



# Fundamentals of Optimization

## Exercise 3 – Solutions

### Remarks

- All questions that are available in the STACK quiz are duly marked. Please solve those using STACK.
- We have added marks for each question. Please note that those are purely for illustrative purposes. The exercise set will not be marked.
- We can derive the inverse of a nonsingular matrix  $A \in \mathbb{R}^{2 \times 2}$  in closed form:

$$A^{-1} = \begin{bmatrix} a & b \\ c & d \end{bmatrix}^{-1} = \frac{1}{\det(A)} \begin{bmatrix} d & -b \\ -c & a \end{bmatrix}.$$

## STACK Problems

### 1 Basic Solutions of Polyhedra in Standard Form (2 marks)

#### (1.1) STACK question

By using the enumeration algorithm presented in Section 13.2.5 of the lecture notes, determine the set of all basic solutions and basic feasible solutions of the following polyhedron:  $\mathcal{P} = \{x \in \mathbb{R}^4 : x_1 + 2x_2 - x_3 + 4x_4 = 10; 2x_1 + 3x_2 - 2x_3 - 2x_4 = 16; x \geq \mathbf{0}\}$ . For each basic solution and basic feasible solution, determine whether it is degenerate or nondegenerate. You can assume that the coefficient matrix  $A$  has full row rank.

[1 mark]

#### Solution

We have  $n = 4$  and  $m = 2$ . Since  $A$  has full row rank, we have  $|B| = m = 2$ . Moreover,

$$A = \begin{bmatrix} 1 & 2 & -1 & 4 \\ 2 & 3 & -2 & -2 \end{bmatrix}, \quad b = \begin{bmatrix} 10 \\ 16 \end{bmatrix}.$$

To solve this question, we consider all possible subsets  $B \subset \{1, 2, 3, 4\}$  such that  $|B| = 2$ , and check whether the conditions for a basic (feasible) solution hold.

- $B = \{1, 2\}$ : The columns  $A^1$  and  $A^2$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 1 & 2 \\ 2 & 3 \end{bmatrix}^{-1} = \begin{bmatrix} -3 & 2 \\ 2 & -1 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 2 \\ 4 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_1 = 2$  and  $\hat{x}_2 = 4$ , and the nonbasic variables are  $\hat{x}_3 = 0$  and  $\hat{x}_4 = 0$ . Therefore,  $\hat{x} = [2, 4, 0, 0]^T$  is a basic feasible solution as  $\hat{x} \geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

- $B = \{1, 3\}$ : The columns  $A^1$  and  $A^3$  are linearly dependent since  $A^1 + A^3 = \mathbf{0} \in \mathbb{R}^2$ . Therefore, this choice does not give rise to a basic solution.

- $B = \{1, 4\}$ : The columns  $A^1$  and  $A^4$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 1 & 4 \\ 2 & -2 \end{bmatrix}^{-1} = \frac{1}{10} \begin{bmatrix} 2 & 4 \\ 2 & -1 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 42/5 \\ 2/5 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_1 = 42/5$  and  $\hat{x}_4 = 2/5$ , and the nonbasic variables are  $\hat{x}_2 = 0$  and  $\hat{x}_3 = 0$ . Therefore,  $\hat{x} = [42/5, 0, 0, 2/5]^T$  is a basic feasible solution as  $\hat{x} \geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

- $B = \{2, 3\}$ : The columns  $A^2$  and  $A^3$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 2 & -1 \\ 3 & -2 \end{bmatrix}^{-1} = \begin{bmatrix} 2 & -1 \\ 3 & -2 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 4 \\ -2 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_2 = 4$  and  $\hat{x}_3 = -2$ , and the nonbasic variables are  $\hat{x}_1 = 0$  and  $\hat{x}_4 = 0$ . Therefore,  $\hat{x} = [0, 4, -2, 0]^T$  is a basic solution but not feasible since  $\hat{x} \not\geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

