

| | |
|------------------------------------------------|-----------------|
| $\alpha, \beta, \alpha, \beta, \gamma, \delta$ | type variables |
| n, m, i, j | index variables |
| x, y, z | term variables |

| | | |
|---------------------------------------------------------------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|---------------------------------------------------------------------------------------|
| | $ \begin{array}{l} \cdot \\ e \\ \hat{\sigma} \backslash vars \\ \hat{\sigma} vars \\ \hat{\sigma}_1 \cup \hat{\sigma}_2 \\ \widehat{\sigma}_i^i \\ (\hat{\sigma}) \quad \text{S} \\ \mathbf{nf}(\hat{\sigma}') \quad \text{M} \\ \hat{\sigma}' vars \quad \text{M} \\ \hat{\sigma}_1 \ \& \ \hat{\sigma}_2 \quad \text{M} \end{array} $ | concatenate |
| $\hat{\tau}, \hat{\rho}$ | $ \begin{array}{l} ::= \\ \cdot \\ \hat{\alpha}^- : \approx N \\ \hat{\alpha}^- : \approx N \\ \overrightarrow{\alpha^-} / \overrightarrow{\alpha^-} \\ \overrightarrow{N} / \overrightarrow{\alpha^-} \\ \hat{\tau}_1 \cup \hat{\tau}_2 \\ \widehat{\tau}_i^i \\ (\hat{\tau}) \quad \text{S} \\ \hat{\tau}' vars \quad \text{M} \\ \hat{\tau}_1 \ \& \ \hat{\tau}_2 \quad \text{M} \end{array} $ | anti-unification substitution concatenate |
| P, Q, R | $ \begin{array}{l} ::= \\ \alpha^+ \\ \downarrow N \\ \exists \alpha^-. P \\ [\sigma] P \quad \text{M} \end{array} $ | positive types |
| N, M, K | $ \begin{array}{l} ::= \\ \alpha^- \\ \uparrow P \\ \forall \alpha^+. N \\ P \rightarrow N \\ [\sigma] N \quad \text{M} \end{array} $ | negative types |
| $\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}, \overrightarrow{\gamma^+}, \overrightarrow{\delta^+}$ | $ \begin{array}{l} ::= \\ \cdot \\ \alpha^+ \\ \overrightarrow{\alpha^+} \\ \overrightarrow{\alpha^+}^i \\ \alpha^+_i \end{array} $ | positive variable list empty list a variable a variable concatenate lists |
| $\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}, \overrightarrow{\gamma^-}, \overrightarrow{\delta^-}$ | $ \begin{array}{l} ::= \\ \cdot \\ \alpha^- \\ \overrightarrow{\alpha^-} \\ \overrightarrow{\alpha^-}^i \\ \alpha^-_i \end{array} $ | negative variables empty list a variable variables concatenate lists |
| $\overrightarrow{\alpha^\pm}, \overrightarrow{\beta^\pm}, \overrightarrow{\gamma^\pm}, \overrightarrow{\delta^\pm}$ | $::= $ | positive or negative variable list |

| | | | |
|--------------------|-----|---------------------------------|---------------------------------|
| | | \cdot | empty list |
| | | α^\pm | a variable |
| | | $\vec{\mathbf{p}}\mathbf{a}$ | variables |
| | | $\overrightarrow{\alpha^\pm}_i$ | concatenate lists |
| P, Q | ::= | | multi-quantified positive types |
| | | α^+ | |
| | | $\downarrow N$ | |
| | | $\exists \alpha^-. P$ | $P \neq \exists \dots$ |
| | | $[\sigma]P$ | M |
| | | $[\hat{\tau}]P$ | M |
| | | $[\hat{\sigma}]P$ | M |
| | | $[\mu]P$ | M |
| | | (P) | S |
| | | $P_1 \vee P_2$ | M |
| | | $\mathbf{nf}(P')$ | M |
| N, M | ::= | | multi-quantified negative types |
| | | α^- | |
| | | $\uparrow P$ | |
| | | $P \rightarrow N$ | |
| | | $\forall \alpha^+. N$ | $N \neq \forall \dots$ |
| | | $[\sigma]N$ | M |
| | | $[\hat{\tau}]N$ | M |
| | | $[\mu]N$ | M |
| | | $[\hat{\sigma}]N$ | M |
| | | (N) | S |
| | | $\mathbf{nf}(N')$ | M |
| \vec{P}, \vec{Q} | ::= | | list of positive types |
| | | \cdot | empty list |
| | | P | a singel type |
| | | \overrightarrow{P}_i | concatenate lists |
| | | $\mathbf{nf}(\vec{P}')$ | M |
| \vec{N}, \vec{M} | ::= | | list of negative types |
| | | \cdot | empty list |
| | | N | a singel type |
| | | \overrightarrow{N}_i | concatenate lists |
| | | $\mathbf{nf}(\vec{N}')$ | M |
| Δ, Γ | ::= | | declarative type context |
| | | \cdot | empty context |
| | | $\overrightarrow{\alpha^+}$ | list of variables |
| | | $\overrightarrow{\alpha^-}$ | list of variables |
| | | $\overrightarrow{\alpha^\pm}$ | list of variables |
| | | $vars$ | |
| | | $\overrightarrow{\Gamma}_i$ | concatenate contexts |
| | | (Γ) | S |
| | | $\Theta(\hat{\alpha}^+)$ | M |

| | | | | |
|-----------------------------|-----|---------------------------------|---|------------------------------------------------|
| | | $\Theta(\hat{\alpha}^-)$ | M | |
| Θ | ::= | | | unification type variable context |
| | | . | | empty context |
| | | $\vec{\alpha}\{\Delta\}$ | | from an ordered list of variables |
| | | $\hat{\alpha}^+\{\Delta\}$ | | from a variable to a list |
| | | $\overline{\Theta_i}^i$ | | concatenate contexts |
| | | (Θ) | S | |
| | | $\Theta _{vars}$ | | leave only those variables that are in the set |
| | | $\Theta_1 \cup \Theta_2$ | | |
| Ξ | ::= | | | anti-unification type variable context |
| | | . | | empty context |
| | | $\vec{\alpha}^-$ | | list of variables |
| | | $\overline{\Xi_i}^i$ | | concatenate contexts |
| | | (Ξ) | S | |
| | | $\Xi_1 \cup \Xi_2$ | | |
| | | $\Xi_1 \cap \Xi_2$ | | |
| | | $\Xi' _{vars}$ | M | |
| $\vec{\alpha}, \vec{\beta}$ | ::= | | | ordered positive or negative variables |
| | | . | | empty list |
| | | $\vec{\alpha}^+$ | | list of variables |
| | | $\vec{\alpha}^-$ | | list of variables |
| | | $\vec{\alpha}^\pm$ | | list of variables |
| | | $\vec{\alpha}^+$ | | list of variables |
| | | $\vec{\alpha}^-$ | | list of variables |
| | | $\vec{\alpha}_1 \setminus vars$ | | setminus |
| | | Γ | | context |
| | | $vars$ | | |
| | | $\vec{\alpha_i}^i$ | | concatenate contexts |
| | | $(\vec{\alpha})$ | S | parenthesis |
| | | $[\mu]\vec{\alpha}$ | M | apply moving to list |
| | | ord vars in P | M | |
| | | ord vars in N | M | |
| | | ord vars in P | M | |
| | | ord vars in N | M | |
| $vars$ | ::= | | | set of variables |
| | | \emptyset | | empty set |
| | | fv P | | free variables |
| | | fv N | | free variables |
| | | fv im P | | free variables |
| | | fv im N | | free variables |
| | | $vars_1 \cap vars_2$ | | set intersection |
| | | $vars_1 \cup vars_2$ | | set union |
| | | $vars_1 \setminus vars_2$ | | set complement |
| | | mv im P | | movable variables |
| | | mv im N | | movable variables |
| | | uv N | | unification variables |

