

SDPNAL+: A MATLAB software package for large-scale SDPs with a user-friendly interface

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- SDP and SDP+ (variable is positive semidefinite and bounded)
- Some examples of SDP+
- User-friendly interface
- **Phase I:** An inexact symmetric Gauss-Seidel (sGS) ADMM for SDP+
- An **sGS decomposition theorem** for convex composite QP
- **Phase II:** An augmented Lagrangian method (ALM) for SDP+
- A semismooth Newton-CG (SNCG) method for solving ALM sub-problems
- SDPNAL+: practical implementation of the 2 phase method
- Numerical experiments

\mathbb{S}_+^n = cone of positive semidefinite matrices. Write $X \succeq 0$ if $X \in \mathbb{S}_+^n$.

$$(\text{SDP}) \quad \min \{ \langle C, X \rangle \mid \mathcal{A}(X) = b, X \in \mathbb{S}_+^n \}$$

where $C \in \mathbb{S}^n$, $b \in \mathbb{R}^m$ are given data; $\mathcal{A} : \mathbb{S}^n \rightarrow \mathbb{R}^m$ is a linear map.

$$(\text{SDP+}) \quad \min \{ \langle C, X \rangle \mid \mathcal{A}(X) = b, X \in \mathbb{S}_+^n, X \in \mathcal{N} \}$$

where $\mathcal{N} = \{X \in \mathbb{S}^n \mid L \leq X \leq U\}$ and L, U are given bounds (entries allow to take $-\infty, \infty$ respectively).

Important case: $\mathcal{N} = \{X \in \mathbb{S}^n \mid X \geq 0\}$, i.e., DNN (doubly nonnegative) SDP.

(SDP) is solvable by powerful interior-point methods if n and m are not too large, say, $n \leq 2000$, $m \leq 10,000$.

m large $\Rightarrow m \times m$ dense “Hessian” matrix cannot be stored explicitly. For $m = 10^5$, needs 100GB RAM memory!

Current research interests focus on $n \leq 5000$ but $m \gg 10,000$.

SDPNAL was developed around 2008/09 for (SDP).

In 2012/13, it was extended to SDPNAL+ for (SDP+) directly without introducing extra equality constraints $X = Y$ to convert $X \in \mathbb{S}_+^n \cap \mathcal{N}$ to $X \in \mathbb{S}_+^n$ and $Y \in \mathcal{N}$.

Now our solver SDPNAL+ can solve general SDP problems:

$$\begin{aligned}
 (\text{genSDP}) \quad & \min \quad \sum_{i=1}^N \langle C_i, X_i \rangle \\
 \text{s.t.} \quad & \sum_{i=1}^N \mathcal{A}_i(X_i) = b \quad (\text{equalities}) \\
 & l \leq \sum_{i=1}^N \mathcal{B}_i(X_i) \leq u \quad (\text{inequalities}) \\
 & X_i \in \mathbb{K}_i \quad (\text{cone}), \quad X_i \in \mathcal{N}_i \quad (\text{bounds}), \quad i = 1 : N
 \end{aligned}$$

where \mathbb{K}_i is either a PSD cone or nonnegative orthant. Currently extending \mathbb{K}_i to other cones such as SOCP.

- Parallel IPM [Benson, Borchers, Fujisawa, ... 03-present]
- First-order gradient methods on NLP formulation (low accuracy) [Burer-Monteiro 03]
- Inexact IPM [Kojima, Toh 04]
- Gen. Lag. method on barrier-penalized dual [Kocvara-Stingl 03]
- ALM on primal SDP from relaxation of lift-and-project scheme [Burer-Vandenbussche 06]
- Boundary-point method: BCD-ALM on dual [Rendl et al. 06]
Reg. methods for SDP \equiv ADMM on dual [Malick-Povh-Rendl 09]
- **SDPNAL**: ADMM+SNCG-ALM on dual [Zhao-Sun-Toh 10]
- SDPAD: ADMM on dual [Wen et al. 10] (used SDPNAL template)
- 2EBD: hybrid proximal extra-gradient method on primal [Monteiro et al. 13] (used SDPNAL template)
- **ADMM+**: convergent sGS-ADMM on SDP+ [Sun-Toh-Yang 15]
- **SDPNAL+**: SNCG-ALM on SDP+ [Yang-Sun-Toh 15]

In **nearest correlation matrix problem**, given data matrix $U \in \mathbb{S}^n$, we want to solve

$$(NCM) \quad \min_X \left\{ \frac{1}{2} \|H \circ (X - U)\|_1 \mid \text{Diag}(X) = \mathbf{1}, X \succeq 0 \right\} \quad \text{► NCM}$$

where $H \in \mathbb{S}^n$ has nonnegative entries and “ \circ ” is the Hadamard product.

