Week 03 Tutorial Solutions

1. On a machine with 16-bit ints, the C expression (30000 + 30000) yields a negative result.

Why the negative result? How can you make it produce the correct result?

Tutorial Recordings

These are *not* a substitute for tutorial attendance. You are expected to attend a tutorial class.

tue13a, with Emily Chen:
BBCollaborate; Mirror@CSE

Answer:

Since the two numbers in the expression are treated as int values, the overall expression is computed as an int value. In C, an int is a signed value, which means that one of the sixteen bits is effectively used as a sign (+ or -) bit, leaving only fifteen bits to represent

the magnitude. The negative result occurs because 60000 cannot be stored in fifteen bits. The maximum value is $7 {\tt FFF}_{16} = 2^{15} - 1 = 32767$.

One approach to resolving the problem is to write the expression as (30000u + 30000u), which will cause C to evaluate the expression as an unsigned int value; unsigned ints can use the whole sixteen bits for magnitude, allowing them to store values up to $FFFF_{16} = 2^{16} - 1 = 65535$.

- 2. Assume that the following hexadecimal values are 16-bit twos-complement. Convert each to the corresponding decimal value.
 - i. 0x0013
 - ii. 0x0444
 - iii. 0x1234
 - iv. 0xffff
 - v. 0x8000

Answer:

For the positive values (i.e., leftmost bit 0), convert to bits and then, for any ith bit which is 1, sum the 2^i values. For the negative values, convert them (using two-complement) to their corresponding value, and then apply the strategy for positive values.

```
i. 0 \times 0013 = 0000\,0000\,0001\,0011_2 = 2^4 + 2^1 + 2^0 = 19 ii. 0 \times 4444 = 0000\,0100\,0100\,0100_2 = 2^{10} + 2^6 + 2^2 = 1092 iii. 0 \times 1234 = 0001\,0010\,0011\,0100_2 = 2^{12} + 2^9 + 2^5 + 2^4 + 2^2 = 4660 iv. 0 \times \text{ffff} \Rightarrow -1 \times (0 \times 0000 + 1) = -1 v. 0 \times 8000 \Rightarrow -1 \times (0 \times 7 \text{fff} + 1) \Rightarrow -((2^{15} - 1) + 1) = -32768
```

- 3. Give a representation for each of the following decimal values in 16-bit twos-complement bit-strings. Show the value in binary, octal and hexadecimal.
 - i. 1
 - ii. 100
 - iii. 1000
 - iv. 10000
 - v. 100000
 - vi. -5
 - vii. -100

Answer:

100000 cannot be represented in 16 bits

- e. 0 01111110 11111111111111111111111
- f. 0 10000000 011000000000000000000000
- h. 0 01101110 10100000101000001010000

Each of the above is a single 32-bit bit-string, but partitioned to show the sign, exponent and fraction parts.

Answer:

All values are computed by the formula

$$\mathsf{sign} \times (1 + \mathsf{frac}) \times 2^{\mathsf{exp} - 127}$$

where

- o sign is 1 if the most significant bit (m.s.b) is 0, or -1 if the m.s.b is 1
- o exp is determined by the 8 bits following the sign bit (as a value in the range 0..255)
- \circ frac is determined by the least significant 23 bits (bit₂₂ \times 2⁻¹ + bit₂₁ \times 2⁻² + \cdots + bit₁ \times 2⁻²² + bit₀ \times 2⁻²³)
- a. $0.0000 = (1+0.0) imes 2^{0-127} = 1.0 imes 2^{-127}$ (... which is close to zero)
- b. -0.0000 ... same as above, but the sign bit is 1, so -ve
- c. $1.5 = (1+0.5) imes 2^{127-127} = 1.5 imes 2^0$
- d. $0.5 = (1+0.0) imes 2^{126-127} = 1.0 imes 2^{-1}$
- e. $0.999999 = (1 + 0.999999) \times 2^{126-127} = 1.999999 \times 2^{-1}$

