

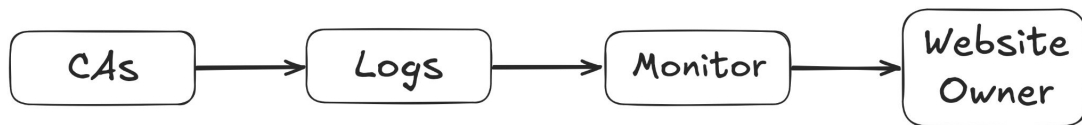
The Design Space between CT and KT

Dennis Jackson

djackson@

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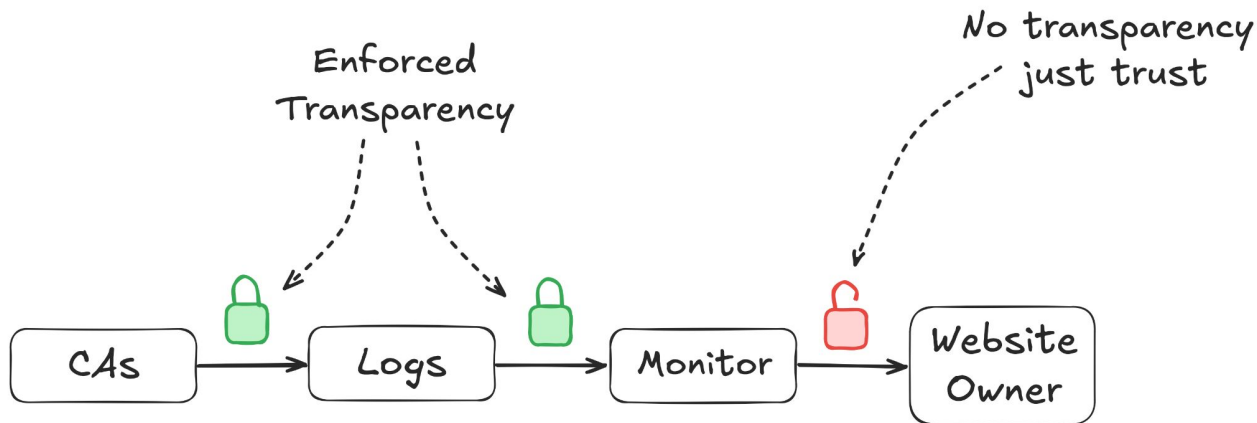
Certificate Transparency



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Certificate Transparency: Pain Points





Certificate Transparency: Pain Points

crt.sh Certificate Search

Enter an **Identity** (Domain Name, Organization Name, etc),
a **Certificate Fingerprint** (SHA-1 or SHA-256) or a **crt.sh ID**:

Search [Advanced...](#)