In **clustering**, given data vectors $\{p_i\}_{i=1}^n$, the goal is to cluster them into k clusters. A possible model [Peng-Wei 07] is:

$$\min \left\{ \langle D, X \rangle \mid \langle I, X \rangle = k, X\mathbf{1} = \mathbf{1}, X \in \mathbb{S}_+^n, X \geq 0 \right\} \quad \text{► Clustering}$$

where $D_{ij} = \|p_i - p_j\|^2$.

Note: D can also be other affinity matrix.

Maximum stable set problem a graph $G = (V, \mathcal{E})$

A stable set S is subset of V such that no vertices in S are adjacent.

Maximum stable set problem: find S with maximum cardinality. Let

$$x_i = \begin{cases} 1 & \text{if } i \in S \\ 0 & \text{otherwise} \end{cases} \quad \Rightarrow \quad |S| = \sum_{i=1}^n x_i.$$

A common formulation of the max-stable-set problem:

$$\begin{aligned} \alpha(G) &:= \max \left\{ |S| = \frac{1}{|S|} \sum_{ij} x_i x_j \mid x_i x_j = 0 \forall (i, j) \in \mathcal{E}, x \in \{0, 1\}^n \right\} \\ &\quad \Downarrow \quad X := xx^T / |S| \\ &\max \left\{ \langle E, X \rangle \mid X_{ij} = 0 \forall (i, j) \in \mathcal{E}, \langle I, X \rangle = 1 \right\} \end{aligned}$$

SDP relaxation: $X = xx^T / |S| \Rightarrow X \succeq 0$, get

$$\theta(G) := \max \left\{ \langle E, X \rangle : X_{ij} = 0 \forall (i, j) \in \mathcal{E}, \langle I, X \rangle = 1, X \succeq 0 \right\}$$

$$\theta_+(G) := n(n+1)/2 \text{ additional constraints } X \geq 0 \quad \text{► theta}$$

Assign n facilities to n locations [Koopmans and Beckmann (1957)]

$A = (a_{ij})$ where a_{ij} = flow from facility i to facility j

$B = (b_{kl})$ where b_{kl} = distance from location k to location l

cost of assignment $\pi = \sum_{i=1}^n \sum_{j=1}^n a_{ij} b_{\pi(i)\pi(j)}$

$$\min_P \left\{ \langle B \otimes A, \text{vec}(P)\text{vec}(P)^T \rangle \mid P \text{ is } n \times n \text{ permutation matrix} \right\}$$

SDP+ relaxation [Povh and Rendl, 09]:

relax $\text{vec}(P)\text{vec}(P)^T$ to the $n^2 \times n^2$ variable $X \in \mathbb{S}_+^{n^2}$ and $X \geq 0$

$$(\text{QAP}) \min \left\{ \langle B \otimes A, X \rangle \mid \mathcal{A}(X) - b = 0, X \in \mathbb{S}_+^{n^2}, X \geq 0 \right\}$$

where the linear constraints (with $m = 3n(n+1)/2$) encode the condition $P^T P = I_n$, $P \geq 0$.

Consider the NCM problem. ► NCM

```
n = 100;
G = randn(n,n);
G = 0.5*(G + G');

model = ccp_model('NCM');
X = var_sdp(n,n);
model.add_variable(X);
model.minimize(l1_norm(X-G));
model.add_affine_constraint(map_diag(X)==ones(n,1));
model.solve;
```

Consider the $\theta+$ problem of a graph with adjacency matrix G .

► theta

```
n = 200;
G = triu(sprand(n,n,0.5),1);
[IE,JE] = find(G);
n = length(G);

model = ccp_model('theta');
X = var_sdp(n,n);
model.add_variable(X);
model.maximize(sum(X));
model.add_affine_constraint(trace(X) == 1);
model.add_affine_constraint(X(IE,JE) == 0);
model.add_affine_constraint(X >= 0);
model.solve;
```

