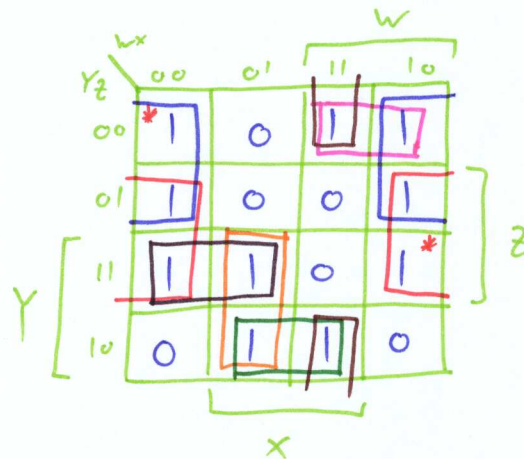


Question 1

i)

$$F = \sum m(0, 1, 3, 6, 7, 8, 9, 11, 12, 14)$$

① a)



② b)

$$\bar{x}\bar{y}, \bar{x}z, \bar{w}yz, \bar{w}xy, xy\bar{z}, wx\bar{z}, w\bar{y}\bar{z}$$

② c)

$$\bar{x}\bar{y}, \bar{x}z$$

① d)

$$F = \bar{x}\bar{y} + \bar{x}z + \bar{w}xy + wx\bar{z}$$

① e)

$$F = \bar{x}(\bar{y} + z) + x(\bar{w}y + w\bar{z})$$

② f)

$$GIC = 18.$$

(ii)  
(5)

$$A\bar{B}\bar{C} + A\bar{B}D + \bar{A}BC + \bar{B}C\bar{D}$$

$$= A\bar{B}\bar{C}(\bar{D}+D) + A\bar{B}D(\bar{C}+C) + \bar{A}BC(\bar{D}+D) + \bar{B}C\bar{D}(\bar{A}+A)$$

$$= \overset{(1)}{A\bar{B}\bar{C}\bar{D}} + \overset{(2)}{A\bar{B}\bar{C}D} + \overset{(2)}{A\bar{B}C\bar{D}} + \overset{(3)}{A\bar{B}CD} + \overset{(4)}{\bar{A}BC\bar{D}} + \overset{(3)}{\bar{A}BCD} \\ + \overset{(4)}{\bar{A}\bar{B}C\bar{D}} + \overset{(1)}{A\bar{B}C\bar{D}}$$

$$= A\bar{B}\bar{D}(\bar{C}+C) + A\bar{C}D(\bar{B}+B) + BCD(\bar{A}+A) + \bar{A}C\bar{D}(\bar{B}+B)$$

$$= A\bar{B}\bar{D} + A\bar{C}D + BCD + \bar{A}C\bar{D}. \quad \checkmark$$

(iii)

$$\begin{array}{r} \overset{1}{0} \overset{0}{0} \overset{0}{0} \overset{1}{1} \quad \overset{1}{0} \overset{0}{0} \overset{1}{1} \quad \overset{0}{0} \overset{1}{0} \overset{0}{0} \quad \overset{0}{0} \overset{0}{1} \overset{0}{0} \\ + \quad 1101 \quad 1001 \quad 0101 \quad 0111 \\ \hline 0110 \quad 0011 \quad 1001 \quad 1001 \end{array}$$

(2) a)  $0x6399$

(2) b) Overflow occurred! This is a Signed sum,  
So inspect the V-flag:  $V = C_{15} \oplus C_{16} = 0 \oplus 1 = 1$ .

iv)

(2) Advantage: The final result is available faster.

Disadvantage: Requires many more logic gates.

## Question 2

③ i)

C	S	R	Q	/Q		Q - next output	q - current output
0	X	X	q	/q	(no change)		
1	0	0	q	/q	(no change)		
1	0	1	0	1	(reset)		
1	1	0	1	0	(set)		
1	1	1	1	1	(undefined state)		

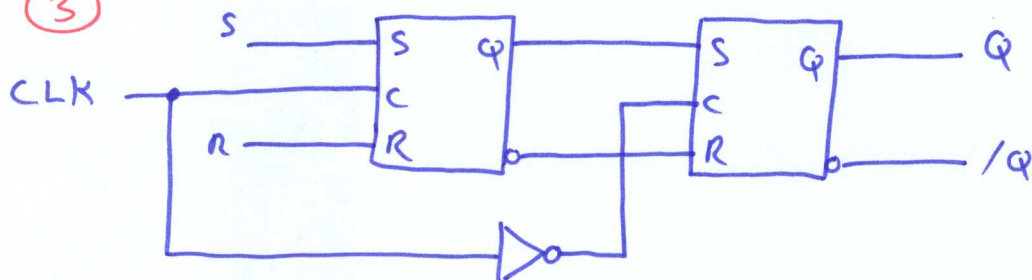
③ ii)

Two possibilities: 1. Initially  $c=s=r=1$  and  $c$  transits from  $1 \rightarrow 0$ . or: 2. Initially  $c=s=r=1$  and both  $s$  and  $r$  transit  $1 \rightarrow 0$  simultaneously.

