

Standard lattices of compatibly embedded finite fields

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JNCF 2019



CONTENTS

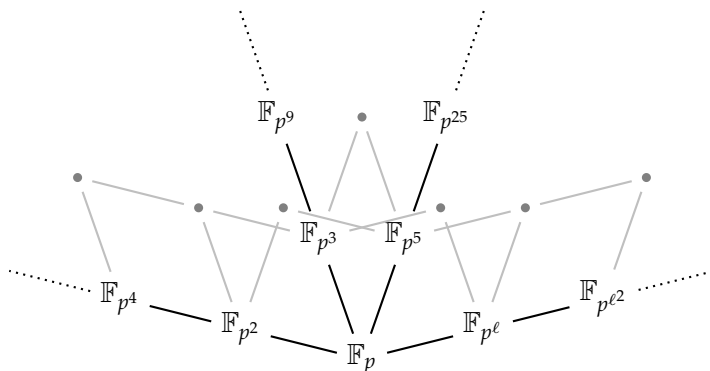
Context

Overview

Standard lattices

CONTEXT

- ▶ Use of Computer Algebra System (CAS)
- ▶ Use of many extensions of a prime finite field \mathbb{F}_p
- ▶ Computations in $\bar{\mathbb{F}}_p$.



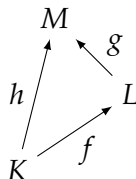
EMBEDDINGS

- ▶ When $l \mid m$, we know $\mathbb{F}_{p^l} \hookrightarrow \mathbb{F}_{p^m}$
 - ▶ How to compute this embedding *efficiently*?
- ▶ Naive algorithm: if $\mathbb{F}_{p^l} = \mathbb{F}_p[x]/(f(x))$, find a root ρ of f in \mathbb{F}_{p^m} and map \bar{x} to ρ . Complexity strictly larger than $\tilde{O}(l^2)$.
- ▶ Lots of other solutions in the literature:
 - ▶ [Lenstra '91]
 - ▶ [Allombert '02] $\tilde{O}(l^2)$
 - ▶ [Rains '96]
 - ▶ [Narayanan '18]

COMPATIBILITY

- ▶ K, L, M three finite fields with $K \hookrightarrow L \hookrightarrow M$
- ▶ $f : K \hookrightarrow L, g : L \hookrightarrow M, h : K \hookrightarrow M$ embeddings

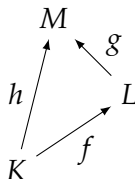
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Compatibility:



$$g \circ f \stackrel{?}{=} h$$

ENSURING COMPATIBILITY: CONWAY POLYNOMIALS

Definition (m -th Conway polynomials C_m)

- ▶ monic
- ▶ irreducible
- ▶ degree m
- ▶ primitive (i.e. its roots generate $\mathbb{F}_{p^m}^\times$)
- ▶ *norm-compatible* (i.e. $C_l \left(X^{\frac{p^m-1}{p^l-1}} = 0 \right) = 0 \pmod{C_m}$ if $l \mid m$)

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- ▶ Standard polynomials
- ▶ Compatible embeddings: $\bar{X} \mapsto \bar{Y}^{\frac{p^m-1}{p^l-1}} \tilde{O}(m^2)$

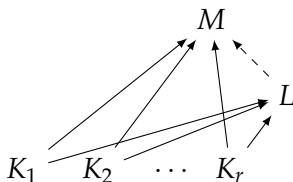
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- ▶ Standard polynomials
- ▶ Compatible embeddings: $\bar{X} \mapsto \bar{Y}^{\frac{p^m-1}{p^l-1}} \tilde{O}(m^2)$
- ▶ Hard to compute (exponential complexity)

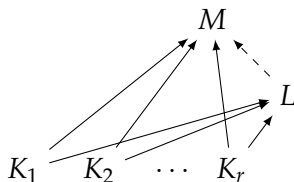
ENSURING COMPATIBILITY: BOSMA, CANNON AND STEEL

- ▶ Framework used in MAGMA
- ▶ Based on the naive embedding algorithm
- ▶ Constraints of the embedding imply that adding a new embedding can be expensive



