$\begin{array}{ll} \alpha,\,\beta,\,\alpha,\,\beta,\,\gamma,\,\delta & \text{type variables} \\ n,\,m,\,i,\,j & \text{index variables} \end{array}$

```
cohort index
n, m
                                                    ::=
                                                                 0
                                                                 n+1
\alpha^+,~\beta^+,~\gamma^+,~\delta^+
                                                                                                           positive variable
                                                                 \alpha^{+n}
\alpha^-,~\beta^-,~\gamma^-,~\delta^-
                                                                                                          negative variable
                                                                 \alpha^{-n}
\alpha^{\pm}, \ \beta^{\pm}
                                                                                                          positive or negative variable
                                                    ::=
                                                                 \alpha^{\pm}
                                                                 \alpha^{\pm n}
                                                    ::=
                                                                                                          substitution
                                                                 id
                                                                 P/\alpha^+
                                                                N/\alpha^-
\overrightarrow{P}/\alpha^+
                                                                 \mathbf{pmas}/\overrightarrow{\alpha^+}

\frac{\mathbf{nmas}}{\widetilde{\alpha}^{+}}/\widetilde{\alpha}^{+}

\overset{\longrightarrow}{\alpha^{+}}/\alpha^{+}

\overset{\longrightarrow}{\alpha^{-}}/\alpha^{-}

                                                                 \mu
                                                                 \sigma_1 \circ \sigma_2
                                                                 \vec{\alpha}_1/\vec{\alpha}_2
                                                                                                S
                                                                  (\sigma)
                                                                                                                 concatenate
                                                                 \mathbf{nf}\left(\sigma'\right)
                                                                                                Μ
                                                                 \sigma'|_{vars}
                                                                                                Μ
                                                                                                           entry of a unification solution
 e
                                                                 \hat{\alpha}^+ :\approx P
                                                                 \widehat{\alpha}^- :\approx N
                                                                 \hat{\alpha}^+ : \geqslant P
                                                                 (e)
                                                                                                S
                                                                 \hat{\sigma}(\hat{\alpha}^+)
                                                                                                Μ
                                                                 \hat{\sigma}(\hat{\alpha}^-)
                                                                                                Μ
                                                                 e_1 \ \& \ e_2
```

unification solution (substitution)

 $\hat{\sigma}$

::=

```
e
                                                    \widehat{\sigma} \backslash vars
                                                    \hat{\sigma}|vars
                                                    \hat{\sigma}_1 \cup \hat{\sigma}_2
                                                                                       concatenate
                                                    (\hat{\sigma})
                                                                          S
                                                    \mathbf{nf}\left(\widehat{\sigma}'\right)
                                                                          Μ
                                                    \hat{\sigma}'|_{vars}
                                                                          Μ
                                                    \hat{\sigma}_1 \& \hat{\sigma}_2
                                                                          Μ
\hat{\tau}
                                                                                   anti-unification substitution
                                         ::=
                                                    \widehat{\alpha}^-:\approx N
                                                    \hat{\alpha}^- :\approx N
                                                    \hat{\tau}_1 \cup \hat{\tau}_2
                                                                                       concatenate
                                                                          S
P, Q
                                                                                   positive types
                                                    \downarrow N
                                                    \exists \alpha^-.P
                                                    [\sigma]P
                                                                          Μ
N, M
                                                                                   negative types
                                                    \alpha^{-}
                                                    \uparrow P
                                                    \forall \alpha^+.N
                                                    P \rightarrow N
                                                    [\sigma]N
                                                                          Μ
                                                                                   positive variable list
                                                                                       empty list
                                                                                       a variable
                                                                                       a variable
                                                                                       concatenate lists
                                                                                  negative variables
                                                                                       empty list
                                                                                       a variable
                                                                                       variables
                                                                                       concatenate lists
\overrightarrow{\alpha^{\pm}}, \ \overrightarrow{\beta^{\pm}}, \ \overrightarrow{\gamma^{\pm}}, \ \overrightarrow{\delta^{\pm}}
                                                                                   positive or negative variable list
                                                                                       empty list
                                                                                       a variable
```

```
variables
                                                 concatenate lists
P, Q
                                              multi-quantified positive types
                                                 P \neq \exists \dots
                         [\sigma]P
                                       Μ
                         [\hat{\tau}]P
                                       Μ
                        [\hat{\sigma}]P
                                       Μ
                        [\mu]P
                                       Μ
                        (P)
                                       S
                        P_1 \vee P_2
                                       Μ
                        \mathbf{nf}(P')
                                       Μ
N, M
                                             multi-quantified negative types
                        \alpha^{-}
                        {\uparrow} P
                        \stackrel{P}{\overrightarrow{\alpha^+}}.N
                                                 N \neq \forall \dots
                        [\sigma]N
                                       Μ
                         [\mu]N
                                       Μ
                        [\hat{\sigma}]N
                                       Μ
                        (N)
                                       S
                        \mathbf{nf}(N')
\vec{P}, \ \vec{Q}
                                             list of positive types
                                                 empty list
                                                 a singel type
                                                 concatenate lists
                                       Μ
\overrightarrow{N}, \overrightarrow{M}
                                             list of negative types
                                                 empty list
                                                 a singel type
                                                 concatenate lists
                                       Μ
\Delta, \Gamma
                                              declarative type context
                                                 empty context
                                                 list of variables
                                                 list of variables
                                                 list of variables
                        vars
                        \overline{\Gamma_i}^{\;i}
                                                 concatenate contexts
                        (\Gamma)
                                       S
                        \Theta(\hat{\alpha}^+)
                                       Μ
```

