$\begin{array}{ll} \alpha,\,\beta,\,\alpha,\,\beta,\,\gamma,\,\delta & \text{type variables} \\ n,\,m,\,i,\,j & \text{index variables} \end{array}$

```
cohort index
n, m
                                                    ::=
                                                                 0
                                                                 n+1
\alpha^+,~\beta^+,~\gamma^+,~\delta^+
                                                                                                           positive variable
                                                                 \alpha^{+n}
\alpha^-,~\beta^-,~\gamma^-,~\delta^-
                                                                                                          negative variable
                                                                 \alpha^{-n}
\alpha^{\pm}, \ \beta^{\pm}
                                                                                                          positive or negative variable
                                                    ::=
                                                                 \alpha^{\pm}
                                                                 \alpha^{\pm n}
                                                    ::=
                                                                                                          substitution
                                                                 id
                                                                 P/\alpha^+
                                                                N/\alpha^-
\overrightarrow{P}/\alpha^+
                                                                 \mathbf{pmas}/\overrightarrow{\alpha^+}

\frac{\mathbf{nmas}}{\widetilde{\alpha}^{+}}/\widetilde{\alpha}^{+}

\overset{\longrightarrow}{\alpha^{+}}/\alpha^{+}

\overset{\longrightarrow}{\alpha^{-}}/\alpha^{-}

                                                                 \mu
                                                                 \sigma_1 \circ \sigma_2
                                                                 \vec{\alpha}_1/\vec{\alpha}_2
                                                                                                S
                                                                  (\sigma)
                                                                                                                 concatenate
                                                                 \mathbf{nf}\left(\sigma'\right)
                                                                                                Μ
                                                                 \sigma'|_{vars}
                                                                                                Μ
                                                                                                           entry of a unification solution
 e
                                                                 \hat{\alpha}^+ :\approx P
                                                                 \widehat{\alpha}^- :\approx N
                                                                 \hat{\alpha}^+ : \geqslant P
                                                                 (e)
                                                                                                S
                                                                 \hat{\sigma}(\hat{\alpha}^+)
                                                                                                Μ
                                                                 \hat{\sigma}(\hat{\alpha}^-)
                                                                                                Μ
                                                                 e_1 \ \& \ e_2
```

unification solution (substitution)

