$\begin{array}{ll} \alpha,\,\beta & \text{type variables} \\ n,\,m,\,i,\,j & \text{index variables} \end{array}$

```
positive variable
                                     \alpha^+
                                                                            negative variable
                                                                            substitution
                                      P/a+

\begin{array}{ccc}
N/a - \\
\overrightarrow{P}/\overrightarrow{\alpha^{+}} \\
\overrightarrow{N}/\overrightarrow{\alpha^{-}} \\
\overrightarrow{\alpha^{+}}/\alpha^{+}
\end{array}

                                     vars_1/vars_2
                                                                                 concatenate
                                                                            entry of a unification solution
e
                                     \widehat{\alpha}^+:\approx P
                                     \widehat{\alpha}^-:\approx N
                                      \widehat{\alpha}^+:\geqslant P
\hat{\sigma}
                                                                            unification solution (substitution)
                                                                                 concatenate
                                      (\hat{\sigma})
                                                                  S
                                     \hat{\sigma}_1 \& \hat{\sigma}_2
                                                                  Μ
P, Q
                                                                            positive types
                                     a+
                                      \downarrow N
                                      \exists \alpha^-.P
                                      [\sigma]P
                                                                  Μ
N, M
                                                                            negative types
                                     a-
                                     \uparrow P
                                     \forall \alpha^+.N
                                      [\sigma]N
                                                                  Μ
```

$\overrightarrow{\alpha^+}, \ \overrightarrow{\beta^+}$ $\rightarrow \rightarrow$::= 	$\overset{\cdot}{\underset{\alpha^{+}_{i}}{\overrightarrow{\alpha^{+}}}}^{i}$		positive variable list empty list a variable concatenate lists
$\overrightarrow{\alpha}^-, \overrightarrow{\beta}^-$::= 	\vdots $\frac{\alpha^{-}}{\alpha^{-}_{i}}^{i}$		negative variables empty list a variable concatenate lists
$P,\ Q$		α^{+} $\downarrow N$ $\exists \alpha^{-}.P$ $[\sigma]P$ $[\mu]P$ $P_{1} \vee P_{2}$ $\mathbf{nf}(P')$		multi-quantified positive types $P \neq \exists \dots$
$N,\ M$::=	α^{-} $\uparrow P$ $P \to N$ $\forall \alpha^{+}.N$ $[\sigma]N$ $[\mu]N$ $\mathbf{nf}(N')$	M M M	multi-quantified negative types $N \neq \forall \dots$ list of positive types empty list
\vec{P}	::= 	$.$ $\overrightarrow{\overrightarrow{P}}_{i}^{i}$		list of positive types empty list a singel type concatenate lists
$ec{N}$::= 	$\vdots \\ \frac{N}{\overrightarrow{N}_i}^i$		list of negative types empty list a singel type concatenate lists
Γ	::=	$\begin{matrix} vars \\ \overrightarrow{\alpha^+} \\ \overrightarrow{\alpha^-} \\ \overline{\Gamma_i}^i \\ (\Gamma) \end{matrix}$	S	declarative type context empty context list of variables list of variables concatenate contexts
$\vec{\alpha}, \vec{\beta}$::= 	$\overrightarrow{\alpha^+}$		ordered positive or negative variables empty list list of variables

```
list of variables
                    \overrightarrow{\alpha}_1 \backslash vars
                                                     setminus
                                                     concatenate contexts
                                            S
                     (\vec{\alpha})
                                                     parenthesis
                     [\mu]\vec{\alpha}
                                                     apply moving to list
                                           Μ
                    ord vars in P
                                            Μ
                    \mathbf{ord}\ vars\mathbf{in}\ \overline{N}
                                           Μ
                    ord vars in P
                                            М
                    \mathbf{ord}\ vars \mathbf{in}\ N
                                            Μ
                                                  set of variables
vars
             ::=
                    Ø
                                                     empty set
                    \mathbf{fv}\,P
                                                     free variables
                    \mathbf{fv}\,N
                                                     free variables
                    \mathbf{fv}\,P
                                                     free variables
                    \mathbf{fv}\,N
                                                     free variables
                                                     set intersection
                    vars_1 \cap vars_2
                                                     set union
                    \mathit{vars}_1 \cup \mathit{vars}_2
                    vars_1 \backslash vars_2
                                                     set complement
                    \mathbf{mv} P
                                                     movable variables
                    \mathbf{mv}\,N
                                                     movable variables
                    \mathbf{u}\mathbf{v} N
                                                     unification variables
                    \mathbf{uv} P
                                                     unification variables
                    \mathbf{fv} N
                                                     free variables
                    \mathbf{fv} P
                                                     free variables
                                           S
                                                     parenthesis
                    (vars)
                    Γ
                                                     context
                                                     ordered list of variables
                     \{\vec{\alpha}\}\
                     [\mu]vars
                                           Μ
                                                     apply moving to varset
\mu
            ::=
                                                     empty moving
                                                     Positive unit substitution
                                                     Positive unit substitution
                                                     Set-like union of movings
                    \mu_1 \cup \mu_2
                                            Μ
                    \overline{\mu_i}^{i}
                                                     concatenate movings
                                            Μ
                                                     restriction on a set
                    \mu|_{vars}
                                            Μ
                                                      inversion
                                                  cohort index
n
            ::=
                    0
                    n+1
\tilde{\alpha}^+
                                                  positive movable variable
                    \widetilde{\alpha}^{+n}
\tilde{\alpha}^-
            ::=
                                                  negative movable variable
                    \tilde{\alpha}^{-n}
```

