

$\alpha, \beta, \alpha, \beta, \gamma, \delta$	type variables
n, m, i, j	index variables
x, y, z	term variables

	$ \begin{array}{l} \cdot \\ e \\ \hat{\sigma} \backslash vars \\ \hat{\sigma} vars \\ \hat{\sigma}_1 \cup \hat{\sigma}_2 \\ \widehat{\sigma}_i^i \\ (\hat{\sigma}) \quad \text{S} \\ \mathbf{nf}(\hat{\sigma}') \quad \text{M} \\ \hat{\sigma}' vars \quad \text{M} \\ \hat{\sigma}_1 \ \& \ \hat{\sigma}_2 \quad \text{M} \end{array} $	concatenate
$\hat{\tau}, \hat{\rho}$	$ \begin{array}{l} \cdot \\ \hat{\alpha}^- : \approx N \\ \hat{\alpha}^- : \approx N \\ \overrightarrow{\alpha^-} / \overrightarrow{\alpha^-} \\ \overrightarrow{N} / \overrightarrow{\alpha^-} \\ \hat{\tau}_1 \cup \hat{\tau}_2 \\ \widehat{\tau}_i^i \\ (\hat{\tau}) \quad \text{S} \\ \hat{\tau}' vars \quad \text{M} \\ \hat{\tau}_1 \ \& \ \hat{\tau}_2 \quad \text{M} \end{array} $	anti-unification substitution concatenate
P, Q, R	$ \begin{array}{l} \alpha^+ \\ \downarrow N \\ \exists \alpha^-. P \\ [\sigma]P \quad \text{M} \end{array} $	positive types
N, M, K	$ \begin{array}{l} \alpha^- \\ \uparrow P \\ \forall \alpha^+. N \\ P \rightarrow N \\ [\sigma]N \quad \text{M} \end{array} $	negative types
$\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}, \overrightarrow{\gamma^+}, \overrightarrow{\delta^+}$	$ \begin{array}{l} \cdot \\ \alpha^+ \\ \overrightarrow{\alpha^+} \\ \overrightarrow{\overrightarrow{\alpha^+}}^i \\ \alpha^+{}_i \end{array} $	positive variable list empty list a variable a variable concatenate lists
$\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}, \overrightarrow{\gamma^-}, \overrightarrow{\delta^-}$	$ \begin{array}{l} \cdot \\ \alpha^- \\ \overrightarrow{\alpha^-} \\ \overrightarrow{\overrightarrow{\alpha^-}}^i \\ \alpha^-{}_i \end{array} $	negative variables empty list a variable variables concatenate lists
$\overrightarrow{\alpha^\pm}, \overrightarrow{\beta^\pm}, \overrightarrow{\gamma^\pm}, \overrightarrow{\delta^\pm}$	$ \begin{array}{l} \cdot \\ \alpha^\pm \\ \overrightarrow{\alpha^\pm} \\ \overrightarrow{\overrightarrow{\alpha^\pm}}^i \\ \alpha^\pm{}_i \end{array} $	positive or negative variable list

	\cdot α^\pm $\vec{\mathbf{p}}\vec{\mathbf{a}}$ $\overrightarrow{\alpha^\pm}_i$	empty list a variable variables concatenate lists
P, Q, R	$::=$ α^+ $\downarrow N$ $\exists \alpha^-.P$ $[\sigma]P$ $[\hat{\tau}]P$ $[\hat{\sigma}]P$ $[\mu]P$ (P) $P_1 \vee P_2$ $\mathbf{nf}(P')$	multi-quantified positive types $P \neq \exists \dots$ M M M M S M M
N, M, K	$::=$ α^- $\uparrow P$ $P \rightarrow N$ $\forall \alpha^+.N$ $[\sigma]N$ $[\hat{\tau}]N$ $[\mu]N$ $[\hat{\sigma}]N$ (N) $\mathbf{nf}(N')$	multi-quantified negative types $N \neq \forall \dots$ M M M M S M
\vec{P}, \vec{Q}	$::=$ \cdot P $[\sigma]\vec{P}$ $\overrightarrow{\vec{P}}_i$ $\mathbf{nf}(\vec{P}')$	list of positive types empty list a singel type M concatenate lists M
\vec{N}, \vec{M}	$::=$ \cdot N $[\sigma]\vec{N}$ $\overrightarrow{\vec{N}}_i$ $\mathbf{nf}(\vec{N}')$	list of negative types empty list a singel type M concatenate lists M
Δ, Γ	$::=$ \cdot $\overrightarrow{\alpha^+}$ $\overrightarrow{\alpha^-}$ $\overrightarrow{\alpha^\pm}$ $vars$ $\overrightarrow{\Gamma}_i$	declarative type context empty context list of variables list of variables list of variables concatenate contexts

		(Γ)	S	
		$\Theta(\hat{\alpha}^+)$	M	
		$\Theta(\hat{\alpha}^-)$	M	
Θ	::=			unification type variable context
		\cdot		empty context
		$\vec{\alpha}\{\Delta\}$		from an ordered list of variables
		$\hat{\alpha}^+\{\Delta\}$		from a variable to a list
		$\overline{\Theta}_i^i$		concatenate contexts
		(Θ)	S	
		$\Theta _{vars}$		leave only those variables that are in the set
		$\Theta_1 \cup \Theta_2$		
Ξ	::=			anti-unification type variable context
		\cdot		empty context
		$\vec{\alpha}^-$		list of variables
		$\overline{\Xi}_i^i$		concatenate contexts
		(Ξ)	S	
		$\Xi_1 \cup \Xi_2$		
		$\Xi_1 \cap \Xi_2$		
		$\Xi' _{vars}$	M	
$\vec{\alpha}, \vec{\beta}$::=			ordered positive or negative variables
		\cdot		empty list
		$\vec{\alpha}^+$		list of variables
		$\vec{\alpha}^-$		list of variables
		$\vec{\alpha}^\pm$		list of variables
		$\vec{\alpha}^+ \vec{\alpha}^-$		list of variables
		$\vec{\alpha}^-$		list of variables
		$\vec{\alpha}_1 \setminus vars$		setminus
		Γ		context
		$vars$		
		$\vec{\alpha}_i^i$		concatenate contexts
		$(\vec{\alpha})$	S	parenthesis
		$[\mu]\vec{\alpha}$	M	apply moving to list
		ord vars in P	M	
		ord vars in N	M	
		ord vars in P	M	
		ord vars in N	M	
$vars$::=			set of variables
		\emptyset		empty set
		fv P		free variables
		fv N		free variables
		fv imP		free variables
		fv imN		free variables
		$vars_1 \cap vars_2$		set intersection
		$vars_1 \cup vars_2$		set union
		$vars_1 \setminus vars_2$		set complement
		mv imP		movable variables

