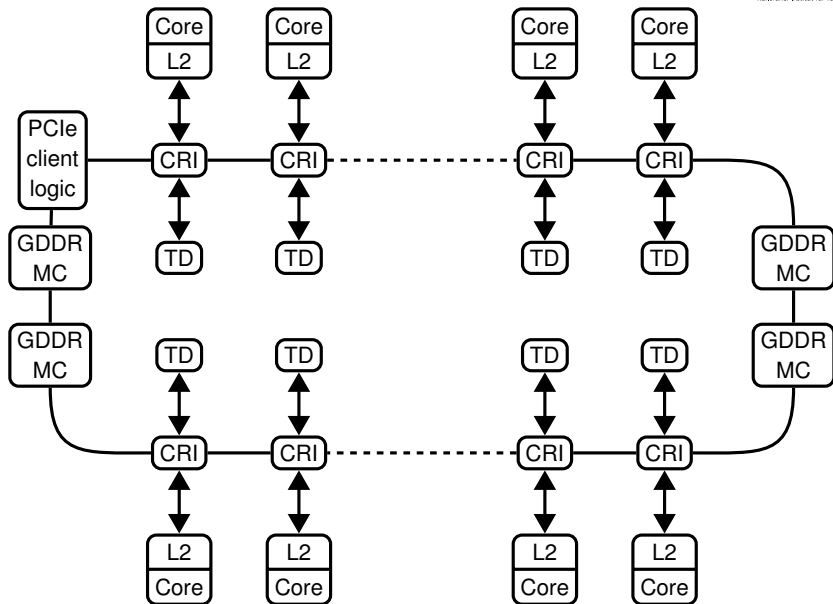


- All graphs small enough to fit into Xeon Phi memory  
⇒ Executed on Xeon Phi directly
- Good speedup for large sparse graphs

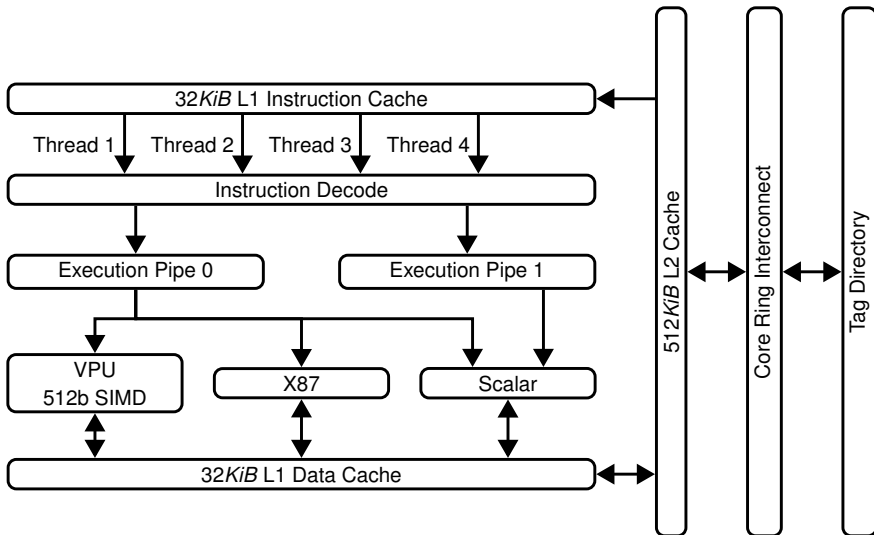
- Complex algorithms and their framework/libraries ported to Xeon Phi
- Good scaling in both algorithms on Xeon Phi
- Offloading large amounts of data too expensive
- Graph drawing algorithm outperforms two-socket Intel Xeon system, especially on sparse graphs
- Future Research: Direct comparison between graph algorithms on GPU and Xeon Phi
- New Xeon Phi “Knights Landing” this year with modern cores and 384 GB of memory



# [Appendix] Xeon Phi: Layout Overview



# [Appendix] Xeon Phi: Core Pipeline



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```
float a[MAX], b[MAX], c[MAX];  
#pragma offload target(mic:0) in(a, b) out(c) signal(c)  
{  
    #pragma omp parallel for  
    for (i = 0; i < MAX; i++)  
        c[i] = a[i] + b[i];  
}  
#pragma offload_wait target(mic:0) wait(c)
```

- Cilk Plus
- Threading Building Blocks
- MPI
- **Vectorization** (Single Instruction Multiple Data) methods

- Presented as a regular computer, stripped down Linux, SSH

- **Parallelization** methods:

- OpenMP
- OpenMP Offloading
- Cilk Plus

