$egin{array}{lll} x,\,y,\,p & {
m term \ variable} \\ \widehat{x},\,\widehat{x} & {
m unification \ term \ variable} \\ n,\,i,\,j,\,k & {
m index \ variables} \\ \end{array}$

```
Potentially non-ground terms
e, v, w
                                   ::=
                                            \boldsymbol{x}
                                            \hat{x}
                                            refl
                                            \langle \rangle
                                            v w
                                                                       \mathsf{bind}\ x\ \mathsf{in}\ v
                                            \lambda x. v
                                            \mathbf{rec}_{eq}^{x.p.\,A}(\,e\,,\,v\,)
                                                                       bind x in A
                                                                       bind p in A
                                            v:A
                                            [w/x]v
                                                                       Μ
                                            [w/\hat{x}]v
                                                                       Μ
                                                                       S
                                            (v)
A, B
                                            \mathbf{eq} \ A \ v \ w
                                            1
                                            \Pi x : A. B
                                                                       bind x in B
                                            \forall x: A. B
                                                                       \mathsf{bind}\;x\;\mathsf{in}\;B
                                                                       Μ
                                            [v/x]A
                                                                       S
                                            (A)
e, v, w, f, g, h
                                                                                             Ground value terms
                                            \boldsymbol{x}
                                            \langle \rangle
                                            v w
                                            \mathbf{refl}
                                            \lambda x. v
                                                                       \mathsf{bind}\ x\ \mathsf{in}\ v
                                            \mathbf{rec}_{eq}^{x.p.A}(e,v)
                                                                       \mathsf{bind}\ x\ \mathsf{in}\ A
                                                                       bind p in A
                                            v:A
                                            \lceil w/x \rceil v
                                                                       Μ
                                            (v)
                                                                       S
X, A, B, C, D
                                            \mathbf{eq}\,A\;v\;w
                                            1
                                                                       \mathsf{bind}\ x\ \mathsf{in}\ B
                                            \Pi x : A.B
                                            \forall x: A. B
                                                                       \mathsf{bind}\ x\ \mathsf{in}\ B
                                            [v/x]A
                                                                       Μ
                                            (A)
                                                                       S
terminals
```

```
\Downarrow
Γ
                                                           declarative variable context
                                                              add to context
                      x:A
                      \overline{\Gamma_i}^{i}
                                                              concatenate contexts
                                                      Μ
                                                              empty context
Θ
                                                           computational variable context
                      x:A
                                                              a variable
                      \hat{x}:A
                                                              a polymorphic variable
                      \hat{x} = v
                                                              instantiate a polymorphic variable
                      \overline{\Theta_i}^{\;i}
                                                              concatenate contexts
                                                      Μ
                                                              empty context
                                                           extra judgements for explicit and inference typing systems
ty_-extra
                      x:A\in\Gamma
                                                              lookup type of x in context \Gamma
formula
                      judgement
                      x:A\in\Gamma
                                                              lookup type of x in context \Gamma
                      x:A\in\Theta
                                                              lookup type of x in context \Theta
                      ut1 \neq ut2
                      v_1 \neq v_2
                      t1 \neq t2
                      v_1 \neq v_2
                      A_1 \neq A_2
                      formula_1 .. formula_n
ok
                      x \in A ok
                      x \in ? A ok
                      \widehat{x} \in ^! v \ ok
                      \hat{x} \in ? v ok
```

De

::=

	$\Gamma \vdash A \unlhd B$	Top-level polymorphic instantiation
İ	$\Gamma \vdash v \Leftarrow A$	Declarative checking
	$\Gamma \vdash v \Rightarrow A$	Declarative inference
	$\Gamma \vdash A$	Declarative type well-formedness
	$\Gamma \vdash A \equiv B$	Declarative type equivalence
	$\Gamma \vdash v \equiv w \Rightarrow A$	Declarative inferring equivalence
	$\Gamma \vdash v \equiv w \Leftarrow A$	Declarative checking equivalence
	$\Gamma \vdash v \equiv_n w \Leftarrow A$	Declarative checking equivalence for terms in WHNF
	$\Gamma \vdash v \equiv_n w \Rightarrow A$	Declarative inferring equivalence for terms in WHNF
	$\Gamma \vdash A \leqslant B$	Declarative subtyping
	$\Gamma \vdash v \searrow w \Rightarrow A$	Declarative inferring CBN reduction to WHNF
	$\Gamma \vdash v \searrow w \Leftarrow A$	Declarative checking reduction to WHNF

