$\begin{array}{ll} \alpha,\,\beta,\,\alpha,\,\beta,\,\gamma,\,\delta & \text{type variables} \\ n,\,m,\,i,\,j & \text{index variables} \end{array}$ 

```
\hat{\sigma} \backslash \hat{\alpha}^+
                                                          \widehat{\sigma} \setminus \widehat{\alpha}^-
\widehat{\sigma}_1 \cup \widehat{\sigma}_2
                                                                                                    concatenate
                                                           (\widehat{\sigma})
\widehat{\sigma}_1 \& \widehat{\sigma}_2
                                                                                     S
                                                                                     Μ
\hat{\tau}
                                                                                               anti-unification substitution
                                               ::=
                                                           \widehat{\alpha}^-:\approx N
                                                         \alpha : \approx 1
\overrightarrow{\alpha} / \overrightarrow{\alpha}^{-}
\overrightarrow{N} / \overrightarrow{\alpha}^{-}
\widehat{\tau}_{1} \cup \widehat{\tau}_{2}
\overline{\widehat{\tau}_{i}}^{i}
                                                                                                    concatenate
                                                                                     S
P, Q
                                                                                               positive types
                                               ::=
                                                           \alpha^+
                                                           \downarrow N
                                                           \exists \alpha^-.P
                                                           [\sigma]P
                                                                                     Μ
N, M
                                                                                               negative types
                                                           \alpha^{-}
                                                           \uparrow P
                                                           \forall \alpha^+.N
                                                           P \rightarrow N
                                                           [\sigma]N
                                                                                     Μ
                                                                                               positive variable list
                                                                                                    empty list
                                                                                                    a variable
                                                                                                    a variable
                                                                                                     concatenate lists
                                                                                               negative variables
                                                                                                    empty list
                                                                                                    a variable
                                                                                                     a variable
                                                                                                     concatenate lists
P, Q
                                                                                               multi-quantified positive types
                                                           \alpha^+
                                                                                                    P \neq \exists \dots
                                                           [\sigma]P
                                                                                     Μ
                                                                                     Μ
                                                           [\hat{\sigma}]P
                                                                                     Μ
```

```
[\mu]P
                                   Μ
                                   S
                      (P)
                     \mathbf{nf}(P')
N, M
                                        multi-quantified negative types
                     \alpha^{-}
                      \uparrow P
                      P \to N
                     \forall \overrightarrow{\alpha^+}.N
                                            N \neq \forall \dots
                     [\sigma]N
                                   Μ
                      [\mu]N
                                   Μ
                      [\hat{\sigma}]N
                                   Μ
                     (N)
                                   S
                      \mathbf{nf}(N')
\vec{P}, \vec{Q}
                                        list of positive types
                                            empty list
                                            a singel type
                                            concatenate lists
                                   Μ
\vec{N}, \vec{M}
                                         list of negative types
                                            empty list
                                            a singel type
                                            concatenate lists
                                   Μ
\Delta, \Gamma
                                         declarative type context
                                            empty context
                                            list of variables
                                            list of variables
                     \overline{\Gamma_i}^i
                                            concatenate contexts
                                   S
Θ
                                         unification type variable context
                                            empty context
                                            list of variables
                                            list of variables
                     \overline{\Theta_i}^{i}
                                            concatenate contexts
                                   S
Ξ
                                         anti-unification type variable context
                                            empty context
                                            list of variables
                                            list of variables
                                            concatenate contexts
                                   S
```

```
\vec{\alpha}, \vec{\beta}
                                                        ordered positive or negative variables
                                                            empty list
                                                            list of variables
                                                            list of variables
                          \overrightarrow{\alpha}_1 \backslash vars
                                                            setminus
                                                            context
                          vars
                          \overline{\overrightarrow{\alpha}_i}^i
                                                            concatenate contexts
                                                  S
                          (\vec{\alpha})
                                                            parenthesis
                          [\mu]\vec{\alpha}
                                                            apply moving to list
                                                  Μ
                          ord vars in P
                                                  Μ
                          ord vars in N
                                                  Μ
                          ord vars in P
                                                  Μ
                          \operatorname{\mathbf{ord}} \operatorname{\mathbf{\mathit{vars}}} \operatorname{\mathbf{in}} N
                                                  Μ
                                                         set of variables
vars
                          Ø
                                                            empty set
                         \mathbf{fv} P
                                                            free variables
                         \mathbf{fv} N
                                                            free variables
                         fv imP
                                                            free variables
                                                            free variables
                         fv imN
                          vars_1 \cap vars_2
                                                            set intersection
                                                            set union
                          vars_1 \cup vars_2
                                                            set complement
                          vars_1 \backslash vars_2
                         mv imP
                                                            movable variables
                          mv imN
                                                            movable variables
                          \mathbf{u}\mathbf{v} N
                                                            unification variables
                                                            unification variables
                          \mathbf{u}\mathbf{v} P
                                                            free variables
                          \mathbf{fv} N
                          \mathbf{fv} P
                                                            free variables
                                                            parenthesis
                                                  S
                          (vars)
                          \{\vec{\alpha}\}\
                                                            ordered list of variables
                                                  Μ
                                                            apply moving to varset
                          [\mu]vars
                                                            empty moving
                         pma1 \mapsto pma2
                                                            Positive unit substitution
                         nma1 \mapsto nma2
                                                            Positive unit substitution
                                                  Μ
                                                            Set-like union of movings
                         \mu_1 \cup \mu_2
                         \overline{\mu_i}^{\ i}
                                                            concatenate movings
                                                  Μ
                          \mu|_{vars}
                                                            restriction on a set
                                                  Μ
                                                            inversion
                         \mathbf{nf}(\mu')
                                                  Μ
                                                         positive unification variable
                         \hat{\alpha}^+
                         \hat{\alpha}^+\{\Delta\}
                                                        negative unification variable
```