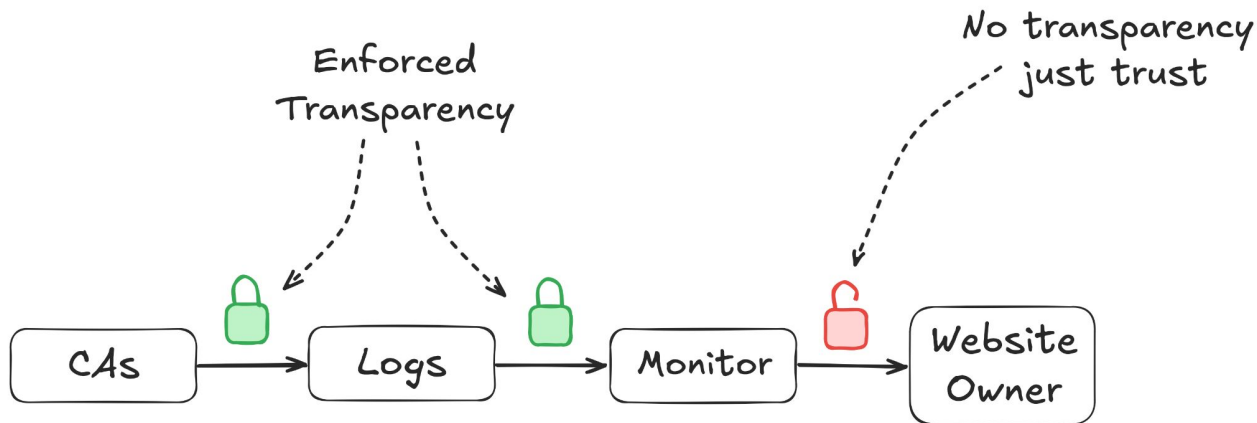
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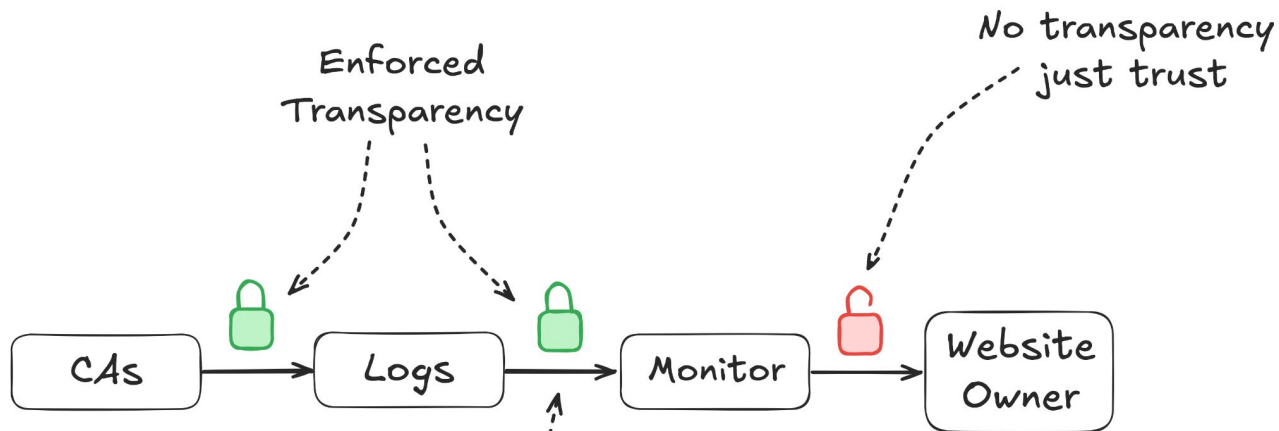
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Certificate Transparency: Pain Points



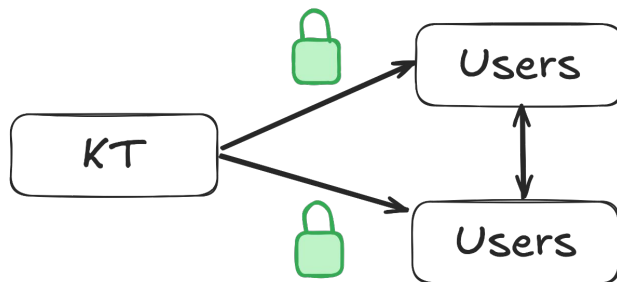
Certificate Transparency: Pain Points



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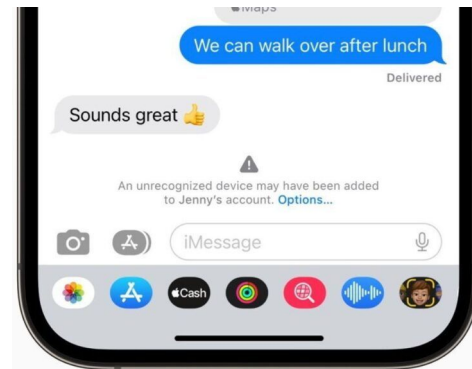
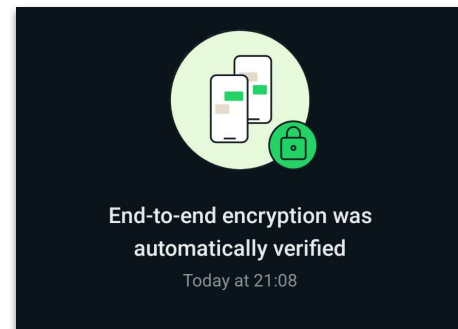
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Key Transparency



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The Space Between

Initial Question:

What ideas could be borrowed from KT and leveraged in a future CT system?