| | | | | |
|-------------------------------------------------------|-----|------------------------------------------------|---|---------------------------------------------|
| | | uv P | | unification variables |
| | | fv N | | free variables |
| | | fv P | | free variables |
| | | $(vars)$ | S | parenthesis |
| | | $\vec{\alpha}$ | | ordered list of variables |
| | | $[\mu]vars$ | M | apply moving to varset |
| | | dom $(\hat{\sigma})$ | M | |
| | | dom $(\hat{\tau})$ | M | |
| | | dom (Θ) | M | |
| μ | ::= | | | |
| | | . | | empty moving |
| | | $pma1 \mapsto pma2$ | | Positive unit substitution |
| | | $nma1 \mapsto nma2$ | | Positive unit substitution |
| | | $\mu_1 \cup \mu_2$ | M | Set-like union of movings |
| | | $\mu_1 \circ \mu_2$ | M | Composition |
| | | $\overline{\mu_i}^i$ | | concatenate movings |
| | | $\mu _{vars}$ | M | restriction on a set |
| | | μ^{-1} | M | inversion |
| | | nf (μ') | M | |
| $\hat{\alpha}^\pm$ | ::= | | | positive/negative unification variable |
| | | $\hat{\alpha}^\pm$ | | |
| $\hat{\alpha}^+$ | ::= | | | positive unification variable |
| | | $\hat{\alpha}^+$ | | |
| | | $\hat{\alpha}^+\{\Delta\}$ | | |
| | | $\hat{\alpha}^\pm$ | | |
| $\hat{\alpha}^-, \hat{\beta}^-$ | ::= | | | negative unification variable |
| | | $\hat{\alpha}^-$ | | |
| | | $\hat{\alpha}_{\{N,M\}}^-$ | | |
| | | $\hat{\alpha}^-\{\Delta\}$ | | |
| | | $\hat{\alpha}^\pm$ | | |
| $\overrightarrow{\alpha^+}, \overrightarrow{\beta^+}$ | ::= | | | positive unification variable list |
| | | . | | empty list |
| | | $\hat{\alpha}^+$ | | a variable |
| | | $\overrightarrow{\hat{\alpha}^+}$ | | from a normal variable, context unspecified |
| | | $\overrightarrow{\overrightarrow{\alpha^+}}^i$ | | concatenate lists |
| $\overrightarrow{\alpha^-}, \overrightarrow{\beta^-}$ | ::= | | | negative unification variable list |
| | | . | | empty list |
| | | $\hat{\alpha}^-$ | | a variable |
| | | Ξ | | from an antiunification context |
| | | $\hat{\alpha}^-\{\Delta\}$ | | from a normal variable |
| | | $\overrightarrow{\hat{\alpha}^-}$ | | from a normal variable, context unspecified |
| | | $\overrightarrow{\overrightarrow{\alpha^-}}^i$ | | concatenate lists |

| | | | |
|-------------|-------|--------------------------------------------------------------|---|
| P, Q | $::=$ | a positive algorithmic type (potentially with metavariables) | |
| | | α^+ | |
| | | pma | |
| | | $\hat{\alpha}^+$ | |
| | | $\downarrow N$ | |
| | | $\xrightarrow{\quad} \exists \alpha^-. P$ | |
| | | $[\sigma]P$ | M |
| | | $[\hat{\tau}]P$ | M |
| | | $[\mu]P$ | M |
| | | (P) | S |
| | | nf (P') | M |
| N, M | $::=$ | a negative algorithmic type (potentially with metavariables) | |
| | | α^- | |
| | | $\hat{\alpha}^-$ | |
| | | $\uparrow P$ | |
| | | $P \rightarrow N$ | |
| | | $\xrightarrow{\quad} \forall \alpha^+. N$ | |
| | | $[\sigma]N$ | M |
| | | $[\hat{\tau}]N$ | M |
| | | $[\mu]N$ | M |
| | | (N) | S |
| | | nf (N') | M |
| $auSol$ | $::=$ | | |
| | | $(\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ | |
| | | $(\Xi, N, \hat{\tau}_1, \hat{\tau}_2)$ | |
| $terminals$ | $::=$ | | |
| | | \exists | |
| | | \forall | |
| | | \uparrow | |
| | | \downarrow | |
| | | \rightarrow | |
| | | \leftrightarrow | |
| | | \in | |
| | | \notin | |
| | | \cdot | |
| | | \top | |
| | | \leq | |
| | | \geq | |
| | | \sqsubset | |
| | | \subset | |
| | | \supset | |
| | | \diagdown | |
| | | \sqcup | |
| | | \mapsto | |
| | | \models^u | |
| | | \models^a | |
| | | \emptyset | |

| | | |
|--------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----------------------------------------------|
| | <div> \circ \Rightarrow \models \models \neq \equiv_n \vee \Downarrow $:\geq$ $:\simeq$ Λ λ \mathbf{let}^\exists \bullet $\Rightarrow\Rightarrow$ </div> | |
| v, w | $::=$ <div> x $\{c\}$ $(v : P)$ (v) </div> | value terms |
| \vec{v} | $::=$ <div> \cdot v \overrightarrow{v}_i^i </div> | list of arguments concatenate |
| c, d | $::=$ <div> $(c : N)$ $\lambda x : P. c$ $\Lambda \alpha^+. c$ $\mathbf{return} v$ $\mathbf{let} x : P = v(\vec{v}); c$ $\mathbf{let} x = v(\vec{v}); c$ $\mathbf{let}^\exists(\alpha^-, x) = v; c$ </div> | computation terms |
| $vctx, \Phi$ | $::=$ <div> \cdot $x : P$ $\overrightarrow{\Phi}_i^i$ </div> | variable context concatenate contexts |
| $formula$ | $::=$ <div> $judgement$ $judgement\ uniquely$ $formula_1 \dots formula_n$ $\mu : vars_1 \leftrightarrow vars_2$ $\mu \mathbf{is\ bijective}$ $\hat{\sigma} \mathbf{is\ functional}$ $\hat{\sigma}_1 \in \hat{\sigma}_2$ $v : P \in \Phi$ </div> | |