 $\Theta(\hat{\alpha}^-)$

Μ

Θ	::=	$ \overrightarrow{\alpha^{+}} \overrightarrow{\alpha^{-}} $ $ \overrightarrow{\alpha^{-}} $ $ vars $ $ \overrightarrow{\Theta_{i}}^{i} $ $ (\Theta) $ $ \Theta _{vars} $ $ \Theta_{1} \cup \Theta_{2} $	S	unification type variable context empty context list of variables list of variables concatenate contexts leave only those variables that are in the set
Ξ	::= 	$ \overrightarrow{\alpha^{+}} \overrightarrow{\alpha^{-}} \overrightarrow{\overline{\Xi}_{i}}^{i} (\Xi) \Xi_{1} \cup \Xi_{2} $	S	anti-unification type variable context empty context list of variables list of variables concatenate contexts
$\vec{\alpha}$, $\vec{\beta}$. $\overrightarrow{\alpha^{+}}$ $\overrightarrow{\alpha^{-}}$ $\overrightarrow{\alpha^{+}}$ $\overrightarrow{\alpha^{-}}$ $\overrightarrow{\alpha^{+}}$ $\overrightarrow{\alpha^{-}}$ $\overrightarrow{\alpha}_{1} \setminus vars$ Γ $vars$ $\overrightarrow{\alpha}_{i}^{i}$ $(\overrightarrow{\alpha})$ $[\mu] \overrightarrow{\alpha}$ ord $vars$ in P	S M M M M	ordered positive or negative variables empty list list of variables list of variables list of variables list of variables setminus context concatenate contexts parenthesis apply moving to list
vars		$egin{array}{ll} \varnothing & & & \text{fv } P & & \\ & & & \text{fv imP} & & & \\ & & & \text{fv imN} & & & \\ & & & vars_1 \cap vars_2 & & \\ & & vars_1 \backslash vars_2 & & \\ & & \text{mv imP} & & \\ & & \text{mv imN} & & \\ & & \text{uv } N & & \\ & & \text{uv } P & & \\ & & \text{fv } N & & \\ \hline \end{array}$		set of variables empty set free variables free variables free variables free variables set intersection set union set complement movable variables movable variables unification variables unification variables free variables free variables