 $\hat{\sigma}$

::=

```
e
                                                          \widehat{\sigma} \backslash vars
                                                          \hat{\sigma}|vars
                                                          \frac{\widehat{\sigma}_1}{\widehat{\sigma}_i} \cup \widehat{\sigma}_2
                                                                                                   concatenate
                                                           (\hat{\sigma})
                                                                                    S
                                                          \mathbf{nf}\left(\widehat{\sigma}'\right)
                                                                                    Μ
                                                          \hat{\sigma}'|_{vars}
                                                                                    Μ
                                                           \hat{\sigma}_1 \& \hat{\sigma}_2
                                                                                    Μ
\hat{\tau}, \ \hat{\rho}
                                                                                             anti-unification substitution
                                               ::=
                                                          \widehat{\alpha}^-:\approx N
                                                          \widehat{\alpha}^- :\approx N
                                                          \vec{N}/\widehat{\alpha^-}
                                                          \hat{\tau}_1 \cup \hat{\tau}_2
\overline{\hat{\tau}_i}^i
                                                                                                   concatenate
                                                           (\hat{\tau})
                                                                                    S
                                                          \hat{\tau}'|_{vars}
                                                                                    Μ
                                                           \hat{\tau}_1 \& \hat{\tau}_2
                                                                                    Μ
P, Q
                                               ::=
                                                                                             positive types
                                                          \alpha^+
                                                          \downarrow N
                                                          \exists \alpha^-.P
                                                           [\sigma]P
                                                                                    Μ
N, M
                                                                                             negative types
                                               ::=
                                                          \alpha^{-}
                                                          \uparrow P
                                                          \forall \alpha^+.N
                                                           P \rightarrow N
                                                          [\sigma]N
                                                                                    Μ
                                                                                             positive variable list
                                                                                                   empty list
                                                                                                   a variable
                                                                                                   a variable
                                                                                                   concatenate lists
                                                                                             negative variables
                                                                                                   empty list
                                                                                                   a variable
                                                                                                   variables
                                                                                                   concatenate lists
\overrightarrow{\alpha^{\pm}}, \overrightarrow{\beta^{\pm}}, \overrightarrow{\gamma^{\pm}}, \overrightarrow{\delta^{\pm}}
                                                                                             positive or negative variable list
```

```
empty list
                                                    a variable
                         \overrightarrow{pa}
                                                    variables
                                                    concatenate lists
P, Q
                                                multi-quantified positive types
                                                    P \neq \exists \dots
                         [\sigma]P
                                         Μ
                         [\hat{\tau}]P
                                         Μ
                         [\hat{\sigma}]P
                                         Μ
                         [\mu]P
                                         Μ
                         (P)
                                         S
                         P_1 \vee P_2
                                         Μ
                         \mathbf{nf}(P')
                                         Μ
N, M
                                                multi-quantified negative types
                         \alpha^{-}

\uparrow P 

P \to N 

\forall \alpha^+. N

                                                   N \neq \forall \dots
                         [\hat{\tau}]N
                                         Μ
                         [\mu]N
                                         Μ
                         [\hat{\sigma}]N
                                         Μ
                         (N)
                                         S
                         \mathbf{nf}\left( N^{\prime}\right)
\vec{P}, \ \vec{Q}
                                                list of positive types
                                                    empty list
                                                    a singel type
                                                    concatenate lists
\overrightarrow{N}, \overrightarrow{M}
                                                list of negative types
                                                    empty list
                                                    a singel type
                                                    concatenate lists
                         \mathbf{nf}(\vec{N}')
\Delta, \Gamma
                                                declarative type context
                                                    empty context
                                                    list of variables
                                                    list of variables
                                                    list of variables
                         vars
                         \overline{\Gamma_i}^{\;i}
                                                    concatenate contexts
                                         S
                         \Theta(\widehat{\alpha}^+)
                                         Μ
```

```
\Theta(\hat{\alpha}^-)
                                           Μ
Θ
                                                  unification type variable context
                                                     empty context
                                                     list of variables
                                                     list of variables
                      vars
                      \overline{\Theta_i}^{i}
                                                     concatenate contexts
                                           S
                      (\Theta)
                      \Theta|_{vars}
                                                     leave only those variables that are in the set
                      \Theta_1 \cup \Theta_2
Ξ
                                                  anti-unification type variable context
                                                     empty context
                                                     list of variables
                                                     concatenate contexts
                                           S
                                           Μ
\vec{\alpha}, \vec{\beta}
                                                  ordered positive or negative variables
                                                     empty list
                                                     list of variables
                                                     list of variables
                                                     list of variables
                                                     list of variables
                                                     list of variables
                      \overrightarrow{\alpha}_1 \backslash vars
                                                     setminus
                                                     context
                      vars
                                                     concatenate contexts
                      (\vec{\alpha})
                                           S
                                                     parenthesis
                      [\mu]\vec{\alpha}
                                                     apply moving to list
                                            Μ
                      ord vars in P
                                           Μ
                      ord vars in N
                                           Μ
                      ord vars in P
                                           Μ
                      \operatorname{\mathbf{ord}} \operatorname{\mathbf{vars}} \operatorname{\mathbf{in}} N
                                           Μ
                                                  set of variables
vars
                      Ø
                                                     empty set
                     \mathbf{fv} P
                                                     free variables
                     \mathbf{fv} N
                                                     free variables
                      fv imP
                                                     free variables
                                                     free variables
                      fv imN
                                                     set intersection
                      vars_1 \cap vars_2
                                                     set union
                      vars_1 \cup vars_2
                      vars_1 \backslash vars_2
                                                     set complement
                     mv imP
                                                     movable variables
                      mv imN
                                                     movable variables
                      \mathbf{u}\mathbf{v} N
                                                     unification variables
```

		$\begin{array}{l} \mathbf{uv} \ P \\ \mathbf{fv} \ N \\ \mathbf{fv} \ P \\ (vars) \\ \overrightarrow{\alpha} \\ [\mu] vars \\ \mathbf{dom} \ (\widehat{\sigma}) \\ \mathbf{dom} \ (\widehat{\tau}) \\ \mathbf{dom} \ (\Theta) \end{array}$	S M M M	unification variables free variables free variables parenthesis ordered list of variables apply moving to varset
μ		$\begin{array}{l} .\\ pma1 \mapsto pma2 \\ nma1 \mapsto nma2 \\ \mu_1 \cup \mu_2 \\ \hline{\mu_1} \circ \mu_2 \\ \hline{\mu_i}^i \\ \mu _{vars} \\ \mu^{-1} \\ \mathbf{nf} \left(\mu' \right) \end{array}$	M M M M	empty moving Positive unit substitution Positive unit substitution Set-like union of movings Composition concatenate movings restriction on a set inversion
\hat{lpha}^{\pm}	::=	\hat{lpha}^{\pm}		positive/negative unification variable
\widehat{lpha}^+	::= 	$\hat{\alpha}^+$ $\hat{\alpha}^+\{\Delta\}$ $\hat{\alpha}^\pm$		positive unification variable
		$\widehat{lpha}^ \widehat{lpha}^{\{N,M\}}$ $\widehat{lpha}^{\{\Delta\}}$ \widehat{lpha}^\pm		negative unification variable
$\overrightarrow{\widetilde{\alpha}^+}, \ \overrightarrow{\widetilde{\beta}^+}$::=	$ \begin{array}{c} \widehat{\alpha}^{+} \\ \widehat{\alpha}^{+} \{\Delta\} \\ \widehat{\alpha}^{+} \\ \widehat{\alpha}^{+}_{i} \end{array} $		positive unification variable list empty list a variable from a normal variable from a normal variable, context unspecified concatenate lists
$\overrightarrow{\widehat{\alpha}^-}$, $\overrightarrow{\widehat{\beta}^-}$::=	$ \frac{\alpha^{+}}{\widehat{\alpha}^{+}i} $ $ \frac{\widehat{\alpha}^{-}}{\widehat{\alpha}^{-}} \underbrace{\frac{\Xi}{\widehat{\alpha}^{-}} \{\Delta\}}_{\widehat{\alpha}^{-}i} $ $ \frac{\widehat{\alpha}^{-}}{\widehat{\alpha}^{-}i} $		negative unification variable list empty list a variable from an antiunification context from a normal variable from a normal variable, context unspecified concatenate lists