$\overrightarrow{\widetilde{\alpha}^+}, \ \overrightarrow{\widetilde{\beta}^+}$::=	$ \overset{\cdot}{\underset{\alpha}{}} \stackrel{\cdot}{\underset{\alpha}{}} \stackrel{\cdot}{\underset{\alpha}{}{\underset{\alpha}{}} \stackrel{\cdot}{\underset{\alpha}{}} \stackrel{}{\underset{\alpha}{}} \stackrel{}{\underset{\alpha}{}}{\underset{\alpha}{}} \stackrel{}{\underset{\alpha}{}} \stackrel{}{\underset{\alpha}{}} \stackrel{}{\underset{\alpha}{$		positive movable variable list empty list a variable from a non-movable variable concatenate lists
$\overrightarrow{\widetilde{\alpha}^-}, \ \overrightarrow{\widetilde{\beta}^-}$::=	$ \widetilde{\alpha}^{-} \xrightarrow{\widetilde{\alpha}^{-n}} \overrightarrow{\alpha}^{i} $		negatiive movable variable list empty list a variable from a non-movable variable
$P,\ Q$::=	$ \overrightarrow{\alpha^{+}}_{i} $ $ \alpha^{+}_{i} $ $ \overrightarrow{\alpha^{+}}_{i} $ $ \downarrow N $ $ \exists \alpha^{-}.P $ $ [\sigma]P $ $ [\mu]P $	M M	multi-quantified positive types with movable variables
$N,\ M$::=	$\begin{array}{c} \alpha^{-} \\ \widetilde{\alpha}^{-} \\ \uparrow P \\ P N \\ \forall \alpha^{+}.N \\ [\sigma] N \\ [\mu] N \end{array}$	M M	multi-quantified negative types with movable variables
$\hat{\alpha}^+$::=	$\hat{\alpha}^+$		positive unification variable
\hat{lpha}^-	::=	$\widehat{lpha}^ \widehat{lpha}^{\{N,M\}}$		negative unification variable
$\widehat{\alpha}^ \overrightarrow{\alpha}^+, \ \overrightarrow{\beta}^+$::=	$ \frac{\widehat{\alpha}^{+}}{\widehat{\alpha}^{+}\{vars\}} $ $ \overrightarrow{\widehat{\alpha}^{+}}_{i} $ $ \overrightarrow{\widehat{\alpha}^{+}}_{i} $		positive unification variable list empty list a variable from a normal variable from a normal variable, context unspecified concatenate lists
$\overrightarrow{\widehat{\alpha^-}}, \overrightarrow{\widehat{\beta^-}}$::=	$ \begin{array}{c} \widehat{\alpha}^{-} \\ \widehat{\alpha}^{-} \{vars\} \\ \widehat{\alpha}^{-} \\ \widehat{\alpha}^{-}_{i} \end{array} $		negative unification variable list empty list a variable from a normal variable from a normal variable, context unspecified concatenate lists

```
P,\ Q \qquad ::= \qquad \text{a positive algorithmic type (potentially with metavariables)} \\ \mid \quad \alpha^+ \\ \mid \quad \widetilde{\alpha}^+ \\ \mid \quad \widehat{\alpha}^+ \{vars\} \\ \mid \quad \downarrow N \\ \mid \quad \exists \alpha^-. P \\ \mid \quad [\sigma] P \qquad \mathsf{M} \\ \\
```

 $[\mu]P$

 $[\mu]N$

 $\mathbf{nf}(P')$

Μ

Μ

Μ

```
formula
                                 judgement
                                 formula_1 .. formula_n
                                 \mu: vars_1 \leftrightarrow vars_2
                                 \mu is bijective
                                  \hat{\sigma} is functional
                                 \hat{\sigma}_1 \in \hat{\sigma}_2
                                  vars_1 \subseteq vars_2
                                  vars_1 = vars_2
                                  vars is fresh
                                  \alpha^- \not\in \mathit{vars}
                                  \alpha^+ \notin vars
                                  \alpha^- \in vars
                                  \alpha^+ \in \mathit{vars}
                                  if any other rule is not applicable
                                  N \neq M
                                 P \neq Q
E1A
                                n \models N \simeq_1^A M = \mun \models P \simeq_1^A Q = \mu
                                                                                                              Negative multi-quantified type equivalence (algorit
                                                                                                              Positive multi-quantified type equivalence (algorith
A
                                 \Gamma \vDash \overline{N} \leqslant M \dashv \widehat{\sigma}
                                                                                                              Negative subtyping
                                 \Gamma \vDash P \geqslant Q \Rightarrow \hat{\sigma}
                                                                                                              Positive supertyping
E1
                          | N \simeq_1^D M 
 | P \simeq_1^D Q 
                                                                                                              Negative multi-quantified type equivalence
                                                                                                              Positive multi-quantified type equivalence
D1
                                                                                                              Negative equivalence on MQ types
                                 \Gamma \vdash N \simeq_1^{\leqslant} M
                                 \Gamma \vdash P \simeq_{1}^{\leq} Q
\Gamma \vdash N \leqslant_{1} M
\Gamma \vdash P \geqslant_{1} Q
                                                                                                              Positive equivalence on MQ types
                                                                                                              Negative subtyping
                                                                                                              Positive supertyping
D\theta
                          \begin{array}{c|c} & \Gamma \vdash N \simeq_0^{\varsigma} M \\ & \Gamma \vdash P \simeq_0^{\varsigma} Q \\ & \Gamma \vdash N \leqslant_0 M \end{array} 
                                                                                                              Negative equivalence
                                                                                                              Positive equivalence
                                                                                                              Negative subtyping
                                                                                                              Positive supertyping
LUBF
                                 P_1 \vee P_2 === Q
                                 \operatorname{ord} \operatorname{varsin} P === \overrightarrow{\alpha}
                                 \operatorname{ord} \operatorname{varsin} N === \overrightarrow{\alpha}
                                 \begin{array}{l} \mathbf{ord} \ vars \mathbf{in} \ P = = = \overrightarrow{\alpha} \\ \mathbf{ord} \ vars \mathbf{in} \ N = = = \overrightarrow{\alpha} \end{array}
                                 \mathbf{nf}(N') === N
```