$$\begin{aligned}
\min \quad & \text{trace}(X^{(1)}) + \text{trace}(X^{(2)}) + \text{sum}(X^{(3)}) \\
\text{s.t.} \quad & -X_{12}^{(1)} + 2X_{33}^{(2)} + 2X_2^{(3)} = 4 \quad (\text{equalities}) \\
& 2X_{23}^{(1)} + X_{42}^{(2)} - X_4^{(3)} = 3 \\
& 2 \leq -X_{12}^{(1)} - 2X_{33}^{(2)} + 2X_2^{(3)} \leq 7 \quad (\text{inequalities}) \\
& X^{(1)} \in \mathbb{S}_+^6, X^{(2)} \in \mathbb{R}^{5 \times 5}, X^{(3)} \in \mathbb{R}_+^7 \quad (\text{cones}) \\
& 0 \leq X^{(1)} \leq 10E_6, \quad 0 \leq X^{(2)} \leq 8E_5 \quad (\text{bounds})
\end{aligned}$$

```

n1 = 6; n2 = 5; n3 = 7;
M = ccp_model('Example_simple');
X1=var_sdp(n1,n1); X2=var_nn(n2,n2); X3=var_nn(n3);
M.add_variable(X1,X2,X3);
M.minimize(trace(X1) + trace(X2) + sum(X3));
M.add_affine_constraint(-X1(1,2)+2*X2(3,3)+2*X3(2)==4);
M.add_affine_constraint(2*X1(2,3)+X2(4,2)-X3(4) == 3);
M.add_affine_constraint(2<=-X1(1,2)-2*X2(3,3)+2*X3(2)<=7);
M.add_affine_constraint(0 <= X1 <= 10);
M.add_affine_constraint(X2 <= 8);
M.solve;

```

For simplicity, consider only $\mathcal{N} = \{X \in \mathbb{S}^n \mid X \geq 0\}$.

Dual of SDP+ and its augmented Lagrangian function are given by:

$$(D) \quad \min\{-\langle b, y \rangle + \delta_{\mathbb{S}_+^n}(S) + \delta_{\mathcal{N}}(Z) \mid \mathcal{A}^*y + S + Z = C\}$$

(a linearly constrained convex problem with 3 blocks of variables);

$$\begin{aligned} \mathcal{L}_\sigma(y, S, Z; X) &= -\langle b, y \rangle + \langle \mathcal{A}^*y + S + Z - C, X \rangle \\ &\quad + \frac{\sigma}{2} \|\mathcal{A}^*y + S + Z - C\|^2 + \delta_{\mathbb{S}_+^n}(S) + \delta_{\mathcal{N}}(Z) \\ &\quad \text{(quadratic in } (y, S, Z) \text{ + nonsmooth terms in } S, Z) \end{aligned}$$

KKT conditions:

$$\mathcal{R}_{\text{KKT}}(y, S, Z; X) := \begin{pmatrix} AX - b \\ S - \Pi_{\mathbb{S}_+^n}(S - X) \\ Z - \Pi_{\mathcal{N}}(Z - X) \\ \mathcal{A}^*y + S + Z - C \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \\ 0 \end{pmatrix}.$$

Input $(y^0, S^0, Z^0; X^0)$. For $k = 0, 1, \dots$, let $\hat{C}^k = C - \sigma^{-1} X^k$

$$(1a) \ y^{k+1} = \operatorname{argmin}_{y \in \mathbb{R}^m} \mathcal{L}_\sigma(y, S^k, Z^k; X^k)$$

$$(1b) \ S^{k+1} = \operatorname{argmin}_{S \in \mathbb{S}_+^n} \mathcal{L}_\sigma(y^{k+1}, S, Z^k; X^k) = \Pi_{\mathbb{S}_+^n}(\hat{C}^k - \mathcal{A}^* y^{k+1} - Z^k)$$

$$(2) \ Z^{k+1} = \operatorname{argmin}_{Z \in \mathcal{N}} \mathcal{L}_\sigma(y^{k+1}, S^{k+1}, Z; X^k) = \Pi_{\mathcal{N}}(\hat{C}^k - \mathcal{A}^* y^{k+1} - S^{k+1})$$

(3) $X^{k+1} = X^k + \tau \sigma (\mathcal{A}^* y^{k+1} + S^{k+1} + Z^{k+1} - C)$, where $\tau \in (0, \frac{1+\sqrt{5}}{2})$ is the step-length.

Direct extension of 2-block ADMM is not guaranteed to converge
[Chen-He-Ye-Yuan, v155, MP 2016]

But sGS-ADMM is guaranteed to converge!

