

# Graph Basics

---

- Graphs
- Properties of Graphs
- Graph Terminology

## ❖ Graphs

---

Many applications require

- a **collection** of **items** (i.e. a set)
- **relationships**/connections between items

Examples:

- **maps**: items are cities, connections are roads
- **web**: items are pages, connections are hyperlinks

Collection types you're familiar with

- lists ... linear sequence of items (COMP1511)
- trees ... branched hierarchy of items (Weeks 02/03)

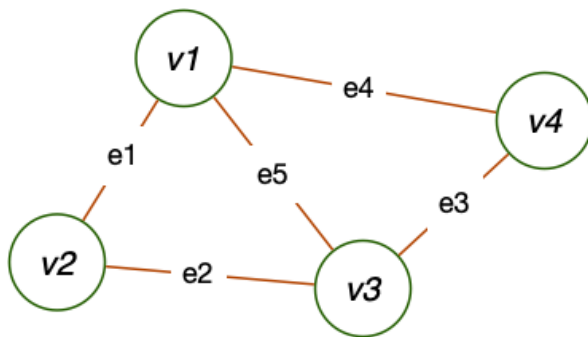
**Graphs** are more general ... allow arbitrary connections

## ❖ ... Graphs

A graph  $G = (V, E)$

- $V$  is a set of **vertices**
- $E$  is a set of **edges** (subset of  $V \times V$ )

Example:



$$V = \{ v1, v2, v3, v4 \}$$

$$E = \{ e1, e2, e3, e4, e5 \}$$

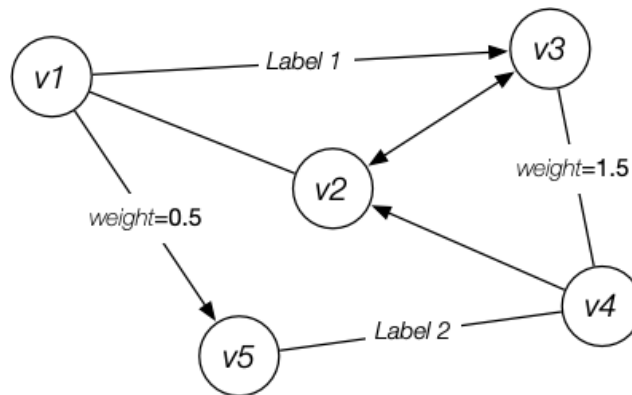
or

$$E = \{ (v1, v2), (v2, v3), (v3, v4), (v1, v4), (v1, v3) \}$$

## ❖ ... Graphs

Nodes are distinguished by a unique identifier

Edges may be (optionally) directed, labelled and/or weighted



## ❖ ... Graphs

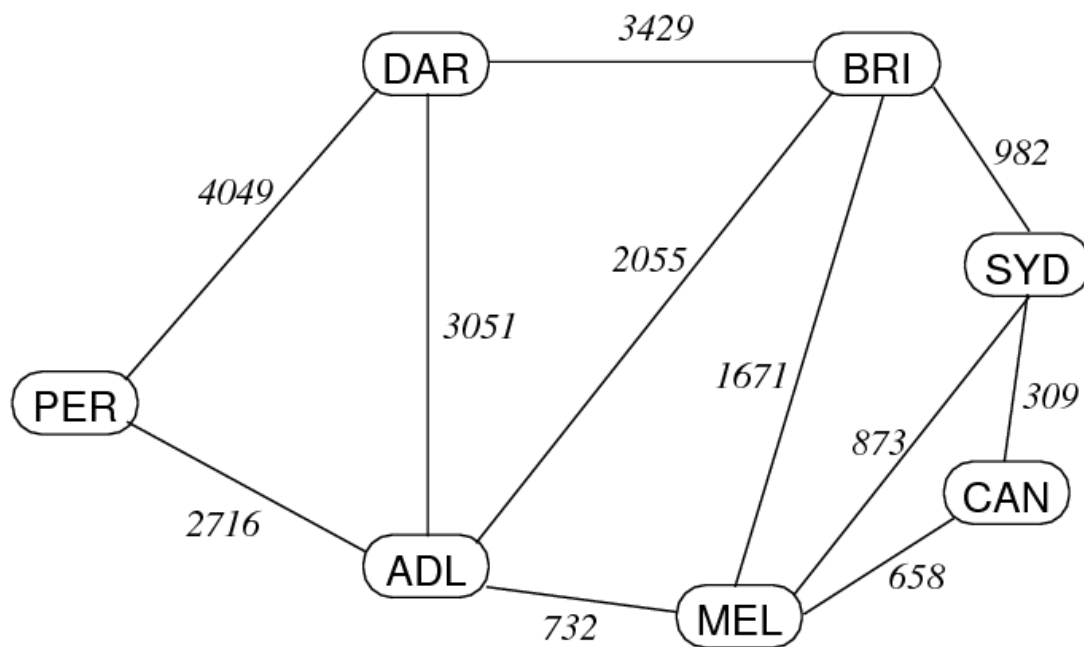
### A real example: Australian road distances

