

# Probability and Computing, 2nd Edition

Solutions to Chapter 5: Balls, Bins, and Random Graphs

Hahndeul Kim

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## 5.1

As  $(1 + 1/n)^n$  increases, we find the smallest  $n$  to reach the threshold.  $(1 + 1/n)^n$  first reaches  $0.99e$  at  $n = 50$ , and  $0.999999e$  at  $n = 499982$ . Since  $(1 - 1/n)^n$  also increases, we solve in a similar way.  $(1 - 1/n)^n$  first reaches  $0.99/e$  at  $n = 51$  and  $0.999999/e$  at  $n = 499991$ .

## 5.2

Recall the formula used in the birthday paradox: If there are  $N$  possibilities, then we solve for the smallest  $n$  that satisfies  $\prod_{i=1}^{n-1} (1 - \frac{i}{N}) \approx \prod_{i=1}^{n-1} e^{-i/n} = e^{-(n-1)n/2N} < 1/2$ . Note that we omitted the final approximation to derive exact numerical answers.

Regardless of whether the number of Social Security number digits is 9 or 13, using the last four digits gives  $N = 10000$  and this gives  $n = 119$ .

In the case where the number of digits is 9 ( $N = 10^9$ ), we get  $n = 37234$ .

In the case where the number of digits is 13 ( $N = 10^{13}$ ), we get  $n = 3723298$ .

## 5.3

Let the number of balls thrown be  $m$ . Then the desired probability is  $\prod_{i=0}^{m-1} (1 - \frac{i}{n})$ .

We first determine  $c_1$ .  $m = c_1\sqrt{n}$  should satisfy  $\prod_{i=0}^{m-1} (1 - \frac{i}{n}) \leq \prod_{i=0}^{m-1} e^{-i/n} = e^{-(m-1)m/2n} \leq e^{-1}$ . Since  $(m-1)m = c_1^2 n - c_1\sqrt{n} \geq 2n$ ,  $(c_1^2 - 2)\sqrt{n} \geq c_1$ .

Therefore, we choose  $c_1$  that is greater than or equal to  $\frac{1}{2} \left( \frac{1}{\sqrt{n}} + \sqrt{\frac{1}{n} + 8} \right)$ .

Now we determine  $c_2$ . To use the given hint, assume that  $2m < n$ .

$\prod_{i=0}^{m-1} (1 - \frac{i}{n}) \geq \prod_{i=0}^{m-1} \exp(-\frac{i}{n} - \frac{i^2}{n^2}) = \exp(-\frac{m(m-1)}{2n} - \frac{(m-1)m(2m-1)}{6n^2})$   
 $= \exp(-\frac{m(m-1)}{2n}(1 + \frac{2m-1}{3n})) \geq \exp(-\frac{m^2}{2n}(1 + \frac{2m}{3n})) \geq \frac{1}{2}$  should be satisfied for  $m = c_2\sqrt{n}$ . This is equivalent to satisfying  $\frac{c_2^2}{2}(1 + \frac{2c_2}{3\sqrt{n}}) \leq \ln 2$ .

Since  $n$  is sufficiently large, choosing  $c_2 = \sqrt{2 \ln 2 - \frac{1}{\ln n}}$  yields the desired result.

## 5.4

Let event  $A$  indicate that there exist two or more people who share a birthday, and event  $B$  indicate that exactly two people share a birthday. Then our desired probability would be  $\Pr(A - B) = \Pr(A) - \Pr(B)$  since  $B \subset A$ .

We first determine  $\Pr(A)$ , which is easy:  $\Pr(A) = 1 - \prod_{i=1}^{100} \frac{366-i}{365}$ .

We now determine  $\Pr(B)$ . If there are  $i$  shared birthdays in which each day is

shared by exactly two people, then the number of possible permutations would be the product of the following terms:

$\binom{365}{i}$  ways to choose  $i$  shared days,  $\binom{100}{2i}$  ways to choose  $2i$  people to share birthdays,  $\prod_{j=1}^i \binom{2j}{2}$  ways to distribute  $i$  birthdays to  $2i$  people and  $\prod_{j=1}^{100-2i} (366 - i - j)$  ways to distribute unique birthdays to the rest.