		mv imN		movable variables
		uv N		unification variables
		uv P		unification variables
		fv N		free variables
		fv P		free variables
		$(vars)$	S	parenthesis
		$\vec{\alpha}$		ordered list of variables
		$[\mu]vars$	M	apply moving to varset
		dom $(\hat{\sigma})$	M	
		dom $(\hat{\tau})$	M	
		dom (Θ)	M	
μ	::=			
		.		empty moving
		$pma1 \mapsto pma2$		Positive unit substitution
		$nma1 \mapsto nma2$		Positive unit substitution
		$\mu_1 \cup \mu_2$	M	Set-like union of movings
		$\mu_1 \circ \mu_2$	M	Composition
		$\overline{\mu_i}^i$		concatenate movings
		$\mu vars$	M	restriction on a set
		μ^{-1}	M	inversion
		nf (μ')	M	
$\hat{\alpha}^\pm$::=			positive/negative unification variable
		$\hat{\alpha}^\pm$		
$\hat{\alpha}^+$::=			positive unification variable
		$\hat{\alpha}^+$		
		$\hat{\alpha}^+\{\Delta\}$		
		$\hat{\alpha}^\pm$		
$\hat{\alpha}^-, \hat{\beta}^-$::=			negative unification variable
		$\hat{\alpha}^-$		
		$\hat{\alpha}^-_{\{N,M\}}$		
		$\hat{\alpha}^- \{\Delta\}$		
		$\hat{\alpha}^\pm$		
$\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}$::=			positive unification variable list
		.		empty list
		$\hat{\alpha}^+$		a variable
		$\overrightarrow{\hat{\alpha}^+}$		from a normal variable, context unspecified
		$\overrightarrow{\overrightarrow{\alpha^+}}^i$		
		α^+_i		concatenate lists
$\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}$::=			negative unification variable list
		.		empty list
		$\hat{\alpha}^-$		a variable
		Ξ		from an antiunification context
		$\hat{\alpha}^- \{\Delta\}$		from a normal variable
		$\overrightarrow{\hat{\alpha}^-}$		from a normal variable, context unspecified

	\xrightarrow{i} α^-_i	concatenate lists
P, Q	$::=$ α^+ \mathbf{pma} $\hat{\alpha}^+$ $\downarrow N$ $\xrightarrow{\quad} \exists \alpha^- . P$ $[\sigma] P$ $[\hat{\tau}] P$ $[\mu] P$ (P) $\mathbf{nf}(P')$	 a positive algorithmic type (potentially with metavariables) M M M S M
N, M	$::=$ α^- $\hat{\alpha}^-$ $\uparrow P$ $P \rightarrow N$ $\xrightarrow{\quad} \forall \alpha^+ . N$ $[\sigma] N$ $[\hat{\tau}] N$ $[\mu] N$ (N) $\mathbf{nf}(N')$	 a negative algorithmic type (potentially with metavariables) M M M S M
$auSol$	$::=$ $(\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ $(\Xi, N, \hat{\tau}_1, \hat{\tau}_2)$	
$terminals$	$::=$ \exists \forall \uparrow \downarrow \rightarrow \leftrightarrow \in \notin \cdot \top \leq \geq \approx \subset \supset \setminus \sqcup \mapsto \approx^u	

	\sim \emptyset \circ \Rightarrow \models \models \neq \equiv_n \vee \Downarrow $:\geq$ $:\simeq$ Λ λ \mathbf{let}^\exists \bullet $\Rightarrow\Rightarrow$	
v, w	$::=$ $ $ $ $ $ $ $ $	value terms M
\vec{v}	$::=$ $ $ $ $ $ $	list of arguments concatenate
c, d	$::=$ $ $ $ $ $ $ $ $ $ $ $ $ $ $	computation terms
$vctx, \Phi$	$::=$ $ $ $ $ $ $	variable context concatenate contexts
$formula$	$::=$ $ $ $ $ $ $ $ $ $ $	<i>judgement</i> <i>judgement uniquely</i> $formula_1 \dots formula_n$ $\mu : vars_1 \leftrightarrow vars_2$ μ is bijective \hat{o} is functional

	$\hat{\sigma}_1 \in \hat{\sigma}_2$ $v : P \in \Phi$ $\hat{\sigma}_1 \subseteq \hat{\sigma}_2$ $vars_1 \subseteq vars_2$ $vars_1 = vars_2$ $vars$ is fresh $\alpha^- \notin vars$ $\alpha^+ \notin vars$ $\alpha^- \in vars$ $\alpha^+ \in vars$ $\hat{\alpha}^+ \in vars$ $\hat{\alpha}^- \in vars$ $\hat{\alpha}^- \in \Theta$ $\hat{\alpha}^+ \in \Theta$ if any other rule is not applicable $\vec{\alpha}_1 = \vec{\alpha}_2$ $e_1 = e_2$ $N = M$ $N \neq M$ $P \neq Q$ $N \neq M$ $P \neq Q$ $P \neq Q$ $N \neq M$	
A	$::=$ $\Gamma; \Theta \models N \leq M \Rightarrow \hat{\sigma}$ $\Gamma; \Theta \models P \geq Q \Rightarrow \hat{\sigma}$	Negative subtyping Positive supertyping
AT	$::=$ $\Gamma; \Phi \models v : P$ $\Gamma; \Phi \models c : N$ $\Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \Rightarrow \hat{\sigma}$	Positive type inference Negative type inference Application type inference
AU	$::=$ $\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \Rightarrow (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ $\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \Rightarrow (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$	
$E1$	$::=$ $N \stackrel{D}{\simeq}_1 M$ $P \stackrel{D}{\simeq}_1 Q$ $P \simeq Q$	Negative multi-quantified type equivalence Positive multi-quantified type equivalence
$D1$	$::=$ $\Gamma \vdash N \simeq_1^{\leq} M$ $\Gamma \vdash P \simeq_1^{\leq} Q$ $\Gamma \vdash N \leq_1 M$ $\Gamma \vdash P \geq_1 Q$ $\Gamma_2 \vdash \sigma_1 \simeq_1^{\leq} \sigma_2 : \Gamma_1$	Negative equivalence on MQ types Positive equivalence on MQ types Negative subtyping Positive supertyping Equivalence of substitutions