Input $(y^0, S^0, Z^0; X^0)$. For $k = 0, 1, \dots$, let $\hat{C}^k = C - \sigma^{-1}X^k$

$$(1a) \hat{y}^{k+1} \approx \operatorname{argmin}_{y \in \mathbb{R}^m} \mathcal{L}_\sigma(y, S^k, Z^k; X^k)$$

$$(1b) S^{k+1} = \operatorname{argmin}_{S \in \mathbb{S}_+^n} \mathcal{L}_\sigma(\hat{y}^{k+1}, S, Z^k; X^k) = \Pi_{\mathbb{S}_+^n}(\hat{C}^k - \mathcal{A}^* \hat{y}^{k+1} - Z^k)$$

$$(1c) \boxed{y^{k+1} \approx \operatorname{argmin}_{y \in \mathbb{R}^m} \mathcal{L}_\sigma(y, S^{k+1}, Z^k; X^k)}$$

$$(2) Z^{k+1} = \operatorname{argmin}_{Z \in \mathcal{N}} \mathcal{L}_\sigma(y^{k+1}, S^{k+1}, Z; X^k) = \Pi_{\mathcal{N}}(\hat{C}^k - \mathcal{A}^* y^{k+1} - S^{k+1})$$

$$(3) X^{k+1} = X^k + \tau\sigma(\mathcal{A}^* y^{k+1} + S^{k+1} + Z^{k+1} - C)$$

In Step 1, the AL function \mathcal{L}_σ for the block (y, S) has the form:

$$\mathcal{L}_\sigma(y, S, Z^k; X^k) \equiv \delta_{\mathbb{S}_+^n}(S) + \frac{\sigma}{2} \|\mathcal{A}^* y + S + Z^k + \hat{C}^k\|^2 - \langle b, y \rangle$$

(QP in (y, S) + nonsmooth term in S)

(1a)–(1c) is equivalent to minimizing $\mathcal{L}_\sigma(y, S) + \text{sGS proximal term}$.
The steps are based on a sGS decomposition theorem.

Theorem Suppose the KKT conditions of (SDP+) has a solution. Let $\{(y^k, S^k, Z^k, X^k)\}$ be the sequence generated by the inexact sGS-ADMM. Then $\{X^k\}$ converges to an optimal solution of (SDP+) and $\{(y^k, S^k, Z^k)\}$ converges to an optimal solution of its dual.

- [1] D.F. Sun, K.C. Toh and L.Q. Yang, [A convergent 3-block semi-proximal ADMM for conic programming with 4-type constraints](#), v25, SIOPT 2015.
- [2] X.D. Li, D.F. Sun, K.C. Toh, [A Schur complement based semiproximal ADMM for convex ...](#), v155, MP 2016. [Schur-complement-ADMM](#)
- [3] X.D. Li, D.F. Sun, K.C. Toh, [QSDPNAL: A two-phase augmented Lagrangian method for convex quadratic SDP](#), arXiv:1512.08872v1, 2015. [Section 2: sGS decomposition theorem, Schur-complement-ADMM = sGS-ADMM](#)
- [4] L. Chen, D.F. Sun, K.C. Toh, [An efficient inexact symmetric Gauss-Seidel based majorized ADMM for ...](#), v161, MP 2017. [inexact sGS-ADMM](#)
- [5] X.D. Li, D.F. Sun, K.C. Toh, [A block sGS decomposition theorem for convex composite quadratic programming and its applications](#), arXiv:1703.06629, 2017. [sGS-ADMM = Schur-complement-ADMM, sSOR-extension](#)

Theorem [Han-Sun-Zhang, MOR 2017: exact version]

Let $\Omega_{\text{KKT}} \neq \emptyset$ be the KKT solution set. Suppose an error bound condition holds for \mathcal{R}_{KKT} at an optimal solution $u^* = (y^*, S^*, Z^*, X^*)$, i.e., $\exists \eta, r > 0$ s.t.

$$\text{dist}(u, \Omega_{\text{KKT}}) \leq \eta \|\mathcal{R}_{\text{KKT}}(u)\| \quad \forall u \in B_r(u^*).$$

Let $u^k = (y^k, S^k, Z^k, X^k)$. Then $\exists \mu \in (0, 1)$ depending on η s.t.

$$\text{dist}(u^{k+1}, \Omega_{\text{KKT}}) \leq \mu \text{dist}(u^k, \Omega_{\text{KKT}}) \quad \forall k \text{ sufficiently large.}$$

Inexact version can be established via the analysis in [Chen-Sun-Toh, MP 2017] and [Han-Sun-Zhang, MOR 2017].

Consider a convex composite QP with 3 blocks:

$$\min \left\{ p(\mathbf{x}_1) + h(x) \mid x = (\mathbf{x}_1; x_2, x_3) \in \mathbb{R}^{n_1} \times \mathbb{R}^{n_2} \times \mathbb{R}^{n_3} \right\}$$

Convex quadratic function $h(x) := \frac{1}{2} \langle x, \mathcal{H}x \rangle - \langle b, x \rangle$

Closed proper convex fun. $p : \mathbb{R}^{n_1} \rightarrow (-\infty, +\infty]$, eg $p(x_1) = \|x_1\|_\infty$

Write $\mathcal{H} = \mathcal{U}^* + \mathcal{D} + \mathcal{U}$, \mathcal{D} diagonal blocks, \mathcal{U} strict upper triangular part. Assume \mathcal{D} invertible.

