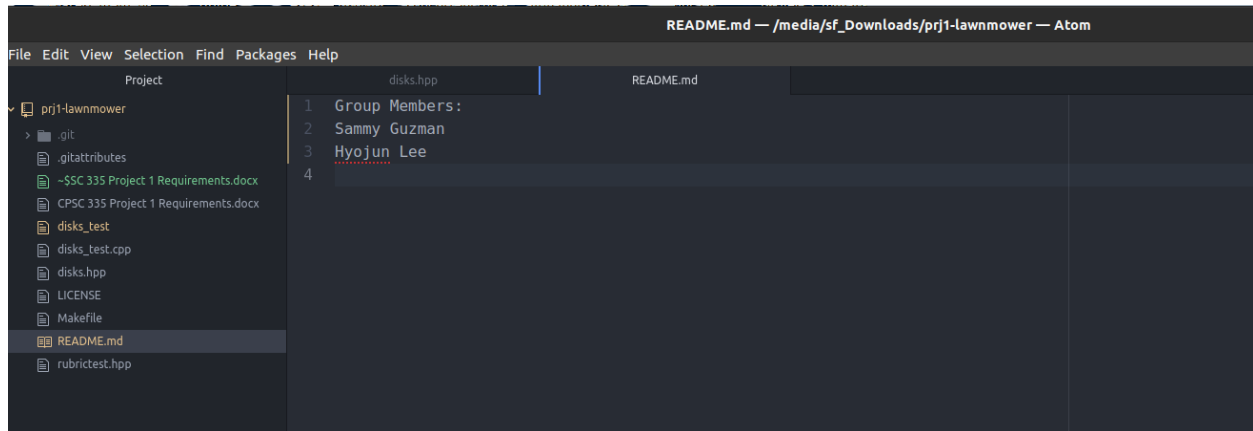


# CPSC 335 Project 1

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```
student@tuffix-vm:/media/sf_Downloads/prj1-lawnmower$ make
g++ -std=c++11 -Wall disks_test.cpp -o disks_test
./disks_test
disk_state still works: passed, score 1/1
sorted_disks still works: passed, score 1/1
disk_state::is_initialized: passed, score 3/3
disk_state::is_sorted: passed, score 3/3
alternate, n=4: passed, score 1/1
alternate, n=3: passed, score 1/1
alternate, other values: passed, score 1/1
lawnmower, n=4: passed, score 1/1
lawnmower, n=3: passed, score 1/1
lawnmower, other values: passed, score 1/1
TOTAL SCORE = 14 / 14
student@tuffix-vm:/media/sf_Downloads/prj1-lawnmower$ S
```

# The lawnmower algorithm

## Algorithm Design:

### Pseudo Code

Input: a positive integer  $n$  and a list of  $2n$  disks of alternating colors -light-disk

Output: List of  $2n$  disks,  $m$  for the number of swaps

```
for i = 0 to n/2 do

    for j = 0 to 2n-1 do
        if list[j] > list[j+1]
            swap
        end
    end for

    for j = 2n-1 to 0 do
        if list[j] < list[j-1]
            swap
        end
    end for
end for
```

### Step Count

```
for i = 0 to n/2 do                                // SC outer loop =
                                                    // (n+1)/2 times
    for j = 0 to 2n-1 do                            // SC forward loop
                                                    // 2n times
        if list[j] > list[j+1]                    // 1 TU
            Swap                                    // 3 TU
            m++                                     // 1 TU
        end
    end for

    for j = 2n-1 to 0 do                            // SC reverse loop
                                                    // 2n times
        if list[j] < list[j-1]                    // 1 TU
            Swap                                    // 3 TU
            M++                                     // 1 TU
        end
    end for
end for
```

$$S.C_{forward} = 2n(1 + \max(4, 0))$$

$$= 2n(1 + 4)$$

$$= 2n(5)$$

$$= 10n$$

$$S.C_{reverse} = 2n(1 + \max(4, 0))$$

$$= 2n(1 + 4)$$

$$= 2n(5)$$

$$= 10n$$

$$S.C. = S.C_{outer} * (S.C_{forward} + S.C_{reverse})$$

$$= \left(\frac{n+1}{2}\right)(10n + 10n)$$

$$= \left(\frac{n+1}{2}\right)(20n)$$

$$= \left(\frac{20n^2 + 20n}{2}\right)$$

$$= (10n^2 + 10n)$$

## Mathematical Analysis

-Big-O efficiency class:  $O(n^2)$

-Definition:

$$10n^2 + 10n \in O(n^2)$$

$$10n^2 + 10n \leq C \cdot n^2$$

$$C = 10 + 10 = 20, no = 1$$

$$10n^2 + 10n \leq 20n^2$$

$$10 + 10 \leq 20$$

$$20 \leq 20 \rightarrow true$$

By Definition, we can say that  $10n^2 + 10n \in O(n^2)$ .

Limit Theorem:

Prove that  $10n^2 + 10n \in O(n^2)$  using the limit theorem

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{10n^2 + 10n}{n^2} \\ &= \lim_{n \rightarrow \infty} \frac{10n^2}{n^2} + \frac{10n}{n^2} \\ &= \lim_{n \rightarrow \infty} 10 + \lim_{n \rightarrow \infty} \frac{10}{n} \end{aligned}$$

$$\lim_{n \rightarrow \infty} = 10 \geq 0,$$

By Limit theorem, we can say that  $10n^2 + 10n \in O(n^2)$ .

# The Alternating Disk Algorithm

## Algorithm Design:

### Pseudo Code

