

$\alpha, \beta, \alpha, \beta, \gamma, \delta$	type variables
$n, m, i, j$	index variables
$x, y, z$	term variables



	$ \begin{array}{l}   \cdot \\   e \\   \hat{\sigma} \backslash vars \\   \hat{\sigma}   vars \\   \hat{\sigma}_1 \cup \hat{\sigma}_2 \\   \widehat{\sigma}_i^i \\   (\hat{\sigma}) \quad \text{S} \\   \mathbf{nf}(\hat{\sigma}') \quad \text{M} \\   \hat{\sigma}'   vars \quad \text{M} \\   \hat{\sigma}_1 \ \& \ \hat{\sigma}_2 \quad \text{M} \end{array} $	concatenate
$\hat{\tau}, \hat{\rho}$	$ \begin{array}{l} ::= \\   \cdot \\   \hat{\alpha}^- : \approx N \\   \hat{\alpha}^- : \approx \boxed{N} \\   \overrightarrow{\alpha^-} / \overrightarrow{\alpha^-} \\   \overrightarrow{N} / \overrightarrow{\alpha^-} \\   \hat{\tau}_1 \cup \hat{\tau}_2 \\   \widehat{\tau}_i^i \\   (\hat{\tau}) \quad \text{S} \\   \hat{\tau}'   vars \quad \text{M} \\   \hat{\tau}_1 \ \& \ \hat{\tau}_2 \quad \text{M} \end{array} $	anti-unification substitution concatenate
$P, Q, R$	$ \begin{array}{l} ::= \\   \alpha^+ \\   \downarrow N \\   \exists \alpha^-. P \\   [\sigma] P \quad \text{M} \end{array} $	positive types
$N, M, K$	$ \begin{array}{l} ::= \\   \alpha^- \\   \uparrow P \\   \forall \alpha^+. N \\   P \rightarrow N \\   [\sigma] N \quad \text{M} \end{array} $	negative types
$\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}, \overrightarrow{\gamma^+}, \overrightarrow{\delta^+}$	$ \begin{array}{l} ::= \\   \cdot \\   \alpha^+ \\   \overrightarrow{\alpha^+} \\   \overrightarrow{\overrightarrow{\alpha^+}}^i \\   \alpha^+_i \end{array} $	positive variable list empty list a variable a variable concatenate lists
$\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}, \overrightarrow{\gamma^-}, \overrightarrow{\delta^-}$	$ \begin{array}{l} ::= \\   \cdot \\   \alpha^- \\   \overrightarrow{\alpha^-} \\   \overrightarrow{\overrightarrow{\alpha^-}}^i \\   \alpha^-_i \end{array} $	negative variables empty list a variable variables concatenate lists
$\overrightarrow{\alpha^\pm}, \overrightarrow{\beta^\pm}, \overrightarrow{\gamma^\pm}, \overrightarrow{\delta^\pm}$	$ ::= $	positive or negative variable list

		$\cdot$	empty list
		$\alpha^\pm$	a variable
		$\vec{\mathbf{p}}\mathbf{a}$	variables
		$\overrightarrow{\alpha^\pm}_i$	concatenate lists
$P, Q$	$::=$		multi-quantified positive types
		$\alpha^+$	
		$\downarrow N$	
		$\exists \alpha^-. P$	$P \neq \exists \dots$
		$[\sigma]P$	M
		$[\hat{\tau}]P$	M
		$[\hat{\sigma}]P$	M
		$[\mu]P$	M
		$(P)$	S
		$P_1 \vee P_2$	M
		$\mathbf{nf}(P')$	M
$N, M$	$::=$		multi-quantified negative types
		$\alpha^-$	
		$\uparrow P$	
		$P \rightarrow N$	
		$\forall \alpha^+. N$	$N \neq \forall \dots$
		$[\sigma]N$	M
		$[\hat{\tau}]N$	M
		$[\mu]N$	M
		$[\hat{\sigma}]N$	M
		$(N)$	S
		$\mathbf{nf}(N')$	M
$\vec{P}, \vec{Q}$	$::=$		list of positive types
		$\cdot$	empty list
		$P$	a singel type
		$\overrightarrow{P}_i$	concatenate lists
		$\mathbf{nf}(\vec{P}')$	M
$\vec{N}, \vec{M}$	$::=$		list of negative types
		$\cdot$	empty list
		$N$	a singel type
		$\overrightarrow{N}_i$	concatenate lists
		$\mathbf{nf}(\vec{N}')$	M
$\Delta, \Gamma$	$::=$		declarative type context
		$\cdot$	empty context
		$\overrightarrow{\alpha^+}$	list of variables
		$\overrightarrow{\alpha^-}$	list of variables
		$\overrightarrow{\alpha^\pm}$	list of variables
		$vars$	
		$\overrightarrow{\Gamma}_i$	concatenate contexts
		$(\Gamma)$	S
		$\Theta(\hat{\alpha}^+)$	M

		$\Theta(\hat{\alpha}^-)$	M	
$\Theta$	::=			unification type variable context
		.		empty context
		$\alpha^+$		list of variables
		$\alpha^-$		list of variables
		$vars$		
		$\overline{\Theta}_i^i$		concatenate contexts
		$(\Theta)$	S	
		$\Theta _{vars}$		leave only those variables that are in the set
		$\Theta_1 \cup \Theta_2$		
$\Xi$	::=			anti-unification type variable context
		.		empty context
		$\alpha^-$		list of variables
		$\Xi_i^i$		concatenate contexts
		$(\Xi)$	S	
		$\Xi_1 \cup \Xi_2$		
		$\Xi_1 \cap \Xi_2$		
		$\Xi' _{vars}$	M	
$\vec{\alpha}, \vec{\beta}$	::=			ordered positive or negative variables
		.		empty list
		$\alpha^+$		list of variables
		$\alpha^-$		list of variables
		$\alpha^\pm$		list of variables
		$\alpha^+$		list of variables
		$\alpha^-$		list of variables
		$\vec{\alpha}_1 \setminus vars$		setminus
		$\Gamma$		context
		$vars$		
		$\overline{\vec{\alpha}}_i^i$		concatenate contexts
		$(\vec{\alpha})$	S	parenthesis
		$[\mu]\vec{\alpha}$	M	apply moving to list
		<b>ord vars in</b> $P$	M	
		<b>ord vars in</b> $N$	M	
		<b>ord vars in</b> $P$	M	
		<b>ord vars in</b> $N$	M	
$vars$	::=			set of variables
		$\emptyset$		empty set
		<b>fv</b> $P$		free variables
		<b>fv</b> $N$		free variables
		<b>fv imP</b>		free variables
		<b>fv imN</b>		free variables
		$vars_1 \cap vars_2$		set intersection
		$vars_1 \cup vars_2$		set union
		$vars_1 \setminus vars_2$		set complement
		<b>mv imP</b>		movable variables
		<b>mv imN</b>		movable variables

