

CS528

**Energy Aware VM Consolidation
in Cloud using Prediction Model**

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Reference

F. Farahnakian et. al, “**Energy Aware VM Consolidation in Cloud using Prediction Model**”, IEEE Trans. On Cloud Computing, June 2019

Energy Aware VM Consolidation in Cloud using Prediction Model

- Virtual Machine (VM) consolidation
 - Promising approach to save energy and
 - improve resource utilization in data centers.
- Many heuristic algorithms for VM consolidation as
 - Vector Bin-Packing Problem.

VM consolidation: Vector Bin Packing

- Given A set of N VMs with resource requirements $(r_c, r_m, r_{db}, r_{nb}, \text{etc})$
 - CPU, memory, disk BW, net BW, etc
- Given a set of homogenous host/machines with capacity $(C_c, C_m, C_{db}, C_{nb}, \text{etc})$
 - CPU Capacity, memory, disk BW, net BW available
- Pack this VMs to minimum number of host
- Simple examples: 2D case or 2 resources case
 - 10 VM with CPU and Mem requirement $vm_i(r_c, r_m)$ need to map to hosts with 4CPU+4GB of RAM
 - Minimize number of CPU

Energy Aware VM Consolidation in Cloud using Prediction Model

- Focused mostly on number of active PM minimization **//Static Problem**
 - According current resource requirements and neglected the future resource demands.
- So, they generate unnecessary VM migrations
 - increase the rate of SLA violations in data centers.
- Needs VM consolidation approach
 - That takes into account both the current and future utilization of resource
- Simple regression-based model may be enough
 - to approximate the future CPU and memory utilization of VMs and PMs

Utilization Prediction-aware VM Consolidation (UP-VMC)

- VM consolidation as 2D vector bin packing
- UP-VMC consider
 - CPU utilization and Memory utilization
 - Also, considers current & future resource utilization
- Approximate the future utilization
 - two regression-based prediction models
 - linear prediction model and k-nearest neighbor.
- Prediction model in order to predict
 - Resource utilization of VMs;
 - resource utilization of PMs;
 - resource utilization of both VMs and PMs.

UtilPred-VM Consolidation

- VM selection methods: affect on performance
 - in terms of the **energy consumption**,
 - the number of **SLA violations** and
 - the **number of migrations**
- Performance of VM consolidation is increased
 - Selects a VM that requires the minimum time for migration to another PM

