do this ligible? Bisection: 1x\_x 1 \le \frac{(b-a)}{2^n} § 2.4 Order of convergence def: We say xn=x+O(Bn) if Bn>O and there efests a constant C such that |x-x| ≥ Cβn, i.e., Xn x with note Bn. del: Suppose n= x = x. el there are constants 270, 270, with  $\lim_{n\to\infty} \frac{|x_{n+1} \times |}{|x_n \times |^{\alpha}} = \lambda$ then we say x > x with order a and asymptotic ernar canatant ). quadratic cutic il >>1 slower x=3 x=2 For a fixed &, sequences with a smaller error constant converge faster. super linear sublinear e.g.)  $\frac{1}{n} \rightarrow 0$  sublinearly because  $\frac{|\vec{n}+1|}{|\vec{1}|^2} = \frac{n}{n+1} \rightarrow 1$ 

$$\frac{\left|\frac{1}{2^{n+1}}-0\right|}{\left|\frac{1}{2^{n}}-0\right|^{\frac{n}{2}}} = \frac{1}{2} - \frac{1}{2}$$

$$\frac{1}{2^{n}} - 0$$
linearly asymptotic error

High a and I low is fast!

thm: Bisection method generates a sequence that converges linearly with asymptotic error constant  $\frac{1}{2}$ .

$$\frac{|X_{n+1}-X|}{|X_n-X|^2} \approx \frac{\frac{(b-a)}{2^{n+1}}}{\frac{(b-a)}{2^n}} = \frac{1}{2}$$

thm: Suppose x=g(x), and suppose g is p times differentiable near x. Assume

$$g'(x) = g''(x) = \dots = g^{(p-1)}(x) = 0$$

ander b and asymptotic error  $\frac{g(b)(x)}{b!}$ .

Pol: Use Taylor's theorem.

Recall ... to solve f(x)=0

$$x^{\nu} = \partial(x^{\nu-1}) := x^{\nu-1} - \frac{t_1(x^{\nu-1})}{t(x^{\nu-1})}$$

Heed point !

$$g(x) = x - \frac{f(x)}{f'(x)}, \text{ any } f(x) = 0$$

$$g'(x) = 1 - \frac{f'(x)f'(x)}{f'(x)} - f(x)f''(x) = 0$$

$$g''(x) = \frac{f(x)f''(x)}{[f'(x)]^{2}} = 0$$

$$g''(x) = \frac{f'(x)f''(x)}{[f'(x)]^{2}} - \frac{f'(x)f''(x)}{[f'(x)]^{2}} = 0$$

$$= \frac{f'(x)f''(x)}{[f'(x)]^{2}} - \frac{f''(x)f''(x)}{[f'(x)]^{2}} = 0$$

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Since g'(x)=0, but g"(x) ≠0, Newton's method converges quadratically with asymptotic error

ence & is known

$$\frac{f''(x)}{2f'(x)}.$$

for large n, if  $e_{n=1} \times \pi \times 1$ ,  $e_{n+1} \approx \lambda e_n^{\alpha}$ 

 $\frac{e_{n+1}}{e_n} \approx \frac{\lambda e_n^{\alpha}}{\lambda e_{n+1}^{\alpha}} = \left(\frac{e_n}{e_{n-1}}\right)^{\alpha}$   $\log \left(\frac{e_{n+1}}{e_n}\right)$ 

 $\alpha \approx \frac{\log (e_{n+1}/e_n)}{\log (e_n/e_{n-1})}$