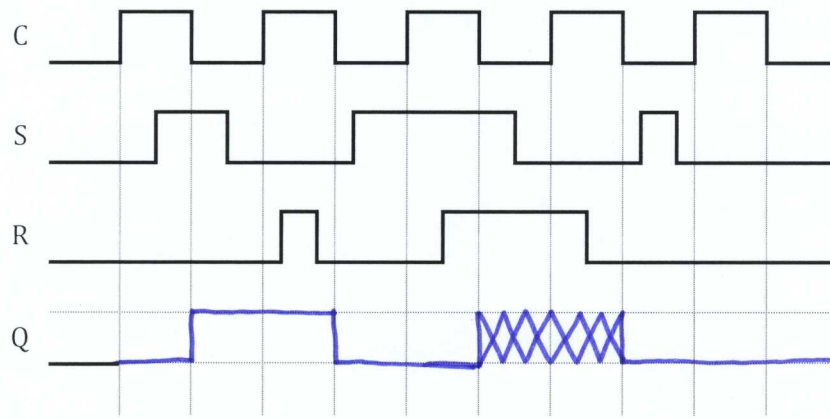
This will cause oscillation behavior at the outputs unless one of the feedback paths is longer than the other. In which case, the outputs will settle to one of the possible states but we cannot know which one (unless we know which is the longer path).

iii)

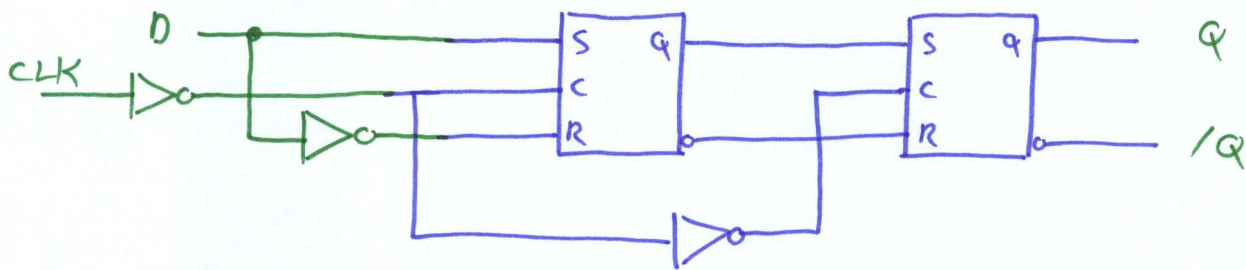
③



(5) iv)



(3) v)



(3) vi)

In the D Flip-Flop, S and R will never be 1 at the same time. Therefore, a race-condition can never occur and we will always be in a known state.



### Question 3

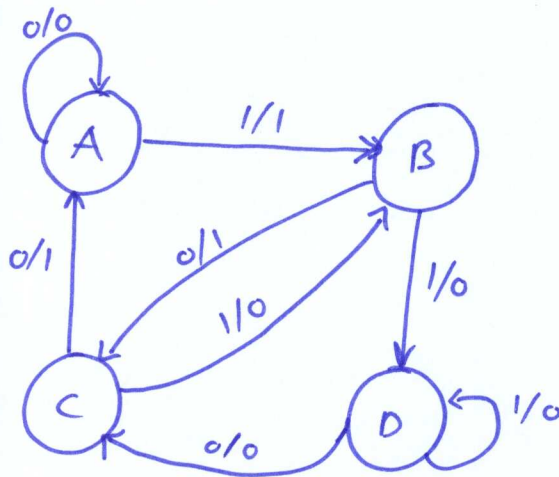
⑧ i) Define states:

A - Last two bits Seen were 00.

B - Last two bits Seen were 01.

C - Last two bits Seen were 10.

D - Last two bits Seen were 11.



This is  
a Mealy  
machine.