 $\mathbf{nf}(P')$

М

a positive algorithmic type (potentially with metavariables) $\,$

N, M ::= a negative algorithmic type (potentially with metavariables)

 $\begin{vmatrix} \alpha^{-} \\ \widehat{\alpha}^{-} \\ \uparrow P \\ | P \rightarrow N \\ | \overrightarrow{\alpha^{+}} \cdot N \\ | [\sigma] N & M \\ | [\widehat{\tau}] N & M \\ | [\mu] N & M \\ | (N) & S \\ | \mathbf{nf} (N') & M \end{vmatrix}$

 $\begin{array}{ccc} auSol & & ::= & \\ & | & (\Xi, Q, \widehat{\tau}_1, \widehat{\tau}_2) \\ & | & (\Xi, N, \widehat{\tau}_1, \widehat{\tau}_2) \end{array}$

terminals

```
:≥
                                       :\simeq
formula
                                       judgement
                                       formula_1 .. formula_n
                                       \mu : vars_1 \leftrightarrow vars_2
                                       \mu is bijective
                                       \hat{\sigma} is functional
                                       \hat{\sigma}_1 \in \hat{\sigma}_2
                                       \hat{\sigma}_1 \subseteq \hat{\sigma}_2
                                       \mathit{vars}_1 \subseteq \mathit{vars}_2
                                       vars_1 = vars_2
                                       vars is fresh
                                       \alpha^- \not\in \mathit{vars}
                                       \alpha^+ \notin vars
                                       \alpha^- \in vars
                                       \alpha^+ \in vars
                                       \widehat{\alpha}^+ \in \mathit{vars}
                                       \widehat{\alpha}^- \in \mathit{vars}
                                       \widehat{\alpha}^- \in \Theta
                                       \widehat{\alpha}^+ \in \Theta
                                       if any other rule is not applicable
                                       \vec{\alpha}_1 = \vec{\alpha}_2
                                       e_1 = e_2
                                       N = M
                                       N \neq M
                                       P \neq Q
\boldsymbol{A}
                           ::=
                                       \begin{array}{l} \Gamma;\,\Theta \vDash N \leqslant M \dashv \widehat{\sigma} \\ \Gamma;\,\Theta \vDash P \geqslant Q \dashv \widehat{\sigma} \end{array}
                                                                                                                               Negative subtyping
                                                                                                                               Positive supertyping
AU
                             ..

\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)

\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 = (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)

E1
                                                                                                                              Negative multi-quantified type equivalence
                                                                                                                              Positive multi-quantified type equivalence
```

```
D1
                   ::=
                           \Gamma \vdash N \simeq_1^{\leqslant} M
                                                                                   Negative equivalence on MQ types
                           \Gamma \vdash P \simeq_1^{\leqslant} Q
                                                                                   Positive equivalence on MQ types
                           \Gamma \vdash N \leqslant_1 M
                                                                                   Negative subtyping
                           \Gamma \vdash P \geqslant_1 Q
                                                                                   Positive supertyping
                           \Gamma_2 \vdash \sigma_1 \simeq_1^{\leqslant} \sigma_2 : \Gamma_1
                                                                                   Equivalence of substitutions
D\theta
                   ::=
                           \Gamma \vdash N \simeq_0^{\leqslant} M
                                                                                   Negative equivalence
                           \Gamma \vdash P \simeq_0^{\mathrm{d}} Q
                                                                                   Positive equivalence
                           \Gamma \vdash N \leqslant_0 M
                                                                                   Negative subtyping
                           \Gamma \vdash P \geqslant_0 Q
                                                                                   Positive supertyping
EQ
                           N = M
                                                                                   Negative type equality (alpha-equivalence)
                           P = Q
                                                                                   Positive type equuality (alphha-equivalence)
                           P = Q
LUBF
                            P_1 \vee P_2 === Q
                           ord vars in P === \vec{\alpha}
                            ord vars in N = = \vec{\alpha}
                           \mathbf{ord}\ vars \mathbf{in}\ P === \overrightarrow{\alpha}
                           \mathbf{ord}\ vars \mathbf{in}\ N = = = \overrightarrow{\alpha}
                           \mathbf{nf}(N') === N
                           \mathbf{nf}(P') === P
                           \mathbf{nf}(N') === N
                           \mathbf{nf}(P') === P
                           \mathbf{nf}(\vec{N}') = = \vec{N}
                           \mathbf{nf}(\vec{P}') === \vec{P}
                           \mathbf{nf}(\sigma') = = \sigma
                           \mathbf{nf}(\mu') === \mu
                           \mathbf{nf}(\widehat{\sigma}') = = = \widehat{\sigma}
                            \sigma'|_{vars}
                            \hat{\sigma}'|_{vars}
                            \hat{\tau}'|_{vars}
                           \Xi'|_{vars}
                            e_1 \& e_2
                            \hat{\sigma}_1 \& \hat{\sigma}_2
                            \hat{\tau}_1 \& \hat{\tau}_2
                            \mathbf{dom}(\hat{\sigma}) === vars
                            \mathbf{dom}(\hat{\tau}) === vars
                            \mathbf{dom}\left(\Theta\right) === vars
LUB
                   ::=
                           \Gamma \vDash P_1 \vee P_2 = Q
                                                                                  Least Upper Bound (Least Common Supertype)
                            \mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q
Nrm
                  ::=
```