		$\begin{array}{l} \textbf{fv} \ P \\ (vars) \\ \overrightarrow{\alpha} \\ [\mu] vars \\ \textbf{dom} \ (\widehat{\sigma}) \\ \textbf{dom} \ (\Theta) \end{array}$	S M M	free variables parenthesis ordered list of variables apply moving to varset
μ	::=	$\begin{array}{l} .\\ pma1 \mapsto pma2 \\ nma1 \mapsto nma2 \\ \mu_1 \cup \mu_2 \\ \hline{\mu_1} \circ \mu_2 \\ \hline{\mu_i}^i \\ \mu _{vars} \\ \mu^{-1} \\ \mathbf{nf} \left(\mu' \right) \end{array}$	M M M M	empty moving Positive unit substitution Positive unit substitution Set-like union of movings Composition concatenate movings restriction on a set inversion
\widehat{lpha}^{\pm}	::=	\hat{lpha}^{\pm}		positive/negative unification variable
$\hat{\alpha}^+$		$\widehat{\alpha}^+$ $\widehat{\alpha}^+$ { Δ } $\widehat{\alpha}^\pm$		positive unification variable
$\hat{\alpha}^-,~\hat{eta}^-$::=	$egin{array}{l} \widehat{lpha}^- \ \widehat{lpha}^{\{N,M\}} \ \widehat{lpha}^{\{\Delta\}} \ \widehat{lpha}^\pm \end{array}$		negative unification variable
$\overrightarrow{\alpha}^+, \ \overrightarrow{\widetilde{\beta}^+}$::= 	$ \begin{array}{c} \widehat{\alpha}^{+} \\ \widehat{\alpha}^{+} \{\Delta\} \\ \widehat{\alpha}^{+} \\ \widehat{\alpha}^{+}_{i} \end{array} $		positive unification variable list empty list a variable from a normal variable from a normal variable, context unspecified concatenate lists
$\overrightarrow{\widehat{\alpha}^-}$, $\overrightarrow{\widehat{\beta}^-}$::=	$ \overrightarrow{\widehat{\alpha}^{+}}_{i}^{i} $ $ \overrightarrow{\widehat{\alpha}^{-}}_{i}^{i} $ $ \overrightarrow{\widehat{\alpha}^{-}}_{i}^{i} $ $ \overrightarrow{\widehat{\alpha}^{-}}_{i}^{i} $		negative unification variable list empty list a variable from an antiunification context from a normal variable from a normal variable, context unspecified concatenate lists
P, Q	::=	α^+		a positive algorithmic type (potentially with metavariables)

```
\mathbf{pma} \\ \widehat{\alpha}^+

\exists \stackrel{N}{\alpha} . P

                                   [\sigma]P
                                                               Μ
                                   [\hat{\tau}]P
                                                                Μ
                                   [\mu]P
                                                               Μ
                                   (P)
                                                               S
                                   \mathbf{nf}(P')
                                                               Μ
                                                                        a negative algorithmic type (potentially with metavariables)
N, M
                                    \alpha^{-}
                                   \hat{\alpha}^-
                                   \uparrow P
                                   P \rightarrow N
                                   \forall \overrightarrow{\alpha^+}.N
                                   [\sigma]N
                                                               Μ
                                   [\mu]N
                                                               Μ
                                                               S
                                    (N)
                                   \mathbf{nf}(N')
                                                               Μ
auSol
                         ::=
                                   (\Xi, Q, \widehat{	au}_1, \widehat{	au}_2)
terminals
                                   \forall
                                    \in
                                   ∉
                                    \leq
                                    \geqslant
                                    \cup
                                   \subseteq
                                    Ø
                                    \Rightarrow
                                    \models
```

```
:≥
formula
                                       judgement
                                       formula_1 .. formula_n
                                       \mu : vars_1 \leftrightarrow vars_2
                                       \mu is bijective
                                       \hat{\sigma} is functional
                                       \hat{\sigma}_1 \in \hat{\sigma}_2
                                       \hat{\sigma}_1 \subseteq \hat{\sigma}_2
                                       vars_1 \subseteq vars_2
                                       vars_1 = vars_2
                                       vars is fresh
                                       \alpha^- \not\in \mathit{vars}
                                       \alpha^+ \not\in \mathit{vars}
                                       \alpha^- \in vars
                                       \alpha^+ \in \mathit{vars}
                                       \widehat{\alpha}^- \in \Theta
                                       \widehat{\alpha}^+ \in \Theta
                                       if any other rule is not applicable
                                       \vec{\alpha}_1 = \vec{\alpha}_2
                                       e_1 = e_2
                                       N \neq M
                                        P \neq Q
\boldsymbol{A}
                            ::=
                                       \Gamma; \Theta \models \overline{N} \leqslant M \dashv \hat{\sigma}
                                                                                                                               Negative subtyping
                                       \Gamma; \Theta \models P \geqslant Q = \hat{\sigma}
                                                                                                                               Positive supertyping
AU
                            ::=
                                      \Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 \rightrightarrows (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)
\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 \rightrightarrows (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)
E1
                                      N \simeq_1^D M
P \simeq_1^D Q
P \simeq Q
                                                                                                                                Negative multi-quantified type equivalence
                                                                                                                                Positive multi-quantified type equivalence
D1
                              \begin{array}{c|c} .- & \\ & \Gamma \vdash N \simeq_1^{\varsigma} M \\ & \Gamma \vdash P \simeq_1^{\varsigma} Q \\ & \Gamma \vdash N \leqslant_1 M \\ & \Gamma \vdash P \geqslant_1 Q \\ & \Gamma_2 \vdash \sigma_1 \simeq_1^{\varsigma} \sigma_2 : \Gamma_1 \end{array} 
                                                                                                                               Negative equivalence on MQ types
                                                                                                                                Positive equivalence on MQ types
                                                                                                                               Negative subtyping
                                                                                                                               Positive supertyping
                                                                                                                               Equivalence of substitutions
```

```
D\theta
                     ::=
                              \Gamma \vdash N \simeq_0^{\leqslant} M
                                                                                           Negative equivalence
                              \Gamma \vdash P \simeq_0^{\scriptscriptstyle \leqslant} Q
                                                                                           Positive equivalence
                              \Gamma \vdash N \leqslant_0 M
                                                                                           Negative subtyping
                              \Gamma \vdash P \geqslant_0 Q
                                                                                           Positive supertyping
EQ
                     ::=
                              N=M
                                                                                           Negative type equality (alpha-equivalence)
                              P = Q
                                                                                           Positive type equuality (alphha-equivalence)
                              P = Q
LUBF
                    ::=
                              P_1 \vee P_2 === Q
                              ord vars in P === \vec{\alpha}
                              ord vars in N = = \vec{\alpha}
                              \operatorname{ord} \operatorname{varsin} P === \overrightarrow{\alpha}
                              ord vars in N === \vec{\alpha}
                              \mathbf{nf}(N') === N
                              \mathbf{nf}(P') === P
                              \mathbf{nf}(N') === N

\mathbf{nf}(P') === P \\
\mathbf{nf}(\vec{N}') === \vec{N}

                              \mathbf{nf}(\overrightarrow{P}') = = = \overrightarrow{P}
                              \mathbf{nf}(\sigma') = = = \sigma
                              \mathbf{nf}(\mu') === \mu
                              \mathbf{nf}(\widehat{\sigma}') ===\widehat{\sigma}
                              \sigma'|_{vars}
                              \hat{\sigma}'|_{vars}
                              e_1 \& e_2
                              \hat{\sigma}_1 \& \hat{\sigma}_2
                              \mathbf{dom}\left(\widehat{\sigma}\right) === vars
                              \mathbf{dom}\left(\Theta\right) === vars
LUB
                    ::=
                              \Gamma \vDash P_1 \vee P_2 = Q
                                                                                          Least Upper Bound (Least Common Supertype)
                              \mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q
Nrm
                     ::=
                              \mathbf{nf}\left( N\right) =M
                              \mathbf{nf}(P) = Q
                              \mathbf{nf}(N) = M
                              \mathbf{nf}(P) = Q
Order
                              \operatorname{ord} \operatorname{varsin} N = \overrightarrow{\alpha}
                              \mathbf{ord}\ vars \mathbf{in}\ P = \overrightarrow{\alpha}
                              ord vars in N = \vec{\alpha}
                              \mathbf{ord}\ vars \mathbf{in}\ P = \overrightarrow{\alpha}
```

```
SM
                                \Gamma \vdash e_1 \& e_2 = e_3
                                                                         Unification Solution Entry Merge
                                \Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3
                                                                         Merge unification solutions
SImp
                                \Gamma \vdash e_1 \Rightarrow e_2
                                                                         Weakening of unification solution entries
                                \Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2
                                                                         Weakening of unification solutions
                                 \Gamma \vdash e_1 \simeq e_2
                                 \Theta \vdash \widehat{\sigma}_1 \simeq \widehat{\sigma}_2
U
                         ::=
                                \Gamma; \Theta \models N \stackrel{u}{\simeq} M = \hat{\sigma}
                                                                         Negative unification
                                \Gamma;\Theta \models P \stackrel{u}{\simeq} Q = \hat{\sigma}
                                                                         Positive unification
WF
                         ::=
                                 \Gamma \vdash N
                                                                         Negative type well-formedness
                                 \Gamma \vdash P
                                                                         Positive type well-formedness
                                \Gamma \vdash N
                                                                         Negative type well-formedness
                                \Gamma \vdash P
                                                                         Positive type well-formedness
                                 \Gamma \vdash \overrightarrow{N}
                                                                         Negative type list well-formedness
                                 \Gamma \vdash \overrightarrow{P}
                                                                         Positive type list well-formedness
                                 \Gamma;\Theta \vdash N
                                                                         Negative unification type well-formedness
                                \Gamma:\Theta \vdash P
                                                                         Positive unification type well-formedness
                                 \Gamma;\Xi \vdash P
                                                                         Positive anti-unification type well-formedness
                                 \Gamma; \Xi_2 \vdash \widehat{\tau} : \Xi_1
                                                                         Antiunification substitution well-formedness
                                 \hat{\sigma}:\Theta
                                                                         Unification substitution well-formedness
                                 \Gamma \vdash^{\supseteq} \Theta
                                                                         Unification context well-formedness
                                 \Gamma_1 \vdash \sigma : \Gamma_2
                                                                         Substitution well-formedness
                                 \Gamma \vdash e
                                                                         Unification solution entry well-formedness
judgement
                                 A
                                 AU
                                 E1
                                 D1
                                 D\theta
                                 EQ
                                 LUB
                                 Nrm
                                 Order
                                 SM
                                 SImp
                                 WF
user\_syntax
                                 n
                                 n
```

