$\alpha, \beta, \alpha, \beta$  type variables n, m, i, j index variables

concatenate lists

$\overrightarrow{N}$	::=     	$.\\ \frac{N}{\overrightarrow{N}_i}^i$		list of negative types empty list a singel type concatenate lists
Δ, Γ	::=	$ \begin{array}{c} \overset{\cdot}{\alpha^{+}} \\ \overset{\cdot}{\alpha^{-}} \\ vars \\ \overline{\Gamma_{i}}^{i} \\ (\Gamma) \end{array} $	S	declarative type context empty context list of variables list of variables concatenate contexts
Θ	::=       	$ \overrightarrow{\widehat{\alpha^{+}}} $ $ \overrightarrow{\widehat{\alpha^{-}}} $ $ \overrightarrow{\Theta_{i}}^{i} $ $ (\Theta) $	S	unification type variable context empty context list of variables list of variables concatenate contexts
Ξ	::=	$ \overrightarrow{\alpha^{+}} \overrightarrow{\widehat{\alpha^{-}}} \overrightarrow{\Xi_{i}}^{i} (\Xi) \Xi_{1} \cup \Xi_{2} $	S	anti-unification type variable context empty context list of variables list of variables concatenate contexts
$\vec{\alpha}$ , $\vec{\beta}$		. $\overrightarrow{\alpha^+}$ $\overrightarrow{\alpha^-}$ $\overrightarrow{\alpha_1} \setminus vars$ $\Gamma$ $vars$ $\overrightarrow{\alpha_i}^i$ $(\overrightarrow{\alpha})$ $[\mu] \overrightarrow{\alpha}$ $ord \ vars \ in \ P$ $ord \ vars \ in \ P$ $ord \ vars \ in \ P$ $ord \ vars \ in \ N$	S M M M M	ordered positive or negative variables empty list list of variables list of variables setminus context  concatenate contexts parenthesis apply moving to list
vars	::=	$\varnothing$ <b>fv</b> $P$ <b>fv</b> $N$ <b>fv</b> $P$ <b>fv</b> $N$ $vars_1 \cap vars_2$		set of variables empty set free variables free variables free variables free variables set intersection 4

		$vars_1 \cup vars_2$ $vars_1 \backslash vars_2$ $\mathbf{mv} P$ $\mathbf{mv} N$ $\mathbf{uv} N$ $\mathbf{uv} P$ $\mathbf{fv} N$ $\mathbf{fv} P$ (vars) $\{\overrightarrow{\alpha}\}$ $[\mu] vars$	S M	set union set complement movable variables movable variables unification variables unification variables free variables free variables parenthesis ordered list of variables apply moving to varset
$\mu$	::=	$\begin{array}{c} \vdots \\ \widetilde{\alpha}_{1}^{+} \mapsto \widetilde{\alpha}_{2}^{+} \\ \widetilde{\alpha}_{1}^{-} \mapsto \widetilde{\alpha}_{2}^{-} \\ \mu_{1} \cup \mu_{2} \\ \overline{\mu_{i}}^{i} \\ \mu _{vars} \\ \mu^{-1} \end{array}$	M M M	empty moving Positive unit substitution Positive unit substitution Set-like union of movings concatenate movings restriction on a set inversion
n	::=   	$0 \\ n+1$		cohort index
$\widetilde{\alpha}^+$	::=	$\widetilde{\alpha}^{+n}$		positive movable variable
$\widetilde{lpha}^-$	::=	$\widetilde{\alpha}^{-n}$		negative movable variable
, ,	::=       	$ \widetilde{\alpha}^{+} \xrightarrow{\widetilde{\alpha}^{+n}} \widetilde{\alpha}^{i} \xrightarrow{\widetilde{\alpha}^{+}_{i}} i $		positive movable variable list empty list a variable from a non-movable variable concatenate lists
$\overrightarrow{\widetilde{\alpha}^-}$ , $\overrightarrow{\widetilde{\beta}^-}$	::=	$\vdots$ $\overbrace{\alpha}^{-}$ $\overrightarrow{\alpha}^{-n}$ $\overrightarrow{\alpha}_{i}$		negatiive movable variable list empty list a variable from a non-movable variable concatenate lists
_ ~		$\alpha^{+}$ $\alpha^{+}$ $\downarrow N$ $\exists \alpha^{-}.P$ $[\sigma]P$	M	multi-quantified positive types with movable variables

```
[\mu]P
                                 Μ
N, M
                                       multi-quantified negative types with movable variables
                     \alpha^{-}
                     \tilde{\alpha}^-
                     \uparrow P
                                 Μ
                      [\mu]N
                                 Μ
\hat{\alpha}^+
                                       positive unification variable
                     \hat{\alpha}^+
\hat{\alpha}^-
                                       negative unification variable
                                       positive unification variable list
                                          empty list
                                          a variable
                                          from a normal variable
                                          from a normal variable, context unspecified
                                          concatenate lists
                                       negative unification variable list
                                          empty list
                                          a variable
                                          from an antiunification context
                                          from a normal variable
                                          from a normal variable, context unspecified
                                          concatenate lists
P, Q
                                       a positive algorithmic type (potentially with metavariables)
                     \alpha^+
                     \tilde{\alpha}^+
                     \hat{\alpha}^+
                                  Μ
                                  Μ
                     [\mu]P
                                  Μ
                     \mathbf{nf}(P')
                                 Μ
N, M
                                       a negative algorithmic type (potentially with metavariables)
                     \alpha^{-}
                     \hat{\alpha}^-
```

```
\uparrow P
                                  P \rightarrow N
                                  \forall \overrightarrow{\alpha^+}.N
                                   [\sigma]N
                                                                                Μ
                                  [\mu]N
                                                                                Μ
                                   \mathbf{nf}(N')
                                                                                Μ
auSol
                         ::=
                                  (\Xi, Q, \widehat{	au}_1, \widehat{	au}_2)
terminals
                         ::=
                                   \exists
                                   \forall
                                   ↑
                                   \in
                                   ∉
                                   \leq
                                   \geqslant
                                   \simeq
                                   \cup
                                   \subseteq
                                   Ø
                                   \models
                                   \dashv
                                   \neq
                                  \equiv_n
                                   \Downarrow
                                   :≥
                                   :≃
formula
                         ::=
                                   judgement
                                  formula_1 .. formula_n
                                  \mu: vars_1 \leftrightarrow vars_2
                                   \mu is bijective
                                  \hat{\sigma} is functional
```

