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Quicksort

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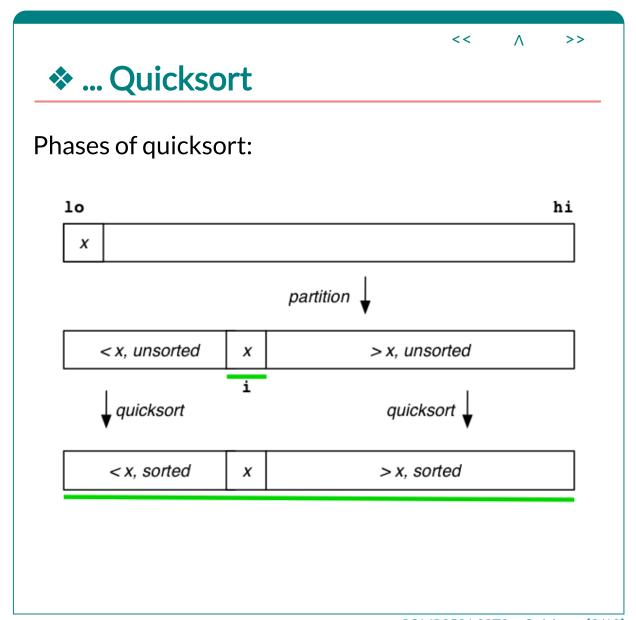
Quicksort

Previous sorts were all $O(n^k)$ (where k>1).

We can do better ...

Quicksort: basic idea

- choose an item to be a "pivot"
- re-arrange (partition) the array so that
 - all elements to left of pivot are smaller than pivot
 - all elements to right of pivot are greater than pivot
- (recursively) sort each of the partitions

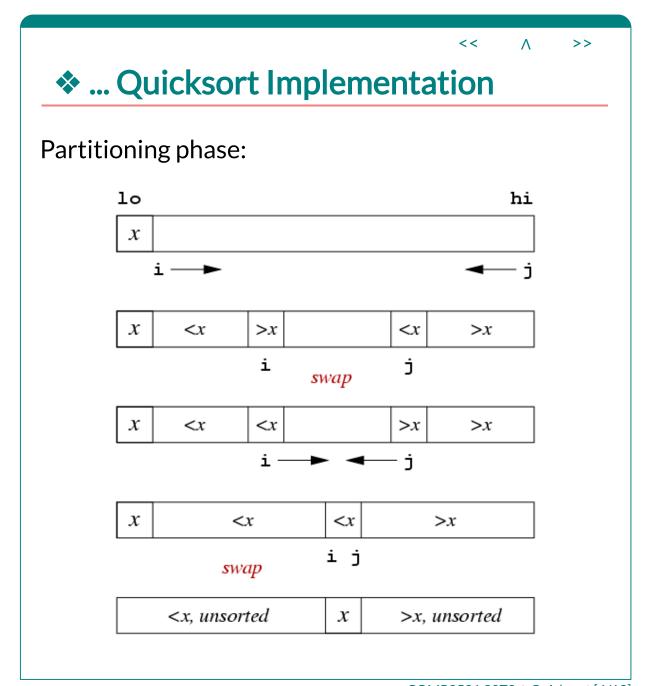


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Quicksort Implementation

Elegant recursive solution ...

```
void quicksort(Item a[], int lo, int hi)
{
   int i; // index of pivot
   if (hi <= lo) return;
   i = partition(a, lo, hi);
   quicksort(a, lo, i-1);
   quicksort(a, i+1, hi);
}</pre>
```



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Partition implementation:

Quicksort Performance

Best case: O(nlogn) comparisons

- choice of pivot gives two equal-sized partitions
- same happens at every recursive level
- each "level" requires approx *n* comparisons
- halving at each level ⇒ log₂n levels

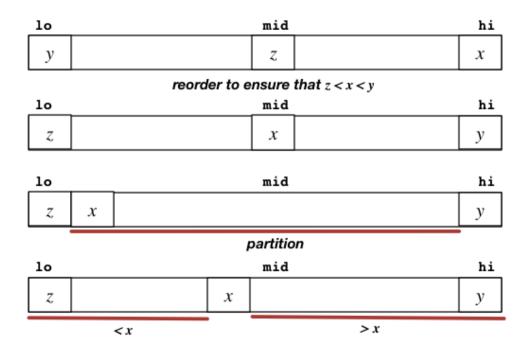
Worst case: $O(n^2)$ comparisons

- always choose lowest/highest value for pivot
- partitions are size 1 and n-1
- each "level" requires approx *n* comparisons
- partitioning to 1 and $n-1 \Rightarrow n$ levels

