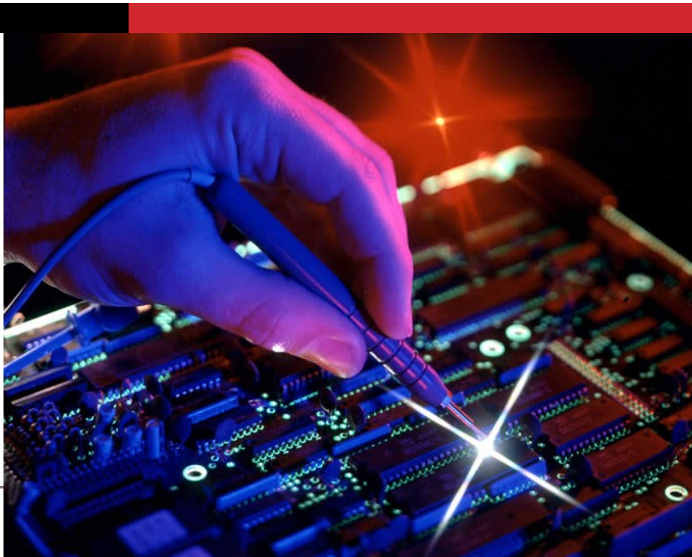
Advanced Computer Architecture

SEQUENTIAL LOGIC DESIGN

Dennis A. N. Gookyi

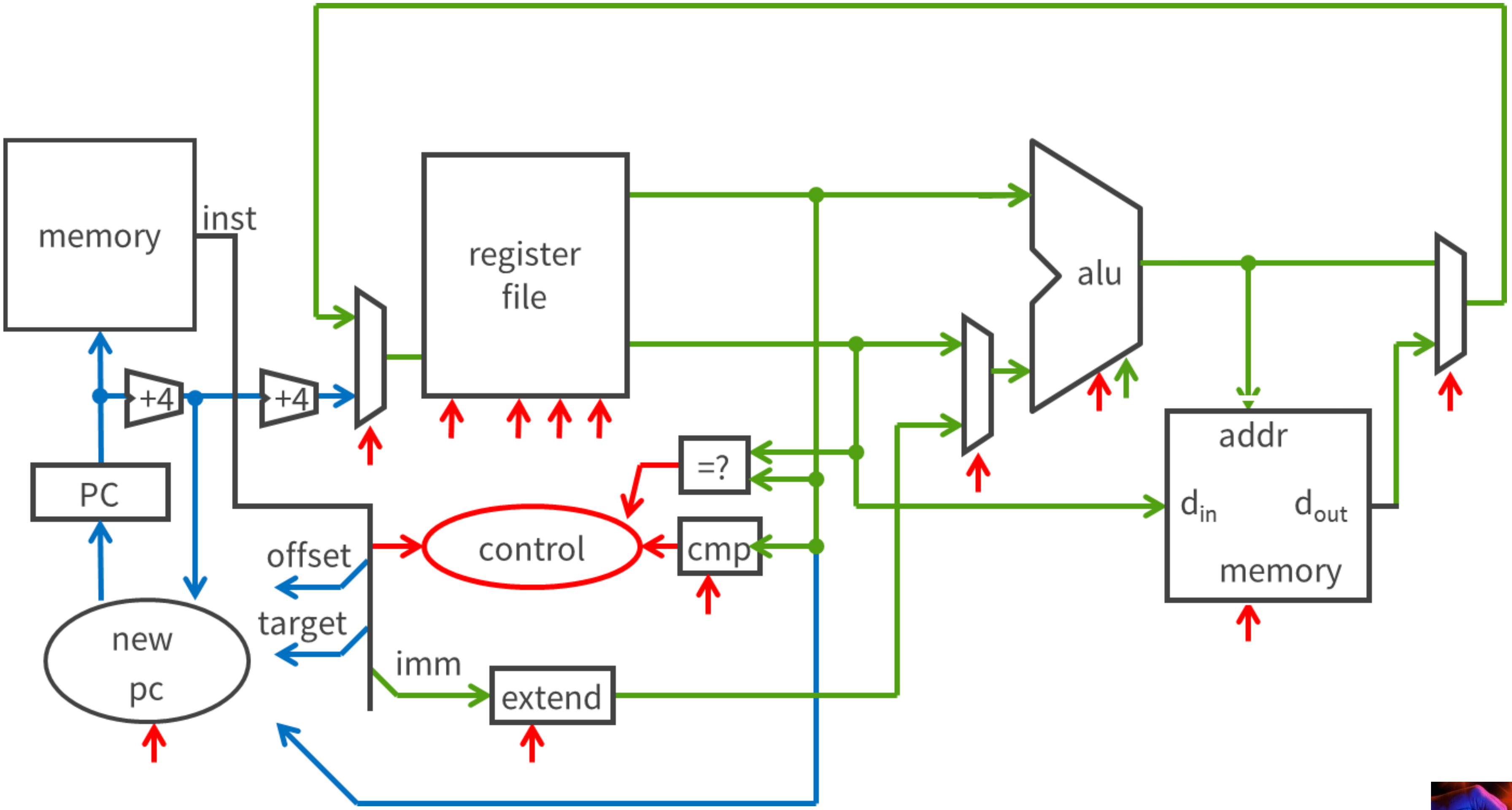


CONTENTS

- Sequential Logic Design

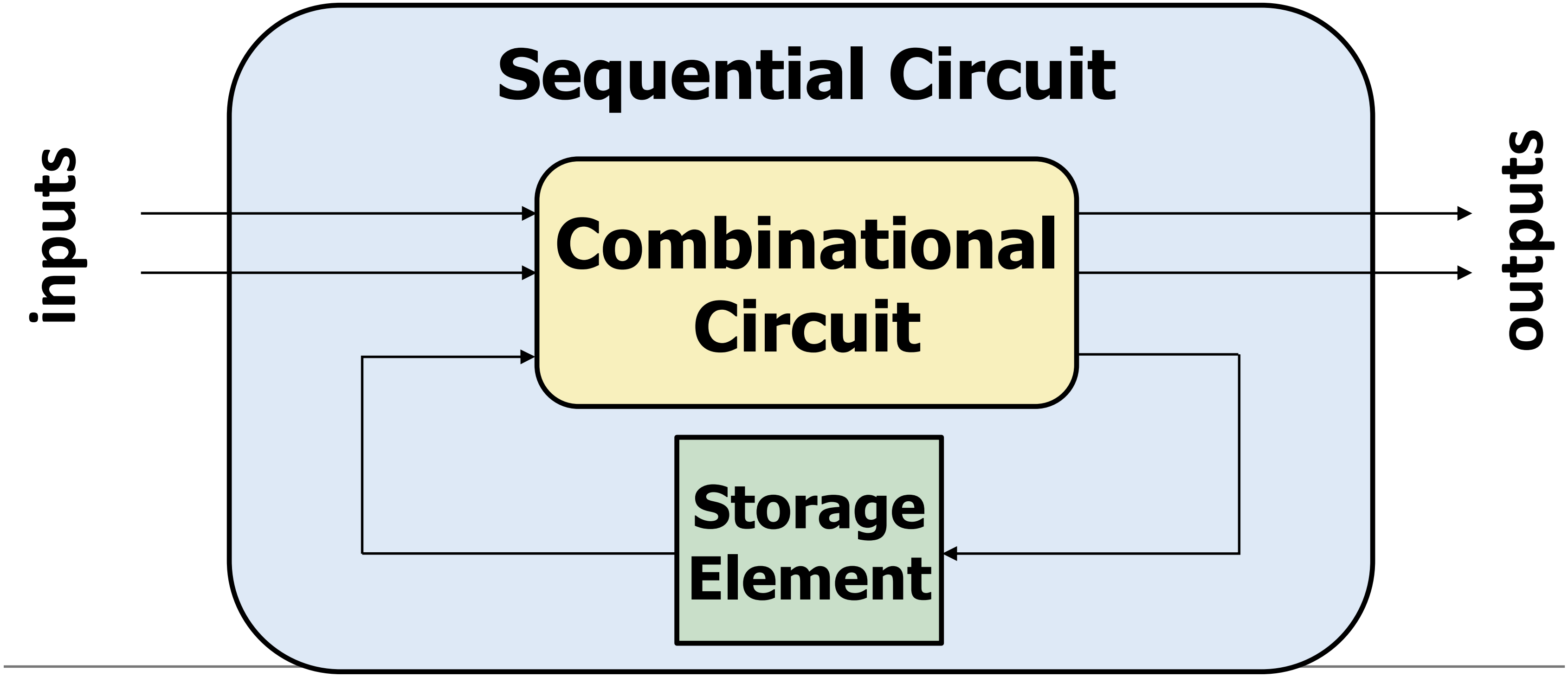
BIG PICTURE: BUILDING A PROCESSOR

- Single cycle processor



INTRODUCTION

- Combinational circuit output depends only on current input
- We want circuits that produce output depending on current and past input values – circuits with memory
- How can we design a circuit that stores information?



SEQUENTIAL LOGIC CIRCUITS

- We have examined designs of circuit elements that can store information
- Now, we will use these elements to build circuits that remember past inputs

Combinational
Only depends on current inputs

Sequential
Opens depending on past inputs

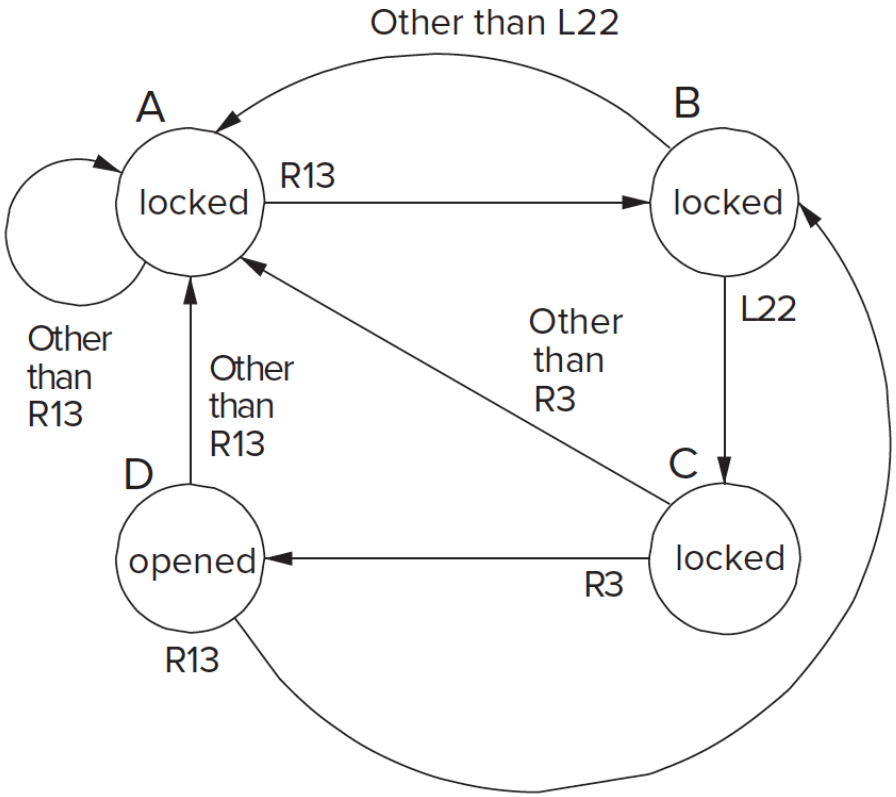
STATE

- **In order for this lock to work, it has to keep track (remember) of the past events!**
- **If passcode is R13-L22-R3, sequence of states to unlock:**
  A. The lock is not open (locked), and no relevant operations have been performed
  B. Locked but user has completed R13
  C. Locked but user has completed R13-L22
  D. Unlocked: user has completed R13-L22-R3

- **The state of a system is a snapshot of all relevant elements of the system at the moment of the snapshot**
  - □ To open the lock, **states A-D must be completed in order**
  - □ If anything else happens (e.g., L5), lock **returns** to state A

STATE DIAGRAM

- Completely describes the operation of the sequential lock



Passcode: **R13-L22-R3**

- We will understand "state diagrams" fully later today

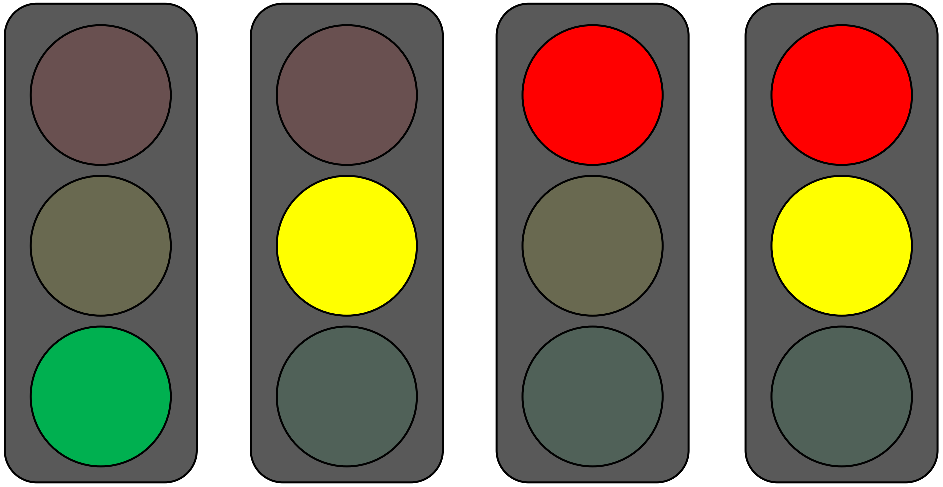
STATE DIAGRAM

- Asynchronous vs. Synchronous State Changes
  - □ Sequential lock we saw is an asynchronous "machine"
    - State transitions occur when they occur
    - There is nothing that synchronizes when each state transition must occur
  - □ Most modern computers are synchronous "machines"
    - State transitions take place after fixed units of time
    - Controlled in part by a clock, as we will see soon
  - □ These are two different design paradigms, with tradeoffs