 $\mathbf{nf}(N) = M$

```
\mathbf{nf}(P) = Q
                                \mathbf{nf}(N) = M
                                \mathbf{nf}(P) = Q
Order
                                \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}
                                \operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}
                                ord vars in \overline{N} = \vec{\alpha}
                                ord vars in P = \vec{\alpha}
SM
                        ::=
                                \Gamma \vdash e_1 \& e_2 = e_3
                                                                           Unification Solution Entry Merge
                                \Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3
                                                                           Merge unification solutions
SImp
                                \Gamma \vdash e_1 \Rightarrow e_2
                                                                           Weakening of unification solution entries
                                \Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2
                                                                           Weakening of unification solutions
                                \Gamma \vdash e_1 \simeq e_2
                                \Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2
U
                        ::=
                                \Gamma;\Theta \models N \stackrel{u}{\simeq} M = \hat{\sigma}
                                                                           Negative unification
                                \Gamma:\Theta \models P \stackrel{u}{\simeq} Q = \hat{\sigma}
                                                                           Positive unification
WF
                                \Gamma \vdash N
                                                                           Negative type well-formedness
                                \Gamma \vdash P
                                                                           Positive type well-formedness
                                \Gamma \vdash N
                                                                           Negative type well-formedness
                                \Gamma \vdash P
                                                                           Positive type well-formedness
                                \Gamma \vdash \overrightarrow{N}
                                                                           Negative type list well-formedness
                                \Gamma \vdash \overrightarrow{P}
                                                                           Positive type list well-formedness
                                \Gamma;\Theta \vdash N
                                                                           Negative unification type well-formedness
                                \Gamma;\Theta \vdash P
                                                                           Positive unification type well-formedness
                                \Gamma;\Xi \vdash N
                                                                           Negative anti-unification type well-formedness
                                \Gamma;\Xi\vdash P
                                                                           Positive anti-unification type well-formedness
                                \Gamma; \Xi_2 \vdash \widehat{\tau} : \Xi_1
                                                                           Antiunification substitution well-formedness
                                \hat{\sigma}:\Theta
                                                                           Unification substitution well-formedness
                                \Gamma \vdash^{\supseteq} \Theta
                                                                           Unification context well-formedness
                                \Gamma_1 \vdash \sigma : \Gamma_2
                                                                           Substitution well-formedness
                                \Gamma \vdash e
                                                                           Unification solution entry well-formedness
judgement
                                A
                                AU
                                E1
                                D1
                                D\theta
                                EQ
```