| | | |
|------|------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|----------------------------------------------------------------------------------------------------------------------------------------------------|
| | $\hat{\sigma}_1 \subseteq \hat{\sigma}_2$ $vars_1 \subseteq vars_2$ $vars_1 = vars_2$ $vars$ is fresh $\alpha^- \notin vars$ $\alpha^+ \notin vars$ $\alpha^- \in vars$ $\alpha^+ \in vars$ $\hat{\alpha}^+ \in vars$ $\hat{\alpha}^- \in vars$ $\hat{\alpha}^- \in \Theta$ $\hat{\alpha}^+ \in \Theta$ if any other rule is not applicable $\vec{\alpha}_1 = \vec{\alpha}_2$ $e_1 = e_2$ $N = M$ $N \neq M$ $P \neq Q$ $N \neq M$ $P \neq Q$ $P \neq Q$ $N \neq M$ | |
| A | $::=$ $\Gamma; \Theta \models N \leq M \Rightarrow \hat{\sigma}$ $\Gamma; \Theta \models P \geq Q \Rightarrow \hat{\sigma}$ | Negative subtyping Positive supertyping |
| AT | $::=$ $\Gamma; \Phi \models v : P$ $\Gamma; \Phi \models c : N$ $\Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \Rightarrow \hat{\sigma}$ | Positive type inference Negative type inference Application type inference |
| AU | $::=$ $\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \Rightarrow (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ $\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \Rightarrow (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$ | |
| $E1$ | $::=$ $N \stackrel{D}{\simeq}_1 M$ $P \stackrel{D}{\simeq}_1 Q$ $P \simeq Q$ | Negative multi-quantified type equivalence Positive multi-quantified type equivalence |
| $D1$ | $::=$ $\Gamma \vdash N \simeq_1^{\leq} M$ $\Gamma \vdash P \simeq_1^{\leq} Q$ $\Gamma \vdash N \leq_1 M$ $\Gamma \vdash P \geq_1 Q$ $\Gamma_2 \vdash \sigma_1 \simeq_1^{\leq} \sigma_2 : \Gamma_1$ | Negative equivalence on MQ types Positive equivalence on MQ types Negative subtyping Positive supertyping Equivalence of substitutions |
| $D0$ | $::=$ | |

| | | | |
|--------|-------|-------------------------------------------------------------|--------------------------------------------|
| | | $\Gamma \vdash N \simeq_0^{\leq} M$ | Negative equivalence |
| | | $\Gamma \vdash P \simeq_0^{\leq} Q$ | Positive equivalence |
| | | $\Gamma \vdash N \leq_0 M$ | Negative subtyping |
| | | $\Gamma \vdash P \geq_0 Q$ | Positive supertyping |
| DT | $::=$ | | |
| | | $\Gamma; \Phi \vdash v : P$ | Positive type inference |
| | | $\Gamma; \Phi \vdash c : N$ | Negative type inference |
| | | $\Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M$ | Spin Application type inference |
| EQ | $::=$ | | |
| | | $N = M$ | Negative type equality (alpha-equivalence) |
| | | $P = Q$ | Positive type equality (alpha-equivalence) |
| | | $\boxed{P} = \boxed{Q}$ | |
| $LUBF$ | $::=$ | | |
| | | $P_1 \vee P_2 === Q$ | |
| | | $\mathbf{ord\ vars\ in\ } \boxed{P} === \vec{\alpha}$ | |
| | | $\mathbf{ord\ vars\ in\ } \boxed{N} === \vec{\alpha}$ | |
| | | $\mathbf{ord\ vars\ in\ } P === \vec{\alpha}$ | |
| | | $\mathbf{ord\ vars\ in\ } N === \vec{\alpha}$ | |
| | | $\mathbf{nf}\ (N') === N$ | |
| | | $\mathbf{nf}\ (P') === P$ | |
| | | $\mathbf{nf}\ (\boxed{N'}) === \boxed{N}$ | |
| | | $\mathbf{nf}\ (\boxed{P'}) === \boxed{P}$ | |
| | | $\mathbf{nf}\ (\vec{N}') === \vec{N}$ | |
| | | $\mathbf{nf}\ (\vec{P}') === \vec{P}$ | |
| | | $\mathbf{nf}\ (\sigma') === \sigma$ | |
| | | $\mathbf{nf}\ (\mu') === \mu$ | |
| | | $\mathbf{nf}\ (\hat{\sigma}') === \hat{\sigma}$ | |
| | | $\sigma' _{vars}$ | |
| | | $\hat{\sigma}' _{vars}$ | |
| | | $\hat{\tau}' _{vars}$ | |
| | | $\Xi' _{vars}$ | |
| | | $e_1 \ \& \ e_2$ | |
| | | $\hat{\sigma}_1 \ \& \ \hat{\sigma}_2$ | |
| | | $\hat{\tau}_1 \ \& \ \hat{\tau}_2$ | |
| | | $\mathbf{dom}\ (\hat{\sigma}) === vars$ | |
| | | $\mathbf{dom}\ (\hat{\tau}) === vars$ | |
| | | $\mathbf{dom}\ (\Theta) === vars$ | |
| LUB | $::=$ | | |
| | | $\Gamma \models P_1 \vee P_2 = Q$ | Least Upper Bound (Least Common Supertype) |
| | | $\mathbf{upgrade}\ \Gamma \vdash P \mathbf{to}\ \Delta = Q$ | |
| Nrm | $::=$ | | |
| | | $\mathbf{nf}\ (N) = M$ | |
| | | $\mathbf{nf}\ (P) = Q$ | |
| | | $\mathbf{nf}\ (\boxed{N}) = \boxed{M}$ | |
| | | $\mathbf{nf}\ (\boxed{P}) = \boxed{Q}$ | |

| | | |
|------------------|----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| <i>Order</i> | $::=$ $\begin{array}{ l} \text{ord vars in } N = \vec{\alpha} \\ \text{ord vars in } P = \vec{\alpha} \\ \text{ord vars in } \boxed{N} = \vec{\alpha} \\ \text{ord vars in } \boxed{P} = \vec{\alpha} \end{array}$ | |
| <i>SM</i> | $::=$ $\begin{array}{ l} \Gamma \vdash e_1 \ \& \ e_2 = e_3 \\ \Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3 \end{array}$ | Unification Solution Entry Merge Merge unification solutions |
| <i>SImp</i> | $::=$ $\begin{array}{ l} \Gamma \vdash e_1 \Rightarrow e_2 \\ \Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2 \\ \Gamma \vdash e_1 \simeq e_2 \\ \Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2 \end{array}$ | Weakening of unification solution entries Weakening of unification solutions |
| <i>U</i> | $::=$ $\begin{array}{ l} \Gamma; \Theta \models \boxed{N} \stackrel{u}{\simeq} M \Rightarrow \hat{\sigma} \\ \Gamma; \Theta \models \boxed{P} \stackrel{u}{\simeq} Q \Rightarrow \hat{\sigma} \end{array}$ | Negative unification Positive unification |
| <i>WF</i> | $::=$ $\begin{array}{ l} \Gamma \vdash N \\ \Gamma \vdash P \\ \Gamma \vdash N \\ \Gamma \vdash P \\ \Gamma \vdash \vec{N} \\ \Gamma \vdash \vec{P} \\ \Gamma; \Theta \vdash \boxed{N} \\ \Gamma; \Theta \vdash \boxed{P} \\ \Gamma; \Xi \vdash \boxed{N} \\ \Gamma; \Xi \vdash \boxed{P} \\ \Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1 \\ \hat{\sigma} : \Theta \\ \Gamma \vdash^{\supset} \Theta \\ \Gamma_1 \vdash \sigma : \Gamma_2 \\ \Gamma \vdash e \end{array}$ | Negative type well-formedness Positive type well-formedness Negative type well-formedness Positive type well-formedness Negative type list well-formedness Positive type list well-formedness Negative unification type well-formedness Positive unification type well-formedness Negative anti-unification type well-formedness Positive anti-unification type well-formedness Antiunification substitution well-formedness Unification substitution well-formedness Unification context well-formedness Substitution well-formedness Unification solution entry well-formedness |
| <i>judgement</i> | $::=$ $\begin{array}{ l} A \\ AT \\ AU \\ E1 \\ D1 \\ D0 \\ DT \\ EQ \\ LUB \\ Nrm \\ Order \\ SM \end{array}$ | |