Define $\text{sGS}(\mathcal{H}) := \mathcal{U}\mathcal{D}^{-1}\mathcal{U}^*$ (symmetric Gauss-Seidel decomp.)

Given \bar{x} , define

$$x^+ := \operatorname{argmin}_x \left\{ p(\mathbf{x}_1) + h(x) + \frac{1}{2} \|x - \bar{x}\|_{\text{sGS}(\mathcal{H})}^2 \right\}$$

Next theorem: can compute x^+ using one sGS cycle!

If $p(\mathbf{x}_1)$ is absent, we get the classical block sGS iteration.

Theorem [Li-Sun-Toh 2015]

It holds that $\mathcal{H} + \text{sGS}(\mathcal{H}) = (\mathcal{D} + \mathcal{U})\mathcal{D}^{-1}(\mathcal{D} + \mathcal{U}^*) \succ 0$.

Backward GS: $3 \rightarrow 2$. Compute

$$x'_3 = \operatorname{argmin} p(\bar{x}_1) + h(\bar{x}_1, \bar{x}_2, x_3) = \mathcal{H}_{33}^{-1}(b_3 - \mathcal{H}_{13}^* \bar{x}_1 - \mathcal{H}_{23}^* \bar{x}_2)$$

$$x'_2 = \operatorname{argmin} p(\bar{x}_1) + h(\bar{x}_1, x_2, x'_3) = \mathcal{H}_{22}^{-1}(b_2 - \mathcal{H}_{12}^* \bar{x}_1 - \mathcal{H}_{23} x'_3)$$

Forward GS: $1 \rightarrow 2 \rightarrow 3$. Compute

$$x_1^+ = \operatorname{argmin} p(x_1) + h(x_1, x'_2, x'_3) \quad (\text{non-smooth/non-quadratic})$$

$$x_2^+ = \operatorname{argmin} p(x_1^+) + h(x_1^+, x_2, x'_3) = \mathcal{H}_{22}^{-1}(b_2 - \mathcal{H}_{12}^* x_1^+ - \mathcal{H}_{23} x'_3)$$

$$x_3^+ = \operatorname{argmin} p(x_1^+) + h(x_1^+, x_2^+, x_3) = \mathcal{H}_{33}^{-1}(b_3 - \mathcal{H}_{13}^* x_1^+ - \mathcal{H}_{23}^* x_2^+)$$

Inexact computation is also allowed! So can use PCG to solve large linear systems.

Theorem [Li-Sun-Toh 2015]

Backward GS: For $i = s, \dots, 2$, compute

$$x'_i = \mathcal{H}_{ii}^{-1} \left(b_i + e'_i - \sum_{j=1}^{i-1} \mathcal{H}_{ji}^* \bar{x}_j - \sum_{j=i+1}^s \mathcal{H}_{ij} x'_j \right).$$

Forward GS: For $i = 2, \dots, s$

$$x_1^+ = \operatorname{argmin} p(x_1) + h(x_1, x'_{\geq 2}) - \langle e_1^+, x_1 \rangle,$$

$$x_i^+ = \mathcal{H}_{ii}^{-1} \left(b_i + e_i^+ - \sum_{j=1}^{i-1} \mathcal{H}_{ji}^* x_j^+ - \sum_{j=i+1}^s \mathcal{H}_{ij} x'_j \right)$$

e^+, e' are error vectors. In this case, x^+ is the exact solution to a slightly perturbed proximal problem:

$$x^+ := \operatorname{argmin}_x \left\{ p(x_1) + h(x) + \frac{1}{2} \|x - \bar{x}\|_{\text{sGS}(\mathcal{H})}^2 - \langle x, \Delta(e', e^+) \rangle \right\}$$

$$\Delta(e', e^+) = e^+ + \mathcal{U} \mathcal{D}^{-1} (e^+ - e').$$

Adding a large proximal term slows the convergence of sGS-ADMM!

With no proximal term added, we consider the ALM for solving dual SDP+.