varsauSolterminalsformula

$\Gamma; \Theta \models N \leqslant M \dashv \hat{\sigma}$ Negative subtyping

Positive supertyping

 $\Gamma; \Theta \models P \geqslant Q \dashv \hat{\sigma}$

$$\frac{\Gamma; \Theta \vDash \alpha^{-} \leqslant \alpha^{-} \dashv \cdot}{\Gamma; \Theta \vDash \mathbf{nf} (P) \stackrel{u}{\simeq} \mathbf{nf} (Q) \dashv \widehat{\sigma}} \quad \text{ASHIFTU}$$

$$\frac{\Gamma; \Theta \vDash \mathbf{nf} (P) \stackrel{u}{\simeq} \mathbf{nf} (Q) \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash P \leqslant \uparrow Q \dashv \widehat{\sigma}} \quad \text{ASHIFTU}$$

$$\frac{\Gamma; \Theta \vDash P \geqslant Q \dashv \widehat{\sigma}_{1} \quad \Gamma; \Theta \vDash N \leqslant M \dashv \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash P \rightarrow N \leqslant Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \quad \text{AARROW}$$

$$\frac{\Gamma; \Theta \vDash P \rightarrow N \leqslant Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}}{\Gamma; \Theta; \widehat{\sigma}^{+}; \Theta, \widehat{\alpha}^{+} \{\Gamma, \widehat{\beta}^{+}\}} \vDash [\widehat{\alpha}^{+}/\widehat{\alpha}^{+}]N \leqslant M \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \forall \widehat{\alpha}^{+}.N \leqslant \forall \widehat{\beta}^{+}.M \dashv \widehat{\sigma} \backslash \widehat{\alpha}^{+}} \quad \text{AFORALL}$$

$$\frac{\Gamma; \Theta \vDash \alpha^{+} \geqslant \alpha^{+} \dashv \cdot}{\Gamma; \Theta \vDash \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}} \quad \text{AShiftD}$$

$$\frac{\Gamma; \Theta \vDash \downarrow N \geqslant \downarrow M \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \downarrow N \geqslant \downarrow M \dashv \widehat{\sigma}} \quad \text{AShiftD}$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}}; \Theta, \widehat{\alpha}^{-}\{\Gamma, \overrightarrow{\beta^{-}}\} \vDash [\overrightarrow{\alpha^{-}}/\overrightarrow{\alpha^{-}}]P \geqslant Q \Rightarrow \widehat{\sigma}}{\Gamma; \Theta \vDash \exists \overrightarrow{\alpha^{-}}.P \geqslant \exists \overrightarrow{\beta^{-}}.Q \Rightarrow \widehat{\sigma} \setminus \widehat{\alpha^{-}}} \quad \text{AEXISTS}$$

$$\frac{\widehat{\alpha}^{+}\{\Delta\} \in \Theta \quad \mathbf{upgrade} \, \Gamma \vdash P \, \mathbf{to} \, \Delta = Q}{\Gamma: \Theta \vDash \widehat{\alpha}^{+} \geqslant P \Rightarrow (\widehat{\alpha}^{+} : \geqslant Q)} \quad \text{APUVAR}$$