- $B = \{2, 4\}$ : The columns  $A^2$  and  $A^4$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 2 & 4 \\ 3 & -2 \end{bmatrix}^{-1} = \frac{1}{16} \begin{bmatrix} 2 & 4 \\ 3 & -2 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 21/4 \\ -1/8 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_2 = 21/4$  and  $\hat{x}_4 = -1/8$ , and the nonbasic variables are  $\hat{x}_1 = 0$  and  $\hat{x}_3 = 0$ . Therefore,  $\hat{x} = [0, 21/4, 0, -1/8]^T$  is a basic solution but not feasible since  $\hat{x} \not\geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

- $B = \{3, 4\}$ : The columns  $A^3$  and  $A^4$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} -1 & 4 \\ -2 & -2 \end{bmatrix}^{-1} = \frac{1}{10} \begin{bmatrix} -2 & -4 \\ 2 & -1 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} -42/5 \\ 2/5 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_3 = -42/5$  and  $\hat{x}_4 = 2/5$ , and the nonbasic variables are  $\hat{x}_1 = 0$  and  $\hat{x}_2 = 0$ . Therefore,  $\hat{x} = [0, 0, -42/5, 2/5]^T$  is a basic solution but not feasible since  $\hat{x} \not\geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

Therefore,  $\mathcal{P}$  has two basic feasible solutions (i.e., vertices) and three basic solutions that are not feasible, all of which are nondegenerate.

## (1.2) STACK question

By using the enumeration algorithm presented in Section 13.2.5 of the lecture notes, determine the set of all basic solutions and basic feasible solutions of the following polyhedron:  $\mathcal{P} = \{x \in \mathbb{R}^3 : 3x_1 + 2x_2 + 4x_3 = 4; -x_1 + x_2 - 2x_3 = 2; x \geq \mathbf{0}\}$ . For each basic solution and basic feasible solution, determine whether it is degenerate or nondegenerate. You can assume that the coefficient matrix  $A$  has full row rank.

[1 mark]

## Solution

We have  $n = 3$  and  $m = 2$ . Since  $A$  has full row rank, we have  $|B| = m = 2$ . Moreover,

$$A = \begin{bmatrix} 3 & 2 & 4 \\ -1 & 1 & -2 \end{bmatrix}, \quad b = \begin{bmatrix} 4 \\ 2 \end{bmatrix}.$$

To solve this question, we consider all possible subsets  $B \subset \{1, 2, 3\}$  such that  $|B| = 2$  and check whether the conditions for a basic (feasible) solution hold.

- $B = \{1, 2\}$ : The columns  $A^1$  and  $A^2$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 3 & 2 \\ -1 & 1 \end{bmatrix}^{-1} = \frac{1}{5} \begin{bmatrix} 1 & -2 \\ 1 & 3 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 0 \\ 2 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_1 = 0$  and  $\hat{x}_2 = 2$ , and the only nonbasic variable is  $\hat{x}_3 = 0$ . Therefore,  $\hat{x} = [0, 2, 0]^T$  is a basic feasible solution as  $\hat{x} \geq \mathbf{0}$ . Note that  $\hat{x}$  is degenerate since there is at least one basic variable that is equal to zero (i.e.,  $\hat{B} \subset B$  or  $1 = |\hat{B}| < |B| = 2$ ).

- $B = \{1, 3\}$ : The columns  $A^1$  and  $A^3$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 3 & 4 \\ -1 & -2 \end{bmatrix}^{-1} = \frac{1}{2} \begin{bmatrix} 2 & 4 \\ -1 & -3 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 8 \\ -5 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_1 = 8$  and  $\hat{x}_3 = -5$ , and the only nonbasic variable is  $\hat{x}_2 = 0$ . Therefore,  $\hat{x} = [8, 0, -5]^T$  is a basic solution but not feasible since  $\hat{x} \not\geq \mathbf{0}$ . Note that  $\hat{x}$  is nondegenerate since both basic variables are different from zero (i.e.,  $\hat{B} = B$ ).