Observation:

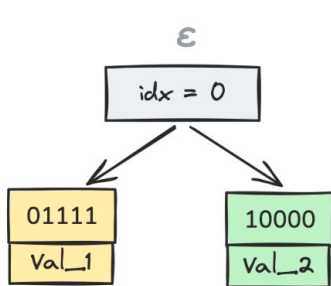
Websites, unlike users, don't need privacy.

The next 7 minutes: Sketching a design that sits between CT and KT.

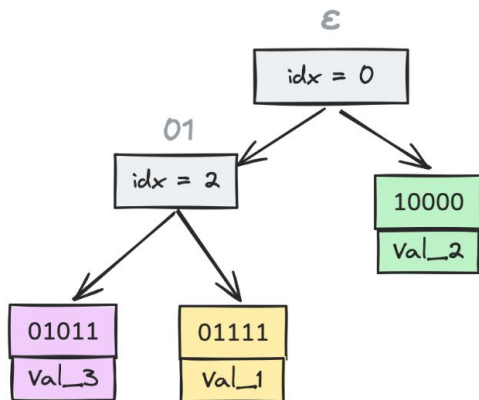
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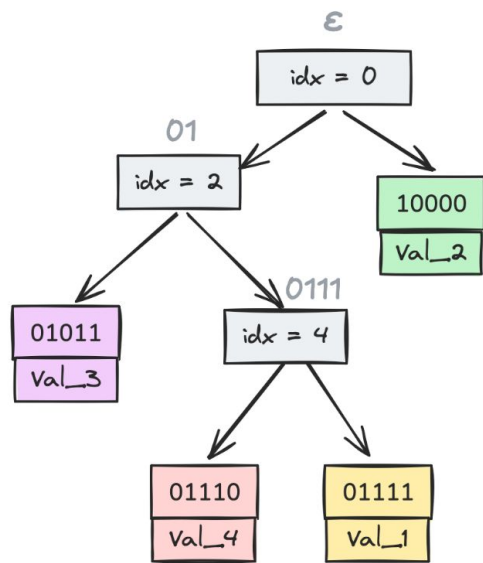
Efficient Verifiable Maps: Merkle Patricia Tries



```
{  
  01111 : Val_1,  
  10000 : Val_2  
}
```



```
{  
  01111 : Val_1,  
  10000 : Val_2,  
  01011 : Val_3  
}
```



```
{  
  01111 : Val_1,  
  10000 : Val_2,  
  01011 : Val_3,  
  01110 : Val_4  
}
```

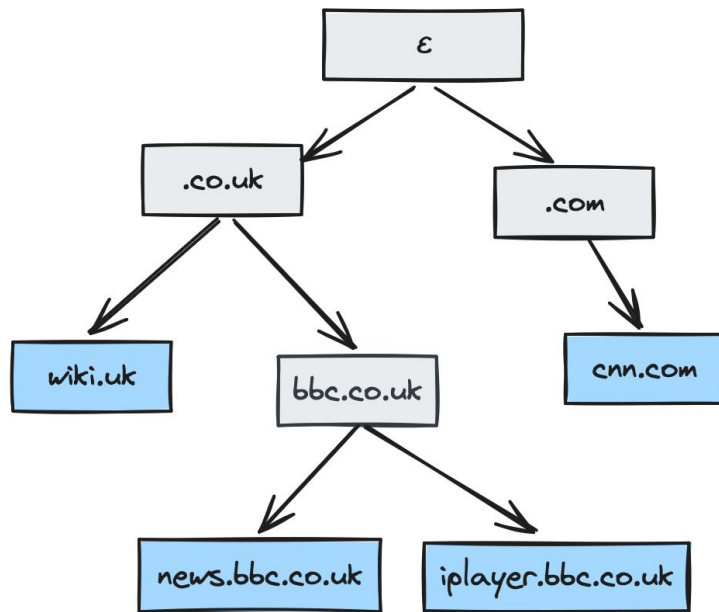
How do we key our MPT?