 $\Gamma \models P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$

$$\frac{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} \alpha^{+} \dashv (\cdot, \alpha^{+}, \cdot, \cdot)}{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \xrightarrow{\text{AUPShift}} \frac{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash \downarrow N_{1} \stackrel{a}{\simeq} \downarrow N_{2} \dashv (\Xi, \downarrow M, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \xrightarrow{\text{AUPShift}}$$

$$\frac{\overrightarrow{\alpha^{-}} \cap \Gamma = \varnothing \quad \Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash \exists \overrightarrow{\alpha^{-}}. P_{1} \stackrel{a}{\simeq} \exists \overrightarrow{\alpha^{-}}. P_{2} \dashv (\Xi, \exists \overrightarrow{\alpha^{-}}. Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \xrightarrow{\text{AUPExists}}$$

 $\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 = (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$

$$\frac{\Gamma \vDash \alpha^{-} \stackrel{a}{\simeq} \alpha^{-} \dashv (\Xi, \alpha^{-}, \cdot, \cdot)}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \quad \text{AUNSHIFT}$$

$$\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} \uparrow P_{2} \dashv (\Xi, \uparrow Q, \widehat{\tau}_{1}, \widehat{\tau}_{2})} \quad \text{AUNSHIFT}$$

$$\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi_{1}, Q, \widehat{\tau}_{1}, \widehat{\tau}_{2}) \quad \Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi_{2}, M, \widehat{\tau}'_{1}, \widehat{\tau}'_{2})}{\Gamma \vDash P_{1} \rightarrow N_{1} \stackrel{a}{\simeq} P_{2} \rightarrow N_{2} \dashv (\Xi_{1} \cup \Xi_{2}, Q \rightarrow M, \widehat{\tau}_{1} \cup \widehat{\tau}'_{1}, \widehat{\tau}_{2} \cup \widehat{\tau}'_{2})} \quad \text{AUNARROW}$$

$$\frac{\text{if any other rule is not applicable} \quad \Gamma \vDash N \quad \Gamma \vDash M}{\Gamma \vDash N \stackrel{a}{\simeq} M \dashv (\widehat{\alpha}^{-}_{\{N,M\}}, \widehat{\alpha}^{-}_{\{N,M\}}, (\widehat{\alpha}^{-}_{\{N,M\}} : \approx N), (\widehat{\alpha}^{-}_{\{N,M\}} : \approx M))} \quad \text{AUNAU}$$

 $N \simeq_1^D M$ Negative multi-quantified type equivalence

$$\frac{\alpha^{-} \simeq_{1}^{D} \alpha^{-}}{\alpha^{-}} \quad \text{E1NVAR}$$

$$\frac{P \simeq_{1}^{D} Q}{\uparrow P \simeq_{1}^{D} \uparrow Q} \quad \text{E1SHIFTU}$$

$$\frac{P \simeq_{1}^{D} Q \quad N \simeq_{1}^{D} M}{P \to N \simeq_{1}^{D} Q \to M} \quad \text{E1ARROW}$$

$$\overrightarrow{\alpha^{+}} \cap \text{fv } M = \varnothing \quad \mu : (\overrightarrow{\beta^{+}} \cap \text{fv } M) \leftrightarrow (\overrightarrow{\alpha^{+}} \cap \text{fv } N) \quad N \simeq_{1}^{D} [\mu] M$$

$$\forall \overrightarrow{\alpha^{+}} . N \simeq_{1}^{D} \forall \overrightarrow{\beta^{+}} . M$$

$$\text{E1FORALL}$$

 $P \simeq_1^D Q$ Positive multi-quantified type equivalence

$$\frac{\alpha^{+} \simeq_{1}^{D} \alpha^{+}}{\alpha^{+}} \quad \text{E1PVAR}$$

$$\frac{N \simeq_{1}^{D} M}{\downarrow N \simeq_{1}^{D} \downarrow M} \quad \text{E1SHIFTD}$$

$$\overrightarrow{\alpha^{-}} \cap \text{fv } Q = \varnothing \quad \mu : (\overrightarrow{\beta^{-}} \cap \text{fv } Q) \leftrightarrow (\overrightarrow{\alpha^{-}} \cap \text{fv } P) \quad P \simeq_{1}^{D} [\mu] Q$$

$$\overrightarrow{\exists \alpha^{-}} . P \simeq_{1}^{D} \overrightarrow{\exists \beta^{-}} . Q$$

$$\text{E1EXISTS}$$

$$P \simeq Q$$

$$\Gamma \vdash N \sim M$$

 $\begin{array}{|c|c|c|c|c|c|}\hline P \simeq Q \\ \hline \Gamma \vdash N \simeq_1^s M \\ \hline \end{array} \quad \text{Negative equivalence on MQ types}$

$$\frac{\Gamma \vdash N \leqslant_{1} M \quad \Gamma \vdash M \leqslant_{1} N}{\Gamma \vdash N \simeq_{1}^{\leqslant} M} \quad \text{D1NDEF}$$

 $\Gamma \vdash P \simeq_1^{\epsilon} Q$ Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geqslant_1 Q \quad \Gamma \vdash Q \geqslant_1 P}{\Gamma \vdash P \simeq_1^{\circ} Q} \quad \text{D1PDEF}$$