LUB

 $user_syntax$

 α nnvarsauSolterminalsformula

 $\boxed{\Gamma;\,\Theta \vDash N \leqslant M \dashv \widehat{\sigma}} \quad \text{Negative subtyping}$

$$\frac{\Gamma; \Theta \vDash \alpha^{-} \leqslant \alpha^{-} \dashv \cdot}{\Gamma; \Theta \vDash \mathbf{nf}(P) \overset{u}{\simeq} \mathbf{nf}(Q) \dashv \widehat{\sigma}} \qquad \text{ASHIFTU}$$

$$\frac{\Gamma; \Theta \vDash \uparrow P \leqslant \uparrow Q \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \uparrow P \leqslant \uparrow Q \dashv \widehat{\sigma}} \qquad \text{ASHIFTU}$$

$$\frac{\Gamma; \Theta \vDash P \geqslant Q \dashv \widehat{\sigma}_{1} \quad \Gamma; \Theta \vDash N \leqslant M \dashv \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash P \to N \leqslant Q \to M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \quad \text{AArrow}$$

$$\frac{\Gamma, \overrightarrow{\beta^{+}}; \Theta, \widehat{\alpha}^{+} \{\Gamma, \overrightarrow{\beta^{+}}\} \vDash [\widehat{\alpha}^{+}/\alpha^{+}] N \leqslant M \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \forall \alpha^{+}. N \leqslant \forall \overrightarrow{\beta^{+}}. M \dashv \widehat{\sigma} \setminus \widehat{\alpha^{+}}} \quad \text{AForall}$$

 $\Gamma; \Theta \models P \geqslant Q \dashv \widehat{\sigma}$ Positive supertyping

$$\frac{\Gamma; \Theta \vDash \alpha^{+} \geqslant \alpha^{+} \Rightarrow \cdot}{\Gamma; \Theta \vDash \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \Rightarrow \widehat{\sigma}} \quad \text{ASHIFTD}$$

$$\frac{\Gamma; \Theta \vDash \mathsf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \Rightarrow \widehat{\sigma}}{\Gamma; \Theta \vDash \mathsf{N} \geqslant \mathsf{M} \Rightarrow \widehat{\sigma}} \quad \text{ASHIFTD}$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}}; \Theta, \widehat{\alpha}^{-} \{\Gamma, \overrightarrow{\beta^{-}}\} \vDash [\widehat{\alpha}^{-}/\alpha^{-}] P \geqslant Q \Rightarrow \widehat{\sigma}}{\Gamma; \Theta \vDash \exists \alpha^{-}. P \geqslant \exists \overrightarrow{\beta^{-}}. Q \Rightarrow \widehat{\sigma} \setminus \widehat{\alpha^{-}}} \quad \text{AEXISTS}$$

$$\frac{\widehat{\alpha}^{+} \{\Delta\} \in \Theta \quad \mathbf{upgrade} \Gamma \vdash P \mathbf{to} \Delta = Q}{\Gamma; \Theta \vDash \widehat{\alpha}^{+} \geqslant P \Rightarrow (\widehat{\alpha}^{+} : \geqslant Q)} \quad \text{APUVAR}$$

 $\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$

$$\frac{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} \alpha^{+} \dashv (\cdot, \alpha^{+}, \cdot, \cdot)}{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUSHIFTD}$$

$$\frac{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \vDash \downarrow N_{1} \stackrel{a}{\simeq} \downarrow N_{2} \dashv (\Xi, \downarrow M, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUSHIFTD}$$

$$\frac{\overrightarrow{\alpha^{-}} \cap \Gamma = \varnothing \quad \Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \vDash \exists \overrightarrow{\alpha^{-}}. P_{1} \stackrel{a}{\simeq} \exists \overrightarrow{\alpha^{-}}. P_{2} \dashv (\Xi, \exists \overrightarrow{\alpha^{-}}. Q, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUEXISTS}$$

 $\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 = (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$

 $N \simeq_1^D M$ Negative multi-quantified type equivalence

$$\frac{\alpha^{-} \simeq_{1}^{D} \alpha^{-}}{\frac{P \simeq_{1}^{D} Q}{\uparrow P \simeq_{1}^{D} \uparrow Q}} \quad \text{E1ShiftU}$$

$$\frac{P \simeq_{1}^{D} Q \quad N \simeq_{1}^{D} M}{P \to N \simeq_{1}^{D} Q \to M} \quad \text{E1Arrow}$$

$$\overrightarrow{\alpha^{+}} \cap \mathbf{fv} M = \varnothing \quad \mu : (\overrightarrow{\beta^{+}} \cap \mathbf{fv} M) \leftrightarrow (\overrightarrow{\alpha^{+}} \cap \mathbf{fv} N) \quad N \simeq_{1}^{D} [\mu] M$$

$$\overrightarrow{\forall \alpha^{+}} \cdot N \simeq_{1}^{D} \forall \overrightarrow{\beta^{+}} \cdot M$$