 $\hat{\sigma}_1 \in \hat{\sigma}_2$   $vars_1 \subseteq vars_2$   $vars_1 = vars_2$ 

```
vars is fresh
                             \alpha^- \notin vars
                             \alpha^+ \not\in \mathit{vars}
                             \alpha^- \in vars
                             \alpha^+ \in vars
                             \widehat{\alpha}^- \in \Theta
                             \hat{\alpha}^+ \in \Theta
                             if any other rule is not applicable
                             N \neq M
                             P \neq Q
A
                    ::=
                             \Gamma; \Theta \models \overline{N} \leq M \Rightarrow \hat{\sigma}
                                                                                                   Negative subtyping
                             \Gamma; \Theta \models P \geqslant Q \dashv \hat{\sigma}
                                                                                                   Positive supertyping
AU
                   ::=
                            \Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)
                            \Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 = (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)
E1
                            N \simeq_1^D M \\ P \simeq_1^D Q
                                                                                                   Negative multi-quantified type equivalence
                                                                                                   Positive multi-quantified type equivalence
                             P \simeq Q
D1
                            \Gamma \vdash N \overset{\mathsf{\scriptstyle \sim}}{}_{\mathsf{1}} M
                                                                                                   Negative equivalence on MQ types
                             \Gamma \vdash P \simeq_1^{\leqslant} Q
                                                                                                   Positive equivalence on MQ types
                            \Gamma \vdash N \leqslant_1 M
                                                                                                   Negative subtyping
                             \Gamma \vdash P \geqslant_1 Q
                                                                                                   Positive supertyping
D\theta
                    ::=
                             \Gamma \vdash N \simeq_0^{\scriptscriptstyle \leqslant} M
                                                                                                   Negative equivalence
                             \Gamma \vdash P \simeq_0^{\scriptscriptstyle \leqslant} Q
                                                                                                   Positive equivalence
                             \Gamma \vdash N \leqslant_0 M
                                                                                                   Negative subtyping
                             \Gamma \vdash P \geqslant_0 Q
                                                                                                   Positive supertyping
EQ
                    ::=
                             N = M
                                                                                                   Negative type equality (alpha-equivalence)
                             P = Q
                                                                                                   Positive type equuality (alphha-equivalence)
                             P = Q
LUBF
                   ::=
                             \operatorname{ord} vars \operatorname{in} P === \overrightarrow{\alpha}
                             ord vars in N === \vec{\alpha}
                             \operatorname{ord} \operatorname{varsin} P === \overrightarrow{\alpha}
                             \mathbf{ord}\ vars \mathbf{in}\ N === \overrightarrow{\alpha}
                             \mathbf{nf}(N') === N
                             \mathbf{nf}(P') === P
                             \mathbf{nf}(N') === N
```

```
\mathbf{nf}(P') === P
                                 e_1 \& e_2
                                 \hat{\sigma}_1 \& \hat{\sigma}_2
LUB
                                 \Gamma \vDash P_1 \vee P_2 = Q
                                                                                      Least Upper Bound (Least Common Supertype)
                                 \mathbf{upgrade}\,\Gamma \vdash P\,\mathbf{to}\,\Delta = Q
Nrm
                        ::=
                                 \mathbf{nf}(N) = M
                                 \mathbf{nf}(P) = Q
                                 \mathbf{nf}(N) = M
                                 \mathbf{nf}(P) = Q
Order
                                 \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}
                                 \mathbf{ord}\ vars \mathbf{in}\ P = \overrightarrow{\alpha}
                                 \mathbf{ord}\ vars \mathbf{in}\ N = \overrightarrow{\alpha}
                                 ord vars in P = \vec{\alpha}
SM
                                 e_1 \& e_2 = e_3\hat{\sigma}_1 \& \hat{\sigma}_2 = \hat{\sigma}_3
                                                                                      Unification Solution Entry Merge
                                                                                      Merge unification solutions
U
                        ::=
                                 \Theta \models N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}
                                                                                      Negative unification
                                 \Theta \vDash P \stackrel{u}{\simeq} Q \rightrightarrows \widehat{\sigma}
                                                                                      Positive unification
WF
                                 \Gamma \vdash N
                                                                                      Negative type well-formedness
                                 \Gamma \vdash P
                                                                                      Positive type well-formedness
                                 \Gamma \vdash N
                                                                                      Negative type well-formedness
                                 \Gamma \vdash P
                                                                                      Positive type well-formedness
                                 \Gamma;\Xi \vdash P
                                                                                      Positive anti-unification type well-formedness
                                 \Gamma; \Xi_2 \vdash \widehat{\tau} : \Xi_1
                                                                                      Antiunification substitution well-formedness
                                 \Theta \vdash \hat{\sigma}
                                                                                      Unification substitution well-formedness
                                 \Gamma \vdash \Theta
                                                                                      Unification context well-formedness
judgement
                                 A
                                 AU
                                 E1
                                 D1
                                 D\theta
                                 EQ
                                 LUB
                                 Nrm
                                 Order
```