$$\overrightarrow{\alpha^{+}}, \overrightarrow{\beta^{+}} \qquad ::= \qquad \qquad \text{positive unification variable list} \\ \overrightarrow{\alpha^{+}}, \overrightarrow{\beta^{+}} \qquad ::= \qquad \qquad \text{empty list} \\ \overrightarrow{\alpha^{+}} + \overrightarrow{\alpha^{$$

```
\in
                                ∉
                                Ø
                                \equiv_n
                                \Downarrow
                                :≽
                                :≃
formula
                                judgement
                                formula_1 .. formula_n
                                \mu : vars_1 \leftrightarrow vars_2
                                \mu is bijective
                                \hat{\sigma} is functional
                                \hat{\sigma}_1 \in \hat{\sigma}_2
                                vars_1 \subseteq vars_2
                                vars_1 = vars_2
                                vars is fresh
                                \alpha^- \not\in \mathit{vars}
                                \alpha^+ \notin vars
                                \alpha^- \in \mathit{vars}
                                \alpha^+ \in vars
                                \widehat{\alpha}^- \in \Theta
                                \widehat{\alpha}^+ \in \Theta
                                if any other rule is not applicable
                                N \neq M
                                P \neq Q
                      ::=
                                \Gamma; \Theta \models N \leqslant M \dashv \hat{\sigma}
                                                                                                      Negative subtyping
                                \Gamma; \Theta \vDash P \geqslant Q \dashv \hat{\sigma}
                                                                                                       Positive supertyping
```

AU::=

 $\boldsymbol{A}$ 

$$| \quad \Gamma \vdash P_1 \stackrel{\circ}{\cong} P_2 \dashv (\Xi, Q, \widehat{\tau}_1, \widehat{\tau}_2) \\ | \quad \Gamma \vdash N_1 \stackrel{\circ}{\cong} N_2 \dashv (\Xi, M, \widehat{\tau}_1, \widehat{\tau}_2) \\ | \quad \Gamma \vdash N_1 \stackrel{\circ}{\cong} N_2 \dashv (\Xi, M, \widehat{\tau}_1, \widehat{\tau}_2) \\ | \quad P \Rightarrow P Q \\ | \quad P \Rightarrow Q \\ | \quad P \vdash N \Rightarrow 1 M \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash N \Rightarrow 1 M \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q \\ | \quad P \vdash P \Rightarrow 1 Q$$