$$\text{E1Forall}$$

 $P \simeq_1^D Q$ Positive multi-quantified type equivalence

$$\frac{\alpha^{+} \simeq_{1}^{D} \alpha^{+}}{\sqrt{N} \simeq_{1}^{D} M} \quad \text{E1ShiftD}$$

$$\frac{\overrightarrow{\alpha^{-}} \cap \text{fv } Q = \varnothing \quad \mu : (\overrightarrow{\beta^{-}} \cap \text{fv } Q) \leftrightarrow (\overrightarrow{\alpha^{-}} \cap \text{fv } P) \quad P \simeq_{1}^{D} [\mu]Q}{\exists \overrightarrow{\alpha^{-}} . P \simeq_{1}^{D} \exists \overrightarrow{\beta^{-}} . Q} \quad \text{E1Exists}$$

$$\frac{\Gamma \vdash N \leqslant_{1} M \quad \Gamma \vdash M \leqslant_{1} N}{\Gamma \vdash N \simeq_{1}^{\varsigma} M} \quad \text{D1NDEF}$$

 $\overline{\Gamma \vdash P \simeq_1^{\epsilon} Q}$ Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geqslant_1 Q \quad \Gamma \vdash Q \geqslant_1 P}{\Gamma \vdash P \simeq_1^{\varsigma} Q} \quad \text{D1PDEF}$$

 $\Gamma \vdash N \leq_1 M$ Negative subtyping

 $\Gamma \vdash P \geqslant_1 Q$ Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{1} \alpha^{+}}{\Gamma \vdash \alpha^{+} \geqslant_{1} \alpha^{+}} \quad D1PVAR$$

$$\frac{\text{>}}{\Gamma \vdash \downarrow N \geqslant_{1} \downarrow M} \quad D1SHIFTD$$

$$\frac{\text{fv } P \cap \overrightarrow{\beta^{-}} = \varnothing \quad \Gamma, \overrightarrow{\beta^{-}} \vdash N_{i} \quad \Gamma, \overrightarrow{\beta^{-}} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^{-}}]P \geqslant_{1} Q}{\Gamma \vdash \exists \overrightarrow{\alpha^{-}}.P \geqslant_{1} \exists \overrightarrow{\beta^{-}}.Q} \quad D1EXISTS$$

 $\begin{array}{|c|c|c|c|c|}\hline \Gamma_2 \vdash \sigma_1 & \simeq_1^{\varsigma} \sigma_2 : \Gamma_1 \\ \hline \Gamma \vdash N & \simeq_0^{\varsigma} M \\ \hline \end{array} \quad \text{Negative equivalence}$

$$\frac{\Gamma \vdash N \leqslant_0 M \quad \Gamma \vdash M \leqslant_0 N}{\Gamma \vdash N \simeq_0^\circ M} \quad \text{D0NDef}$$

 $\Gamma \vdash P \simeq_0^{\epsilon} Q$ Positive equivalence

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash Q \geqslant_0 P}{\Gamma \vdash P \simeq_0^{\varsigma} Q} \quad \text{D0PDEF}$$

 $\Gamma \vdash N \leqslant_0 M$ Negative subtyping

$$\frac{\Gamma \vdash \alpha^- \leqslant_0 \alpha^-}{\Gamma \vdash P \stackrel{\sim}{=} 0} \frac{D0NVAR}{Q}$$

$$\frac{\Gamma \vdash P \stackrel{\sim}{=} Q}{\Gamma \vdash P \leqslant_0 \uparrow Q} D0SHIFTU$$

$$\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leqslant_0 M \quad M \neq \forall \beta^+.M'}{\Gamma \vdash \forall \alpha^+.N \leqslant_0 M} D0FORALLL$$

$$\frac{\Gamma, \alpha^+ \vdash N \leqslant_0 M}{\Gamma \vdash N \leqslant_0 \forall \alpha^+.M} D0FORALLR$$

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash N \leqslant_0 M}{\Gamma \vdash P \rightarrow N \leqslant_0 Q \rightarrow M} D0ARROW$$

 $\overline{|\Gamma \vdash P \geqslant_0 Q|}$ Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{0} \alpha^{+}}{\Gamma \vdash \lambda^{-} \geqslant_{0} M} \quad D0PVAR$$

$$\frac{\Gamma \vdash N \simeq_{0}^{\leq} M}{\Gamma \vdash \downarrow N \geqslant_{0} \downarrow M} \quad D0SHIFTD$$

$$\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^{-}]P \geqslant_{0} Q \quad Q \neq \exists \alpha^{-}.Q'}{\Gamma \vdash \exists \alpha^{-}.P \geqslant_{0} Q} \quad D0EXISTSL$$

$$\frac{\Gamma, \alpha^{-} \vdash P \geqslant_{0} Q}{\Gamma \vdash P \geqslant_{0} \exists \alpha^{-}.Q} \quad D0EXISTSR$$

