

3.

a. $Opt(i)$:
 if $i = 0$ then
 $result \leftarrow 0$
 else
 $result \leftarrow \infty$
 for k in $[0 .. i)$ do
 $penaltyForLastDay \leftarrow (200 - (a_i - a_k))^2$
 $penaltyForPreviousDays \leftarrow Opt(k)$
 $result \leftarrow \min(result, penaltyForLastDay + penaltyForPreviousDays)$
 return $result$

b. $Opt(i)$:
 if $T[i] = \perp$ then
 if $i = 0$ then
 $T[i] \leftarrow 0$
 else
 $T[i] \leftarrow \infty$
 for k in $[0 .. i)$ do
 $penaltyForLastDay \leftarrow (200 - (a_i - a_k))^2$
 $penaltyForPreviousDays \leftarrow Opt(k)$
 $T[i] \leftarrow \min(result, penaltyForLastDay + penaltyForPreviousDays)$
 return $T[i]$

This algorithm will execute operations up to n for each n , making run time without memoization $n(n+1)/2$ or just $O(n^2)$.

c. $Opt(n)$:
 for i in $[0 .. n)$
 $T[i] \leftarrow 0$ // default value
 for i in $[0 .. n)$
 $result \leftarrow \infty$
 for k in $[0..i)$ do
 $penaltyForLastDay \leftarrow (200 - (a_i - a_k))^2$
 $penaltyForPreviousDays \leftarrow T[i]$
 if $penaltyForLastDay + penaltyForPreviousDays < result$
 $result = penaltyForLastDay + penaltyForPreviousDays$
 $T[i] = result$
 return $T[i]$

For the same reason as in (b), we can see that this is also $O(n^2)$

d. $Sol(i)$:
 if $i = 0$ then
 return `emptyList()`

else

for k in $[0 .. i)$ do

if $T[i - k] + (200 - (a_i - a_k))^2 = T[i]$ then

return *concat*(*Sol*($i - k$), k)

This algorithm can be implemented in linear time ($O(n)$), if the concat operation is also linear. This can be achieved with a linked list data structure storing out answer.