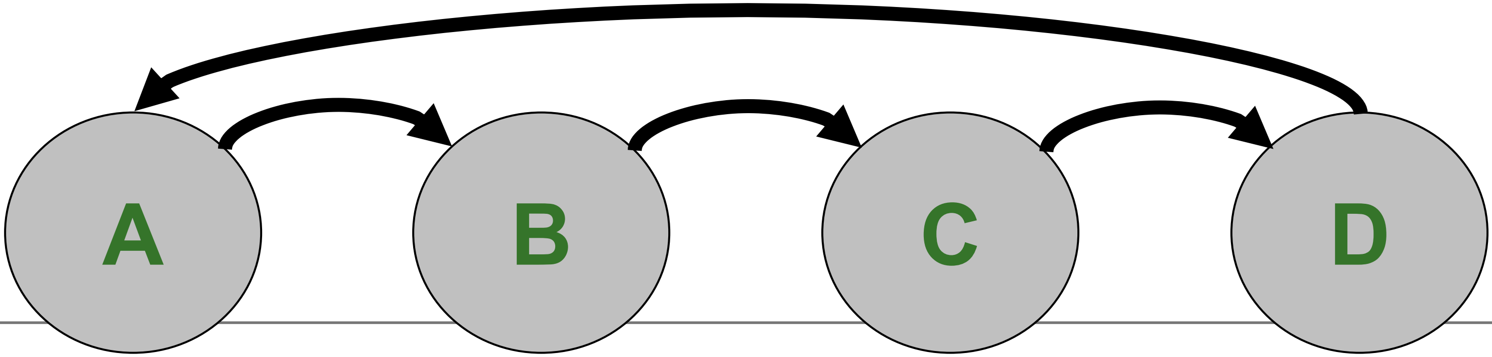
STATE DIAGRAM

- A standard Swiss traffic light has 4 states

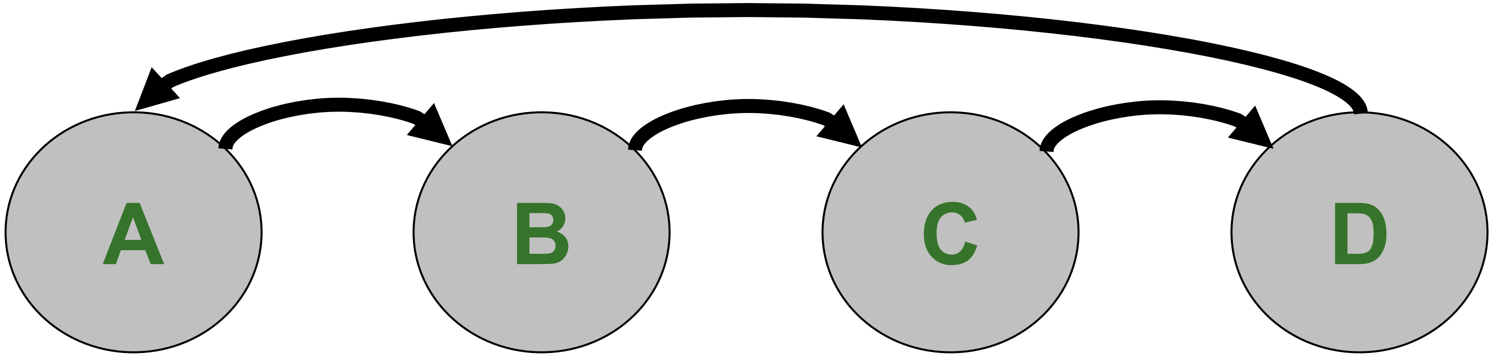
A. Green
B. Yellow
C. Red
D. Red and Yellow



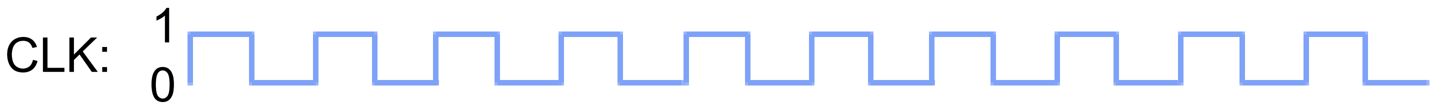
- The sequence of these states are always as follows



STATE DIAGRAM - CLOCK



- When should the light change from one state to another?
- We need a clock to dictate when to change state
  - □ Clock signal alternates between 0 & 1



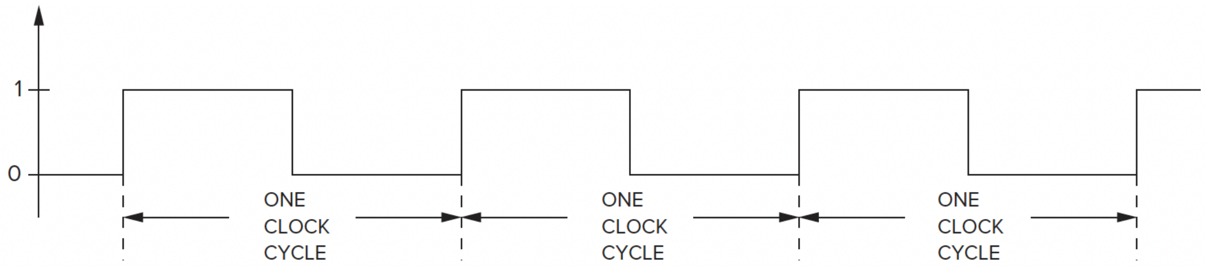
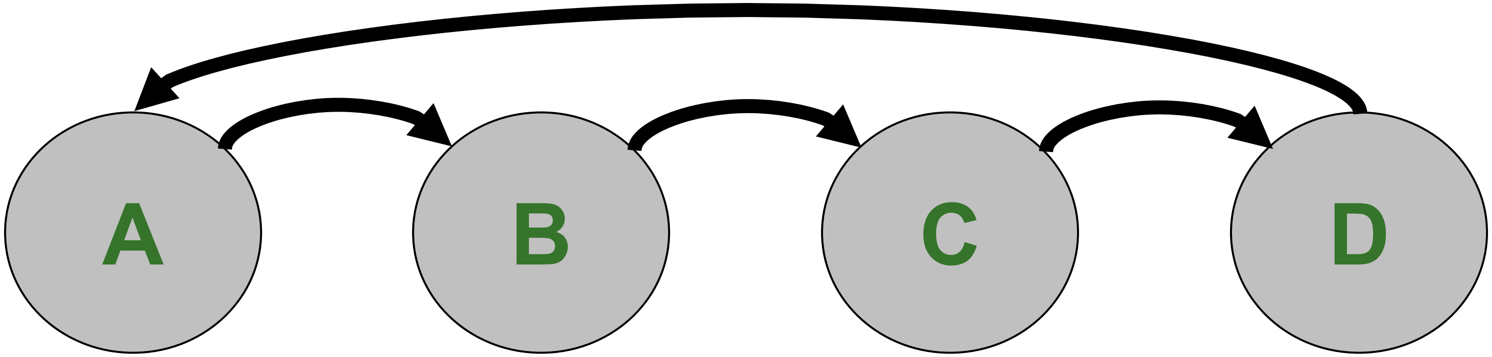


Figure 3.28 A clock signal.

STATE DIAGRAM - CLOCK



- When should the light change from one state to another?
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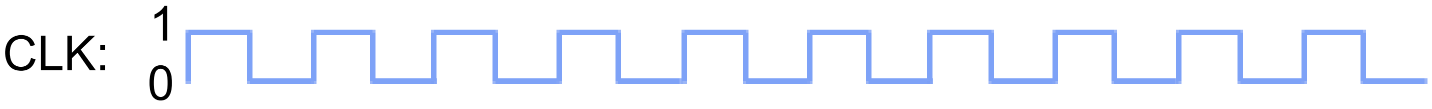
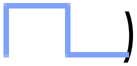


- At the start of a clock cycle (), system state changes
  - □ During a clock cycle, the state stays constant
  - □ In this traffic light example, we are assuming the traffic light stays in each state an equal amount of time

STATE DIAGRAM - CLOCK

- Clock is a general mechanism that triggers transition from one state to another in a (synchronous) sequential circuit

- Clock synchronizes state changes across many sequential circuit elements

- Combinational logic evaluates for the length of the clock cycle

- Clock cycle should be chosen to accommodate maximum combinational circuit delay

STATE DIAGRAM

- Asynchronous vs. Synchronous State Changes
  - □ Sequential lock we saw is an asynchronous "machine"
    - State transitions occur when they occur
    - There is nothing that synchronizes when each state transition must occur
  - □ Most modern computers are synchronous "machines"
    - State transitions take place after fixed units of time
    - Controlled in part by a clock, as we will see soon
  - □ These are two different design paradigms, with tradeoffs
    - Synchronous control can be easier to get correct when the system consists of many components and many states
    - Asynchronous control can be more efficient (no clock overheads)

FINITE STATE MACHINE (FSM)

- What is a Finite State Machine (FSM)?
  - A discrete-time model of a stateful system
  - Each state represents a snapshot of the system at a given time

- An FSM pictorially shows
  - The set of all possible states that a system can be in
  - How the system transitions from one state to another

- An FSM can model
  - A traffic light, an elevator, fan speed, a microprocessor, etc.

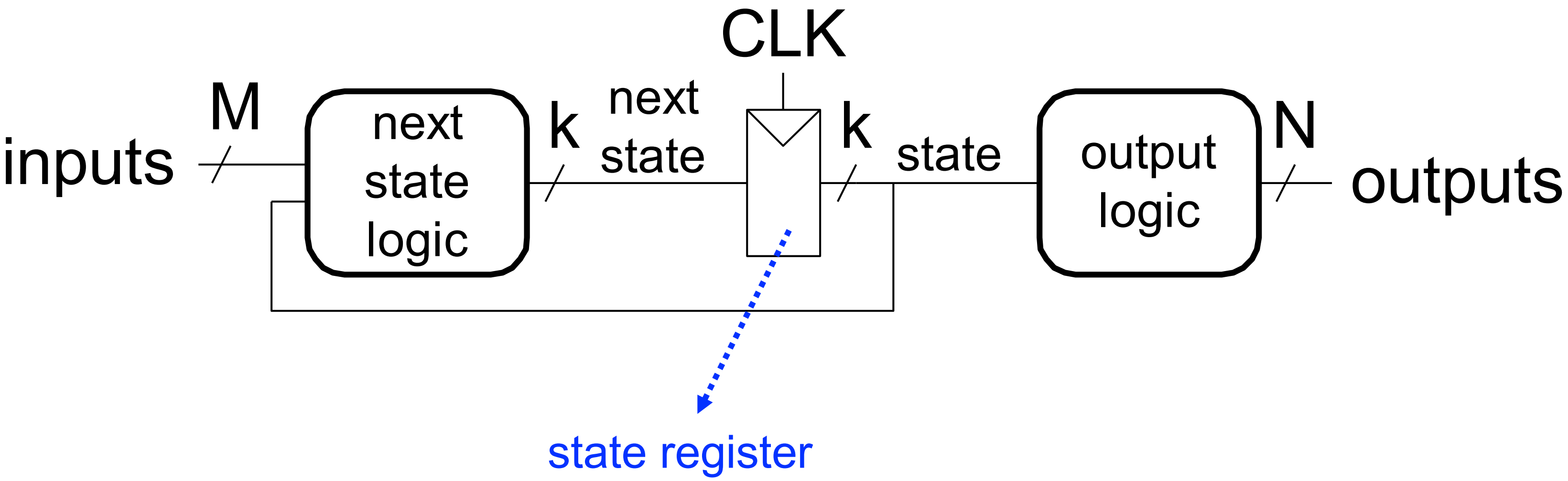
- An FSM enables us to pictorially think of a stateful system using simple diagrams

FINITE STATE MACHINE (FSM)

- ❖ FSM consist of five elements:
  - □ A finite number of states
    - State: snapshot of all relevant elements of the system at the time of the snapshot
  - □ A finite number of external inputs
  - □ A finite number of external outputs
  - □ An explicit specification of all state transitions
    - How to get from one state to another
  - □ An explicit specification of what determines each external output value

FINITE STATE MACHINE (FSM)

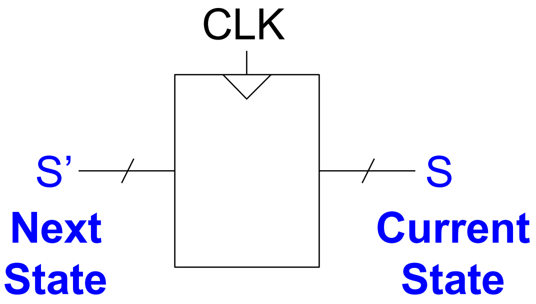
- Each FSM consists of three separate parts:
  - □ next state logic
  - □ state register
  - □ output logic



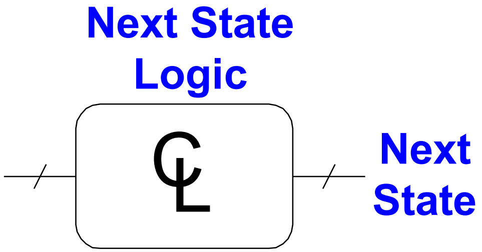
At the beginning of the clock cycle, next state is latched into the state register

FINITE STATE MACHINE (FSM)

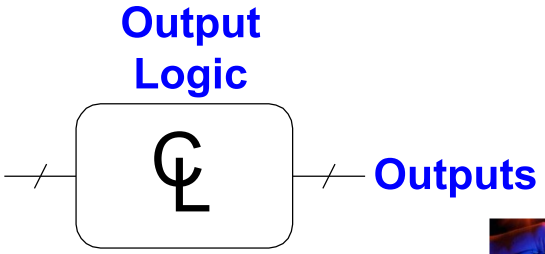
- Sequential Circuits
  - □ State register(s)
    - Store the current state and
    - Load the next state at the clock edge



- Combinational Circuits
  - □ Next state logic
    - Determines what the next state will be



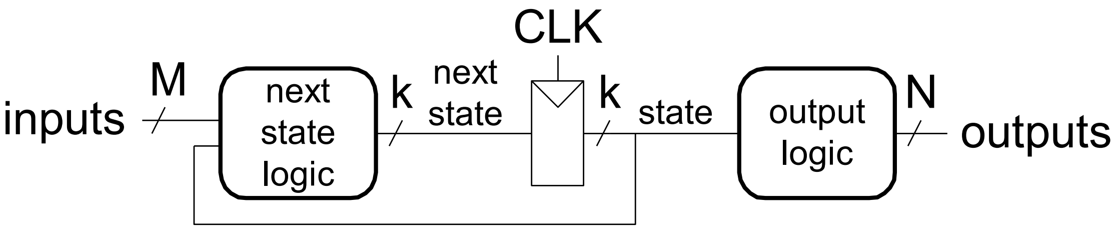
  - □ Output logic
    - Generates the outputs



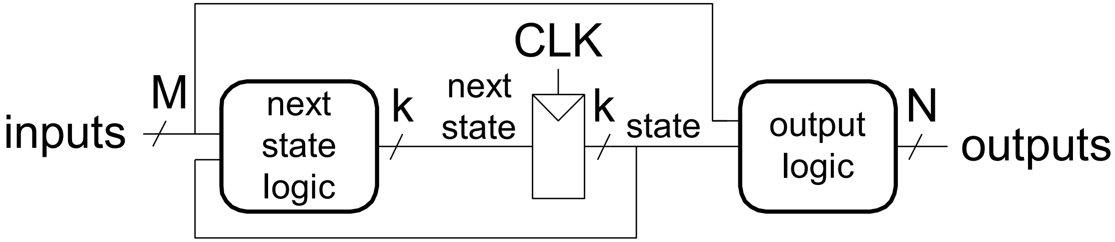
FINITE STATE MACHINE (FSM)

- Next state is determined by the current state and the inputs
- Two types of finite state machines differ in the output logic:
  - □ Moore FSM: outputs depend only on the current state
  - □ Mealy FSM: outputs depend on the current state and the inputs