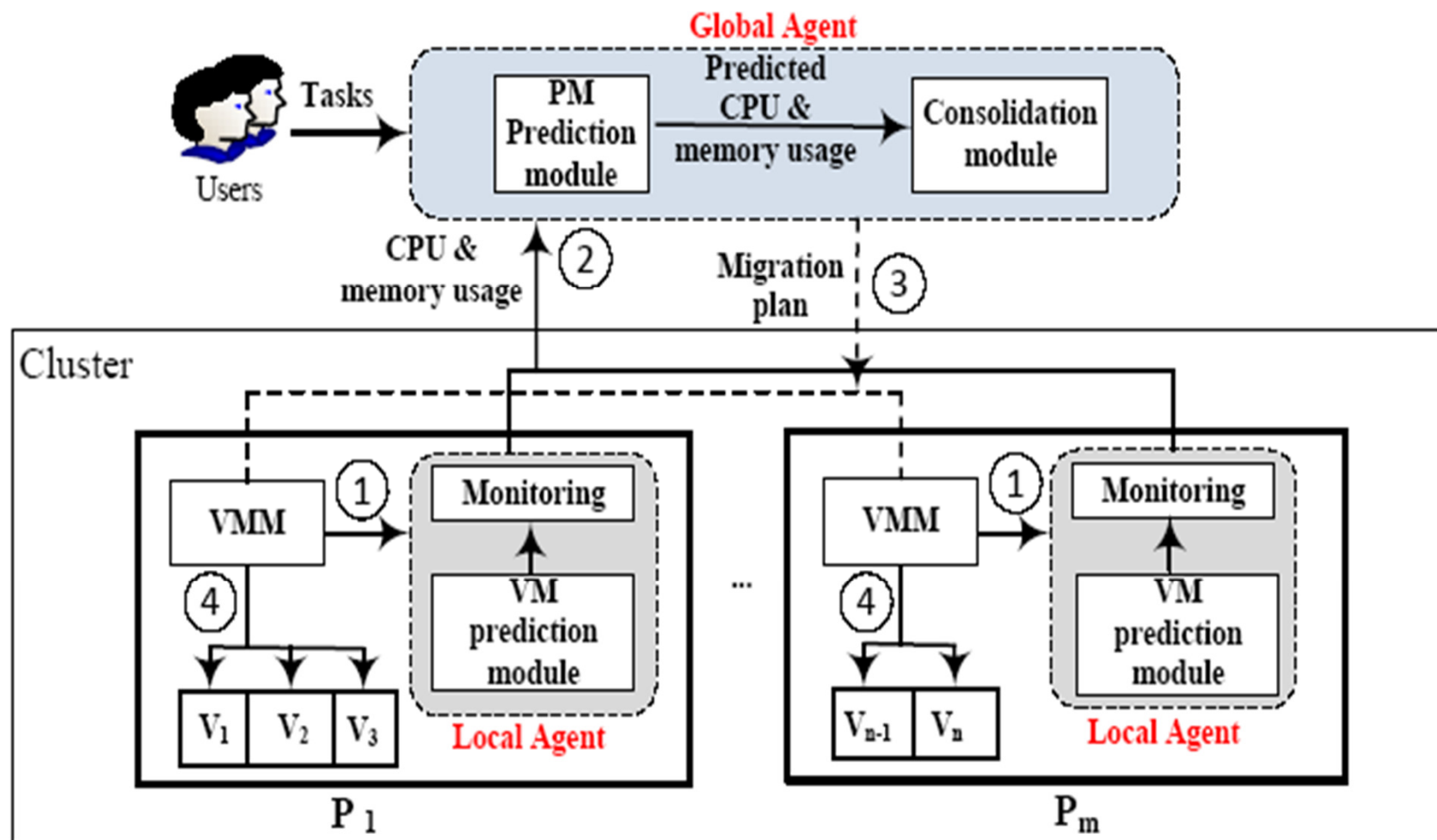
System Architecture

- A data center consists of m heterogeneous PMs $\mathbf{P} = \{ p_1, p_2, \dots, p_m \}$.
- Each PM with D type of resources
 - CPU, memory, network I/O and storage capacity.
- Multiple VMs can be allocated to each PM
 - through Virtual Machine Monitor (VMM).
- At any given time, users submit their requests
 - for provisioning of n VMs, $\mathbf{V} = \{ v_1, v_2, \dots, v_n \}$
 - which are allocated to the PMs.

System Architecture

- As requested utilization of VMs and PMs vary over time
 - An initial efficient allocation approach needs to be augmented
 - with a VM consolidation algorithm that can be applied periodically.
- In order to adapt and optimize
 - the VM placement periodically according to the workload.
 - We need to monitor and predict the workload of all the VMs and PMs
 - **Local agent at PM, Global agent at scheduler level**

System Architecture



System Architecture

- **System Arch. consists of two kind of agents:**
 - Fully distributed Local Agents (LAs) in PMs
 - Global Agent (GA) resides in a master node
- **Each LA monitor and predict**
 - Monitors the current resource utilization of all VMs in a PM periodically.
 - Approximates the future utilization of all VMs in a PM using a regression-based prediction model
- **GA collects info. LAs**
 - to maintain the overall view of current and future resource utilization of VMs.
 - GA builds a **global best migration plan** to optimize VM placement