$$\left|\begin{array}{c} \mathbf{nf}\left(P'\right) ===P \\ \mathbf{nf}\left(N'\right) ===N \\ \mathbf{nf}\left(P'\right) ===P \\ \partial_{1}k\widehat{\sigma}_{2} ===\widehat{\sigma} \end{array} \right|$$

$$LUB \qquad ::= \\ \left|\begin{array}{c} P_{1} \vee P_{2} = Q \\ P_{1} \vee P_{2} = Q \end{array} \right|$$
 Least Upper Bound (Least Common Supertype)
$$AU \qquad ::= \\ \left|\begin{array}{c} \Gamma \vDash P_{1} \stackrel{\mathcal{S}}{\sim} P_{2} = \left(Q,\widehat{\sigma}_{1},\widehat{\sigma}_{2}\right) \\ \Gamma \vDash N_{1} \stackrel{\mathcal{S}}{\sim} N_{2} = \left(M,\widehat{\sigma}_{1},\widehat{\sigma}_{2}\right) \\ \end{array} \right|$$

$$\Gamma \vDash N_{1} \stackrel{\mathcal{S}}{\sim} N_{2} = \left(M,\widehat{\sigma}_{1},\widehat{\sigma}_{2}\right)$$
 Order
$$::= \\ \left|\begin{array}{c} \operatorname{ord} \operatorname{vars in} N = \overrightarrow{\sigma} \\ \operatorname{ord} \operatorname{vars in} P = \overrightarrow{\sigma} \\ \end{array} \right|$$
 ord $\operatorname{vars in} P = \overrightarrow{\sigma}$
$$\operatorname{ord} \operatorname{vars in} P = \overrightarrow{\sigma} \\ \end{aligned}$$

$$\operatorname{ord} \operatorname{vars in} P = \overrightarrow{\sigma} \\ \operatorname{ord} \operatorname{vars in} P = \overrightarrow{\sigma} \\ \end{aligned}$$

$$\operatorname{ord} \operatorname{var in} P = \overrightarrow{\sigma} \\ \end{aligned}$$

$$\operatorname{ord} \operatorname{va$$

 $\frac{Nrm}{SM}$

```
U
                                                            WF
  user\_syntax
                                                           \alpha
                                                           terminals
                                                           formula
n \models N \simeq^A_1 M = \mu
                                                     Negative multi-quantified type equivalence (algorithmic)
                                                                              \frac{n \vDash \alpha^{-} \simeq_{1}^{A} \alpha^{-} \dashv \cdot}{n \vDash P \simeq_{1}^{A} Q \dashv \mu} \quad \text{E1ANVAR}
\frac{n \vDash P \simeq_{1}^{A} Q \dashv \mu}{n \vDash \uparrow P \simeq_{1}^{A} \uparrow Q \dashv \mu} \quad \text{E1ASHIFTU}
```

$$n \vDash \uparrow P \simeq_1^A \uparrow Q = \mu$$

$$n \vDash P \simeq_1^A Q = \mu_1 \quad n \vDash N \simeq_1^A M = \mu_2 \quad \mu_1 \cup \mu_2 \text{ is bijective}$$

$$n \vDash P \to N \simeq_1^A Q \to M = \mu_1 \cup \mu_2$$

$$E1AARROW$$

$$\frac{n+1 \vDash [\overrightarrow{\alpha^{+n}}/\overrightarrow{\alpha^{+}}]N \simeq_{1}^{A} [\overrightarrow{\beta^{+n}}/\overrightarrow{\beta^{+}}]M \dashv \mu}{n \vDash \forall \overrightarrow{\alpha^{+}}.N \simeq_{1}^{A} \forall \overrightarrow{\beta^{+}}.M \dashv \mu|_{\mathbf{mv}\,M}} \qquad \text{E1AFORALL}$$

$$\overline{n \vDash \widetilde{\alpha}^{-n} \simeq_{1}^{A} \widetilde{\beta}^{-n} \dashv \widetilde{\beta}^{-n} \mapsto \widetilde{\alpha}^{-n}} \qquad \text{E1ANMVAR}$$

 $n \models P \simeq_1^A Q = \mu$ Positive multi-quantified type equivalence (algorithmic)

$$\frac{n \vDash \alpha^{+} \simeq_{1}^{A} \alpha^{+} \dashv \cdot}{n \vDash \sqrt{N} \simeq_{1}^{A} \sqrt{M} \dashv \mu} \qquad \text{E1ASHIFTD}$$

$$\frac{n \vDash N \simeq_{1}^{A} M \dashv \mu}{n \vDash \sqrt{N} \simeq_{1}^{A} \sqrt{M} \dashv \mu} \qquad \text{E1ASHIFTD}$$

