

MultiMap

implementation on a hash table, collision resolution by separate chaining

- Project no. 1 -

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ADT MultiMap

A multimap is a container which stores $\langle \text{key}, \text{value} \rangle$ pairs. The keys are not unique, therefore a key can have multiple associated values. The order in which the pairs are kept is not important.

MultiMap Domain

$$\text{MM} = \{ \text{mm} \mid \text{mm is a multimap with } \langle k, v \rangle \text{ pairs, where } k \in \text{TKey and } v \in \text{TValue} \}$$

MultiMap Interface

init(mm):

descr: creates a new empty multimap
pre: true
post: $\text{mm} \in \text{MM}$, mm is an empty multimap

destroy(mm):

descr: destroys a multimap
pre: $\text{mm} \in \text{MM}$
post: the multimap was destroyed

add(mm, k, v):

descr: add a new key-value pair to the multimap
pre: $\text{mm} \in \text{MM}$, $k \in \text{TKey}$, $v \in \text{TValue}$
post: $\text{mm}' \in \text{MM}$, $\text{mm}' = \text{mm} \cup \langle k, v \rangle$

remove(mm, k, v):

descr: removes a $\langle \text{key}, \text{value} \rangle$ pair from the multimap
pre: $\text{mm} \in \text{MM}$, $k \in \text{TKey}$, $v \in \text{TValue}$
post: $\text{remove} \leftarrow \begin{cases} \text{true,} & \text{if } \langle k, v \rangle \in \text{mm, mm}' \in \text{MM, mm}' = \text{mm} - \langle k, v \rangle \\ \text{false,} & \text{otherwise} \end{cases}$

search(mm, k, l):

descr: returns a list with all the values associated to a key
pre: $\text{mm} \in \text{MM}$, $k \in \text{Tkey}$
post: $l \in L$, l is the list of values associated to the key k. If k is not in the multimap, l is an empty list

iterator(mm, it):

descr: returns an iterator over the multimap
pre: $\text{mm} \in \text{MM}$
post: $it \in I$, it is an iterator over mm, the current element from it is the first pair from mm, or it is invalid if mm is empty

size(mm):

descr: returns the number of pairs from the multimap
pre: $\text{mm} \in \text{MM}$
post: $\text{size} \leftarrow$ the number of pairs from mm

isEmpty(m):

descr: verifies if the map is empty
pre: $m \in M$
post: $\text{isEmpty} \leftarrow \begin{cases} \text{true,} & \text{if mm contains no pairs} \\ \text{false,} & \text{otherwise} \end{cases}$

Iterator Domain

$I = \{ it \mid it \text{ is an iterator over a multimap} \}$

Iterator Interface

`init(it, mm):`

descr: creates a new iterator for a multimap

pre: mm is a multimap

post: $it \in I$ and it points to the first element in mm if mm is not empty or it is not valid

`getCurrent(it):`

descr: returns the current element from the iterator

pre: $it \in I$, it is valid

post: $getCurrent \leftarrow \langle k, v \rangle$, $k \in TKey$, $v \in TValue$, $\langle k, v \rangle$ is the current pair from it

throws: an exception if it is invalid

`next(it):`

descr: moves the current element from the container to the next pair or makes the iterator invalid if no pairs are left

pre: $it \in I$, it is valid

post: the current element from it points to the next element from the multimap

throws: an exception if it is invalid

`valid(it):`

descr: verifies if the iterator is valid

pre: $it \in I$

post: $valid \leftarrow \begin{cases} \text{True,} & \text{if it points to a valid element from the container} \\ \text{False,} & \text{otherwise} \end{cases}$

MultiMap representation

Pair:

key: TKey
value: TValue

Node:

info: Pair
next: ↑ Node

Multimap:

table: ↑Node[]
m: Integer
h: TFunction

Iterator representation

Iterator:

mm: MultiMap
currentPos: Integer
currentNode: ↑ Node

Multimap implementation

```
subalgorithm init(mm) is:
    mm.m ← @initial capacity
    mm.table ← @allocate an array with mm.m positions
    mm.hf ← hf
    for i ← 0, mm.m execute
        mm.table[i] ← NIL
    end-for
end-subalgorithm
```