| | | |
|----------------|-------|----------------------------------------------|
| | | $SImp$ |
| | | U |
| | | WF |
| $user_syntax$ | $::=$ | |
| | | α |
| | | n |
| | | x |
| | | n |
| | | α^+ |
| | | α^- |
| | | α^\pm |
| | | σ |
| | | e |
| | | $\hat{\sigma}$ |
| | | $\hat{\tau}$ |
| | | P |
| | | N |
| | | $\overrightarrow{\alpha^+}$ |
| | | $\overrightarrow{\alpha^-}$ |
| | | $\overrightarrow{\alpha^\pm}$ |
| | | P |
| | | N |
| | | \vec{P} |
| | | \vec{N} |
| | | Γ |
| | | Θ |
| | | Ξ |
| | | $\vec{\alpha}$ |
| | | $vars$ |
| | | μ |
| | | $\hat{\alpha}^\pm$ |
| | | $\hat{\alpha}^+$ |
| | | $\hat{\alpha}^-$ |
| | | $\widetilde{\overrightarrow{\alpha^+}}$ |
| | | $\overrightarrow{\overrightarrow{\alpha^+}}$ |
| | | $\overrightarrow{\overrightarrow{\alpha^-}}$ |
| | | \boxed{P} |
| | | \boxed{N} |
| | | $auSol$ |
| | | $terminals$ |
| | | v |
| | | \vec{v} |
| | | c |
| | | $vctx$ |
| | | $formula$ |

$\boxed{\Gamma; \Theta \models N \leqslant M \Rightarrow \hat{\sigma}}$

Negative subtyping

$\overline{\Gamma; \Theta \models \alpha^- \leqslant \alpha^- \Rightarrow} \quad \text{ANVAR}$

$$\begin{array}{c}
\frac{\Gamma; \Theta \models \mathbf{nf}(P) \stackrel{u}{\simeq} \mathbf{nf}(Q) \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \uparrow P \leq \uparrow Q \Rightarrow \hat{\sigma}} \quad \text{AShiftU} \\
\\
\frac{\Gamma; \Theta \models P \geq Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \Theta \models N \leq M \Rightarrow \hat{\sigma}_2 \quad \Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}}{\Gamma; \Theta \models P \rightarrow N \leq Q \rightarrow M \Rightarrow \hat{\sigma}} \quad \text{AArrow} \\
\\
\frac{\text{<<multiple parses>>}}{\Gamma; \Theta \models \forall \alpha^+. N \leq \forall \beta^+. M \Rightarrow \hat{\sigma} \setminus \hat{\alpha}^+} \quad \text{AForall} \\
\\
\boxed{\Gamma; \Theta \models P \geq Q \Rightarrow \hat{\sigma}} \quad \text{Positive supertyping} \\
\\
\frac{}{\Gamma; \Theta \models \alpha^+ \geq \alpha^+ \Rightarrow \cdot} \quad \text{APVar} \\
\\
\frac{\Gamma; \Theta \models \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \downarrow N \geq \downarrow M \Rightarrow \hat{\sigma}} \quad \text{AShiftD} \\
\\
\frac{\Gamma, \vec{\beta}^-; \Theta, \hat{\alpha}^- \{ \Gamma, \vec{\beta}^- \} \models [\hat{\alpha}^- / \alpha^-] P \geq Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \exists \alpha^-. P \geq \exists \beta^-. Q \Rightarrow \hat{\sigma} \setminus \hat{\alpha}^-} \quad \text{AExists} \\
\\
\frac{\hat{\alpha}^+ \{ \Delta \} \in \Theta \quad \mathbf{upgrade} \Gamma \vdash P \text{ to } \Delta = Q}{\Gamma; \Theta \models \hat{\alpha}^+ \geq P \Rightarrow (\hat{\alpha}^+ : \geq Q)} \quad \text{APUVar} \\
\\
\boxed{\Gamma; \Phi \models v : P} \quad \text{Positive type inference} \\
\\
\frac{v : P \in \Phi}{\Gamma; \Phi \models v : P} \quad \text{ATVar} \\
\\
\frac{\Gamma; \Phi \models c : N}{\Gamma; \Phi \models \{c\} : \downarrow N} \quad \text{ATThunk} \\
\\
\frac{\Gamma; \Phi \models v : P \quad \Gamma; \cdot \models Q \geq P \Rightarrow \cdot}{\Gamma; \Phi \models (v : Q) : Q} \quad \text{ATAnnot} \\
\\
\boxed{\Gamma; \Phi \models c : N} \quad \text{Negative type inference} \\
\\
\frac{\Gamma; \Phi \models c : N \quad \Gamma; \cdot \models N \leq M \Rightarrow \cdot}{\Gamma; \Phi \models (c : M) : M} \quad \text{ATAnnotN} \\
\\
\frac{\Gamma; \Phi, x : P \models c : N}{\Gamma; \Phi \models \lambda x : P. c : P \rightarrow N} \quad \text{ATTLam} \\
\\
\frac{\Gamma, \alpha^+; \Phi \models c : N}{\Gamma; \Phi \models \Lambda \alpha^+. c : \forall \alpha^+. N} \quad \text{ATTlam} \\
\\
\frac{\Gamma; \Phi \models v : P}{\Gamma; \Phi \models \mathbf{return} v : \uparrow P} \quad \text{ATReturn} \\
\\
\frac{\Gamma; \Phi \models v : \downarrow M \quad \Gamma; \Phi; \cdot \models M \bullet \vec{v} \Rightarrow \uparrow Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \mathbf{uv} Q \{ \Gamma \} \models \uparrow Q \leq \uparrow P \Rightarrow \hat{\sigma}_2 \quad \mathbf{uv} Q \{ \Gamma \} \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma} \quad \Gamma; \Phi, x : P \models c : N}{\Gamma; \Phi \models \mathbf{let} x : P = v(\vec{v}); c : N} \quad \text{ATLetAnn} \\
\\
\frac{\Gamma; \Phi \models v : \downarrow M \quad \Gamma; \Phi; \cdot \models M \bullet \vec{v} \Rightarrow \uparrow Q \Rightarrow \hat{\sigma} \quad \mathbf{uv}(Q) = \emptyset \quad \Gamma; \Phi, x : Q \models c : N}{\Gamma; \Phi \models \mathbf{let} x = v(\vec{v}); c : N} \quad \text{ATLet} \\
\\
\frac{\Gamma, \alpha^-; \Phi \models v : \exists \alpha^-. P \quad \Gamma, \alpha^-; \Phi, x : P \models c : N \quad \Gamma \vdash N}{\Gamma; \Phi \models \mathbf{let}^3(\alpha^-, x) = v; c : N} \quad \text{ATUnpack} \\
\\
\boxed{\Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \Rightarrow \hat{\sigma}} \quad \text{Application type inference}
\end{array}$$

