

CSE 450 Operating Systems  
Homework 4

Name: partial solutions

Question 2. Consider the following snapshot of a system (P=Process, R=Resource) :

| Available      |                |                |                |
|----------------|----------------|----------------|----------------|
| R <sub>a</sub> | R <sub>b</sub> | R <sub>c</sub> | R <sub>d</sub> |
| 1              | 5              | 2              | 0              |

| Maximum Demand |                |                |                |                |
|----------------|----------------|----------------|----------------|----------------|
|                | R <sub>a</sub> | R <sub>b</sub> | R <sub>c</sub> | R <sub>d</sub> |
| P <sub>0</sub> | 0              | 3              | 1              | 2              |
| P <sub>1</sub> | 1              | 7              | 5              | 0              |
| P <sub>2</sub> | 2              | 3              | 5              | 6              |
| P <sub>3</sub> | 0              | 6              | 5              | 2              |
| P <sub>4</sub> | 0              | 6              | 5              | 6              |

| Current Allocation |                |                |                |                |
|--------------------|----------------|----------------|----------------|----------------|
|                    | R <sub>a</sub> | R <sub>b</sub> | R <sub>c</sub> | R <sub>d</sub> |
| P <sub>0</sub>     | 0              | 0              | 1              | 2              |
| P <sub>1</sub>     | 1              | 0              | 0              | 0              |
| P <sub>2</sub>     | 1              | 3              | 5              | 4              |
| P <sub>3</sub>     | 0              | 6              | 3              | 2              |
| P <sub>4</sub>     | 0              | 0              | 1              | 4              |

Answer the following questions using banker's algorithm:

- a) [5 points] Calculate the *Needs* matrix: (= maximum demand – current allocation)

| Needs          |                |                |                |                |
|----------------|----------------|----------------|----------------|----------------|
|                | R <sub>a</sub> | R <sub>b</sub> | R <sub>c</sub> | R <sub>d</sub> |
| P <sub>0</sub> | 0              | 3              | 0              | 0              |
| P <sub>1</sub> | 0              | 7              | 5              | 0              |
| P <sub>2</sub> | 1              | 0              | 0              | 2              |
| P <sub>3</sub> | 0              | 0              | 2              | 0              |
| P <sub>4</sub> | 0              | 6              | 4              | 2              |

- b) [10 points] Is the system in a safe state? If so, show how you derive a safe order with Safety Algorithm in which the processes can run. Show the different values of the *work* vector after each iteration. What is the sequence of processes that the algorithm implicitly created?

Work = available = [1, 5, 2, 0]

Flag = [0, 0, 0, 0, 0, 0]

For i = 0; i < 5; i++

Find i that satisfies:

- 1) Need[i] <= work
- 2) Flag[i] == 0

Find i = 3, that [0, 0, 2, 0] <= work [1,5,2,0], so work = work + allocation[3] = [1, 11, 3, 2]

Next, find  $i = 4$  satisfies the conditions, so  $work = work + allocation[4] = [1, 11, 4, 6]$ ,  
 Next,  $i$  can go with either 0 or 2, assume we go with  $i=0$ , then  $work = work + allocation[0]$   
 $= [1, 11, 5, 8]$ ;  
 Next,  $i = 1$ ,  $work = work + allocation[1] = [2, 11, 5, 8]$ ;  
 Next  $i = 2$ ,  $work = work + allocation[2] = [3, 14, 10, 12]$ .

So the sequence we found from safety algorithm is  $P_3, P_4, P_0, P_1, P_2$ ; [note: there could be other sequences available as well].

- c) If a request from process  $P_0$  arrives for  $(0, 3, 0, 0)$ , can the request be granted immediately? Justify your answer, using only the knowledge of the sequence you found at sub-question (b).

~~A request  $[0, 3, 0, 0]$  from  $P_0$  won't be granted immediately, since  $[0, 3, 0, 0] \not\leq available$   
 $[1, 5, 2, 0]$~~

1. Request  $[0, 3, 0, 0] < Available = [1, 5, 2, 0]$

2. Request  $[0, 3, 0, 0] \leq need[0] = [0, 3, 0, 0]$

We assume the request will be granted, then

Work = Available = available – request =  $[1, 2, 2, 0]$ , and allocation  $[0] = allocation[0] +$   
 request =  $[0, 3, 1, 2]$ , needs $[0] = [0, 0, 0, 0]$

Run the safety check as follows:

For  $i = 0; i < 5; i++$

Find  $i$  that satisfies:

- 1) Need $[i] \leq work$
- 2) Flag $[i] == 0$

First, we found  $i = 3$ , needs $[3] = [0, 0, 2, 0] < work = [1, 2, 2, 0]$ , then  $work = work +$   
 allocation  $[3] = [1, 8, 5, 2]$ , flag $[3] = true$ ;

We next try  $i = 4$ , needs $[4] = [0, 6, 4, 2] < work = [1, 8, 5, 2]$ , so now  $work = work +$   
 allocation $[4] = [1, 8, 6, 6]$ , flag $[4] = true$  ... similarly, we could find  $i = 0, 1, 2$  satisfying the  
 conditions, so we could find a sequence of processes,  $P_3, P_4, P_0, P_1, P_2$ , which supports the  
 immediate allocation of  $[0, 3, 0, 0]$  to process  $P_0$ .