Thus,  $\Pr(B) = \sum_{i=0}^{50} \frac{365!100!}{i!(100-2i)!(265+i)!2^i} \times \frac{1}{365^{100}}$ .

Therefore, we can determine our desired probability  $\Pr(A-B) = \Pr(A) - \Pr(B)$ .

## 5.5

Let  $X \sim \text{Poisson}(\lambda)$ . Then  $M_X(t) = \mathbf{E}[e^{tX}] = e^{\lambda(e^t-1)}$  holds. By computing the second derivative of  $M_X(t)$  with respect to  $t$  and plugging  $t = 0$  in, we get  $\mathbf{E}[X^2] = \lambda + \lambda^2$ . Thus,  $\mathbf{Var}[X] = \lambda$  follows.

## 5.6

We first show that  $Y \sim \text{Poisson}(\mu p)$ .

$$\begin{aligned} \Pr(Y = k) &= \sum_{i=k}^{\infty} \Pr(X = i) \binom{i}{k} p^k (1-p)^{i-k} = \sum_{i=k}^{\infty} \frac{e^{-\mu} \mu^i}{i!} \frac{i!}{k!(i-k)!} p^k (1-p)^{i-k} \\ &= \frac{e^{-\mu} (\mu p)^k}{k!} \sum_{i=k}^{\infty} \frac{(\mu(1-p))^{i-k}}{(i-k)!} = \frac{e^{-\mu} (\mu p)^k}{k!} e^{\mu(1-p)} = \frac{e^{-\mu p} (\mu p)^k}{k!}. \end{aligned}$$

We can also similarly show that  $Z \sim \text{Poisson}(\mu(1-p))$ .

Now we show that  $\Pr(Y = i, Z = j) = \Pr(Y = i) \Pr(Z = j)$ . Note that  $X = Y + Z$  by definition. This allows us to write  $\Pr(Y = i, Z = j)$  as  $\Pr(Y = i, X = i + j) = \Pr(X = i + j) \binom{i+j}{i} p^i (1-p)^j = \frac{e^{-\mu} \mu^{i+j}}{i!j!} p^i (1-p)^j$ .

Since  $\Pr(Y = i) \Pr(Z = j) = \frac{e^{-\mu p} (\mu p)^i}{i!} \frac{e^{-(1-p)\mu} ((1-p)\mu)^j}{j!} = \frac{e^{-\mu} \mu^{i+j} p^i (1-p)^j}{i!j!} = \Pr(Y = i, X = i + j)$ ,  $Y \perp\!\!\!\perp Z$ . ■

## 5.7

We first prove that  $\ln(1+x) \leq x$ , which is equivalent to  $1+x \leq e^x$ .

Since  $\ln(1+x) - x = -\frac{x^2}{2} + \frac{x^3}{3} - \dots$ , this can be seen as an alternating series as  $\frac{x^n}{n}$  is monotonically decreasing in  $|x| \leq 1$ . We can apply rearrangements to the alternating series as  $\ln(1+x) - x = -\sum_{n=1}^{\infty} \left( \frac{x^{2n}}{2n} - \frac{x^{2n+1}}{2n+1} \right)$ , since the Taylor expansion of  $\ln(1+x)$  is absolutely convergent (to  $e^x - 1$ ). The rearrangement gives  $\ln(1+x) - x \leq 0$ , which is the desired result.

We now prove  $x + \ln(1-x^2) \leq \ln(1+x)$ , which is equivalent to  $e^x(1-x^2) \leq 1+x$ . Since  $\ln(1-x^2) = \ln(1+x) + \ln(1-x)$ ,  $x + \ln(1-x^2) \leq \ln(1+x)$  is reduced to  $\ln(1-x) \leq -x$ . At  $|x| \leq 1$ , this is equivalent to  $\ln(1+x) \leq x$ , which we have previously proved. ■

## 5.8

(a) Since the ball is equally likely to fall in one of the three bins, the desired probability is  $1/3$ .