SM

UWF $user\_syntax$  $\alpha$ nvars $\begin{array}{c} n \\ \widetilde{\alpha}^{+} \\ \widetilde{\alpha}^{-} \\ \widetilde{\alpha}^{+} \\ \widetilde{\alpha}^{-} \\ P \\ N \\ \widehat{\alpha}^{+} \\ \widehat{\alpha}^{-} \\ \widetilde{\alpha}^{-} \\ \widetilde{\alpha}^{-} \end{array}$ 

 $\boxed{\Gamma;\,\Theta \vDash N \leqslant M \dashv \widehat{\sigma}} \quad \text{Negative subtyping}$ 

auSol terminals formula

$$\frac{\Gamma; \; \Theta \vDash \alpha^{-} \leqslant \alpha^{-} \dashv \cdot}{\Gamma; \; \Theta \vDash \mathbf{nf} \; (P) \overset{u}{\simeq} \mathbf{nf} \; (Q) \dashv \widehat{\sigma}}$$
$$\frac{\Theta \vDash \mathbf{nf} \; (P) \overset{u}{\simeq} \mathbf{nf} \; (Q) \dashv \widehat{\sigma}}{\Gamma; \; \Theta \vDash \uparrow P \leqslant \uparrow Q \dashv \widehat{\sigma}} \quad \text{AShiftU}$$

$$\frac{\Gamma; \Theta \vDash P \geqslant Q \dashv \widehat{\sigma}_{1} \quad \Gamma; \Theta \vDash N \leqslant M \dashv \widehat{\sigma}_{2}}{\Gamma; \Theta \vDash P \to N \leqslant Q \to M \dashv \widehat{\sigma}_{1} \& \widehat{\sigma}_{2}} \quad \text{AArrow}$$

$$\frac{\Gamma, \overrightarrow{\beta^{+}}; \Theta, \widehat{\alpha^{+}} \{\Gamma, \overrightarrow{\beta^{+}}\} \vDash [\widehat{\alpha^{+}}/\alpha^{+}] N \leqslant M \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \forall \alpha^{+}. N \leqslant \forall \overrightarrow{\beta^{+}}. M \dashv \widehat{\sigma} \setminus \widehat{\alpha^{+}}} \quad \text{AForall}$$