$$\frac{n+1 \vDash [\overrightarrow{\alpha^{-n}}/\overrightarrow{\alpha^{-}}]P \simeq_{1}^{A} [\overrightarrow{\beta^{-n}}/\overrightarrow{\beta^{-}}]Q \dashv \mu}{n \vDash \overrightarrow{\alpha^{-}} \cdot P \simeq_{1}^{A} \overrightarrow{\beta^{+}} \cdot Q \dashv \mu|_{\mathbf{mv} Q}} \qquad \text{E1AEXISTS}$$

$$\frac{n+1 \vDash [\overrightarrow{\alpha^{-n}}/\overrightarrow{\alpha^{-}}]P \simeq_{1}^{A} (\overrightarrow{\beta^{-n}}/\overrightarrow{\beta^{-}})Q \dashv \mu|_{\mathbf{mv} Q}}{n \vDash \overrightarrow{\alpha^{+n}} \simeq_{1}^{A} \widetilde{\beta^{+n}} \dashv \widetilde{\beta^{+n}} \mapsto \widetilde{\alpha^{+n}}} \qquad \text{E1APMVAR}$$

 $\Gamma \models N \leqslant M \dashv \widehat{\sigma}$ Negative subtyping

$$\frac{\mathbf{nf}(P) \stackrel{u}{\simeq} \mathbf{nf}(Q) = \widehat{\sigma}}{\Gamma \models \uparrow P \leqslant \uparrow Q = \widehat{\sigma}} \quad \text{ASHIFTU}$$

$$\frac{\mathbf{nf}(P) \stackrel{u}{\simeq} \mathbf{nf}(Q) = \widehat{\sigma}}{\Gamma \models \uparrow P \leqslant \uparrow Q = \widehat{\sigma}} \quad \text{ASHIFTU}$$

$$\frac{\Gamma \models P \geqslant Q = \widehat{\sigma}_1 \quad \Gamma \models N \leqslant M = \widehat{\sigma}_2}{\Gamma \models P \rightarrow N \leqslant Q \rightarrow M = \widehat{\sigma}_1 \& \widehat{\sigma}_2} \quad \text{AARROW}$$

$$\frac{\Gamma, \overrightarrow{\beta^+} \models [\widehat{\alpha}^+ \{\Gamma, \overrightarrow{\beta^+}\} / \overrightarrow{\alpha^+}] N \leqslant M = \widehat{\sigma}}{\Gamma \models \forall \overrightarrow{\alpha^+}. N \leqslant \forall \overrightarrow{\beta^+}. M = \widehat{\sigma} \backslash \overrightarrow{\widehat{\alpha}^+}} \quad \text{AFORALL}$$

 $\Gamma \models P \geqslant Q \dashv \widehat{\sigma}$ Positive supertyping

$$\frac{\mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}}{\Gamma \vDash \lambda N \geqslant \lambda M \dashv \widehat{\sigma}} \quad \text{ASHIFTD}$$

$$\frac{\mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}}{\Gamma \vDash \lambda N \geqslant \lambda M \dashv \widehat{\sigma}} \quad \text{ASHIFTD}$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}} \vDash [\widehat{\alpha}^{-}\{\Gamma, \overrightarrow{\beta^{-}}\}/\widehat{\alpha^{-}}]P \geqslant Q \dashv \widehat{\sigma}}{\Gamma \vDash \overrightarrow{\beta\alpha^{-}}.P \geqslant \overrightarrow{\beta\beta^{-}}.Q \dashv \widehat{\sigma}} \quad \text{AEXISTS}$$

$$\frac{\mathbf{nf}(P) = P' \quad vars_1 = \mathbf{fv} P' \setminus vars \quad vars_2 \mathbf{is} \mathbf{fresh}}{\Gamma \vDash \widehat{\alpha}^{+}\{vars\} \geqslant P \dashv (\widehat{\alpha}^{+} : \geqslant P' \vee [vars_2/vars_1]P')} \quad \text{APUVAR}$$

 $|N \simeq_1^D M|$ Negative multi-quantified type equivalence

$$\frac{\alpha^{-} \simeq_{1}^{D} \alpha^{-}}{\alpha^{-} \simeq_{1}^{D} Q} \quad \text{E1ShiftU}$$

$$\frac{P \simeq_{1}^{D} Q}{\uparrow P \simeq_{1}^{D} \uparrow Q} \quad \text{E1ShiftU}$$

$$\frac{P \simeq_{1}^{D} Q \quad N \simeq_{1}^{D} M}{P \to N \simeq_{1}^{D} Q \to M} \quad \text{E1Arrow}$$

$$\frac{\overrightarrow{\alpha^{+}} \cap \mathbf{fv} \, M = \varnothing \quad \mu : (\overrightarrow{\beta^{+}} \cap \mathbf{fv} \, M) \leftrightarrow (\overrightarrow{\alpha^{+}} \cap \mathbf{fv} \, N) \quad N \simeq_{1}^{D} [\mu] M}{\forall \overrightarrow{\alpha^{+}} . N \simeq_{1}^{D} \forall \overrightarrow{\beta^{+}} . M} \quad \text{E1Forall}$$

 $P \simeq_1^D Q$ Positive multi-quantified type equivalence

$$\frac{\alpha^{+} \simeq_{1}^{D} \alpha^{+}}{\alpha^{+}} \quad \text{E1PVAR}$$

$$\frac{N \simeq_{1}^{D} M}{\downarrow N \simeq_{1}^{D} \downarrow M} \quad \text{E1SHIFTD}$$

$$\overrightarrow{\alpha^{-}} \cap \text{fv } Q = \varnothing \quad \mu : (\overrightarrow{\beta^{-}} \cap \text{fv } Q) \leftrightarrow (\overrightarrow{\alpha^{-}} \cap \text{fv } P) \quad P \simeq_{1}^{D} [\mu]Q$$

$$\overrightarrow{\exists \alpha^{-}} . P \simeq_{1}^{D} \overrightarrow{\exists \beta^{-}} . Q$$

$$\text{E1EXISTS}$$

 $\Gamma \vdash N \simeq M$ Negative equivalence on MQ types

$$\frac{\Gamma \vdash N \leqslant_{1} M \quad \Gamma \vdash M \leqslant_{1} N}{\Gamma \vdash N \simeq_{1}^{\leqslant} M} \quad \text{D1NDEF}$$

