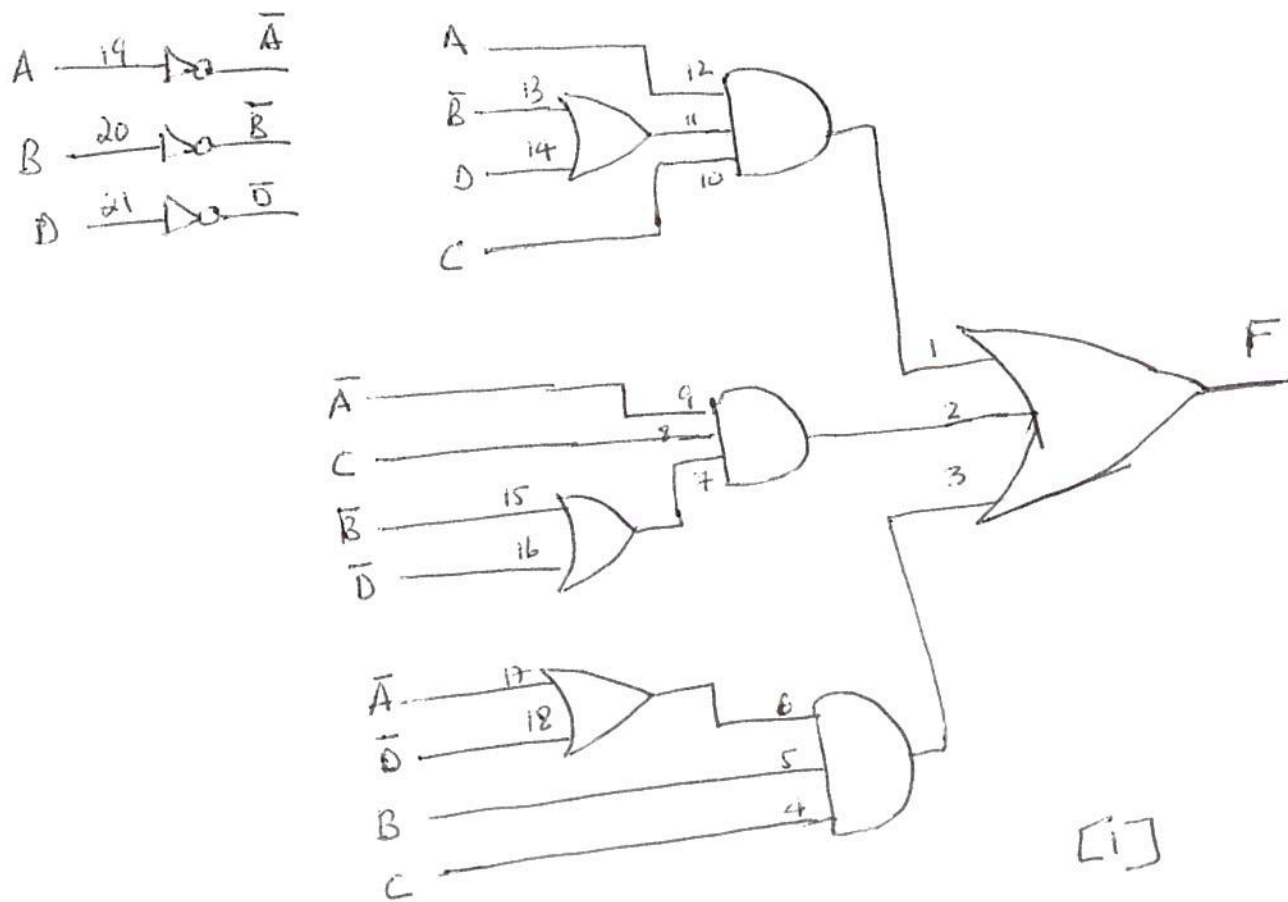


Question 1.

$$F(A, B, C, D) = AC(\bar{B} + D) + \bar{A}C(\bar{B} + \bar{D}) + BC(\bar{A} + \bar{D})$$

- a) No of terms = 6
 No of literals = $4 + 4 + 4 = 12$
 No of Complimented literals = 3
 Total gate input cost = $6 + 12 + 3 = 21$



$$GIC = 21$$

b) $F = A\bar{B}C + ACD + \bar{A}\bar{B}C + \bar{A}CD + \bar{A}BC + BC\bar{D}$

Truth table

	A	B	C	D	F
0	0	0	0	0	0
1	0	0	0	1	0
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	0
5	0	1	0	1	0
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	0
9	1	0	0	1	0
10	1	0	1	0	1
11	1	0	1	1	1
12	1	1	0	0	0
13	1	1	0	1	0
14	1	1	1	0	1
15	1	1	1	1	1

$$F = \sum m(2, 3, 6, 7, 10, 11, 14, 15)$$

[2]

c) $F = \sum m(2, 3, 6, 7, 10, 11, 14, 15)$ [2]

d) $F = \prod M(0, 1, 4, 5, 8, 9, 12, 13)$ [2]

e) Four Variable K-map

AB \ CD		C			
		D		D	
A	B	0	0	1	1
	0	0	0	1	1
	0	0	0	1	1
	0	0	0	1	1

$$F = C$$

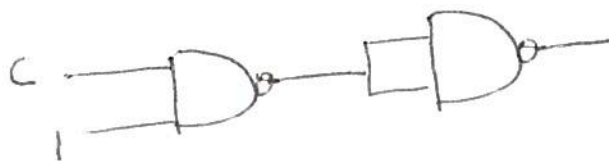
There is only one prime implicant, which is also essential. [4]

f) product of sums

Combine the zeros to find \bar{F}

$$\bar{F} = \bar{C} \Rightarrow F = C \text{ is also product of sums [3]}$$

g) $F = C \cdot 1$



possible answer
no implementation is needed as the function is a single variable [2]

h) GIC for implementation in (g) is 4 \rightarrow

[But if no implementation, total reduction in]
GIC = 21

otherwise, total reduction in gic is $21 - 4 = 17$ [1]

$$\begin{aligned}
 \text{ii) } (23B4.A5)_{12} &= 2 \times 12^3 + 3 \times 12^2 + 11 \times 12^1 + 4 \times 12^0 \\
 &\quad + 10 \times 12^{-1} + 5 \times 12^{-2} \\
 &= 3456 + 432 + 132 + 4 + 0.8333\ldots + 0.0347222\ldots \\
 &\quad + 0.26805222\ldots \\
 &= (4024.86805222\ldots)_{10}
 \end{aligned}$$

