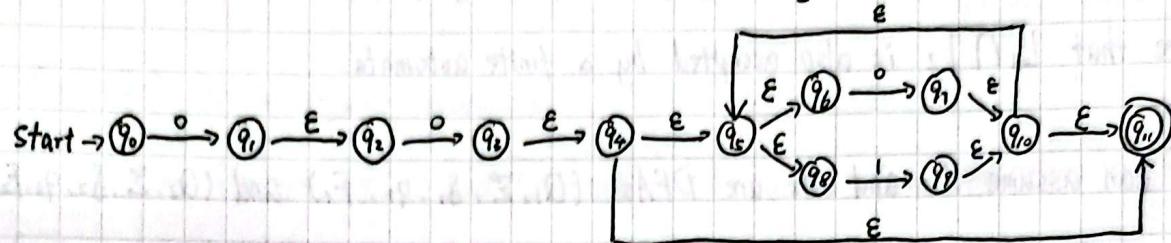


Week 3 · By Deng Yufan

Prob 1. Convert $(010+)^*$ to ϵ -NFA using Thompson's method.



Prob 2. Let $A = \{Q, \Sigma, \delta, q_0, q_f\}$ be a ϵ -NFA. st. no trans into q_0 or out of q_f .

Describe the language using $L = L(A)$.

- a. A adding an ϵ -trans from q_f to q_0 .

LL^* .

- b. A adding an ϵ -trans from q_0 to every state reachable from q_0 .

A suffix of ϵ . every string in L , which can be ϵ .

- c. A adding an ϵ -trans from every state reachable ~~to~~ to q_f to q_f .

A prefix of every string in L , which can be ϵ .

- d. A doing both b. and c.

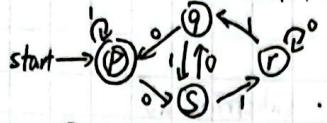
A range (or continuous substring) of every string in L , which can be ϵ .

Prob 3. Convert the DFA to RE using state-elimination.

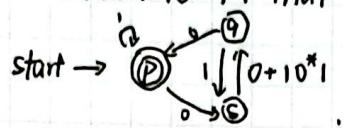
	0	1
$\rightarrow *p$	s	P
q	P	s
r	r	q
s	q	r



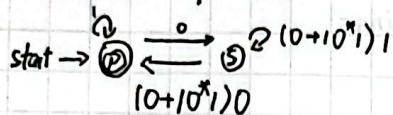
Step 1. Use RE as labels, we have



Step 2. Eliminate r , that is



Step 3. Eliminate q , that is



Step 4. Eliminate s , that is

$$\text{start} \rightarrow P \xrightarrow{b} 1 + 0 [(0+10^*1)]^* (0+10^*1) 0$$

Step 5. Final answer.

$$[1 + 0 [(0+10^*1)]^* (0+10^*1) 0]^*$$

Prob 4. Prove or disprove the following statements.

a. $(\epsilon + R)^* = R^*$. (true).

We are about to prove $(\epsilon + a)^* = a^*$. First we have $\epsilon + a = a + \epsilon$, so $(\epsilon + a)^n = \bigcup_{i=0}^{n-1} \epsilon^i a^{n-i}$

$$\text{Then } (\epsilon + R)^* = \bigcup_{n=0}^{+\infty} (\epsilon + R)^n = \bigcup_{n=0}^{+\infty} \bigcup_{i=0}^n (\cancel{\epsilon + R})^i a^{n-i} = \bigcup_{i=0}^{+\infty} a^i = a^*.$$

b. $(RS + R)^* R = R(SR + R)^*$. (true).

We are about to prove $(ab+a)^* a = a(ab+a)^*$.

First we have $(b+\epsilon)(ab+a) = bab + ba + ab + a = (ba+a)(b+\epsilon)$. So $(b+\epsilon)(ab+a)^n = (ba+a)(b+\epsilon)^n$

$$\text{Then for } i \geq 1, (ab+a)^i a = a(b+\epsilon)(ab+a)^{i-1} a = a(ab+a)^{i-1} (b+\epsilon)a = a(ab+a)^{i-1} \cdots = (ba+a)^{i-1} (b+\epsilon)$$

$$\text{So } (ab+a)^* a = \bigcup_{i=0}^{+\infty} (ab+a)^i a = \bigcup_{i=0}^{+\infty} a(ab+a)^{i-1} = a(ab+a)^*.$$

c. $(R+S)^* S = (R^* S)^*$. (~~false~~). (false).

We are about to prove $(ab)^* b \neq (a^* b)^*$.

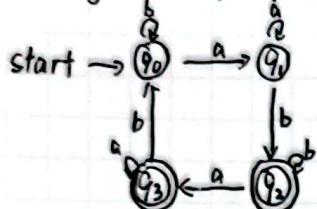
~~We are about to prove $(ab)^* b$ and $(a^* b)^*$ both refer to all strings end with~~

Obviously $\epsilon \notin (a+b)^* b$, since ϵ don't end with b , but $\epsilon \in (a^* b)^*$. So $(a+b)^* b \neq (a^* b)^*$,

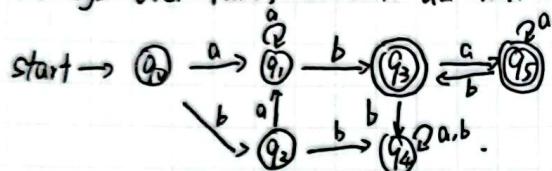


Prob 5. Construct DFA.

a. Strings over $\{a,b\}$ have and only have odd number of substring ab.



b. Strings over $\{a,b\}$ contain ab but not bb as substring.



Prob 6. Construct RE.

a. $\{xwx^2 \mid x, w \in (a+b)^*\}$

$$\cancel{a(a+b)(a+b)^*b} \cdot a(a+b)(a+b)^*a + b(a+b)(a+b)^*b.$$

b. $\{w \mid w \in \{0,1\}^*, \text{ has at least 3 } 1 \text{ and the third from the last position is a } 1\}$.

$$(0+1)^*111 + (0+1)^*10^*1(01+10) + (0+1)^*10^*10^*100.$$



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