Complexity: $\Theta(m)$

```
subalgorithm destroy(mm) is:
    for i ← 0, mm.m execute
        currentNode ← mm.table[i]
        while (currentNode != NIL) execute
            prevNode ← currentNode
            currentNode ← [currentNode].next
            free(prevNode)
        end-while
    mm.m ← 0
    free(mm.table)
    end-for
end-subalgorithm
```

Complexity: $\Theta(m)$

```
subalgorithm add(mm, k, v) is:
    allocate (newNode)
    [newNode].info.key ← k
    [newNode].info.value ← v
    [newNode].next ← mm.table[mm.hf(k, mm.m)]
    mm.table[mm.hf(k, mm.m)] ← newNode
end-subalgorithm
```

Complexity: $\Theta(1)$

```
function remove(mm, k, v) is:
    currentNode ← mm.table[mm.hf(k, mm.m)]
    prevNode ← NIL

    if currentNode = NIL then
        remove ← false
    end-if
    found ← false
    while currentNode != NIL and found = false execute
        if [currentNode].info.key = k and [currentNode].info.value = v then
            found ← true
        else
            prevNode ← currentNode
            currentNode ← [currentNode].next
        end-if
    end-while
```

```

    if found = false then
        remove ← false
    end-if
    if prevNode = NIL and [currentNode].next = NIL then
        mm.table[mm.hf(k, mm.m)] ← NIL
    else
        if prevNode = NIL and [currentNode].next != NIL then
            mm.table[mm.hf(k, mm.m)] ← [currentNode].next
        else
            [prevNode].next ← [currentNode].next
        end-if
    end-if
    remove ← true
end-function

```

Complexity: $\Theta(1 + \alpha)$

```

function search(mm, k, l) is:
    currentNode ← mm.table[mm.hf(k, mm.m)]
    i ← 0
    while currentNode != NIL execute
        if [currentNode].info.key = k then
            l[i] ← [currentNode].info.value
            i++
        end-if
        currentNode ← [currentNode].next
    end-while
    search ← l
end-function

```

Complexity: $\Theta(1 + \alpha)$

```

function size(mm) is:
    count ← 0
    for i ← 0, mm.m execute
        currentNode ← mm.table[i]
        while currentNode != NIL execute
            currentNode ← [currentNode].next
            count = count + 1
        end-while
    end-for
    size ← count
end-function

```

Complexity: $\Theta(m + \text{number of pairs in the table})$

```

function isEmpty(mm) is:
    for i=0, mm.m execute
        if mm.table[i] != NIL then
            isEmpty ← false
        end-if
    end-for
    isEmpty ← true
end-function

```

Complexity: $\Theta(m)$

```

function iterator(mm, it) is:
    iterator ← init(it, mm) //init from iterator init
end-function

```

Complexity: $O(m)$

Where $\alpha = n / m$ is the load factor of the table with m slots containing n elements, that is, the average number of elements stored in a chain. In case of separate chaining α can be less than, equal to, or greater than 1.

```

function hashFunction(s, m) is:
    //descr: is a function that maps a key k to a slot in the table
    //pre: s is a string, m is an integer representing the multimap's capacity
    //post: returns an integer between 0 and m-1 representing the allocated slot for the string in the multimap's table
    pos ← 0
    for i ← 0, length(s) execute
        pos ← pos + int(s[i]) * 31 ^ (length(s) - i - 1)
    end-for
    if pos < 0 then
        pos ← -pos
    end-if
    hashFunction ← pos % m
end-function

```

Complexity: Θ (length of string)

Iterator implementation

```
subalgorithm init(it, mm) is:
    it.mm ← mm
    it.currentPos ← 0
    it.currentNode ← NIL
    while it.currentPos < it.mm.m and it.currentNode = NIL execute
        it.currentNode ← it.mm.table[currentPos]
        it.currentPos = it.currentPos + 1
    end-while
    if valid(it) then
        it.currentPos ← it.currentPos - 1
    else
        it.currentNode ← NIL
    end-if
end-subalgorithm
```

Complexity: $O(m)$

```
function getCurrent(it) is:
    if valid(it) then
        getCurrent ← [it.currentNode].info
    else
        @throw exception for invalid iterator
    end-if
end-function
```