$D0$	$::=$ $\mid \Gamma \vdash N \simeq_0^< M$ $\mid \Gamma \vdash P \simeq_0^< Q$ $\mid \Gamma \vdash N \leq_0 M$ $\mid \Gamma \vdash P \geq_0 Q$	Negative equivalence Positive equivalence Negative subtyping Positive supertyping
DT	$::=$ $\mid \Gamma; \Phi \vdash v : P$ $\mid \Gamma; \Phi \vdash c : N$ $\mid \Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M$	Positive type inference Negative type inference Spin Application type inference
EQ	$::=$ $\mid N = M$ $\mid P = Q$ $\mid P = \boxed{Q}$	Negative type equality (alpha-equivalence) Positive type equality (alpha-equivalence)
$LUBF$	$::=$ $\mid P_1 \vee P_2 === Q$ $\mid \mathbf{ord\,vars\,in}\, P === \vec{\alpha}$ $\mid \mathbf{ord\,vars\,in}\, N === \vec{\alpha}$ $\mid \mathbf{ord\,vars\,in}\, P === \vec{\alpha}$ $\mid \mathbf{ord\,vars\,in}\, N === \vec{\alpha}$ $\mid \mathbf{nf}\,(N') === N$ $\mid \mathbf{nf}\,(P') === P$ $\mid \mathbf{nf}\,(N') === \boxed{N}$ $\mid \mathbf{nf}\,(P') === \boxed{P}$ $\mid \mathbf{nf}\,(\vec{N}') === \vec{N}$ $\mid \mathbf{nf}\,(\vec{P}') === \vec{P}$ $\mid \mathbf{nf}\,(\sigma') === \sigma$ $\mid \mathbf{nf}\,(\mu') === \mu$ $\mid \mathbf{nf}\,(\hat{\sigma}') === \hat{\sigma}$ $\mid \sigma' _{vars}$ $\mid \hat{\sigma}' _{vars}$ $\mid \hat{\tau}' _{vars}$ $\mid \Xi' _{vars}$ $\mid e_1 \ \&\; e_2$ $\mid \hat{\sigma}_1 \ \&\; \hat{\sigma}_2$ $\mid \hat{\tau}_1 \ \&\; \hat{\tau}_2$ $\mid \mathbf{dom}\,(\hat{\sigma}) === vars$ $\mid \mathbf{dom}\,(\hat{\tau}) === vars$ $\mid \mathbf{dom}\,(\Theta) === vars$	
LUB	$::=$ $\mid \Gamma \models P_1 \vee P_2 = Q$ $\mid \mathbf{upgrade}\,\Gamma \vdash P \mathbf{to}\, \Delta = Q$	Least Upper Bound (Least Common Supertype)
Nrm	$::=$ $\mid \mathbf{nf}\,(N) = M$ $\mid \mathbf{nf}\,(P) = Q$ $\mid \mathbf{nf}\,(\boxed{N}) = \boxed{M}$	

		$\mathbf{nf}(P) = Q$	
<i>Order</i>	$::=$	$\mathbf{ord\ vars\ in\ } N = \vec{\alpha}$ $\mathbf{ord\ vars\ in\ } P = \vec{\alpha}$ $\mathbf{ord\ vars\ in\ } N = \vec{\alpha}$ $\mathbf{ord\ vars\ in\ } P = \vec{\alpha}$	
<i>SM</i>	$::=$	$\Gamma \vdash e_1 \ \& \ e_2 = e_3$ $\Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3$	Unification Solution Entry Merge Merge unification solutions
<i>SImp</i>	$::=$	$\Gamma \vdash e_1 \Rightarrow e_2$ $\Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2$ $\Gamma \vdash e_1 \simeq e_2$ $\Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2$	Weakening of unification solution entries Weakening of unification solutions
<i>U</i>	$::=$	$\Gamma; \Theta \vdash N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}$ $\Gamma; \Theta \vdash P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}$	Negative unification Positive unification
<i>WF</i>	$::=$	$\Gamma \vdash N$ $\Gamma \vdash P$ $\Gamma \vdash N$ $\Gamma \vdash P$ $\Gamma \vdash \vec{N}$ $\Gamma \vdash \vec{P}$ $\Gamma; \Theta \vdash N$ $\Gamma; \Theta \vdash P$ $\Gamma; \Xi \vdash N$ $\Gamma; \Xi \vdash P$ $\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1$ $\hat{\sigma} : \Theta$ $\Gamma \vdash^= \Theta$ $\Gamma_1 \vdash \sigma : \Gamma_2$ $\Gamma \vdash e$	Negative type well-formedness Positive type well-formedness Negative type well-formedness Positive type well-formedness Negative type list well-formedness Positive type list well-formedness Negative unification type well-formedness Positive unification type well-formedness Negative anti-unification type well-formedness Positive anti-unification type well-formedness Antiunification substitution well-formedness Unification substitution well-formedness Unification context well-formedness Substitution well-formedness Unification solution entry well-formedness
<i>judgement</i>	$::=$	A AT AU $E1$ $D1$ $D0$ DT EQ LUB	