Quicksort Improvements

Choice of pivot can have significant effect:

- always choosing largest/smallest ⇒worst case
- try to find "intermediate" value by median-ofthree



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... Quicksort Improvements

Median-of-three partitioning:

```
void medianOfThree(Item a[], int lo, int hi)
   int mid = (lo+hi)/2;
   if (less(a[mid],a[lo])) swap(a, lo, mid);
   if (less(a[hi],a[mid])) swap(a, mid, hi);
   if (less(a[mid],a[lo])) swap(a, lo, mid);
   // now, we have a[lo] < a[mid] < a[hi]
   // swap a[mid] to a[lo+1] to use as pivot
   swap(a, mid, lo+1);
void quicksort(Item a[], int lo, int hi)
{
   if (hi <= lo) return;</pre>
   medianOfThree(a, lo, hi);
   int i = partition(a, lo+1, hi-1);
   quicksort(a, lo, i-1);
   quicksort(a, i+1, hi);
}
```

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Another source of inefficiency:

- pushing recursion down to very small partitions
- overhead in recursive function calls
- little benefit from partitioning when size < 5

Solution: handle small partitions differently

- switch to insertion sort on small partitions, or
- don't sort yet; use post-quicksort insertion sort

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Quicksort with thresholding ...

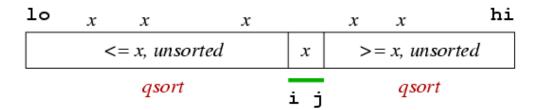
```
void quicksort(Item a[], int lo, int hi)
{
    if (hi-lo < Threshhold) {
        insertionSort(a, lo, hi);
        return;
    }
    medianOfThree(a, lo, hi);
    int i = partition(a, lo+1, hi-1);
    quicksort(a, lo, i-1);
    quicksort(a, i+1, hi);
}</pre>
```



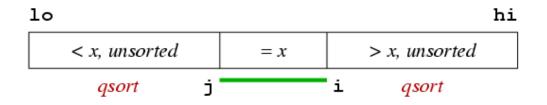
... Quicksort Improvements

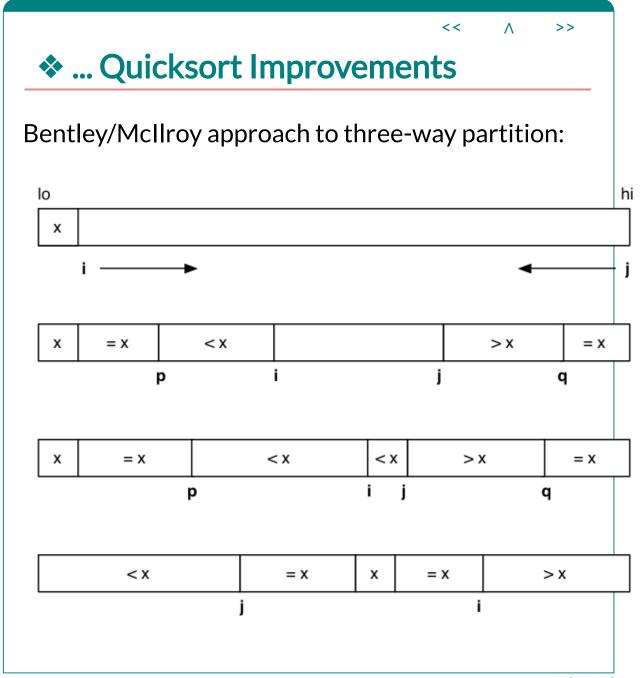
If the array contains many duplicate keys

standard partitioning does not exploit this



 can improve performance via three-way partitioning





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Non-recursive Quicksort

Quicksort can be implemented using an explicit stack:

```
void quicksortStack (Item a[], int lo, int hi)
{
   Stack s = newStack();
   StackPush(s,hi); StackPush(s,lo);
   while (!StackEmpty(s)) {
      lo = StackPop(s);
      hi = StackPop(s);
      if (hi > lo) {
        int i = partition (a,lo,hi);
        StackPush(s,hi); StackPush(s,i+1);
        StackPush(s,i-1); StackPush(s,lo);
      }
   }
}
```

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Produced: 20 Jul 2020