Key := Reverse Domain Name Notation

KT designs traditionally use a VRF to obscure user identities but we can use structured keys.

Benefits:

- Path length to root is proportional to number of eTLD+1s
- Subdomains have a shared path to the root (enables proof compression)

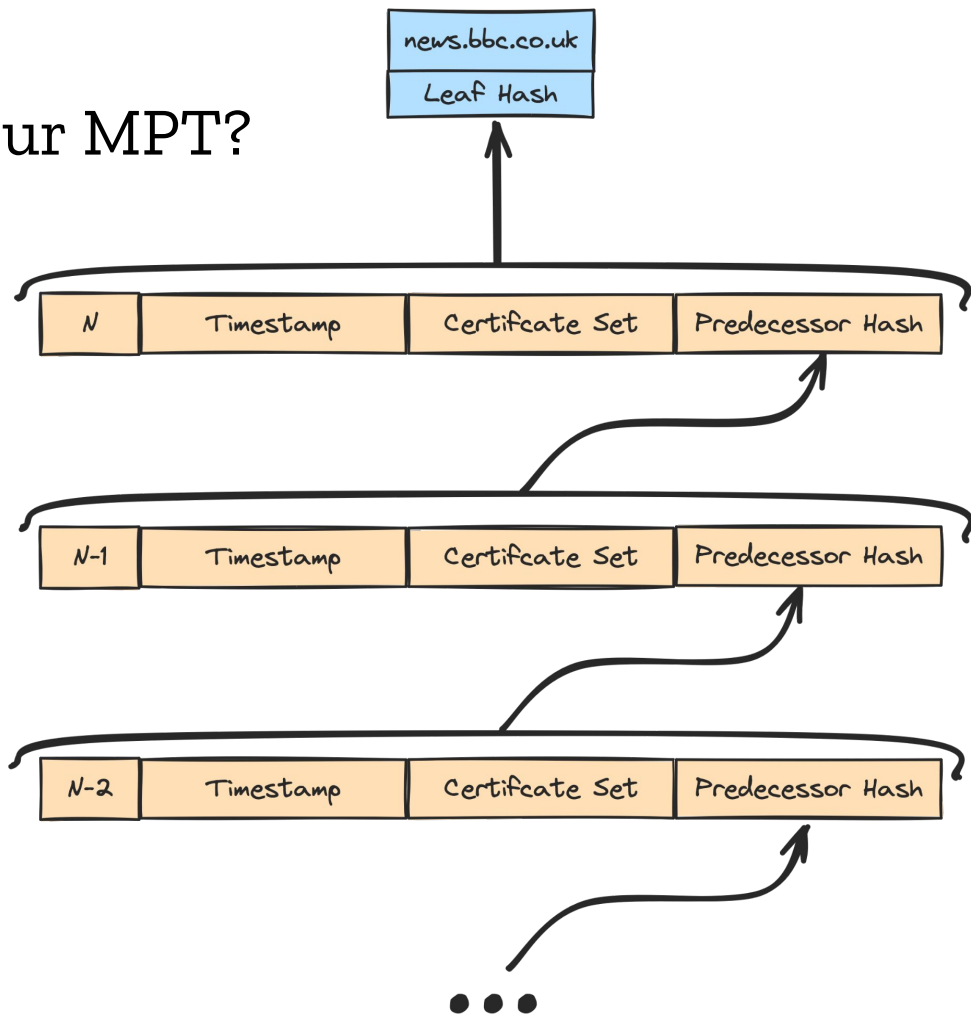


What to put in the leaves of our MPT?