```
float a[MAX], b[MAX], c[MAX];  
cilk_for (i = 0; i < MAX; i++)  
    c[i] = a[i] + b[i];
```

- Threading Building Blocks
- MPI

- **Vectorization** (Single Instruction Multiple Data) methods



- Presented as a regular computer, stripped down Linux, SSH

- **Parallelization** methods:

- OpenMP
- OpenMP Offloading
- Cilk Plus
- Threading Building Blocks

```
float a[MAX], b[MAX], c[MAX];  
parallel_for(size_t(0), MAX, size_t(1), [=](size_t i) {  
    c[i] = a[i] + b[i];  
});
```

- MPI

- **Vectorization** (Single Instruction Multiple Data) methods

- Presented as a regular computer, stripped down Linux, SSH
- **Parallelization** methods
- **Vectorization** (Single Instruction Multiple Data) methods:
  - Manual Vectorization: `c = _mm512_add_ps(a, b)`

```
1  #include <immintrin.h>
2  float a[MAX] __attribute__((aligned(64)));
3  float b[MAX] __attribute__((aligned(64)));
4  float c[MAX] __attribute__((aligned(64)));
5  __m512 _a, _b, _c;
6  for (i = 0; i < MAX; i += 16) {
7      _a = _mm512_load_ps(&a[i]);
8      _b = _mm512_load_ps(&b[i]);
9      _c = _mm512_add_ps(_a, _b);
10     _mm512_store_ps(&c[i], _c);
11 }
```

- Auto-Vectorization
- Cilk Plus

- Presented as a regular computer, stripped down Linux, SSH
- **Parallelization** methods
- **Vectorization** (Single Instruction Multiple Data) methods:
  - Manual Vectorization
  - Auto-Vectorization

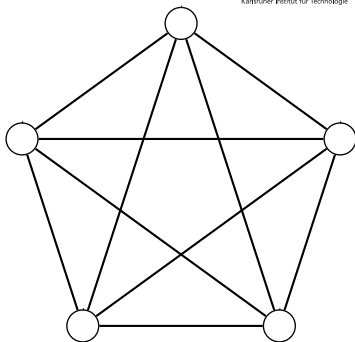
```
float a[MAX] __attribute__((aligned(64)));  
float b[MAX] __attribute__((aligned(64)));  
float c[MAX] __attribute__((aligned(64)));  
#pragma simd  
for (i = 0; i < MAX; i++)  
    c[i] = a[i] + b[i];
```

- Cilk Plus

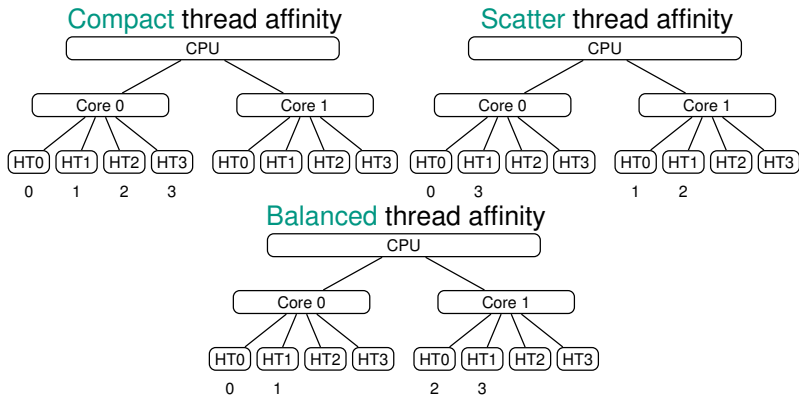
- Presented as a regular computer, stripped down Linux, SSH
- **Parallelization** methods
- **Vectorization** (Single Instruction Multiple Data) methods:
  - Manual Vectorization
  - Auto-Vectorization
  - Cilk Plus

```
float a[MAX] __attribute__((aligned(64)));  
float b[MAX] __attribute__((aligned(64)));  
float c[MAX] __attribute__((aligned(64)));  
c[i:MAX] = a[i:MAX] + b[i:MAX];
```

- Graph  $G = (V, E)$  consists of
  - Set of vertices  $V$
  - Set of edges  $E \subseteq V \times V$
- Edge  $e = (u, v) \in E$ : connection from source  $u$  to target  $v$
- Number of vertices  $n = |V|$
- Number of edges  $m = |E|$
- Undirected graphs only:  $(u, v) \in E$  iff  $(v, u) \in E$
- Neighborhood  $N(u) = \{v : (u, v) \in E\}$
- Loop  $(u, u) \in E$
- Degree  $\deg(v)$ : number of incident edges, counting loops twice
- Distance: number of edges in shortest path connecting vertices
- Diameter  $d$ : greatest distance between any pair of vertices



## ■ OpenMP Scheduling:



⇒ Scatter and balanced 1.4 times faster