$$\begin{array}{c}
\frac{}{\Gamma; \Theta \models \alpha^- \leq \alpha^- =} \text{ANVAR} \\
\frac{\Gamma; \Theta \models \mathbf{nf}(P) \stackrel{u}{\simeq} \mathbf{nf}(Q) = \hat{\sigma}}{\Gamma; \Theta \models \uparrow P \leq \uparrow Q = \hat{\sigma}} \text{AShiftU} \\
\frac{\Gamma; \Theta \models P \geq Q = \hat{\sigma}_1 \quad \Gamma; \Theta \models N \leq M = \hat{\sigma}_2 \quad \Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}}{\Gamma; \Theta \models P \rightarrow N \leq Q \rightarrow M = \hat{\sigma}} \text{AArrow} \\
\frac{\text{<<multiple parses>>}}{\Gamma; \Theta \models \forall \alpha^+. N \leq \forall \beta^+. M = \hat{\sigma} \setminus \hat{\alpha}^+} \text{Aforall} \\
\boxed{\Gamma; \Theta \models P \geq Q = \hat{\sigma}} \quad \text{Positive supertyping}
\end{array}$$

$$\begin{array}{c}
\frac{}{\Gamma; \Theta \models \alpha^+ \geq \alpha^+ =} \text{APVAR} \\
\frac{\Gamma; \Theta \models \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) = \hat{\sigma}}{\Gamma; \Theta \models \downarrow N \geq \downarrow M = \hat{\sigma}} \text{AShiftD} \\
\frac{\Gamma, \vec{\beta}^-; \Theta, \hat{\alpha}^- \{ \Gamma, \vec{\beta}^- \} \models [\hat{\alpha}^- / \alpha^-] P \geq Q = \hat{\sigma}}{\Gamma; \Theta \models \exists \alpha^+. P \geq \exists \beta^+. Q = \hat{\sigma} \setminus \hat{\alpha}^-} \text{Aexists} \\
\frac{\hat{\alpha}^+ \{ \Delta \} \in \Theta \quad \mathbf{upgrade} \Gamma \vdash P \text{ to } \Delta = Q}{\Gamma; \Theta \models \hat{\alpha}^+ \geq P = (\hat{\alpha}^+ : \geq Q)} \text{APUVar} \\
\boxed{\Gamma; \Phi \models v : P} \quad \text{Positive type inference}
\end{array}$$

$$\begin{array}{c}
\frac{v : P \in \Phi}{\Gamma; \Phi \models v : P} \text{ATVar} \\
\frac{\Gamma; \Phi \models c : N}{\Gamma; \Phi \models \{c\} : \downarrow N} \text{ATTThunk} \\
\frac{\Gamma; \Phi \models v : P \quad \Gamma; \cdot \models Q \geq P = \cdot}{\Gamma; \Phi \models (v : Q) : Q} \text{ATAnnot}
\end{array}$$

$\boxed{\Gamma; \Phi \models c : N}$ Negative type inference

$$\begin{array}{c}
\frac{\Gamma; \Phi \models c : N \quad \Gamma; \cdot \models N \leq M = \cdot}{\Gamma; \Phi \models (c : M) : M} \text{ATAnnotN} \\
\frac{\Gamma; \Phi, x : P \models c : N}{\Gamma; \Phi \models \lambda x : P. c : P \rightarrow N} \text{ATTLam} \\
\frac{\Gamma, \alpha^+; \Phi \models c : N}{\Gamma; \Phi \models \Lambda \alpha^+. c : \forall \alpha^+. N} \text{ATTlam} \\
\frac{\Gamma; \Phi \models v : P}{\Gamma; \Phi \models \mathbf{return} v : \uparrow P} \text{ATReturn} \\
\frac{\Gamma; \Phi \models v : \downarrow M \quad \Gamma; \Phi; \cdot \models M \bullet \vec{v} \Rightarrow \uparrow Q = \hat{\sigma}_1 \quad \Gamma; \mathbf{uv} Q \{ \Gamma \} \models \uparrow Q \leq \uparrow P = \hat{\sigma}_2 \quad \mathbf{uv} Q \{ \Gamma \} \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma} \quad \Gamma; \Phi, x : P \models c : N}{\Gamma; \Phi \models \mathbf{let} x : P = v(\vec{v}); c : N} \text{ATLetAnn} \\
\frac{\Gamma; \Phi \models v : \downarrow M \quad \Gamma; \Phi; \cdot \models M \bullet \vec{v} \Rightarrow \uparrow Q = \hat{\sigma} \quad \mathbf{uv}(Q) = \emptyset \quad \Gamma; \Phi, x : Q \models c : N}{\Gamma; \Phi \models \mathbf{let} x = v(\vec{v}); c : N} \text{ATLet}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma, \alpha^-; \Phi \models v : \exists \alpha^-. P \quad \Gamma, \alpha^-; \Phi, x : P \models c : N \quad \Gamma \vdash N}{\Gamma; \Phi \models \mathbf{let}^\exists(\alpha^-, x) = v; c : N} \text{ ATUNPACK} \\
\\
\boxed{\Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \models \hat{\sigma}} \quad \text{Application type inference} \\
\\
\frac{N \neq \forall \alpha^+. M}{\Gamma; \Phi; \Theta \models N \bullet \cdot \Rightarrow N \models \cdot} \text{ ATEMTPTY} \\
\\
\frac{\Gamma; \Phi \models v : P \quad \Gamma; \Theta \models Q \succcurlyeq P \models \hat{\sigma}_1 \quad \Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \models \hat{\sigma}_2}{\Gamma; \Phi; \Theta \models Q \rightarrow N \bullet v, \vec{v} \Rightarrow M \models \hat{\sigma}_1 \& \hat{\sigma}_2} \text{ ATARROW} \\
\\
\frac{\text{<<multiple parses>>}}{\Gamma; \Phi; \Theta \models \forall \alpha^+. N \bullet \vec{v} \Rightarrow M \models \hat{\sigma}} \text{ ATFORALL} \\
\\
\boxed{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)} \\
\\
\frac{}{\Gamma \models \alpha^+ \stackrel{a}{\simeq} \alpha^+ \models (\cdot, \alpha^+, \cdot, \cdot)} \text{ AUPVAR} \\
\\
\frac{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N_1 \stackrel{a}{\simeq} \downarrow N_2 \models (\Xi, \downarrow M, \hat{\tau}_1, \hat{\tau}_2)} \text{ AUSHIFTD} \\
\\
\frac{\vec{\alpha}^+ \cap \Gamma = \emptyset \quad \Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \exists \alpha^-. P_1 \stackrel{a}{\simeq} \exists \alpha^-. P_2 \models (\Xi, \exists \alpha^-. Q, \hat{\tau}_1, \hat{\tau}_2)} \text{ AUEXISTS} \\
\\
\boxed{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)} \\
\\
\frac{}{\Gamma \models \alpha^- \stackrel{a}{\simeq} \alpha^- \models (\cdot, \alpha^-, \cdot, \cdot)} \text{ AUNVAR} \\
\\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \uparrow P_1 \stackrel{a}{\simeq} \uparrow P_2 \models (\Xi, \uparrow Q, \hat{\tau}_1, \hat{\tau}_2)} \text{ AUSHIFTU} \\
\\
\frac{\vec{\alpha}^+ \cap \Gamma = \emptyset \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \forall \alpha^+. N_1 \stackrel{a}{\simeq} \forall \alpha^+. N_2 \models (\Xi, \forall \alpha^+. M, \hat{\tau}_1, \hat{\tau}_2)} \text{ AUFORALL} \\
\\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi_1, Q, \hat{\tau}_1, \hat{\tau}_2) \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi_2, M, \hat{\tau}'_1, \hat{\tau}'_2)}{\Gamma \models P_1 \rightarrow N_1 \stackrel{a}{\simeq} P_2 \rightarrow N_2 \models (\Xi_1 \cup \Xi_2, Q \rightarrow M, \hat{\tau}_1 \cup \hat{\tau}'_1, \hat{\tau}_2 \cup \hat{\tau}'_2)} \text{ AUARROW} \\
\\
\frac{\text{if any other rule is not applicable} \quad \Gamma \vdash N \quad \Gamma \vdash M}{\Gamma \models N \stackrel{a}{\simeq} M \models (\hat{\alpha}_{\{N, M\}}^-, \hat{\alpha}_{\{N, M\}}^-, (\hat{\alpha}_{\{N, M\}}^- : \approx N), (\hat{\alpha}_{\{N, M\}}^- : \approx M))} \text{ AUAU} \\
\\
\boxed{N \simeq_1^D M} \quad \text{Negative multi-quantified type equivalence} \\
\\
\frac{}{\alpha^- \simeq_1^D \alpha^-} \text{ E1NVAR} \\
\\
\frac{P \simeq_1^D Q}{\uparrow P \simeq_1^D \uparrow Q} \text{ E1SHIFTU} \\
\\
\frac{P \simeq_1^D Q \quad N \simeq_1^D M}{P \rightarrow N \simeq_1^D Q \rightarrow M} \text{ E1ARROW} \\
\\
\frac{\vec{\alpha}^+ \cap \mathbf{fv} M = \emptyset \quad \mu : (\vec{\beta}^+ \cap \mathbf{fv} M) \leftrightarrow (\vec{\alpha}^+ \cap \mathbf{fv} N) \quad N \simeq_1^D [\mu]M}{\forall \alpha^+. N \simeq_1^D \forall \beta^+. M} \text{ E1FORALL}
\end{array}$$