 $\Gamma \vdash P \simeq Q$ Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geqslant_1 Q \quad \Gamma \vdash Q \geqslant_1 P}{\Gamma \vdash P \simeq_1^{\varsigma} Q} \quad \text{D1PDEF}$$

 $\overline{|\Gamma \vdash N \leq_1 M|}$ Negative subtyping

$$\frac{\Gamma \vdash \alpha^{-} \leqslant_{1} \alpha^{-}}{\Gamma \vdash P \leqslant_{1}^{-} Q} \quad D1SHIFTU$$

$$\frac{\Gamma \vdash P \leqslant_{1}^{-} Q}{\Gamma \vdash P \leqslant_{1} \uparrow Q} \quad D1SHIFTU$$

$$\frac{\Gamma \vdash P \geqslant_{1} Q \quad \Gamma \vdash N \leqslant_{1} M}{\Gamma \vdash P \to N \leqslant_{1} Q \to M} \quad D1ARROW$$

$$\frac{\Gamma, \overrightarrow{\beta^{+}} \vdash P_{i} \quad \Gamma, \overrightarrow{\beta^{+}} \vdash [\overrightarrow{P}/\overrightarrow{\alpha^{+}}]N \leqslant_{1} M}{\Gamma \vdash \forall \overrightarrow{\alpha^{+}}, N \leqslant_{1} \forall \overrightarrow{\beta^{+}}, M} \quad D1FORALL$$

 $\overline{|\Gamma \vdash P \geqslant_1 Q|}$ Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{1} \alpha^{+}}{\Gamma \vdash N \cong_{1}^{s} M} \quad D1PVAR$$

$$\frac{\Gamma \vdash N \cong_{1}^{s} M}{\Gamma \vdash \downarrow N \geqslant_{1} \downarrow M} \quad D1SHIFTD$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}} \vdash N_{i} \quad \Gamma, \overrightarrow{\beta^{-}} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^{-}}]P \geqslant_{1} Q'}{\Gamma \vdash \exists \overrightarrow{\alpha^{-}}.P \geqslant_{1} \exists \overrightarrow{\beta^{-}}.Q} \quad D1EXISTSL$$

 $\Gamma \vdash N \simeq_0^{\leqslant} M$ Negative equivalence

$$\frac{\Gamma \vdash N \leqslant_0 M \quad \Gamma \vdash M \leqslant_0 N}{\Gamma \vdash N \simeq_0^{\epsilon} M} \quad \text{D0NDEF}$$

 $\Gamma \vdash P \simeq_0^{\leqslant} Q$ Positive equivalence

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash Q \geqslant_0 P}{\Gamma \vdash P \simeq_0^{\varsigma} Q} \quad \text{D0PDEF}$$

 $\Gamma \vdash N \leqslant_0 M$ Negative subtyping

$$\frac{\Gamma \vdash a - \leqslant_0 a -}{\Gamma \vdash P = \circ_0 Q} \quad D0 \text{NVAR}$$

$$\frac{\Gamma \vdash P = \circ_0 Q}{\Gamma \vdash \uparrow P \leqslant_0 \uparrow Q} \quad D0 \text{SHIFTU}$$

$$\frac{\Gamma \vdash P \quad \Gamma \vdash [P/a +] N \leqslant_0 M \quad M \neq \forall \beta^+ . M'}{\Gamma \vdash \forall \alpha^+ . N \leqslant_0 M} \quad D0 \text{FORALLL}$$

$$\frac{\Gamma, \alpha^+ \vdash N \leqslant_0 M}{\Gamma \vdash N \leqslant_0 \forall \alpha^+ . M} \quad D0 \text{FORALLR}$$

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash N \leqslant_0 M}{\Gamma \vdash P \to N \leqslant_0 Q \to M} \quad D0 \text{ARROW}$$

 $\overline{|\Gamma \vdash P \geqslant_0 Q|}$ Positive supertyping

$$\frac{\Gamma \vdash a + \geqslant_0 a +}{\Gamma \vdash a + \geqslant_0 a +} \quad \text{D0PVar}$$

$$\frac{\Gamma \vdash N \simeq_0^{\leq} M}{\Gamma \vdash \downarrow N \geqslant_0 \downarrow M} \quad \text{D0ShiftD}$$

$$\frac{\Gamma \vdash N \quad \Gamma \vdash [N/a -] P \geqslant_0 Q \quad Q \neq \exists \alpha^-. Q'}{\Gamma \vdash \exists \alpha^-. P \geqslant_0 Q} \quad \text{D0ExistsL}$$

$$\frac{\Gamma, \alpha^- \vdash P \geqslant_0 Q}{\Gamma \vdash P \geqslant_0 \exists \alpha^-. Q} \quad \text{D0ExistR}$$