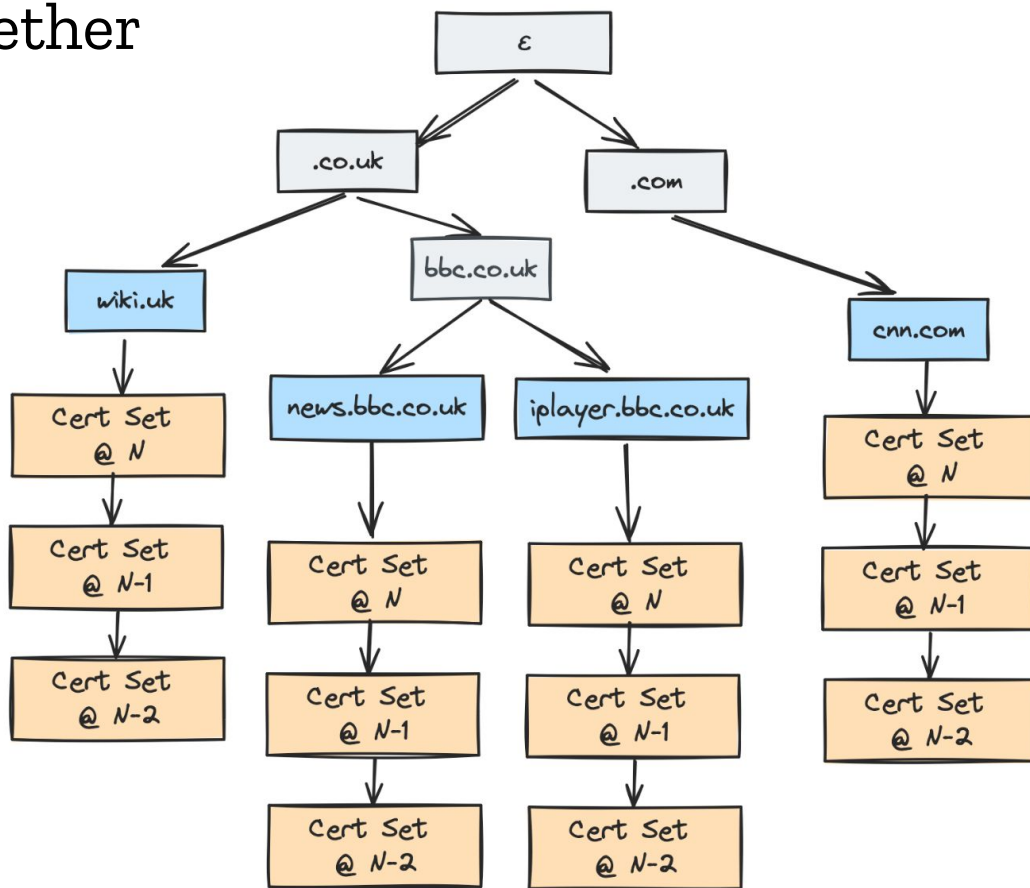
A Hash Chain of Certificate Sets

Leaf Values :=

- Set of Valid Certs
- + Timestamp
- + SeqNo
- + Hash of Predecessor



All Together



Merkle Patricia Trie

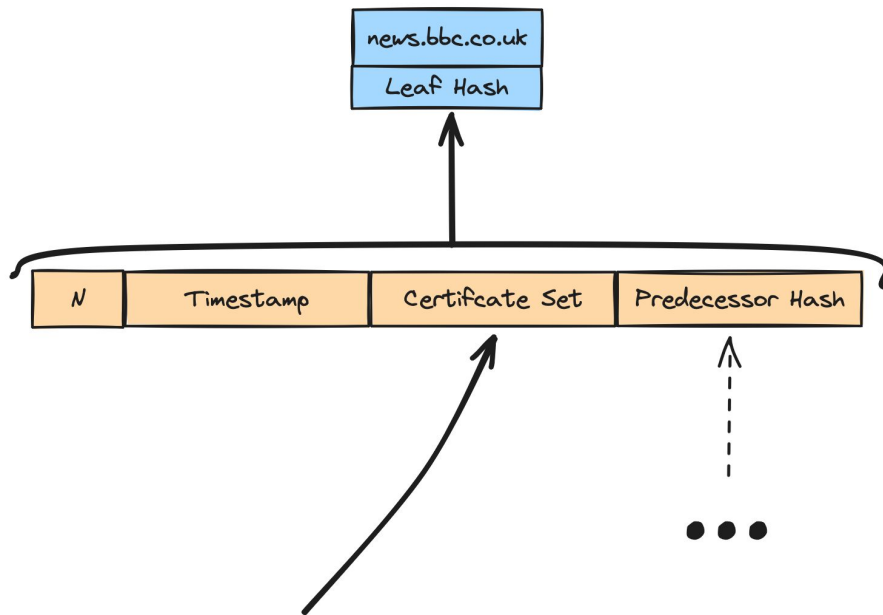
Hash Chains