- $B = \{2, 3\}$ : The columns  $A^2$  and  $A^3$  are linearly independent and we get

$$(A_B)^{-1} = \begin{bmatrix} 2 & 4 \\ 1 & -2 \end{bmatrix}^{-1} = \frac{1}{8} \begin{bmatrix} 2 & 4 \\ 1 & -2 \end{bmatrix} \quad \text{and} \quad \hat{x}_B = (A_B)^{-1}b = \begin{bmatrix} 2 \\ 0 \end{bmatrix}.$$

Hence, the basic variables are  $\hat{x}_2 = 2$  and  $\hat{x}_3 = 0$ , and the only nonbasic variable is  $\hat{x}_1 = 0$ . Therefore,  $\hat{x} = [0, 2, 0]^T$  is a basic feasible solution as  $\hat{x} \geq \mathbf{0}$ . Note that  $\hat{x}$  is degenerate since there is at least one basic variable that is equal to zero (i.e.,  $\hat{B} \subset B$  or  $1 = |\hat{B}| < |B| = 2$ ). Note that this is the same basic feasible solution as the one given by  $B = \{1, 2\}$ .

It follows that  $\mathcal{P}$  has one degenerate basic feasible solution (i.e., vertex) and one nondegenerate basic solution that is not feasible.

## 2 Optimality Conditions and Degeneracy (3 marks)

### (2.1) STACK question

Consider the following linear program in standard form

$$\min\{-x_1 - 4x_2 - x_3 + 2x_4 : x_1 + 4x_2 + x_3 = 8; x_1 + 2x_2 + x_4 = 4; x \geq \mathbf{0}\}$$

and the vertices

- (a)  $\hat{x} = [4, 0, 4, 0]^T$ .
- (b)  $\hat{x} = [0, 2, 0, 0]^T$ ,
- (c)  $\hat{x} = [0, 0, 8, 4]^T$ .

You can assume that the coefficient matrix  $A$  has full row rank. For each vertex, decide whether the vertex is optimal or not, and whether it is degenerate or not.

**[3 marks]**

Write down a valid choice for the index set  $B$ , the reduced costs  $\bar{c}_j$ ,  $j \in \{1, \dots, n\}$ , for that basis, and a “candidate” improving direction  $d \in \mathbb{R}^n$  if one exists. For the latter, if  $\bar{c} \not\geq \mathbf{0}$ , use  $d_{j^*} = 1$  and  $d_j = 0$ ,  $j \in N \setminus \{j^*\}$  to derive the direction, where  $j^* \in N$  is the index with the smallest reduced cost  $\bar{c}_j$ . Verify whether the candidate improving direction  $d$  is indeed an improving feasible direction at that vertex. If  $\bar{c} \geq \mathbf{0}$ , then enter  $d = \mathbf{0}$ .

If the vertex is degenerate, write down all possible choices of the indices for the index set  $B$ , together with the corresponding reduced costs  $\bar{c}$  and candidate improving directions.

## Solution

First observe that  $n = 4$  and  $m = 2$ . Since  $A$  has full row rank,  $|B| = 2$ . Moreover,

$$A = \begin{bmatrix} 1 & 4 & 1 & 0 \\ 1 & 2 & 0 & 1 \end{bmatrix}, \quad b = \begin{bmatrix} 8 \\ 4 \end{bmatrix}, \quad c = [-1, -4, -1, 2]^T.$$