Integral part to base 5

$$\begin{array}{rcl}
 4024 \div 5 & = & 804 & 4 \\
 804 \div 5 & = & 160 & 4 \\
 160 \div 5 & = & 32 & 0 \\
 32 \div 5 & = & 6 & 2 \\
 6 \div 5 & = & 1 & 1 \\
 1 \div 5 & = & 0 & 1
 \end{array}$$

$$(112044)_5$$

Fractional part

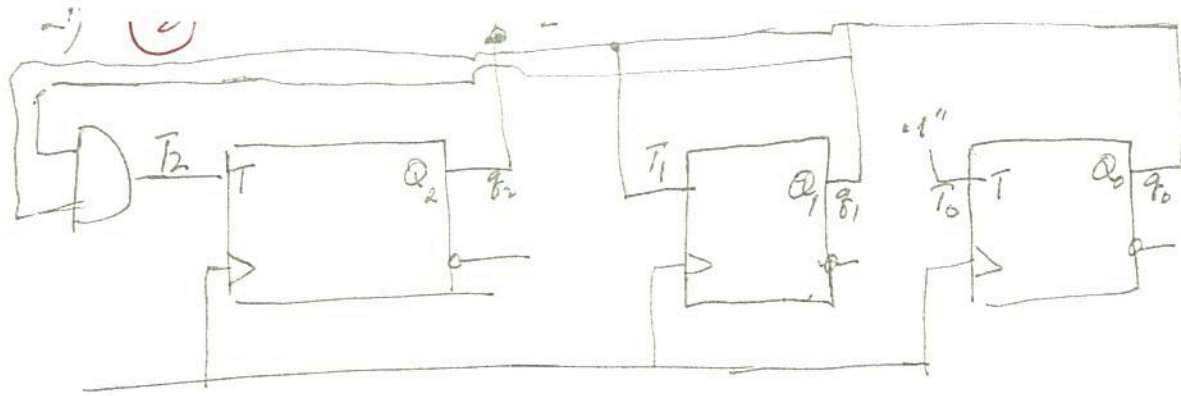
$$\begin{array}{rcl}
 0.86805222\ldots \times 5 & = & 4 \\
 0.34026111\ldots \times 5 & = & 1 \\
 0.70130555\ldots \times 5 & = & 3 \\
 0.50652777\ldots \times 5 & = & 2 \\
 0.53263888\ldots \times 5 & = & 2
 \end{array}$$

$$(0.4132)_5$$

Truncating into 4 digits after radix point

$$(112044.4132)_5$$

[3]



Excitation equations

i)

$$T_0 = 1$$

$$T_1 = Q_0$$

$$T_2 = Q_0 \cdot Q_1$$

~~$C_0 = Q_0$~~
 ~~$C_1 = Q_1$~~
 ~~$C_2 = Q_2$~~

$$C_0 = Q_0$$

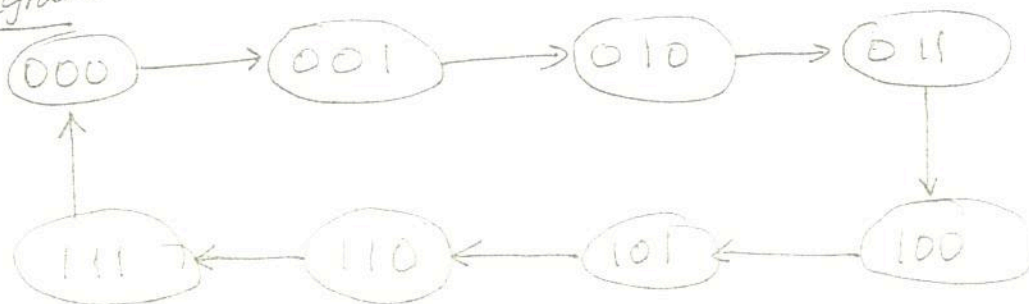
$$C_1 = Q_1$$

$$C_2 = Q_2$$

State table

Present state			Next state			Excitation inputs		
Q_2	Q_1	Q_0	Q_2	Q_1	Q_0	T_2	T_1	T_0
0	0	0	0	0	1	0	0	1
0	0	1	0	1	0	0	1	1
0	1	0	0	1	1	0	0	1
0	1	1	1	0	0	1	1	1
1	0	0	1	0	1	0	0	1
1	0	1	1	1	0	0	1	1
1	1	0	1	1	1	0	0	1
1	1	1	0	0	0	1	1	1

State diagram



ii) Moore type (Moore counter)

iii) Three bit binary Counter

iv) Characteristic table for JK flip flop

J	K	$Q(t+1)$	$Q(t)$	$Q(t+1)$	J	K
0	0	$Q(t)$	0	0	0	X
0	1	0	0	1	1	X
1	0	1	1	0	X	1
1	1	$\bar{Q}(t)$	1	1	X	0

Present state	Next state	Flip-flop inputs
$q_2 \ q_1 \ q_0$	$Q_2 \ Q_1 \ Q_0$	$J_2 \ K_2 \ J_1 \ K_1 \ J_0 \ K_0$
0 0 0	0 0 1	0 X 0 X 1 X
0 0 1	0 1 0	0 X 1 X X 1
0 1 0	0 1 1	0 X X 0 1 X
0 1 1	1 0 0	1 X X 1 X 1
1 0 0	1 0 1	X 0 0 X 1 X
1 0 1	1 1 0	X 0 1 X 1 X
1 1 0	1 1 1	X 1 X 1 X 1
1 1 1	0 0 0	

K_0 :

$q_1 \ q_0$	q_1
q_2	
0	X 1 1 X
1	X 1 1 X

$$K_0 = 1$$

K_1 :

$q_1 \ q_0$	q_1
q_2	
0	X X 1 0
1	X X 1 0

$$K_1 = q_0$$

J_0 :

$q_1 \ q_0$	q_1
q_2	
0	1 X X 1
1	1 X X 1

$$J_0 = 1$$

J_1 :

