

Practical Analysis

Assume array A is indexed from 1 to n .

INEFFICIENT SORT(A, n)

1. **for** $i = 1$ **to** $n!$ **do**
2. Boolean $sortedSoFar = \text{TRUE}$;
3. $j = 1$;
4. $P = \text{nextPermutation}(A)$; // $\theta(n)$
5. **while** $j < n$ **and** $sortedSoFar$ **do**
6. **if** $P[j] > P[j + 1]$
7. **then** $sortedSoFar = \text{FALSE}$ // $O(n)$
8. $j++$
9. **if** ($sortedSoFar$) **then output** P // $\theta(n)$

Analyze the worst-case complexity of INEFFICIENT SORT assuming that the nextPermutation function always takes $\Theta(n)$ time.

Complexity: $\theta(n \cdot n!) = \theta(n!)$

$\lim(n!/(n \cdot n!)) = \lim(1/n) = 0$

Your answer should fit above the line!