 $\mathbf{nf}(N) = M$ 

 $n \\ n$ 

NauSolterminalsformula

# $\Gamma; \Theta \models N \leqslant M \dashv \hat{\sigma}$ Negative subtyping

$$\frac{\Gamma; \Theta \vDash \alpha^{-} \leqslant \alpha^{-} \dashv \cdot}{\Gamma; \Theta \vDash \mathsf{nf} (P) \stackrel{u}{\simeq} \mathsf{nf} (Q) \dashv \widehat{\sigma}} \qquad \mathsf{ASHIFTU}$$

$$\frac{\Theta \vDash \mathsf{nf} (P) \stackrel{u}{\simeq} \mathsf{nf} (Q) \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \uparrow P \leqslant \uparrow Q \dashv \widehat{\sigma}} \qquad \mathsf{ASHIFTU}$$

$$\frac{\Gamma; \Theta \vDash P \geqslant Q \dashv \widehat{\sigma}_{1} \quad \Gamma; \Theta \vDash N \leqslant M \dashv \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash P \rightarrow N \leqslant Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \qquad \mathsf{AARROW}$$

$$\frac{\Gamma; \Theta \vDash P \rightarrow N \leqslant Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash \nabla \widehat{\sigma}^{+} \{\Gamma, \widehat{\beta}^{+}\}} \vDash [\widehat{\alpha}^{+}/\widehat{\alpha}^{+}] N \leqslant M \dashv \widehat{\sigma}$$

$$\Gamma; \Theta \vDash \forall \widehat{\alpha}^{+}. N \leqslant \forall \widehat{\beta}^{+}. M \dashv \widehat{\sigma} \backslash \widehat{\alpha}^{+}$$

$$\mathsf{AFORALL}$$

 $\Gamma; \Theta \models P \geqslant Q \dashv \widehat{\sigma}$  Positive supertyping

$$\frac{\Gamma; \Theta \vDash \alpha^{+} \geqslant \alpha^{+} \dashv \cdot}{\Gamma; \Theta \vDash \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}}$$

$$\frac{\Theta \vDash \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \bot N \geqslant \bot M \dashv \widehat{\sigma}} \quad \text{ASHIFTD}$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}}; \Theta, \overrightarrow{\widehat{\alpha}^{-}} \{\Gamma, \overrightarrow{\beta^{-}}\} \vDash [\overrightarrow{\widehat{\alpha}^{-}}/\overrightarrow{\alpha^{-}}]P \geqslant Q \Rightarrow \widehat{\sigma}}{\Gamma; \Theta \vDash \overrightarrow{\beta\alpha^{-}}.P \geqslant \overrightarrow{\beta\beta^{-}}.Q \Rightarrow \widehat{\sigma}} \quad \text{AEXISTS}$$

$$\frac{\mathbf{upgrade} \Gamma \vdash \mathbf{nf} (P) \mathbf{to} \Delta = Q}{\Gamma: \Theta \vDash \widehat{\alpha}^{+} \{\Delta\} \geqslant P \Rightarrow (\Delta \vdash \widehat{\alpha}^{+} : \geqslant Q)} \quad \text{APUVAR}$$

 $\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ 

$$\frac{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} \alpha^{+} \Rightarrow (\cdot, \alpha^{+}, \cdot, \cdot)}{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \Rightarrow (\Xi, M, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \qquad \text{AUPShift}$$

$$\frac{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \Rightarrow (\Xi, M, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash \downarrow N_{1} \stackrel{a}{\simeq} \downarrow N_{2} \Rightarrow (\Xi, \downarrow M, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \qquad \text{AUPShift}$$

$$\frac{\{\overrightarrow{\alpha^{-}}\} \cap \{\Gamma\} = \varnothing \quad \Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \Rightarrow (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash \overrightarrow{\beta\alpha^{-}} \cdot P_{1} \stackrel{a}{\simeq} \overrightarrow{\beta\alpha^{-}} \cdot P_{2} \Rightarrow (\Xi, \overrightarrow{\beta\alpha^{-}}, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \qquad \text{AUPEXISTS}$$