$$\begin{array}{c}
\frac{N \neq \forall \alpha^+ . M}{\Gamma; \Phi; \Theta \models N \bullet \cdot \Rightarrow N \models \cdot} \text{ATEMPTY} \\
\frac{\Gamma; \Phi \models v : P \quad \Gamma; \Theta \models Q \succcurlyeq P \models \hat{\sigma}_1 \quad \Gamma; \Phi; \Theta \models N \bullet \vec{v} \Rightarrow M \models \hat{\sigma}_2}{\Gamma; \Phi; \Theta \models Q \rightarrow N \bullet v, \vec{v} \Rightarrow M \models \hat{\sigma}_1 \& \hat{\sigma}_2} \text{ATARROW} \\
\frac{\text{<<multiple parses>>}}{\Gamma; \Phi; \Theta \models \forall \alpha^+ . N \bullet \vec{v} \Rightarrow M \models \hat{\sigma}} \text{ATFORALL} \\
\boxed{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)} \\
\frac{}{\Gamma \models \alpha^+ \stackrel{a}{\simeq} \alpha^+ \models (\cdot, \alpha^+, \cdot, \cdot)} \text{AUPVAR} \\
\frac{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N_1 \stackrel{a}{\simeq} \downarrow N_2 \models (\Xi, \downarrow M, \hat{\tau}_1, \hat{\tau}_2)} \text{AUSHIFTD} \\
\frac{\vec{\alpha}^- \cap \Gamma = \emptyset \quad \Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \exists \alpha^+ . P_1 \stackrel{a}{\simeq} \exists \alpha^+ . P_2 \models (\Xi, \exists \alpha^+ . Q, \hat{\tau}_1, \hat{\tau}_2)} \text{AUEXISTS} \\
\boxed{\Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)} \\
\frac{}{\Gamma \models \alpha^- \stackrel{a}{\simeq} \alpha^- \models (\cdot, \alpha^-, \cdot, \cdot)} \text{AUNVAR} \\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \uparrow P_1 \stackrel{a}{\simeq} \uparrow P_2 \models (\Xi, \uparrow Q, \hat{\tau}_1, \hat{\tau}_2)} \text{AUSHIFTU} \\
\frac{\vec{\alpha}^+ \cap \Gamma = \emptyset \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \forall \alpha^+ . N_1 \stackrel{a}{\simeq} \forall \alpha^+ . N_2 \models (\Xi, \forall \alpha^+ . M, \hat{\tau}_1, \hat{\tau}_2)} \text{AUFORALL} \\
\frac{\Gamma \models P_1 \stackrel{a}{\simeq} P_2 \models (\Xi_1, Q, \hat{\tau}_1, \hat{\tau}_2) \quad \Gamma \models N_1 \stackrel{a}{\simeq} N_2 \models (\Xi_2, M, \hat{\tau}'_1, \hat{\tau}'_2)}{\Gamma \models P_1 \rightarrow N_1 \stackrel{a}{\simeq} P_2 \rightarrow N_2 \models (\Xi_1 \cup \Xi_2, Q \rightarrow M, \hat{\tau}_1 \cup \hat{\tau}'_1, \hat{\tau}_2 \cup \hat{\tau}'_2)} \text{AUARROW} \\
\frac{\text{if any other rule is not applicable} \quad \Gamma \vdash N \quad \Gamma \vdash M}{\Gamma \models N \stackrel{a}{\simeq} M \models (\hat{\alpha}_{\{N, M\}}^-, \hat{\alpha}_{\{N, M\}}^-, (\hat{\alpha}_{\{N, M\}}^- : \approx N), (\hat{\alpha}_{\{N, M\}}^- : \approx M))} \text{AUAU} \\
\boxed{N \simeq_1^D M} \quad \text{Negative multi-quantified type equivalence}
\end{array}$$

$$\begin{array}{c}
\frac{}{\alpha^- \simeq_1^D \alpha^-} \text{E1NVAR} \\
\frac{P \simeq_1^D Q}{\uparrow P \simeq_1^D \uparrow Q} \text{E1SHIFTU} \\
\frac{P \simeq_1^D Q \quad N \simeq_1^D M}{P \rightarrow N \simeq_1^D Q \rightarrow M} \text{E1ARROW} \\
\frac{\vec{\alpha}^+ \cap \mathbf{fv} M = \emptyset \quad \mu : (\vec{\beta}^+ \cap \mathbf{fv} M) \leftrightarrow (\vec{\alpha}^+ \cap \mathbf{fv} N) \quad N \simeq_1^D [\mu]M}{\forall \alpha^+ . N \simeq_1^D \forall \beta^+ . M} \text{E1FORALL}
\end{array}$$

$$\boxed{P \simeq_1^D Q} \quad \text{Positive multi-quantified type equivalence}$$

