

# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
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Computational golden era of:

logic & type rich  
computation

Statistical  
computation

# Computational golden era of:

logic & type rich  
computation

Expressive type systems:

Haskell, OCaml, Idris

Mechanised mathematics:

Agda, Coq, Isabelle/Hol, Lean

Verification:

SMT-powered, realistic  
systems

Statistical  
computation

generative modelling  
+

efficient inference:

Monte-Carlo simulation  
or gradient-based  
optimisation

"AI"

Computational golden era of:

logic & type rich  
computation

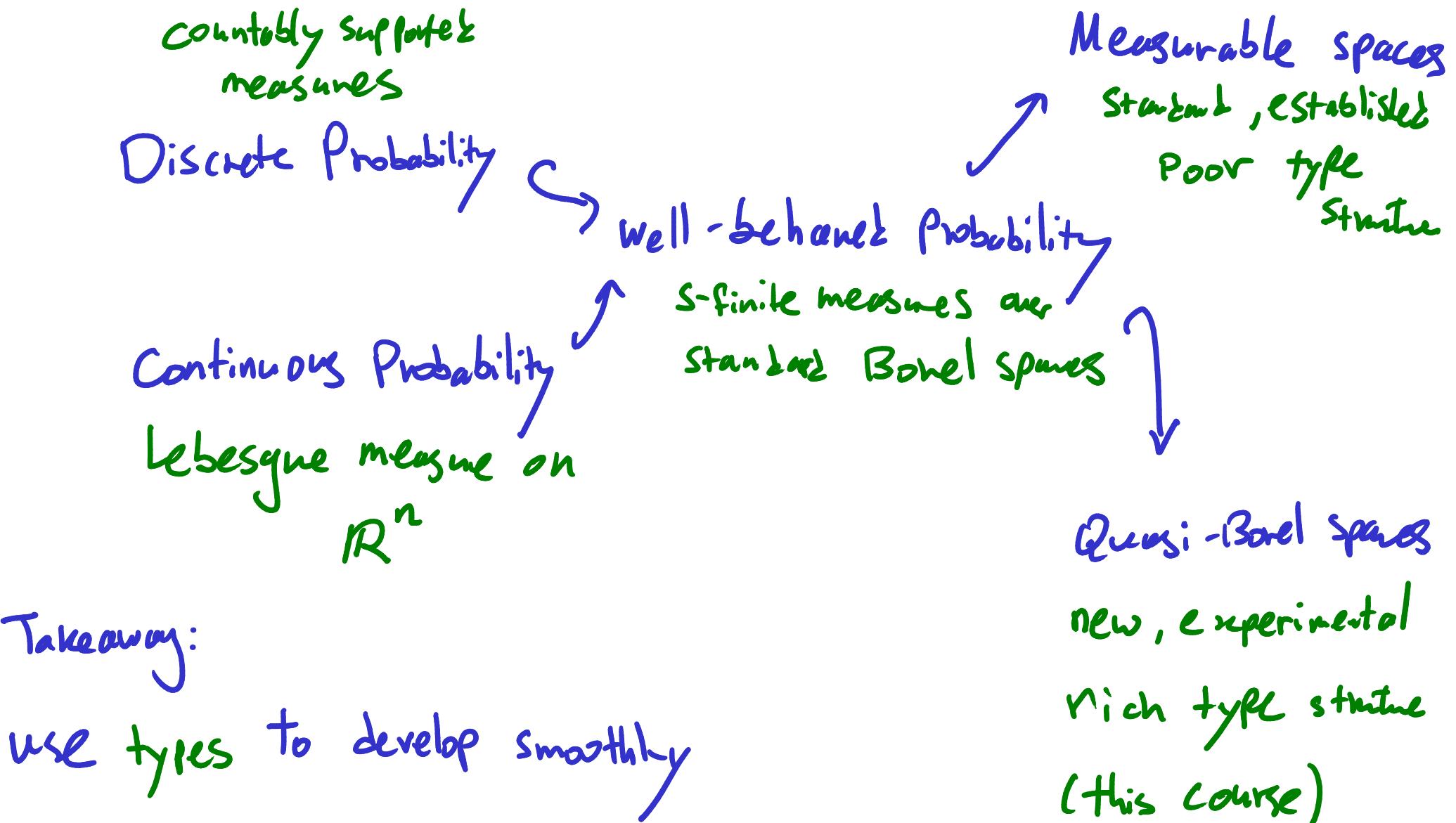
Statistical  
computation

Clear connection to

Foundations:

- Ralf's
- John's courses
- Michael's
- Dominik's
- this course

# Why foundations?



Plan:

- 1) type-driven Probability: discrete case (Mon + Tue (?) )
- 2) Borel sets & measurable spaces (Tue)
- 3) Quasi Borel spaces, Simple type structure (Wed)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

Please ask questions!

Smibble



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# Language of distribution & Probability

$X$  type (=space) of values / outcomes

$\mathcal{D}X$  type of distributions / measures over  $X$

$\mathcal{P}X \subseteq \mathcal{D}X$  sub type of probability measures (total measure)

$\mathcal{B}X$  type of measurable events - Subsets of  $X$  we wish to measure

$\mathbb{W}$  type of weights :  $[0, \infty]$

→ type judgment

$\mu : \mathcal{D}X, E : \mathcal{B}X \vdash c_e[E] : \mathbb{W}$

↳ measure  $\mu$  assigns to  $E$

# Axioms for measures

---

Empty event :  $\emptyset : \mathcal{B}X$

Its measure is  $0 : \mathbb{W}$  :

$$\mu : \mathcal{D}X \vdash \underset{\mu}{\text{Ce}}[\emptyset] = 0 : \mathbb{W}$$

# Axioms for measures

---

$BX$  is a Boolean Sub-algebra:

$$E : BX \vdash E^c : BX$$

$$E, F : BX \vdash E \cup F, E \cap F : BX$$

$$E, C : BX, \mu : DX \vdash \quad (\text{disjoint additivity})$$

$$\underset{\mu}{\text{Ce}}[E] = \underset{\mu}{\text{Ce}}[E \cap C] + \underset{\mu}{\text{Ce}}[E \cap C^c] : W$$

# Axioms for measures

$\omega := (\mathbb{N}, \leq)$      $(B, \subseteq)$      $(W, \leq)$     posets

$$(BX, \subseteq)^\omega := \left\{ (E_n)_{n \in \mathbb{N}} \in (BX)^\mathbb{N} \mid E_0 \subseteq E_1 \subseteq E_2 \subseteq \dots \right\}$$

$(BX, \subseteq)$  and  $(W, \leq)$  are  $\omega$ -chain-closed:

$$E_- : (BX, \subseteq)^\omega \vdash \bigvee_n E_n : BX \quad \alpha_- : (W, \leq)^\omega \vdash \sup_n \alpha_n : W$$

$$E_- : (BX, \subseteq)^\omega, \mu : D_X \vdash \quad \text{(Scott Continuity)}$$

$$\text{Ce}_{\mu} \left[ \bigvee_n E_n \right] = \sup_n \text{Ce}_{\mu} [E_n] : W$$

# Axiom for Probability

$$\text{Cast} : \text{PX} \xleftarrow{\leq} \text{DX}$$

$$1 : \mathbb{W}$$

$$\mu : \text{PX} \vdash \text{Ce}[X] = 1 : \mathbb{W}$$

*Cast  $\mu$*

Avoid casting:

$$E : BX, \mu : \text{PX} \vdash \Pr_{\Gamma}[E] := \text{Ce}[E] : [0,1] \subseteq \mathbb{W}$$

*Cast  $\mu$*

# Axioms for measures

---

Integration:

$$\mu : \mathbf{DX}, \varphi : \mathbb{W}^X \vdash \int_\mu \varphi : \mathbb{W} \quad (\text{Lebesgue integral})$$

Again, avoid casting:

$$\mu : \mathbf{PX}, \varphi : \mathbb{W}^X \vdash \mathbb{E}_{\mu}[\varphi] := \int_{\mu} (\text{cast } \mu) \varphi : \mathbb{W} \quad (\text{Expectation})$$

More structure & notation later (...technical...)