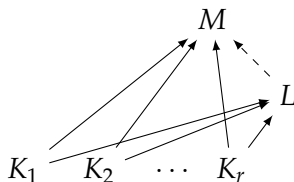
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- ▶ Non standard polynomials

IDEAS

- ▶ Plugging Allombert's embedding algorithm in Bosma, Cannon, and Steel
- ▶ Generalizing Bosma, Cannon, and Steel
- ▶ Generalizing Conway polynomials

Goal: bring the best of both worlds

ALLOMBERT'S EMBEDDING ALGORITHM I

- ▶ Based on an extension of *Kummer theory*
- ▶ For $p \nmid l$, we work in $A_l = \mathbb{F}_{p^l} \otimes \mathbb{F}_p(\zeta_l)$, and study

$$(\sigma \otimes 1)(x) = (1 \otimes \zeta_l)x \quad (\text{H90})$$

- ▶ Solutions of (H90) form a $\mathbb{F}_p(\zeta_l)$ -vector space of dimension 1
- ▶ $\alpha_l = \sum_{j=0}^{a-1} x_j \otimes \zeta_l^j$ solution of (H90), then x_0 generates \mathbb{F}_{p^l} .
 - ▶ Let $[\zeta_l] = x_0$ the projection on the first coordinate
- ▶ $(\alpha_l)^l = 1 \otimes c \in 1 \otimes \mathbb{F}_p(\zeta_l)$

ALLOMBERT'S EMBEDDING ALGORITHM II

Input: $\mathbb{F}_{p^l}, \mathbb{F}_{p^m}$, with $l \mid m$, ζ_l and ζ_m with $(\zeta_m)^{m/l} = \zeta_l$

Output: $s \in \mathbb{F}_{p^l}, t \in \mathbb{F}_{p^m}$, such that $s \mapsto t$ defines an embedding $\phi : \mathbb{F}_{p^l} \rightarrow \mathbb{F}_{p^m}$

1. Construct A_l and A_m
2. Find $\alpha_l \in A_l$ and $\alpha_m \in A_m$, nonzero solutions of (H90) for the roots ζ_l and ζ_m
3. Compute $(\alpha_l)^l = 1 \otimes c_l$ and $(\alpha_m)^m = 1 \otimes c_m$
4. Compute $\kappa_{l,m}$ a l -th root of c_l/c_m
5. Return $\lfloor \alpha_l \rfloor$ and $\lfloor (1 \otimes \kappa_{l,m})(\alpha_m)^{m/l} \rfloor$

ALLOMBERT AND BOSMA, CANON, AND STEEL

- ▶ Need to store one constant $\kappa_{l,m}$ for each pair $(\mathbb{F}_{p^l}, \mathbb{F}_{p^m})$
- ▶ The constant $\kappa_{l,m}$ depends on α_l and α_m

We would like to:

- ▶ get rid of the constants $\kappa_{l,m}$ (e.g. have $\kappa_{l,m} = 1$)
- ▶ equivalently, get "standard" solutions of (H90)
 - ▶ select solutions α_l, α_m that always define the same embedding
 - ▶ such that the constants $\kappa_{l,m}$ are well understood (e.g. $\kappa_{l,m} = 1$)

THE CASE $l \mid m \mid p - 1$

Let $l \mid m \mid p - 1$

- ▶ $A_l = \mathbb{F}_{p^l} \otimes \mathbb{F}_p \cong \mathbb{F}_{p^l}$
- ▶ $A_m = \mathbb{F}_{p^m}$
- ▶ $\sigma(\alpha_l) = \zeta_l \alpha_l$ and $\sigma(\alpha_m) = \zeta_m \alpha_m$
- ▶ $(\alpha_l)^l = c_l \in \mathbb{F}_p$ and $(\alpha_m)^m = c_m \in \mathbb{F}_p$
- ▶ $\kappa_{l,m} = \sqrt[l]{c_l/c_m}$
- ▶ $\kappa_{l,m} = 1$ implies $c_l = c_m$