		<b>uv</b> $N$	unification variables
		<b>uv</b> $P$	unification variables
		<b>fv</b> $N$	free variables
		<b>fv</b> $P$	free variables
		$(vars)$	S parenthesis
		$\vec{\alpha}$	ordered list of variables
		$[\mu]vars$	M apply moving to varset
		<b>dom</b> $(\hat{\sigma})$	M
		<b>dom</b> $(\hat{\tau})$	M
		<b>dom</b> $(\Theta)$	M
$\mu$	::=		
		.	empty moving
		$pma1 \mapsto pma2$	Positive unit substitution
		$nma1 \mapsto nma2$	Positive unit substitution
		$\mu_1 \cup \mu_2$	M Set-like union of movings
		$\mu_1 \circ \mu_2$	M Composition
		$\overline{\mu_i}^i$	concatenate movings
		$\mu vars$	M restriction on a set
		$\mu^{-1}$	M inversion
		<b>nf</b> $(\mu')$	M
$\hat{\alpha}^\pm$	::=		positive/negative unification variable
		$\hat{\alpha}^\pm$	
$\hat{\alpha}^+$	::=		positive unification variable
		$\hat{\alpha}^+$	
		$\hat{\alpha}^+\{\Delta\}$	
		$\hat{\alpha}^\pm$	
$\hat{\alpha}^-, \hat{\beta}^-$	::=		negative unification variable
		$\hat{\alpha}^-$	
		$\hat{\alpha}_{\{N,M\}}^-$	
		$\hat{\alpha}^-\{\Delta\}$	
		$\hat{\alpha}^\pm$	
$\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}$	::=		positive unification variable list
		.	empty list
		$\hat{\alpha}^+$	a variable
		$\overrightarrow{\hat{\alpha}^+\{\Delta\}}$	from a normal variable
		$\overrightarrow{\hat{\alpha}^+}$	from a normal variable, context unspecified
		$\overrightarrow{\overrightarrow{\alpha^+}}^i$	concatenate lists
		$\alpha^+_i$	
$\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}$	::=		negative unification variable list
		.	empty list
		$\hat{\alpha}^-$	a variable
		$\Xi$	from an antiunification context
		$\overrightarrow{\hat{\alpha}^-\{\Delta\}}$	from a normal variable
		$\overrightarrow{\hat{\alpha}^-}$	from a normal variable, context unspecified

	$\xrightarrow{i}$ $\alpha^-_i$	concatenate lists
$P, Q$	$::=$ $\alpha^+$ $\mathbf{pma}$ $\hat{\alpha}^+$ $\downarrow N$ $\xrightarrow{\quad} \exists \alpha^- . P$ $[\sigma] P$ $[\hat{\tau}] P$ $[\mu] P$ $(P)$ $\mathbf{nf}(P')$	a positive algorithmic type (potentially with metavariables)       M M M S M
$N, M$	$::=$ $\alpha^-$ $\hat{\alpha}^-$ $\uparrow P$ $P \rightarrow N$ $\xrightarrow{\quad} \forall \alpha^+ . N$ $[\sigma] N$ $[\hat{\tau}] N$ $[\mu] N$ $(N)$ $\mathbf{nf}(N')$	a negative algorithmic type (potentially with metavariables)       M M M S M
$auSol$	$::=$ $(\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ $(\Xi, N, \hat{\tau}_1, \hat{\tau}_2)$	
$terminals$	$::=$ $\exists$ $\forall$ $\uparrow$ $\downarrow$ $\mapsto$ $\leftrightarrow$ $\in$ $\notin$ $\cdot$ $\top$ $\leq$ $\geq$ $\approx$ $\subset$ $\supset$ $\setminus$ $\sqcup$ $\mapsto$ $\approx^u$	

	$\sim$ $\emptyset$ $\circ$ $\Rightarrow$ $\models$ $\models$ $\neq$ $\equiv_n$ $\vee$ $\Downarrow$ $:\geq$ $:\simeq$ $\Lambda$ $\lambda$ $\mathbf{let}^\exists$ $\bullet$ $\Rightarrow\Rightarrow$	
$v, w$	$::=$ $ $ $x$ $ $ $\{c\}$ $ $ $(v : P)$ $ $ $(v)$	value terms      M
$\vec{v}$	$::=$ $ $ $\cdot$ $ $ $v$ $ $ $\overrightarrow{v}_i^i$	list of arguments   concatenate
$c, d$	$::=$ $ $ $\lambda x : P. c$ $ $ $\Lambda \alpha^+. c$ $ $ $\mathbf{return} v$ $ $ $\mathbf{let} x : P = v(\vec{v}); c$ $ $ $\mathbf{let} x = v(\vec{v}); c$ $ $ $\mathbf{let}^\exists(\alpha^-, x) = v; c$	computation terms
$vctx, \Upsilon$	$::=$ $ $ $\cdot$ $ $ $x : P$ $ $ $\overline{\Upsilon}_i^i$	variable context   concatenate contexts
<i>formula</i>	$::=$ $ $ $judgement$ $ $ $judgement \text{ uniquely}$ $ $ $formula_1 \dots formula_n$ $ $ $\mu : vars_1 \leftrightarrow vars_2$ $ $ $\mu \text{ is bijective}$ $ $ $\hat{\sigma} \text{ is functional}$ $ $ $\hat{\sigma}_1 \in \hat{\sigma}_2$	