④ ii)

X	0	1	0	0	1	0	1	1	0	0	1	0	
State	A	A	B	C	A	B	C	B	D	C	A	B	C
Z	0	1	1	1	1	1	0	0	0	1	1	1	

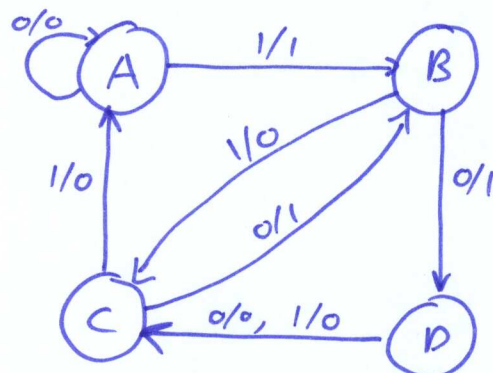
⑧ iii) Define states:

A - Last two outputs were 00.

B - Last two outputs were 01.

C - Last two outputs were 10.

D - Last two outputs were 11.



## Question 4

i)

$$F = \overline{AB + BC + AC}$$

③ a)

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

③ b)

PDN:  $F = \overline{AB + BC + AC}$

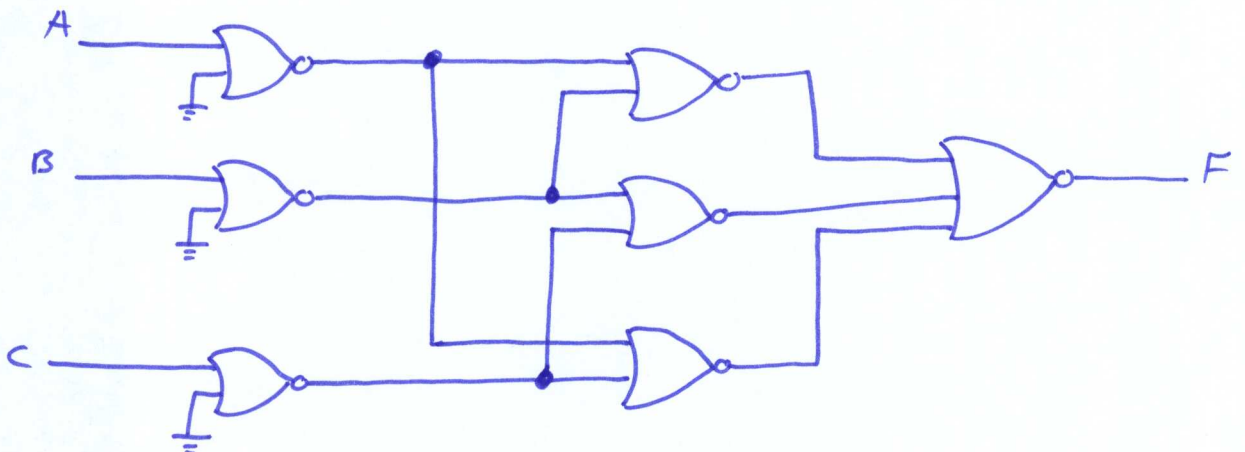
PUN:  $F = (\bar{A} + \bar{B})(\bar{B} + \bar{C})(\bar{A} + \bar{C})$

⇓

$$PDN = \overline{AB + BC + AC} = \overline{AB} \cdot \overline{BC} \cdot \overline{AC}$$

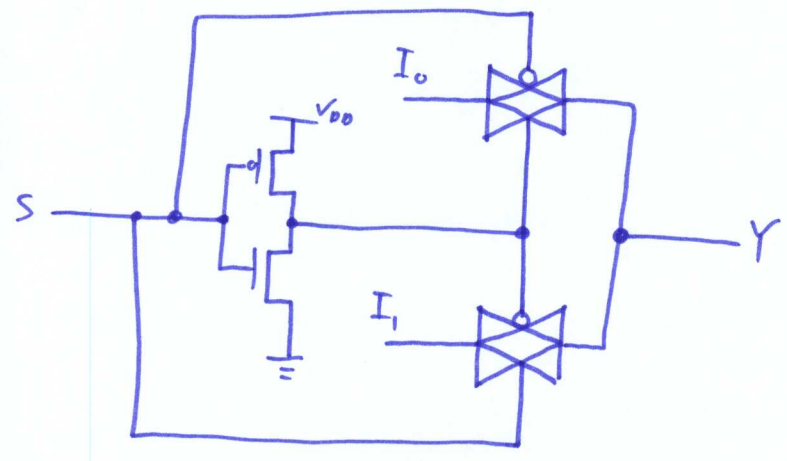
$$= (\bar{A} + \bar{B})(\bar{B} + \bar{C})(\bar{A} + \bar{C}) = PUN.$$

⑥ c)