$q_1 \ q_0$	q_1
q_2	
0	0 1 X X
1	0 1 X X

$$J_1 = q_0$$

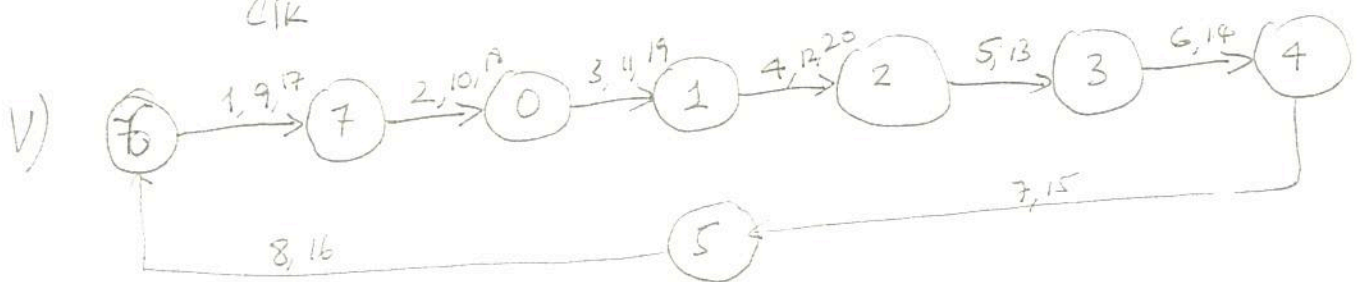
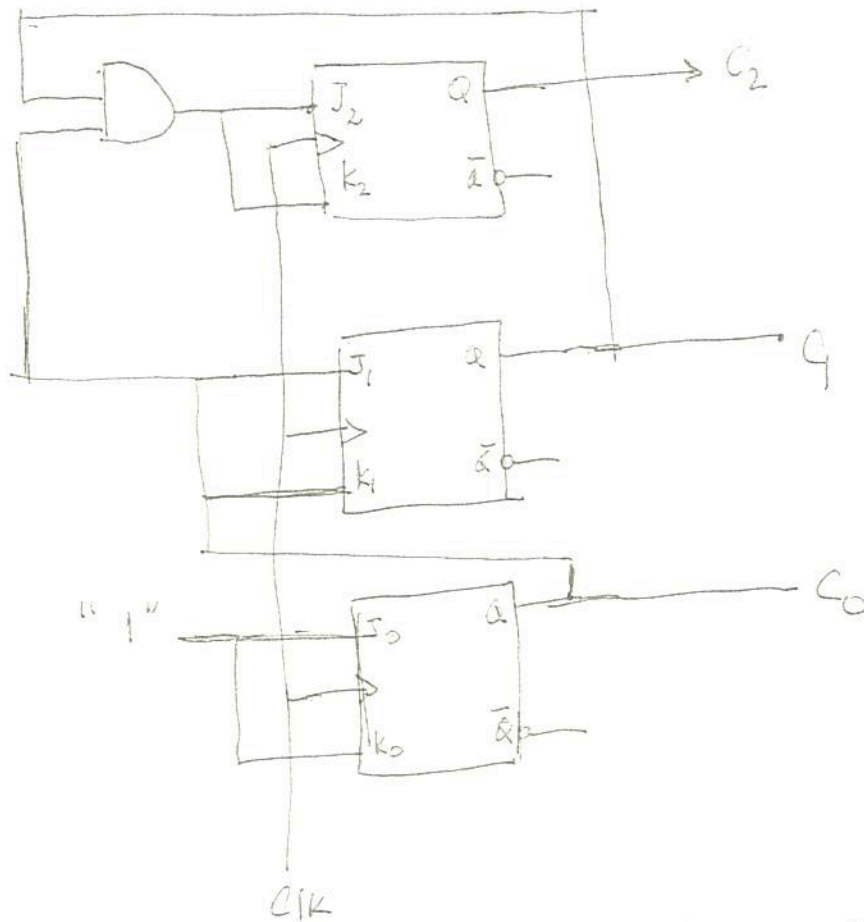
52.

q_2	q_1	q_0	
0	0	1	0
X	X	X	X

$J_2 = q_0 q_1$

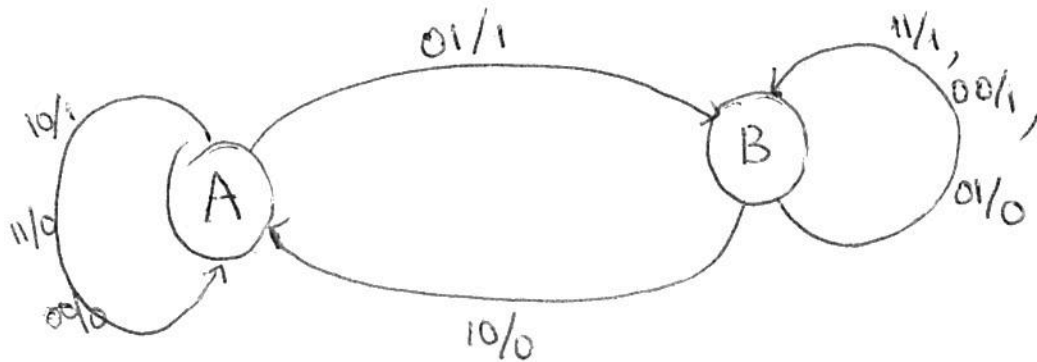
q_2	q_1	q_0	
X	X	X	X
0	0	1	0

$K_2 = q_0 q_1$



The output of the system will be 010 after 20 clock cycle.