- (a) We have  $\hat{x}_1 > 0$  and  $\hat{x}_3 > 0$  and, hence,  $\hat{B} = \{1, 3\}$ . As  $|\hat{B}| = m = 2$ ,  $\hat{x}$  is nondegenerate and we obtain  $B = \hat{B}$  and  $N = \{2, 4\}$ . Furthermore, we obtain

$$(A_B)^{-1} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}^{-1} = \begin{bmatrix} 0 & 1 \\ 1 & -1 \end{bmatrix}$$

and

$$\begin{aligned} \bar{c}_2 &= c_2 - c_B^T (A_B)^{-1} A^2 = -4 - [-1, -1] \begin{bmatrix} 0 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 4 \\ 2 \end{bmatrix} = 0 \\ \bar{c}_4 &= c_4 - c_B^T (A_B)^{-1} A^4 = 2 - [-1, -1] \begin{bmatrix} 0 & 1 \\ 1 & -1 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = 2 \end{aligned}$$

Since  $\bar{c} \geq \mathbf{0}$ , we conclude that  $\hat{x}$  is an optimal nondegenerate vertex by Corollary 15.4 and  $\hat{d} = \mathbf{0}$ .

- (b) We only have  $\hat{x}_2 > 0$  and, hence,  $\hat{B} = \{2\}$ . As  $|\hat{B}| = 1 < m = 2$ ,  $\hat{x}$  is a degenerate vertex. There are now three possibilities to extend  $\hat{B}$  to a basis:  $B = \{1, 2\}$ ,  $B = \{2, 3\}$ , and  $B = \{2, 4\}$ .

- $B = \{1, 2\}$  and  $N = \{3, 4\}$ : We obtain

$$(A_B)^{-1} = \begin{bmatrix} 1 & 4 \\ 1 & 2 \end{bmatrix}^{-1} = \frac{1}{2} \begin{bmatrix} -2 & 4 \\ 1 & -1 \end{bmatrix}$$

and

$$\begin{aligned} \bar{c}_3 &= c_3 - c_B^T (A_B)^{-1} A^3 = -1 - [-1, -4] \begin{bmatrix} -1 & 2 \\ 1/2 & -1/2 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = 0 \\ \bar{c}_4 &= c_4 - c_B^T (A_B)^{-1} A^4 = 2 - [-1, -4] \begin{bmatrix} -1 & 2 \\ 1/2 & -1/2 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = 2 \end{aligned}$$

Since  $\bar{c} \geq \mathbf{0}$ , we conclude that  $\hat{x}$  is an optimal vertex by Corollary 15.4 and  $d = \mathbf{0}$ .

- $B = \{2, 3\}$  and  $N = \{1, 4\}$ : We obtain

$$(A_B)^{-1} = \begin{bmatrix} 4 & 1 \\ 2 & 0 \end{bmatrix}^{-1} = \begin{bmatrix} 0 & 1/2 \\ 1 & -2 \end{bmatrix}$$

and

$$\begin{aligned} \bar{c}_1 &= c_1 - c_B^T (A_B)^{-1} A^1 = -1 - [-4, -1] \begin{bmatrix} 0 & 1/2 \\ 1 & -2 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = 0 \\ \bar{c}_4 &= c_4 - c_B^T (A_B)^{-1} A^4 = 2 - [-4, -1] \begin{bmatrix} 0 & 1/2 \\ 1 & -2 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = 2 \end{aligned}$$

Since  $\bar{c} \geq \mathbf{0}$ , we conclude that  $\hat{x}$  is an optimal vertex by Corollary 15.4 and  $d = \mathbf{0}$ .

- $B = \{2, 4\}$  and  $N = \{1, 3\}$ : We obtain

$$(A_B)^{-1} = \begin{bmatrix} 4 & 0 \\ 2 & 1 \end{bmatrix}^{-1} = \begin{bmatrix} 1/4 & 0 \\ -1/2 & 1 \end{bmatrix}$$

and

$$\begin{aligned}\bar{c}_1 &= c_1 - c_B^T(A_B)^{-1}A^1 = -1 - [-4, 2] \begin{bmatrix} 1/4 & 0 \\ -1/2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = -1 \\ \bar{c}_3 &= c_3 - c_B^T(A_B)^{-1}A^3 = -1 - [-4, 2] \begin{bmatrix} 1/4 & 0 \\ -1/2 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = 1\end{aligned}$$