 $\Gamma; \Theta \models P \geqslant Q \dashv \hat{\sigma}$  Positive supertyping

$$\frac{\Gamma; \Theta \vDash \alpha^{+} \geqslant \alpha^{+} \dashv \cdot}{\Gamma; \Theta \vDash \Lambda^{+} \geqslant \Lambda^{+} \dashv \widehat{\sigma}} \quad APVAR$$

$$\frac{\Theta \vDash \mathbf{nf}(N) \stackrel{u}{\simeq} \mathbf{nf}(M) \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \downarrow N \geqslant \downarrow M \dashv \widehat{\sigma}} \quad ASHIFTD$$

$$\frac{\Gamma, \overrightarrow{\beta^{-}}; \Theta, \widehat{\alpha^{-}} \{\Gamma, \overrightarrow{\beta^{-}}\} \vDash [\widehat{\alpha^{-}}/\widehat{\alpha^{-}}]P \geqslant Q \dashv \widehat{\sigma}}{\Gamma; \Theta \vDash \exists \alpha^{-}.P \geqslant \exists \overrightarrow{\beta^{-}}.Q \dashv \widehat{\sigma}} \quad AEXISTS$$

$$\frac{\mathbf{upgrade} \Gamma \vdash \mathbf{nf}(P) \mathbf{to} \Delta = Q}{\Gamma; \Theta \vDash \widehat{\alpha}^{+} \{\Delta\} \geqslant P \dashv (\Delta \vdash \widehat{\alpha}^{+} : \geqslant Q)} \quad APUVAR$$

 $\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 = (\Xi, Q, \hat{\tau}_1, \hat{\tau}_2)$ 

$$\frac{\Gamma \vDash \alpha^{+} \stackrel{a}{\simeq} \alpha^{+} \dashv (\cdot, \alpha^{+}, \cdot, \cdot)}{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUPShift}$$

$$\frac{\Gamma \vDash N_{1} \stackrel{a}{\simeq} N_{2} \dashv (\Xi, M, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \vDash \downarrow N_{1} \stackrel{a}{\simeq} \downarrow N_{2} \dashv (\Xi, \downarrow M, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUPShift}$$

$$\frac{\{\overrightarrow{\alpha^{-}}\} \cap \{\Gamma\} = \varnothing \quad \Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \vDash \exists \overrightarrow{\alpha^{-}} . P_{1} \stackrel{a}{\simeq} \exists \overrightarrow{\alpha^{-}} . P_{2} \dashv (\Xi, \exists \overrightarrow{\alpha^{-}} . Q, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUPEXISTS}$$

 $\Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 = (\Xi, M, \hat{\tau}_1, \hat{\tau}_2)$ 

$$\frac{\Gamma \vDash \alpha^{-} \stackrel{a}{\simeq} \alpha^{-} \dashv (\Xi, \alpha^{-}, \cdot, \cdot)}{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUNVAR}$$

$$\frac{\Gamma \vDash P_{1} \stackrel{a}{\simeq} P_{2} \dashv (\Xi, Q, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \vDash \uparrow P_{1} \stackrel{a}{\simeq} \uparrow P_{2} \dashv (\Xi, \uparrow Q, \hat{\tau}_{1}, \hat{\tau}_{2})} \quad \text{AUNSHIFT}$$

$$\frac{\Gamma \vDash P_1 \stackrel{a}{\simeq} P_2 \dashv (\Xi_1, Q, \hat{\tau}_1, \hat{\tau}_2) \quad \Gamma \vDash N_1 \stackrel{a}{\simeq} N_2 \dashv (\Xi_2, M, \hat{\tau}_1', \hat{\tau}_2')}{\Gamma \vDash P_1 \to N_1 \stackrel{a}{\simeq} P_2 \to N_2 \dashv (\Xi_1 \cup \Xi_2, Q \to M, \hat{\tau}_1 \cup \hat{\tau}_1', \hat{\tau}_2 \cup \hat{\tau}_2')} \quad \text{AUNARROW}$$

$$\frac{\text{if any other rule is not applicable} \quad \Gamma \vdash N \quad \Gamma \vdash M}{\Gamma \vDash N \stackrel{a}{\simeq} M \dashv (\widehat{\alpha}^-_{\{N,M\}}, \widehat{\alpha}^-_{\{N,M\}}, (\widehat{\alpha}^-_{\{N,M\}} :\approx N), (\widehat{\alpha}^-_{\{N,M\}} :\approx M))} \quad \text{AUNAU}$$