(1) Compute

$$\begin{aligned}(y^{k+1}, S^{k+1}, Z^{k+1}) &\approx \operatorname{argmin} \left\{ \mathcal{L}_k(y, S, Z) := \mathcal{L}_{\sigma_k}(y, S, Z; X^k) \right\} \\ &= \operatorname{argmin} \left\{ -\langle b, y \rangle + \frac{\sigma}{2} \|\mathcal{A}^* y + S + Z + \widehat{C}^k\|^2 + \delta_{\mathbb{S}_+^n}(S) + \delta_{\mathcal{N}}(Z) \right\}\end{aligned}$$

(2) Update $X^{k+1} = X^k + \sigma_k(\mathcal{A}^* y^{k+1} + S^{k+1} + Z^{k+1} - C)$;
update $\sigma_{k+1} \uparrow \sigma_\infty \leq \infty$.

Define $X^{k+1} = X^k + \sigma_k R_D(y^{k+1}, S^{k+1}, Z^{k+1})$,

$$e^{k+1} = \begin{bmatrix} \mathcal{A}X^{k+1} - b \\ X^{k+1} - \Pi_{\mathbb{S}_+^n}(X^{k+1} - S^{k+1}) \\ X^{k+1} - \Pi_{\mathcal{N}^n}(X^{k+1} - Z^{k+1}) \end{bmatrix}.$$

In Step 1, we use the following **easy-to-check** stopping conditions:

$$(A) \quad \|e^{k+1}\| \leq \frac{\epsilon_k^2}{1 + \|(X, y, S, Z)^{k+1}\|} \min \left\{ \frac{1}{\sigma_k}, \frac{1}{1 + \|X^{k+1} - X^k\|} \right\}$$

$$(B) \quad \|e^{k+1}\| \leq \frac{\eta_k^2 \|X^{k+1} - X^k\|^2}{1 + \|(X, y, S, Z)^{k+1}\|} \min \left\{ \frac{1}{\sigma_k}, \frac{1}{1 + \|X^{k+1} - X^k\|} \right\}$$

where $\{\epsilon_k\}$ and $\{\delta_k\}$ are nonnegative summable sequences.

Theorem [Rockafellar 76] Let $\Omega_P \neq \emptyset$ be the primal optimal solution set and Slater's condition holds for primal problem (P). Under stopping condition (A), we have $X^k \rightarrow X^*$ and $(y^{k+1}, S^{k+1}, Z^{k+1})$ converges to a dual optimal solution.

Theorem [Cui-Sun-Toh] If in addition, the blue stopping conditions are added, and the essential primal objective function P^{obj} satisfies a quadratic growth condition at X^* , i.e., \exists a neighborhood \mathcal{U} of X^* and $\kappa > 0$ s.t.

$$P^{\text{obj}}(X) \geq P^{\text{obj}}(X^*) + \kappa^{-1} \text{dist}^2(X, \Omega_P) \quad \forall X \in \mathcal{U}$$

Then for k large, we have

$$\text{dist}(X^{k+1}, \Omega_P) \leq \theta_k \text{dist}(X^k, \Omega_P)$$

$$\text{dual feasibility at } (y^{k+1}, S^{k+1}, Z^{k+1}) \leq \tau_k \text{dist}(X^k, \Omega_P)$$

$$\text{dual objective gap at } (y^{k+1}, S^{k+1}, Z^{k+1}) \leq \tau'_k \text{dist}(X^k, \Omega_P)$$

$$\text{where } \theta_k \approx \frac{\kappa}{\sqrt{\kappa^2 + \sigma_k^2}}, \quad \tau_k \approx \frac{1}{\sigma_k}, \quad \tau'_k \approx \frac{\|X^k\| + \|X^{k+1}\|}{2\sigma_k}$$

Larger σ_k gives faster convergence, but inner problem is harder to solve.

For simplicity, assume $\mathcal{N} = \mathbb{S}^n$ and hence the variable Z is absent.

$$\begin{aligned} & \operatorname{argmin}_{y, S} \left\{ \mathcal{L}_\sigma(y, S) \equiv \delta_{\mathbb{S}_+^n}(S) + \frac{\sigma}{2} \|\mathcal{A}^* y + S - \widehat{C}^k\|^2 - \langle b, y \rangle \right\} \\ & \equiv \operatorname{argmin}_y \left\{ \Phi^k(y) := -\langle b, y \rangle + \frac{\sigma}{2} \|\Pi_{\mathbb{S}_+^n}(\mathcal{A}^* y - \widehat{C}^k)\|^2 \right\} \text{ (project out } S) \end{aligned}$$

Optimality condition of **unconstrained subproblem in y** is:

$$\nabla \Phi^k(y) = -b + \sigma \mathcal{A} \Pi_{\mathbb{S}_+^n}(\mathcal{A}^* y - \widehat{C}^k) = 0.$$

Solve for solution y^{k+1} by semismooth Newton-CG (SNCG) method. Then compute $S^{k+1} = \Pi_{\mathbb{S}_+^n}(\widehat{C}^k - \mathcal{A}^* y^{k+1})$.