 $\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 \rightrightarrows (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$ 

$$\frac{\Gamma \vDash \alpha^{-} \stackrel{a}{\simeq} \alpha^{-} \dashv (\Xi, \alpha^{-}, \cdot, \cdot)}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \quad \text{AUNSHIFT}$$

$$\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} \uparrow P_{2} \dashv (\Xi, \uparrow Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \quad \text{AUNSHIFT}$$

$$\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi_{1}, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2}) \quad \Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi_{2}, M, \widehat{\tau}'_{1}, \widehat{\tau}'_{2})}{\Gamma \vDash P_{1} \rightarrow N_{1} \stackrel{a}{\simeq} P_{2} \rightarrow N_{2} \dashv (\Xi_{1} \cup \Xi_{2}, Q \rightarrow M, \widehat{\tau}_{1} \cup \widehat{\tau}'_{1}, \widehat{\tau}_{2} \cup \widehat{\tau}'_{2})} \quad \text{AUNARROW}$$

$$\frac{\text{if any other rule is not applicable} \quad \Gamma \vDash N \quad \Gamma \vDash M}{\Gamma \vDash N \stackrel{a}{\simeq} M \dashv (\widehat{\alpha}^{-}_{\{N,M\}}, \widehat{\alpha}^{-}_{\{N,M\}}, (\widehat{\alpha}^{-}_{\{N,M\}} : \approx N), (\widehat{\alpha}^{-}_{\{N,M\}} : \approx M))} \quad \text{AUNAU}$$

 $N \simeq_1^D M$  Negative multi-quantified type equivalence

$$\frac{\alpha^{-} \simeq_{1}^{D} \alpha^{-}}{\alpha^{-}} \quad \text{E1NVAR}$$

$$\frac{P \simeq_{1}^{D} Q}{\uparrow P \simeq_{1}^{D} \uparrow Q} \quad \text{E1SHIFTU}$$

$$\frac{P \simeq_{1}^{D} Q \quad N \simeq_{1}^{D} M}{P \to N \simeq_{1}^{D} Q \to M} \quad \text{E1ARROW}$$

$$\frac{\{\overrightarrow{\alpha^{+}}\} \cap \text{fv } M = \varnothing \quad \mu : (\{\overrightarrow{\beta^{+}}\} \cap \text{fv } M) \leftrightarrow (\{\overrightarrow{\alpha^{+}}\} \cap \text{fv } N) \quad N \simeq_{1}^{D} [\mu] M}{\forall \overrightarrow{\alpha^{+}} . N \simeq_{1}^{D} \forall \overrightarrow{\beta^{+}} . M}$$

$$\text{E1FORALL}$$

 $P \simeq_1^D Q$  Positive multi-quantified type equivalence

$$\frac{\alpha^{+} \simeq_{1}^{D} \alpha^{+}}{\alpha^{+} \simeq_{1}^{D} M} \quad \text{E1ShiftD}$$

$$\frac{N \simeq_{1}^{D} M}{\downarrow N \simeq_{1}^{D} \downarrow M} \quad \text{E1ShiftD}$$

$$\overrightarrow{\{\alpha^{-}\}} \cap \text{fv } Q = \varnothing \quad \mu : (\overrightarrow{\{\beta^{-}\}} \cap \text{fv } Q) \leftrightarrow (\overrightarrow{\{\alpha^{-}\}} \cap \text{fv } P) \quad P \simeq_{1}^{D} [\mu]Q \quad \text{E1Exists}$$

$$\overrightarrow{\exists \alpha^{-}} \cdot P \simeq_{1}^{D} \overrightarrow{\exists \beta^{-}} \cdot Q$$

$$P \simeq Q$$

$$\Gamma \vdash N \sim^{\$} M$$

 $\begin{array}{|c|c|c|c|c|c|}\hline P \simeq Q \\ \hline \Gamma \vdash N \simeq_1^s M \\ \hline \end{array} \quad \text{Negative equivalence on MQ types}$ 

$$\frac{\Gamma \vdash N \leqslant_{1} M \quad \Gamma \vdash M \leqslant_{1} N}{\Gamma \vdash N \simeq_{1}^{\leqslant} M} \quad \text{D1NDEF}$$