Solution Approaches

- Each PM with **d** type of resources
 - **CPU, memory**, network I/O and storage capacity.
- Each PM **p** has a **d**-dimensional total capacity vector
$$\mathbf{C_p} = \langle \mathbf{C_p^1}, \mathbf{C_p^2}, \dots, \mathbf{C_p^d} \rangle$$
 - where $\mathbf{C_p^d}$ is total $\mathbf{d_{th}}$ resource capacity of PM **p**.
- Used capacity vector of the PM **p** as
$$\mathbf{U_p} = \langle \mathbf{U_p^1}, \mathbf{U_p^2}, \dots, \mathbf{U_p^d} \rangle$$
 - where $\mathbf{U_p^d}$ denotes used capacity of the resource type **d** :
simplicity **d=2, CPU, memory**
- For instance, used CPU capacity of a PM is estimated
 - as the sum of the CPU utilization of the three VMs if three VMs are hosted by the same PM

Solution Approaches

- Each VM v has a d -dimensional total capacity vector

$$Cv = \langle C^1_v, C^2_v, \dots, C^d_v \rangle$$

– where C^d_v is total d_{th} resource capacity of VM v .

- Used capacity vector of the VM v as

$$U_v = \langle U^1_v, U^2_v, \dots, U^d_v \rangle$$

– where U^d_v denotes used capacity of the resource type d : simplicity $d=2$, CPU, memory

- As the resource utilization of VMs
 - vary over time due to dynamic workloads,
 - the VM placement need to be optimized periodically

Solution Approaches : Step 1

- Aims to migrate some VMs from
 - **over-loaded PMs** and
 - **predicated over-loaded PMs**.
- If at least one resource (i.e., CPU or memory)
 - Exceeds total capacity, **PM is over-loaded**
 - Belongs to **overloaded PMs set** (P_{over}).
- If at least one resource predicted utilization value
 - is larger than capacity, **PM is predicted overloaded**
 - Belongs to **Predicted over-loaded PMs set** (P^{\wedge}_{over})

Prediction Model

- Predicted Util Vector of PM: $PU_{p_{de}} = \alpha + \beta U_{p_{de}}$
 - Current used capacity vector : $U_{p_{de}}$
 - α and β derived using linear regression
- Regression coefficients can be estimated

$$\beta = \frac{\sum_{i=1}^n (X_i - X^b)(Y_i - Y^b)}{\sum_{i=1}^n (X_i - X^b)^2}$$

$$\alpha = Y^b - \beta X^b$$

- where X^b is the mean value of X_1, X_2, \dots, X_n , and
- Y^b is the mean value of Y_1, Y_2, \dots, Y_n

Prediction Model

- Predicted Util Vector of VM: $PU_v = \alpha + \beta U_v$
 - Current used Util vector : U_v
 - α and β derived using linear regression
- PM load $Load_p = \sum_{d \in \{1,2,...,|D|\}} R_p^d$
 - where $R_p^d = \frac{U_p^d}{C_p^d}$ where U, C are Utilized & Capacity
- VM load $Load_v = \sum_{d \in \{1,2,...,|D|\}} R_v^d$ where
$$R_v^d = \frac{U_v^d}{C_v^d}$$

Constraints on Consolidation

- Constraint 1 : Used Capacity of destination and added with used capacity of VM should be less than threshold

$$U_{p_{de}} + U_v \leq T \cdot C_{p_{de}}$$

- Constraint 2 : Predicted Capacity of destination and Predicted capacity of VM should be less than threshold

$$PU_{p_{de}} + PU_v \leq T \cdot C_{p_{de}}$$

- Both should hold

Scale up and Scale Down

- Scale UP: demand is high
 - If required switch on more PM to serve better
- Scale down : Demand is less
 - If require power off some PM to save power