- 3) a) Two possible states are available
 A - with no borrow
 B - with borrow



[3]

- b) state assignment

A = 0 B = 1

State table

Present state	next state, output			
	xy = 00	xy = 01	xy = 11	xy = 10
A	A, 0	B, 1	A, 0	A, 1
B	B, 1	B, 0	B, 1	A, 0

with state assignment

Present state	Next state, output			
	xy = 00	xy = 01	xy = 11	xy = 10
0	0, 0	1, 1	0, 0	0, 1
1	1, 1	1, 0	1, 1	0, 0

Q_0^+

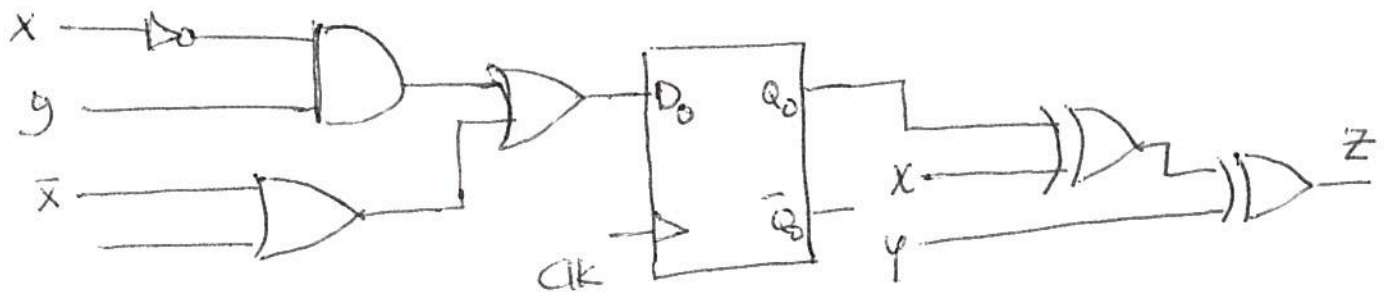
	$\overline{x}y$	$x\overline{y}$	xy	$\overline{x}\overline{y}$
Q_0	0	1	0	0
Q_0	1	1	1	0
	\overline{y}	y	\overline{x}	x

$$Q_0^+ = \overline{x}y + Q_0(\overline{x} + y)$$

Truth Table for Z:

	$\overline{x}y$		x	
Q_0	0	1	0	1
Q_0	1	0	1	0
	y			

$$Z = Q_0 \oplus X \oplus Y$$



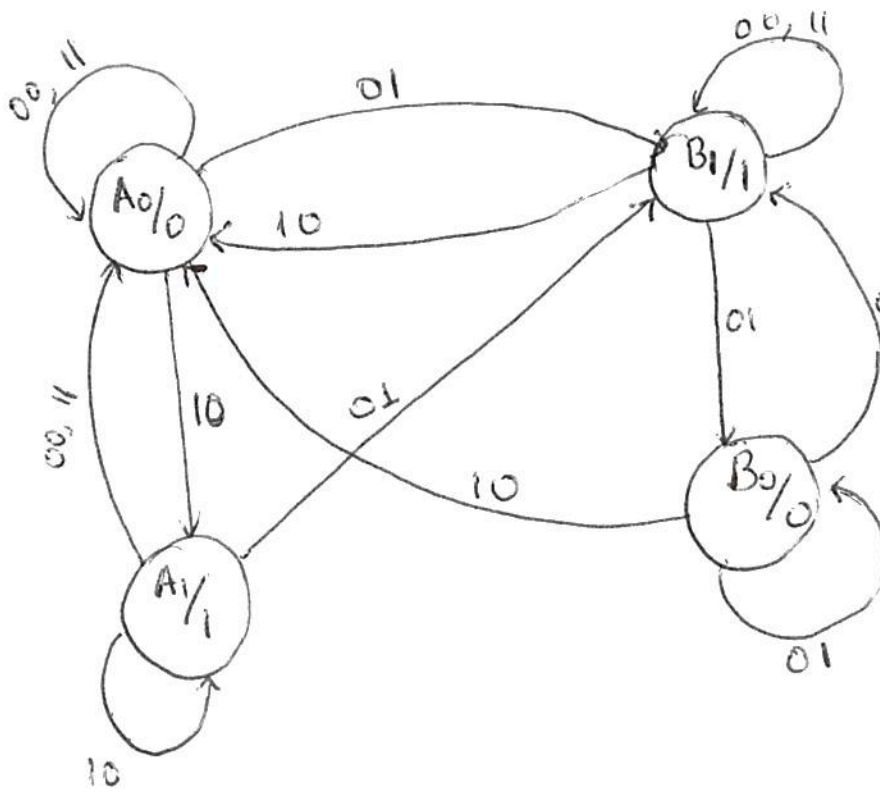
[4]

c) module Seq-binary-sub(x,y,clk);
 input x,y,clk;
 output z;
 wire A;
 reg B;
 assign A = (~x & y) | (B & (~x | y));
 assign z = x ^ y ^ B;
 DFF D(B,A,clk);
endmodule

module DFF(Q,DO,clk);
 input DO,clk;
 output reg Q;
 always @(posedge clk)
 Q <= A;
endmodule

[3]

d)



A_0 = no borrow and sub is zero
 A_1 = no borrow and sub one
 B_0 = borrow ~~and~~ sub is 0
 B_1 = borrow but sub is 1

[4]

ii)

a)

A	B	C	y
0	0	0	0
0	1	0	1
1	0	0	0
1	1	0	0
0	0	1	1
0	1	1	1
1	0	1	0
1	1	1	1

$A < B$ y or $C = 1$
 $A \geq B$ y or $C = 0$

Starting value should be low to ~~assure~~ take into account the possible equality

[2]