 $\Gamma \vdash P \simeq_1^{\leqslant} Q$ Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geqslant_1 Q \quad \Gamma \vdash Q \geqslant_1 P}{\Gamma \vdash P \simeq_1^{\circ} Q} \quad \text{D1PDEF}$$

 $\Gamma \vdash N \leqslant_1 M$ Negative subtyping

 $\overline{|\Gamma \vdash P \geqslant_1 Q|}$  Positive supertyping

 $\begin{array}{|c|c|c|c|c|}\hline \Gamma_2 \vdash \sigma_1 \simeq_1^{\leftarrow} \sigma_2 : \Gamma_1 \\\hline \Gamma \vdash N \simeq_0^{\leftarrow} M \\\hline \end{array} \quad \begin{array}{|c|c|c|c|c|}\hline \text{Equivalence of substitutions}\\\hline \end{array}$ 

$$\frac{\Gamma \vdash N \leqslant_0 M \quad \Gamma \vdash M \leqslant_0 N}{\Gamma \vdash N \simeq_0^\circ M} \quad \text{D0NDef}$$

 $\Gamma \vdash P \simeq_0^{\leqslant} Q$ Positive equivalence

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash Q \geqslant_0 P}{\Gamma \vdash P \simeq_0^6 Q} \quad \text{D0PDEF}$$

 $\Gamma \vdash N \leqslant_0 M$ Negative subtyping

$$\frac{\Gamma \vdash \alpha^- \leqslant_0 \alpha^-}{\Gamma \vdash P \stackrel{\sim_0}{=} Q} \quad D0NVAR$$
 
$$\frac{\Gamma \vdash P \stackrel{\sim_0}{=} Q}{\Gamma \vdash P \leqslant_0 \uparrow Q} \quad D0SHIFTU$$
 
$$\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leqslant_0 M \quad M \neq \forall \beta^+.M'}{\Gamma \vdash \forall \alpha^+.N \leqslant_0 M} \quad D0FORALLL$$
 
$$\frac{\Gamma, \alpha^+ \vdash N \leqslant_0 M}{\Gamma \vdash N \leqslant_0 \forall \alpha^+.M} \quad D0FORALLR$$

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash N \leqslant_0 M}{\Gamma \vdash P \to N \leqslant_0 Q \to M} \quad \text{D0Arrow}$$

 $\Gamma \vdash P \geqslant_0 Q$  Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{0} \alpha^{+}}{\Gamma \vdash N \simeq_{0}^{\leq} M} \quad D0PVAR$$

$$\frac{\Gamma \vdash N \simeq_{0}^{\leq} M}{\Gamma \vdash \downarrow N \geqslant_{0} \downarrow M} \quad D0SHIFTD$$

$$\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^{-}]P \geqslant_{0} Q \quad Q \neq \exists \alpha^{-}.Q'}{\Gamma \vdash \exists \alpha^{-}.P \geqslant_{0} Q} \quad D0EXISTSL$$

$$\frac{\Gamma, \alpha^{-} \vdash P \geqslant_{0} Q}{\Gamma \vdash P \geqslant_{0} \exists \alpha^{-}.Q} \quad D0EXISTSR$$

N = M Negative type equality (alpha-equivalence) P = Q Positive type equivalence) P = Q ord vars in P

 $\mathbf{ord}\ vars\mathbf{in}\ N$ 

ord vars in P

 $\mathbf{ord}\ vars\mathbf{in}\ N$ 

 $\mathbf{nf}(N')$ 

 $\mathbf{nf}(P')$ 

 $\mathbf{nf}(N')$ 

 $\mathbf{nf}(P')$ 

 $|\mathbf{nf}(\vec{N}')|$ 

 $|\mathbf{nf}(\vec{P}')|$ 

 $\mathbf{nf}\left(\sigma'\right)$ 

 $\mathbf{nf}(\mu')$ 

 $|\sigma'|_{vars}$ 

 $e_1 \& e_2$ 

 $\widehat{\sigma}_1 \& \widehat{\sigma}_2$ 

 $\overline{|\Gamma \models P_1 \lor P_2 = Q|}$  Least Upper Bound (Least Common Supertype)