Since  $\bar{c} \not\geq \mathbf{0}$  and  $\hat{x}$  is a degenerate vertex, we cannot infer any information about the optimality of  $\hat{x}$ . The candidate improving direction is obtained by setting  $d_1 = 1$  and  $d_3 = 0$ , and

$$d_B = -(A_B)^{-1}A_N d_N = - \begin{bmatrix} 1/4 & 0 \\ -1/2 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} -1/4 \\ -1/2 \end{bmatrix}$$

Therefore,

$$d = [1, -1/4, 0, -1/2]^T.$$

Note that  $d$  is an improving direction since  $c^T d = [-1, -4, -1, 2]^T [1, -1/4, 0, -1/2]^T = -1 = \bar{c}_1 < 0$ . Let us check if  $d$  is indeed a feasible direction at  $\hat{x}$ . Note that  $A(\hat{x} + \lambda d) = b$  for any  $\lambda \in \mathbb{R}$  since  $Ad = \mathbf{0}$ . Consider  $\hat{x} + \lambda d = [0, 2, 0, 0]^T + \lambda[1, -1/4, 0, -1/2]^T = [\lambda, 2 - (1/4)\lambda, 0, (-1/2)\lambda]^T$ . Note that there does not exist a real number  $\lambda^* > 0$  such that  $\hat{x} + \lambda d \geq \mathbf{0}$  if  $\lambda \in [0, \lambda^*]$  since the fourth component becomes negative whenever  $\lambda > 0$ . Therefore,  $d$  is not a feasible direction at  $\hat{x}$ . Note that we have  $d \in \mathcal{D} \setminus \hat{\mathcal{D}}$ .

- (c) We have  $\hat{x}_3 > 0$  and  $\hat{x}_4 > 0$  and, hence,  $\hat{B} = \{3, 4\}$ . As  $|\hat{B}| = m = 2$ ,  $\hat{x}$  is nondegenerate and we obtain  $B = \hat{B}$  and  $N = \{1, 2\}$ . Furthermore, we obtain

$$(A_B)^{-1} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}^{-1} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

and

$$\begin{aligned}\bar{c}_1 &= c_1 - c_B^T(A_B)^{-1}A^1 = -1 - [-1, 2] \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = -2 \\ \bar{c}_2 &= c_2 - c_B^T(A_B)^{-1}A^2 = -4 - [-1, 2] \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 4 \\ 2 \end{bmatrix} = -4\end{aligned}$$

Since  $\bar{c} \not\geq \mathbf{0}$  and  $\hat{x}$  is a nondegenerate vertex, we can conclude that  $\hat{x}$  is not optimal by Proposition 17.1. The candidate improving direction is obtained by setting  $d_1 = 0$  and  $d_2 = 1$ , and

$$d_B = -(A_B)^{-1}A_N d_N = - \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & 4 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} 0 \\ 1 \end{bmatrix} = \begin{bmatrix} -4 \\ -2 \end{bmatrix}$$

Therefore,

$$d = [0, 1, -4, -2]^T.$$

Note that  $d$  is an improving direction since  $c^T d = [-1, -4, -1, 2]^T [0, 1, -4, -2]^T = -4 = \bar{c}_2 < 0$ . Let us check if  $d$  is indeed a feasible direction at  $\hat{x}$ . Note that  $A(\hat{x} + \lambda d) = b$  for any  $\lambda \in \mathbb{R}$  since  $Ad = \mathbf{0}$ . Consider  $\hat{x} + \lambda d = [0, 0, 8, 4]^T + \lambda[0, 1, -4, -2]^T = [0, \lambda, 8 - 4\lambda, 4 - 2\lambda]^T$ . Note that letting  $\lambda^* = 2 > 0$ , we obtain  $\hat{x} + \lambda d \geq \mathbf{0}$  if  $\lambda \in [0, \lambda^*]$ . Therefore,  $d$  is a feasible direction at  $\hat{x}$ . Note that we have  $d \in \hat{\mathcal{D}}$  since  $\mathcal{D} = \hat{\mathcal{D}}$  for a nondegenerate vertex by Section 17.4.1 in the lecture notes.