Have: language + axioms

Want: model

today: discrete measures

rest of course: discrete + continuous

## Discrete model

type  $X$  : set

$D_X := \{ \mu : X \rightarrow \mathbb{W} \mid \mu \text{ is Countably Supported} \}$   
(next slide)

Support

Power set

$\mu : \mathbb{W}^X, S : \mathcal{P}X \vdash S \text{ supports } \mu :=$

$\forall x : X. \mu x > 0 \Rightarrow x \in S : \text{Prop}$

$\mu : \mathbb{W}^X \vdash \text{Supp } \mu := \{x \in X \mid \mu x > 0\} : \mathcal{P}X$

$\text{Supp } \mu$  is the smallest set supporting  $\mu$

## Discrete model

type  $X$  : set

$$DX := \{ \mu : X \rightarrow \mathbb{W} \mid \mu \text{ is Countably Supported} \}$$

$$:= \{ \mu : X \rightarrow \mathbb{W} \mid \text{Supp } \mu \text{ is Countable} \}$$

## Ex. measures

- $X$  ctbl. Counting measure  $\#_X : DX$   
 $\#_X := \lambda x : X. 1$  (NB:  $\text{Supp } \#_X = X \sqrt{\text{ctbl}}$ )
- Dirac measure:  
 $\sigma : X \vdash \delta_x := \lambda x'. \begin{cases} x = x' : 1 \\ \text{o.w.} : 0 \end{cases} : DX$   
NB:  $\text{Supp } \delta_x = \{x\} \sqrt{\text{ctbl}}$
- Zero measure  $\underline{0} := \lambda x. 0 : DX$   
NB:  $\text{Supp } \underline{0} = \emptyset \sqrt{\text{ctbl}}$

## Discrete model

type  $X$ : set

$DX := \{ \mu : X \rightarrow \mathbb{W} \mid \mu \text{ is Countably Supported} \}$

$$\mu : DX, E : BX \vdash C_E[\mu] := \sum_{x \in E} \mu x$$

$$:= \sum_{x \in E \cap \text{Supp } \mu} \mu x$$

Lemma:  $\mu : DX, S \in \mathcal{P}_{\text{ctbl}}^X, S \text{ supports } \mu, E : BX \vdash$

$$C_E[\mu] = \sum_{x \in E \cap S} \mu x$$

Ex:

- $E : B X \vdash$   $C_e[E] = |\underset{\#_x}{E}| := \begin{cases} E \text{ has } n \text{ elements: } n \\ E \text{ infinite: } \infty \end{cases}$

- $E : B X, n : X \vdash C_e[E] = \underset{\delta_n}{\delta_{\in E}} = \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases} =: [x \in E] : \mathbb{W}$

NB:  $E : B X \vdash [- \in E] : X \rightarrow \mathbb{W}$

indicator  
function

- $E : B X \vdash C_e[E] = \underset{\Omega}{\Omega}$

## Validate axioms

$$\mu : \text{DX} \vdash \underset{\mu}{\text{Ce}}[\emptyset] = 0 : \mathbb{W}$$

$$E, C : \text{BX}, \mu : \text{DX} \vdash$$

$$\underset{\mu}{\text{Ce}}[E] = \underset{\mu}{\text{Ce}}[E \cap C] + \underset{\mu}{\text{Ce}}[E \cap C^c] : \mathbb{W}$$

$$E_- : (\text{BX}, \subseteq)^\omega, \mu : \text{DX} \vdash$$

$$\underset{\mu}{\text{Ce}}[\bigvee E_n] = \sup_n \underset{\mu}{\text{Ce}}[E_n] : \mathbb{W}$$

Kernels  $\kappa$  from  $\Gamma$  to  $X$ :

$$\kappa : (DX)^\Gamma$$

kernels are "open/parameterised" measures

Ex: Dirac kernel:  $\delta_+ : (DX)^X$

## Kock Integral

$$\mu : D\Gamma, \kappa : DX \vdash \int^\Gamma \mu \kappa : DX$$

In discrete model:

$$\int^\Gamma \mu \kappa := \lambda x : X. \sum_{n \in \Gamma} \underbrace{\mu n \cdot k(n; x)}_{:= h \vdash x}$$

## (Weak) disintegration problem:

Input:  $\mu: D\Gamma$      $V: DX$

Output: a kernel  $k:(DX)^{\Gamma}$  s.t.

$$\oint \mu k = V$$

Call such  $k \stackrel{a}{=} (\text{weak}) \text{ disintegration of } V$

w.r.t.  $\mu$ .

(non-standard  
terminology)

Ex disintegration:

$$\underline{n} := \{0, 1, 2, \dots, n-1\}$$

disintegrate  $\#_{\geq \frac{n+1}{2}}$  w.r.t.  $\#_{\geq}$

$$k: \left(D(\#_{\geq \frac{n+1}{2}})\right)^2$$
$$k(x; f) := \begin{cases} f(n) = x & : 1 \\ \text{o.w.} & : 0 \end{cases}$$

$$\left(\#_{\geq} k\right) f = \sum_{x \in \underline{n}} \#_{\geq}^1 x \cdot k(x; f)$$

NB:  $\text{Supp}(k)$   
 $\sqrt{c+b}$

$$= k(0; f) + k(1; f) = \#_{\geq \frac{n+1}{2}}(f) = 1$$

# Probability measures

$$P_X := \left\{ \mu : D_X \mid \underset{\mu}{\text{C}_e}[X] = 1 \right\} \hookrightarrow^{\subseteq} D_X$$

Lemma:  $\delta_- : X \rightarrow D_X$  and  $\oint : D\Gamma \times (D_X)^r \rightarrow D_X$

lift along the inclusion      cast:  $P \hookrightarrow^{\subseteq} D$ :

$$\begin{array}{ccc} X & \xrightarrow{\delta_-} & P_X \\ & \dashv & \downarrow \text{cast} \\ & \xrightarrow{\delta_-} & D_X \end{array}$$

$$\begin{array}{ccc} P\Gamma \times (P_X)^r & \dashv \oint \dashrightarrow & P_X \\ \text{cast} \times (\text{cast}) \downarrow & & \downarrow \text{cast} \\ D\Gamma \times (D_X)^r & \xrightarrow{\oint} & D_X \end{array}$$

Prop (discrete Giry):

(Michèle Giry '82)

$(P, \delta_-, \oint)$  is a monad i.e.

$$m : \Gamma, n : (Dx)^\Gamma \vdash \oint \delta_n k = k \ r$$

$$\mu : D X \vdash \oint \mu(\lambda x) \delta_x = \mu : D X$$

$$\mu : D\Gamma, \kappa : (Dx)^\Gamma, t : (DY)^X \vdash$$

$$\oint \mu(\lambda x) \left( \oint (\kappa r) t \right) = \oint \left( \oint \mu \kappa \right) (\lambda x) t(x)$$

Corollary:  $(P, \delta_-, \oint)$  is a monad.