$$\frac{\Gamma \models \alpha^{+} \vee \alpha^{+} = \alpha^{+}}{\Gamma \models \alpha^{+} \vee \alpha^{+} = \alpha^{+}} \quad \text{LUBVAR}$$

$$\frac{\Gamma, \cdot \models \downarrow N \stackrel{a}{\simeq} \downarrow M \Rightarrow (\Xi, P, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \models \downarrow N \vee \downarrow M = \exists \alpha^{-}. [\alpha^{-}/\Xi] P} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \alpha^{-}, \beta^{-}}{\Gamma, \alpha^{-}, \beta^{-}} \models P_{1} \vee P_{2} = Q$$

$$\frac{\Gamma, \alpha^{-}, \beta^{-}}{\Gamma, \beta^{-}} \models P_{1} \vee P_{2} = Q$$

$$\Gamma \models \exists \alpha^{-}. P_{1} \vee \exists \beta^{-}. P_{2} = Q$$

$$\text{LUBEXISTS}$$

 $\frac{|\mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q|}{|\mathbf{nf}\,(N) = M|}$ 

 $\mathbf{nf}(P) = Q$ 

 $\mathbf{nf}(N) = M$ 

$$\overline{\mathbf{nf}(\hat{\alpha}^{-}) = \hat{\alpha}^{-}}$$
 NRMNUVAR

$$\mathbf{nf}(P) = Q$$

$$\overline{\mathbf{nf}(\widehat{\alpha}^{+}) = \widehat{\alpha}^{+}}$$
 NRMPUVAR

## $\mathbf{ord} \ vars \mathbf{in} \ N = \vec{\alpha}$

$$\frac{\alpha^{-} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \alpha^{-}} \quad \text{ONVARIN}$$

$$\frac{\alpha^{-} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \uparrow P = \overrightarrow{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}_1 \quad \mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}_2}{\mathbf{ord} \ vars \mathbf{in} \ P \to N = \overrightarrow{\alpha}_1, (\overrightarrow{\alpha}_2 \setminus \{\overrightarrow{\alpha}_1\})} \quad \text{OARROW}$$

$$\frac{vars \cap \{\overrightarrow{\alpha^+}\} = \varnothing \quad \mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}}{\mathbf{ord} \ vars \mathbf{in} \ \forall \overrightarrow{\alpha^+}. N = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

#### $\operatorname{ord} \operatorname{varsin} P = \overrightarrow{\alpha}$

$$\frac{\alpha^{+} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \alpha^{+}} \quad \operatorname{OPVarIn}$$

$$\frac{\alpha^{+} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \cdot} \quad \operatorname{OPVarNIn}$$

$$\frac{\operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \sqrt{N} = \overrightarrow{\alpha}} \quad \operatorname{OSHIFTD}$$

$$\frac{vars \cap \{\overrightarrow{\alpha^{-}}\} = \varnothing \quad \operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \exists \overrightarrow{\alpha^{-}}.P = \overrightarrow{\alpha}} \quad \operatorname{OEXISTS}$$

## $\operatorname{\mathbf{ord}} vars \operatorname{\mathbf{in}} N = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}} \operatorname{\mathit{vars}} \operatorname{\mathbf{in}} \widehat{\alpha}^- = \cdot}$$
 ONUVAR

## $\operatorname{\mathbf{ord}} vars \operatorname{\mathbf{in}} P = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}}\operatorname{\mathit{vars}}\operatorname{\mathbf{in}}\widehat{\alpha}^{+}=\cdot}$$
 OPUVAR