Moore FSM

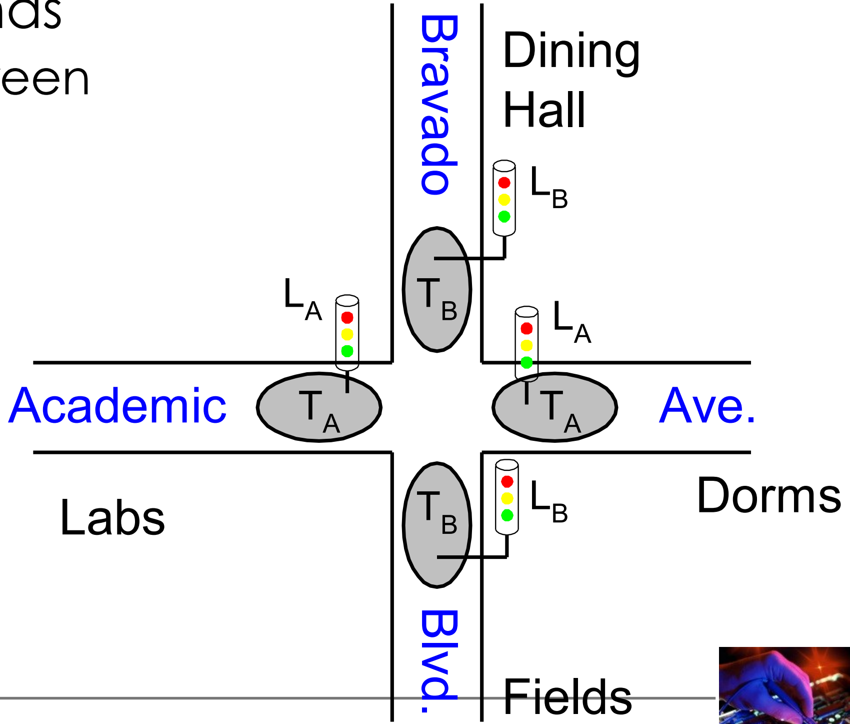


Mealy FSM



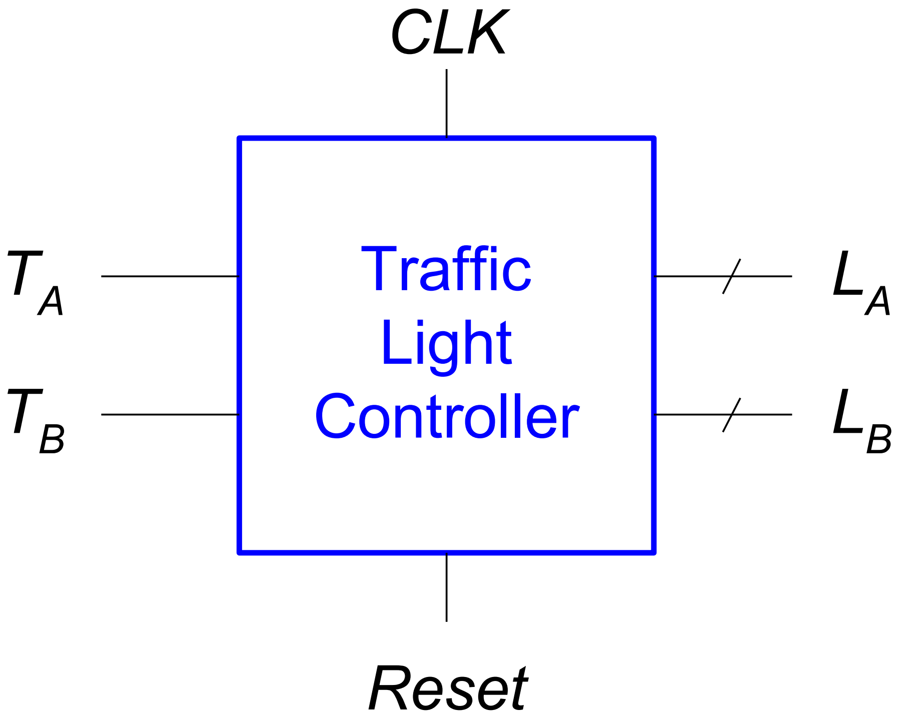
FINITE STATE MACHINE (FSM)

- Example: "Smart" traffic light controller
  - □ 2 inputs:
    - Traffic sensors: T_A , T_B (TRUE when there's traffic)
  - □ 2 outputs:
    - Lights: L_A , L_B (Red, Yellow, Green)
  - □ State can change every 5 seconds
    - Except if green and traffic, stay green



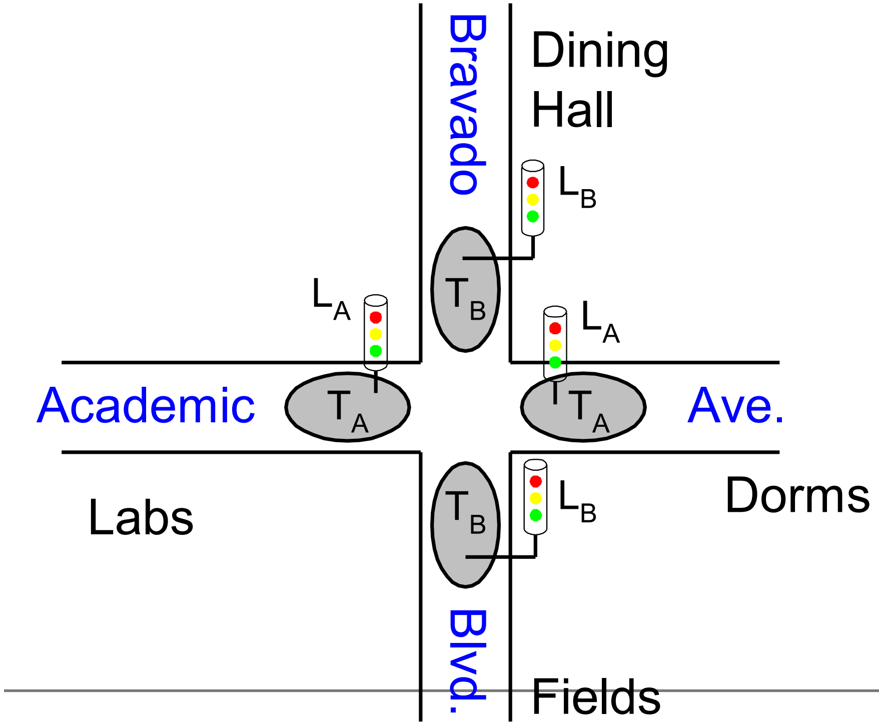
FINITE STATE MACHINE (FSM)

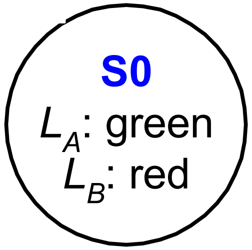
- Inputs: CLK, Reset, T_A , T_B
- Outputs: L_A , L_B



FINITE STATE MACHINE (FSM)

- Moore FSM: outputs labeled in each state
  - □ States: Circles
  - □ Transitions: Arcs





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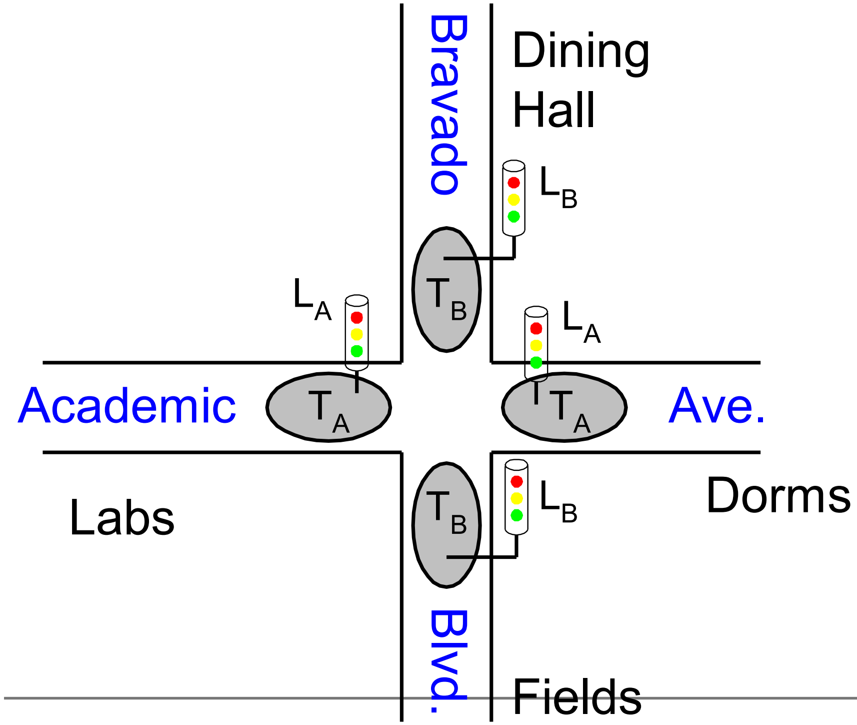

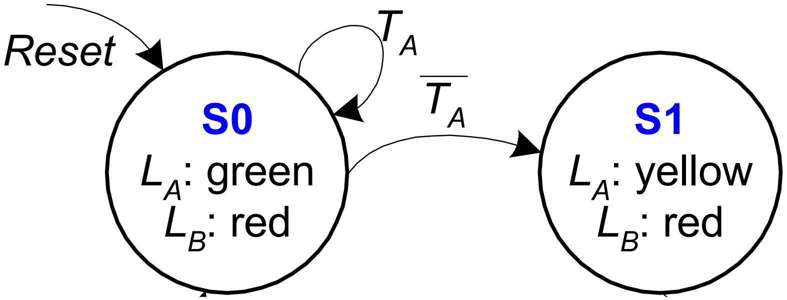

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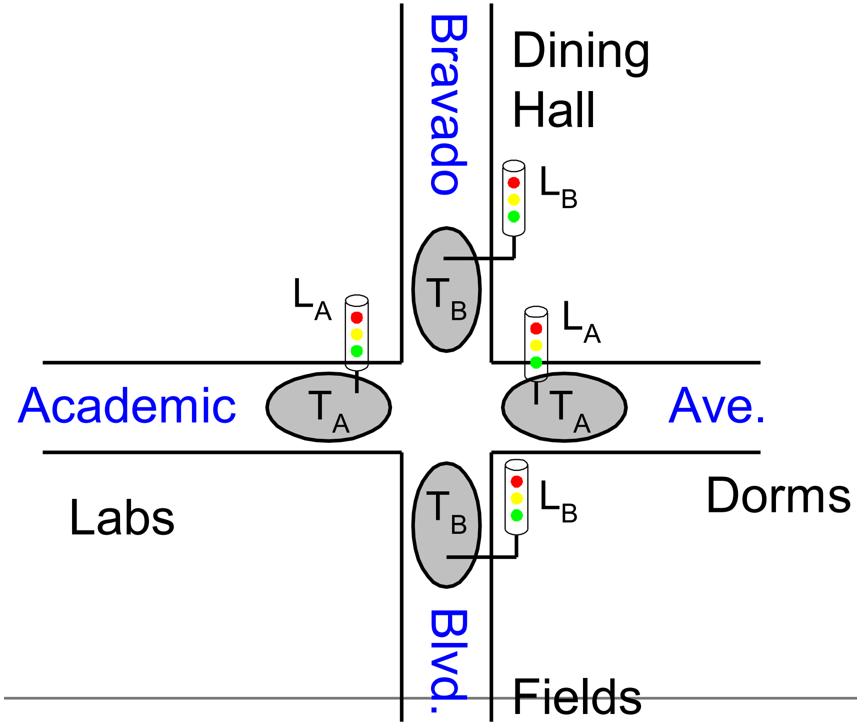

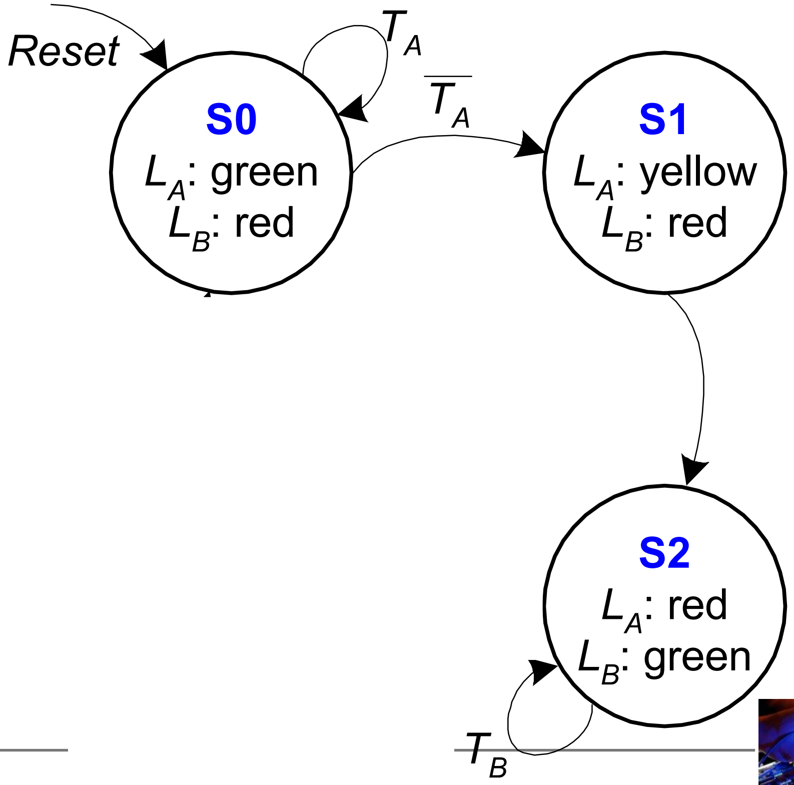

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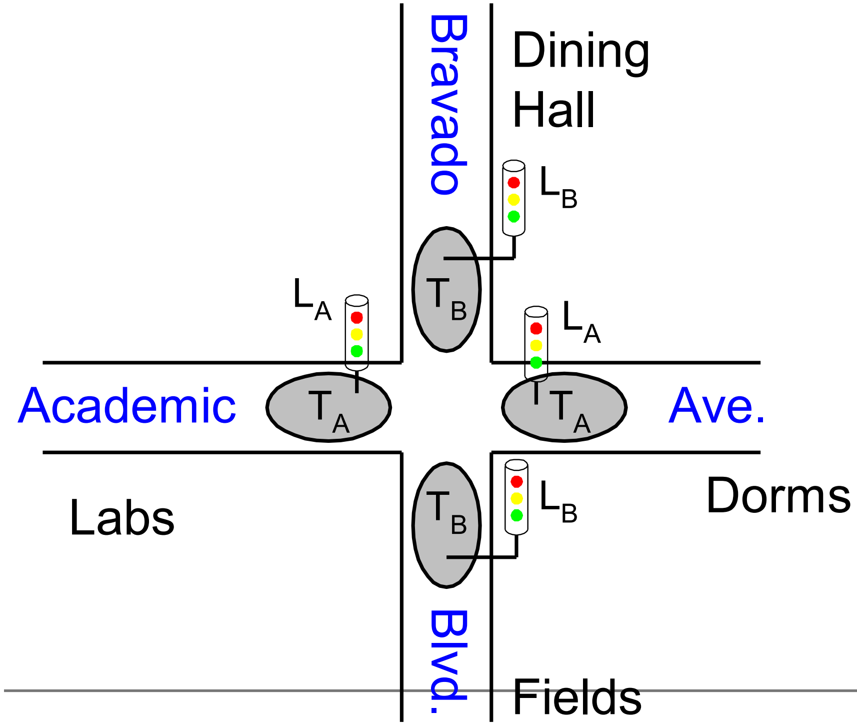

