CSEP 590 Data Compression Autumn 2007

Adaptive Huffman Coding

Adaptive Huffman Coding

- · One pass
- · During the pass calculate the frequencies
- · Update the Huffman tree accordingly
 - Coder new Huffman tree computed after transmitting the symbol
 - Decoder new Huffman tree computed after receiving the symbol
- Symbol set and their initial codes must be known ahead of time.
- Need NYT (not yet transmitted symbol) to indicate a new leaf is needed in the tree.

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Optimal Tree Numbering

• a:5, b:2, c:1, d:3



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Weight the Nodes

a:5, b:2, c:1, d:3



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Number the Nodes

• a:5, b:2, c:1, d:3



Number the nodes as they are removed from the priority queue.

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Adaptive Huffman Principle

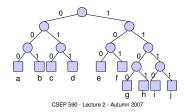
- In an optimal tree for n symbols there is a numbering of the nodes y₁<y₂<...<y_{2n-1} such that their corresponding weights x₁,x₂,..., x_{2n-1} satisfy:
 - $-\ \mathsf{X}_1 \leq \mathsf{X}_2 \leq \ldots \leq \mathsf{X}_{2\mathsf{n}\text{-}1}$
 - siblings are numbered consecutively
- And vice versa
 - That is, if there is such a numbering then the tree is optimal. We call this the node number invariant.

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Initialization

- Symbols a_1, a_2, \ldots, a_m have a basic prefix code, used when symbols are first encountered.
- Example: a, b ,c, d, e, f, g, h, i, j



Initialization

- The tree will encode up to m + 1 symbols including NYT.
- We reserve numbers 1 to 2m + 1 for node numbering.
- The initial Huffman tree consists of a single node
 weight

0 2m + 1 node number

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Coding Algorithm

- If a new symbol is encountered then output the code for NYT followed by the fixed code for the symbol. Add the new symbol to the tree.
- If an old symbol is encountered then output its code.
- Update the tree to preserve the node number invariant.

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Decoding Algorithm

- Decode the symbol using the current tree.
- If NYT is encountered then use the fixed code to decode the symbol. Add the new symbol to the tree
- Update the tree to preserve the node number invariant.

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Updating the Tree

- 1. Let y be leaf (symbol) with current weight x.*
- 2. If y the root update x by 1, otherwise,
- 4. Update x by 1
- 5. Let y be the parent with its weight x and go to 2.

*We never update the weight of NYT

** This exchange will preserve the node number invariant

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Example

• aabcdad in alphabet {a,b,..., j}

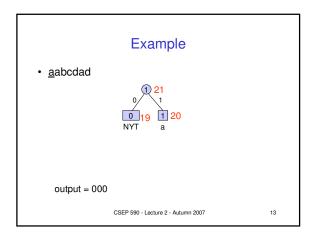
0 21 NYT

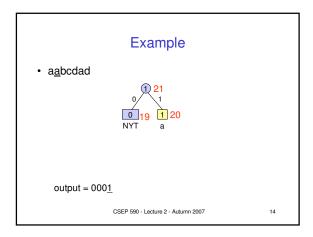
fixed code for a

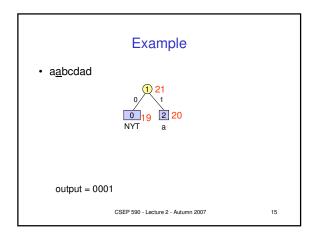
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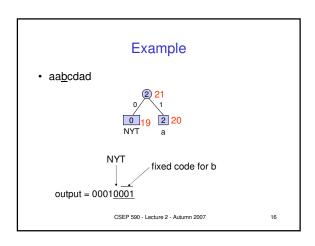
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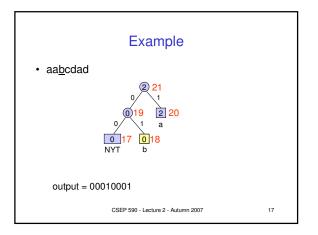
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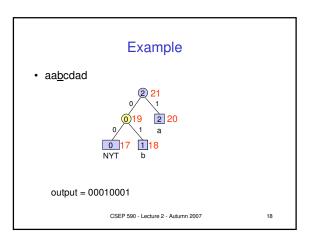


