

## Linear Partition Worksheet

Define  $M[n, k]$  to be the minimum possible cost over all partitions of  $(s_1, \dots, s_n)$  into  $k$  ranges.

1. Consider the input (100, 200, 300, 400, 500, 600, 700) with  $k = 3$ . What is  $M[7, 3]$ ?  
(Hint: It should be possible to answer this by visual inspection.)

100 200 300 | 500 600 | 700

$$M[7, 3] = 1,100$$

2. What are  $M[1, 2]$ ,  $M[2, 2]$ ,  $M[3, 2]$ ,  $M[4, 2]$ ,  $M[5, 2]$ ,  $M[6, 2]$ , and  $M[7, 2]$ ? (again, use visual inspection).

$$M[1, 2] = 100$$

$$M[2, 2] = 200$$

$$M[3, 2] = 300$$

$$M[4, 2] = 600$$

$$M[5, 2] = 900$$

$$M[6, 2] = 1,100$$

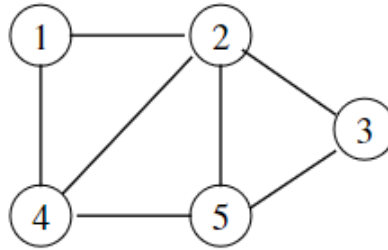
$$M[7, 2] = 1,500$$

3. Can you write a formula for  $M[7, 3]$  in terms of  $M[1, 2]$ ,  $M[2, 2]$ ,  $M[3, 2]$ ,  $M[4, 2]$ ,  $M[5, 2]$ ,  $M[6, 2]$ , and  $M[7, 2]$ ?

$$M[n, k] = \min\{\max(M[i, k-1], \sum_{j=i+1}^n s_j)\}$$

$$M[7, 3] = \max(M[6, 2], s[7]) = 1,100$$

## Graph Representation Worksheet



- What are the storage requirements assuming an adjacency matrix is used. Assume each element of the adjacency matrix requires four bytes.

	1	2	3	4	5
1	0	1	0	1	0
2	1	0	1	1	1
3	0	1	0	0	1
4	1	1	0	0	1
5	0	1	1	1	0

$$25 \times 4 = 100 \text{ bytes}$$

- Repeat for an adjacency list representation. Assume that an int requires 4 bytes and that a pointer also requires 4 bytes.

1	<input type="text"/>	→	2 -> 4
2	<input type="text"/>	→	1 -> 4 -> 5 -> 3
3	<input type="text"/>	→	2 -> 5
4	<input type="text"/>	→	1 -> 2 -> 5
5	<input type="text"/>	→	4 -> 2 -> 3

$$5 \times 4 = 20$$

$$14 \text{ nodes} \times 8 = 112$$

$$20 + 112 = 132$$

- Now, consider an undirected graph with 100 vertices and 1000 edges. What are the storage requirements for the adjacent matrix and adjacency list data structures?

$$\text{AM: } 4n^2 \Rightarrow 4 \times 100^2 = 40,000$$

$$\text{AL: } 4n + 8(2m) = 4n + 16m \Rightarrow 4 \times 100 + 16 \times 1000 = 16,400$$

Adjacency list is better than adjacency matrices