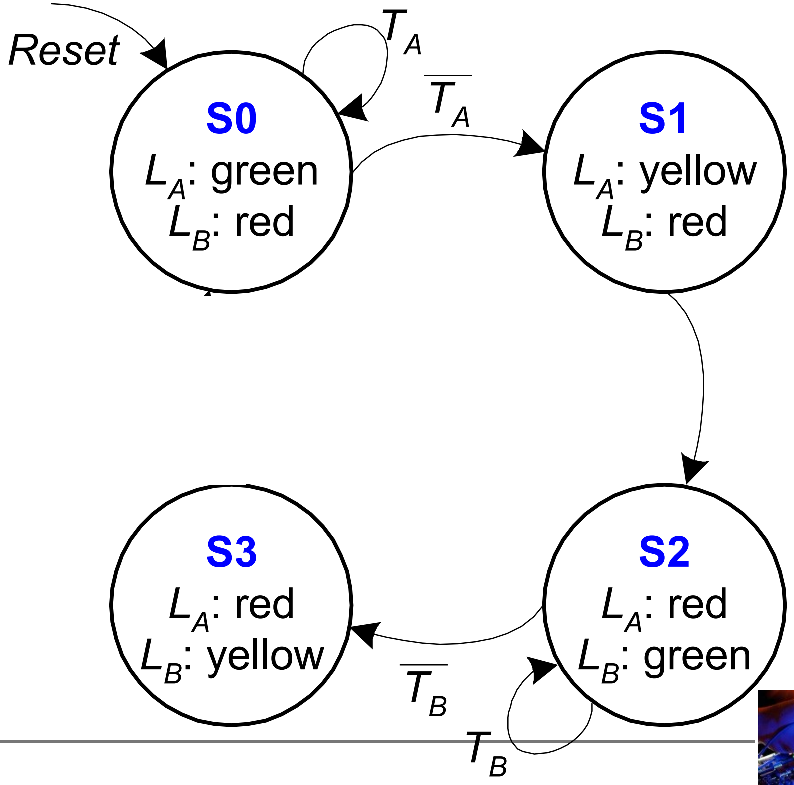

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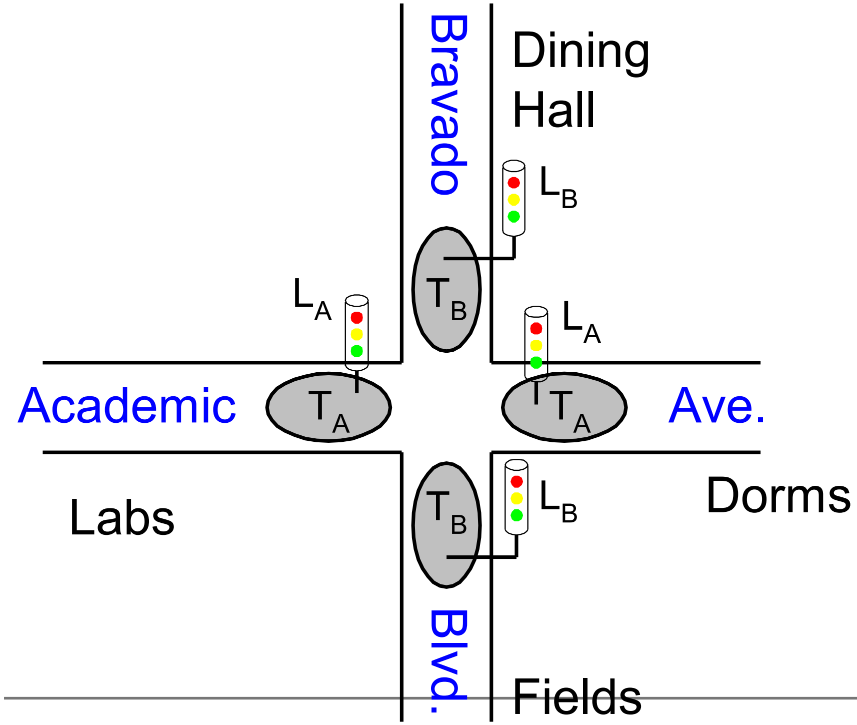
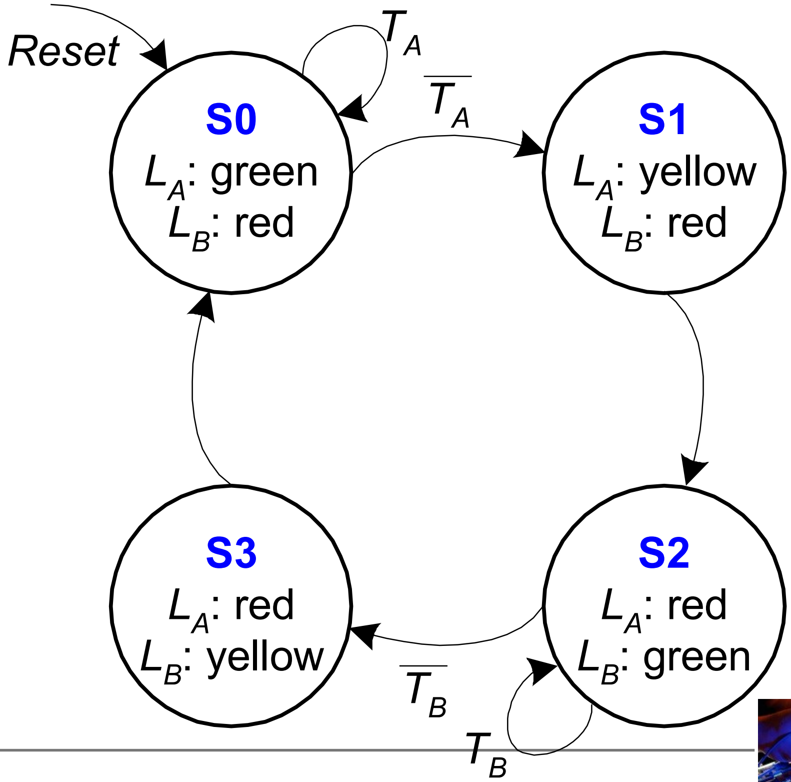
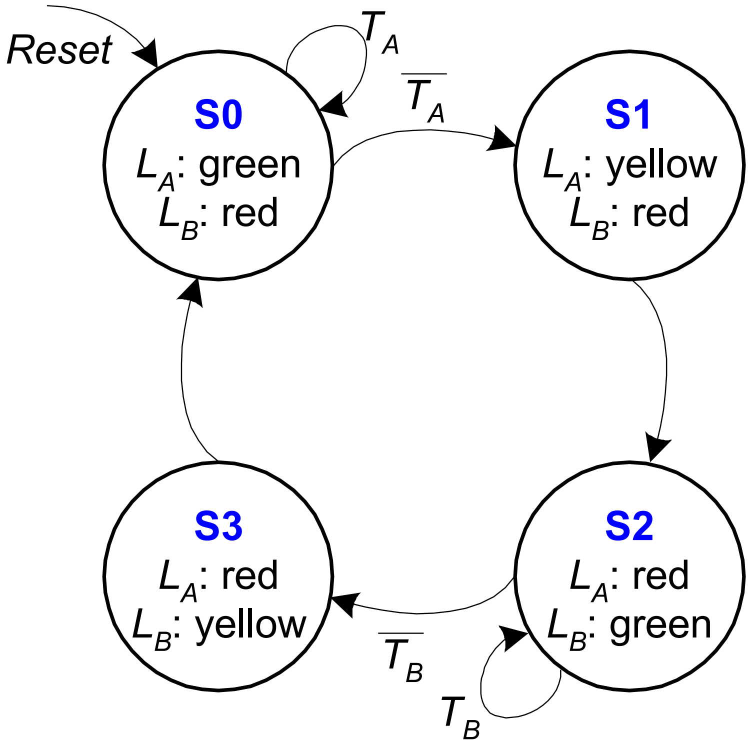


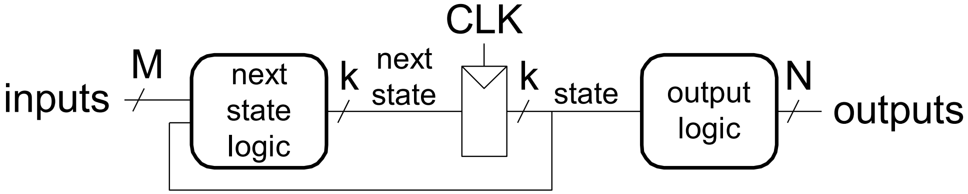


FSM – STATE TRANSITION TABLE

- State transition table



Current State	Inputs		Next State
S	T_A	T_B	S'
S0	0	X	
S0	1	X	
S1	X	X	
S2	X	0	
S2	X	1	
S3	X	X	



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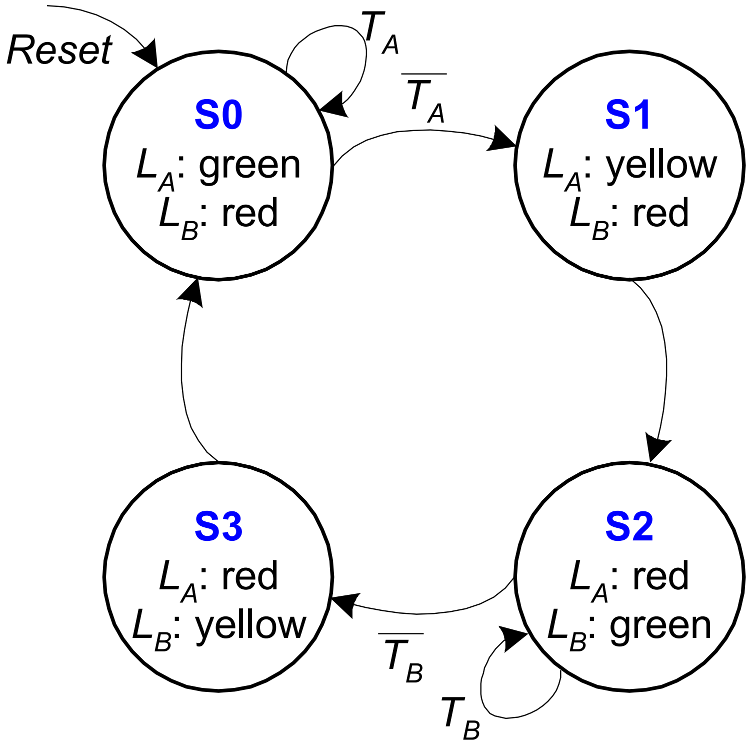
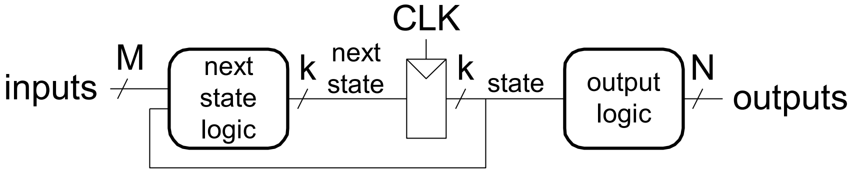


Current State	Inputs		Next State
S	T_A	T_B	S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0



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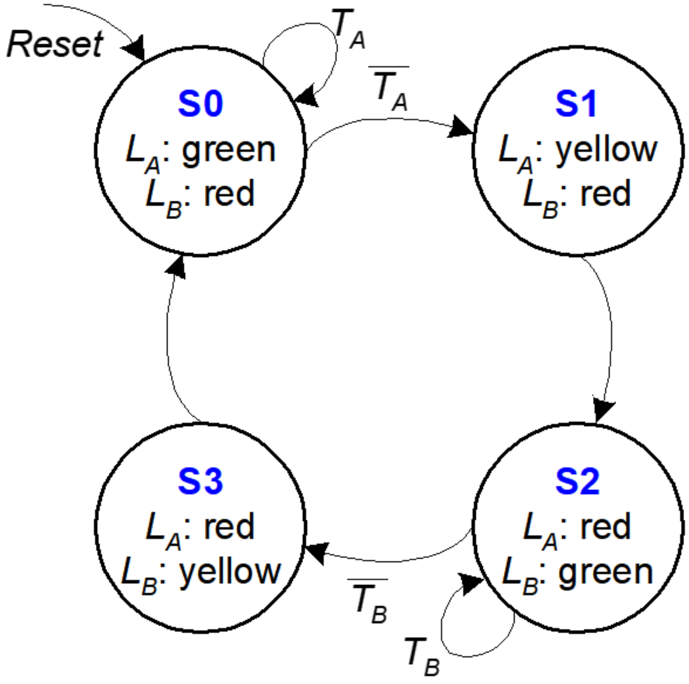
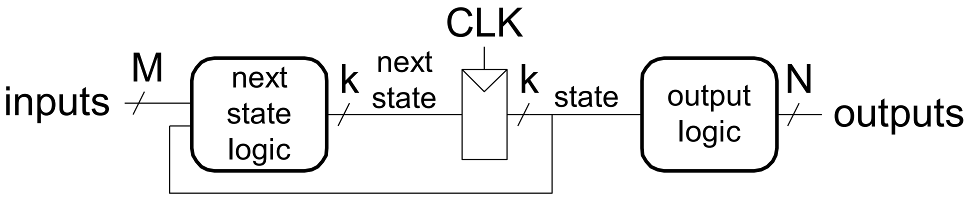


Current State	Inputs		Next State
S	T_A	T_B	S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0

State	Encoding
S0	00
S1	01
S2	10
S3	11



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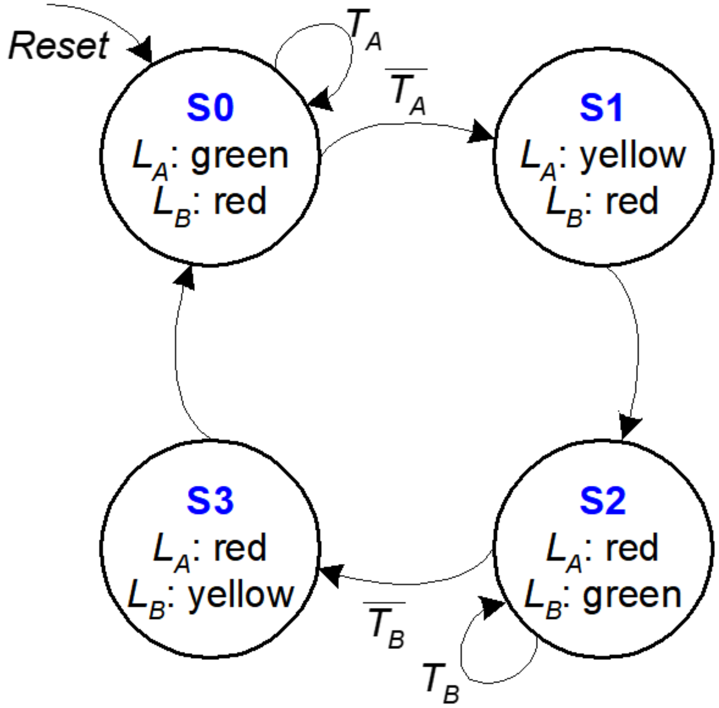
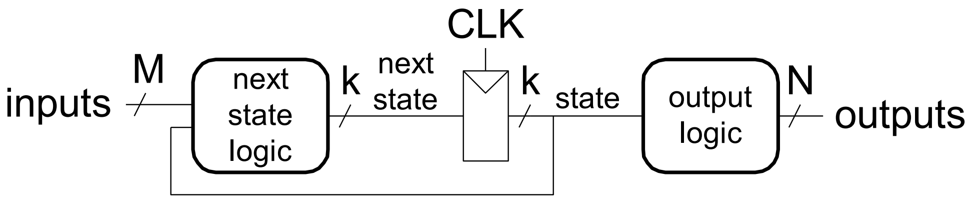


Current State		Inputs		Next State	
S_1	S_0	T_A	T_B	S'_1	S'_0
0	0	0	X	0	1
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0



State	Encoding
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FSM – STATE TRANSITION TABLE