$\boxed{P \simeq_1^D Q}$ Positive multi-quantified type equivalence

$$\begin{array}{c}
\frac{}{\alpha^+ \simeq_1^D \alpha^+} \text{E1PVAR} \\
\frac{N \simeq_1^D M}{\downarrow N \simeq_1^D \downarrow M} \text{E1SHIFTD} \\
\frac{\overrightarrow{\alpha^-} \cap \mathbf{fv} Q = \emptyset \quad \mu : (\overrightarrow{\beta^-} \cap \mathbf{fv} Q) \leftrightarrow (\overrightarrow{\alpha^-} \cap \mathbf{fv} P) \quad P \simeq_1^D [\mu]Q}{\exists \overrightarrow{\alpha^-}. P \simeq_1^D \exists \overrightarrow{\beta^-}. Q} \text{E1EXISTS}
\end{array}$$

$\boxed{P \simeq Q}$
 $\boxed{\Gamma \vdash N \simeq_1^{\leq} M}$ Negative equivalence on MQ types

$$\frac{\Gamma \vdash N \leq_1 M \quad \Gamma \vdash M \leq_1 N}{\Gamma \vdash N \simeq_1^{\leq} M} \text{D1NDEF}$$

$\boxed{\Gamma \vdash P \simeq_1^{\leq} Q}$ Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash Q \geq_1 P}{\Gamma \vdash P \simeq_1^{\leq} Q} \text{D1PDEF}$$

$\boxed{\Gamma \vdash N \leq_1 M}$ Negative subtyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^- \leq_1 \alpha^-} \text{D1NVAR} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \uparrow P \leq_1 \uparrow Q} \text{D1SHIFTU} \\
\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash N \leq_1 M}{\Gamma \vdash P \rightarrow N \leq_1 Q \rightarrow M} \text{D1ARROW} \\
\frac{\mathbf{fv} N \cap \overrightarrow{\beta^+} = \emptyset \quad \Gamma, \overrightarrow{\beta^+} \vdash P_i \quad \Gamma, \overrightarrow{\beta^+} \vdash [\overrightarrow{P}/\overrightarrow{\alpha^+}]N \leq_1 M}{\Gamma \vdash \forall \overrightarrow{\alpha^+}. N \leq_1 \forall \overrightarrow{\beta^+}. M} \text{D1FORALL}
\end{array}$$

$\boxed{\Gamma \vdash P \geq_1 Q}$ Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^+ \geq_1 \alpha^+} \text{D1PVAR} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \downarrow N \geq_1 \downarrow M} \text{D1SHIFTD} \\
\frac{\mathbf{fv} P \cap \overrightarrow{\beta^-} = \emptyset \quad \Gamma, \overrightarrow{\beta^-} \vdash N_i \quad \Gamma, \overrightarrow{\beta^-} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^-}]P \geq_1 Q}{\Gamma \vdash \exists \overrightarrow{\alpha^-}. P \geq_1 \exists \overrightarrow{\beta^-}. Q} \text{D1EXISTS}
\end{array}$$