 $N \simeq_1^D M$  Negative multi-quantified type equivalence

$$\frac{\alpha^{-} \simeq_{1}^{D} \alpha^{-}}{\alpha^{-} \simeq_{1}^{D} Q} \quad \text{E1NVar}$$

$$\frac{P \simeq_{1}^{D} Q}{\uparrow P \simeq_{1}^{D} \uparrow Q} \quad \text{E1ShiftU}$$

$$\frac{P \simeq_{1}^{D} Q \quad N \simeq_{1}^{D} M}{P \to N \simeq_{1}^{D} Q \to M} \quad \text{E1Arrow}$$

$$\frac{\{\overrightarrow{\alpha^{+}}\} \cap \mathbf{fv} \, M = \varnothing \quad \mu : (\{\overrightarrow{\beta^{+}}\} \cap \mathbf{fv} \, M) \leftrightarrow (\{\overrightarrow{\alpha^{+}}\} \cap \mathbf{fv} \, N) \quad N \simeq_{1}^{D} [\mu] M}{\forall \overrightarrow{\alpha^{+}} . N \simeq_{1}^{D} \forall \overrightarrow{\beta^{+}} . M} \quad \text{E1Forall}$$

 $P \simeq_1^D Q$  Positive multi-quantified type equivalence

$$\frac{\alpha^{+} \simeq_{1}^{D} \alpha^{+}}{\alpha^{+}} \quad \text{E1PVAR}$$

$$\frac{N \simeq_{1}^{D} M}{\downarrow N \simeq_{1}^{D} \downarrow M} \quad \text{E1SHIFTD}$$

$$\frac{\{\overrightarrow{\alpha^{-}}\} \cap \text{fv } Q = \varnothing \quad \mu : (\{\overrightarrow{\beta^{-}}\} \cap \text{fv } Q) \leftrightarrow (\{\overrightarrow{\alpha^{-}}\} \cap \text{fv } P) \quad P \simeq_{1}^{D} [\mu]Q}{\exists \overrightarrow{\alpha^{-}} . P \simeq_{1}^{D} \exists \overrightarrow{\beta^{-}} . Q} \quad \text{E1EXISTS}$$

 $\begin{array}{|c|c|c|c|c|}\hline P \simeq Q \\ \hline \Gamma \vdash N & \cong_1^{\leq} M \\ \hline \end{array} \quad \text{Negative equivalence on MQ types}$ 

$$\frac{\Gamma \vdash N \leqslant_{1} M \quad \Gamma \vdash M \leqslant_{1} N}{\Gamma \vdash N \simeq_{1}^{\leqslant} M} \quad \text{D1NDEF}$$

 $\Gamma \vdash P \simeq_1^{\varsigma} Q$  Positive equivalence on MQ types

$$\frac{\Gamma \vdash P \geqslant_1 Q \quad \Gamma \vdash Q \geqslant_1 P}{\Gamma \vdash P \simeq_1^{\circ} Q} \quad \text{D1PDEF}$$

 $\Gamma \vdash N \leq_1 M$  Negative subtyping

 $\Gamma \vdash P \geqslant_1 Q$  Positive supertyping

 $\Gamma \vdash N \simeq_0^{\leq} M$  Negative equivalence

$$\frac{\Gamma \vdash N \leqslant_0 M \quad \Gamma \vdash M \leqslant_0 N}{\Gamma \vdash N \simeq_0^{\leqslant} M} \quad \text{D0NDEF}$$

 $\Gamma \vdash P \simeq_0^{\leqslant} Q$  Positive equivalence

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash Q \geqslant_0 P}{\Gamma \vdash P \simeq_0^{\varsigma} Q} \quad \text{D0PDEF}$$