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$\mathbb{W}$  type of weights :  $[0, \infty]$

→ type judgment

$\mu : \mathcal{D}X, E : \mathcal{B}X \vdash c_e[E] : \mathbb{W}$

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# Language of distribution & Probability

Recap

$X$  type (=space) of values / outcomes

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$\mathbb{W}$  type of weights :  $[0, \infty]$

→ type judgment

$\mu : \mathcal{D}X, E : \mathcal{B}X \vdash c_e[E] : \mathbb{W}$

↳ measure  $\mu$  assigns to  $E$

## Axioms for measures/distributions

Recap

$$\mu : \mathbf{D}X \vdash \underset{\mu}{\text{Ce}}[\emptyset] = 0 : \mathbb{W}$$

$$E, C : \mathbf{B}X, \mu : \mathbf{D}X \vdash$$

$$\underset{\mu}{\text{Ce}}[E] = \underset{\mu}{\text{Ce}}[E \cap C] + \underset{\mu}{\text{Ce}}[E \cap C^c] : \mathbb{W}$$

$$E_- : (\mathbf{B}X, \subseteq)^\omega, \mu : \mathbf{D}X \vdash$$

$$\underset{\mu}{\text{Ce}}\left[\bigvee_n E_n\right] = \sup_n \underset{\mu}{\text{Ce}}[E_n] : \mathbb{W}$$

# Kernels & their Koch integral

Recap

kernel from  $\Gamma$  to  $X$ :  $k: (DX)^\Gamma$  or  $k: \Gamma \rightarrow DX$

Dirac kernel:  $\delta_- : X \rightarrow DX$

Koch integral:  $\mu: D\Gamma$ ,  $k: (DX)^\Gamma \vdash \oint \mu k : DX$   
or  $\oint \mu(dx) \kappa(x)$  (*dx binding occurs in  $\kappa(x)$* )

Giry monads:  $(D, \delta_-, \oint) \dashv (P, \mathcal{S}_-, \oint)$ .

## Discrete model

Recap

$$\text{type} : \text{set} \quad W := [0, \infty] \quad \mathcal{B}X := P_X$$

$$DX := \{\mu : X \rightarrow W \mid \text{Supp } \mu \text{ countable}\}$$

$$P_X := \left\{ \mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1 \right\}$$

$$\underset{\mu}{\text{Ce}}[E] := \sum_{x \in E} \mu_x \quad \delta_x := \lambda x'. \begin{cases} x = x': 0 \\ x \neq x': 1 \end{cases}$$

$$\phi \mu k := \lambda x. \sum_{m \in \Gamma} \mu^m \cdot k(m; x)$$

## Ex distributions

Recap

Counting measure ( $\lambda_{ctbl}$ ):  $\#_X := \lambda_X \cdot 1$

Dirac measure  $\delta_x$  (prev slide)

Zero measure  $\underline{\varnothing} := \lambda_X \cdot 0$

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## Product measures

$$\mu: D X, \nu: D Y \vdash \mu \otimes \nu := \int \mu(dx) \int \nu(dy) \delta_{(x,y)} : D(X \times Y)$$

( $\otimes$  lifts along  $P \hookrightarrow D$ )

$$= \lambda(x,y). \mu x \cdot \nu y$$

discrete model

$$E_{\#} : \#_{X \times Y} = \#_X \otimes \#_Y$$

Indeed:

$$(\# \otimes \#)(x,y) = \#x \cdot \#y = 1 \cdot 1 = 1 = \#(x,y)$$

build measures  
compositionally

$$\text{Notation: } \lambda : D(X \times Y), \kappa : (DZ)^{X \times Y} \vdash \oint \lambda(\Delta x, \Delta y) \kappa(x, y) \\ := \oint \lambda \kappa$$

Fubini - Tonelli Thm:

---

Integrate in any order:

$$\mu : DX, \nu : DY, \kappa : (DZ)^{X \times Y} \vdash$$

$$\oint \mu(dx) \oint \nu(dy) \kappa(x, y) = \oint (\mu \otimes \nu)(dx, dy) \\ = \oint \nu(dy) \oint \mu(dx) \kappa(x, y)$$

## Pushing a measure forward

$$\mu: D_{\Omega}, d: X^{\Omega} \vdash \mu_f := \phi \mu(d\omega) \delta_{\alpha\omega} : DX$$

$$= \lambda x. \sum_{\omega \in \Omega} \mu \omega$$

$$\alpha\omega = x$$

$\alpha: X^{\Omega}$ : random element

(w.r.t.  $\mu$ )

$\mu_{\alpha}: DX$ : the law of  $\alpha$

Ex: We can represent configurations of 2 dice using  $\underline{6} \times \underline{6}$

Letting  $(+): \underline{6}^2 \rightarrow \mathbb{N}^2 \xrightarrow{(+)} \mathbb{N}$

we have that the law of  $(+)$ :

$$(\#_{\underline{6}} \otimes \#_{\underline{6}})_{(+)} : \mathbb{D}/\mathbb{N}$$

is the number of rolls whose sum is given

build measures  
compositionally

## Scaling a measure

$$(\cdot) : \mathbb{W} \times D_X \longrightarrow D_X$$

$$a \cdot \mu := \lambda x. a \cdot \mu x$$

$$\boxed{NB: \text{Supp}(a \cdot \mu) = \begin{cases} a=0: \emptyset \\ a \neq 0: \text{Supp } \mu \end{cases}}$$

$\checkmark_{c+61}$

$(\cdot) : \mathbb{W} \times D_X \rightarrow D_X$  is an action of monoid  $(\mathbb{W}, (\cdot), 1)$  on  $D_X$ :

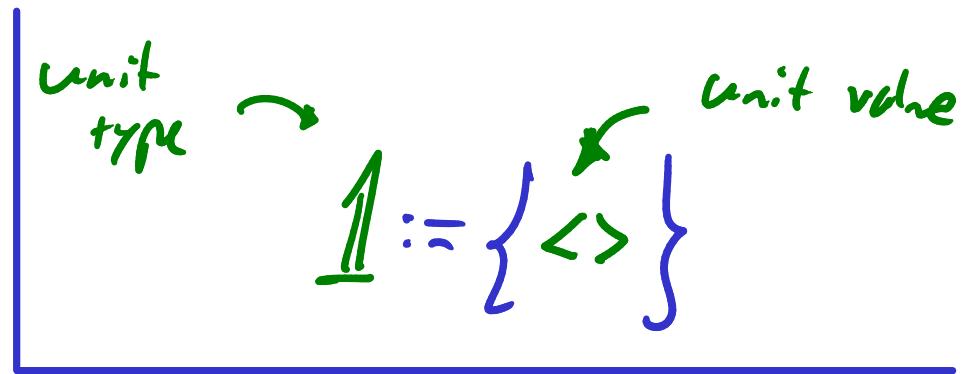
$$\mu : D_X \vdash$$

$$1 \cdot \mu = \mu$$

$$a, b : \mathbb{W}, \mu : D_X \vdash$$

$$a \cdot (b \cdot \mu) = (a \cdot b) \cdot \mu$$

## Normalisation



$\mu : D X, C_C[X] \neq 0, \infty +$

$$\|\mu\| := \left( \frac{1}{C_C[X]} \right) \cdot \mu : P X$$

Ex:

$$\emptyset \neq A \subseteq_{fin} X : U_{A \subseteq X} := \|\#_A\|_{A \subseteq X} : P X$$

$$1 \xrightarrow{\#_A} D A \xrightarrow{(-)_{A \subseteq X}} D X \xrightarrow{\|\cdot\|} P X$$

I.e.