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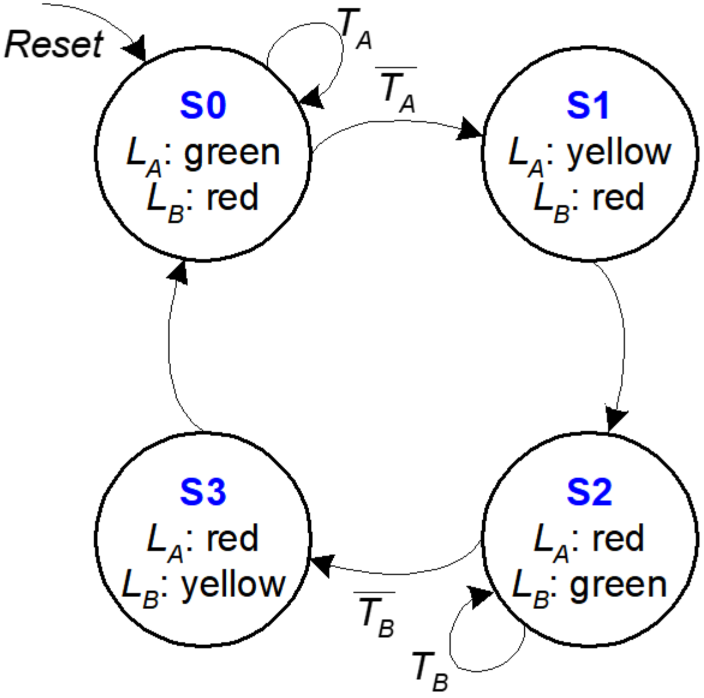

Current State		Inputs		Next State	
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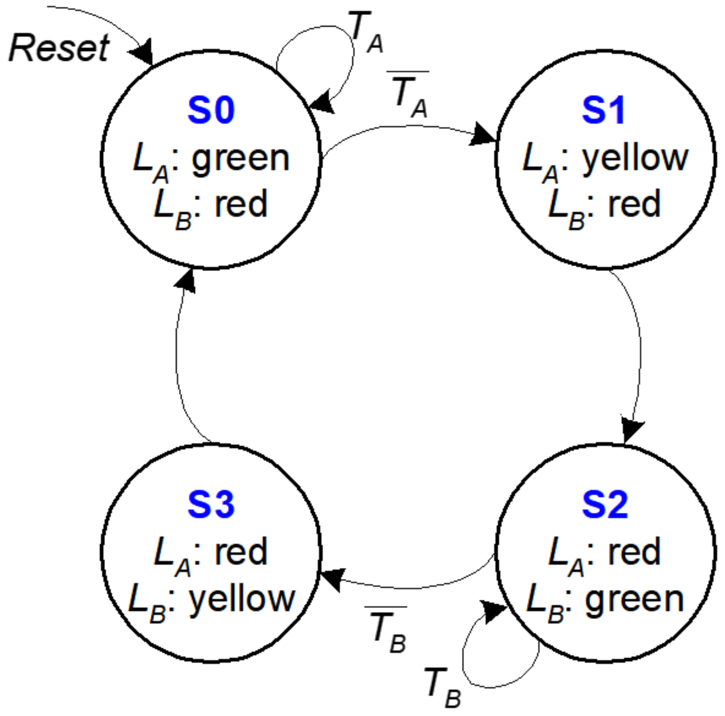
0	0	0	X	0	1
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

State	Encoding
S0	00
S1	01
S2	10
S3	11

$S'_1 = ?$

FSM – STATE TRANSITION TABLE

- State transition table



Current State		Inputs		Next State	
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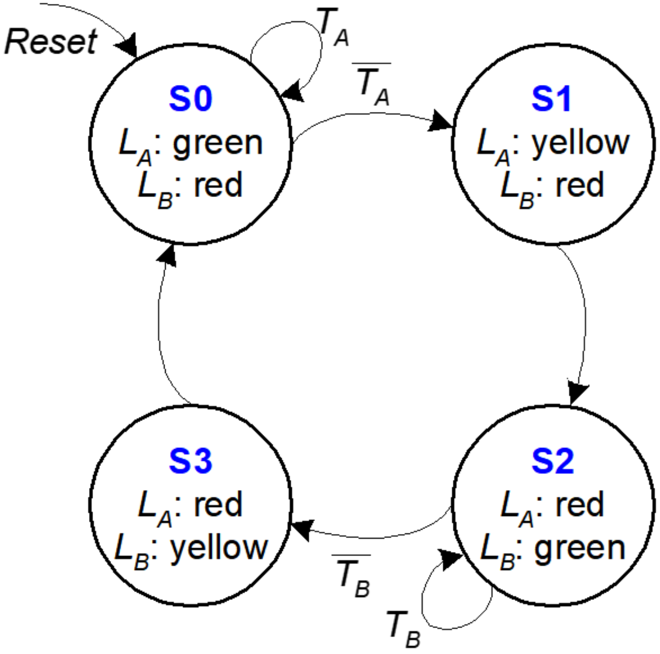
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

State	Encoding
S0	00
S1	01
S2	10
S3	11

$$S'_1 = (\overline{S_1} \cdot S_0) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B}) + (S_1 \cdot \overline{S_0} \cdot T_B)$$

FSM – STATE TRANSITION TABLE

- State transition table



Current State		Inputs		Next State	
S_1	S_0	T_A	T_B	S'_1	S'_0
0	0	0	X	0	1
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

State	Encoding
S0	00
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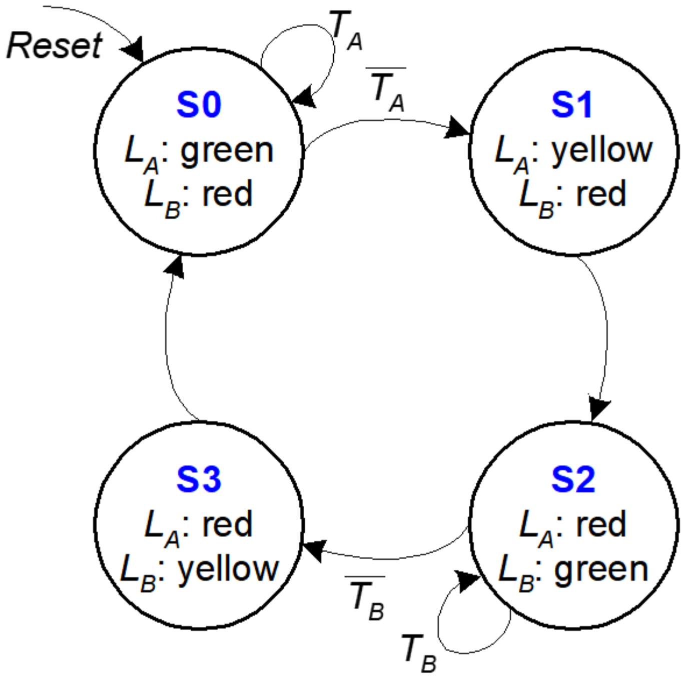
S2	10
S3	11

$$S'_1 = (\overline{S_1} \cdot S_0) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B}) + (S_1 \cdot \overline{S_0} \cdot T_B)$$

$$S'_0 = ?$$

FSM – STATE TRANSITION TABLE

- State transition table



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0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

State	Encoding
S0	00
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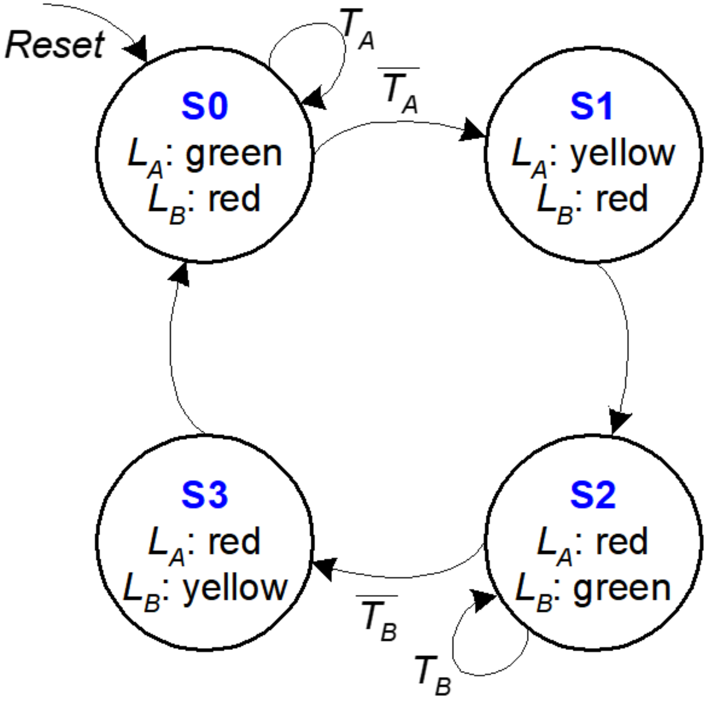
$$S'_1 = (\overline{S_1} \cdot S_0) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B}) + (S_1 \cdot \overline{S_0} \cdot T_B)$$

$$S'_0 = (\overline{S_1} \cdot \overline{S_0} \cdot \overline{T_A}) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B})$$

FSM – STATE TRANSITION TABLE

- State transition table



Current State		Inputs		Next State	
S_1	S_0	T_A	T_B	S'_1	S'_0
0	0	0	X	0	1
0	0	1	X	0	0
0	1	X	X	1	0
1	0	X	0	1	1
1	0	X	1	1	0
1	1	X	X	0	0

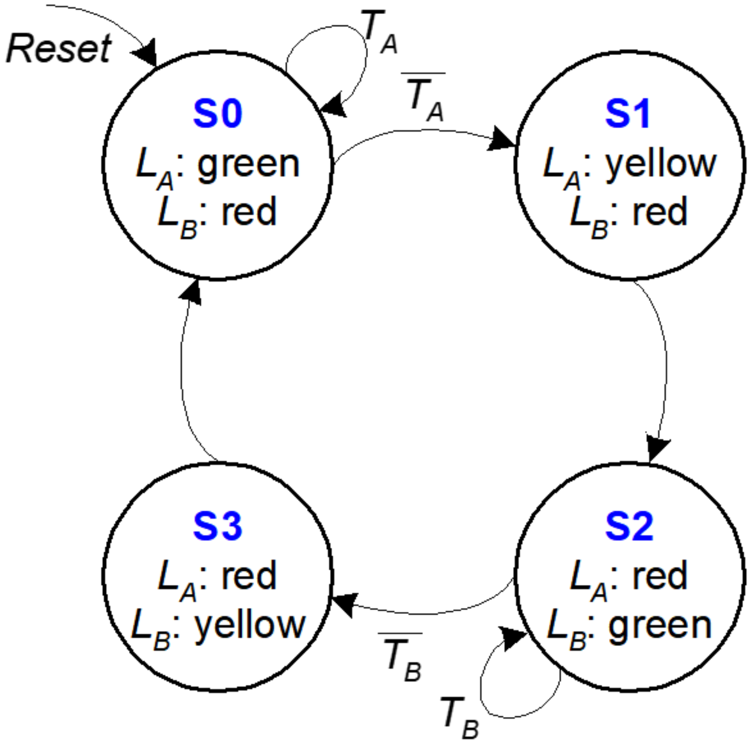
State	Encoding
S0	00
S1	01
S2	10
S3	11

$S'_1 = S_1 \text{ xor } S_0$ **(Simplified)**

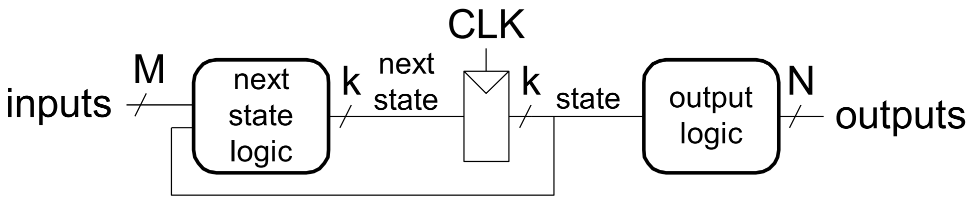
$S'_0 = (\overline{S_1} \cdot \overline{S_0} \cdot \overline{T_A}) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B})$

FSM – OUTPUT TABLE

- Output table

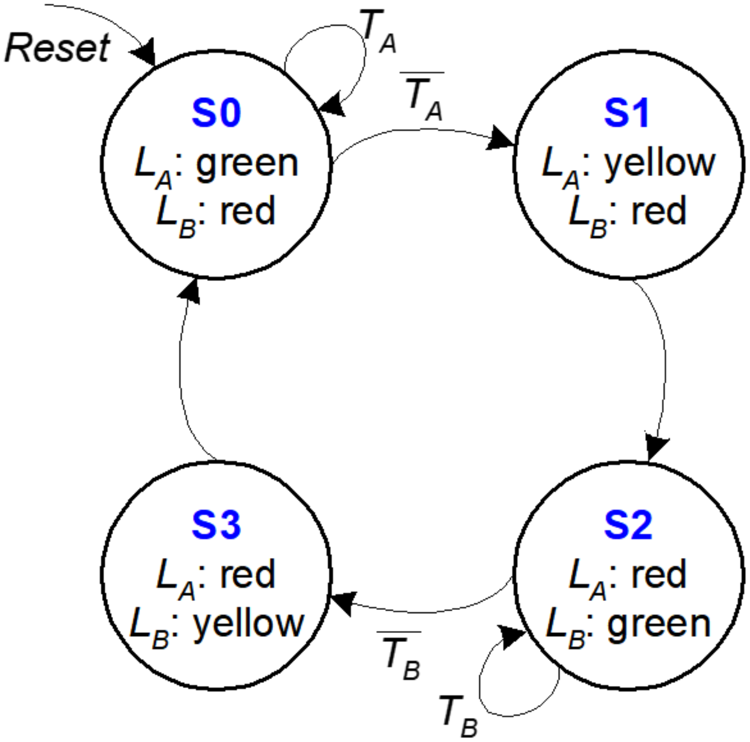


Current State		Outputs	
S_1	S_0	L_A	L_B
0	0	green	red
0	1	yellow	red
1	0	red	green
1	1	red	yellow

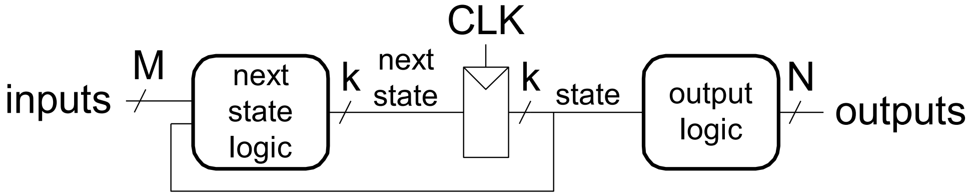


FSM – OUTPUT TABLE

- Output table



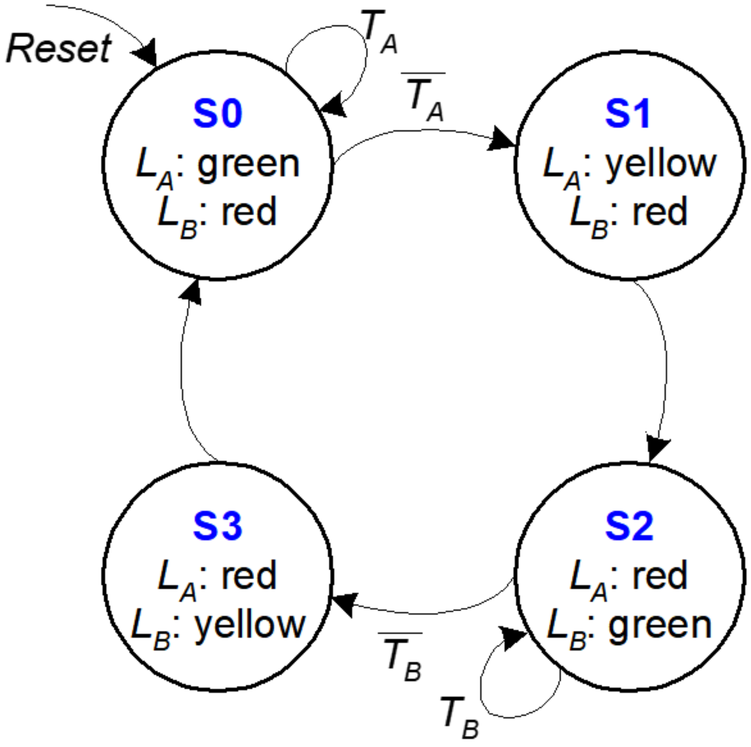
Current State		Outputs	
S_1	S_0	L_A	L_B
0	0	green	red
0	1	yellow	red
1	0	red	green
1	1	red	yellow