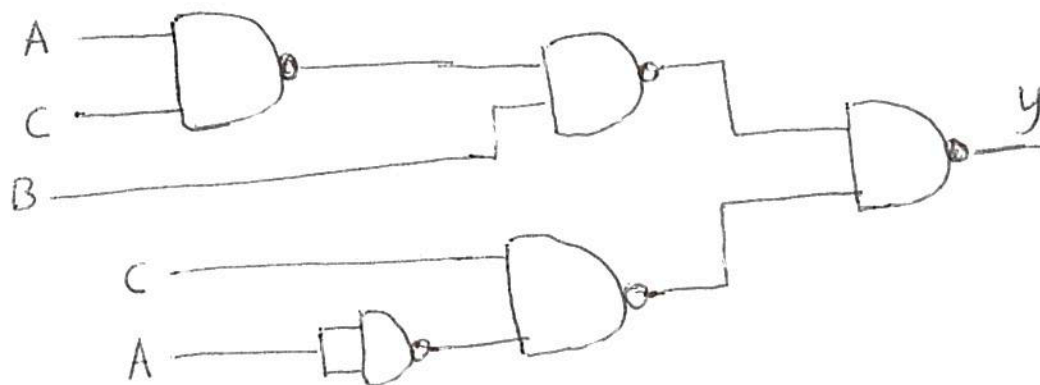
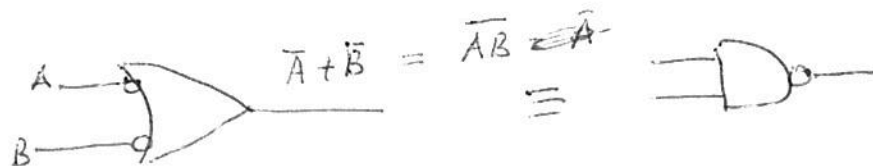
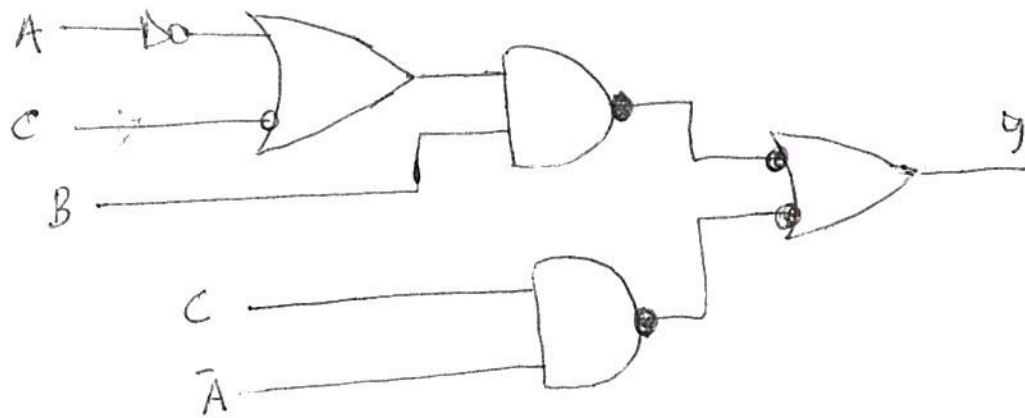
C should be Low

b)

$y =$

				A
		AB		
	C	0	1	0
	C	1	0	1
				B

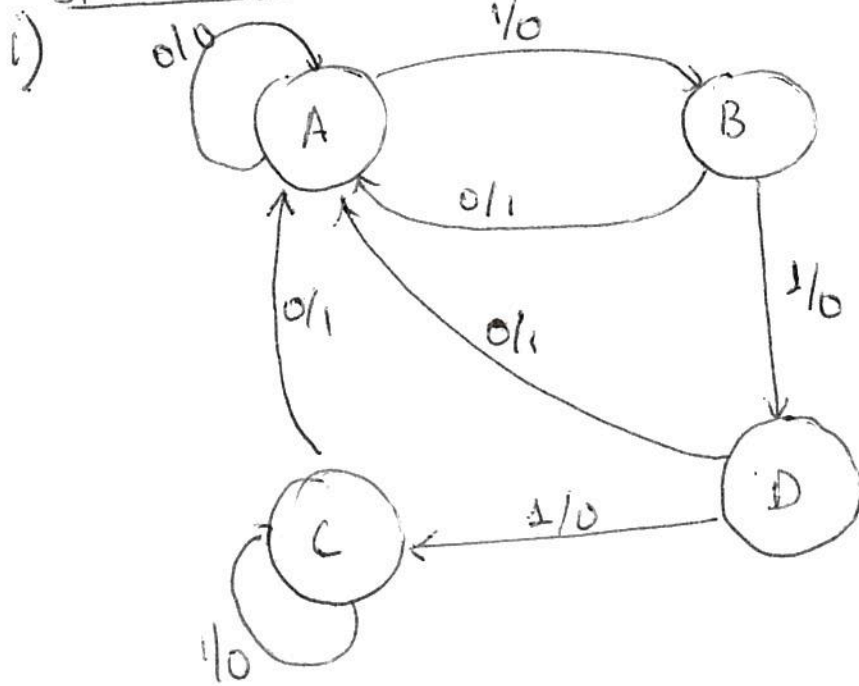
$$y = B(\bar{A} + C) + C\bar{A}$$



[4]

Questi 7

state diagram

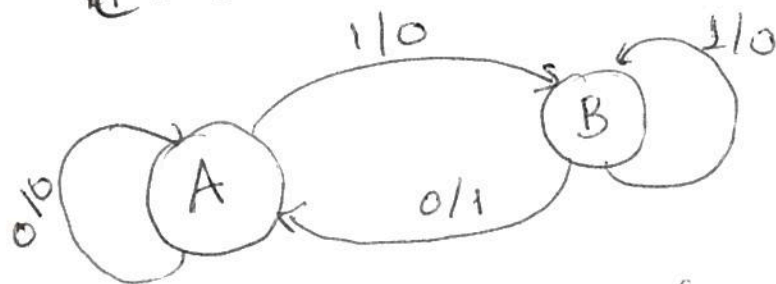


State table

Present State	next state, output	
	X=0	X=1
A	A, 0	B, 0
B	A, 1	D, 0
C	A, 0	C, 0
D	A, 1	C, 0

State reduction

~~A~~ = D = B, A reduces to two states



a) Mealy, output is a function of X [2]

b)

Present State	Next State, output	
	$X=0$	$X=1$
Q_0	Q_0^+, y	Q_0^+, y
0	0, 0	1, 0
1	0, 1	1, 0

A = 0
B = 1

Q_0^+

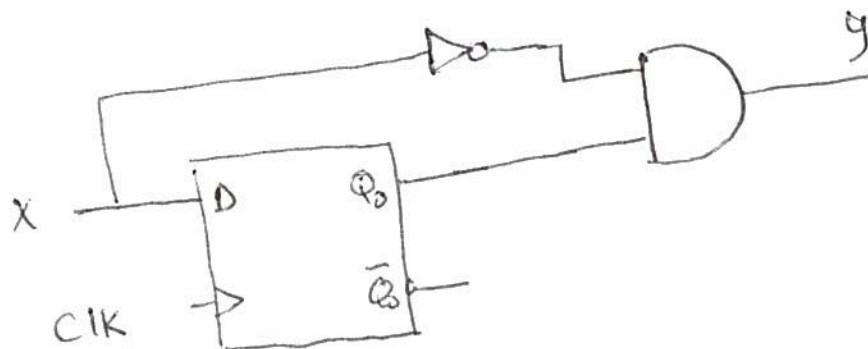
Q_0	X	
0	0	0
0	1	1
1	0	0
1	1	1

$$Q_0^+ = X$$

y

Q_0	X	
0	0	0
0	1	0
1	0	1
1	1	0

$$y = Q_0 \bar{X}$$



[4]