	$v : P \in \Upsilon$ $\hat{\sigma}_1 \subseteq \hat{\sigma}_2$ $vars_1 \subseteq vars_2$ $vars_1 = vars_2$ <b><math>vars</math> is fresh</b> $\alpha^- \notin vars$ $\alpha^+ \notin vars$ $\alpha^- \in vars$ $\alpha^+ \in vars$ $\hat{\alpha}^+ \in vars$ $\hat{\alpha}^- \in vars$ $\hat{\alpha}^- \in \Theta$ $\hat{\alpha}^+ \in \Theta$ if any other rule is not applicable $\vec{\alpha}_1 = \vec{\alpha}_2$ $e_1 = e_2$ $\boxed{N} = \boxed{M}$ $N \neq M$ $P \neq Q$ $N \neq M$ $P \neq Q$	
$A$	$::=$ $\Gamma; \Theta \vdash \boxed{N} \leqslant M \Rightarrow \hat{\sigma}$ $\Gamma; \Theta \vdash \boxed{P} \geqslant Q \Rightarrow \hat{\sigma}$	Negative subtyping Positive supertyping
$AU$	$::=$ $\Gamma \vdash P_1 \stackrel{a}{\simeq} P_2 \Rightarrow (\Xi, \boxed{Q}, \hat{\tau}_1, \hat{\tau}_2)$ $\Gamma \vdash N_1 \stackrel{a}{\simeq} N_2 \Rightarrow (\Xi, \boxed{M}, \hat{\tau}_1, \hat{\tau}_2)$	
$E1$	$::=$ $\boxed{N} \stackrel{\textcolor{brown}{D}}{\simeq}_1 M$ $\boxed{P} \stackrel{\textcolor{brown}{D}}{\simeq}_1 Q$ $\boxed{P} \simeq \boxed{Q}$	Negative multi-quantified type equivalence Positive multi-quantified type equivalence
$D1$	$::=$ $\Gamma \vdash N \simeq_1^{\leqslant} M$ $\Gamma \vdash P \simeq_1^{\leqslant} Q$ $\Gamma \vdash N \leqslant_1 M$ $\Gamma \vdash P \geqslant_1 Q$ $\Gamma_2 \vdash \sigma_1 \simeq_1^{\leqslant} \sigma_2 : \Gamma_1$	Negative equivalence on MQ types Positive equivalence on MQ types Negative subtyping Positive supertyping Equivalence of substitutions
$D0$	$::=$ $\Gamma \vdash N \simeq_0^{\leqslant} M$ $\Gamma \vdash P \simeq_0^{\leqslant} Q$ $\Gamma \vdash N \leqslant_0 M$ $\Gamma \vdash P \geqslant_0 Q$	Negative equivalence Positive equivalence Negative subtyping Positive supertyping
$DT$	$::=$	

	$\begin{array}{ l} \Gamma; \Upsilon \vdash v : P \\ \Gamma; \Upsilon \vdash c : N \\ \Gamma; \Upsilon \vdash N \bullet \vec{v} \Rightarrow M \end{array}$	Positive type inference Negative type inference Spin Application type inference
$EQ$	$\begin{array}{ l} N = M \\ P = Q \\ \boxed{P} = \boxed{Q} \end{array}$	Negative type equality (alpha-equivalence) Positive type equality (alpha-equivalence)
$LUBF$	$\begin{array}{ l} P_1 \vee P_2 === Q \\ \text{ord vars in } \boxed{P} === \vec{\alpha} \\ \text{ord vars in } \boxed{N} === \vec{\alpha} \\ \text{ord vars in } P === \vec{\alpha} \\ \text{ord vars in } N === \vec{\alpha} \\ \text{nf}(N') === N \\ \text{nf}(P') === P \\ \text{nf}(\boxed{N'}) === \boxed{N} \\ \text{nf}(\boxed{P'}) === \boxed{P} \\ \text{nf}(\vec{N}') === \vec{N} \\ \text{nf}(\vec{P}') === \vec{P} \\ \text{nf}(\sigma') === \sigma \\ \text{nf}(\mu') === \mu \\ \text{nf}(\hat{\sigma}') === \hat{\sigma} \\ \sigma' _{vars} \\ \hat{\sigma}' _{vars} \\ \hat{\tau}' _{vars} \\ \Xi' _{vars} \\ e_1 \ \& \ e_2 \\ \hat{\sigma}_1 \ \& \ \hat{\sigma}_2 \\ \hat{\tau}_1 \ \& \ \hat{\tau}_2 \\ \text{dom}(\hat{\sigma}) === vars \\ \text{dom}(\hat{\tau}) === vars \\ \text{dom}(\Theta) === vars \end{array}$	
$LUB$	$\begin{array}{ l} \Gamma \models P_1 \vee P_2 = Q \\ \text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q \end{array}$	Least Upper Bound (Least Common Supertype)
$Nrm$	$\begin{array}{ l} \text{nf}(N) = M \\ \text{nf}(P) = Q \\ \text{nf}(N) = \boxed{M} \\ \text{nf}(P) = \boxed{Q} \end{array}$	
$Order$	$\begin{array}{ l} \text{ord vars in } N = \vec{\alpha} \\ \text{ord vars in } P = \vec{\alpha} \\ \text{ord vars in } \boxed{N} = \vec{\alpha} \\ \text{ord vars in } \boxed{P} = \vec{\alpha} \end{array}$	