```
For i = 0 to N+1 step
  if (i % 2 == 0)
    for j = 0 to N-1 step + 2
      if (list[j] > list[j + 1])
        Swap disks
        m++;
      end if
    end for
  else do
    for j = 1 to N-2 step + 2
      if (list[j] > list[j + 1])
        Swap disks
        m++;
      end if
    end for
  end if
end for
```

### Step Count

```
For i = 0 to N // (n-0)+1/1= n+1
  if (i % 2 == 0) // 2 TU
    for j = 0 to N-1 step + 2 // (n-1-0+1)/2 = n/2
      if (list[j] > list[j + 1]) // 1 TU
        Swap disks // 3 TU
        M++; // 1 TU
      end
    end
  else
    for j = 1 to N-2 step + 2 // ((n-2)-1+1)/2= n-2/2
      if (list[j] > list[j + 1]) // 1 TU
        Swap disks // 3 TU
        M++; // 1 TU
      end
    end
  end if
End
```

$$S.C. = S.C_{outer} * (S.C_{if\ block})$$

$$S.C_{outer\ loop} = n + 1$$

$$S.C_{if\ block} = 2 + \max(S.C_{even\ loop}, S.C_{odd\ loop})$$

$$\begin{aligned} S.C_{even\ loop} &= \frac{n}{2} (1 + \max(4, 0)) \\ &= \frac{n}{2} (1 + 4) \\ &= \frac{n}{2} (5) \\ &= \frac{5n}{2} \end{aligned}$$

$$\begin{aligned} S.C_{odd\ loop} &= \frac{n-2}{2} (1 + \max(4, 0)) \\ &= \frac{n-2}{2} (1 + 4) \\ &= \frac{n-2}{2} (5) \\ &= \frac{5n-10}{2} \end{aligned}$$

$$S.C = (n + 1)(2 + \max(\frac{5n}{2}, \frac{5n-10}{2}))$$

$$= (n + 1)(2 + \frac{5n}{2})$$

$$= \frac{5n^2}{2} + \frac{5n}{2} + 2n + 2$$

$$= \frac{5n^2}{2} + \frac{9n}{2} + 2$$

# Mathematical Analysis

-Big-O efficiency class:

-Definition:

$$\frac{5n^2}{2} + \frac{9n}{2} + 2 \in O(n^2)$$

$$\frac{5n^2}{2} + \frac{9n}{2} + 2 \leq C \cdot n^2$$

$$C = \frac{5}{2} + \frac{9}{2} + 2 = 9, \text{ no } = 1$$

$$\frac{5n^2}{2} + \frac{9n}{2} + 2 \leq 9n^2$$

$$\frac{5}{2} + \frac{9}{2} + 2 \leq 9$$

$$9 \leq 9 \rightarrow \text{true}$$

By definition, we can say that  $\frac{5n^2}{2} + \frac{9n}{2} + 2 \in O(n^2)$ .

Limit Theorem:

Prove that  $\frac{5n^2}{2} + \frac{9n}{2} + 2 \in O(n^2)$  using the limit theorem

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{\frac{5}{2}n^2 + \frac{9}{2}n + 2}{n^2} \\ &= \lim_{n \rightarrow \infty} \frac{\frac{5}{2}n^2}{n^2} + \frac{\frac{9}{2}n}{n^2} + \frac{2}{n^2} \\ &= \lim_{n \rightarrow \infty} \frac{5}{2} + \lim_{n \rightarrow \infty} \frac{\frac{9}{2}}{n} + \lim_{n \rightarrow \infty} \frac{2}{n^2} \\ &= \frac{5}{2} + 0 + 0 \\ &= \frac{5}{2} \\ & \lim_{n \rightarrow \infty} = \frac{5}{2} \geq 0, \end{aligned}$$

By Limit Theorem, we can say that  $\frac{5n^2}{2} + \frac{9n}{2} + 2 \in O(n^2)$ .