 $P_1 \vee P_2$

ord varsin P

ord vars in N

ord vars in P

 $\overline{\mathbf{ord} \ vars \mathbf{in} \ N}$

 $\mathbf{nf}(N')$

 $\mathbf{nf}(P')$

 $\mathbf{nf}(N')$

 $\mathbf{nf}(P')$

 $\hat{\sigma}_1 \& \hat{\sigma}_2$

 $\overline{|P_1 \vee P_2 = Q|}$ Least Upper Bound (Least Common Supertype)

$$\frac{\overline{\alpha^{+} \vee \alpha^{+} = \alpha^{+}}}{(\mathbf{f} \mathbf{v} \, N \cup \mathbf{f} \mathbf{v} \, M) \vDash \downarrow N \stackrel{a}{\simeq} \downarrow M = (P, \hat{\sigma}_{1}, \hat{\sigma}_{2})} \quad \text{LUBShift}$$

$$\frac{1}{\sqrt{N} \vee \sqrt{M} = \exists \alpha^{-}. [\alpha^{-}/\mathbf{u} \mathbf{v} \, P] P} \qquad \qquad \text{LUBShift}$$

$$\frac{1}{\sqrt{N} \vee \sqrt{M} = \exists \alpha^{-}. [\alpha^{-}/\mathbf{u} \mathbf{v} \, P] P} \qquad \qquad \text{LUBEXISTS}$$

$$\frac{1}{\sqrt{N} \vee \sqrt{M} = \exists \alpha^{-}. P_{1} \vee \exists \beta^{-}. P_{2} = P_{1} \vee P_{2}} \qquad \qquad \text{LUBEXISTS}$$

 $\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (Q, \hat{\sigma}_1, \hat{\sigma}_2)$

$$\frac{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} \alpha^{+} \Rightarrow (\alpha^{+}, \cdot, \cdot)}{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} N_{2} \Rightarrow (M, \hat{\sigma}_{1}, \hat{\sigma}_{2})} \quad \text{AUPShift}$$

$$\frac{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \Rightarrow (M, \hat{\sigma}_{1}, \hat{\sigma}_{2})}{\Gamma \vDash \downarrow N_{1} \stackrel{a}{\simeq} \downarrow N_{2} \Rightarrow (\downarrow M, \hat{\sigma}_{1}, \hat{\sigma}_{2})} \quad \text{AUPShift}$$

$$\frac{\overrightarrow{\alpha^{-}} \cap \Gamma = \varnothing \quad \Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \Rightarrow (Q, \hat{\sigma}_{1}, \hat{\sigma}_{2})}{\Gamma \vDash \overrightarrow{\beta \alpha^{-}} . P_{1} \stackrel{a}{\simeq} \overrightarrow{\beta \alpha^{-}} . P_{2} \Rightarrow (\overrightarrow{\beta \alpha^{-}} . Q, \hat{\sigma}_{1}, \hat{\sigma}_{2})} \quad \text{AUPExists}$$

 $\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 \Rightarrow (M, \widehat{\sigma}_1, \widehat{\sigma}_2)$

$$\frac{\Gamma \vDash \alpha^{-\frac{a}{2}} \alpha^{-} \dashv (\alpha^{-}, \cdot, \cdot)}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (Q, \widehat{\sigma}_{1}, \widehat{\sigma}_{2})} \quad \text{AUNSHIFT}
\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (Q, \widehat{\sigma}_{1}, \widehat{\sigma}_{2})}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} \uparrow P_{2} \dashv (\uparrow Q, \widehat{\sigma}_{1}, \widehat{\sigma}_{2})} \quad \text{AUNSHIFT}
\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (Q, \widehat{\sigma}_{1}, \widehat{\sigma}_{2}) \quad \Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (M, \widehat{\sigma}'_{1}, \widehat{\sigma}'_{2})}{\Gamma \vDash P_{1} \rightarrow N_{1} \stackrel{a}{\simeq} P_{2} \rightarrow N_{2} \dashv (Q \rightarrow M, \widehat{\sigma}_{1} \cup \widehat{\sigma}'_{1}, \widehat{\sigma}_{2} \cup \widehat{\sigma}'_{2})} \quad \text{AUNARROW}
\frac{\text{if any other rule is not applicable} \quad \Gamma \vDash N \quad \Gamma \vDash M \quad \text{AUNAU}}{\Gamma \vDash N \stackrel{a}{\simeq} M \dashv (\widehat{\alpha}^{-}_{\{N,M\}}, (\widehat{\alpha}^{-}_{\{N,M\}} : \approx N), (\widehat{\alpha}^{-}_{\{N,M\}} : \approx M))} \quad \text{AUNAU}$$

 $\mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}$

$$\frac{\alpha^- \in vars}{\mathbf{ord} \ vars \ \mathbf{in} \ \alpha^- = \alpha^-} \quad \text{ONVARIN}$$

$$\frac{\alpha^- \notin vars}{\mathbf{ord} \ vars \ \mathbf{in} \ \alpha^- = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\text{ord } vars \text{ in } \widehat{\alpha}^{-}\{vars'\} = \cdot}{\text{ord } vars \text{ in } P = \overrightarrow{\alpha}} \quad \text{ONUVAR}$$

$$\frac{\text{ord } vars \text{ in } P = \overrightarrow{\alpha}}{\text{ord } vars \text{ in } N = \overrightarrow{\alpha}_{2}} \quad \text{OSHIFTU}$$

$$\frac{\text{ord } vars \text{ in } P = \overrightarrow{\alpha}_{1} \quad \text{ord } vars \text{ in } N = \overrightarrow{\alpha}_{2}}{\text{ord } vars \text{ in } P \to N = \overrightarrow{\alpha}_{1}, (\overrightarrow{\alpha}_{2} \setminus \{\overrightarrow{\alpha}_{1}\})} \quad \text{OARROW}$$

$$\frac{vars \cap \overrightarrow{\alpha^{+}} = \emptyset \quad \text{ord } vars \text{ in } N = \overrightarrow{\alpha}}{\text{ord } vars \text{ in } \bigvee \overrightarrow{\alpha^{+}}. N = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