 $\overline{N=M}$ Negative type equality (alpha-equivalence) $\overline{P=Q}$ Positive type equality (alphha-equivalence) $\overline{P=Q}$ $\overline{P_1 \vee P_2}$

 $\mathbf{ord}\ vars\mathbf{in}\ P$

 $\mathbf{ord}\ vars\mathbf{in}\ N$

ord vars in P

 $\overline{\mathbf{ord}\ vars\mathbf{in}\ N}$

 $\mathbf{nf}(N')$

 $\mathbf{nf}\left(P'\right)$

 $\mathbf{nf}\left(N'
ight)$

 $\mathbf{nf}\left(P'\right)$

 $\mathbf{nf}\,(\overrightarrow{\vec{N}'})$

 $\mathbf{nf}\,(\overrightarrow{\vec{P}}')$

 $\mathbf{nf}\left(\sigma'\right)$

 $\mathbf{nf}\left(\mu'\right)$

 $\mathbf{nf}\left(\widehat{\sigma}'\right)$

 $\sigma'|_{vars}$

 $[\hat{\sigma}'|_{vars}]$

 $|\hat{ au}'|_{vars}$

 $\Xi'|_{vars}$

 $e_1 \& e_2$

 $\hat{\sigma}_1 \& \hat{\sigma}_2$

 $[\hat{\tau}_1 \& \hat{\tau}_2]$

 $\operatorname{dom}(\widehat{\sigma})$

 $\operatorname{\mathbf{dom}}(\widehat{\tau})$

 $\operatorname{\mathbf{dom}}(\Theta)$

 $\overline{|\Gamma \models P_1 \lor P_2 = Q|}$ Least Upper Bound (Least Common Supertype)

$$\frac{\Gamma, \cdot \models \mathbf{nf} (\downarrow N) \stackrel{a}{\simeq} \mathbf{nf} (\downarrow M) \dashv (\Xi, P, \widehat{\tau}_1, \widehat{\tau}_2)}{\Gamma \models \downarrow N \lor \downarrow M = \exists \widehat{\alpha}^-. [\widehat{\alpha}^-/\Xi]P} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \widehat{\alpha}^-, \widehat{\beta}^- \models P_1 \lor P_2 = Q}{\Gamma \models \exists \widehat{\alpha}^-. P_1 \lor \exists \widehat{\beta}^-. P_2 = Q} \quad \text{LUBEXISTS}$$

 $\mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q$

$$\begin{array}{cccc} \Gamma = \Delta, \overrightarrow{\alpha^{\pm}} & \overrightarrow{\beta^{\pm}} \text{ is fresh} & \overrightarrow{\gamma^{\pm}} \text{ is fresh} \\ \Delta, \overrightarrow{\beta^{\pm}}, \overrightarrow{\gamma^{\pm}} \vDash [\overrightarrow{\beta^{\pm}}/\overrightarrow{\alpha^{\pm}}]P \vee [\overrightarrow{\gamma^{\pm}}/\overrightarrow{\alpha^{\pm}}]P = Q \\ \hline & \mathbf{upgrade} \, \Gamma \vdash P \, \mathbf{to} \, \Delta = Q \end{array} \quad \text{LUBUPGRADE}$$

 $\mathbf{nf}\left(N\right) =M$

 $\mathbf{nf}(P) = Q$

 $\mathbf{nf}(N) = M$

$$\overline{\mathbf{nf}(\widehat{\alpha}^{-}) = \widehat{\alpha}^{-}}$$
 NRMNUVAR

$$\mathbf{nf}(P) = Q$$

$$\overline{\mathbf{nf}(\hat{\alpha}^+) = \hat{\alpha}^+} \quad N_{RM}PUV_{AR}$$

$\mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}$

$$\frac{\alpha^{-} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \alpha^{-}} \quad \text{ONVARIN}$$

$$\frac{\alpha^{-} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \uparrow P = \overrightarrow{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}_1 \quad \mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}_2}{\mathbf{ord} \ vars \mathbf{in} \ P \to N = \overrightarrow{\alpha}_1, (\overrightarrow{\alpha}_2 \backslash \overrightarrow{\alpha}_1)} \quad \text{OARROW}$$

$$\frac{vars \cap \overrightarrow{\alpha^+} = \varnothing \quad \mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}}{\mathbf{ord} \ vars \mathbf{in} \ \forall \overrightarrow{\alpha^+}. N = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