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Revocation Transparency

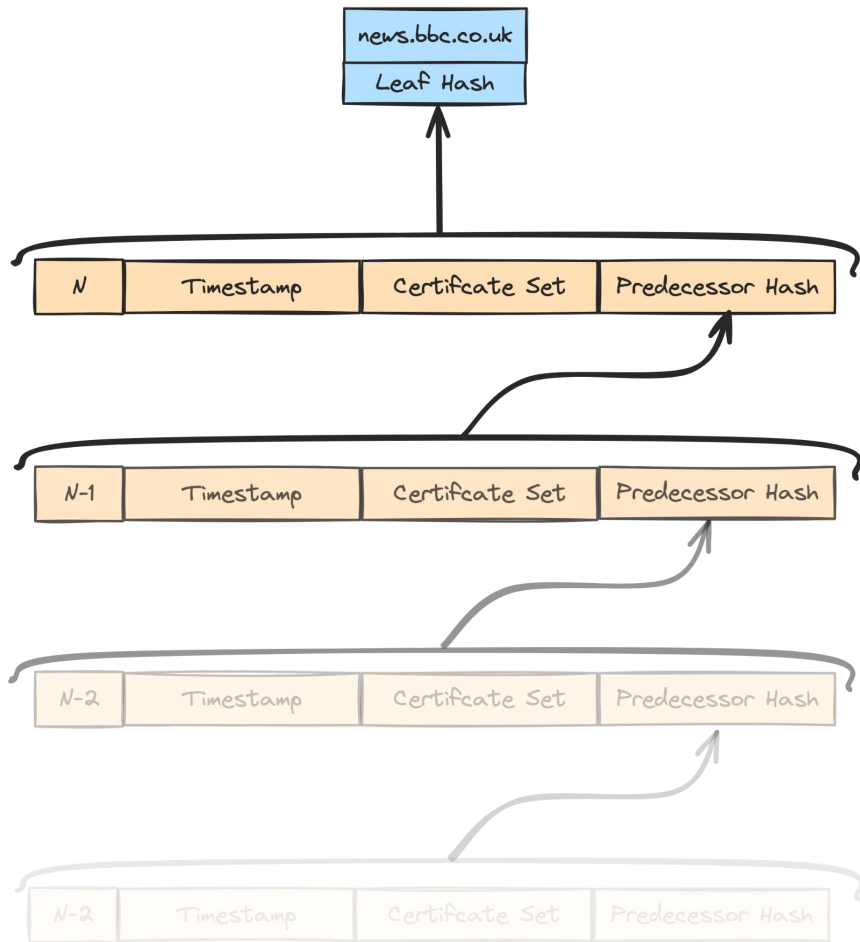
- Revocation status maintained as a structural invariant.
- When a certificate gets added or revoked, we
 - Insert a new node at the top of the hash chain
 - Update its path to the root through the MPT.



