

9.

- a. $i=0 \rightarrow \text{vec.append}(1)$ will call $\text{resize}(1)$ which will set `newsz` to 1. Line 13 was executed 0 times.
 $i=1 \rightarrow \text{resize}(2)$ will be called which will have line 13 be executed once as size is not yet updated.
 $i=2 \rightarrow \text{resize}(3)$ will be called which will have line 13 be executed twice.
 $i=3 \rightarrow \text{resize}(4)$ will be called which will have line 13 be executed three times.
It is clear to see that there is a pattern. vec.append is called n times and for every i th iteration, line 13 is called i times. This means that the formula for the number of times line 13 is called is essentially $n * n(n-1)/2$ (number of loops in the outer for-loop * number of loops in the inner for-loop). This means that the number of times line 13 is executed is $\Theta(n^2)$.
- b. $i=0 \rightarrow \text{vec.append}(1)$ will call $\text{resize}(1)$ which will set `newsz` to 1. Line 13 was executed 0 times.
 $i=1 \rightarrow \text{resize}(2)$ will be called which will have line 13 be executed once.
 $i=2 \rightarrow \text{resize}(3)$ will be called which will have line 13 be executed twice.
 $i=3 \rightarrow \text{resize}(4)$ will be called which will have line 13 not be executed.
 $i=4 \rightarrow \text{resize}(5)$ will be called which will have line 13 be executed four times.
When i is a power of 2, only then will the nested loop execute.
 $2^{\lfloor \log_2(n) \rfloor} \rightarrow$ Which is essentially $O(n)$

Outer For Loop $i=0$	Inner For Loop line 13	You can rewrite this to make it so	Outer For Loop $i=0$	Redistributed series
1	1	that each 'i'	1	1
2	2	will have	2	1
3	0	only 1 other	3	1
4	4	execution.	4	1
5	0	This makes	5	1
6	0	$O(n)$	6	1
7	0		7	1
8	8		8	1