$\boxed{\Gamma_2 \vdash \sigma_1 \simeq_1^{\leq} \sigma_2 : \Gamma_1}$ Equivalence of substitutions
 $\boxed{\Gamma \vdash N \simeq_0^{\leq} M}$ Negative equivalence

$$\frac{\Gamma \vdash N \leq_0 M \quad \Gamma \vdash M \leq_0 N}{\Gamma \vdash N \simeq_0^{\leq} M} \text{D0NDEF}$$

$\boxed{\Gamma \vdash P \simeq_0^{\leq} Q}$ Positive equivalence

$$\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash Q \geq_0 P}{\Gamma \vdash P \simeq_0^{\leq} Q} \text{D0PDEF}$$

$\boxed{\Gamma \vdash N \leq_0 M}$ Negative subtyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^- \leq_0 \alpha^-} \text{D0NVAR} \\
\frac{\Gamma \vdash P \simeq_0^{\leq} Q}{\Gamma \vdash \uparrow P \leq_0 \uparrow Q} \text{D0SHIFTU} \\
\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leq_0 M \quad M \neq \forall \beta^+. M'}{\Gamma \vdash \forall \alpha^+. N \leq_0 M} \text{D0FORALLL} \\
\frac{\Gamma, \alpha^+ \vdash N \leq_0 M}{\Gamma \vdash N \leq_0 \forall \alpha^+. M} \text{D0FORALLR} \\
\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash N \leq_0 M}{\Gamma \vdash P \rightarrow N \leq_0 Q \rightarrow M} \text{D0ARROW}
\end{array}$$

$\boxed{\Gamma \vdash P \geq_0 Q}$ Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^+ \geq_0 \alpha^+} \text{D0PVAR} \\
\frac{\Gamma \vdash N \simeq_0^{\leq} M}{\Gamma \vdash \downarrow N \geq_0 \downarrow M} \text{D0SHIFTD} \\
\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^-]P \geq_0 Q \quad Q \neq \exists \alpha^-. Q'}{\Gamma \vdash \exists \alpha^-. P \geq_0 Q} \text{D0EXISTSL} \\
\frac{\Gamma, \alpha^- \vdash P \geq_0 Q}{\Gamma \vdash P \geq_0 \exists \alpha^-. Q} \text{D0EXISTSR}
\end{array}$$

$\boxed{\Gamma; \Phi \vdash v : P}$ Positive type inference

$$\begin{array}{c}
\frac{v : P \in \Phi}{\Gamma; \Phi \vdash v : P} \text{DTVAR} \\
\frac{\Gamma; \Phi \vdash c : N}{\Gamma; \Phi \vdash \{c\} : \downarrow N} \text{DTTHUNK} \\
\frac{\Gamma; \Phi \vdash v : P \quad \Gamma \vdash Q \geq_1 P}{\Gamma; \Phi \vdash (v : Q) : Q} \text{DTANNOTP}
\end{array}$$

$\boxed{\Gamma; \Phi \vdash c : N}$ Negative type inference

$$\begin{array}{c}
\frac{\Gamma; \Phi, x : P \vdash c : N}{\Gamma; \Phi \vdash \lambda x : P. c : P \rightarrow N} \text{DTTLAM} \\
\frac{\Gamma, \alpha^+; \Phi \vdash c : N}{\Gamma; \Phi \vdash \Lambda \alpha^+. c : \forall \alpha^+. N} \text{DTTLAM} \\
\frac{\Gamma; \Phi \vdash v : P}{\Gamma; \Phi \vdash \mathbf{return} v : \uparrow P} \text{DTRETURN} \\
\frac{\Gamma; \Phi \vdash v : \downarrow M \quad \Gamma; \Phi \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \quad \Gamma \vdash \uparrow Q \leq_1 \uparrow P \quad \Gamma; \Phi, x : P \vdash c : N}{\Gamma; \Phi \vdash \mathbf{let} x : P = v(\vec{v}); c : N} \text{DTLETANN} \\
\frac{\Gamma; \Phi \vdash v : \downarrow M \quad \Gamma; \Phi \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \text{ uniquely} \quad \Gamma; \Phi, x : Q \vdash c : N}{\Gamma; \Phi \vdash \mathbf{let} x = v(\vec{v}); c : N} \text{DTLET} \\
\frac{\Gamma, \alpha^-; \Phi \vdash v : \exists \alpha^-. P \quad \Gamma, \alpha^-; \Phi, x : P \vdash c : N \quad \Gamma \vdash N}{\Gamma; \Phi \vdash \mathbf{let}^{\exists}(\alpha^-, x) = v; c : N} \text{DTUNPACK}
\end{array}$$

$$\frac{\Gamma; \Phi \vdash c : N \quad \Gamma \vdash N \leqslant_1 M}{\Gamma; \Phi \vdash (c : M) : M} \quad \text{DT}_{\text{ANNOTN}}$$

$$\boxed{\Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M} \quad \text{Spin Application type inference}$$

$$\frac{N \neq \forall \alpha^+ . M}{\Gamma; \Phi \vdash N \bullet \cdot \Rightarrow N} \quad \text{DT}_{\text{EMTPTY}}$$

$$\frac{\Gamma; \Phi \vdash v : P \quad \Gamma \vdash Q \geqslant_1 P \quad \Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M}{\Gamma; \Phi \vdash Q \rightarrow N \bullet v, \vec{v} \Rightarrow M} \quad \text{DT}_{\text{ARROW}}$$

$$\frac{\Gamma \vdash \vec{P} \quad \Gamma; \Phi \vdash [\vec{P}/\alpha^+] N \bullet \vec{v} \Rightarrow M}{\Gamma; \Phi \vdash \forall \alpha^+ . N \bullet \vec{v} \Rightarrow M} \quad \text{DT}_{\text{FORALL}}$$