Unrevoked by nature
of its position as direct
child of leaf node.

Graceful Expiry

- We can also gracefully forget old entries without having to update the tree.
- Log storage is proportional to number of active domains rather than total number of issued certificates.



Proof Lengths

- $\mathbf{D} :=$ | All eTLD+1s |
- $\mathbf{S} :=$ | Site's Subdomains |
- $\mathbf{C} :=$ | Site's Valid Certs |
- $\mathbf{E} :=$ | Site's Expired / Revoked Certs |

Proof a certificate is not revoked (or absent):

$\text{Log } \mathbf{D} + \text{log } \mathbf{S} + \text{log } \mathbf{C}$

Proof of Certificate History:

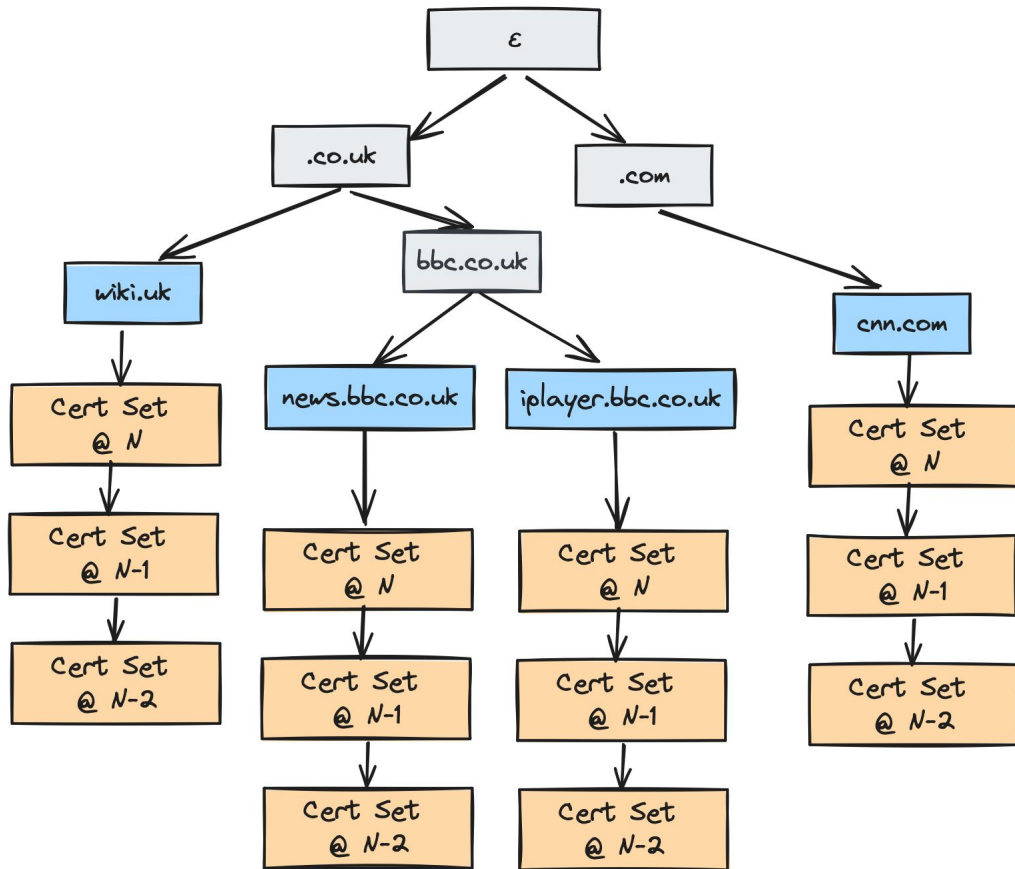
$\text{Log } \mathbf{D} + \text{log } \mathbf{S} + \mathbf{C} + \mathbf{E}$