 $\Gamma \vdash N \leq_0 M$  Negative subtyping

$$\frac{\Gamma \vdash \alpha^- \leqslant_0 \alpha^-}{\Gamma \vdash P \stackrel{<}{\sim}_0^{\varsigma} Q} \quad D0\text{NVar}$$

$$\frac{\Gamma \vdash P \stackrel{\sim}{\sim}_0^{\varsigma} Q}{\Gamma \vdash \uparrow P \leqslant_0 \uparrow Q} \quad D0\text{ShiftU}$$

$$\frac{\Gamma \vdash P \quad \Gamma \vdash [P/\alpha^+]N \leqslant_0 M \quad M \neq \forall \beta^+.M'}{\Gamma \vdash \forall \alpha^+.N \leqslant_0 M} \quad D0\text{ForallL}$$

$$\frac{\Gamma, \alpha^+ \vdash N \leqslant_0 M}{\Gamma \vdash N \leqslant_0 \forall \alpha^+.M} \quad D0\text{ForallR}$$

$$\frac{\Gamma \vdash P \geqslant_0 Q \quad \Gamma \vdash N \leqslant_0 M}{\Gamma \vdash P \rightarrow N \leqslant_0 Q \rightarrow M} \quad D0\text{Arrow}$$

 $\overline{|\Gamma \vdash P \geqslant_0 Q|}$  Positive supertyping

$$\frac{\Gamma \vdash \alpha^{+} \geqslant_{0} \alpha^{+}}{\Gamma \vdash \lambda^{+} \geqslant_{0} M} \quad D0PVAR$$

$$\frac{\Gamma \vdash N \simeq_{0}^{\leq} M}{\Gamma \vdash \downarrow N \geqslant_{0} \downarrow M} \quad D0SHIFTD$$

$$\frac{\Gamma \vdash N \quad \Gamma \vdash [N/\alpha^{-}]P \geqslant_{0} Q \quad Q \neq \exists \alpha^{-}.Q'}{\Gamma \vdash \exists \alpha^{-}.P \geqslant_{0} Q} \quad D0EXISTSL$$

$$\frac{\Gamma, \alpha^{-} \vdash P \geqslant_{0} Q}{\Gamma \vdash P \geqslant_{0} \exists \alpha^{-}.Q} \quad D0EXISTSR$$

N = M Negative type equality (alpha-equivalence) P = Q Positive type equivalence) P = Q ord vars in P

ord vars in N

ord vars in P

 $\overline{\mathbf{ord} \ vars \mathbf{in} \ N}$ 

 $\mathbf{nf}(N')$ 

 $|\mathbf{nf}(P')|$ 

 $\mathbf{nf}(N')$ 

 $\mathbf{nf}\left(P'
ight)$ 

 $e_1 \& e_2$ 

 $\widehat{\sigma}_1 \& \widehat{\sigma}_2$ 

 $\overline{\Gamma \models P_1 \lor P_2 = Q}$  Least Upper Bound (Least Common Supertype)

$$\frac{\Gamma \models \alpha^{+} \vee \alpha^{+} = \alpha^{+}}{\Gamma \models \lambda^{+} \vee \alpha^{+} = \alpha^{+}} \quad \text{LUBVAR}$$

$$\frac{\Gamma, \cdot \models \downarrow N \stackrel{a}{\simeq} \downarrow M \Rightarrow (\Xi, P, \hat{\tau}_{1}, \hat{\tau}_{2})}{\Gamma \models \downarrow N \vee \downarrow M = \exists \alpha^{-}. [\alpha^{-}/\Xi] P} \quad \text{LUBSHIFT}$$

$$\frac{\Gamma, \alpha^{-}, \beta^{-}}{\Gamma, \alpha^{-}, \beta^{-}} \models P_{1} \vee P_{2} = Q$$

$$\frac{\Gamma, \alpha^{-}, \beta^{-}}{\Gamma, \beta^{-}} \models P_{1} \vee P_{2} = Q$$

$$\Gamma \models \exists \alpha^{-}. P_{1} \vee \exists \beta^{-}. P_{2} = Q$$

$$LUBEXISTS$$

$$\overline{\mathbf{nf}(\alpha^{-}) = \alpha^{-}} \quad \text{NrmNVar}$$

$$< < \text{multiple parses} >> \\ \overline{\mathbf{nf}(\uparrow P) = \uparrow Q} \quad \text{NrmShiftU}$$

$$< < \text{multiple parses} >> \\ \overline{\mathbf{nf}(P \to N) = Q \to M} \quad \text{NrmArrow}$$

$$< < \text{multiple parses} >> \\ \overline{\mathbf{nf}(\forall \alpha^{+}.N) = \forall \alpha^{+'}.N'} \quad \text{NrmForall}$$

 $\mathbf{nf}(P) = Q$ 

$$\overline{\mathbf{nf}(\alpha^{+}) = \alpha^{+}} \quad \text{NRMPVAR}$$

$$< < \mathbf{multiple parses} >>$$

$$\mathbf{nf}(\downarrow N) = \downarrow M \qquad \qquad \text{NRMSHIFTD}$$

$$< < \mathbf{multiple parses} >>$$

$$\overline{\mathbf{nf}(\exists \alpha^{-}.P) = \exists \alpha^{-'}.P'} \qquad \text{NRMEXISTS}$$