$$\frac{}{\alpha^+ \simeq_1^D \alpha^+} \text{E1PVAR}$$

$$\begin{array}{c}
\frac{N \simeq_1^D M}{\downarrow N \simeq_1^D \downarrow M} \quad \text{E1SHIFTD} \\
\frac{\overrightarrow{\alpha^-} \cap \mathbf{fv} Q = \emptyset \quad \mu : (\overrightarrow{\beta^-} \cap \mathbf{fv} Q) \leftrightarrow (\overrightarrow{\alpha^-} \cap \mathbf{fv} P) \quad P \simeq_1^D [\mu]Q}{\exists \overrightarrow{\alpha^-}. P \simeq_1^D \exists \overrightarrow{\beta^-}. Q} \quad \text{E1EXISTS} \\
\\
\boxed{P \simeq Q} \\
\boxed{\Gamma \vdash N \simeq_1^{\leq} M} \quad \text{Negative equivalence on MQ types} \\
\\
\frac{\Gamma \vdash N \leq_1 M \quad \Gamma \vdash M \leq_1 N}{\Gamma \vdash N \simeq_1^{\leq} M} \quad \text{D1NDEF} \\
\\
\boxed{\Gamma \vdash P \simeq_1^{\leq} Q} \quad \text{Positive equivalence on MQ types} \\
\\
\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash Q \geq_1 P}{\Gamma \vdash P \simeq_1^{\leq} Q} \quad \text{D1PDEF} \\
\\
\boxed{\Gamma \vdash N \leq_1 M} \quad \text{Negative subtyping} \\
\\
\frac{}{\Gamma \vdash \alpha^- \leq_1 \alpha^-} \quad \text{D1NVAR} \\
\\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \uparrow P \leq_1 \uparrow Q} \quad \text{D1SHIFTU} \\
\\
\frac{\Gamma \vdash P \geq_1 Q \quad \Gamma \vdash N \leq_1 M}{\Gamma \vdash P \rightarrow N \leq_1 Q \rightarrow M} \quad \text{D1ARROW} \\
\\
\frac{\mathbf{fv} N \cap \overrightarrow{\beta^+} = \emptyset \quad \Gamma, \overrightarrow{\beta^+} \vdash P_i \quad \Gamma, \overrightarrow{\beta^+} \vdash [\overrightarrow{P}/\overrightarrow{\alpha^+}]N \leq_1 M}{\Gamma \vdash \forall \overrightarrow{\alpha^+}. N \leq_1 \forall \overrightarrow{\beta^+}. M} \quad \text{D1FORALL} \\
\\
\boxed{\Gamma \vdash P \geq_1 Q} \quad \text{Positive supertyping} \\
\\
\frac{}{\Gamma \vdash \alpha^+ \geq_1 \alpha^+} \quad \text{D1PVAR} \\
\\
\frac{\text{<<multiple parses>>}}{\Gamma \vdash \downarrow N \geq_1 \downarrow M} \quad \text{D1SHIFTD} \\
\\
\frac{\mathbf{fv} P \cap \overrightarrow{\beta^-} = \emptyset \quad \Gamma, \overrightarrow{\beta^-} \vdash N_i \quad \Gamma, \overrightarrow{\beta^-} \vdash [\overrightarrow{N}/\overrightarrow{\alpha^-}]P \geq_1 Q}{\Gamma \vdash \exists \overrightarrow{\alpha^-}. P \geq_1 \exists \overrightarrow{\beta^-}. Q} \quad \text{D1EXISTS} \\
\\
\boxed{\Gamma_2 \vdash \sigma_1 \simeq_1^{\leq} \sigma_2 : \Gamma_1} \quad \text{Equivalence of substitutions} \\
\boxed{\Gamma \vdash N \simeq_0^{\leq} M} \quad \text{Negative equivalence} \\
\\
\frac{\Gamma \vdash N \leq_0 M \quad \Gamma \vdash M \leq_0 N}{\Gamma \vdash N \simeq_0^{\leq} M} \quad \text{D0NDEF} \\
\\
\boxed{\Gamma \vdash P \simeq_0^{\leq} Q} \quad \text{Positive equivalence} \\
\\
\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash Q \geq_0 P}{\Gamma \vdash P \simeq_0^{\leq} Q} \quad \text{D0PDEF} \\
\\
\boxed{\Gamma \vdash N \leq_0 M} \quad \text{Negative subtyping} \\
\\
\frac{}{\Gamma \vdash \alpha^- \leq_0 \alpha^-} \quad \text{D0NVAR} \\
\\
\frac{\Gamma \vdash P \simeq_0^{\leq} Q}{\Gamma \vdash \uparrow P \leq_0 \uparrow Q} \quad \text{D0SHIFTU}
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leq_0 M \quad M \neq \forall \beta^+.M'}{\Gamma \vdash \forall \alpha^+.N \leq_0 M} \quad \text{D0FORALLL} \\
\frac{\Gamma, \alpha^+ \vdash N \leq_0 M}{\Gamma \vdash N \leq_0 \forall \alpha^+.M} \quad \text{D0FORALLR} \\
\frac{\Gamma \vdash P \geq_0 Q \quad \Gamma \vdash N \leq_0 M}{\Gamma \vdash P \rightarrow N \leq_0 Q \rightarrow M} \quad \text{D0ARROW}
\end{array}$$

$\boxed{\Gamma \vdash P \geq_0 Q}$ Positive supertyping

$$\begin{array}{c}
\frac{}{\Gamma \vdash \alpha^+ \geq_0 \alpha^+} \quad \text{D0PVAR} \\
\frac{\Gamma \vdash N \simeq_0^{\leq} M}{\Gamma \vdash \downarrow N \geq_0 \downarrow M} \quad \text{D0SHIFTD} \\
\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^-]P \geq_0 Q \quad Q \neq \exists \alpha^-.Q'}{\Gamma \vdash \exists \alpha^-.P \geq_0 Q} \quad \text{D0EXISTS L} \\
\frac{\Gamma, \alpha^- \vdash P \geq_0 Q}{\Gamma \vdash P \geq_0 \exists \alpha^-.Q} \quad \text{D0EXISTS R}
\end{array}$$

$\boxed{\Gamma; \Phi \vdash v : P}$ Positive type inference

$$\begin{array}{c}
\frac{v : P \in \Phi}{\Gamma; \Phi \vdash v : P} \quad \text{DTVAR} \\
\frac{\Gamma; \Phi \vdash c : N}{\Gamma; \Phi \vdash \{c\} : \downarrow N} \quad \text{DTTHUNK} \\
\frac{\Gamma; \Phi \vdash v : P \quad \Gamma \vdash Q \geq_1 P}{\Gamma; \Phi \vdash (v : Q) : Q} \quad \text{DTANNOTP}
\end{array}$$

$\boxed{\Gamma; \Phi \vdash c : N}$ Negative type inference

$$\begin{array}{c}
\frac{\Gamma; \Phi, x : P \vdash c : N}{\Gamma; \Phi \vdash \lambda x : P. c : P \rightarrow N} \quad \text{DTTLAM} \\
\frac{\Gamma, \alpha^+; \Phi \vdash c : N}{\Gamma; \Phi \vdash \Lambda \alpha^+. c : \forall \alpha^+. N} \quad \text{DTTLAM} \\
\frac{\Gamma; \Phi \vdash v : P}{\Gamma; \Phi \vdash \text{return } v : \uparrow P} \quad \text{DTRETURN} \\
\frac{\Gamma; \Phi \vdash v : \downarrow M \quad \Gamma; \Phi \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \quad \Gamma \vdash \uparrow Q \leq_1 \uparrow P \quad \Gamma; \Phi, x : P \vdash c : N}{\Gamma; \Phi \vdash \text{let } x : P = v(\vec{v}); c : N} \quad \text{DTLETANN} \\
\frac{\Gamma; \Phi \vdash v : \downarrow M \quad \Gamma; \Phi \vdash M \bullet \vec{v} \Rightarrow \uparrow Q \text{ uniquely} \quad \Gamma; \Phi, x : Q \vdash c : N}{\Gamma; \Phi \vdash \text{let } x = v(\vec{v}); c : N} \quad \text{DTLET} \\
\frac{\Gamma, \alpha^-; \Phi \vdash v : \exists \alpha^-. P \quad \Gamma, \alpha^-; \Phi, x : P \vdash c : N \quad \Gamma \vdash N}{\Gamma; \Phi \vdash \text{let}^{\exists}(\alpha^-, x) = v; c : N} \quad \text{DTUNPACK} \\
\frac{\Gamma; \Phi \vdash c : N \quad \Gamma \vdash N \leq_1 M}{\Gamma; \Phi \vdash (c : M) : M} \quad \text{DTANNOTN}
\end{array}$$

