Study Report: Narrowing

Deivid Vale

November 7, 2018

Abstract

Narrowing...

1 Introduction

2 Preliminaries

3 Narrowing

In this section, we discuss the use of narrowing techniques on solving equations over an equational theory E. By solving an equation we mean finding a substitution σ such that $s\sigma =_E t\sigma$. These solutions can be found by unification *modulo* E (if such a unification algorithm exists for this theory), but here we are interested in applications of rewriting theory to solving such equations.

We start the presentation by the so-called syntactic unification, i.e. E-unification with empty E. The transformation rules given below is due to Martelli and Montanari [citation], the idea is to transform sets of equations to other sets of equations until a termination state is reached; that is, a solution state or a failure one. We also extend this same set of rules for solving E-equations.

Definition 1. Let Σ be a signature. An equational goal is a finite set of Σ-equations.

Table 1 Martelli-Montanari rules

---

(1) Trivial

$$\{x \stackrel{?}{=} x\} \cup G \implies G$$

Delete trivial equations.

(2) Decompose

$$\{f(s_1, \cdots, s_n) \stackrel{?}{=} f(t_1, \ldots, t_n)\} \cup G \implies \{s_1 \stackrel{?}{=} t_1, \ldots, s_n \stackrel{?}{=} t_n\} \cup G$$

(3) Symbol Clash

$$\{f(s_1, \ldots, s_n) \stackrel{?}{=} g(t_1, \ldots, t_n)\} \cup G \implies \bot \text{ if } f \neq g$$

(4) Orient

$$\{t \stackrel{?}{=} x\} \cup G \implies \{x \stackrel{?}{=} t\} \cup G \text{ if } t \notin V$$

(5) Occurs Check

$$\{x \stackrel{?}{=} t\} \cup G \implies \bot \text{ se } x \in \texttt{vars}(t) \text{ and } x \neq t$$

(6) Variable Elimination

$$\{x \stackrel{?}{=} t\} \cup G \implies G\{x \mapsto t\} \text{ if } x \notin \texttt{vars}(t)$$

The application of the above rules non-deterministically transforms goals into goals:

$$G_0 \implies \cdots \implies G_n$$

Each application of a rule will then called a *elementary derivation step*. As for the case of (Variable Elimination), we may get some substitution on the way, we make them explicit by writing:

$$G_0 \implies G_1 \implies_{\sigma_1} G_2 \implies_{\sigma_2} \cdots \implies_{\sigma_i} \cdots \implies_{\sigma_{n-1}} G_n$$

The computed solution of the derivation chain is then the composition of such substitutions in their order of appearance.

Definition 2. A *successful* derivation chain is a finite sequence of equational goals $G_0, G_1, \ldots, G_n$ such that the last goal is empty. We also say a derivation chain has *failed* if it end is the fail symbol $\perp$.

Example 1. 1. We want to determine an mgu of the terms $f(g(x), h(x, u))$ and $f(z, h(f(y, y), z))$. That is, solving the equational goal:

$$\begin{aligned}
&\{f(g(x), h(x, u)) = f(z, h(f(y, y), z))\} \\
&\quad \overset{(2)}{\Longrightarrow} \{g(x) = z, h(x, u) = h(f(y, y), z)\} \\
&\quad \overset{(4)}{\Longrightarrow} \{z = g(x), h(x, u) = h(f(y, y), z)\} \\
&\quad \overset{(6)}{\Longrightarrow}_{[z/g(x)]} \{h(x, u) = h(f(y, y), g(x))\} \\
&\quad \overset{(2)}{\Longrightarrow} \{x = f(y, y), u = g(x)\} \\
&\quad \overset{(6)}{\Longrightarrow}_{[x/f(y,y)]} \{u = g(f(y, y))\} \\
&\quad \overset{(6)}{\Longrightarrow}_{[u/f(y,y)]} \emptyset
\end{aligned}$$

We get as the computed solution to the problem the composition $[z/g(x)][x/f(y, y)][u/g(f(y, y))]$.

2. A failing unification derivation:

$$\begin{aligned}
&\{h(x, y, x) = h(y, g(x), x)\} \\
&\quad \overset{(2)}{\Longrightarrow} \{x = y, x = g(x), x = x\} \\
&\quad \overset{(5)}{\Longrightarrow} \perp
\end{aligned}$$

It can be proved that this set of rules derive a correct and terminating procedure for syntactic unification. (add citation here).

The situation changes if we want to solve equations for a non-empty equation theory E.

4 Narrowing

Definition 3. Say a term s *narrows* to a term t if there exists a non-variable position $p \in \texttt{pos}(s)$, a variant $l \to r$ of a rewrite rule in $\mathcal{R}$, and a substitution σ satisfying two conditions:

1. σ is a *mgu* of $s|_p$ and l,
2. $t = (s[r]_p)\,\sigma$.

The relation $\rightsquigarrow$ is called *narrowing relation.* We write $s \rightsquigarrow_\sigma^* t$ if there exists a narrowing derivation

$$s = t_1 \rightsquigarrow_{\sigma_1} t_2 \rightsquigarrow_{\sigma_2} t_3 \rightsquigarrow \cdots \rightsquigarrow_{\sigma_{n-1}} t_n = t$$

Remark 1. Renaming of rewrite rules will be mandatory to ensure completeness of the narrowing approach. We always use a simple renaming such that $\texttt{vars}(l) \cap \texttt{vars}(s) = \emptyset$. This also extend to chains of narrowing derivations.

Narrowing was first introduced in the context of E-unification. Fay [cite] and Hullot [23] shows that narrowing is a complete method for solving equations in the theory defined by a confluent and terminating term rewriting system. In fact, narrowing is a first attempt for solving equations in arbitrary equational theories E, with the requirement that they can be represented as a convergent rewrite system. This approach also shows an important application for the Knuth-Bendix completion procedure: it prepares the way for solving equations over E, by delivering a complete TRS for E (if possible).

For this set of technology we give the name narrowing, we now present the first application (of solving equations modulo E) in the framework of transformation rules on sets of equational goals.

Consider an equational theory

References