 $\Gamma \vdash N \leqslant_1 M$ Negative subtyping

 $\Gamma \vdash P \geqslant_1 Q$ Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{1} \alpha^{+}}{\Gamma \vdash \alpha^{+} \geqslant_{1} \alpha^{+}} \quad D1PVAR$$

$$\frac{\text{>}}{\Gamma \vdash \downarrow N \geqslant_{1} \downarrow M} \quad D1SHIFTD$$

$$\frac{\text{fv } P \cap \overrightarrow{\beta^{-}} = \varnothing \quad \Gamma, \overrightarrow{\beta^{-}} \vdash N_{i} \quad \Gamma, \overrightarrow{\beta^{-}} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^{-}}]P \geqslant_{1} Q}{\Gamma \vdash \exists \overrightarrow{\alpha^{-}}.P \geqslant_{1} \exists \overrightarrow{\beta^{-}}.Q} \quad D1EXISTS$$

 $\begin{array}{|c|c|c|c|}\hline \Gamma_2 \vdash \sigma_1 \simeq_1^{\leftarrow} \sigma_2 : \Gamma_1 \\\hline \Gamma \vdash N \simeq_0^{\leftarrow} M \\\hline \end{array} \quad \begin{array}{|c|c|c|c|c|}\hline \text{Equivalence of substitutions}\\\hline \end{array}$

$$\frac{\Gamma \vdash N \leqslant_0 M \quad \Gamma \vdash M \leqslant_0 N}{\Gamma \vdash N \simeq_0^\circ M} \quad \text{D0NDef}$$

 $\Gamma \vdash P \simeq_0^{\leqslant} Q$ Positive equivalence

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash Q \geqslant_0 P}{\Gamma \vdash P \simeq_0^{\varsigma} Q} \quad \text{D0PDEF}$$

 $\Gamma \vdash N \leqslant_0 M$ Negative subtyping

$$\frac{\Gamma \vdash \alpha^- \leqslant_0 \alpha^-}{\Gamma \vdash \alpha^- \leqslant_0 \alpha^-} \quad D0NVAR$$

$$\frac{\Gamma \vdash P \simeq_0^{\varsigma} Q}{\Gamma \vdash \uparrow P \leqslant_0 \uparrow Q} \quad D0SHIFTU$$

$$\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leqslant_0 M \quad M \neq \forall \beta^+.M'}{\Gamma \vdash \forall \alpha^+.N \leqslant_0 M} \quad D0FORALLL$$

$$\frac{\Gamma, \alpha^+ \vdash N \leqslant_0 M}{\Gamma \vdash N \leqslant_0 \forall \alpha^+.M} \quad D0FORALLR$$

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash N \leqslant_0 M}{\Gamma \vdash P \to N \leqslant_0 Q \to M} \quad \text{D0Arrow}$$

 $\Gamma \vdash P \geqslant_0 Q$ Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{0} \alpha^{+}}{\Gamma \vdash N \simeq_{0}^{\leq} M} \quad D0PVAR$$

$$\frac{\Gamma \vdash N \simeq_{0}^{\leq} M}{\Gamma \vdash \downarrow N \geqslant_{0} \downarrow M} \quad D0SHIFTD$$

$$\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^{-}]P \geqslant_{0} Q \quad Q \neq \exists \alpha^{-}.Q'}{\Gamma \vdash \exists \alpha^{-}.P \geqslant_{0} Q} \quad D0EXISTSL$$

$$\frac{\Gamma, \alpha^{-} \vdash P \geqslant_{0} Q}{\Gamma \vdash P \geqslant_{0} \exists \alpha^{-}.Q} \quad D0EXISTSR$$

 $\mathbf{ord} \ vars \mathbf{in} \ P$

 $\mathbf{ord} \ vars \mathbf{in} \ N$

 $\mathbf{ord} \ vars \mathbf{in} \ P$

 $\mathbf{ord}\ vars\mathbf{in}\ N$

 $\mathbf{nf}(N')$

 $\mathbf{nf}(P')$

 $\mathbf{nf}(N')$

 $\mathbf{nf}(P')$

 $\mathbf{nf}(\overrightarrow{N}')$

 $\mathbf{nf}(\overrightarrow{P}')$

 $|\overline{\mathbf{nf}}(\sigma')|$

 $|\mathbf{nf}(\mu')|$

 $\mathbf{nf}(\widehat{\sigma}')$

 $|\sigma'|_{vars}$

 $|\hat{\sigma}'|_{vars}|$

 $e_1 \& e_2$

 $\hat{\sigma}_1 \& \hat{\sigma}_2$

 $\operatorname{\mathbf{dom}}\left(\widehat{\sigma}\right)$

 $\operatorname{\mathbf{dom}}\left(\Theta\right)$

 $\boxed{\Gamma \vDash P_1 \vee P_2 = Q} \quad \text{Least Upper Bound (Least Common Supertype)}$

$$\frac{\Gamma \models \alpha^{+} \lor \alpha^{+} = \alpha^{+}}{\Gamma \models \mathbf{nf} (\downarrow N) \stackrel{a}{\simeq} \mathbf{nf} (\downarrow M) \dashv (\Xi, P, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \cdot \models \mathbf{nf} (\downarrow N) \stackrel{a}{\simeq} \mathbf{nf} (\downarrow M) \dashv (\Xi, P, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \models \downarrow N \lor \downarrow M = \exists \alpha^{-}. [\alpha^{-}/\Xi]P} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \alpha^{-}, \beta^{-}}{\Gamma \models \exists \alpha^{-}. P_{1} \lor \exists \beta^{-}. P_{2} = Q} \quad \text{LUBEXISTS}$$