Complexity: $\Theta(1)$

```
subalgorithm next(it) is:
    it.currentNode ← [it.currentNode].next
    while it.currentPos < it.mm.m and it.currentNode = NIL execute
        it.currentPos ← it.currentPos + 1
        it.currentNode ← it.mm.table[it.currentPos]
    end-while
    if valid(it) = false then
        @throw exception for invalid iterator
    end-if
end-subalgorithm
```

Complexity: $O(m)$

```
function valid(it) is:
    if it.currentNode != NIL then
        valid ← true
    else
        valid ← false
    end-if
end-function
```

Complexity: $\Theta(1)$

Problem statement

The local bookstore of a small town had become increasingly popular and the business is thriving. To be more efficient the bookstore now needs an app that stores their stock of books. Each book has an author and title. An author can write more than one book. The app should allow the staff to search by author, add a new book and delete an existing book really fast. Two or more occurrences of the same book (same author and title) will be kept separately as different entrances (also, when we delete a book only one occurrence will be deleted).

Justification

In this case the multimap is a great choice for a container because each book consists of a <key, value> pair where the key is the author of the book and the value represents its title. Furthermore, since an author can write more than one book and we can also have more than one book having the same author and title in stock the keys are not unique.

Implementation

main.cpp

subalgorithm populate(mm) is:

//descr: populates the multimap with some given pairs

//pre: mm is a multimap

//post: the multimap was populated with some given elements

 add(mm, author, title) *//multiple times for different values*

end-subalgorithm

Complexity: $\Theta(1)$

subalgorithm printMenu() is:

//descr: prints the menu for the user interface

end-subalgorithm

Complexity: $\Theta(1)$

function main() is:

 init (mm)

 populate(mm)

Complexity: $\Theta(m)$

Complexity: $\Theta(1)$

 while true execute:

 printMenu()

 read command

 if command = 0 then

 @ exit console

 end-if

 else if command = 1 then

 read author

 read title

 add(mm, author, title)

 end-if

Complexity: $\Theta(1)$

```

else if command = 2 then
    read author
    read title
    if remove(mm, author, title) = true
        print "The book was removed"
    else
        print "The book did not exist"
    end-if
end-if

```

Complexity: $\Theta(1 + \alpha)$

```

else if command = 3 then
    read author
    search(mm, author, books)
    for i ← 0, length(books) execute
        print books[i].title
    end-for
end-if

```

Complexity: $\Theta(1 + \alpha)$

```

else if command = 4 then
    @try
    imm ← iterator(imm, mm)
    while valid(imm) execute
        b ← getCurrent(imm)
        print b
        next(imm)
    end-while
    @catch iterator error
end-if

```

Complexity: $\Theta(m + \text{number of pairs in the table})$

end-while

main ← 0

end-function

Where $\alpha = n / m$ is the load factor of the table with m slots containing n elements, that is, the average number of elements stored in a chain. In case of separate chaining α can be less than, equal to, or greater than 1.

Tests

Test.h

#pragma once

```
class Test
{
public:
    Test() {};
    ~Test() {};

    void createTest();
    void addTest();
    void removeTest();
    void searchTest();
    void sizeTest();
    void isEmptyTest();
};
```