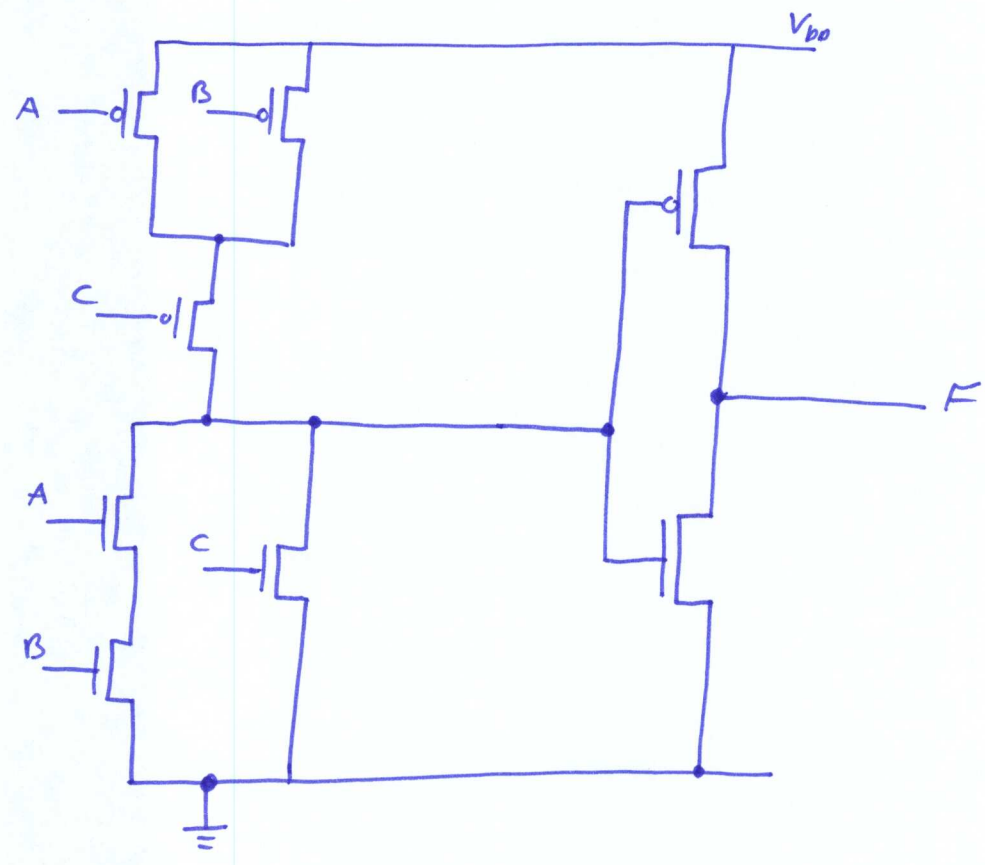
4 (ii)

2-to-1-Line Multiplexer:



4 (iii)

$$F = AB + C$$



## Question 5

① i)

The code implements a Moore machine.

④ ii)

current state	next state		output z
	x=0	x=1	
A	A	B	0
B	E	D	1
C	D	B	0
D	C	B	0
E	B	F	1
F	E	C	0

⑥ iii)

B	X				
C	A~D B~B	X			
D	A~C B~B	X	C~D B~B		
E	X	B~E D~F	X	X	
F	A~E B~C	X	D~E B~C	C~E B~C	X
	A	B	C	D	E

New states:

$A \sim C \sim D$

B

E

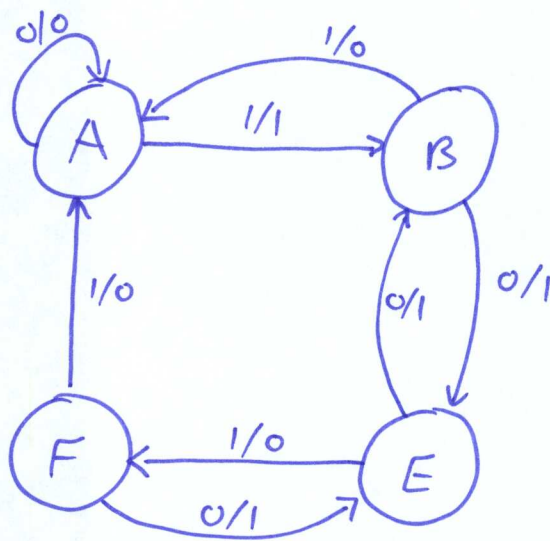
F

current state	next state		output z
	x=0	x=1	
A	A	B	0
B	E	A	1
E	B	F	1
F	E	A	0



⑥ iv) Mealy Machine:

current state	next state, output	
	x=0	x=1
A	A, 0	B, 1
B	E, 1	A, 0
E	B, 1	F, 0
F	E, 1	A, 0



③ v)

X	1	0	1	0	1	0	0	1	0	0
Z	1	1	0	1	0	1	1	0	0	0