$$U_{A \subseteq X} := \lambda n. \begin{cases} n \in A : \frac{1}{|A|} \\ n \notin A : 0 \end{cases}$$

so

$$\bigcup_{n \in A} = \delta_n$$

## Standard vocabulary

Joint distributions:

$$\mu : D(X_1 \times X_2)$$

Marginal distribution:

$$X_1 \xleftarrow{\pi_1} X_1 \times X_2 \xrightarrow{\pi_2} X_2$$

law of projection

$$\mu_{\pi_i} : D X_i$$

Marginalisation:  $\mu_{\pi_i} = \iint \mu(dx, dy) S_x$

integrate out  $y$

Exercise:  $\mu : P X, V : D x \vdash (\mu \otimes V)_{\pi_2} = V$

## independence

Pairing R.E.S:

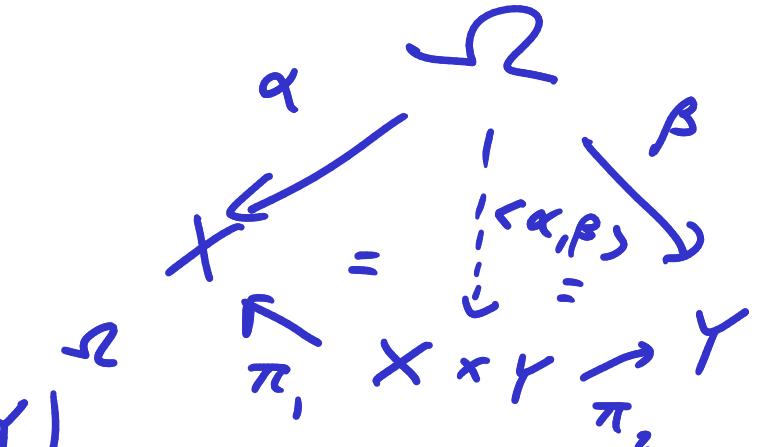
$$\alpha : X^{\Omega}, \beta : Y^{\Omega} \vdash$$

$$\langle \alpha, \beta \rangle := \lambda w. \langle \alpha w, \beta w \rangle : (X \times Y)^{\Omega}$$

$$\lambda : D\Omega, \alpha : X^{\Omega}, \beta : Y^{\Omega} \vdash \alpha \perp \beta := \lambda_{\langle \alpha, \beta \rangle} = \lambda_{\alpha} \oplus \lambda_{\beta}$$

: Prop

$\alpha, \beta$  independent w.r.t.  $\lambda$



Ex<sup>(Durrett)</sup> represent Outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda : \bigcup_{c \in C} \bigcup_{c \in C} \bigcup_{c \in C} : P_\Omega$$

$$\pi_i : \Omega \rightarrow C \quad \text{Outcome of } i^{\text{th}} \text{ toss}$$

$$\text{Same}_{ij} : \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{? (=)} \mathbb{B}$$

where :  $(?) : C^2 \rightarrow \mathbb{B} := \{\text{True}, \text{False}\}$

$$?_{x=y} := \begin{cases} x=y : \text{True} \\ x \neq y : \text{False} \end{cases}$$

Ex <sup>(Durrett)</sup> represent outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda : U_C \otimes U_C \otimes U_C : P_{\Omega}$$

$$\pi_i : \Omega \rightarrow C \quad \text{Outcome of } i^{\text{th}} \text{ toss}$$

$$\text{Same}_{ij} : \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{?} B$$

marginalisation

$$\lambda_{\text{Same}_{12}}^T = (U_C \otimes U_C)^T \stackrel{?}{=} \frac{1}{4} + \frac{1}{4} = \frac{1}{2}$$

$$\begin{matrix} U_C(T) \cdot U_C(T) \\ \downarrow \\ \frac{1}{4} \\ \uparrow \\ U_C(H) \cdot U_C(H) \end{matrix}$$

$$\text{so } \lambda_{\text{Same}_{12}}^F = \frac{1}{2} \text{ too}$$

Ex <sup>(Durrett)</sup> represent Outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda : \bigcup_{C^3} \otimes \bigcup_{C^3} \otimes \bigcup_{C^3} : P_\Omega$$

$$\pi_i : \Omega \rightarrow C \quad \text{Outcome of } i^{\text{th}} \text{ toss}$$

$$\underline{i \neq j} : \lambda_{\text{Same}_{ij}} = V_{\mathbb{B}}$$

$$\text{Same}_{ij} : \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{?} \mathbb{B}$$

$$\lambda : \begin{matrix} (T, T) \mapsto \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \\ \langle \text{Same}_{12}, \text{Same}_{23} \rangle \end{matrix} \hookrightarrow \lambda(H, H, H) \quad \hookrightarrow \lambda(T, T, T)$$

$$(T, F) \mapsto \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4}$$

$$\hookrightarrow \lambda(H, H, T) \quad \hookrightarrow \lambda(T, T, H)$$

Ex<sup>(Durrett)</sup> represent Outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega := C \times C \times C \quad \lambda : U_C \otimes U_C \otimes U_C : P_{\Omega}$$

$$\pi_i : \Omega \rightarrow C \quad \text{Outcome of } i^{\text{th}} \text{ toss}$$

$$\underline{i \neq j} \quad \lambda_{\text{Same}_{ij}} = V_{IB}$$

$$\text{Same}_{ij} : \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{?} IB$$

$$\lambda_{\langle \text{Same}_{12}, \text{Same}_{23} \rangle} = V_{IB \times IB} = V_{IB} \otimes V_{IB} = \lambda_{\text{Same}_{12}} \otimes \lambda_{\text{Same}_{13}}$$

$$\text{So } \text{Same}_{12} \perp \lambda \text{ Same}_{13}$$

## independence

Pairing R.E.S:

$$\alpha : X^{\Omega}, \beta : Y^{\Omega} \vdash$$

$$\langle \alpha, \beta \rangle := \lambda w. \langle \alpha w, \beta w \rangle : (X + Y)^{\Omega}$$

$$\lambda : D\Omega, \alpha : X^{\Omega}, \beta : Y^{\Omega} \vdash \alpha \perp_{\lambda} \beta := \lambda_{\langle \alpha, \beta \rangle} = \lambda_{\alpha} \otimes \lambda_{\beta} : \text{Prop}$$

$\alpha, \beta$  independent w.r.t.  $\lambda$

I-ary version:

$$\lambda : D\Omega, \alpha_i : \prod_{i \in I} X_i^{\Omega} \vdash \perp_{\lambda, i \in I}^{\alpha_i} :=$$

$\alpha_i$  independent  
w.r.t.  $\lambda$

$$\forall J \subseteq_{\text{fin}} I. \quad \lambda_{\langle \alpha_j \rangle_{j \in J}} = \bigotimes_{j \in J} \lambda_{\alpha_j} : \text{Prop}$$

Ex <sup>(Durrett)</sup> represent outcomes of 3 coin tosses:

$$C := \{T, H\} \quad \Omega = C \times C \times C \quad \lambda : \bigcup_{C \times C \times C} \times P_{\Omega}$$

$$\pi_i : \Omega \rightarrow C \quad \text{Outcome of } i^{\text{th}} \text{ toss}$$

$$\underline{i \neq j} : \lambda_{\text{Same}_{ij}} = V_{\mathbb{B}}$$

$$\text{Same}_{ij} : \Omega \xrightarrow{\langle \pi_i, \pi_j \rangle} C \times C \xrightarrow{?} \mathbb{B}$$

$$\begin{matrix} i \neq j \\ * \\ n \end{matrix} : \text{Same}_{ij} \perp \text{Same}_{jk}$$

$$\frac{1}{\lambda} \left\{ \text{Same}_{12}, \text{Same}_{23}, \text{Same}_{13} \right\}$$

$$\text{Intuition: Same}_{13} = \text{IFF} (\text{Same}_{12}, \text{Same}_{23})$$