Distance	Adelaide	Brisbane	Canberra	Darwin	Melbourne	Perth	Sydney
Adelaide	-	2055	1390	3051	732	2716	1605
Brisbane	2055	-	1291	3429	1671	4771	982
Canberra	1390	1291	-	4441	658	4106	309
Darwin	3051	3429	4441	-	3783	4049	4411
Melbourne	732	1671	658	3783	-	3448	873
Perth	2716	4771	4106	4049	3448	-	3972
Sydney	1605	982	309	4411	873	3972	-

Notes: vertices are cities, edges are distance between cities, symmetric

## ❖ ... Graphs

Alternative representation of above:



COMP2521 20T2 ♦ Graph Basics [5/15]

## ❖ ... Graphs

---

Questions we might ask about a graph:

- is there a way to get from item A to item B?
- what is the best way to get from A to B?
- which items are directly connected ( $A \leftrightarrow B$ )?

Graph algorithms are generally more complex than tree/list ones:

- no implicit order of items
- graphs may contain cycles
- concrete representation is less obvious
- algorithm complexity depends on connection complexity

## ❖ Properties of Graphs

---

Terminology:  $|V|$  and  $|E|$  (cardinality) normally written just as  $V$  and  $E$ .

A graph with  $V$  vertices has at most  $V(V-1)/2$  edges.

The ratio  $E:V$  can vary considerably.

- if  $E$  is closer to  $V^2$ , the graph is **dense**
- if  $E$  is closer to  $V$ , the graph is **sparse**
  - Example: web pages and hyperlinks

Knowing whether a graph is sparse or dense is important

- may affect choice of data structures to represent graph
- may affect choice of algorithms to process graph



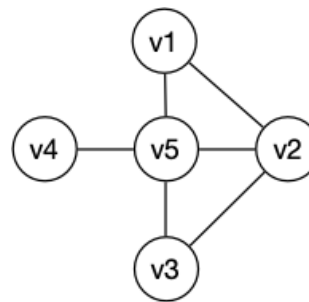
## ❖ Graph Terminology

For an edge  $e$  that connects vertices  $v$  and  $w$

- $v$  and  $w$  are **adjacent** (neighbours)
- $e$  is **incident** on both  $v$  and  $w$

**Degree** of a vertex  $v$

- number of edges incident on  $e$



$\text{degree}(v1) = 2$   
 $\text{degree}(v2) = 3$   
 $\text{degree}(v3) = 2$   
 $\text{degree}(v4) = 1$   
 $\text{degree}(v5) = 4$

Synonyms:

- vertex = node
- edge = arc = link (Note: some people use arc for *directed* edges)