(b) Since the bin 2 did not receive balls, we can simply think of this as throwing balls  $n$  into  $n - 1$  bins. The conditional expectation would be  $n/(n - 1)$ .

(c) Note that the probability that bin 1 receives more balls than bin 2 is the same as that of bin 2 receiving more balls than bin 1. Thus, we first compute the probability that two bins receive the same number of balls, which is

$$P = \sum_{k=0}^{\lfloor n/2 \rfloor} \binom{n}{2k} \binom{2k}{k} \left(\frac{1}{n}\right)^{2k} \left(1 - \frac{2}{n}\right)^{n-2k}. \text{ The desired probability would be } (1 - P)/2.$$

## 5.9

In the given condition, the expected number of elements in a single bucket is at most  $a$ . Since  $a = O(1)$ , sorting all buckets can still be done in linear time.

## 5.10

(a) By the Poisson approximation, the probability  $p$  is bounded as  $p \leq e\sqrt{n}(\frac{1}{e})^n$ .

(b)  $\frac{n!}{n^n}$ .

(c) Since  $\Pr(Z = n) = \frac{e^{-n}n^n}{n!}$  when  $Z \sim \text{Poisson}(n)$ ,  $\frac{n!}{n^n} \times \frac{e^{-n}n^n}{n!} = e^{-n}$  shows the claim. Theorem 5.6 states that the distribution  $(Y_1, \dots, Y_n)$  constrained on  $\sum_{i=1}^n Y_i = n$  is equivalent to the balls and bins model. Note that each  $Y_i$  follows  $\text{Poisson}(1)$  and each  $X_i$  denotes the load of the  $i$ th bin in the balls and bins model. Then using theorem 5.6,  $(1/e)^n / (\frac{e^{-n}n^n}{n!}) = \frac{\Pr(\forall i, Y_i=1)}{\Pr(\sum_i Y_i=n)} = \Pr(\forall i, Y_i = 1 | \sum_i Y_i = n) = \Pr(\forall i, X_i = 1) = \frac{n!}{n^n}$ .

## 5.11

Let  $X_i$  be the indicator variables that are 1 if there is a  $k$ -gap starting at  $i$ .

$$\text{Let } X = \sum_{i=0}^{n-k} X_i.$$

(a) The expected number of  $k$ -gaps would be  $\sum_{i=0}^{n-k} \mathbf{E}[X_i] = (n - k + 1)(1 - \frac{k}{n})^m$ .

(b) First, we assume that the bin loads follow the Poisson distribution to derive the Poisson-approximated Chernoff bound. We divide  $\{X_i\}$  into  $k$  subsets, so that all indicator variables in the same subset are independent. First, we derive the Chernoff bound for  $Y_0 = \sum_{i \geq 0} X_{ik}$  where  $0 < \delta < 1$ .

$\Pr(Y_0 \geq (1+\delta)\mathbf{E}[Y_0]) \leq \exp(-\frac{1}{3}\frac{\mathbf{E}[X]}{k}\delta^2)$ , and  $\Pr(Y_0 \leq (1-\delta)\mathbf{E}[Y_0]) \leq \exp(-\frac{1}{2}\frac{\mathbf{E}[X]}{k}\delta^2)$  holds. Therefore,  $\Pr(|Y_0 - \mathbf{E}[Y_0]| \geq \delta\mathbf{E}[Y_0]) \leq 2\exp(-\frac{1}{3}\frac{\mathbf{E}[X]}{k}\delta^2)$ .

With the union bound for all  $Y_i$ , we get  $\Pr(|X - \mathbf{E}[X]| \geq \delta \mathbf{E}[X]) \leq 2k \exp(-\frac{1}{3} \frac{\mathbf{E}[X]}{k} \delta^2)$ . Since we have used the Poisson approximation to compute the Chernoff bound, the computed upper bound should be multiplied by  $e\sqrt{m}$ .

## 5.12