Test.cpp

```
#include "Test.h"
#include "MultiMap.h"
#include <assert.h>
```

```
void Test::createTest()
{
    MultiMap mm{ };
    Node* l[100];
    mm.getLibrary(l);
    for (int i = 0; i < mm.getCapacity(); i++)
        assert(l[i] == nullptr);
}
```

```
void Test::addTest()
{
    MultiMap mm{ };
    assert(mm.size() == 0);
    mm.add("F. Scott Fitzgerald", "The Great Gatsby");
    assert(mm.size() == 1);
    mm.add("Margaret Atwood", "The Handmaid's Tale");
    assert(mm.size() == 2);
    mm.add("Margaret Atwood", "The Handmaid's Tale");
    assert(mm.size() == 3);
    mm.add("Daniel Keyes", "Flowers For Algernon");
    mm.add("Amy Harmon", "From Sand and Ash");
    mm.add("Jane Austen", "Emma");
    mm.add("Daniel Keyes", "The Minds of Billy Milligan");
    mm.add("Amy Harmon", "From Sand and Ash");
}
```

```

        mm.add("Jane Austen", "Persuasion");
        mm.add("Charlotte Bronte", "Jane Eyre");
        assert(mm.size() == 10);
        mm.add("Daniel Keyes", "Flowers For Algernon");
        mm.add("Amy Harmon", "From Sand and Ash");
        mm.add("Jane Austen", "Emma");
        mm.add("Daniel Keyes", "The Minds of Billy Milligan");
        mm.add("Amy Harmon", "From Sand and Ash");
        mm.add("Jane Austen", "Persuasion");
        mm.add("Charlotte Bronte", "Jane Eyre");
        assert(mm.size() == 17);
    }

void Test::removeTest()
{
    MultiMap mm{ };
    assert(mm.size() == 0);
    mm.add("F. Scott Fitzgerald", "The Great Gatsby");
    assert(mm.size() == 1);
    assert(mm.remove("F. Scott Fitzgerald", "The Great Gatsby")==true);
    assert(mm.size() == 0);
    assert(mm.remove("F. Scott Fitzgerald", "The Great Gatsby")==false);
    assert(mm.size() == 0);
    mm.add("Margaret Atwood", "The Handmaid's Tale");
    assert(mm.size() == 1);
    mm.add("Margaret Atwood", "The Handmaid's Tale");
    assert(mm.size() == 2);
    assert(mm.remove("F. Scott Fitzgerald", "The Great Gatsby")==false);
    assert(mm.size() == 2);
    assert(mm.remove("Margaret Atwood", "The Handmaid's Tale")==true);
    assert(mm.size() == 1);
    assert(mm.remove("Margaret Atwood", "The Handmaid's Tale")==true);
    assert(mm.size() == 0);
}

void Test::searchTest()
{
    MultiMap mm{ };
    assert(mm.size() == 0);
    mm.add("Daniel Keyes", "Flowers For Algernon");
    mm.add("Amy Harmon", "From Sand and Ash");
    mm.add("Jane Austen", "Emma");
    mm.add("Daniel Keyes", "The Minds of Billy Milligan");
    mm.add("Amy Harmon", "From Sand and Ash");
    mm.add("Jane Austen", "Persuasion");
    mm.add("Charlotte Bronte", "Jane Eyre");
    mm.add("Daniel Keyes", "Flowers For Algernon");
    mm.add("Amy Harmon", "From Sand and Ash");
    mm.add("Jane Austen", "Emma");
    mm.add("Daniel Keyes", "The Minds of Billy Milligan");
    mm.add("Amy Harmon", "From Sand and Ash");
    mm.add("Jane Austen", "Persuasion");
}

```

```
mm.add("Charlotte Bronte", "Jane Eyre");
mm.add("F. Scott Fitzgerald", "The Great Gatsby");
mm.add("Jane Austen", "Persuasion");
```

```
vector <string> v;
mm.search("Kristin Hannah",v);
assert(v.size() == 0);
```

```
mm.search("Daniel Keyes",v);
assert(v.size() == 4);
```

```
mm.search("Charlotte Bronte",v);
assert(v.size() == 2);
```

```
mm.search("Jane Austen",v);
assert(v.size() == 5);
```

```
mm.search("F. Scott Fitzgerald",v);
assert(v.size() == 1);
```

```
}
```

```
void Test::sizeTest()
```

```
{
    //contained in the others
}
```

```
void Test::isEmptyTest()
```

```
{
    MultiMap mm{ };
    assert(mm.size() == 0);
    assert(mm.isEmpty() == 1);
    mm.add("F. Scott Fitzgerald", "The Great Gatsby");
    assert(mm.size() == 1);
    assert(mm.isEmpty() == 0);
    mm.remove("F. Scott Fitzgerald", "The Great Gatsby");
    assert(mm.size() == 0);
    assert(mm.isEmpty() == 1);
    mm.add("Margaret Atwood", "The Handmaid's Tale");
    assert(mm.size() == 1);
    assert(mm.isEmpty() == 0);
}
```