$SM$	$::=$ $\mid$ $\Gamma \vdash e_1 \ \& \ e_2 = e_3$ $\mid$ $\Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3$	Unification Solution Entry Merge Merge unification solutions
$SImp$	$::=$ $\mid$ $\Gamma \vdash e_1 \Rightarrow e_2$ $\mid$ $\Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2$ $\mid$ $\Gamma \vdash e_1 \simeq e_2$ $\mid$ $\Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2$	Weakening of unification solution entries Weakening of unification solutions
$U$	$::=$ $\mid$ $\Gamma; \Theta \models N \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma}$ $\mid$ $\Gamma; \Theta \models P \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma}$	Negative unification Positive unification
$WF$	$::=$ $\mid$ $\Gamma \vdash N$ $\mid$ $\Gamma \vdash P$ $\mid$ $\Gamma \vdash N$ $\mid$ $\Gamma \vdash P$ $\mid$ $\Gamma \vdash \vec{N}$ $\mid$ $\Gamma \vdash \vec{P}$ $\mid$ $\Gamma; \Theta \vdash N$ $\mid$ $\Gamma; \Theta \vdash P$ $\mid$ $\Gamma; \Xi \vdash N$ $\mid$ $\Gamma; \Xi \vdash P$ $\mid$ $\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1$ $\mid$ $\hat{\sigma} : \Theta$ $\mid$ $\Gamma \vdash^{\supset} \Theta$ $\mid$ $\Gamma_1 \vdash \sigma : \Gamma_2$ $\mid$ $\Gamma \vdash e$	Negative type well-formedness Positive type well-formedness Negative type well-formedness Positive type well-formedness Negative type list well-formedness Positive type list well-formedness Negative unification type well-formedness Positive unification type well-formedness Negative anti-unification type well-formedness Positive anti-unification type well-formedness Antiunification substitution well-formedness Unification substitution well-formedness Unification context well-formedness Substitution well-formedness Unification solution entry well-formedness
$judgement$	$::=$ $\mid$ $A$ $\mid$ $AU$ $\mid$ $E1$ $\mid$ $D1$ $\mid$ $D0$ $\mid$ $DT$ $\mid$ $EQ$ $\mid$ $LUB$ $\mid$ $Nrm$ $\mid$ $Order$ $\mid$ $SM$ $\mid$ $SImp$ $\mid$ $U$ $\mid$ $WF$	
$user\_syntax$	$::=$ $\mid$ $\alpha$ $\mid$ $n$	

	$x$
	$n$
	$\alpha^+$
	$\alpha^-$
	$\alpha^\pm$
	$\sigma$
	$e$
	$\hat{\sigma}$
	$\hat{\tau}$
	$P$
	$N$
	$\overrightarrow{\alpha^+}$
	$\overrightarrow{\alpha^-}$
	$\overrightarrow{\alpha^\pm}$
	$P$
	$N$
	$\overrightarrow{P}$
	$\overrightarrow{N}$
	$\Gamma$
	$\Theta$
	$\Xi$
	$\vec{\alpha}$
	$vars$
	$\mu$
	$\hat{\alpha}^\pm$
	$\hat{\alpha}^+$
	$\hat{\alpha}^-$
	$\widetilde{\alpha^+}$
	$\widetilde{\alpha^-}$
	$\alpha^+$
	$\alpha^-$
	$\overline{P}$
	$\overline{N}$
	$auSol$
	$terminals$
	$v$
	$\vec{v}$
	$c$
	$vctx$
	$formula$

$\boxed{\Gamma; \Theta \models N \leq M \Rightarrow \hat{\sigma}}$

Negative subtyping

$$\begin{array}{c}
\overline{\Gamma; \Theta \models \alpha^- \leq \alpha^- \Rightarrow} \quad \text{ANVAR} \\
\frac{\Gamma; \Theta \models \mathbf{nf}(P) \stackrel{u}{\simeq} \mathbf{nf}(Q) \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \uparrow P \leq \uparrow Q \Rightarrow \hat{\sigma}} \quad \text{AShiftU} \\
\frac{\Gamma; \Theta \models P \geq Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \Theta \models N \leq M \Rightarrow \hat{\sigma}_2}{\Gamma; \Theta \models P \rightarrow N \leq Q \rightarrow M \Rightarrow \hat{\sigma}_1 \ \& \ \hat{\sigma}_2} \quad \text{AArrow} \\
\frac{\Gamma, \vec{\beta}^+; \Theta, \hat{\alpha}^+ \{\Gamma, \vec{\beta}^+\} \models [\hat{\alpha}^+ / \alpha^+] N \leq M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \forall \vec{\alpha}^+. N \leq \forall \vec{\beta}^+. M \Rightarrow \hat{\sigma} \setminus \hat{\alpha}^+} \quad \text{AForall}
\end{array}$$

$\boxed{\Gamma; \Theta \models P \succcurlyeq Q \models \hat{\sigma}}$  Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma; \Theta \models \alpha^+ \succcurlyeq \alpha^+ \models \cdot} \text{APVAR} \\
\frac{\Gamma; \Theta \models \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \models \hat{\sigma}}{\Gamma; \Theta \models \downarrow N \succcurlyeq \downarrow M \models \hat{\sigma}} \text{ASHIFTD} \\
\frac{\Gamma, \vec{\beta}^-; \Theta, \hat{\alpha}^- \{ \Gamma, \vec{\beta}^- \} \models [\hat{\alpha}^- / \alpha^-] P \succcurlyeq Q \models \hat{\sigma}}{\Gamma; \Theta \models \exists \vec{\alpha}^- . P \succcurlyeq \exists \vec{\beta}^- . Q \models \hat{\sigma} \setminus \hat{\alpha}^-} \text{AEXISTS} \\
\frac{\hat{\alpha}^+ \{ \Delta \} \in \Theta \quad \mathbf{upgrade} \Gamma \vdash P \text{ to } \Delta = Q}{\Gamma; \Theta \models \hat{\alpha}^+ \succcurlyeq P \models (\hat{\alpha}^+ : \geq Q)} \text{APUVAR}
\end{array}$$