Calc:

$$\begin{aligned} \lambda_{\langle \text{Same}_{12}, \text{Same}_{23}, \text{Same}_{13} \rangle} (T, T, T) &= \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \neq \frac{1}{2^3} = \lambda_{\text{Same}_{12}} \odot \lambda_{\text{Same}_{23}} \odot \lambda_{\text{Same}_{13}} \\ &\hookrightarrow \lambda(H, H, H) \quad \hookrightarrow \lambda(T, T, T) \end{aligned}$$

## Vocabulary

(Discrete) Measure Space  $(X, \mu : D_X)$

measure preserving  $f : (X, \mu) \rightarrow (Y, \nu)$

function  $f : X \rightarrow Y$  s.t.  $\mu_f = \nu$

$\mu : D_X$ ,  $f : X \rightarrow Y \vdash \mu$  invariant under  $f :=$

$f : (X, \mu) \rightarrow (Y, \nu)$

Ex:

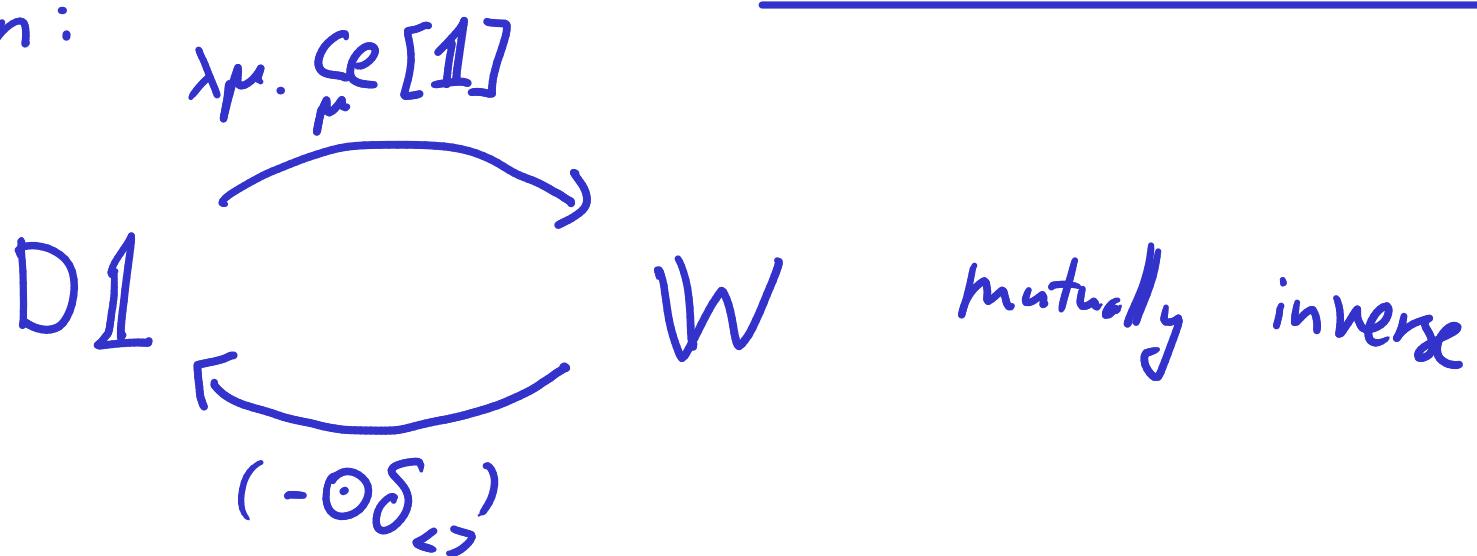
$\mu : D_X, \nu : D_Y \vdash$

Swap :  $(X \times Y, \mu \otimes \nu) \longrightarrow (Y \times X, \nu \otimes \mu)$  so

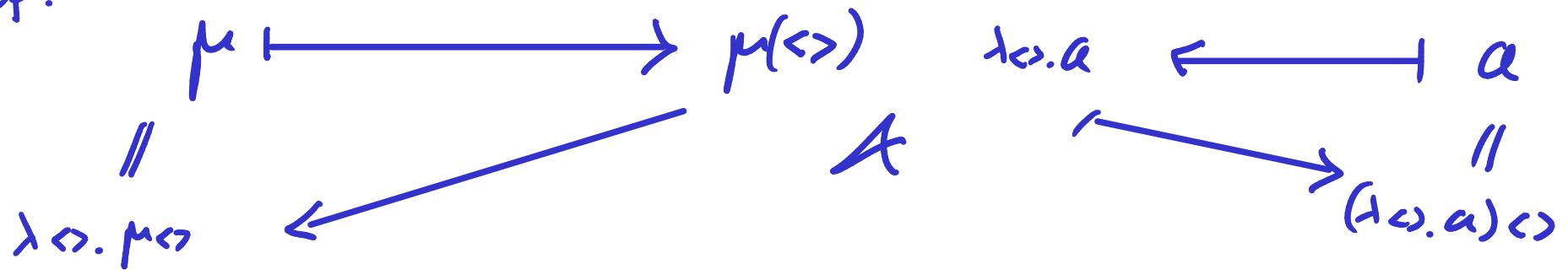
$\mu : D_X \vdash \mu \otimes \mu$  invariant under Swap

## Weights as measures

Observation:



Proof:



□

NB: unit type  $\rightarrow$   $\mathbf{1} := \{\langle\rangle\}$  unit value

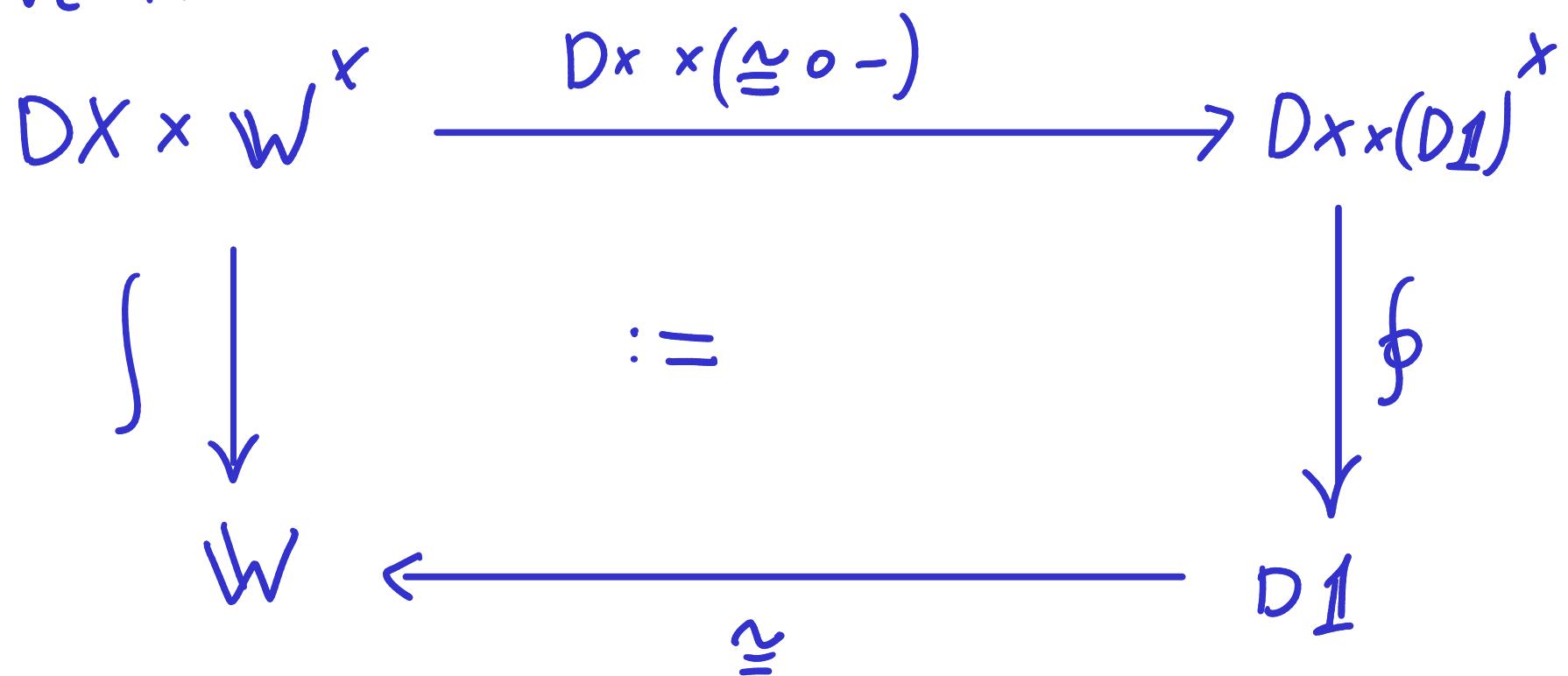
# Integration

$$\mu: \mathcal{D}X, \varphi: \mathbb{W}^X \vdash \int^\mu \varphi : \mathbb{W}$$

$$:= \sum_{x \in X} \mu x \cdot \varphi x$$

(Lebesgue integral)

Can derive it:



## Additivity:

$$\text{I ctsl, } \mu_-(DX)^I \vdash \sum_{i \in I} \mu_i : DX$$

$$:= \lambda x. \sum_{i \in I} \mu_i x$$

NB:

$$\text{supp} \sum_i \mu_i \subseteq$$

$$\bigcup_i \text{supp } \mu_i$$

✓ctsll

Ex: Bernoulli distribution

$$p:[0,1] \vdash B(p) := p \cdot \delta_{\text{True}} + (1-p) \cdot \delta_{\text{False}} : P/B$$

$$\text{i.e. } \beta_p : \begin{aligned} \text{True} &\mapsto p \\ \text{False} &\mapsto 1-p \end{aligned}$$

Thm (affine-linearity):

$\phi$  is affine-linear in each argument:

$I \vdash b : I$

$$M : (\mathbf{D}\Gamma)^I, k : (\mathbf{D}x)^I \vdash \phi\left(\sum_{i \in I} a_i \cdot \mu_i\right) k = \sum_{i \in I} a_i \cdot \phi \mu_i k$$

$I \vdash b : I$ ,  $\mu : \mathbf{D}\Gamma$ ,  $a_i : W^I$ ,  $k_i : \mathbf{D}x^I$

$$\int \mu(dx) \left( \sum_{i \in I} a_i \cdot k_i(x) \right) = \sum_{i \in I} a_i \cdot \phi \mu k_i$$

Prop:  $\mathbb{W} \cong D1$  is a  $\sigma$ -semi-ring isomorphism:

$$(\mathbb{W}, \Sigma, (\cdot), 1) \cong (D1, \Sigma, (\cdot), \delta_{\leq})$$

and  $(\cdot) : \mathbb{W} \times Dx \rightarrow Dx$  makes  $Dx$  into a module:

$$\left( \sum_{i \in I} a_i \right) \cdot \mu = \sum_{i \in I} (a_i \cdot \mu) \quad a \cdot \sum_{i \in I} \mu_i = \sum_{i \in I} a \cdot \mu_i$$

Corollary:  $\int$  is affine-linear in each argument.

## Random variable :

NB:  $\bar{\mathbb{R}} := [-\infty, \infty]$

A random element  $\alpha: \bar{\mathbb{R}}^\Omega$  (wrt some  $\mu: D\mathcal{L}$ )

Can add, multiply r.v.'s.

To integrate r.v.'s:

$$(-)^+: \bar{\mathbb{R}}^\Omega \longrightarrow \mathbb{W}^\Omega$$

$$\alpha^+ := \lambda w. \begin{cases} \alpha \cdot w \geq 0 : \alpha w \\ 0.w : 0 \end{cases} = [\alpha \geq 0] \cdot |\alpha|$$

$$\alpha^- := \lambda w. \begin{cases} \alpha \cdot w \leq 0 : |\alpha w| \\ 0.w : 0 \end{cases} = [\alpha \leq 0] \cdot |\alpha|$$

So  $\alpha = \alpha^+ - \alpha^-$

$\mu: D\Omega, \alpha: \overline{\mathbb{R}}^n, \int \mu \alpha^+ < \infty \text{ or } \int \mu \alpha^- < \infty +$

$$\int \mu \alpha := \int \mu \alpha^+ - \int \mu \alpha^- : \overline{\mathbb{R}}$$

Ex. The (discrete) Lebesgue  $p$ -space:

$$p \in [1, \infty), \mu: P\Omega \vdash L_p(\Omega, \mu) :=$$

$$\left\{ \alpha: \overline{\mathbb{R}}^n \mid \underset{\mu}{\mathbb{E}}[|\alpha|^p] < \infty \right\}$$

$L_p(\Omega, \mu)$  has a norm  $\|\alpha\| := \sqrt[p]{\underset{\mu}{\mathbb{E}}[|\alpha|^p]}$  almost Banach

$L_2(\Omega, \mu)$  has an inner product  $\langle \alpha, \beta \rangle := \underset{\mu}{\mathbb{E}}[\alpha \cdot \beta]$  almost Hilbert