$\boxed{\Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M}$ Spin Application type inference

$$\frac{N \neq \forall \alpha^+. M}{\Gamma; \Phi \vdash N \bullet \cdot \Rightarrow N} \quad \text{DTEMTPTY}$$

$$\frac{\Gamma; \Phi \vdash v : P \quad \Gamma \vdash Q \geq_1 P \quad \Gamma; \Phi \vdash N \bullet \vec{v} \Rightarrow M}{\Gamma; \Phi \vdash Q \rightarrow N \bullet v, \vec{v} \Rightarrow M} \text{DTARROW}$$

$$\frac{\Gamma \vdash \vec{P} \quad \Gamma; \Phi \vdash [\vec{P}/\alpha^+] N \bullet \vec{v} \Rightarrow M}{\Gamma; \Phi \vdash \overrightarrow{\forall \alpha^+}. N \bullet \vec{v} \Rightarrow M} \text{DTFORALL}$$

$\boxed{N = M}$ Negative type equality (alpha-equivalence)

$\boxed{P = Q}$ Positive type equality (alpha-equivalence)

$\boxed{P = Q}$

$\boxed{P_1 \vee P_2}$

$\boxed{\text{ord vars in } P}$

$\boxed{\text{ord vars in } N}$

$\boxed{\text{ord vars in } P}$

$\boxed{\text{ord vars in } N}$

$\boxed{\text{nf } (N')}$

$\boxed{\text{nf } (P')}$

$\boxed{\text{nf } (N')}$

$\boxed{\text{nf } (P')}$

$\boxed{\text{nf } (\vec{N}')$

$\boxed{\text{nf } (\vec{P}')$

$\boxed{\text{nf } (\sigma')}$

$\boxed{\text{nf } (\mu')}$

$\boxed{\mathbf{nf}(\hat{\sigma}')}$
 $\boxed{\sigma'|_{vars}}$
 $\boxed{\hat{\sigma}'|_{vars}}$
 $\boxed{\hat{\tau}'|_{vars}}$
 $\boxed{\Xi'|_{vars}}$
 $\boxed{e_1 \ \& \ e_2}$
 $\boxed{\hat{\sigma}_1 \ \& \ \hat{\sigma}_2}$
 $\boxed{\hat{\tau}_1 \ \& \ \hat{\tau}_2}$
 $\boxed{\mathbf{dom}(\hat{\sigma})}$
 $\boxed{\mathbf{dom}(\hat{\tau})}$
 $\boxed{\mathbf{dom}(\Theta)}$
 $\boxed{\Gamma \models P_1 \vee P_2 = Q}$ Least Upper Bound (Least Common Supertype)

$$\begin{array}{c}
\overline{\Gamma \models \alpha^+ \vee \alpha^+ = \alpha^+} \quad \text{LUBVAR} \\
\frac{\Gamma, \cdot \models \mathbf{nf}(\downarrow N) \stackrel{a}{\simeq} \mathbf{nf}(\downarrow M) \models (\Xi, \mathbf{P}, \hat{\tau}_1, \hat{\tau}_2)}{\Gamma \models \downarrow N \vee \downarrow M = \exists \overrightarrow{\alpha^-}. [\overrightarrow{\alpha^-} / \Xi] \mathbf{P}} \quad \text{LUBSHIFT} \\
\frac{\Gamma, \overrightarrow{\alpha^-}, \overrightarrow{\beta^-} \models P_1 \vee P_2 = Q}{\Gamma \models \exists \overrightarrow{\alpha^-}. P_1 \vee \exists \overrightarrow{\beta^-}. P_2 = Q} \quad \text{LUBEXISTS}
\end{array}$$

 $\boxed{\mathbf{upgrade} \Gamma \vdash P \text{ to } \Delta = Q}$

$$\frac{\Gamma = \Delta, \vec{\alpha}^\pm \quad \vec{\beta}^\pm \text{ is fresh} \quad \vec{\gamma}^\pm \text{ is fresh} \quad \Delta, \vec{\beta}^\pm, \vec{\gamma}^\pm \models [\vec{\beta}^\pm / \vec{\alpha}^\pm]P \vee [\vec{\gamma}^\pm / \vec{\alpha}^\pm]P = Q}{\text{upgrade } \Gamma \vdash P \text{ to } \Delta = Q} \quad \text{LUBUPGRADE}$$

$$\boxed{\text{nf}(N) = M}$$

$$\frac{}{\text{nf}(\alpha^-) = \alpha^-} \quad \text{NRMNVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\uparrow P) = \uparrow Q} \quad \text{NRMSHIFTU}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(P \rightarrow N) = Q \rightarrow M} \quad \text{NRMARROW}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\forall \alpha^+. N) = \forall \alpha^{+'}. N'} \quad \text{NRMFORALL}$$

$$\boxed{\text{nf}(P) = Q}$$

$$\frac{}{\text{nf}(\alpha^+) = \alpha^+} \quad \text{NRMPVAR}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\downarrow N) = \downarrow M} \quad \text{NRMSHIFTD}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\text{nf}(\exists \alpha^-. P) = \exists \alpha^{-'}. P'} \quad \text{NRME EXISTS}$$

$$\boxed{\text{nf}(N) = M}$$

$$\frac{}{\text{nf}(\hat{\alpha}^-) = \hat{\alpha}^-} \quad \text{NRMNUVAR}$$

$$\boxed{\text{nf}(P) = Q}$$

$$\frac{}{\text{nf}(\hat{\alpha}^+) = \hat{\alpha}^+} \quad \text{NRMPUVAR}$$

$$\boxed{\text{ord vars in } N = \vec{\alpha}}$$

$$\frac{\alpha^- \in \text{vars}}{\text{ord vars in } \alpha^- = \alpha^-} \quad \text{ONVARIN}$$

$$\frac{\alpha^- \notin \text{vars}}{\text{ord vars in } \alpha^- = .} \quad \text{ONVARNIN}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}}{\text{ord vars in } \uparrow P = \vec{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\text{ord vars in } P = \vec{\alpha}_1 \quad \text{ord vars in } N = \vec{\alpha}_2}{\text{ord vars in } P \rightarrow N = \vec{\alpha}_1, (\vec{\alpha}_2 \setminus \vec{\alpha}_1)} \quad \text{OARROW}$$

$$\frac{\text{vars} \cap \vec{\alpha}^+ = \emptyset \quad \text{ord vars in } N = \vec{\alpha}}{\text{ord vars in } \forall \alpha^+. N = \vec{\alpha}} \quad \text{OFORALL}$$

$$\boxed{\text{ord vars in } P = \vec{\alpha}}$$

$$\frac{\alpha^+ \in \text{vars}}{\text{ord vars in } \alpha^+ = \alpha^+} \quad \text{OPVARIN}$$

$$\frac{\alpha^+ \notin \text{vars}}{\mathbf{ord\,vars\,in}\,\alpha^+ = \cdot} \quad \text{OPVARNIN}$$

$$\frac{\mathbf{ord\,vars\,in}\,N = \vec{\alpha}}{\mathbf{ord\,vars\,in}\,\downarrow N = \vec{\alpha}} \quad \text{OSHIFTD}$$