## ❖ ... Graph Terminology

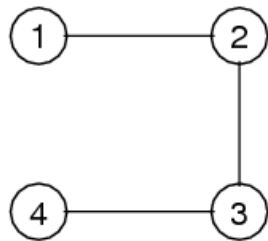
**Path:** a sequence of vertices where

- each vertex has an edge to its predecessor

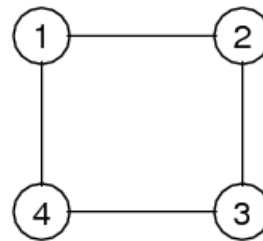
**Cycle:** a path where

- last vertex in path is same as first vertex in path

**Length** of path or cycle = #edges



*Path: 1-2, 2-3, 3-4*



*Cycle: 1-2, 2-3, 3-4, 4-1*

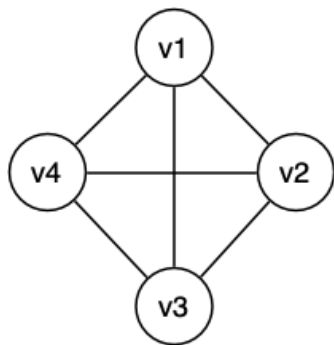
## ❖ ... Graph Terminology

### Connected graph

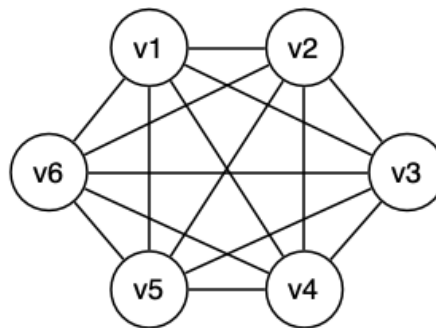
- there is a *path* from each vertex to every other vertex
- if a graph is not connected, it has  $\geq 2$  **connected components**

### Complete graph $K_V$

- there is an *edge* from each vertex to every other vertex
- in a complete graph,  $E = V(V-1)/2$



Complete  
Graphs



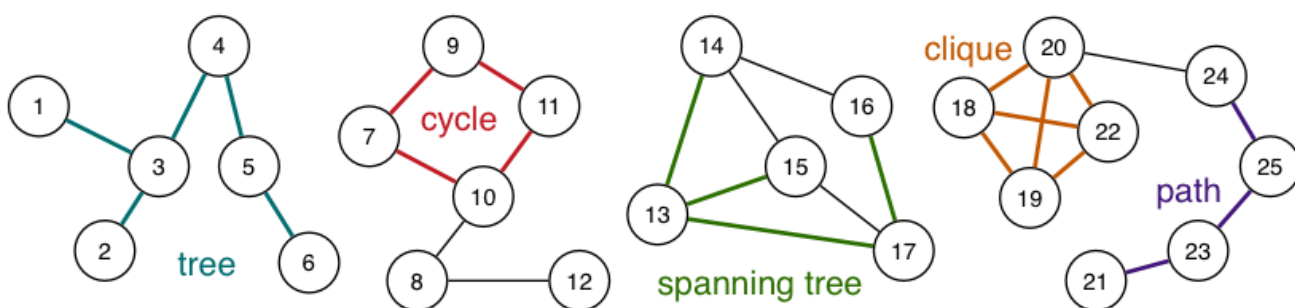
## ❖ ... Graph Terminology

**Tree:** connected (sub)graph with no cycles

**Spanning tree:** tree containing all vertices

**Clique:** complete subgraph

Consider the following *single graph*:



This graph has 25 vertices, 32 edges, and 4 connected components

Note: The entire graph has no spanning tree; what is shown in green is a spanning tree of the third connected component