$$\boxed{N = M} \quad \text{Negative type equality (alpha-equivalence)}$$

$$\boxed{P = Q} \quad \text{Positive type equality (alphha-equivalence)}$$

$$\boxed{P = Q}$$

$$\boxed{P_1 \vee P_2}$$

$$\boxed{\text{ord vars in } P}$$

$$\boxed{\text{ord vars in } N}$$

$$\boxed{\text{ord vars in } P}$$

$$\boxed{\text{ord vars in } N}$$

$$\boxed{\text{nf } (N')}$$

$$\boxed{\text{nf } (P')}$$

$$\boxed{\text{nf } (N')}$$

$$\boxed{\text{nf } (P')}$$

$$\boxed{\text{nf } (\vec{N}')}$$

$$\boxed{\text{nf } (\vec{P}')}$$

$\boxed{\mathbf{nf}(\sigma')}$ $\boxed{\mathbf{nf}(\mu')}$ $\boxed{\mathbf{nf}(\hat{\sigma}')}$ $\boxed{\sigma'|_{vars}}$ $\boxed{\hat{\sigma}'|_{vars}}$ $\boxed{\hat{\tau}'|_{vars}}$ $\boxed{\Xi'|_{vars}}$ $\boxed{e_1 \ \& \ e_2}$ $\boxed{\hat{\sigma}_1 \ \& \ \hat{\sigma}_2}$ $\boxed{\hat{\tau}_1 \ \& \ \hat{\tau}_2}$ $\boxed{\mathbf{dom}(\hat{\sigma})}$ $\boxed{\mathbf{dom}(\hat{\tau})}$ $\boxed{\mathbf{dom}(\Theta)}$ $\boxed{\Gamma \models P_1 \vee P_2 = Q}$ Least Upper Bound (Least Common Supertype) $\overline{\Gamma \models \alpha^+ \vee \alpha^+ = \alpha^+} \quad \text{LUBVAR}$

$$\frac{\Gamma, \cdot \models \mathbf{nf}(\downarrow N) \stackrel{a}{\simeq} \mathbf{nf}(\downarrow M) \Rightarrow (\Xi, P, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N \vee \downarrow M = \exists \alpha^-. [\alpha^- / \Xi] P} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \vec{\alpha}^-, \vec{\beta}^- \models P_1 \vee P_2 = Q}{\Gamma \models \exists \alpha^-. P_1 \vee \exists \beta^-. P_2 = Q} \quad \text{LUBEXISTS}$$

$$\boxed{\text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q}$$

$$\frac{\begin{array}{l} \Gamma = \Delta, \vec{\alpha}^\pm \quad \vec{\beta}^\pm \text{ is fresh} \quad \vec{\gamma}^\pm \text{ is fresh} \\ \Delta, \vec{\beta}^\pm, \vec{\gamma}^\pm \models [\vec{\beta}^\pm / \vec{\alpha}^\pm] P \vee [\vec{\gamma}^\pm / \vec{\alpha}^\pm] P = Q \end{array}}{\text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q} \quad \text{LUBUPGRADE}$$

$$\boxed{\mathbf{nf}(N) = M}$$

$$\frac{}{\mathbf{nf}(\alpha^-) = \alpha^-} \quad \text{NRMNVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\mathbf{nf}(\uparrow P) = \uparrow Q} \quad \text{NRMSHIFTU}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\mathbf{nf}(P \rightarrow N) = Q \rightarrow M} \quad \text{NRMARROW}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\mathbf{nf}(\vec{\forall} \alpha^+. N) = \vec{\forall} \alpha^{+'}. N'} \quad \text{NRMFORALL}$$

$$\boxed{\mathbf{nf}(P) = Q}$$

$$\frac{}{\mathbf{nf}(\alpha^+) = \alpha^+} \quad \text{NRMPVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\mathbf{nf}(\downarrow N) = \downarrow M} \quad \text{NRMSHIFTD}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\mathbf{nf}(\exists \vec{\alpha}^-. P) = \exists \vec{\alpha}^{-'}. P'} \quad \text{NRME EXISTS}$$

$$\boxed{\mathbf{nf}(N) = M}$$

$$\frac{}{\mathbf{nf}(\hat{\alpha}^-) = \hat{\alpha}^-} \quad \text{NRMNUVAR}$$

$$\boxed{\mathbf{nf}(P) = Q}$$

$$\frac{}{\mathbf{nf}(\hat{\alpha}^+) = \hat{\alpha}^+} \quad \text{NRMPUVAR}$$

$$\boxed{\text{ord vars in } N = \vec{\alpha}}$$

$$\frac{\alpha^- \in \text{vars}}{\text{ord vars in } \alpha^- = \alpha^-} \quad \text{ONVARIN}$$

$$\frac{\alpha^- \notin \text{vars}}{\text{ord vars in } \alpha^- = \cdot} \quad \text{ONVARININ}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}}{\text{ord vars in } \uparrow P = \vec{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}_1 \quad \text{ord vars in } N = \vec{\alpha}_2}{\text{ord vars in } P \rightarrow N = \vec{\alpha}_1, (\vec{\alpha}_2 \setminus \vec{\alpha}_1)} \quad \text{OARROW}$$

$$\frac{vars \cap \vec{\alpha}^+ = \emptyset \quad \mathbf{ord} \, vars \mathbf{in} \, N = \vec{\alpha}}{\mathbf{ord} \, vars \mathbf{in} \, \forall \vec{\alpha}^+. N = \vec{\alpha}} \quad \text{OForALL}$$

$$\boxed{\mathbf{ord} \, vars \mathbf{in} \, P = \vec{\alpha}}$$

$$\frac{\alpha^+ \in vars}{\mathbf{ord} \, vars \mathbf{in} \, \alpha^+ = \alpha^+} \quad \text{OPVarIN}$$

$$\frac{\alpha^+ \notin vars}{\mathbf{ord} \, vars \mathbf{in} \, \alpha^+ = .} \quad \text{OPVarNIN}$$

$$\frac{\mathbf{ord} \, vars \mathbf{in} \, N = \vec{\alpha}}{\mathbf{ord} \, vars \mathbf{in} \, \downarrow N = \vec{\alpha}} \quad \text{OShiftD}$$

$$\frac{vars \cap \vec{\alpha}^- = \emptyset \quad \mathbf{ord} \, vars \mathbf{in} \, P = \vec{\alpha}}{\mathbf{ord} \, vars \mathbf{in} \, \exists \vec{\alpha}^-. P = \vec{\alpha}} \quad \text{OExists}$$