$\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}$

$$\frac{\alpha^{+} \in vars}{\operatorname{ord} \, vars \operatorname{in} \alpha^{+} = \alpha^{+}} \quad \text{OPVARIN}$$

$$\frac{\alpha^{+} \notin vars}{\operatorname{ord} \, vars \operatorname{in} \alpha^{+} = \cdot} \quad \text{OPVARNIN}$$

$$\frac{\operatorname{ord} \, vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} \, vars \operatorname{in} \setminus N = \overrightarrow{\alpha}} \quad \text{OSHIFTD}$$

$$\frac{vars \cap \overrightarrow{\alpha^{-}} = \varnothing \quad \operatorname{ord} \, vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} \, vars \operatorname{in} \exists \overrightarrow{\alpha^{-}} . P = \overrightarrow{\alpha}} \quad \text{OEXISTS}$$

$$\operatorname{ord} \, vars \operatorname{in} \exists \overrightarrow{\alpha^{-}} . P = \overrightarrow{\alpha}$$

$\mathbf{ord} \ vars \mathbf{in} \ N = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}} \operatorname{\mathit{vars}} \operatorname{\mathbf{in}} \widehat{\alpha}^- = \cdot} \quad \operatorname{ONUVAR}$$

 $\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}} \operatorname{\mathit{vars}} \operatorname{\mathbf{in}} \widehat{\alpha}^{+} = \cdot} \quad \operatorname{OPUVar}$$

 $\Gamma \vdash e_1 \& e_2 = e_3$ Unification Solution Entry Merge

$$\frac{\Gamma \vDash P_1 \lor P_2 = Q}{\Gamma \vdash (\widehat{\alpha}^+ : \geqslant P_1) \& (\widehat{\alpha}^+ : \geqslant P_2) = (\widehat{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vDash P \geqslant Q \dashv \widehat{\sigma}'}{\Gamma \vdash (\widehat{\alpha}^+ : \approx P) \& (\widehat{\alpha}^+ : \geqslant Q) = (\widehat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vDash Q \geqslant P \dashv \widehat{\sigma}'}{\Gamma \vdash (\widehat{\alpha}^+ : \geqslant P) \& (\widehat{\alpha}^+ : \approx Q) = (\widehat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\begin{array}{c} & < \mathsf{multiple parses} > \\ & \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \approx P') = (\hat{\alpha}^+ : \approx P) \\ & < \mathsf{multiple parses} > \\ & \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \approx P') = (\hat{\alpha}^+ : \approx P) \\ & \vdash (\hat{\alpha}^+ : \approx P_1) \& (\hat{\alpha}^- : \approx N') = (\hat{\alpha}^- : \approx N) \\ \hline & \vdash (\hat{\alpha}^- : \approx N_1) \& (\hat{\alpha}^- : \approx N') = (\hat{\alpha}^- : \approx N) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2) \\ \hline & \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P$$

$$\frac{\Gamma, \overrightarrow{\alpha^{-}}; \Theta \vDash P \overset{u}{\simeq} Q \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \overrightarrow{\beta\alpha^{-}}. P \overset{u}{\simeq} \overrightarrow{\beta\alpha^{-}}. Q \dashv \widehat{\sigma}} \quad \text{UEXISTS}$$

$$\frac{\widehat{\alpha}^{+} \{\Delta\} \in \Theta \quad \Delta \vdash P}{\Gamma; \Theta \vDash \widehat{\alpha}^{+} \overset{u}{\simeq} P \dashv (\widehat{\alpha}^{+} : \approx P)} \quad \text{UPUVAR}$$

 $\overline{\Gamma \vdash N}$ Negative type well-formedness

 $\overline{\Gamma \vdash P}$ Positive type well-formedness

 $\overline{\Gamma \vdash N}$ Negative type well-formedness

 $\overline{\Gamma \vdash P}$ Positive type well-formedness

 $\Gamma \vdash \overrightarrow{N}$ Negative type list well-formedness

 $\Gamma \vdash \overrightarrow{P}$ Positive type list well-formedness

 $\Gamma; \Theta \vdash N$ Negative unification type well-formedness

 $\Gamma; \Theta \vdash P$ Positive unification type well-formedness

 $\Gamma;\Xi \vdash N$ Negative anti-unification type well-formedness

 $\Gamma;\Xi\vdash P$ Positive anti-unification type well-formedness

 $\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1$ Antiunification substitution well-formedness

 $\widehat{\sigma} : \Theta$ Unification substitution well-formedness

 $\Gamma \vdash^{\supseteq} \Theta$ Unification context well-formedness

 $\Gamma_1 \vdash \sigma : \Gamma_2$ Substitution well-formedness

 $\Gamma \vdash e$ Unification solution entry well-formedness

Definition rules: 74 good 14 bad Definition rule clauses: 144 good 14 bad