$$\boxed{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}$$

$$\begin{array}{c}
\frac{}{\Gamma \models \alpha^+ \stackrel{a}{\simeq} \alpha^+ \models (\cdot, \alpha^+, \cdot, \cdot)} \text{AUPVAR} \\
\frac{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N_1 \stackrel{a}{\simeq} \downarrow N_2 \models (\Xi, \downarrow M, \hat{\tau}_1, \hat{\tau}_2)} \text{AUSHIFTD} \\
\frac{\vec{\alpha}^- \cap \Gamma = \emptyset \quad \Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \exists \vec{\alpha}^- . P_1 \stackrel{a}{\simeq} \exists \vec{\alpha}^- . P_2 \models (\Xi, \exists \vec{\alpha}^- . Q, \hat{\tau}_1, \hat{\tau}_2)} \text{AUEXISTS}
\end{array}$$

$$\boxed{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}$$

$$\begin{array}{c}
\frac{}{\Gamma \models \alpha^- \stackrel{a}{\simeq} \alpha^- \models (\cdot, \alpha^-, \cdot, \cdot)} \text{AUNVAR} \\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \uparrow P_1 \stackrel{a}{\simeq} \uparrow P_2 \models (\Xi, \uparrow Q, \hat{\tau}_1, \hat{\tau}_2)} \text{AUSHIFTU} \\
\frac{\vec{\alpha}^+ \cap \Gamma = \emptyset \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \forall \vec{\alpha}^+ . N_1 \stackrel{a}{\simeq} \forall \vec{\alpha}^+ . N_2 \models (\Xi, \forall \vec{\alpha}^+ . M, \hat{\tau}_1, \hat{\tau}_2)} \text{AUFORALL} \\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi_1, Q, \hat{\tau}_1, \hat{\tau}_2) \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi_2, M, \hat{\tau}'_1, \hat{\tau}'_2)}{\Gamma \models P_1 \rightarrow N_1 \stackrel{a}{\simeq} P_2 \rightarrow N_2 \models (\Xi_1 \cup \Xi_2, Q \rightarrow M, \hat{\tau}_1 \cup \hat{\tau}'_1, \hat{\tau}_2 \cup \hat{\tau}'_2)} \text{AUARROW} \\
\frac{\text{if any other rule is not applicable} \quad \Gamma \vdash N \quad \Gamma \vdash M}{\Gamma \models N \stackrel{a}{\simeq} M \models (\hat{\alpha}^-_{\{N, M\}}, \hat{\alpha}^-_{\{N, M\}}, (\hat{\alpha}^-_{\{N, M\}} : \approx N), (\hat{\alpha}^-_{\{N, M\}} : \approx M))} \text{AUAU}
\end{array}$$

$\boxed{N \simeq_1^D M}$  Negative multi-quantified type equivalence

$$\begin{array}{c}
\frac{}{\alpha^- \simeq_1^D \alpha^-} \text{E1NVAR} \\
\frac{P \simeq_1^D Q}{\uparrow P \simeq_1^D \uparrow Q} \text{E1SHIFTU} \\
\frac{P \simeq_1^D Q \quad N \simeq_1^D M}{P \rightarrow N \simeq_1^D Q \rightarrow M} \text{E1ARROW} \\
\frac{\vec{\alpha}^+ \cap \mathbf{fv} M = \emptyset \quad \mu : (\vec{\beta}^+ \cap \mathbf{fv} M) \leftrightarrow (\vec{\alpha}^+ \cap \mathbf{fv} N) \quad N \simeq_1^D [\mu] M}{\forall \vec{\alpha}^+ . N \simeq_1^D \forall \vec{\beta}^+ . M} \text{E1FORALL}
\end{array}$$

$\boxed{P \simeq_1^D Q}$  Positive multi-quantified type equivalence

$$\begin{array}{c}
\frac{}{\alpha^+ \simeq_1^D \alpha^+} \text{E1PVAR} \\
\frac{N \simeq_1^D M}{\downarrow N \simeq_1^D \downarrow M} \text{E1SHIFTD} \\
\frac{\overrightarrow{\alpha^-} \cap \mathbf{fv} Q = \emptyset \quad \mu : (\overrightarrow{\beta^-} \cap \mathbf{fv} Q) \leftrightarrow (\overrightarrow{\alpha^-} \cap \mathbf{fv} P) \quad P \simeq_1^D [\mu]Q}{\exists \overrightarrow{\alpha^-}. P \simeq_1^D \exists \overrightarrow{\beta^-}. Q} \text{E1EXISTS}
\end{array}$$

$\boxed{P \simeq Q}$   
 $\boxed{\Gamma \vdash N \simeq_1^{\leq} M}$  Negative equivalence on MQ types

$$\frac{\Gamma \vdash N \leq_1 M \quad \Gamma \vdash M \leq_1 N}{\Gamma \vdash N \simeq_1^{\leq} M} \text{D1NDEF}$$

$\boxed{\Gamma \vdash P \simeq_1^{\leq} Q}$  Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash Q \geq_1 P}{\Gamma \vdash P \simeq_1^{\leq} Q} \text{D1PDEF}$$

$\boxed{\Gamma \vdash N \leq_1 M}$  Negative subtyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^- \leq_1 \alpha^-} \text{D1NVAR} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \uparrow P \leq_1 \uparrow Q} \text{D1SHIFTU} \\
\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash N \leq_1 M}{\Gamma \vdash P \rightarrow N \leq_1 Q \rightarrow M} \text{D1ARROW} \\
\frac{\mathbf{fv} N \cap \overrightarrow{\beta^+} = \emptyset \quad \Gamma, \overrightarrow{\beta^+} \vdash P_i \quad \Gamma, \overrightarrow{\beta^+} \vdash [\overrightarrow{P}/\overrightarrow{\alpha^+}]N \leq_1 M}{\Gamma \vdash \forall \overrightarrow{\alpha^+}. N \leq_1 \forall \overrightarrow{\beta^+}. M} \text{D1FORALL}
\end{array}$$