Al ::=

 $\begin{array}{ccc} judgement & ::= & \\ & | & ok \\ & | & De \end{array}$

 $user_syntax$::=

 $\begin{array}{c|cccc} & x & \\ & \widehat{x} & \\ & n & \\ & e & \\ & A & \\ & e & \\ & X & \\ & terminals & \\ & \Gamma & \\ & \Theta & \\ & ty_extra & \\ & formula & \end{array}$

 $x \in A \ ok$ $x \in A \ ok$ $\hat{x} \in v \ ok$ $\hat{x} \in v \ ok$

 $\overline{\Gamma \vdash A \subseteq B}$ Top-level polymorphic instantiation

$$\frac{\Gamma \vdash w \Leftarrow A \quad \Gamma \vdash [w/x]B \unlhd C}{\Gamma \vdash \forall x : A. B \unlhd C} \quad \text{De_inst_Forall}$$

$$\frac{A \neq \forall x : B. C}{\Gamma \vdash A \unlhd A} \quad \text{De_inst_Base}$$

 $\Gamma \vdash v \leftarrow A$ Declarative checking

$$\begin{array}{ll} \Gamma \vdash A & \Gamma, x : A \vdash v \Leftarrow B \\ \hline \Gamma \vdash v \Leftarrow \forall x : A . \, B \\ \hline \\ \frac{\Gamma \vdash A & \Gamma, x : A \vdash v \Leftarrow B}{\Gamma \vdash \lambda x . \, v \Leftarrow \Pi x : A . \, B} \end{array} \text{ De_Check_Pi}$$

$$\frac{\Gamma \vdash v \equiv w \Leftarrow A}{\Gamma \vdash \mathbf{refl} \Leftarrow \mathbf{eq} \ A \ v \ w} \quad \text{De_check_Refl}$$

$$\frac{\Gamma \vdash v \Rightarrow A \quad \Gamma \vdash A \leqslant B}{\Gamma \vdash v \Leftarrow B} \quad \text{De_check_Sub}$$

 $\Gamma \vdash v \Rightarrow A$ Declarative inference

$$\frac{x:A\in\Gamma}{\Gamma\vdash x\Rightarrow A} \quad \text{De_Inf_Var}$$

$$\frac{x:A\in\Gamma}{\Gamma\vdash x\Rightarrow A} \quad \text{De_Inf_Var}$$

$$\frac{\Gamma\vdash v\Rightarrow A \quad \Gamma\vdash A\trianglelefteq\Pi x:B.C}{\Gamma\vdash v\ w\Rightarrow [w/x]B} \quad \text{De_Inf_App}$$

$$\frac{\Gamma\vdash A \quad \Gamma\vdash v\Leftarrow A}{\Gamma\vdash v:A\Rightarrow A} \quad \text{De_Inf_Ann}$$

$$\frac{\Gamma \vdash e \Rightarrow \operatorname{eq} A \, v \, w \quad \Gamma, x : A, p : \operatorname{eq} A \, v \, x \vdash X \quad \Gamma \vdash f \Leftarrow [\operatorname{refl}/p][v/x]X}{\Gamma \vdash \operatorname{rec}_{eq}^{x.p.X}(e,f) \Rightarrow [e/p][w/x]X} \quad \text{Delinf_Receq}$$