In particular, for $m = p - 1$ we obtain

$$\sigma(\alpha_{p-1}) = (\alpha_{p-1})^p = \zeta_{p-1} \alpha_{p-1}$$

- ▶ $(\alpha_{p-1})^{p-1} = c_{p-1} = \zeta_{p-1}$
- ▶ this implies $\forall l \mid p - 1, c_l = \zeta_{p-1}$

COMPLETE ALGEBRA

Let $A_l = \mathbb{F}_{p^l} \otimes \mathbb{F}_p(\zeta_l)$

Definition (degree, level)

- ▶ *degree* of A_l : l
- ▶ *level* of A_l : $a = [\mathbb{F}_p(\zeta_l) : \mathbb{F}_p]$

Idea: consider the largest algebra for a given level

Definition (Complete algebra of level a)

- ▶ $A_{p^a-1} = \mathbb{F}_{p^{p^a-1}} \otimes \mathbb{F}_p(\zeta_{p^a-1}) \cong \mathbb{F}_{p^{p^a-1}} \otimes \mathbb{F}_{p^a}$

STANDARD SOLUTIONS

How to define **standard solutions** of (H90)?

Lemma

If α_{p^a-1} is a solution of (H90) for ζ_{p^a-1} , then $c_{p^a-1} = (\zeta_{p^a-1})^a$.

Definition (Standard solution)

Let A_l an algebra of level a , $\alpha_l \in A_l$ a solution of (H90) for

$\zeta_l = (\zeta_{p^a-1})^{\frac{p^a-1}{l}}$, α_l is **standard** if $c_l = (\zeta_{p^a-1})^a$

Definition (Standard polynomial)

All standard solutions α_l define the same irreducible polynomial of degree l , we call it the **standard polynomial** of degree l .

STANDARD EMBEDDINGS (SAME LEVEL)

Let $l \mid m$ and A_l, A_m algebras with the **same level** a , $\zeta_l = (\zeta_m)^{m/l}$

- ▶ α_l and α_m **standard** solutions of (H90) for ζ_l and ζ_m

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 - ▶ $c_l = c_m = (\zeta_{p^a-1})^a$
 - ▶ $\kappa_{l,m} = 1$
- ▶ The embedding $[\alpha_l] \mapsto [(\alpha_m)^{m/l}]$ is **standard** too (only depends on ζ_{p^a-1}).

STANDARD EMBEDDINGS (DIFFERENT LEVEL)

Let $l \mid m$ and A_l of level a , A_m of level b , $a \neq b$.

- Natural norm-compatibility condition, we want:

$$(\zeta_{p^b-1})^{\frac{p^b-1}{p^a-1}} = N(\zeta_{p^b-1}) = \phi_{\mathbb{F}_{p^a} \hookrightarrow \mathbb{F}_{p^b}}(\zeta_{p^a-1})$$

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$$(\alpha_{p^b-1})^{\frac{p^b-1}{p^a-1}} = (1 \otimes \kappa_{p^a-1, p^b-1}) \Phi_{A_{p^a-1} \hookrightarrow A_{p^b-1}}(\alpha_{p^a-1}) \text{ with}$$

$$\kappa_{p^a-1, p^b-1} = (\zeta_{p^b-1})^{\frac{(a-b)p^a+b+bp^b-ap^a}{(p^a-1)^2}}$$

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$$\kappa_{p^a-1, p^b-1} = (\zeta_{p^b-1})^{\frac{(a-b)p^a+b+bp^b-ap^a}{(p^a-1)^2}}$$

- If α_l and α_m are **standard solutions**, then

$$\kappa_{l,m} = (\zeta_{p^b-1})^{\frac{(a-b)p^a+b+bp^b-ap^a}{(p^a-1)l}}$$

STANDARD EMBEDDINGS (DIFFERENT LEVEL)