$\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}$

$$\frac{\alpha^{+} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \alpha^{+}} \quad \operatorname{OPVarIn}$$

$$\frac{\alpha^{+} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \cdot} \quad \operatorname{OPVarNIn}$$

$$\overline{\operatorname{ord} vars \operatorname{in} \widehat{\alpha}^{+} \{ vars' \} = \cdot} \quad \operatorname{OPUVar}$$

$$\frac{\operatorname{ord} vars \operatorname{in} \widehat{\alpha}^{+} \{ vars' \} = \cdot}{\operatorname{ord} vars \operatorname{in} \widehat{\lambda}^{-} = \overrightarrow{\alpha}} \quad \operatorname{OSHIFTD}$$

$$\frac{vars \cap \overrightarrow{\alpha^{-}} = \varnothing \quad \operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \overrightarrow{\beta} \overrightarrow{\alpha^{-}} . P = \overrightarrow{\alpha}} \quad \operatorname{OEXISTS}$$

$$\overline{\operatorname{ord} vars \operatorname{in} \exists \overrightarrow{\alpha^{-}} . P = \overrightarrow{\alpha}} \quad \operatorname{OEXISTS}$$

$$\begin{aligned}
\mathbf{nf}(N) &= M \\
\mathbf{nf}(P) &= Q \\
\mathbf{nf}(N) &= M
\end{aligned}$$

$$\overline{\mathbf{nf}(\alpha^{-}) = \alpha^{-}} \quad \text{NrmNVar}$$

$$\overline{\mathbf{nf}(\widehat{\alpha}^{-}\{vars\}) = \widehat{\alpha}^{-}\{vars\}} \quad \text{NrmNUVar}$$

$$\frac{\mathbf{nf}(P) = Q}{\mathbf{nf}(\uparrow P) = \uparrow Q} \quad \text{NrmShiftU}$$

$$\underline{\mathbf{nf}(P) = Q \quad \mathbf{nf}(N) = M} \quad \text{NrmArrow}$$

$$\underline{\mathbf{nf}(P) = Q \quad \mathbf{nf}(N) = M} \quad \text{NrmArrow}$$

$$\underline{\mathbf{nf}(P) = N' \quad \mathbf{ord} \stackrel{\rightarrow}{\alpha^{+}} \text{in} \quad N' = \stackrel{\rightarrow}{\alpha^{+'}}} \quad \text{NrmForall}$$

$$\underline{\mathbf{nf}(N) = N' \quad \mathbf{ord} \stackrel{\rightarrow}{\alpha^{+'}} . N'} \quad \text{NrmForall}$$

 $\mathbf{nf}(P) = Q$

$$\frac{\mathbf{nf}(\alpha^{+}) = \alpha^{+}}{\mathbf{nf}(\widehat{\alpha}^{+}\{vars\}) = \widehat{\alpha}^{+}\{vars\}} \quad \text{NRMPUVAR}$$

$$\frac{\mathbf{nf}(N) = M}{\mathbf{nf}(\downarrow N) = \downarrow M} \quad \text{NRMSHIFTD}$$

$$\frac{\mathbf{nf}(P) = P' \quad \mathbf{ord} \stackrel{\longrightarrow}{\alpha^{-}} \mathbf{in} \quad P' = \stackrel{\longrightarrow}{\alpha^{-'}}}{\mathbf{nf}(\exists \widehat{\alpha^{-}}.P) = \exists \widehat{\alpha^{-'}}.P'} \quad \text{NRMEXISTS}$$

 $e_1 \& e_2 = e_3$ Unification Solution Entry Merge

$$\overline{\hat{\alpha}^{+}} : \geqslant P \& \hat{\alpha}^{+} : \geqslant Q = \hat{\alpha}^{+} : \geqslant P \lor Q \qquad \text{SMEPSUPSUP}$$

$$\frac{\mathbf{fv} \, P \cup \mathbf{fv} \, Q \vDash P \geqslant Q \dashv \hat{\sigma}'}{\hat{\alpha}^{+}} : \approx P \& \hat{\alpha}^{+} : \geqslant Q = \hat{\alpha}^{+} : \approx P \qquad \text{SMEPEQSUP}$$

$$\frac{\mathbf{fv} \, P \cup \mathbf{fv} \, Q \vDash Q \geqslant P \dashv \hat{\sigma}'}{\hat{\alpha}^{+}} : \geqslant P \& \hat{\alpha}^{+} : \approx Q = \hat{\alpha}^{+} : \approx Q \qquad \text{SMEPSUPEQ}$$

$$\overline{\hat{\alpha}^{+}} : \approx P \& \hat{\alpha}^{+} : \approx P = \hat{\alpha}^{+} : \approx P \qquad \text{SMEPEQEQ}$$

$$\overline{\hat{\alpha}^{+}} : \approx N \& \hat{\alpha}^{-} : \approx N = \hat{\alpha}^{-} : \approx N \qquad \text{SMENEQEQ}$$