 $\mathbf{nf}(N) = M$ 

$$\underline{\mathbf{nf}(\widehat{\alpha}^{-}) = \widehat{\alpha}^{-}}$$
 NRMNUVAR

 $\mathbf{nf}(P) = Q$ 

$$\overline{\mathbf{nf}(\widehat{\alpha}^+) = \widehat{\alpha}^+}$$
 NRMPUVAR

## $\mathbf{ord} \ vars \mathbf{in} \ N = \vec{\alpha}$

$$\frac{\alpha^{-} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \alpha^{-}} \quad \text{ONVARIN}$$

$$\frac{\alpha^{-} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{-} = \cdot} \quad \text{ONVARNIN}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \uparrow P = \overrightarrow{\alpha}} \quad \text{OSHIFTU}$$

$$\frac{\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}_{1} \quad \operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}_{2}}{\operatorname{ord} vars \operatorname{in} P \to N = \overrightarrow{\alpha}_{1}, (\overrightarrow{\alpha}_{2} \setminus \{\overrightarrow{\alpha}_{1}\})} \quad \text{O.}$$

$$\frac{\text{ord } vars \text{ in } P = \alpha_1 \quad \text{ord } vars \text{ in } N = \alpha_2}{\text{ord } vars \text{ in } P \to N = \overrightarrow{\alpha}_1, (\overrightarrow{\alpha}_2 \setminus \{\overrightarrow{\alpha}_1\})} \quad \text{OARROW}$$

$$\frac{vars \cap \{\overrightarrow{\alpha}^+\} = \varnothing \quad \text{ord } vars \text{ in } N = \overrightarrow{\alpha}}{\text{ord } vars \text{ in } \forall \overrightarrow{\alpha}^+. N = \overrightarrow{\alpha}} \quad \text{OFORALL}$$

## $\operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}$

$$\frac{\alpha^{+} \in vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \alpha^{+}} \quad \operatorname{OPVarIn}$$

$$\frac{\alpha^{+} \notin vars}{\operatorname{ord} vars \operatorname{in} \alpha^{+} = \cdot} \quad \operatorname{OPVarNIn}$$

$$\frac{\operatorname{ord} vars \operatorname{in} N = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \sqrt{N} = \overrightarrow{\alpha}} \quad \operatorname{OSHIFTD}$$

$$\frac{vars \cap \{\overrightarrow{\alpha^{-}}\} = \varnothing \quad \operatorname{ord} vars \operatorname{in} P = \overrightarrow{\alpha}}{\operatorname{ord} vars \operatorname{in} \exists \overrightarrow{\alpha^{-}}.P = \overrightarrow{\alpha}} \quad \operatorname{OEXISTS}$$

$$\operatorname{ord} vars \operatorname{in} \exists \overrightarrow{\alpha^{-}}.P = \overrightarrow{\alpha}$$

## $\operatorname{\mathbf{ord}} vars \operatorname{\mathbf{in}} N = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}} \operatorname{\mathit{vars}} \operatorname{\mathbf{in}} \widehat{\alpha}^- = \cdot}$$
 ONUVAR

## $\mathbf{ord} \ vars \mathbf{in} \ P = \overrightarrow{\alpha}$

$$\frac{}{\operatorname{\mathbf{ord}} \operatorname{\mathit{vars}} \operatorname{\mathbf{in}} \widehat{\alpha}^{+} = \cdot} \quad \operatorname{OPUVar}$$

 $\overline{|e_1 \& e_2 = e_3|}$  Unification Solution Entry Merge

$$\frac{\Gamma \vDash P_1 \lor P_2 = Q}{(\Gamma \vdash \widehat{\alpha}^+ : \geqslant P_1) \& (\Gamma \vdash \widehat{\alpha}^+ : \geqslant P_2) = (\Gamma \vdash \widehat{\alpha}^+ : \geqslant Q)} \quad \text{SMESUPSUP}$$

$$\frac{\Gamma; \cdot \vDash P \geqslant Q \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \geqslant Q) = (\Gamma \vdash \widehat{\alpha}^+ : \approx P)} \quad \text{SMEEQSUP}$$