 $\boxed{\mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q}$

$$\begin{array}{ccc} \Gamma = \Delta, \overrightarrow{\alpha^{\pm}} & \overrightarrow{\beta^{\pm}} \text{ is fresh} & \overrightarrow{\gamma^{\pm}} \text{ is fresh} \\ \underline{\Delta, \overrightarrow{\beta^{\pm}}, \overrightarrow{\gamma^{\pm}} \vDash [\overrightarrow{\beta^{\pm}}/\overrightarrow{\alpha^{\pm}}]P \vee [\overrightarrow{\gamma^{\pm}}/\overrightarrow{\alpha^{\pm}}]P = Q} \\ & \textbf{upgrade} \ \Gamma \vdash P \textbf{ to} \ \Delta = Q \end{array} \quad \text{LUBUPGRADE}$$

 $\mathbf{nf}\left(N\right) = M$

$$\overline{\mathbf{nf}(\alpha^{-})} = \alpha^{-} \qquad \text{NrmNVar}$$

$$< < \mathbf{multiple parses} >> \\ \overline{\mathbf{nf}(\uparrow P)} = \uparrow Q \qquad \qquad \text{NrmShiftU}$$

$$< < \mathbf{multiple parses} >> \\ \overline{\mathbf{nf}(P \to N)} = Q \to M \qquad \qquad \text{NrmArrow}$$

$$< < \mathbf{multiple parses} >> \\ \overline{\mathbf{nf}(\forall \alpha^{+}.N)} = \forall \alpha^{+'}.N' \qquad \qquad \text{NrmForall}$$

$$\overline{\mathbf{nf}(\forall \alpha^{+}.N)} = \forall \alpha^{+'}.N'$$

 $\mathbf{nf}(P) = Q$

$$\overline{\mathbf{nf}(\alpha^{+}) = \alpha^{+}} \quad \text{NRMPVAR}$$

$$<<\mathbf{multiple parses}>> \\ \overline{\mathbf{nf}(\downarrow N) = \downarrow M} \quad \text{NRMSHIFTD}$$

$$<<\mathbf{multiple parses}>> \\ \overline{\mathbf{nf}(\exists \alpha^{-}.P) = \exists \alpha^{-\prime}.P'} \quad \text{NRMEXISTS}$$

 $\mathbf{nf}(N) = M$

$$\underline{\mathbf{nf}(\widehat{\alpha}^{-}) = \widehat{\alpha}^{-}}$$
 NRMNUVAR

 $\mathbf{nf}\left(P\right) = Q$

$$\frac{1}{\mathbf{nf}(\hat{\alpha}^+) = \hat{\alpha}^+} \quad \text{NRMPUVAR}$$

 $\mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}$

$$\frac{\alpha^{-} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \alpha^{-}} \quad \text{ONVARIN}$$

$$\frac{\alpha^{-} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \uparrow P = \overrightarrow{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}_{1} \quad \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}_{2}}{\operatorname{ord} vars \operatorname{in} P \to N = \overrightarrow{\alpha}_{1}, (\overrightarrow{\alpha}_{2} \backslash \overrightarrow{\alpha}_{1})} \quad \text{OARROW}$$

$$\frac{vars \cap \overrightarrow{\alpha^{+}} = \varnothing \quad \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} V = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

$$\frac{vars \cap \overrightarrow{\alpha^{+}} = \varnothing \quad \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} V = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

 $\mathbf{ord}\,vars\,\mathbf{in}\,P=\vec{\alpha}$

$$\frac{\alpha^{+} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \alpha^{+}} \quad \operatorname{OPVarIn}$$

$$\frac{\alpha^{+} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \cdot} \quad \operatorname{OPVarNIn}$$

$$\frac{\operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \downarrow N = \overrightarrow{\alpha}} \quad \operatorname{OSHIFTD}$$

$$\frac{vars \cap \overrightarrow{\alpha^-} = \varnothing \quad \text{ord } vars \text{in } P = \overrightarrow{\alpha}}{\text{ord } vars \text{in } \overrightarrow{\alpha^-} = \overrightarrow{\alpha}} \quad \text{OEXISTS} }$$

$$\frac{\overrightarrow{\text{ord } vars \text{in } \overrightarrow{\alpha^-} = \overrightarrow{\alpha}}}{\text{ord } vars \text{in } \overrightarrow{\alpha^+} = \overrightarrow{\alpha}} \quad \text{OFUVAR} }$$

$$\frac{\overrightarrow{\text{ord } vars \text{in } \overrightarrow{\alpha^+} = \overrightarrow{\alpha}}}{\text{ord } vars \text{in } \overrightarrow{\alpha^+} = \overrightarrow{\alpha}} \quad \text{OPUVAR} }$$

$$\frac{\Gamma \vdash e_1 \& e_2 = e_3}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P_1) \& (\overrightarrow{\alpha}^+ : \geqslant P_2) = (\overrightarrow{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPSUP} }$$

$$\frac{\Gamma \vdash P_1 \lor P_2 = Q}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P_1) \& (\overrightarrow{\alpha}^+ : \geqslant P_2) = (\overrightarrow{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPSUP} }$$