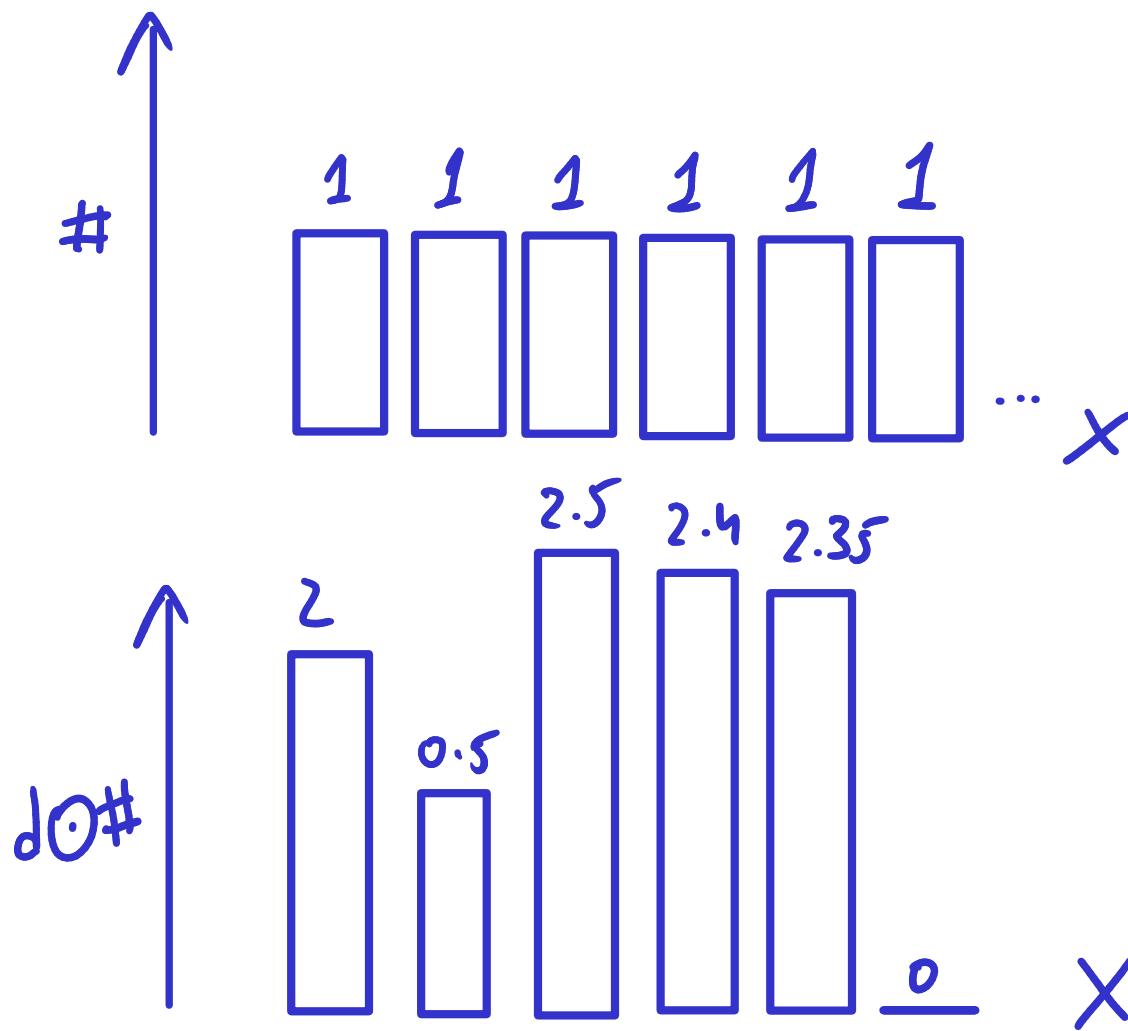
## Density

a density over  $X$ :  $d : X \rightarrow W$

$$d : W^X, \mu : D_X \vdash d \odot \mu : D_X \\ := \oint \mu(dx) (dx \cdot \delta_x)$$

Warning The types of measures & densities in the discrete model are close, but still different. They coincide on countable sets, so people often confuse them. Types help us keep them separate.

Intuition:



## Almost certain Properties

$E: BX, \mu: DX \vdash \mu(dx)$ -almost certainly  $x \in E : \text{Prop}$   
 $\quad := [x \in E] \odot \mu = \mu$

$$\text{NB: } [x \in E] = \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases} : \mathbb{W}$$

When  $\mu: Px$  we say instead

$\mu(dx)$ -almost surely  $x \in E$

Exercise Look up the def. of a normed space

and modify the definition so that  $L_p(\Omega, \mu)$  is a normed space up-to almost sure equality.

## Absolute continuity

$\mu, \nu : D^X, d : W^X \vdash d = \frac{d\mu}{d\nu} : \text{Prop}$

$\quad := \mu = d \odot \nu$

$\mu, \nu : D^X \vdash \mu \ll \nu := \mu \text{ is absolutely continuous w.r.t. } \nu : \text{Props}$

$\quad := \exists d : W^X. \quad d = \frac{d\mu}{d\nu}.$

$\quad =: \mu \text{ has a density w.r.t. } \nu$

Lemma:  $\mu, \nu : D^X,$   
 $\mu \ll \nu,$   
 $k : (D^Y)^X$

$$\oint V(dx) \frac{d\mu}{d\nu}(x) \cdot k_x = \oint \mu(dx) k_x$$

$$\underline{Ex}: \bigcup_{A \subseteq X} \ll (\#_A)_{\text{Cost}: A \subseteq X}$$

$$\frac{dV_{A \subseteq X}}{d(\#_A)_{\text{Cost}}} = \lambda x. \begin{cases} x \in A : & \frac{1}{|A|} \\ \text{D.W.} : & 0 \end{cases}$$

but also:

$$\frac{dV_{A \subseteq X}}{d(\#_A)_{\text{Cost}}} = \lambda x. \frac{1}{|A|}$$

Radon-Nikodym Thm: (discrete version)

$\mu, \nu : P X \vdash \mu \ll \nu$  iff  $\forall x. \nu x = 0 \Rightarrow \mu x = 0$   
i.e.  $\text{Supp } \mu \subseteq \text{Supp } \nu$

In that case, if  $d_1, d_2 = \frac{d\mu}{d\nu}$  then

$$\nu(dx)\text{-a.s. } d_1 x = d_2 x$$

Ex: for ctbl  $X$ ,  $\forall \mu : D X . \mu \ll \#_X$ . Proof: vacuously, as  $\#_X x \neq 0$ .

Then  $\lambda x. \mu x = \frac{d\mu}{d\#} .$

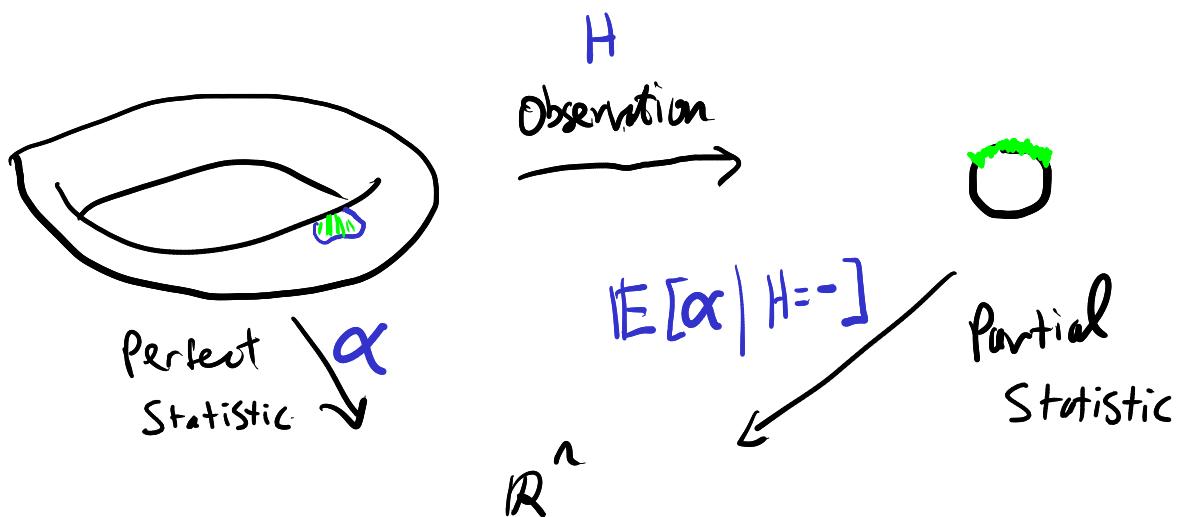
# Conditional expectation

$\beta$  is a conditional expectation of  $\alpha$  wrt.  $\mu$  along  $H$

$$\mu: D\Omega, H: X^\Omega, \alpha: L_1(\Omega, \mu), \beta: L_1(X, \mu_H)$$

$$\vdash \beta = \mathbb{E}[\alpha | H = -] \quad : \text{Prop}$$

$$:= \forall \varphi: L_1(Y, \mu_H^M). \int \mu_H(d\omega) \beta(\omega) \cdot \varphi(\omega) = \int \mu(d\omega) \alpha(\omega) \cdot \varphi(H\omega)$$



Thm (Kolmogorov): (discrete version)

There is a function

$$\mathbb{E}_{\mu}[-|-\vdash -] \in \prod_{\mu: P_{\Omega}} \prod_{H: X^{\Omega}} \mathcal{L}_1(\Omega, \mu) \rightarrow \mathcal{L}_1(X, \mu_H)$$

s.t.  $\mathbb{E}_{\mu}[\alpha | H = -]$  is a conditional expectation of  $\alpha$  w.r.t.  $\mu$   
along  $H$ .