c)

input	0	1	0	1	1	0	0
State	B	A	B	A	B	B	A
Output	1	0	1	0	0	1	0

[2]

ii)

$$\bar{A}(A+C)(\bar{A}+B)(A+B) = BC$$

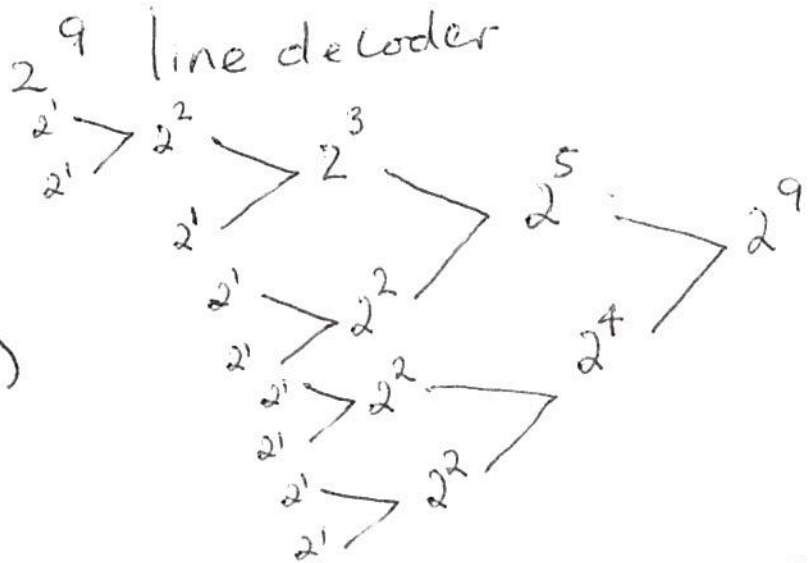
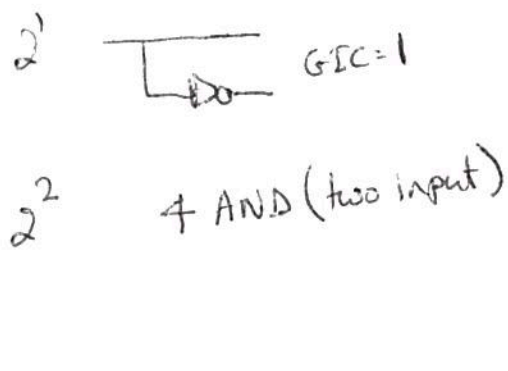
~~$$B(A+C)B = BA+BC$$~~

$$A+B=1 \quad \text{and} \quad A=B=0 \Rightarrow A=\bar{B}$$

$$\begin{aligned} B(A+C)(B+B) &= B(A+C) = BA+BC \\ &= AB+BC = 0+BC = BC \end{aligned} \quad [4]$$

iii)

9 to 2^9 line decoder



$$\begin{aligned} \text{GIC} &= 2 \cdot (2^9 + 2^4 + 2^5 + 2^2 + 2^2 + 2^2 + 2^3 + 2^2) + 9 \\ &= 2 (512 + 32 + 16 + 4 + 4 + 4 + 8 + 4) + 9 \\ &= \underline{\underline{1177}} \end{aligned} \quad [4]$$

iv)

a) Bidirectional Shift register

[2]

b) Moore type System

[2]

Question 5

1331
2200

i)

A	B	C	D		B ₃	B ₂	B ₁	B ₀
A ₃	A ₂	A ₁	A ₀					
0	0	0	0		0	0	0	0
0	0	0	1		1	1	1	1
0	0	1	0		1	1	1	0
0	0	1	1		1	1	0	1
0	1	0	0		1	1	0	0
0	1	0	1		1	0	1	1
0	1	1	0		1	0	1	0
0	1	1	1		1	0	0	1
1	0	0	0		1	0	0	0
1	0	0	1		0	1	1	1
1	0	1	0		0	1	1	0
1	0	1	1		0	1	0	1
1	1	0	0		0	1	0	0
1	1	0	1		0	0	1	1
1	1	1	0		0	0	1	0
1	1	1	1		0	0	0	1

			<u>A₃</u>	
<u>B₀</u>	0	1	1	0
	0	1	1	0
<u>A₃</u>	0	1	1	0
	0	1	1	0
			<u>A₀</u>	

$$B_0 = A_0$$

B_1

	$A_1 A_0$		
$A_3 A_2$	0	1	0
	0	1	0
A_3	0	1	0
	0	1	0

A_1
 A_2
 A_0

$$B_1 = A_1 \bar{A}_0 + A_0 \bar{A}_1$$

$$= A_0 \oplus A_1$$

B_2

	A_1		
A_3	0	1	1
	1	0	0
A_3	1	0	0
	0	1	1

A_2
 A_0

$$B_2 = \bar{A}_3 \bar{A}_1 (A_0 \oplus A_2) + A_3 \bar{A}_1 (A_0 \oplus A_2) + A_1 \bar{A}_2$$