As illustrated by this example, a linear programming problem may have multiple vertices that are optimal (i.e., (a) and (b)) such that one of them is nondegenerate and another one is degenerate.

## Open Ended Problems

### 3 Feasible Directions and Optimality Conditions (2.5 marks)

Consider the following linear programming problem (P) in standard form:

$$(P) \quad \min\{c^T x : Ax = b, \quad x \geq \mathbf{0}\},$$

where  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ , and  $c \in \mathbb{R}^n$ .

Let  $\bar{x} \in \mathbb{R}^n$  be an optimal solution of (P) such that  $\bar{x}$  is not a vertex.

(3.1) Prove that there exists a feasible direction  $\bar{d} \in \mathbb{R}^n$  at  $\bar{x}$  such that  $\bar{d} \neq \mathbf{0}$  and  $c^T \bar{d} = 0$ .

[1.5 marks]

#### Solution

Note that (P) is a linear programming problem (P) in standard form. By the hypothesis, the optimal value is finite since  $\bar{x} \in \mathbb{R}^n$  is an optimal solution (i.e.,  $z^* = c^T \bar{x}$ ). Then, by the Fundamental Theorem of Linear Programming, there exists a basic feasible solution (i.e., a vertex), say  $x^*$ , which is also an optimal solution and we know that  $x^* - \bar{x} \neq 0$  since  $\bar{x}$  is not a vertex, and that  $z^* = c^T \bar{x} = c^T x^*$ . Let  $\bar{d} = x^* - \bar{x}$ . Then, since  $\bar{x}$  and  $x^* = \bar{x} + \bar{d}$  are feasible solutions,  $\bar{d}$  is clearly a feasible direction at  $\bar{x}$  (using  $\lambda^* = 1$ ). Now, observe that  $\bar{d} = x^* - \bar{x} \neq 0$  and  $c^T \bar{d} = c^T(x^* - \bar{x}) = c^T x^* - c^T \bar{x} = 0$ , which proves the claim.

(3.2) Prove that (P) has an infinite number of optimal solutions.

[1 mark]

#### Solution

From (3.1), we know that there exists a vertex  $x^*$  which is also an optimal solution and  $x^* \neq \bar{x}$  since  $\bar{x}$  is not a vertex. Since, the feasible region of a linear programming problem is always a polyhedron, and every polyhedron is a convex set by Remark 6.2,  $(\lambda x^* + (1 - \lambda)\bar{x})$  is also a feasible solution of (P) for all  $\lambda \in [0, 1]$ . Now we want to show that  $(\lambda x^* + (1 - \lambda)\bar{x})$  is an optimal solutions for each  $\lambda \in [0, 1]$ . By the following equality,

$$c^T(\lambda x^* + (1 - \lambda)\bar{x}) = \lambda c^T x^* + (1 - \lambda)c^T \bar{x} = \lambda z^* + (1 - \lambda)z^* = z^*.$$

Clearly  $\lambda x^* + (1 - \lambda)\bar{x}$  is also an optimal solution for any  $\lambda \in [0, 1]$ . Since we have infinitely many real numbers in any interval  $[a, b] \subseteq \mathbb{R}$  such that  $b > a$ , we conclude that we can find infinitely many  $\lambda$  values such that  $\lambda \in [0, 1]$ . Therefore, we have infinitely many optimal solutions.

### 4 Reduced Costs and Optimality Conditions (2.5 marks)

Consider the following linear programming problem in standard form:

$$(P) \quad \min\{c^T x : Ax = b, \quad x \geq \mathbf{0}\},$$

where  $A \in \mathbb{R}^{m \times n}$ ,  $b \in \mathbb{R}^m$ , and  $c \in \mathbb{R}^n$ . Assume that  $A$  has full row rank. Let  $x^* \in \mathbb{R}^n$  be a vertex with the corresponding index sets  $B$  and  $N$ .