$\boxed{\Gamma \vdash P \geq_1 Q}$  Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^+ \geq_1 \alpha^+} \text{D1PVAR} \\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \downarrow N \geq_1 \downarrow M} \text{D1SHIFTD} \\
\frac{\mathbf{fv} P \cap \overrightarrow{\beta^-} = \emptyset \quad \Gamma, \overrightarrow{\beta^-} \vdash N_i \quad \Gamma, \overrightarrow{\beta^-} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^-}]P \geq_1 Q}{\Gamma \vdash \exists \overrightarrow{\alpha^-}. P \geq_1 \exists \overrightarrow{\beta^-}. Q} \text{D1EXISTS}
\end{array}$$

$\boxed{\Gamma_2 \vdash \sigma_1 \simeq_1^{\leq} \sigma_2 : \Gamma_1}$  Equivalence of substitutions  
 $\boxed{\Gamma \vdash N \simeq_0^{\leq} M}$  Negative equivalence

$$\frac{\Gamma \vdash N \leq_0 M \quad \Gamma \vdash M \leq_0 N}{\Gamma \vdash N \simeq_0^{\leq} M} \text{D0NDEF}$$

$\boxed{\Gamma \vdash P \simeq_0^{\leq} Q}$  Positive equivalence

$$\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash Q \geq_0 P}{\Gamma \vdash P \simeq_0^{\leq} Q} \text{D0PDEF}$$

$\boxed{\Gamma \vdash N \leq_0 M}$  Negative subtyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^- \leq_0 \alpha^-} \text{D0NVAR} \\
\frac{\Gamma \vdash P \simeq_0^{\leq} Q}{\Gamma \vdash \uparrow P \leq_0 \uparrow Q} \text{D0SHIFTU} \\
\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leq_0 M \quad M \neq \forall \beta^+. M'}{\Gamma \vdash \forall \alpha^+. N \leq_0 M} \text{D0FORALLL} \\
\frac{\Gamma, \alpha^+ \vdash N \leq_0 M}{\Gamma \vdash N \leq_0 \forall \alpha^+. M} \text{D0FORALLR} \\
\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash N \leq_0 M}{\Gamma \vdash P \rightarrow N \leq_0 Q \rightarrow M} \text{D0ARROW}
\end{array}$$

$\boxed{\Gamma \vdash P \geq_0 Q}$  Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^+ \geq_0 \alpha^+} \text{D0PVAR} \\
\frac{\Gamma \vdash N \simeq_0^{\leq} M}{\Gamma \vdash \downarrow N \geq_0 \downarrow M} \text{D0SHIFTD} \\
\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^-]P \geq_0 Q \quad Q \neq \exists \alpha^-. Q'}{\Gamma \vdash \exists \alpha^-. P \geq_0 Q} \text{D0EXISTSL} \\
\frac{\Gamma, \alpha^- \vdash P \geq_0 Q}{\Gamma \vdash P \geq_0 \exists \alpha^-. Q} \text{D0EXISTSR}
\end{array}$$

$\boxed{\Gamma; \Upsilon \vdash v : P}$  Positive type inference

$$\begin{array}{c}
\frac{v : P \in \Upsilon}{\Gamma; \Upsilon \vdash v : P} \text{DTVAR} \\
\frac{\Gamma; \Upsilon \vdash c : N}{\Gamma; \Upsilon \vdash \{c\} : \downarrow N} \text{DTTHUNK} \\
\frac{\Gamma; \Upsilon \vdash v : P \quad \Gamma \vdash Q \geq_1 P}{\Gamma; \Upsilon \vdash (v : Q) : Q} \text{DTANNOT}
\end{array}$$

$\boxed{\Gamma; \Upsilon \vdash c : N}$  Negative type inference

$$\begin{array}{c}
\frac{\Gamma; \Upsilon, x : P \vdash c : N}{\Gamma; \Upsilon \vdash \lambda x : P. c : P \rightarrow N} \text{DTTLAM} \\
\frac{\Gamma, \alpha^+; \Upsilon \vdash c : N}{\Gamma; \Upsilon \vdash \Lambda \alpha^+. c : \forall \alpha^+. N} \text{DTTLAM} \\
\frac{\Gamma; \Upsilon \vdash v : P}{\Gamma; \Upsilon \vdash \mathbf{return} v : \uparrow P} \text{DTRETURN} \\
\frac{\Gamma; \Upsilon \vdash v : \downarrow M \quad \Gamma; \Upsilon \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \quad \Gamma \vdash \uparrow Q \leq_1 \uparrow P \quad \Gamma; \Upsilon, x : P \vdash c : N}{\Gamma; \Upsilon \vdash \mathbf{let} x : P = v(\vec{v}); c : N} \text{DTLETANN} \\
\frac{\Gamma; \Upsilon \vdash v : \downarrow M \quad \Gamma; \Upsilon \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \text{ uniquely} \quad \Gamma; \Upsilon, x : Q \vdash c : N}{\Gamma; \Upsilon \vdash \mathbf{let} x = v(\vec{v}); c : N} \text{DTLET} \\
\frac{\Gamma, \alpha^-; \Upsilon \vdash v : \exists \alpha^-. P \quad \Gamma, \alpha^-; \Upsilon, x : P \vdash c : N \quad \Gamma \vdash N}{\Gamma; \Upsilon \vdash \mathbf{let}^{\exists}(\alpha^-, x) = v; c : N} \text{DTUNPACK}
\end{array}$$

$\boxed{\Gamma; \Upsilon \vdash N \bullet \vec{v} \Rightarrow M}$  Spin Application type inference

$$\frac{N \neq \forall \alpha^+. M}{\Gamma; \Upsilon \vdash N \bullet \cdot \Rightarrow N} \text{ DTEMTPTY}$$

$$\frac{\Gamma; \Upsilon \vdash v : P \quad \Gamma \vdash Q \geqslant_1 P \quad \Gamma; \Upsilon \vdash N \bullet \vec{v} \Rightarrow M}{\Gamma; \Upsilon \vdash Q \rightarrow N \bullet v, \vec{v} \Rightarrow M} \text{ DTARROW}$$

$$\frac{\Gamma \vdash \vec{P} \quad \Gamma; \Upsilon \vdash [\vec{P}/\alpha^+] N \bullet \vec{v} \Rightarrow M}{\Gamma; \Upsilon \vdash \forall \alpha^+. N \bullet \vec{v} \Rightarrow M} \text{ DTFORALL}$$