Let $l \mid m$ and A_l of level a , A_m of level b , $a \neq b$ and

- ▶ $(\zeta_{p^b-1})^{\frac{p^b-1}{p^a-1}} = N(\zeta_{p^b-1}) = \phi_{\mathbb{F}_{p^a} \hookrightarrow \mathbb{F}_{p^b}}(\zeta_{p^a-1}),$
- ▶ $\zeta_l = (\zeta_{p^a-1})^{\frac{p^a-1}{l}}$
- ▶ $\zeta_m = (\zeta_{p^b-1})^{\frac{p^b-1}{m}}$
- ▶ α_l and α_m **standard solutions** of (H90) for ζ_l and ζ_m
- ▶ $\kappa_{l,m}$ only depends on ζ_{p^b-1} and **is easy to compute**
- ▶ The embedding $[\alpha_l] \mapsto [(1 \otimes \kappa_{l,m})(\alpha_m)^{m/l}]$ is **standard** too (only depends on $\zeta_{p^a-1}, \zeta_{p^b-1}$).

COMPATIBILITY AND COMPLEXITY

Proposition (Compatibility)

Let $l \mid m \mid n$ and $f : \mathbb{F}_{p^l} \hookrightarrow \mathbb{F}_{p^m}$, $g : \mathbb{F}_{p^m} \hookrightarrow \mathbb{F}_{p^n}$, $h : \mathbb{F}_{p^l} \hookrightarrow \mathbb{F}_{p^n}$ the standard embeddings. Then we have $g \circ f = h$.

Proposition (Complexity)

Given a collection of Conway polynomials of degree up to d , for any $l \mid m \mid p^i - 1$, $i \leq d$

- ▶ Computing a standard solution α_l takes $\tilde{O}(l^2)$
- ▶ Given α_l and α_m , computing the standard embedding $f : \mathbb{F}_{p^l} \hookrightarrow \mathbb{F}_{p^m}$ takes $\tilde{O}(m^2)$

IMPLEMENTATION

Implementation using Flint/C and Nemo/Julia.

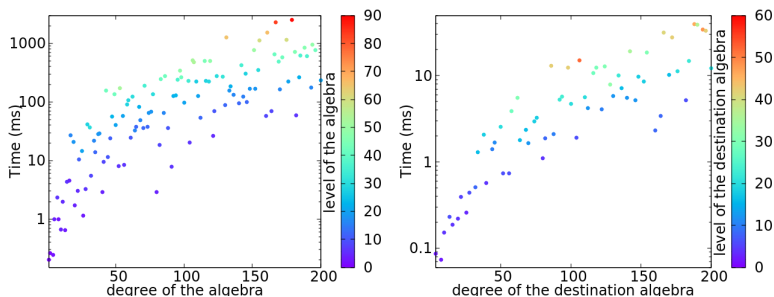


Figure: Timings for computing α_l (left, logscale), and for computing $\mathbb{F}_{p^2} \hookrightarrow \mathbb{F}_{p^l}$ (right, logscale) for $p = 3$.

STANDARD POLYNOMIALS

$$\begin{array}{r} x + 1 \\ x^3 + x + 1 \\ x^5 + x^3 + 1 \\ x^7 + x + 1 \\ x^9 + x^7 + x^4 + x^2 + 1 \\ x^{11} + x^8 + x^7 + x^6 + x^2 + x + 1 \\ x^{13} + x^{10} + x^5 + x^3 + 1 \\ x^{15} + x + 1 \\ x^{17} + x^{11} + x^{10} + x^8 + x^7 + x^6 + x^4 + x^3 + x^2 + x + 1 \\ x^{19} + x^{17} + x^{16} + x^{15} + x^{14} + x^{13} + x^{12} + x^8 + x^7 + x^6 + x^5 + x^3 + 1 \end{array}$$

Table: The ten first standard polynomials derived from Conway polynomials for $p = 2$.

CONCLUSION, FUTURE WORKS

- ▶ We implicitly assume that we have **compatible roots** ζ (i.e. $\zeta_l = (\zeta_m)^{m/l}$ for $l \mid m$)
 - ▶ In practice, this is done using **Conway polynomials**
- ▶ With Conway polynomials up to degree d , we can compute embeddings to finite fields up to any degree $l \mid p^i - 1, i \leq d$
 - ▶ quasi-quadratic complexity

Future works:

- ▶ Make this less standard, but more practical

Thank you!
Merci !