## ❖ ... Graph Terminology

---

A **spanning tree** of connected graph  $G = (V, E)$

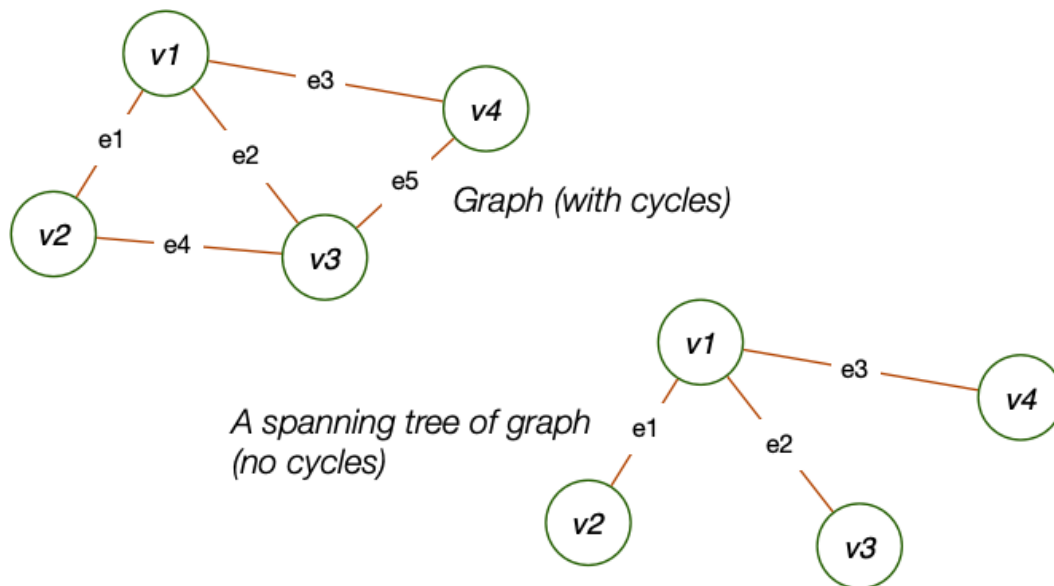
- is a subgraph of  $G$  containing all of  $V$
- and is a single tree (connected, no cycles)

A **spanning forest** of non-connected graph  $G = (V, E)$

- is a subgraph of  $G$  containing all of  $V$
- and is a set of trees (not connected, no cycles),
  - with one tree for each *connected component*

## ❖ ... Graph Terminology

Can form spanning tree from graph by removing edges



Many possible spanning trees can be formed. Which is "best"?

## ❖ ... Graph Terminology

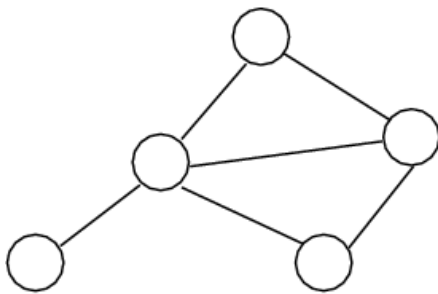
### Undirected graph

- $edge(u,v) = edge(v,u)$ , no self-loops (i.e. no  $edge(v,v)$ )

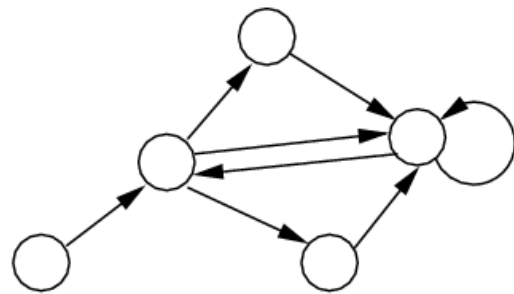
### Directed graph

- $edge(u,v) \neq edge(v,u)$ , can have self-loops (i.e.  $edge(v,v)$ )

Examples:



*Undirected graph*



*Directed graph*

## ❖ ... Graph Terminology

---

Other types of graphs ...

### Weighted graph

- each edge has an associated value (weight)
- e.g. road map (weights on edges are distances between cities)

### Multi-graph

- allow multiple edges between two vertices
- e.g. function call graph ( $f()$  calls  $g()$  in several places)

### Labelled graph

- edges have associated labels
- can be used to add semantic information



Produced: 21 Jun 2020