 $\overline{\Gamma \vdash A}$ Declarative type well-formedness

$$\frac{\Gamma \vdash A \quad \Gamma \vdash v \Leftarrow A \quad \Gamma \vdash w \Leftarrow A}{\Gamma \vdash \mathbf{eq} \, A \, v \, w} \quad \text{De_wfT_EQ}$$

$$\frac{\Gamma \vdash A \quad \Gamma, x : A \vdash B}{\Gamma \vdash \Pi x : A . \, B} \quad \text{De_wfT_PI}$$

$$\frac{\Gamma \vdash A \quad \Gamma, x : A \vdash B \quad x \in B \quad De_wfT_PI}{\Gamma \vdash \Pi x : A . \, B} \quad \text{De_wfT_Forall}$$

 $\Gamma \vdash A \equiv B$ Declarative type equivalence

$$\begin{array}{c} \Gamma \vdash 1 \equiv 1 \\ \hline \Gamma \vdash A_1 \equiv A_2 \quad \Gamma, x : A_1 \vdash B_1 \equiv B_2 \\ \hline \Gamma \vdash \Pi x : A_1. \ B_1 \equiv \Pi x : A_2. \ B_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash A_1 \equiv A_2 \quad \Gamma, x : A_1 \vdash B_1 \equiv B_2 \\ \hline \Gamma \vdash \Pi x : A_1. \ B_1 \equiv \Pi x : A_2. \ B_2 \\ \hline \hline \Gamma \vdash A_1 \equiv A_2 \quad \Gamma, x : A_1 \vdash B_1 \equiv B_2 \\ \hline \Gamma \vdash \forall x : A_1. \ B_1 \equiv \forall x : A_2. \ B_2 \\ \hline \hline \end{array} \quad \begin{array}{c} \Gamma \vdash A_1 \equiv A_2 \quad \Gamma \vdash v_1 \equiv v_2 \Leftarrow A_1 \quad \Gamma \vdash w_1 \equiv w_2 \Leftarrow A_1 \\ \hline \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ w_1 \equiv \mathbf{eq} \ A_2 \ v_2 \ w_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_2 \ v_2 \ v_3 \ v_1 \ v_2 \ v_2 \ v_2 \\ \hline \end{array} \quad \begin{array}{c} \Gamma \vdash \mathbf{eq} \ A_1 \ v_1 \ v_2 \ v_2 \ v_3 \ v_3 \ v_2 \ v_3 \\ \hline \end{array}$$

 $\Gamma \vdash v \equiv w \Rightarrow A$ Declarative inferring equivalence

$$\frac{\Gamma \vdash v \searrow v' \Rightarrow A \quad \Gamma \vdash w \searrow w' \Rightarrow A \quad \Gamma \vdash v' \equiv_n w' \Rightarrow A}{\Gamma \vdash v \equiv w \Rightarrow A} \quad \text{De_equivInf_Red}$$

 $\Gamma \vdash v \equiv w \leftarrow A$ Declarative checking equivalence

$$\frac{\Gamma \vdash v \searrow v' \Leftarrow A \quad \Gamma \vdash w \searrow w' \Leftarrow A \quad \Gamma \vdash v' \equiv_n w' \Leftarrow A}{\Gamma \vdash v \equiv w \Leftarrow A} \quad \text{De_equivChe_Red}$$

 $\overline{\Gamma \vdash v \equiv_n w \Leftarrow A}$ Declarative checking equivalence for terms in WHNF

$$\frac{\Gamma \vdash e \searrow e' \Rightarrow \operatorname{eq} A \, v \, w \quad e' \neq \operatorname{refl} : \operatorname{eq} A \, v \, w}{\Gamma \vdash \operatorname{rec}_{eq}^{x.p.X}(e,f) \searrow \operatorname{rec}_{eq}^{x.p.X}(e',f) \Rightarrow [e/p][w/x]X} \quad \text{De_RedInf_Rec}$$

 $\Gamma \vdash v \setminus w \leftarrow A$ Declarative checking reduction to WHNF

Definition rules: 45 good 0 bad Definition rule clauses: 81 good 0 bad