$\boxed{N = M}$  Negative type equality (alpha-equivalence)

$\boxed{P = Q}$  Positive type equality (alphha-equivalence)

$\boxed{P = Q}$

$\boxed{P_1 \vee P_2}$

$\boxed{\text{ord vars in } P}$

$\boxed{\text{ord vars in } N}$

$\boxed{\text{ord vars in } P}$

$\boxed{\text{ord vars in } N}$

$\boxed{\text{nf } (N')}$

$\boxed{\text{nf } (P')}$

$\boxed{\text{nf } (N')}$

$\boxed{\text{nf } (P')}$

$\boxed{\text{nf } (\vec{N}')}$

$\boxed{\text{nf } (\vec{P}')}$

$\boxed{\text{nf } (\sigma')}$



$\mathbf{nf}(\mu')$  $\mathbf{nf}(\hat{\sigma}')$  $\sigma'|_{vars}$  $\hat{\sigma}'|_{vars}$  $\hat{\tau}'|_{vars}$  $\Xi'|_{vars}$  $e_1 \ \& \ e_2$  $\hat{\sigma}_1 \ \& \ \hat{\sigma}_2$  $\hat{\tau}_1 \ \& \ \hat{\tau}_2$  $\mathbf{dom}(\hat{\sigma})$  $\mathbf{dom}(\hat{\tau})$  $\mathbf{dom}(\Theta)$  $\Gamma \models P_1 \vee P_2 = Q$     Least Upper Bound (Least Common Supertype)

$$\begin{array}{c}
\overline{\Gamma \models \alpha^+ \vee \alpha^+ = \alpha^+} \quad \text{LUBVAR} \\
\frac{\Gamma, \cdot \models \mathbf{nf}(\downarrow N) \stackrel{a}{\simeq} \mathbf{nf}(\downarrow M) \Rightarrow (\Xi, P, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N \vee \downarrow M = \overrightarrow{\exists \alpha^-} . [\overrightarrow{\alpha^-} / \Xi] P} \quad \text{LUBSHIFT}
\end{array}$$

$$\frac{\Gamma, \vec{\alpha}^-, \vec{\beta}^- \vdash P_1 \vee P_2 = Q}{\Gamma \vdash \exists \vec{\alpha}^-. P_1 \vee \exists \vec{\beta}^-. P_2 = Q} \text{ LUBEXISTS}$$

$$\boxed{\text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q}$$

$$\frac{\begin{array}{l} \Gamma = \Delta, \vec{\alpha}^\pm \quad \vec{\beta}^\pm \text{ is fresh} \quad \vec{\gamma}^\pm \text{ is fresh} \\ \Delta, \vec{\beta}^\pm, \vec{\gamma}^\pm \vdash [\vec{\beta}^\pm / \vec{\alpha}^\pm] P \vee [\vec{\gamma}^\pm / \vec{\alpha}^\pm] P = Q \end{array}}{\text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q} \text{ LUBUPGRADE}$$

$$\boxed{\text{nf}(N) = M}$$

$$\overline{\text{nf}(\alpha^-) = \alpha^-} \quad \text{NRMNVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\uparrow P) = \uparrow Q} \quad \text{NRMSHIFTU}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(P \rightarrow N) = Q \rightarrow M} \quad \text{NRMARROW}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\forall \vec{\alpha}^+. N) = \forall \vec{\alpha}^+. N'} \quad \text{NRMFORALL}$$

$$\boxed{\text{nf}(P) = Q}$$

$$\overline{\text{nf}(\alpha^+) = \alpha^+} \quad \text{NRMPVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\downarrow N) = \downarrow M} \quad \text{NRMSHIFTD}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\exists \vec{\alpha}^-. P) = \exists \vec{\alpha}^-. P'} \quad \text{NRME EXISTS}$$

$$\boxed{\text{nf}(N) = M}$$

$$\overline{\text{nf}(\hat{\alpha}^-) = \hat{\alpha}^-} \quad \text{NRMNUVAR}$$

$$\boxed{\text{nf}(P) = Q}$$

$$\overline{\text{nf}(\hat{\alpha}^+) = \hat{\alpha}^+} \quad \text{NRMPUVAR}$$

$$\boxed{\text{ord vars in } N = \vec{\alpha}}$$

$$\frac{\alpha^- \in \text{vars}}{\text{ord vars in } \alpha^- = \alpha^-} \quad \text{ONVARIN}$$

$$\frac{\alpha^- \notin \text{vars}}{\text{ord vars in } \alpha^- = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}}{\text{ord vars in } \uparrow P = \vec{\alpha}} \quad \text{OShiftU}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}_1 \quad \text{ord vars in } N = \vec{\alpha}_2}{\text{ord vars in } P \rightarrow N = \vec{\alpha}_1, (\vec{\alpha}_2 \setminus \vec{\alpha}_1)} \quad \text{OARROW}$$

$$\frac{\text{vars} \cap \vec{\alpha}^+ = \emptyset \quad \text{ord vars in } N = \vec{\alpha}}{\text{ord vars in } \forall \vec{\alpha}^+. N = \vec{\alpha}} \quad \text{Oforall}$$

$$\boxed{\text{ord vars in } P = \vec{\alpha}}$$

$$\frac{\alpha^+ \in \text{vars}}{\text{ord vars in } \alpha^+ = \alpha^+} \quad \text{OPVARIN}$$

$$\frac{\alpha^+ \notin \text{vars}}{\text{ord vars in } \alpha^+ = \cdot} \quad \text{OPVARNIN}$$