(4.1) Suppose that  $\bar{x} \in \mathbb{R}^n$  is a feasible solution of (P) such that  $\bar{x} \neq x^*$ . Prove that there exists an index  $k \in N$  such that  $\bar{x}_k > 0$ .

[1 mark]

### Solution

Let  $x^* \in \mathbb{R}^n$  be a vertex of (P) with the corresponding index sets  $B$  and  $N$ . We have  $x_B^* = (A_B)^{-1}b \geq \mathbf{0}$  and  $x_N^* = \mathbf{0} \in \mathbb{R}^{n-m}$  by Proposition 13.2. Let  $\bar{x} \in \mathbb{R}^n$  be a feasible solution of (P) such that  $\bar{x} \neq x^*$ . Suppose, for a contradiction, that the conclusion is false. Then, we have  $\bar{x}_j = x_j^* = 0$  for each  $j \in N$ . Therefore using the same index sets  $B$  and  $N$  for the vertex  $x^*$ , we conclude that  $\bar{x}_N = \mathbf{0} \in \mathbb{R}^{n-m}$ . Since  $A\bar{x} = A_B\bar{x}_B + A_N\bar{x}_N = A_B\bar{x}_B = b$ , we obtain  $\bar{x}_B = (A_B)^{-1}b$ , which implies that  $\bar{x}_B = x_B^*$ . Since  $\bar{x}_B = x_B^*$  and  $\bar{x}_N = x_N^*$ , we then conclude that  $x^* = \bar{x}$ , which is a contradiction. Therefore, we conclude that there exists an index  $k \in N$  such that  $\bar{x}_k > 0$ .

- (4.2) Consider the vertex  $x^*$  again. Suppose that reduced costs of all nonbasic variables are strictly positive, i.e.,

$$\bar{c}_j = c_j - c_B^T(A_B)^{-1}A^j > 0, \quad j \in N.$$

Prove, using (4.1), that  $x^*$  is the unique optimal solution.

[1.5 marks]

### Solution

Let  $x^* \in \mathbb{R}^n$  be an optimal vertex with the corresponding index sets  $B$  and  $N$ . Suppose that reduced costs of all nonbasic variables are strictly positive.

Suppose, for a contradiction, that  $x^*$  is not the unique optimal solution. Then, there exists a feasible solution  $\bar{x} \in \mathbb{R}^n$  such that  $\bar{x} \neq x^*$  and  $\bar{x}$  is also an optimal solution. In the following, we want to derive a contradiction by considering the vector  $d = \bar{x} - x^*$ .

By (4.1), there exists an index  $k \in N$  such that  $\bar{x}_k > 0$  whereas  $x_k^* = 0$  since  $k \in N$ . Therefore,  $d_k = \bar{x}_k - x_k^* = \bar{x}_k > 0$ .

Note that  $Ad = A(\bar{x} - x^*) = b - b = \mathbf{0}$ , i.e.  $A_B d_B + A_N d_N = \mathbf{0}$ , which implies that  $d_B = -(A_B)^{-1}A_N d_N$ . Since  $c^T \bar{x} = c^T x^*$ , we obtain

$$0 = c^T d = c_B^T d_B + c_N^T d_N = -c_B^T(A_B)^{-1}A_N d_N + c_N^T d_N = \sum_{j \in N} \bar{c}_j d_j > 0,$$

where the strict inequality follows from  $d_j = \bar{x}_j - x_j^* = \bar{x}_j \geq 0$  and  $\bar{c}_j > 0$  for each  $j \in N$ , and the observation that  $d_k = \bar{x}_k > 0$  and  $\bar{c}_k > 0$ . We obtain a contradiction. Therefore, we conclude that  $x^*$  is the unique optimal solution.