$$\boxed{\mathbf{ord} \, vars \mathbf{in} \, N = \vec{\alpha}}$$

$$\frac{}{\mathbf{ord} \, vars \mathbf{in} \, \hat{\alpha}^- = .} \quad \text{ONUVar}$$

$$\boxed{\mathbf{ord} \, vars \mathbf{in} \, P = \vec{\alpha}}$$

$$\frac{}{\mathbf{ord} \, vars \mathbf{in} \, \hat{\alpha}^+ = .} \quad \text{OPUVar}$$

$$\boxed{\Gamma \vdash e_1 \& e_2 = e_3} \quad \text{Unification Solution Entry Merge}$$

$$\frac{\Gamma \vdash P_1 \vee P_2 = Q}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \& (\hat{\alpha}^+ : \geq P_2) = (\hat{\alpha}^+ : \geq Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vdash P \geq Q \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \geq Q) = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vdash Q \geq P \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \geq P) \& (\hat{\alpha}^+ : \approx Q) = (\hat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\frac{\ll\text{multiple parses}\gg}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \approx P') = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEPEQEQ}$$

$$\frac{\ll\text{multiple parses}\gg}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \& (\hat{\alpha}^- : \approx N') = (\hat{\alpha}^- : \approx N)} \quad \text{SMENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3} \quad \text{Merge unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \Rightarrow e_2} \quad \text{Weakening of unification solution entries}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPESUPSUP}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUP}$$

$$\frac{\ll\text{multiple parses}\gg}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2)} \quad \text{SIMPEPEQEQ}$$

$$\frac{\ll\text{multiple parses}\gg}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \Rightarrow (\hat{\alpha}^- : \approx N_2)} \quad \text{SIMPENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2} \quad \text{Weakening of unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \simeq e_2}$$

$$\begin{array}{c}
\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \simeq (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUPSUP} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \simeq (\hat{\alpha}^+ : \approx P_2)} \quad \text{SIMPEEQPEQEQ} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \simeq (\hat{\alpha}^- : \approx N_2)} \quad \text{SIMPEEQNEQEQ}
\end{array}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2} \\
\boxed{\Gamma; \Theta \models N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}} \quad \text{Negative unification}$$

$$\begin{array}{c}
\frac{}{\Gamma; \Theta \models \alpha^- \stackrel{u}{\simeq} \alpha^- \Rightarrow \cdot} \quad \text{UNVAR} \\
\frac{\Gamma; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \uparrow P \stackrel{u}{\simeq} \uparrow Q \Rightarrow \hat{\sigma}} \quad \text{USHIFTU} \\
\frac{\Gamma; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \Theta \models N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}_2}{\Gamma; \Theta \models P \rightarrow N \stackrel{u}{\simeq} Q \rightarrow M \Rightarrow \hat{\sigma}_1 \ \& \ \hat{\sigma}_2} \quad \text{UARROW} \\
\frac{\Gamma, \overrightarrow{\alpha^+}; \Theta \models N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \forall \alpha^+. N \stackrel{u}{\simeq} \forall \alpha^+. M \Rightarrow \hat{\sigma}} \quad \text{UFORALL} \\
\frac{\hat{\alpha}^- \{\Delta\} \in \Theta \quad \Delta \vdash N}{\Gamma; \Theta \models \hat{\alpha}^- \stackrel{u}{\simeq} N \Rightarrow (\hat{\alpha}^- : \approx N)} \quad \text{UNUVAR}
\end{array}$$

$$\boxed{\Gamma; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}} \quad \text{Positive unification}$$

$$\begin{array}{c}
\frac{}{\Gamma; \Theta \models \alpha^+ \stackrel{u}{\simeq} \alpha^+ \Rightarrow \cdot} \quad \text{UPVAR} \\
\frac{\Gamma; \Theta \models N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \downarrow N \stackrel{u}{\simeq} \downarrow M \Rightarrow \hat{\sigma}} \quad \text{USHIFTD} \\
\frac{\Gamma, \overrightarrow{\alpha^-}; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \exists \alpha^-. P \stackrel{u}{\simeq} \exists \alpha^-. Q \Rightarrow \hat{\sigma}} \quad \text{UEXISTS} \\
\frac{\hat{\alpha}^+ \{\Delta\} \in \Theta \quad \Delta \vdash P}{\Gamma; \Theta \models \hat{\alpha}^+ \stackrel{u}{\simeq} P \Rightarrow (\hat{\alpha}^+ : \approx P)} \quad \text{UPUVAR}
\end{array}$$

$$\begin{array}{ll}
\boxed{\Gamma \vdash N} & \text{Negative type well-formedness} \\
\boxed{\Gamma \vdash P} & \text{Positive type well-formedness} \\
\boxed{\Gamma \vdash N} & \text{Negative type well-formedness} \\
\boxed{\Gamma \vdash P} & \text{Positive type well-formedness} \\
\boxed{\Gamma \vdash \vec{N}} & \text{Negative type list well-formedness} \\
\boxed{\Gamma \vdash \vec{P}} & \text{Positive type list well-formedness} \\
\boxed{\Gamma; \Theta \vdash N} & \text{Negative unification type well-formedness} \\
\boxed{\Gamma; \Theta \vdash P} & \text{Positive unification type well-formedness} \\
\boxed{\Gamma; \Xi \vdash N} & \text{Negative anti-unification type well-formedness} \\
\boxed{\Gamma; \Xi \vdash P} & \text{Positive anti-unification type well-formedness} \\
\boxed{\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1} & \text{Antiunification substitution well-formedness} \\
\boxed{\hat{\sigma} : \Theta} & \text{Unification substitution well-formedness} \\
\boxed{\Gamma \vdash^\supset \Theta} & \text{Unification context well-formedness}
\end{array}$$

$\Gamma_1 \vdash \sigma : \Gamma_2$	Substitution well-formedness
$\Gamma \vdash e$	Unification solution entry well-formedness

Definition rules:	98 good	16 bad
Definition rule clauses:	195 good	16 bad