$$\frac{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P) \& (\overrightarrow{\alpha}^+ : \geqslant Q) = (\overrightarrow{\alpha}^+ : \geqslant Q)}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P) \& (\overrightarrow{\alpha}^+ : \geqslant Q) = (\overrightarrow{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPEQ} }$$

$$\frac{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P) \& (\overrightarrow{\alpha}^+ : \geqslant P) = (\overrightarrow{\alpha}^+ : \geqslant P)}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P) \& (\overrightarrow{\alpha}^+ : \geqslant P)} \quad \text{SMEPEQEQ} }$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^- : \geqslant N_1) \& (\overrightarrow{\alpha}^- : \geqslant N') = (\overrightarrow{\alpha}^- : \geqslant N)}} \quad \text{SMENEQEQ}$$

$$\frac{P \vdash P_1 \geqslant P_2}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPESUPSUP} }$$

$$\frac{\Gamma \vdash P_1 \geqslant P_2}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)}} \quad \text{SIMPEEQSUP}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \geqslant P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^- : \approx N_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)}} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \geqslant P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \approx P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{(<\text{multiple parses})}{\Gamma \vdash (\overrightarrow{\alpha}^+ : \approx P_1) \Rightarrow (\overrightarrow{\alpha}^+ : \approx P_2)} \quad \text{SIMPEPQEQ}$$

$$\frac{\Gamma; \Theta \vDash \alpha^{-\frac{u}{\simeq}} \alpha^{-} \dashv \cdot}{\Gamma; \Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}} \quad \text{UNVAR}$$

$$\frac{\Gamma; \Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \uparrow P \stackrel{u}{\simeq} \uparrow Q \dashv \widehat{\sigma}} \quad \text{USHIFTU}$$

$$\frac{\Gamma; \Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}_{1} \quad \Gamma; \Theta \vDash N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash P \rightarrow N \stackrel{u}{\simeq} Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \quad \text{UARROW}$$

$$\frac{\Gamma; \Theta \vDash P \rightarrow N \stackrel{u}{\simeq} Q \rightarrow M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash V \stackrel{u}{\simeq} N \stackrel{u}{\simeq} V \stackrel{u}{\Rightarrow} M \dashv \widehat{\sigma}} \quad \text{UFORALL}$$

$$\frac{\Gamma; \Theta \vDash V \stackrel{u}{\sim} N \stackrel{u}{\simeq} V \stackrel{u}{\Rightarrow} M \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash V \stackrel{u}{\sim} N \stackrel{u}{\simeq} V \stackrel{u}{\Rightarrow} N \dashv \widehat{\sigma}} \quad \text{UNUVAR}$$

$$\frac{\widehat{\alpha}^{-} \{\Delta\} \in \Theta \quad \Delta \vdash N}{\Gamma; \Theta \vDash \widehat{\alpha}^{-} \stackrel{u}{\simeq} N \dashv \widehat{\sigma}^{-} : \approx N} \quad \text{UNUVAR}$$

 $\Gamma; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \widehat{\sigma}$ Positive unification

$$\frac{\Gamma;\Theta \vDash \alpha^{+} \stackrel{u}{\simeq} \alpha^{+} \dashv \cdot}{\Gamma;\Theta \vDash N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}} \quad \text{UPVAR}$$

$$\frac{\Gamma;\Theta \vDash N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}}{\Gamma;\Theta \vDash N \stackrel{u}{\simeq} \downarrow M \dashv \widehat{\sigma}} \quad \text{USHIFTD}$$

$$\frac{\Gamma,\alpha^{-};\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \widehat{\sigma}}{\Gamma;\Theta \vDash \exists \alpha^{-}.P \stackrel{u}{\simeq} \exists \alpha^{-}.Q \dashv \widehat{\sigma}} \quad \text{UEXISTS}$$

$$\frac{\widehat{\alpha}^{+}\{\Delta\} \in \Theta \quad \Delta \vdash P}{\Gamma;\Theta \vDash \widehat{\alpha}^{+} \stackrel{u}{\simeq} P \dashv (\widehat{\alpha}^{+} : \approx P)} \quad \text{UPUVAR}$$

 $\overline{\Gamma \vdash N}$ Negative type well-formedness

 $\Gamma \vdash P$ Positive type well-formedness

 $\overline{\Gamma \vdash N}$ Negative type well-formedness

 $\Gamma \vdash P$ Positive type well-formedness

 $\Gamma \vdash N$ Negative type list well-formedness

 $\Gamma \vdash \overrightarrow{P}$ Positive type list well-formedness

 $\Gamma; \Theta \vdash N$ Negative unification type well-formedness

 $\Gamma; \Theta \vdash P$ Positive unification type well-formedness

 $\Gamma;\Xi\vdash P$ Positive anti-unification type well-formedness

 $\Gamma;\Xi_2 \vdash \hat{\tau}:\Xi_1$ Antiunification substitution well-formedness

 $\widehat{\sigma}: \Theta$ Unification substitution well-formedness

 $\Gamma \vdash^{\supseteq} \Theta$ Unification context well-formedness

 $\Gamma_1 \vdash \sigma : \Gamma_2$ Substitution well-formedness

 $\overline{\Gamma \vdash e}$ Unification solution entry well-formedness

Definition rules: 73 good 14 bad Definition rule clauses: 142 good 14 bad