$$\frac{\text{ord vars in } N = \vec{\alpha}}{\text{ord vars in } \downarrow N = \vec{\alpha}} \quad \text{OSHIFTD}$$

$$\frac{\text{vars} \cap \vec{\alpha}^- = \emptyset \quad \text{ord vars in } P = \vec{\alpha}}{\text{ord vars in } \exists \alpha^-. P = \vec{\alpha}} \quad \text{OEXISTS}$$

$$\boxed{\text{ord vars in } N = \vec{\alpha}}$$

$$\frac{}{\text{ord vars in } \hat{\alpha}^- = \cdot} \quad \text{ONUVAR}$$

$$\boxed{\text{ord vars in } P = \vec{\alpha}}$$

$$\frac{}{\text{ord vars in } \hat{\alpha}^+ = \cdot} \quad \text{OPUVAR}$$

$$\boxed{\Gamma \vdash e_1 \& e_2 = e_3} \quad \text{Unification Solution Entry Merge}$$

$$\frac{\Gamma \vdash P_1 \vee P_2 = Q}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \& (\hat{\alpha}^+ : \geq P_2) = (\hat{\alpha}^+ : \geq Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vdash P \geq Q \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \geq Q) = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vdash Q \geq P \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \geq P) \& (\hat{\alpha}^+ : \approx Q) = (\hat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \approx P') = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEPEQEQ}$$

$$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \& (\hat{\alpha}^- : \approx N') = (\hat{\alpha}^- : \approx N)} \quad \text{SMENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3} \quad \text{Merge unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \Rightarrow e_2} \quad \text{Weakening of unification solution entries}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPESUPSUP}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUP}$$

$$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2)} \quad \text{SIMPEPEQEQ}$$

$$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \Rightarrow (\hat{\alpha}^- : \approx N_2)} \quad \text{SIMPENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2} \quad \text{Weakening of unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \simeq e_2}$$

$$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \simeq (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUPSUP}$$

$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \simeq (\hat{\alpha}^+ : \approx P_2)}$	SIMPEEQPEQEQ
$\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \simeq (\hat{\alpha}^- : \approx N_2)}$	SIMPEEQNEQEQ
$\boxed{\Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2}$	
$\boxed{\Gamma; \Theta \vdash N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}$	Negative unification
$\frac{}{\Gamma; \Theta \vdash \alpha^- \overset{u}{\simeq} \alpha^- \Rightarrow \cdot}$	UNVAR
$\frac{\Gamma; \Theta \vdash P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \vdash \uparrow P \overset{u}{\simeq} \uparrow Q \Rightarrow \hat{\sigma}}$	USHIFTU
$\frac{\Gamma; \Theta \vdash P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \Theta \vdash N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}_2}{\Gamma; \Theta \vdash P \rightarrow N \overset{u}{\simeq} Q \rightarrow M \Rightarrow \hat{\sigma}_1 \ \& \ \hat{\sigma}_2}$	UARROW
$\frac{\Gamma, \overrightarrow{\alpha^+}; \Theta \vdash N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \vdash \forall \overrightarrow{\alpha^+}. N \overset{u}{\simeq} \forall \overrightarrow{\alpha^+}. M \Rightarrow \hat{\sigma}}$	UFORALL
$\frac{\hat{\alpha}^- \{\Delta\} \in \Theta \quad \Delta \vdash N}{\Gamma; \Theta \vdash \hat{\alpha}^- \overset{u}{\simeq} N \Rightarrow (\hat{\alpha}^- : \approx N)}$	UNUVAR
$\boxed{\Gamma; \Theta \vdash P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}$	Positive unification
$\frac{}{\Gamma; \Theta \vdash \alpha^+ \overset{u}{\simeq} \alpha^+ \Rightarrow \cdot}$	UPVAR
$\frac{\Gamma; \Theta \vdash N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \vdash \downarrow N \overset{u}{\simeq} \downarrow M \Rightarrow \hat{\sigma}}$	USHIFTD
$\frac{\Gamma, \overrightarrow{\alpha^-}; \Theta \vdash P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \vdash \exists \overrightarrow{\alpha^-}. P \overset{u}{\simeq} \exists \overrightarrow{\alpha^-}. Q \Rightarrow \hat{\sigma}}$	UEXISTS
$\frac{\hat{\alpha}^+ \{\Delta\} \in \Theta \quad \Delta \vdash P}{\Gamma; \Theta \vdash \hat{\alpha}^+ \overset{u}{\simeq} P \Rightarrow (\hat{\alpha}^+ : \approx P)}$	UPUVAR
$\boxed{\Gamma \vdash N}$	Negative type well-formedness
$\boxed{\Gamma \vdash P}$	Positive type well-formedness
$\boxed{\Gamma \vdash N}$	Negative type well-formedness
$\boxed{\Gamma \vdash P}$	Positive type well-formedness
$\boxed{\Gamma \vdash \vec{N}}$	Negative type list well-formedness
$\boxed{\Gamma \vdash \vec{P}}$	Positive type list well-formedness
$\boxed{\Gamma; \Theta \vdash N}$	Negative unification type well-formedness
$\boxed{\Gamma; \Theta \vdash P}$	Positive unification type well-formedness
$\boxed{\Gamma; \Xi \vdash N}$	Negative anti-unification type well-formedness
$\boxed{\Gamma; \Xi \vdash P}$	Positive anti-unification type well-formedness
$\boxed{\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1}$	Antiunification substitution well-formedness
$\boxed{\hat{\sigma} : \Theta}$	Unification substitution well-formedness
$\boxed{\Gamma \vdash^\supset \Theta}$	Unification context well-formedness
$\boxed{\Gamma_1 \vdash \sigma : \Gamma_2}$	Substitution well-formedness
$\boxed{\Gamma \vdash e}$	Unification solution entry well-formedness

Definition rules: 86 good 14 bad  
Definition rule clauses: 168 good 14 bad