Algorithm : Consolidation and Scale Up

Set $M_1 = \Phi$;

for p_{so} in $P_{over} \cup P_{over}^{\wedge}$ do **//Over loaded**

Sort VMs V_m on PM p_{so} in ascending order
based on U_{mem}

for v in V_m do

for p_{de} in $P - (P_{over} \cup P_{over}^{\wedge})$ do **//Non-overloaded**

if $U_{p_{de}} + U_v \leq T.C_{p_{de}}$ & $PU_{p_{de}} + PU_v \leq T.C_{p_{de}}$

$M_1 = M_1 \cup [(p_{so}, v, p_{de})]$

Update Up_{so} and Up_{de} ; break

if not able to find any destination PM

Switch on the dormant PM p ; **//Scale UP**

Algorithm : Consolidation-Scale Down

Sort P_{active} in descending of Load_p ;

for $\text{PM}_{\text{so}} = |P_{\text{active}}|$ to 1 do **//Start from Light loaded one**

V_m = sort VMs on PM_{so} in descending order of Load_v ;

Set $M_2 = \Phi$;

for v in V_m do

 success=false;

 for p_{de} in $P_{\text{active}} - \text{PM}_{\text{so}}$ do

 if $U_{p_{\text{de}}} + U_v \leq T.C_{p_{\text{de}}}$ & $PU_{p_{\text{de}}} + PU_v \leq T.C_{p_{\text{de}}}$

$M_2 = M_2 \cup [(p_{\text{so}}, v, p_{\text{de}})]$; success=True;

 Update $U_{p_{\text{so}}}$ and $U_{p_{\text{de}}}$; break

if success = false; Recover $U_{p_{\text{so}}}$ and $U_{p_{\text{de}}}$; $M_2 = \Phi$;

else **Switch PMso to the sleep mode**

Performance Metrics

- **SLA Violation (SLAV)** **$SLAV = SLAVO * SLAVM$**
 - due to Overload (SLAVO),
 - due to Migration (SLAVM)
- **$SLAVO = \frac{1}{M} \sum_{i=1}^m \frac{T_{s_i}}{T_{a_i}}$**
 - M number of PM, T_{s_i} total time PM I experienced CPU/Mem utilization above 100%
- **$SLAVM = \frac{1}{N} \sum_{j=1}^n \frac{C_{d_j}}{C_{r_j}}$**
 - Experience performance degradation of j the VM by migration C_{r_j} total capacity requested by VM

Reference:

Wu et.al, *End to End Delay Minimization for Scientific Workflow in cloud under Budget Constraints*, IEEE Trans. On Cloud Computing. 2015.

Introduction

- With emergence of cloud computing and rapid deployment of cloud infrastructures
 - Number of **scientific workflows** have been shifted to cloud environments.
- Challenges:
 - Reducing financial cost in addition
 - to meeting the traditional goal : performance
- Quick evaluation of scientific workflow
 - to minimize the workflow end-to-end delay under a user-specified financial constraint

Scientific Workflow : SWF

- Large-scale scientific computing tasks
 - for data generation, processing, and analysis are
 - often assembled and constructed as **Workflows**
 - comprised of many interdependent **modules**
- Workflow (WF) module communicates
 - with others through the sharing of data sets,
 - which are either stored in shared file system or
 - transferred from node to node by WF management system
- Scientific Workflows are
 - typically executed in a distributed manner
 - in heterogeneous network environments

WF in Clouds System

- It is essential construct analytical models
 - to quantify the network performance of scientific workflows
 - in IaaS cloud environments,
 - and formulate a task scheduling problem
- Scheduling Problems: WF on Cloud
 - to minimize the workflow end-to-end delay
 - under a user-specified financial cost constraint,
- Referred to as Minimum End-to-end Delay under Cost Constraint (MED-CC)

Workflow Execution in Cloud

- Workflow is represented as
 - a directed acyclic graph (DAG),
- Submitted to the workflow engine
 - for executing, scheduling, tracking and reporting
- Workflow (WF) have independent tasks (Work/W), and
 - Execution model with inter-module dependencies
 - Identify and quantify the key financial and time cost