$$= (A_0 \oplus A_2) (\bar{A}_3 \bar{A}_1 + A_3 \bar{A}_1) + A_1 \bar{A}_2$$

$$= \bar{A}_1 (A_0 \oplus A_2) + A_1 \bar{A}_2$$

$$= \bar{A}_1 (A_0 \oplus A_2) \oplus A_1 \bar{A}_2$$

B_3

	A_1		
A_3	0	1	1
	1	1	1
A_3	0	0	0
	1	0	0

A_2
 A_0

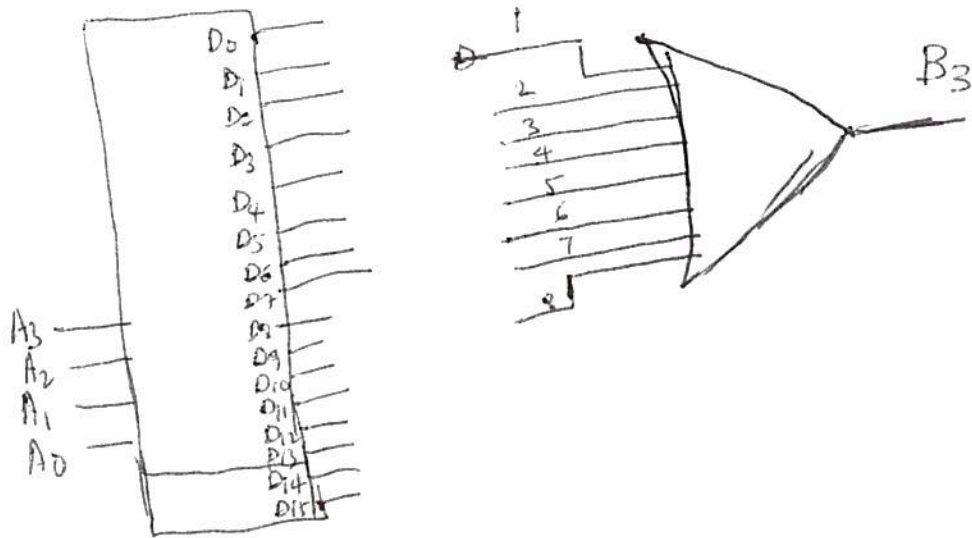
$$B_3 = (A_3 \oplus A_0) \bar{A}_1 \bar{A}_2 + A_1 \bar{A}_3 + A_2 \bar{A}_3$$

$$= (\bar{A}_1 \bar{A}_2 (A_0 \oplus A_3) \oplus A_1 \bar{A}_3) + A_2 \bar{A}_3$$

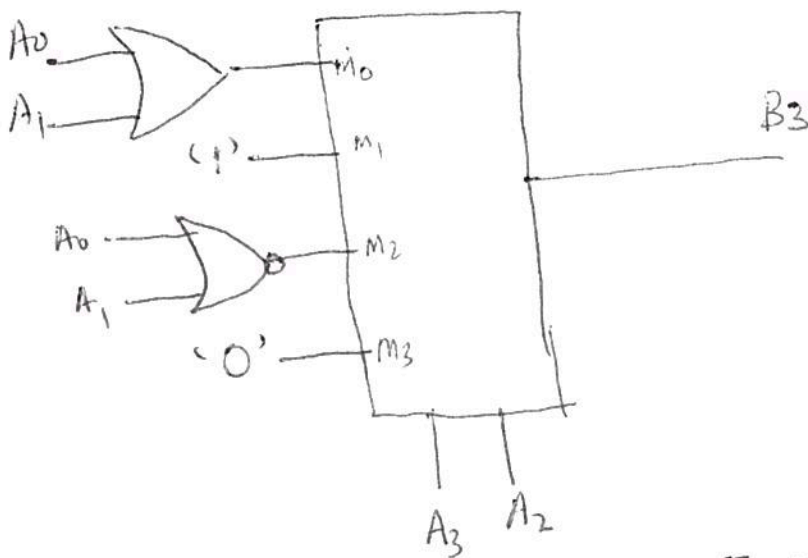
$$= (\bar{A}_1 \bar{A}_2 (A_0 \oplus A_3) \oplus A_1 \bar{A}_3) \oplus A_2 \bar{A}_3$$

b) $B_3 = \sum m(1, 2, 3, 4, 5, 6, 7, 8)$

using decoder



Multiplexer

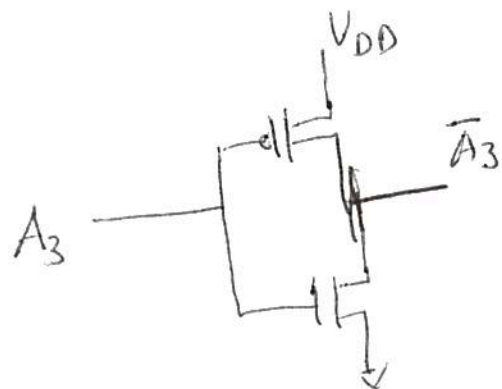
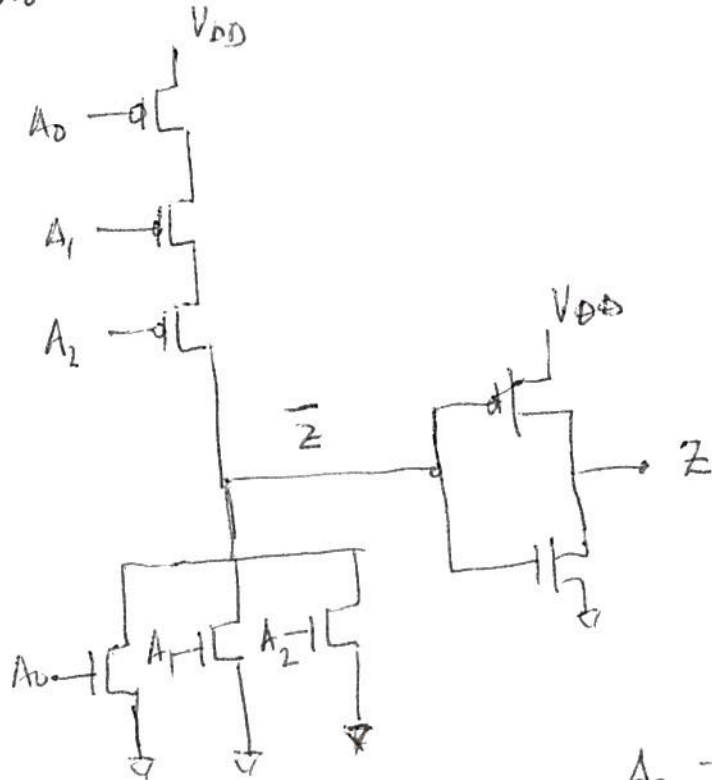


[2]