Privacy Preserving KT solutions:

- Tree size scales a multiple of all certificates.
- Proof lengths are $3 \times \text{Log} (|\text{Tree size}|)$

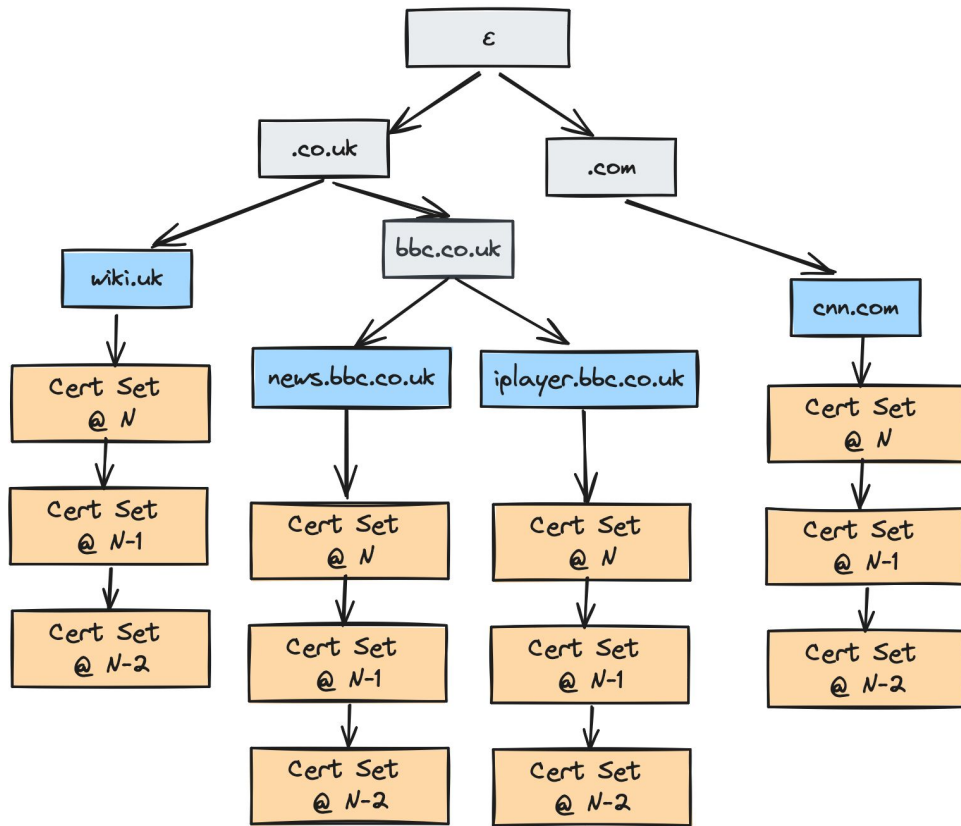
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Observations

- Cheap E2E proof that a given certificate is present and unrevoked, e.g. in an MTC-style design.
- Enables succinct proofs of all certificates issued for a given domain and its subdomains.
- Lower egress costs - can use a quorum of auditors rather than general public monitors.
- Storage costs grow proportional to valid certificates, not total issued certificates.



Thoughts

- KT-like designs might help solve challenges in the CT ecosystem.
- Loosening the privacy constraint of KT unlocks a rich design space with opportunities for much greater efficiencies.
- **Ideal Monitoring Story?** `certbot --certificate-report *.example.com`
- Did this sketch pique your interest? Do you know of other work in this area?
Come say hi!

Credits: Kevin Lewi in particular, but also many conversations at HACS & RWC 2024, including Bas Westerbaan, Sophie Schmieg, Esha Ghosh, Alexander Scheel, Kevin Milner, Richard Barnes, and Brendan McMillion.

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