$\nabla \Phi^k(y)$ is not differentiable, but is strongly semismooth [Sun-Sun, 2002]. Thus SNCG is expected to have at least superlinear convergence.

Solve $\nabla\Phi^k(y) = -b + \sigma\mathcal{A}\Pi_{\mathbb{S}_+^n}(U) = 0$, $U = \mathcal{A}^*y - \hat{C}^k$.

At the current iteration, y_l , we solve a generalized Newton equation:

$$\mathcal{H}\Delta y \approx \nabla\Phi^k(y_l), \quad \text{where } \mathcal{H}\Delta y = \sigma\mathcal{A}\Pi'_{\mathbb{S}_+^n}(U)[\mathcal{A}^*\Delta y] \quad (1)$$

From eigenvalue decomp: $U = QDQ^T$ with $d_1 \geq \dots \geq d_r \geq 0 > d_{r+1} \geq \dots \geq d_n$, we choose

$$\Pi'_{\mathbb{S}_+^n}(U)[M] = Q(\Omega \circ (Q^T M Q))Q^T, \quad (2)$$

$\Omega_{ij} = (d_i^+ - d_j^+)/(d_i - d_j)$. Let $\gamma = \{1, \dots, r\}$, $\bar{\gamma} = \{r+1, \dots, n\}$,

$$\Omega = \begin{bmatrix} E_{\gamma\gamma} & \Omega_{\gamma\bar{\gamma}} \\ \Omega_{\bar{\gamma}\gamma} & 0 \end{bmatrix}.$$

When problem is primal nondegenerate, $\text{cond}(\mathcal{H})$ is bounded:

$$\text{cond}(\mathcal{H}) \leq \sigma \Theta(1) \text{cond}([\mathcal{A}Q_\gamma \otimes Q_\gamma, \mathcal{A}Q_\gamma \otimes Q_{\bar{\gamma}}])^2$$

The structure in Ω allows for efficient computation of matrix-vector multiply for CG in solving (1). Direct evaluation of

$$Y := \Pi'_{\mathbb{S}^n_+}(U)[M] = Q(\Omega \circ (Q^T M Q))Q^T$$

needs 4 matrix-matrix multiplications = $8n^3$ operations. But with the structure of Ω , can compute Y as follows:

$$Y = H + H^T, \quad H = Q_\gamma \left[\frac{1}{2} (U Q_\gamma) Q_\gamma^T + (\Omega_{\gamma\bar{\gamma}} \circ (U Q_{\bar{\gamma}})) Q_{\bar{\gamma}}^T \right]$$

where $U = Q_\gamma M$. The cost is at most $6rn^2$.

If $r \approx n$, then use

$$\begin{aligned} Y &= Q(E \circ (Q^T M Q))Q^T - Q(\bar{\Omega} \circ (Q^T M Q))Q^T \\ &= M - Q(\bar{\Omega} \circ (Q^T M Q))Q^T \end{aligned}$$

where $\bar{\Omega} = E - \Omega$ has a similar structure as Ω but with a large block of 0. The cost is $6(n-r)n^2$.

Let ADMM+ denote the sGS-ADMM.

1. Generate a good starting point to **warm-start** SNCG-ALM:
 $(y^0, S^0, Z^0, X^0, \sigma_0) \leftarrow \text{ADMM+}(\bar{y}^0, \bar{S}^0, \bar{Z}^0, \bar{X}^0, \bar{\sigma}_0)$

2. For $k = 0, 1, \dots$

Generate $(y^{k+1}, S^{k+1}, Z^{k+1})$ in ALM-subproblem via SNCG

Compute X^{k+1} based on $(y^{k+1}, S^{k+1}, Z^{k+1})$, update σ_{k+1}

If progress of SNCG-ALM is slow,

Rescale data

Let $(\bar{y}^k, \bar{S}^k, \bar{Z}^k, \bar{X}^k, \bar{\sigma}_k)$ denote rescaled $(y^k, S^k, Z^k, X^k, \sigma_k)$

Rescaling is chosen such that $\|\bar{X}^k\| \approx \max\{\|\bar{S}^k\|, \|\bar{Z}^k\|\}$

Goto Step 1: Restart with ADMM+ $(\bar{y}^k, \bar{S}^k, \bar{Z}^k, \bar{X}^k, \bar{\sigma}_k)$

$$\eta \equiv \frac{\|\mathcal{R}_{\text{KKT}}(y^{k+1}, S^{k+1}, Z^{k+1}, X^{k+1})\|}{1 + \|(y^{k+1}, S^{k+1}, Z^{k+1}, X^{k+1})\|} \leq 10^{-6}.$$