Output	Encoding
green	00
yellow	01
red	10

FSM – OUTPUT TABLE

- Output table



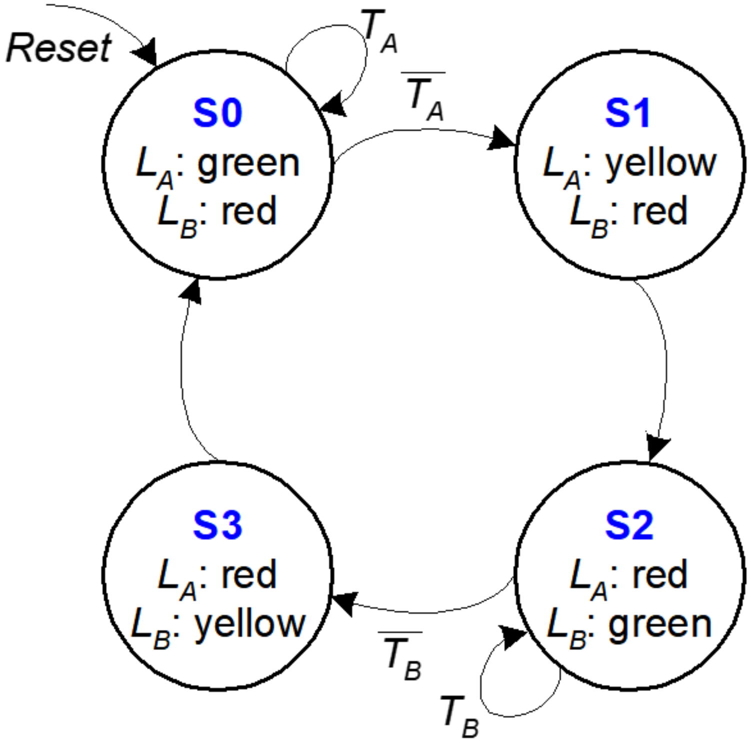
Current State		Outputs			
S_1	S_0	L_{A1}	L_{A0}	L_{B1}	L_{B0}
0	0	0	0	1	0
0	1	0	1	1	0
1	0	1	0	0	0
1	1	1	0	0	1

$$L_{A1} = S_1$$

Output	Encoding
green	00
yellow	01
red	10

FSM – OUTPUT TABLE

- Output table



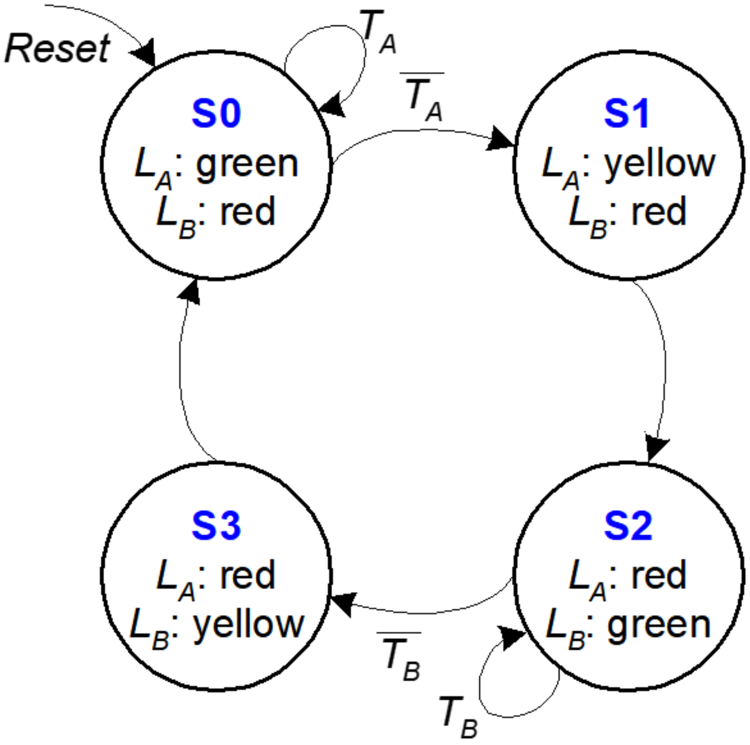
Current State		Outputs			
S_1	S_0	L_{A1}	L_{A0}	L_{B1}	L_{B0}
0	0	0	0	1	0
0	1	0	1	1	0
1	0	1	0	0	0
1	1	1	0	0	1

Output	Encoding
green	00
yellow	01
red	10

$$L_{A1} = S_1$$
$$L_{A0} = \overline{S_1} \cdot S_0$$

FSM – OUTPUT TABLE

- Output table



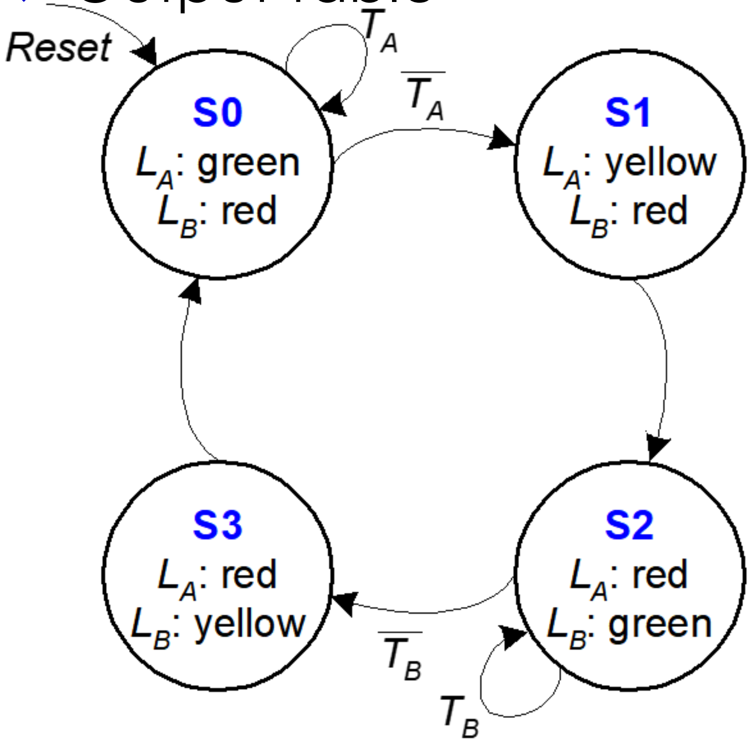
Current State		Outputs			
S_1	S_0	L_{A1}	L_{A0}	L_{B1}	L_{B0}
0	0	0	0	1	0
0	1	0	1	1	0
1	0	1	0	0	0
1	1	1	0	0	1

$$L_{A1} = S_1$$
$$L_{A0} = \overline{S_1} \cdot S_0$$
$$L_{B1} = \overline{S_1}$$

Output	Encoding
green	00
yellow	01
red	10

FSM – OUTPUT TABLE

- Output table



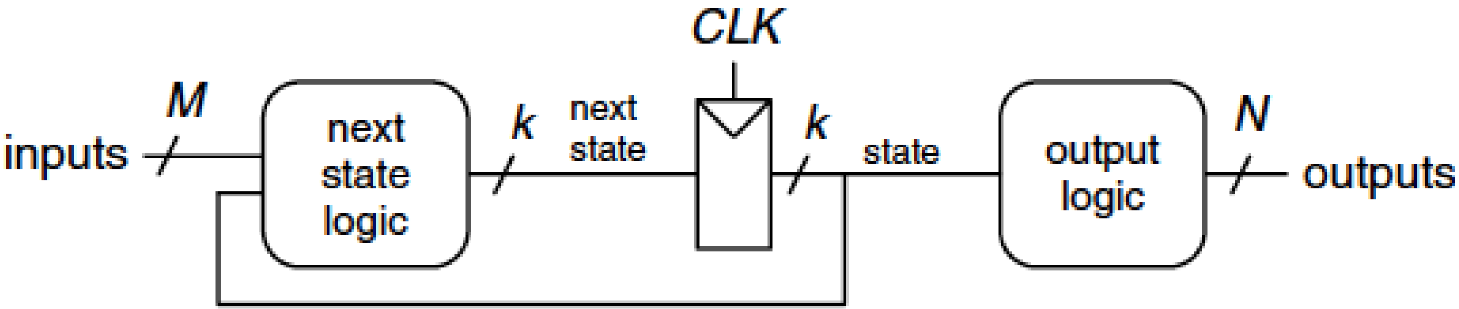
Current State		Outputs			
S_1	S_0	L_{A1}	L_{A0}	L_{B1}	L_{B0}
0	0	0	0	1	0
0	1	0	1	1	0
1	0	1	0	0	0
1	1	1	0	0	1

Output	Encoding
green	00
yellow	01
red	10

$$L_{A1} = S_1$$
$$L_{A0} = \overline{S_1} \cdot S_0$$
$$L_{B1} = \overline{S_1}$$
$$L_{B0} = S_1 \cdot S_0$$

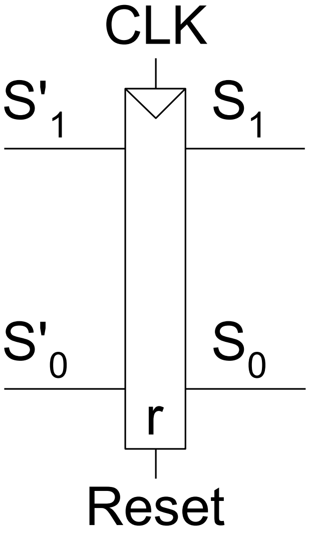
FSM – SCHEMATIC

- Overview



FSM – SCHEMATIC

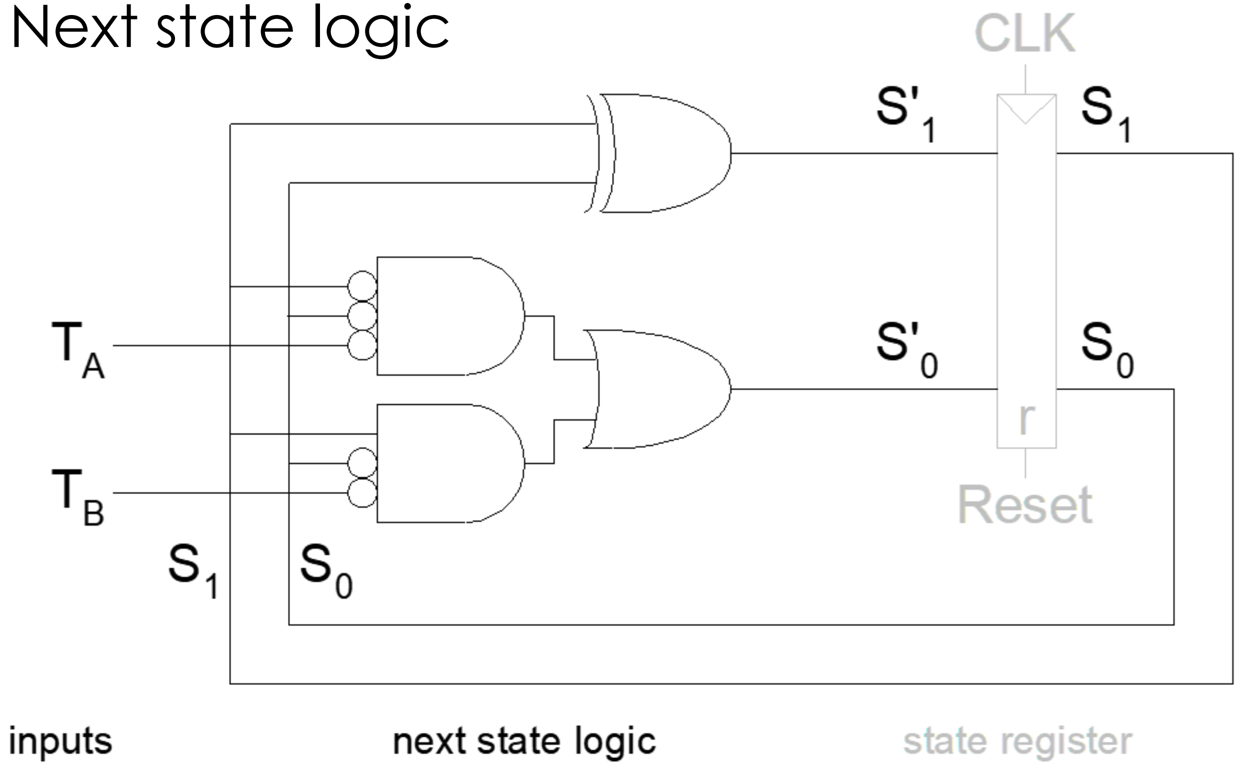
- State register



state register

FSM – SCHEMATIC

- Next state logic

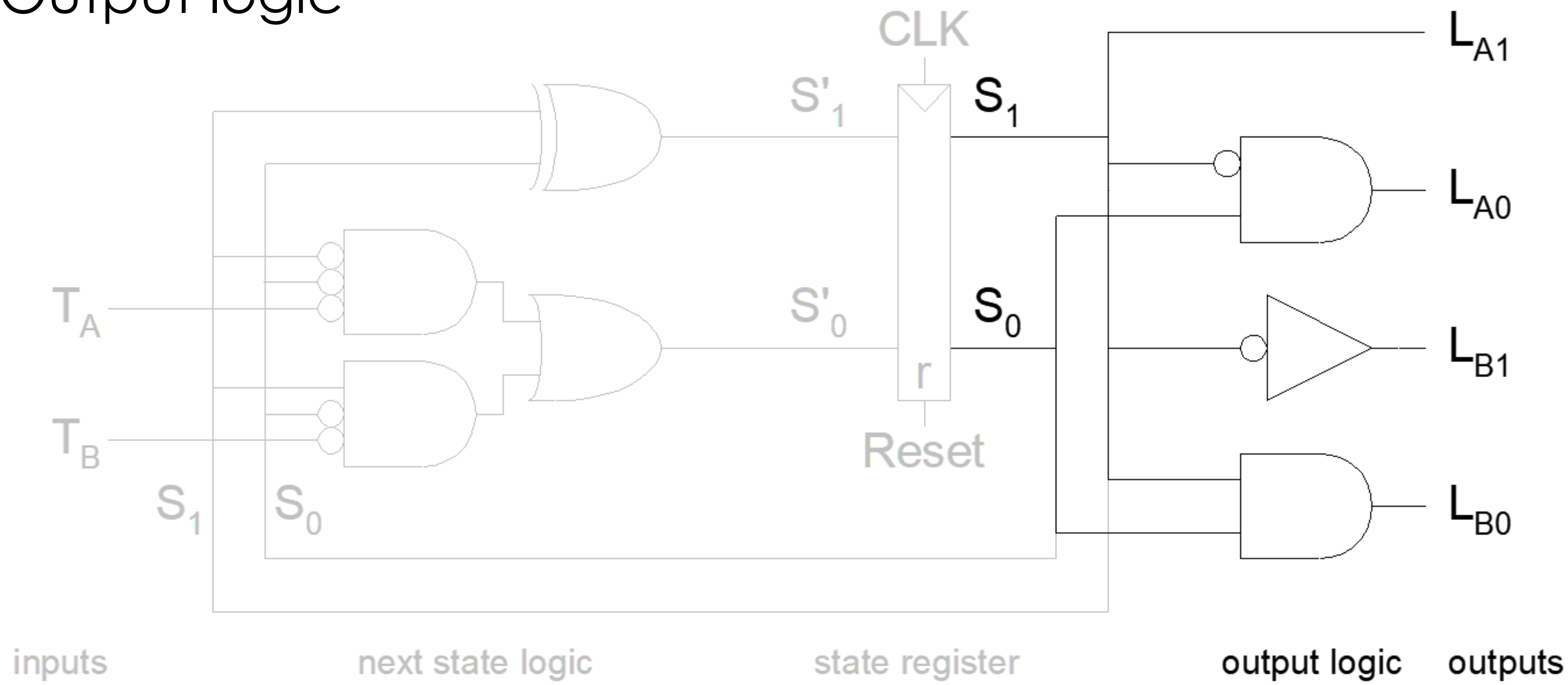


$$S'_1 = S_1 \text{ xor } S_0$$

$$S'_0 = (\overline{S_1} \cdot \overline{S_0} \cdot \overline{T_A}) + (S_1 \cdot \overline{S_0} \cdot \overline{T_B})$$