$$\frac{\Gamma; \cdot \vDash Q \geqslant P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \geqslant P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx Q) = (\Gamma \vdash \widehat{\alpha}^+ : \approx Q)} \quad \text{SMESUPEQ}$$

$$\frac{\Gamma; \cdot \vDash Q \geqslant P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx Q) = (\Gamma \vdash \widehat{\alpha}^+ : \approx Q)} \quad \text{SMEPEQEQ}$$

$$\frac{\Gamma; \cdot \vDash Q \Rightarrow P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx P)} \quad \text{SMEPEQEQ}$$

$$\frac{\Gamma; \cdot \vDash Q \Rightarrow P \dashv \widehat{\sigma}'}{(\Gamma \vdash \widehat{\alpha}^+ : \approx P) \& (\Gamma \vdash \widehat{\alpha}^+ : \approx P)} \quad \text{SMEPEQEQ}$$

 $\widehat{\sigma}_1 \& \widehat{\sigma}_2 = \widehat{\sigma}_3$  Merge unification solutions

 $\Theta \models N \stackrel{u}{\simeq} M \dashv \widehat{\sigma}$  Negative unification

$$\frac{}{\Theta \vDash \alpha^{-} \stackrel{u}{\simeq} \alpha^{-} \dashv \cdot} \quad \text{UNVAR}$$

$$\frac{\Theta \vDash P \overset{u}{\simeq} Q \dashv \widehat{\sigma}}{\Theta \vDash \uparrow P \overset{u}{\simeq} \uparrow Q \dashv \widehat{\sigma}} \quad \text{USHIFTU}$$

$$\frac{\Theta \vDash P \overset{u}{\simeq} Q \dashv \widehat{\sigma}_1 \quad \Theta \vDash N \overset{u}{\simeq} M \dashv \widehat{\sigma}_2}{\Theta \vDash P \to N \overset{u}{\simeq} Q \to M \dashv \widehat{\sigma}_1 \& \widehat{\sigma}_2} \quad \text{UARROW}$$

$$\frac{\Theta \vDash N \overset{u}{\simeq} M \dashv \widehat{\sigma}}{\Theta \vDash \forall \alpha^+. N \overset{u}{\simeq} \forall \alpha^+. M \dashv \widehat{\sigma}} \quad \text{UFORALL}$$

$$\frac{\widehat{\alpha}^- \{\Delta\} \in \Theta \quad \Delta \vdash N}{\Theta \vDash \widehat{\alpha}^- \overset{u}{\simeq} N \dashv (\Delta \vdash \widehat{\alpha}^- : \approx N)} \quad \text{UNUVAR}$$

 $\Theta \models P \stackrel{u}{\simeq} Q \dashv \hat{\sigma}$  Positive unification

$$\frac{\Theta \vDash \alpha^{+} \stackrel{u}{\simeq} \alpha^{+} \dashv \cdot}{\Theta \vDash N \stackrel{u}{\simeq} M \dashv \hat{\sigma}} \quad \text{UPVAR}$$

$$\frac{\Theta \vDash N \stackrel{u}{\simeq} M \dashv \hat{\sigma}}{\Theta \vDash \downarrow N \stackrel{u}{\simeq} \downarrow M \dashv \hat{\sigma}} \quad \text{USHIFTD}$$

$$\frac{\Theta \vDash P \stackrel{u}{\simeq} Q \dashv \hat{\sigma}}{\Theta \vDash \exists \alpha^{-}.P \stackrel{u}{\simeq} \exists \alpha^{-}.Q \dashv \hat{\sigma}} \quad \text{UEXISTS}$$

$$\frac{\hat{\alpha}^{+} \{\Delta\} \in \Theta \quad \Delta \vdash P}{\Theta \vDash \hat{\alpha}^{+} \stackrel{u}{\simeq} P \dashv (\Delta \vdash \hat{\alpha}^{+} : \approx P)} \quad \text{UPUVAR}$$

 $\Gamma \vdash N$ Negative type well-formedness

 $\Gamma \vdash P$  Positive type well-formedness

Negative type well-formedness

 $\Gamma \vdash P$ Positive type well-formedness

 $\Gamma;\Xi\vdash P$  Positive anti-unification type well-formedness  $\Gamma;\Xi_2\vdash \hat{\tau}:\Xi_1$  Antiunification substitution well-formedn Antiunification substitution well-formedness

Unification substitution well-formedness  $\Theta \vdash \hat{\sigma}$ 

 $\Gamma \vdash \Theta$ Unification context well-formedness

Definition rules: 72 good 7 bad Definition rule clauses: 130 good 7 bad