 $\widehat{\sigma}_1 \& \widehat{\sigma}_2 = \widehat{\sigma}_3$ Merge unification solutions

$$\overline{\cdot \& \widehat{\sigma} = \widehat{\sigma} } \quad \text{SMEMPTY}$$

$$(\widehat{\alpha}^+ :\approx P) \in \widehat{\sigma}_2$$

$$\widehat{\sigma}_1 \& (\widehat{\sigma}_2 \backslash \widehat{\alpha}^+) = \widehat{\sigma}_3$$

$$\overline{(\widehat{\alpha}^+ :\approx P, \widehat{\sigma}_1)} \& \widehat{\sigma}_2 = (\widehat{\alpha}^+ :\approx P, \widehat{\sigma}_3) \quad \text{SMPEQEQ}$$

$$\frac{(\widehat{\alpha}^+ :\geqslant Q) \in \widehat{\sigma}_2 \quad \widehat{\sigma}_1 \& (\widehat{\sigma}_2 \backslash \widehat{\alpha}^+) = \widehat{\sigma}_3 }{(\widehat{\alpha}^+ :\geqslant P, \widehat{\sigma}_1)} \& \widehat{\sigma}_2 = (\widehat{\alpha}^+ :\geqslant P \vee Q, \widehat{\sigma}_3) \quad \text{SMPSUPSUP}$$

$$(\widehat{\alpha}^+ :\approx Q) \in \widehat{\sigma}_2 \quad \text{fv } Q \cup \text{fv } P \vDash Q \geqslant P \dashv \widehat{\sigma}'$$

$$\widehat{\sigma}_1 \& (\widehat{\sigma}_2 \backslash \widehat{\alpha}^+) = \widehat{\sigma}_3 \quad \text{SMPSUPEQ}$$

$$(\widehat{\alpha}^+ :\geqslant P, \widehat{\sigma}_1) \& \widehat{\sigma}_2 = (\widehat{\alpha}^+ :\approx Q, \widehat{\sigma}_3) \quad \text{SMPSUPEQ}$$

$$(\widehat{\alpha}^+ :\geqslant Q) \in \widehat{\sigma}_2 \quad \text{fv } Q \cup \text{fv } P \vDash P \geqslant Q \dashv \widehat{\sigma}'$$

$$\widehat{\sigma}_1 \& (\widehat{\sigma}_2 \backslash \widehat{\alpha}^+) = \widehat{\sigma}_3 \quad \text{SMPEQSUP}$$

$$(\widehat{\alpha}^+ :\approx P, \widehat{\sigma}_1) \& \widehat{\sigma}_2 = (\widehat{\alpha}^+ :\approx P, \widehat{\sigma}_3) \quad \text{SMPEQSUP}$$

$$(\widehat{\alpha}^- :\approx N) \in \widehat{\sigma}_2 \quad \widehat{\sigma}_1 \& (\widehat{\sigma}_2 \backslash \widehat{\alpha}^-) = \widehat{\sigma}_3 \quad \text{SMNEQEQ}$$

$$\widehat{(\widehat{\alpha}^- :\approx N, \widehat{\sigma}_1)} \& \widehat{\sigma}_2 = (\widehat{\alpha}^- :\approx N, \widehat{\sigma}_3) \quad \text{SMNEQEQ}$$

 $|N \stackrel{u}{\simeq} M = \widehat{\sigma}|$ Negative unification

$$\frac{A^{-\frac{u}{2}}\alpha^{-} \dashv \cdot}{\alpha^{-\frac{u}{2}}\alpha^{-} \dashv \cdot} \quad \text{UNVAR}$$

$$\frac{P \overset{u}{\cong} Q \dashv \widehat{\sigma}}{\uparrow P \overset{u}{\cong} \uparrow Q \dashv \widehat{\sigma}} \quad \text{USHIFTU}$$

$$\frac{P \overset{u}{\cong} Q \dashv \widehat{\sigma}_{1} \quad N \overset{u}{\cong} M \dashv \widehat{\sigma}_{2}}{P \to N \overset{u}{\cong} Q \to M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \quad \text{UARROW}$$

$$\frac{N \overset{u}{\cong} M \dashv \widehat{\sigma}}{\forall \alpha^{+}.N \overset{u}{\cong} \forall \alpha^{+}.M \dashv \widehat{\sigma}} \quad \text{UFORALL}$$

$$\frac{\mathbf{fv} N \subseteq vars}{\widehat{\alpha}^{-} \{vars\} \overset{u}{\cong} N \dashv \widehat{\alpha}^{-} :\approx N} \quad \text{UNUVAR}$$

 $P \stackrel{u}{\simeq} Q = \widehat{\sigma}$ Positive unification

$$\frac{\alpha^{+} \stackrel{u}{\simeq} \alpha^{+} \dashv \cdot}{\frac{N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}}{\downarrow N \stackrel{u}{\simeq} \downarrow M \dashv \widehat{\sigma}}} \quad \text{USHIFTD}$$

$$\frac{P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}}{\exists \alpha^{-}.P \stackrel{u}{\simeq} \exists \alpha^{-}.Q \dashv \widehat{\sigma}} \quad \text{UEXISTS}$$

$$\frac{\mathbf{fv} P \subseteq vars}{\widehat{\alpha}^{+} \{vars\} \stackrel{u}{\simeq} P \dashv \widehat{\alpha}^{+} : \approx P} \quad \text{UPUVAR}$$

 $\begin{array}{|c|c|c|c|}\hline \Gamma \vdash N & \text{Negative type well-formedness} \\ \hline \Gamma \vdash P & \text{Positive type well-formedness} \\ \hline \Gamma \vdash N & \text{Negative type well-formedness} \\ \end{array}$

 $\overline{\Gamma \vdash P}$ Positive type well-formedness

Definition rules: 94 good 0 bad Definition rule clauses: 165 good 0 bad