

COMP3121 Assignment 3 - Q5

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Answer

We first initialise a 2D array with size $|V| \times |V|$, $weight[|V|][|V|]$, to store the weights between two vertices in $G(V, E)$. The diagonal entries in $weight[|V|][|V|]$ will be 0, as visiting a node from itself will have no weight. Nodes not directly reachable will be infinity.

Second initialise a 3D array of size $|V| \times |V| \times K$, $dp[s][d][k]$, to represent the maximum total weight path from node s to d using exactly k edges in total. Initialise the whole array to infinity except for $d = s$, which sets to 0.

We now start finding the the maximum total weight path. For each s in V (V is the set of vertex in graph G), we loop K times, in each loop we examine every d in V , and update dp by do the following:

$$dp[s][d][k+1] = \max(arr[s][i][k] + weight[i][d])$$

for all i which has an edge from i to d , i.e. $weight[i][d]$ is not infinity or 0.

Here $arr[s][i][k]$ is a copy of $dp[s][d][k]$ before updating dp .

After finishing above, we loop to find the maximum number in $dp[s][d][K]$, which is the the maximum total weight of exactly length K we can get from G .

To get path, we use a recursive function and an array $path[]$ to store the path,

$$opt(s, d, k) = \max\{opt(s, i, k) + weight[i][d] : weight[i][d] \text{ is not infinity or } 0\}$$

We add all i chosen from $\max\{\}$ to array $path$, where i represents each node visited, and reverse path when finishing the whole recursive process, then we get the maximum total weight of exactly length K path from the array $path$.

Time complexity is $O(K * |V|^2)$, as we have $|V|$ vertices in G , we have chosen every vertex as source node s , and for each source node s we loop K times to check every destination node d . Hence the time complexity is $O(|V| * K * |V|) = O(K * |V|^2)$.