$$\frac{\text{vars} \cap \vec{\alpha}^- = \emptyset \quad \mathbf{ord\,vars\,in}\,P = \vec{\alpha}}{\mathbf{ord\,vars\,in}\,\exists \alpha^-. P = \vec{\alpha}} \quad \text{OEXISTS}$$

$$\boxed{\mathbf{ord\,vars\,in}\,N = \vec{\alpha}}$$

$$\frac{}{\mathbf{ord\,vars\,in}\,\hat{\alpha}^- = \cdot} \quad \text{ONUVAR}$$

$$\boxed{\mathbf{ord\,vars\,in}\,P = \vec{\alpha}}$$

$$\frac{}{\mathbf{ord\,vars\,in}\,\hat{\alpha}^+ = \cdot} \quad \text{OPUVAR}$$

$$\boxed{\Gamma \vdash e_1 \& e_2 = e_3} \quad \text{Unification Solution Entry Merge}$$

$$\frac{\Gamma \vdash P_1 \vee P_2 = Q}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \& (\hat{\alpha}^+ : \geq P_2) = (\hat{\alpha}^+ : \geq Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vdash P \succcurlyeq Q \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \geq Q) = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vdash Q \succcurlyeq P \Rightarrow \hat{\sigma}'}{\Gamma \vdash (\hat{\alpha}^+ : \geq P) \& (\hat{\alpha}^+ : \approx Q) = (\hat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^+ : \approx P) \& (\hat{\alpha}^+ : \approx P') = (\hat{\alpha}^+ : \approx P)} \quad \text{SMEPEQEQ}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \& (\hat{\alpha}^- : \approx N') = (\hat{\alpha}^- : \approx N)} \quad \text{SMENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3} \quad \text{Merge unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \Rightarrow e_2} \quad \text{Weakening of unification solution entries}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPESUPSUP}$$

$$\frac{\Gamma \vdash P_1 \geq_1 P_2}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUP}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \Rightarrow (\hat{\alpha}^+ : \approx P_2)} \quad \text{SIMPEPEQEQ}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \Rightarrow (\hat{\alpha}^- : \approx N_2)} \quad \text{SIMPENEQEQ}$$

$$\boxed{\Theta \vdash \hat{\sigma}_1 \Rightarrow \hat{\sigma}_2} \quad \text{Weakening of unification solutions}$$

$$\boxed{\Gamma \vdash e_1 \simeq e_2}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^+ : \geq P_1) \simeq (\hat{\alpha}^+ : \geq P_2)} \quad \text{SIMPEEQSUPSUP}$$

$$\frac{\langle\langle \text{multiple parses} \rangle\rangle}{\Gamma \vdash (\hat{\alpha}^+ : \approx P_1) \simeq (\hat{\alpha}^+ : \approx P_2)} \quad \text{SIMPEEQPEQEQ}$$

| | | |
|-----------------------------------------------------------------------------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|--------------|
| | $\frac{\text{<<multiple parses>>}}{\Gamma \vdash (\hat{\alpha}^- : \approx N_1) \simeq (\hat{\alpha}^- : \approx N_2)}$ | SIMPEEQNEQEQ |
| $\boxed{\Theta \vdash \hat{\sigma}_1 \simeq \hat{\sigma}_2}$ | | |
| $\boxed{\Gamma; \Theta \models N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}$ | Negative unification | |
| | $\frac{}{\Gamma; \Theta \models \alpha^- \overset{u}{\simeq} \alpha^- \Rightarrow \cdot}$ | UNVAR |
| | $\frac{\Gamma; \Theta \models P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \uparrow P \overset{u}{\simeq} \uparrow Q \Rightarrow \hat{\sigma}}$ | USHIFTU |
| | $\frac{\Gamma; \Theta \models P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}_1 \quad \Gamma; \Theta \models N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}_2}{\Gamma; \Theta \models P \rightarrow N \overset{u}{\simeq} Q \rightarrow M \Rightarrow \hat{\sigma}_1 \ \& \ \hat{\sigma}_2}$ | UARROW |
| | $\frac{\Gamma, \overrightarrow{\alpha^+}; \Theta \models N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \forall \alpha^+. N \overset{u}{\simeq} \forall \alpha^+. M \Rightarrow \hat{\sigma}}$ | UFORALL |
| | $\frac{\hat{\alpha}^- \{ \Delta \} \in \Theta \quad \Delta \vdash N}{\Gamma; \Theta \models \hat{\alpha}^- \overset{u}{\simeq} N \Rightarrow (\hat{\alpha}^- : \approx N)}$ | UNUVAR |
| $\boxed{\Gamma; \Theta \models P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}$ | Positive unification | |
| | $\frac{}{\Gamma; \Theta \models \alpha^+ \overset{u}{\simeq} \alpha^+ \Rightarrow \cdot}$ | UPVAR |
| | $\frac{\Gamma; \Theta \models N \overset{u}{\simeq} M \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \downarrow N \overset{u}{\simeq} \downarrow M \Rightarrow \hat{\sigma}}$ | USHIFTD |
| | $\frac{\Gamma, \overrightarrow{\alpha^-}; \Theta \models P \overset{u}{\simeq} Q \Rightarrow \hat{\sigma}}{\Gamma; \Theta \models \exists \alpha^-. P \overset{u}{\simeq} \exists \alpha^-. Q \Rightarrow \hat{\sigma}}$ | UEXISTS |
| | $\frac{\hat{\alpha}^+ \{ \Delta \} \in \Theta \quad \Delta \vdash P}{\Gamma; \Theta \models \hat{\alpha}^+ \overset{u}{\simeq} P \Rightarrow (\hat{\alpha}^+ : \approx P)}$ | UPUVAR |
| $\boxed{\Gamma \vdash N}$ | Negative type well-formedness | |
| $\boxed{\Gamma \vdash P}$ | Positive type well-formedness | |
| $\boxed{\Gamma \vdash N}$ | Negative type well-formedness | |
| $\boxed{\Gamma \vdash P}$ | Positive type well-formedness | |
| $\boxed{\Gamma \vdash \overrightarrow{N}}$ | Negative type list well-formedness | |
| $\boxed{\Gamma \vdash \overrightarrow{P}}$ | Positive type list well-formedness | |
| $\boxed{\Gamma; \Theta \vdash N}$ | Negative unification type well-formedness | |
| $\boxed{\Gamma; \Theta \vdash P}$ | Positive unification type well-formedness | |
| $\boxed{\Gamma; \Xi \vdash N}$ | Negative anti-unification type well-formedness | |
| $\boxed{\Gamma; \Xi \vdash P}$ | Positive anti-unification type well-formedness | |
| $\boxed{\Gamma; \Xi_2 \vdash \hat{\tau} : \Xi_1}$ | Antiunification substitution well-formedness | |
| $\boxed{\hat{\sigma} : \Theta}$ | Unification substitution well-formedness | |
| $\boxed{\Gamma \vdash^\supset \Theta}$ | Unification context well-formedness | |
| $\boxed{\Gamma_1 \vdash \sigma : \Gamma_2}$ | Substitution well-formedness | |
| $\boxed{\Gamma \vdash e}$ | Unification solution entry well-formedness | |

Definition rules: 98 good 16 bad
Definition rule clauses: 195 good 16 bad