Performance of our **SDPNAL+** and **ADMM+** versus

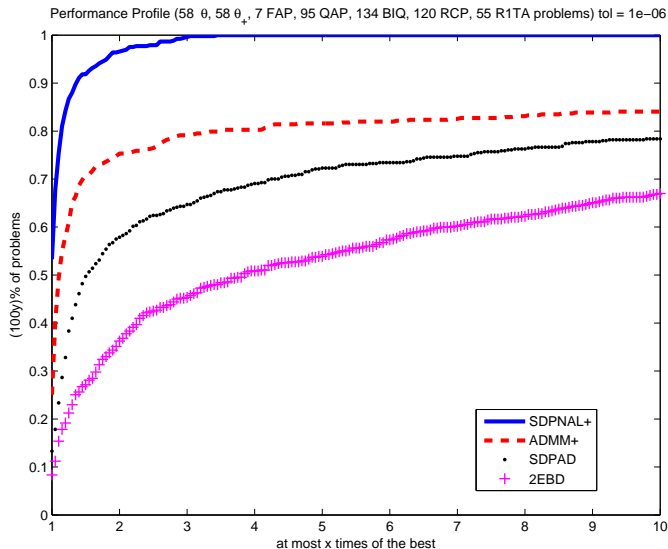
SDPAD: the directly extended ADMM implemented in [Wen et al.]

2EBD-HPE [Monteiro et al.]

Numbers of problems which are solved to the accuracy $\eta \leq 10^{-6}$

| problem set (No.) | SDPNAL+ | ADMM+ | SDPAD | 2EBD |
|--------------------|----------------|--------------|------------|------------|
| θ (58) | 58 | 56 | 53 | 53 |
| θ_+ (58) | 58 | 58 | 58 | 56 |
| FAP (7) | 7 | 7 | 7 | 7 |
| QAP (95) | 95 | 39 | 30 | 16 |
| BIQ (134) | 134 | 134 | 134 | 134 |
| RCP (120) | 120 | 120 | 114 | 109 |
| R1TA (55) | 55 | 42 | 47 | 18 |
| Total (527) | 527 | 456 | 443 | 393 |

Performance profiles of SDPNAL+, ADMM+, SDPAD and 2EBD



Implemented the algorithms in MATLAB.

Runs perform on PC with (12 cores) Intel Xeon CPU E5-2680 @ 2.50 GHz and 128 GB RAM.

Stop SDPAD and 2EBD after 25000 iterations or 20 hours.

| Prob | $m; n$ | η | | | time (hour:minute) |
|-------------------|-----------------------------|--------|-------|---------|---------------------|
| | | SDPAD | 2EBD | SDPNAL+ | |
| 1dc.2048 | 58368+ \mathcal{N} ; 2048 | 9.9-7 | 9.9-7 | 9.9-7 | 3:56 2:10 1:08 |
| fap25 | 2118+ \mathcal{N} ; 2118 | 9.9-7 | 9.9-7 | 9.5-7 | 3:26 0:54 0:43 |
| nug30 | 1393+ \mathcal{N} ; 900 | 1.1-5 | 1.7-5 | 9.6-7 | 2:10 1:46 0:09 |
| tai30a | 1393+ \mathcal{N} ; 900 | 4.6-6 | 1.3-5 | 9.9-7 | 2:34 1:47 0:10 |
| nsym_rd[40,40,40] | 672399; 1600 | 1.5-3 | 2.0-3 | 8.6-7 | 2:48 4:54 0:04 |
| nonsym(14,4) | 1.16M; 2744 | 1.4-2 | 5.2-3 | 1.3-7 | 7:39 14:01 0:20 |

Results show that it is essential to use **second-order information** and **second-order structured sparsity** to solve hard problems!

- We have tested SDPNAL+ on about 520 SDPs from θ, θ_+ , QAP, binary QP, rank-1 tensor approximation, etc
- When the problems are primal-dual nondegenerate, SDPNAL+ can efficiently solve large SDPs to high accuracy. SDPAD and 2EDB also performed well, though SDPNAL+ is often much more efficient.
- Many of the tested SDPs are degenerate, but SDPNAL+ can still solve them accurately with $\eta < 10^{-6}$. On the other hand, SDPAD and 2EDB were not able to solve many such problems.

Currently under development:

- ① sparse SDPNAL+ so as to handle larger matrix variable when the data has conducive sparsity structure
- ② a more advanced user-friendly interface

Thank you for your attention!