(... but we're limiting this to the approximately 7 digits we can represent here)

f.
$$2.75 = (1+0.375) imes 2^{128-127} = 1.375 imes 2^1$$

g.
$$3145728.00 = (1+0.5) imes 2^{148-127} = 1.5 imes 2^{21}$$

h.
$$0.0000124165 = (1+0.627451) imes 2^{110-127} = 1.627451 imes 2^{-17}$$

(... again, limited to approximately 7 digits)

- 5. Convert the following decimal numbers into IEEE 754-encoded bit-strings:
 - a. 2.5
 - b. 0.375
 - c. 27.0
 - d. 100.0

Answer:

We need to first express the number k as $(1 + \mathsf{frac}) \times 2^n$. To work out the fraction, we divide k by the largest 2^n that is smaller than k.

$$=1.25\times 2^1=(1+0.25)\times 2^{128-127}$$

$$=1.5 imes 2^{-2}=(1+0.5) imes 2^{125-127}$$

$$=1.6875 imes 2^4=(1+0.6875) imes 2^{131-127}$$
 where $0.6875=2^{-1}+2^{-3}+2^{-4}$

d. 0 10000101 1001000000000000000000000

$$=1.5625 imes 2^6=(1+0.5625) imes 2^{133-127}$$
 where $0.5625=2^{-1}+2^{-4}$

6. Write a C function, six_middle_bits, which, given a uint32_t, extracts and returns the middle six bits.

Answer:

```
uint32_t six_middle_bits(uint32_t u) {
   return (u >> 13) & 0x3F;
}
```

7. Draw diagrams to show the difference between the following two data structures:

```
struct {
    int a;
    float b;
} x1;
union {
    int a;
    float b;
} x2;
```

If x1 was located at &x1 == 0x1000 and x2 was located at &x2 == 0x2000, what would be the values of &x1.a, &x1.b, &x2.a, and &x2.b?

The struct contains two separate components. The union contains two components that occupy the same memory space. Struct ... x1 a int b float union ... x2 a,b int or float

8. How large (#bytes) is each of the following C unions?

```
a. union { int a; int b; } u1;
b. union { unsigned short a; char b; } u2;
c. union { int a; char b[12]; } u3;
d. union { int a; char b[14]; } u4;
e. union { unsigned int a; int b; struct { int x; int y; } c; } u5;
```

You may assume sizeof(char) == 1, sizeof(short) == 2, sizeof(int) == 4.

&x1.a==0x1000, &x1.b==0x1004, &x2.a==0x2000, &x2.b==0x2000

Answer:

The size of a union is the size of its largest variant:

```
a. sizeof(u1) == 4
b. sizeof(u2) == 2
c. sizeof(u3) == 12
d. sizeof(u4) == 16, with padding on string
e. sizeof(u5) == 8
```

Note that the above results may vary depending on the machine architecture and the compiler.

9. Consider the following C union

```
union _all {
   int ival;
   char cval;
   char sval[4];
   float fval;
   unsigned int uval;
};
```

If we define a variable union _all var; and assign the following value var.uval = 0×00313233 ;, then what will each of the following printf(3)s produce:

```
a. printf("%x\n", var.uval);b. printf("%d\n", var.ival);c. printf("%c\n", var.cval);d. printf("%s\n", var.sval);
```

```
e. printf("%f\n", var.fval);f. printf("%e\n", var.fval);
```

You can assume that bytes are arranged from right-to-left in increasing address order.

Answer:

This is just a matter of interpreting the bytes/words in terms of the appropriate data type:

```
a. printf("%x\n", var.uval); gives 313233b. printf("%d\n", var.ival); gives 3224115
```

- c. printf("%c\n", var.cval); gives 3 (based on the byte ordering)
- d. printf("%s\n", var.sval); gives **321** (based on the byte ordering)
- e. printf("%f\n", var.fval); gives **0.000000** (actually a number very close to zero)
- f. printf("%e\n", var.fval); gives 4.517947e-39

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For all enquiries, please email the class account at <u>cs1521@cse.unsw.edu.au</u>

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