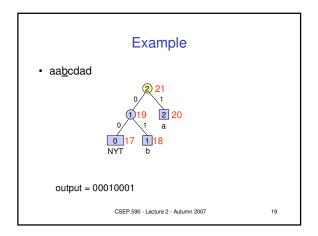


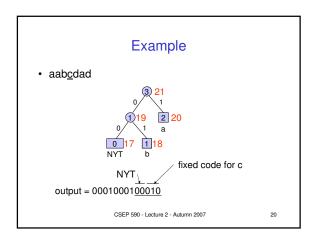


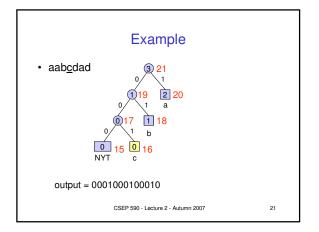


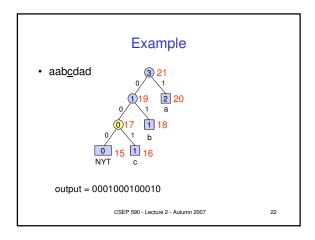


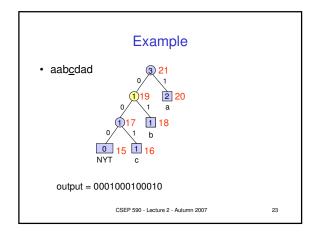


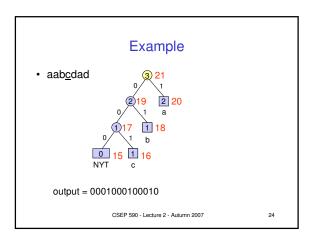


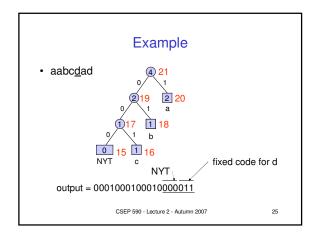


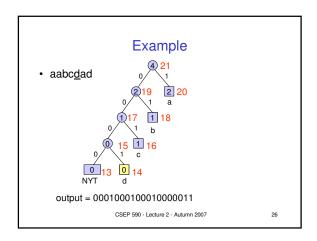


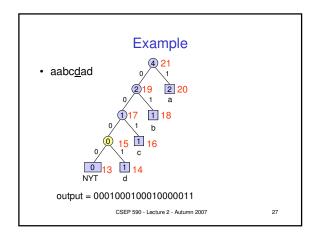


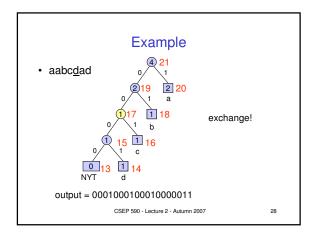


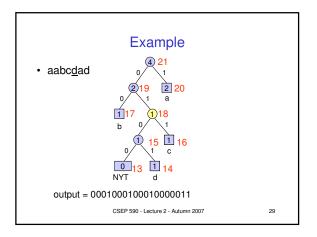


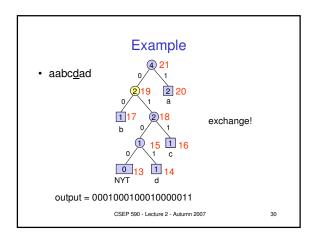


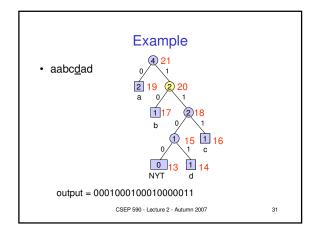


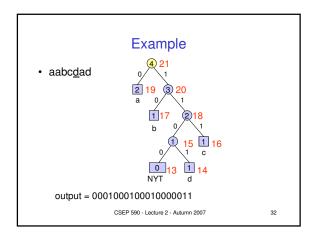


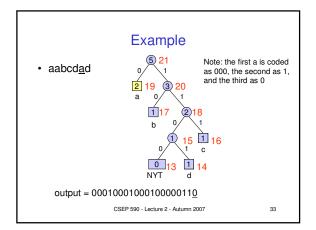


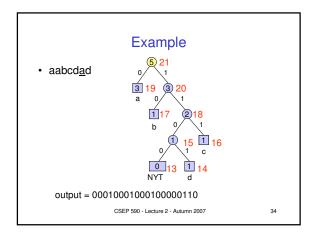


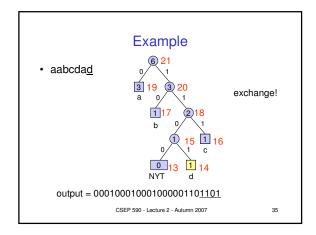


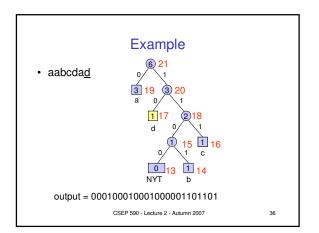


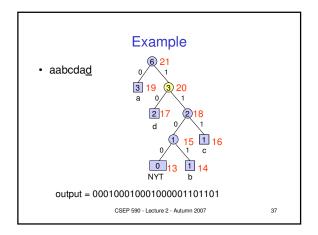


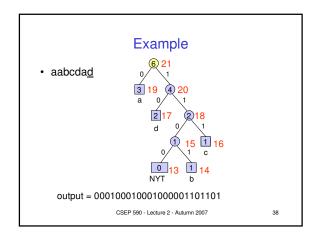


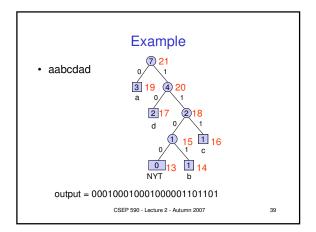


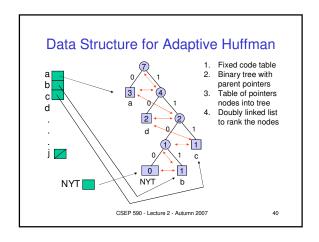


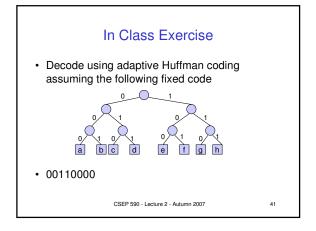












Huffman Summary

- Statistical compression algorithm
- Prefix code
- Fixed-to-variable length code
- Optimization to create a best code
- Symbol merging
- Context
- Adaptive coding
- Decoder and encoder behave almost the same
- Need for data structures and algorithms

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