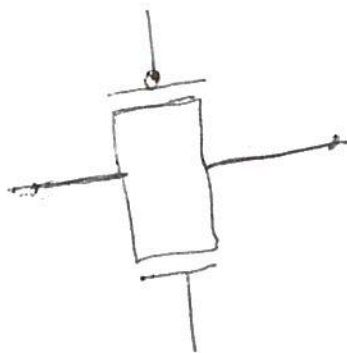
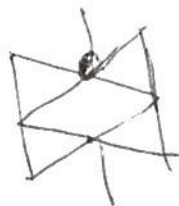
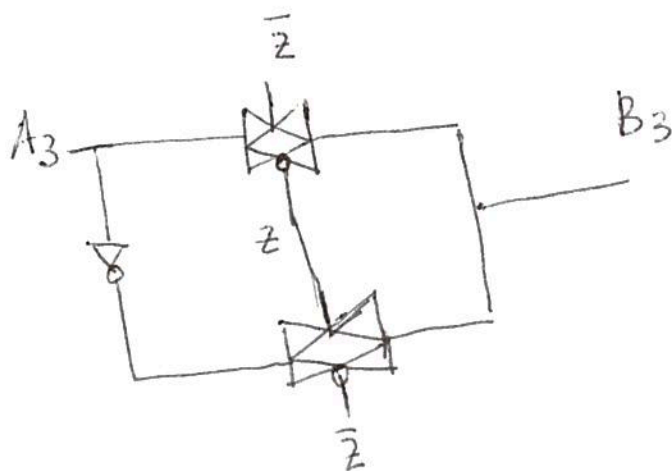
$$\begin{aligned}
 c) \quad B_3 &= A_1 \bar{A}_3 + A_2 \bar{A}_3 + \bar{A}_3 A_0 + \bar{A}_0 \bar{A}_1 \bar{A}_2 A_3 \\
 &= \bar{A}_3 (A_1 + A_2 + A_0) + A_3 (\bar{A}_0 \bar{A}_1 \bar{A}_2) \\
 &= \bar{A}_3 (A_1 + A_2 + A_0) + A_3 (\overline{A_0 + A_1 + A_2}) \\
 &= A_3 \oplus (A_0 + A_1 + A_2) \\
 &= A_3 \oplus Z
 \end{aligned}$$

where $Z = A_0 + A_1 + A_2$

$$\bar{Z} = \bar{A}_0 \cdot \bar{A}_1 \cdot \bar{A}_2$$



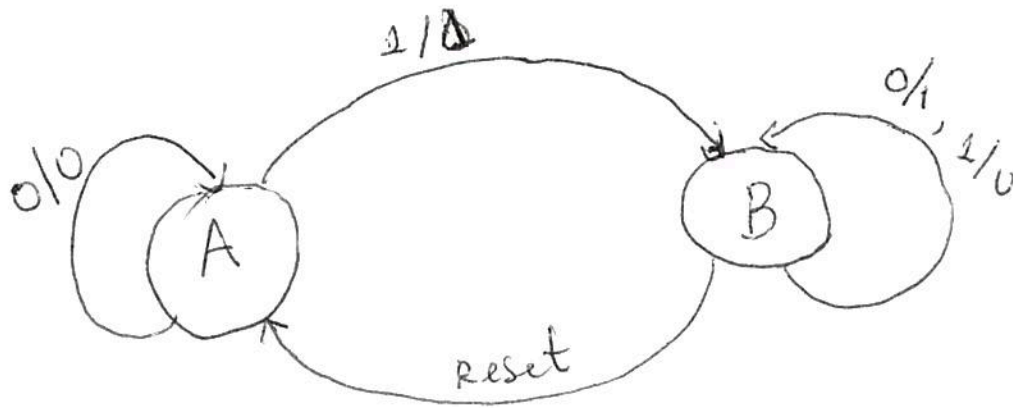
$$B_3 = A_3 \oplus Z$$



Transmission gate

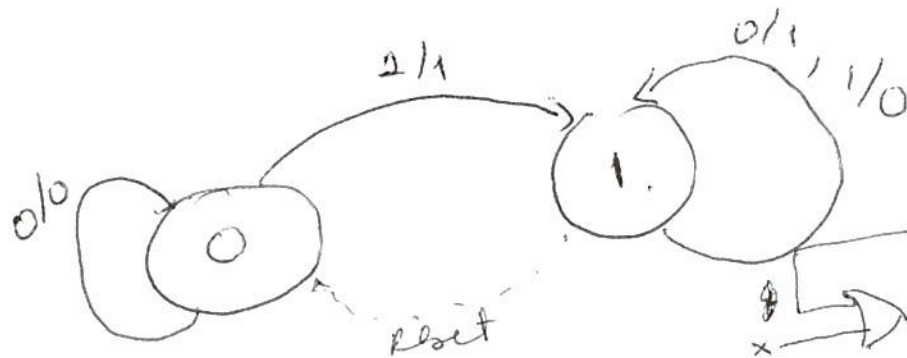
ii) There will be two states

A \rightarrow state in which arriving bit is unchanged
 B \rightarrow state in which arriving bit is complemented



state assignment
 A = 0 B = 1

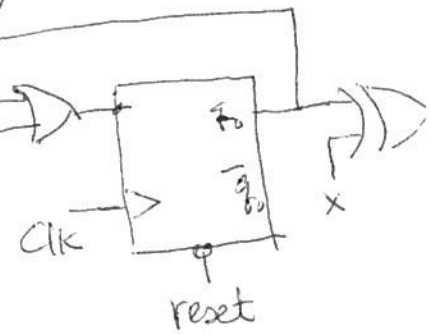
[4]



state table
 present state

q_0
 0
 1

next state, output
 $x=0$
 q_0^+
 0, 0
 1, 1
 $x=1$
 q_0^+
 1, 1
 1, 0



q_0^+	q_0	x	
		0	1
0	0	0	1
1	1	1	1

$$q_0^+ = q_0 + x$$

q_0	x	
	0	1
0	0	1
1	1	0

$$z = q_0 \oplus x$$

ii)

$$\begin{array}{r} R_1 = 1000111100010101 \\ R_2 = 1010101110010111 \\ \hline 0011101010101100 \end{array}$$

$$V = C_{out} \oplus C_{in} = 1 \oplus 0 = 1$$

$$C_o = C_{out} = 1$$

$$N = 0$$

$$Z = 0$$

[4]

iv) a) When C is low and A is low

Q_6 is off

Q_1 is off

Q_5 is ON

Q_7 is off

Q_2 is off

Q_4 is ON

Q_8 is off

Q_3 is off

Y is High

b) When C is low and A is high

Q_6 is off

Q_1 is ON

Q_5 is off

Q_7 is off

Q_2 is ON

Q_4 is off

Q_8 is off

Q_3 is ON

Y is Low

c) When C is high

Q_6 is ON

Q_5 is OFF

Q_2 is OFF

Q_7 is ON

Q_4 is OFF

Q_3 is OFF

Q_8 is ON

Q_1 is OFF

Y is high impedance

Tristate Buffer

[3]