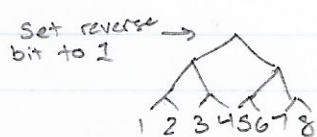


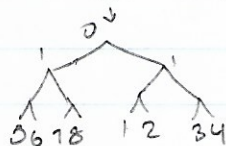
- 6) In order to say that we will have this tree represented in reverse order, we put a 1 bit onto the root level. This is simply one operation, making the complexity $O(1)$. The reverse bits in a node would then be pushed down to each of its children and then swapping the children. Then, by clearing the parent, you could consider finding the k th value the same way as in 2.5.
- For example, given a tree representing leaves with values 1-8



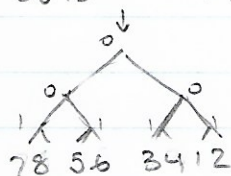
Assuming $k=5$

↳ This is finding the k th number in the reversed tree

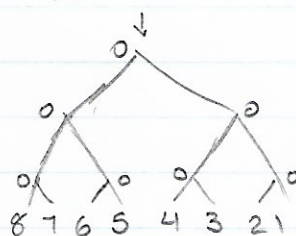
Zero bit and Push down to children and swap



$5 > 4 \rightarrow$ Search path for k is to the right
counter = $5 - 4 \rightarrow 1$



$1 < 2 \rightarrow$ Path is to the left



Path of k is to the left again

Making the element at $k=4$ of the reverse

algorithm _{to find}
Because the path to k is preserved, we can derive split and concat in the same runtime