FSM – SCHEMATIC

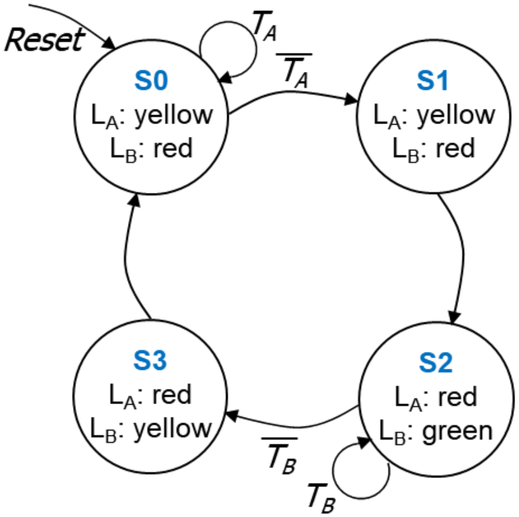
- Output logic

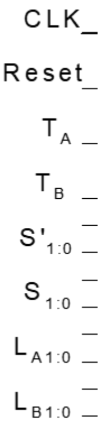


$$L_{A1} = S_1$$
$$L_{A0} = \overline{S_1} \cdot S_0$$
$$L_{B1} = \overline{S_1}$$
$$L_{B0} = S_1 \cdot S_0$$