#### $e_1 \& e_2 = e_3$ Unification Solution Entry Merge

$$\frac{\Gamma \vDash P_1 \lor P_2 = Q}{(\Gamma \vdash \widehat{\alpha}^+ : \geqslant P_1) \& (\Gamma \vdash \widehat{\alpha}^+ : \geqslant P_2) = (\Gamma \vdash \widehat{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vDash P \geqslant Q \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \geqslant Q) = (\Gamma \vdash \widehat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vDash Q \geqslant P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \geqslant P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx Q) = (\Gamma \vdash \widehat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\frac{\Gamma; \cdot \vDash Q \geqslant P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx Q)} \quad \text{SMEPEQEQ}$$

$$\frac{1}{(\Gamma \vdash \widehat{\alpha}^{-} :\approx N) \& (\Gamma \vdash \widehat{\alpha}^{-} :\approx N) = (\Gamma \vdash \widehat{\alpha}^{-} :\approx N)} \quad \text{SMENEQEQ}$$

 $\widehat{\sigma}_1 \& \widehat{\sigma}_2 = \widehat{\sigma}_3$  Merge unification solutions

 $\Theta \models N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}$  Negative unification

$$\frac{\Theta \vDash \alpha^{-\frac{u}{\simeq}} \alpha^{-} \dashv \cdot \text{UNVAR}}{\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma} \qquad \text{USHIFTU}}$$

$$\frac{\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}}{\Theta \vDash \uparrow P \stackrel{u}{\simeq} \uparrow Q \dashv \widehat{\sigma}} \qquad \text{USHIFTU}$$

$$\frac{\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}_{1} \quad \Theta \vDash N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}_{2}}{\Theta \vDash P \rightarrow N \stackrel{u}{\simeq} Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \qquad \text{UARROW}$$

$$\frac{\Theta \vDash N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}}{\Theta \vDash \forall \alpha^{+}. N \stackrel{u}{\simeq} \forall \alpha^{+}. M \dashv \widehat{\sigma}} \qquad \text{UFORALL}$$

$$\frac{\widehat{\alpha}^{-} \{\Delta\} \in \Theta \quad \Delta \vdash N}{\Theta \vDash \widehat{\alpha}^{-} \stackrel{u}{\simeq} N \dashv (\Delta \vdash \widehat{\alpha}^{-} : \approx N)} \qquad \text{UNUVAR}$$

 $\Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \widehat{\sigma}$  Positive unification

$$\frac{\Theta \vDash \alpha^{+} \stackrel{u}{\simeq} \alpha^{+} \dashv \cdot}{\Theta \vDash N \stackrel{u}{\simeq} M \dashv \hat{\sigma}} \quad \text{USHIFTD}$$

$$\frac{\Theta \vDash N \stackrel{u}{\simeq} M \dashv \hat{\sigma}}{\Theta \vDash \downarrow N \stackrel{u}{\simeq} \downarrow M \dashv \hat{\sigma}} \quad \text{USHIFTD}$$

$$\frac{\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \hat{\sigma}}{\Theta \vDash \exists \alpha^{-}.P \stackrel{u}{\simeq} \exists \alpha^{-}.Q \dashv \hat{\sigma}} \quad \text{UEXISTS}$$

$$\frac{\hat{\alpha}^{+} \{\Delta\} \in \Theta \quad \Delta \vdash P}{\Theta \vDash \hat{\alpha}^{+} \stackrel{u}{\simeq} P \dashv (\Delta \vdash \hat{\alpha}^{+} : \approx P)} \quad \text{UPUVAR}$$

 $\Gamma \vdash N$  Negative type well-formedness

 $\overline{\Gamma \vdash P}$  Positive type well-formedness

 $\overline{\overline{\Gamma \vdash N}}$  Negative type well-formedness

 $\Gamma \vdash P$  Positive type well-formedness

 $\Gamma \vdash N$  Negative type list well-formedness

 $\Gamma \vdash \overrightarrow{P}$  Positive type list well-formedness

 $\Gamma;\Xi\vdash P$  Positive anti-unification type well-formedness

 $\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1$  Antiunification substitution well-formedness

 $|\Theta \vdash \hat{\sigma}|$  Unification substitution well-formedness

 $\Gamma \vdash \Theta$  Unification context well-formedness

 $\Gamma_1 \vdash \sigma : \Gamma_2$  Substitution well-formedness

Definition rules: 72 good 7 bad Definition rule clauses: 130 good 7 bad