FSM – TIMING DIAGRAM

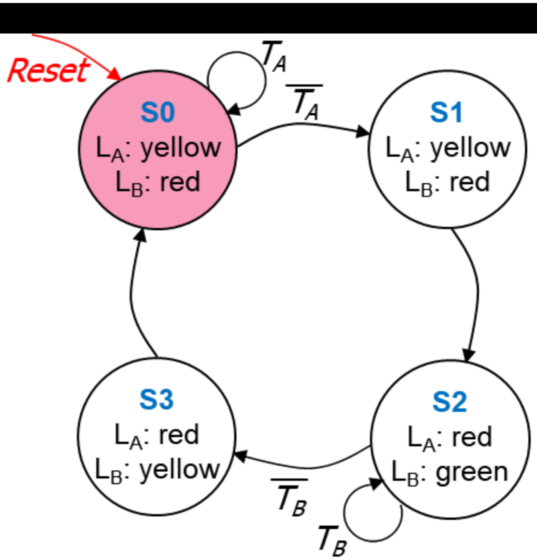
- Timing diagram

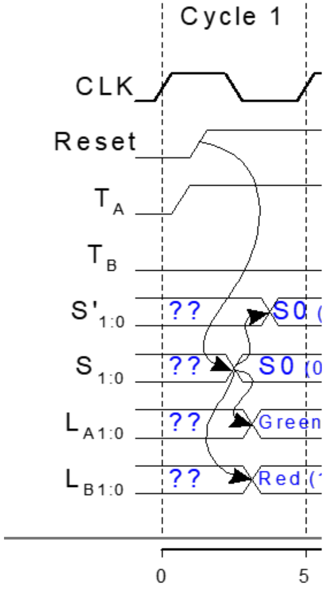




FSM – TIMING DIAGRAM

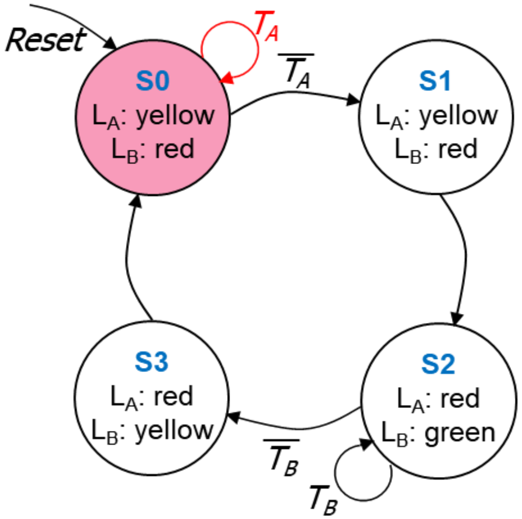
- Timing diagram

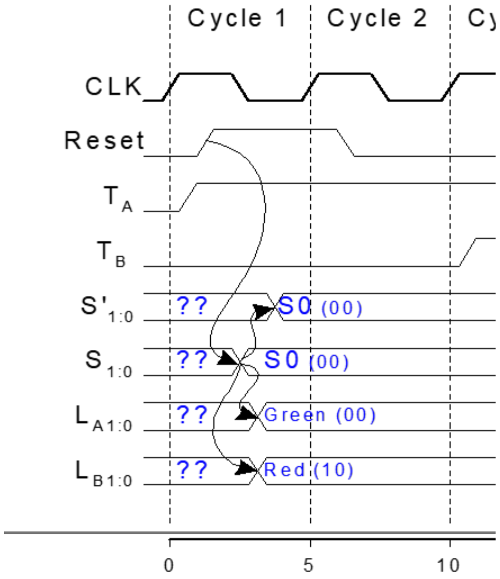




FSM – TIMING DIAGRAM

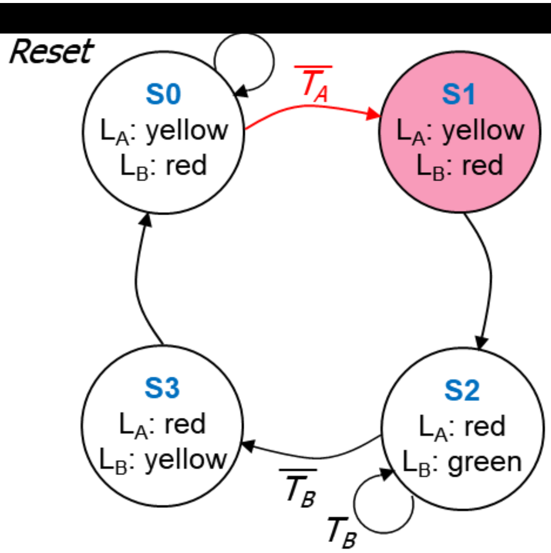
- Timing diagram

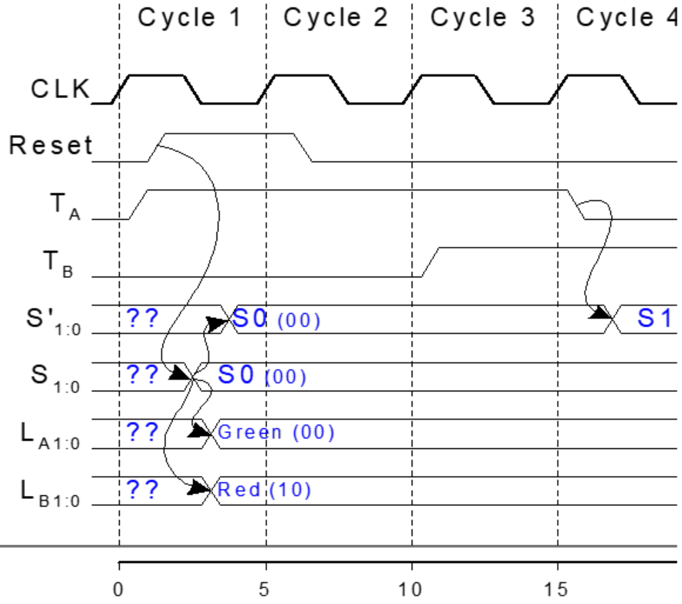




FSM – TIMING DIAGRAM

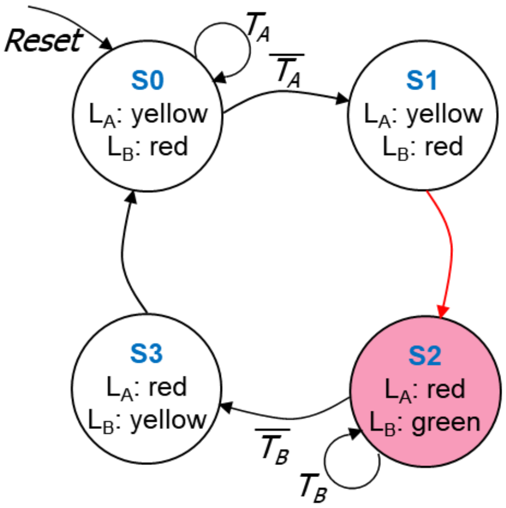
- Timing diagram

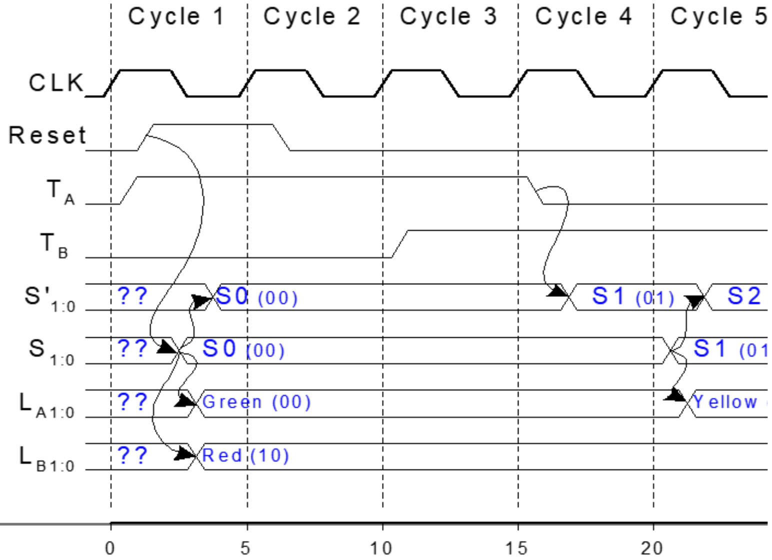




FSM – TIMING DIAGRAM

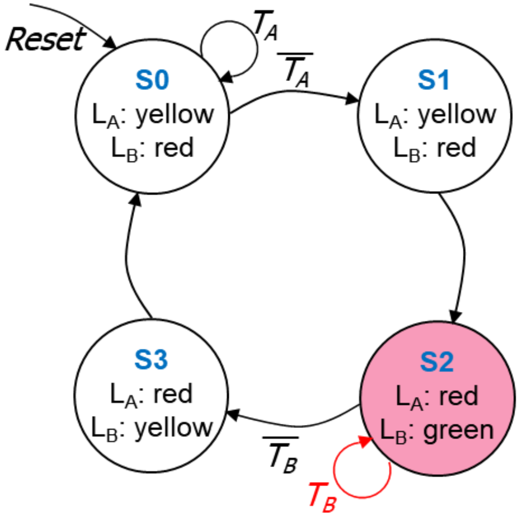
- Timing diagram

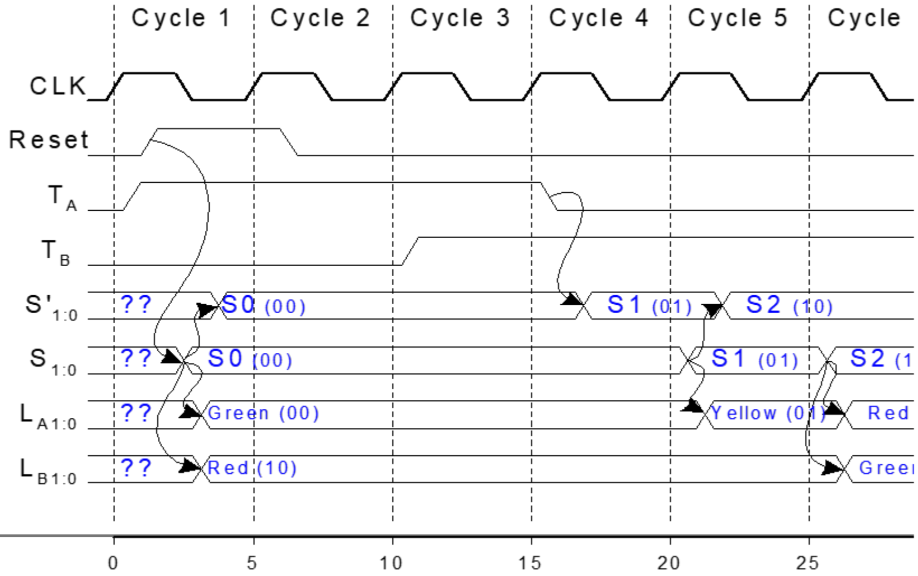




FSM – TIMING DIAGRAM

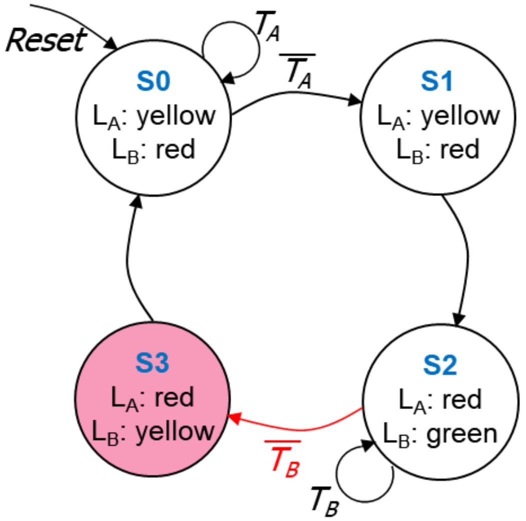
- Timing diagram

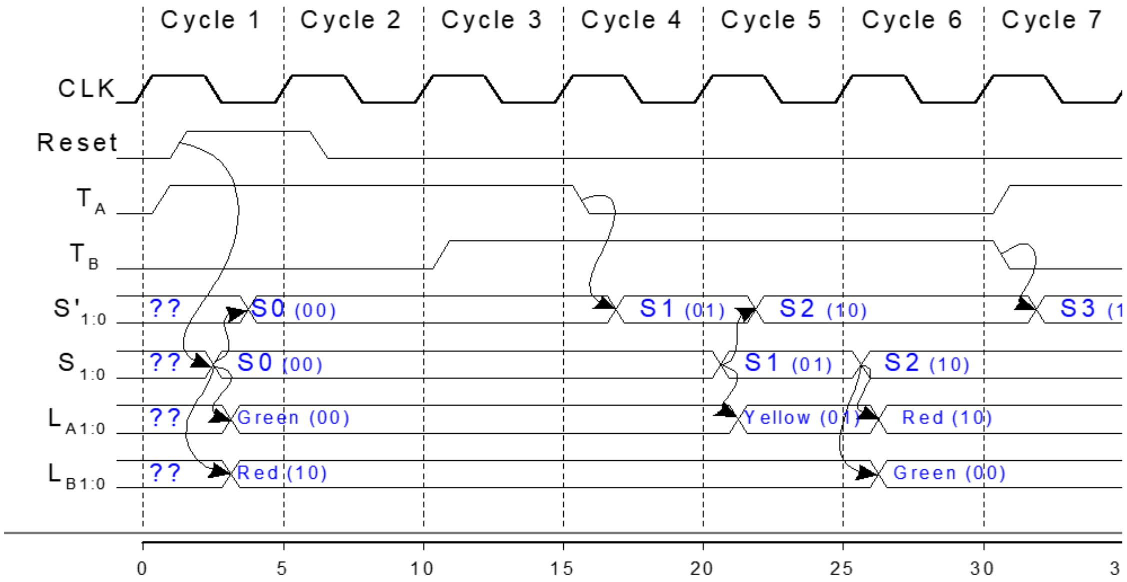




FSM – TIMING DIAGRAM

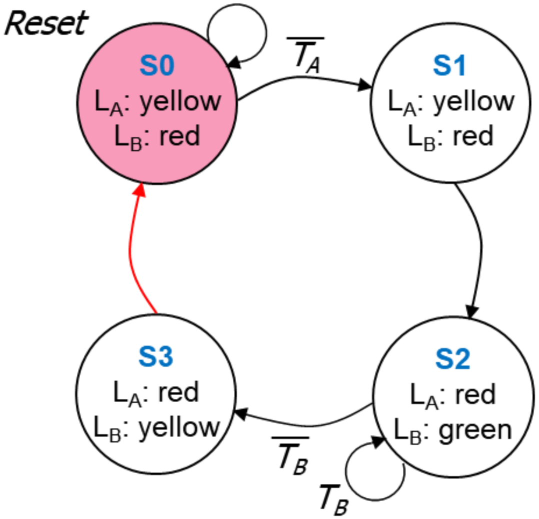
- Timing diagram

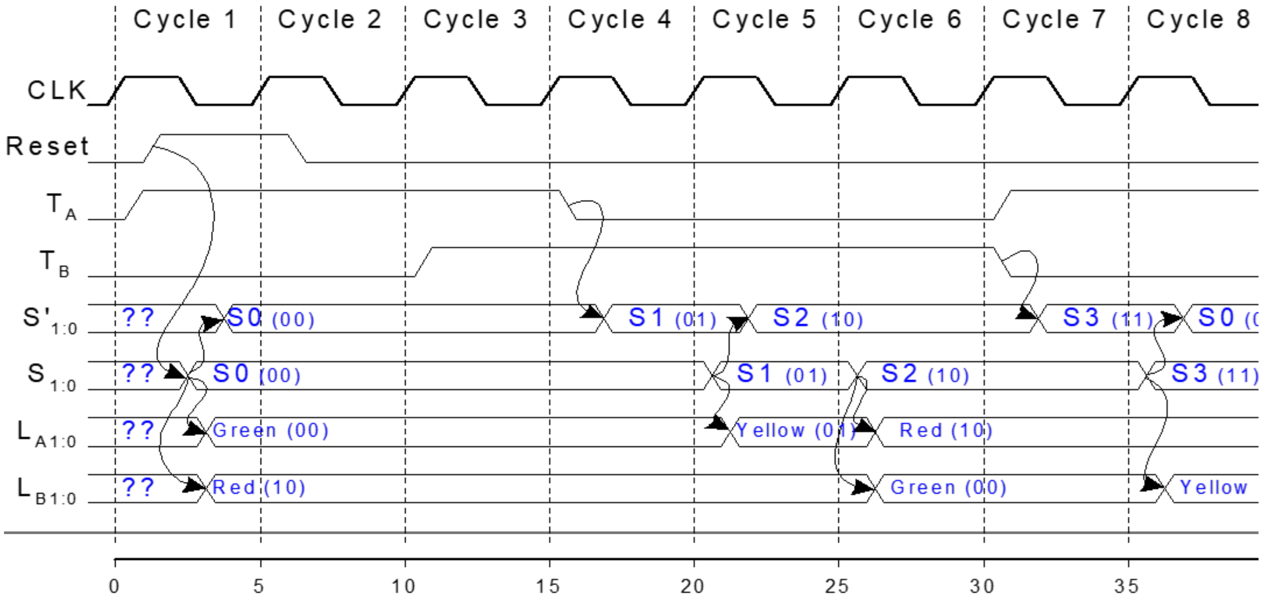




FSM – TIMING DIAGRAM

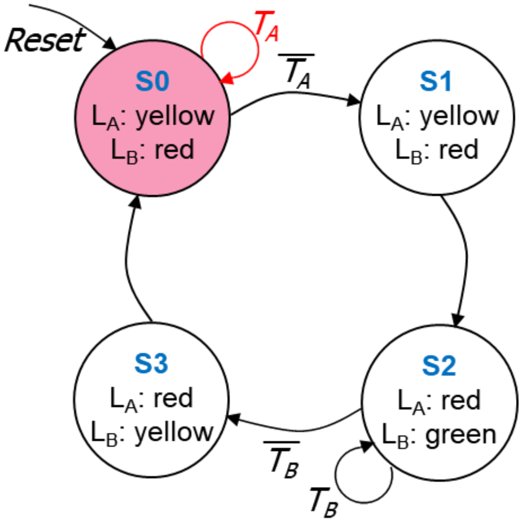
- Timing diagram

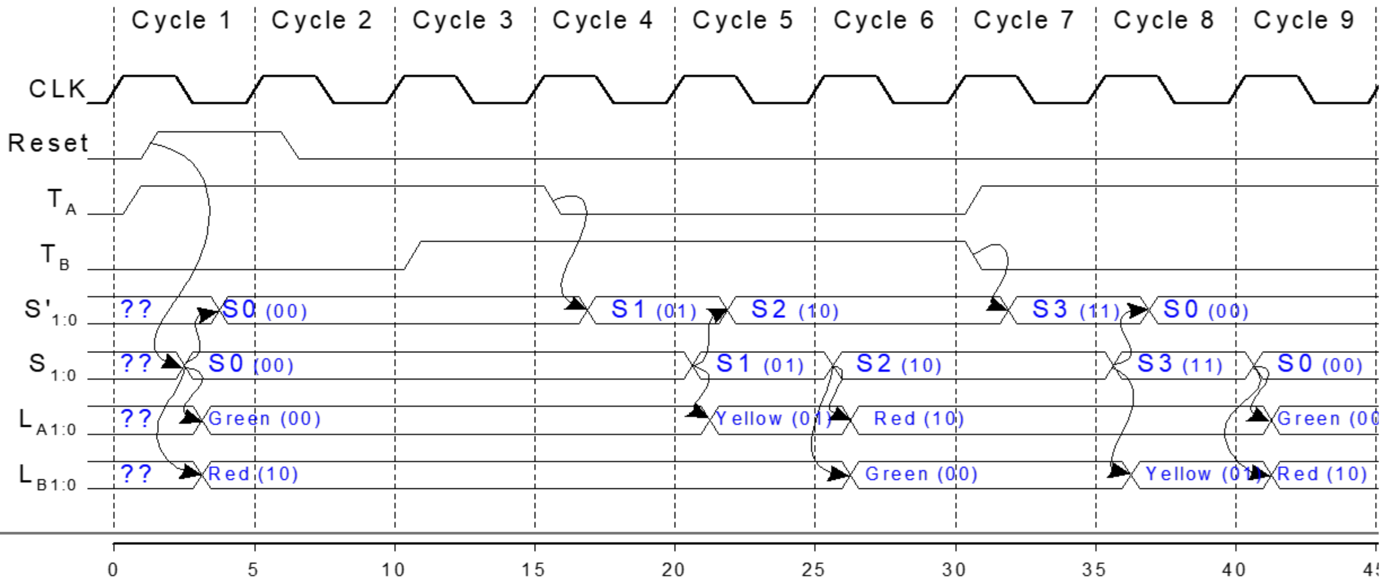




FSM – TIMING DIAGRAM

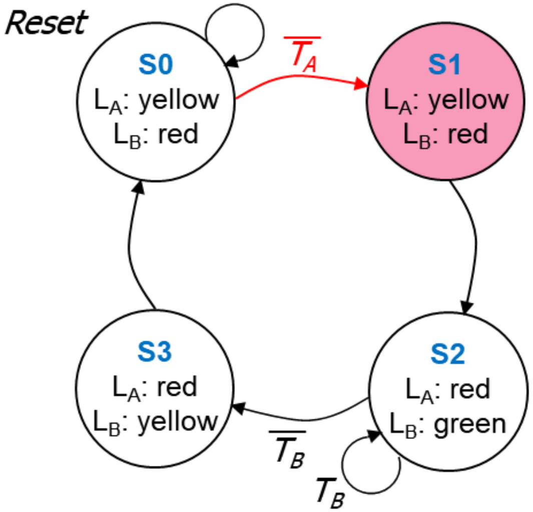
- Timing diagram

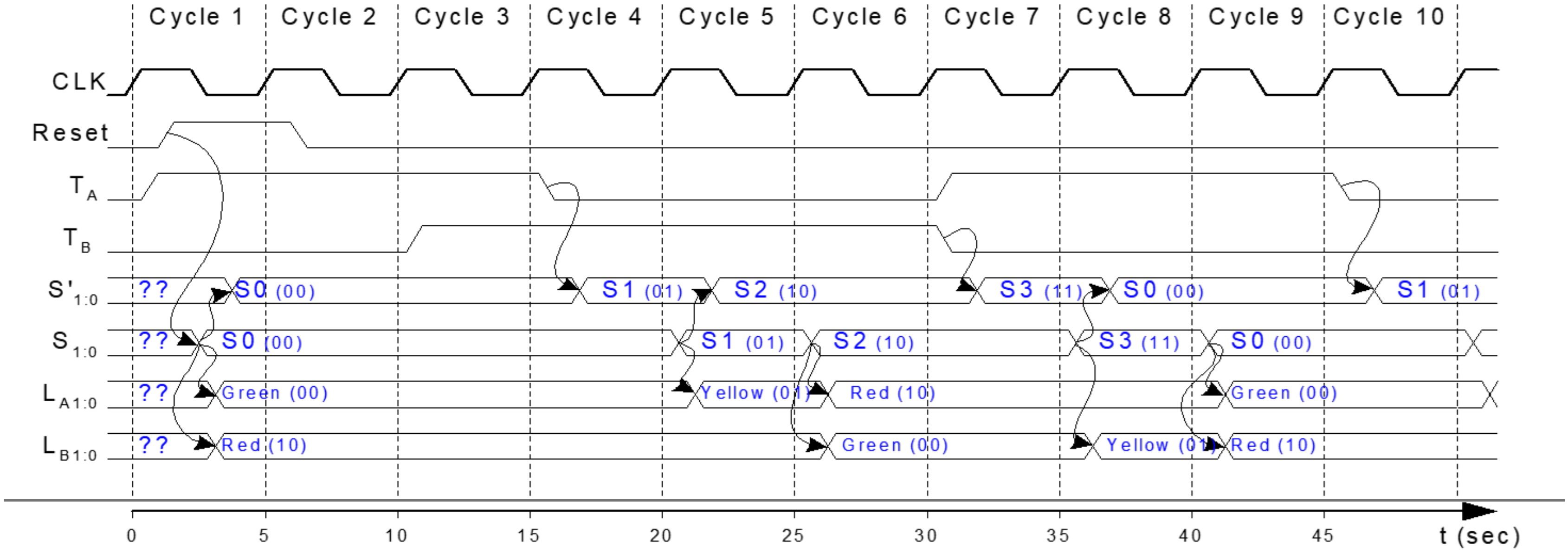




FSM – TIMING DIAGRAM

- Timing diagram

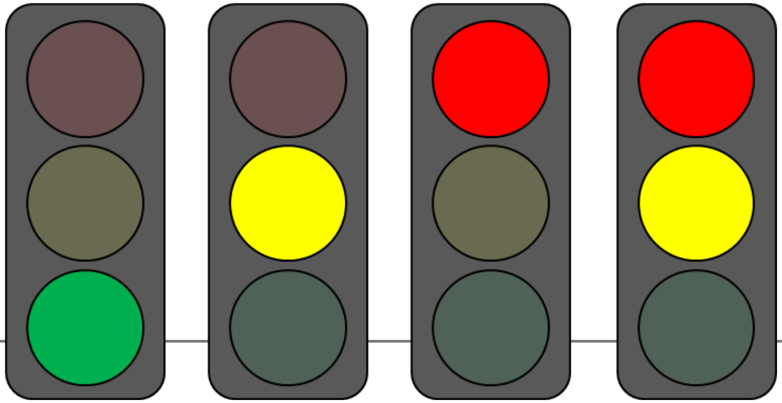




FSM – STATE ENCODING

- How do we encode the state bits?
- Three common state binary encodings with different tradeoffs
  - □ Fully Encoded
  - □ 1-Hot Encoded
  - □ Output Encoded

- Let’s see an example Swiss traffic light with 4 states
  - □ Green, Yellow, Red, Yellow+Red



FSM – STATE ENCODING

- Binary Encoding (Full Encoding):
  - Use the minimum possible number of bits
    - Use $\log_2$(num_states) bits to represent the states
  - Example state encodings: 00, 01, 10, 11
  - Minimizes # flip-flops, but not necessarily output logic or next state logic

FSM – STATE ENCODING

- One-Hot Encoding:
  - Each bit encodes a different state
    - Uses num_states bits to represent the states
    - Exactly 1 bit is “hot” for a given state
  - Example state encodings: 0001, 0010, 0100, 1000
  - Simplest design process – very automatable
  - Maximizes # flip-flops, minimizes next state logic

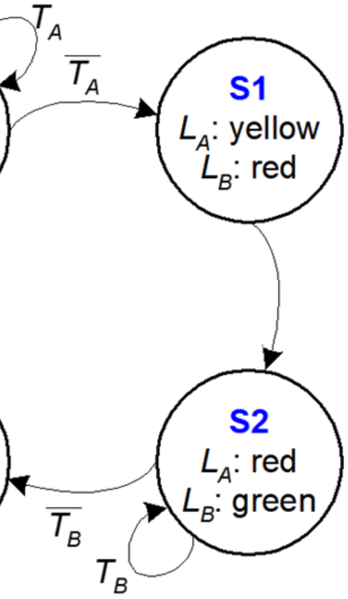
FSM – STATE ENCODING

- Output Encoding:
  - □ Outputs are directly accessible in the state encoding
  - □ For example, since we have 3 outputs (light color), encode state with 3 bits, where each bit represents a color
  - □ Example states: 001, 010, 100, 110
    - Bit0 encodes green light output,
    - Bit1 encodes yellow light output
    - Bit2 encodes red light output
  - □ Minimizes output logic
  - □ Only works for Moore Machines (output function of state)

FSM – STATE ENCODING

- The designer must carefully choose an encoding scheme to optimize the design under given constraints

- sdsd



State Transition Table

Current State	Inputs		Next State
S	T_A	T_B	S'
S0	0	X	S1
S0	1	X	S0
S1	X	X	S2
S2	X	0	S3
S2	X	1	S2
S3	X	X	S0

State	Encoding
S0	0001
S1	0010
S2	0100
S3	1000

Output T

Current State	
	LA
S0	gree
S1	yello
S2	red
S3	red

Output	Encodin
green	0001
yellow	0010
red	0100