## Conditional Probability (discrete version):

$$H: X^{\Omega}, \mu: P_X \vdash \underset{\mu}{\text{Pr}}[- \mid H = -] : (P_{\Omega})^X$$
$$:= \lambda x_0 : X. \lambda \omega_0 : \Omega. \underset{\omega \sim \mu}{\mathbb{E}} [\llbracket \omega_0 = w \rrbracket \mid H_w = x_0]$$

## Bayes's Thm (discrete version, adapted from Williams):

Let  $\lambda : P(X \times \Theta)$  joint probability distribution.

Assume  $\mu : D_X$ ,  $V : D_\Theta$  s.t.  $\lambda \ll \mu \otimes V$ .

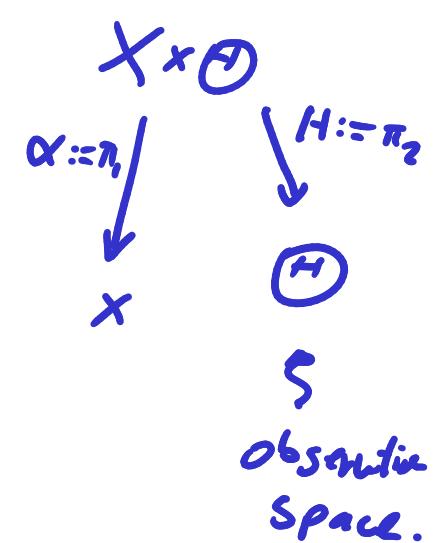
with  $d_{X,\Theta} = \frac{d\lambda}{d(\mu \otimes V)}$ .

OBS 1:  $d_X : W^X$

$$d_X := \lambda_{\Theta} \int V(d\Theta) d_{X,\Theta}(x, \theta)$$

then  $d_X = \frac{d\lambda}{d\mu}$

& similarly  $(d_{\Theta} : W^\Theta) := \lambda_{\Theta} \int \mu(dx) d_{X,\Theta}(x, \theta) = \frac{d\lambda_{\Theta}}{d\mu}$



## Bayes's Thm (discrete version, adapted from Williams):

Let  $\lambda : P(X \times \Theta)$  joint probability distribution.

Assume  $\mu : D_X, V : D_\Theta$  s.t.  $\lambda \ll \mu \otimes V$ .

with  $d_{X,H} = \frac{d\lambda}{d(\mu \otimes V)}$  .  $d_X = \frac{d\lambda}{d\mu}$   $d_\Theta = \frac{d\lambda_H}{dV}$

Let  $d_{X|H}^{(-|\cdot)} : X \times \Theta \rightarrow W$

$$d_{X|H}^{(-|\cdot)}(x|\theta) := \begin{cases} d_H \theta \neq 0: & \\ & \\ \text{o.w.:} & \end{cases}$$

$$\frac{d_{X,H}(x,\theta)}{d_H \theta}$$
  

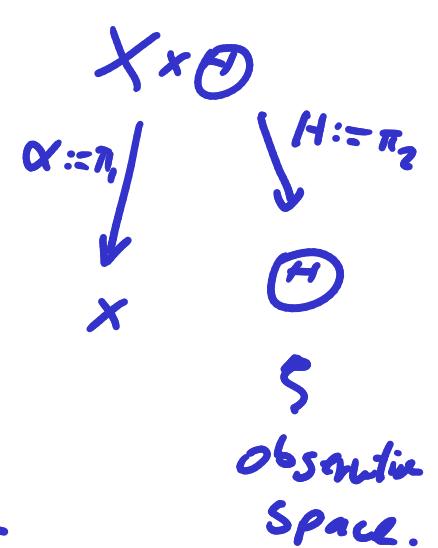
$$0$$

$$\lambda_{X|H=-} : \Theta \rightarrow P_X$$

$$\lambda_{X|H=\theta} := d_{X|H}^{(-|\theta)} \odot \mu$$

Bayes's formula:

$$P_r[-|H=-] = \lambda_{X|H=-}$$



## Summary

$\mu \otimes \nu$  Product measures & Fubini-Tonelli;

$\mu_H$  Push-forward / law

$(D^X, \Sigma, (\cdot))$  module structure and affine linearity of  $\phi$

} Lebesgue integration

Standard vocabulary: joint dist., marginalisation, independence, invariance

density & Radon-Nikodym derivatives (heed the Warning)

almost certain properties

Conditional expectation & Probability

with Bayes's Thm.

Plan:

- 1) type-driven Probability: discrete case (Mon + Tue) ✓
- 2) Borel sets & measurable spaces (Wed) (Tue)
- 3) Quasi Borel spaces, Simple type structure (Wed)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

Please ask questions!

Smibble



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# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
4–16 December, 2023  
Canberra, ACT, Australia



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Institute

Facebook Research NCSC

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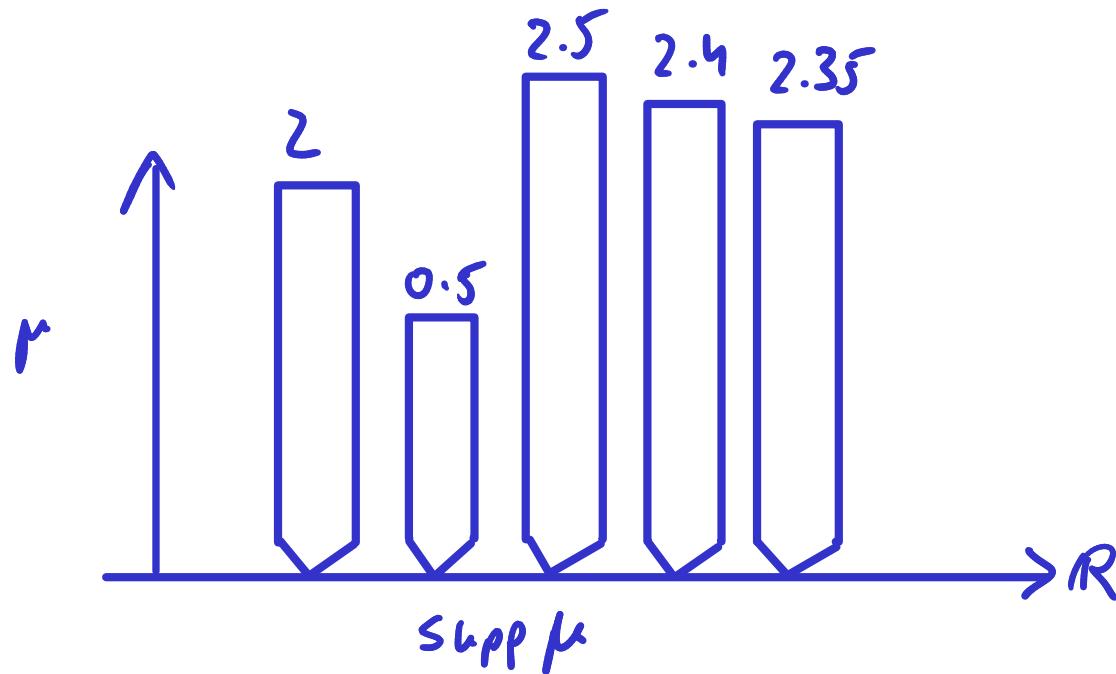
Please ask questions!

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discrete model measure only histograms:



Want :

- lengths
- areas
- volumes .

Continuous *Caveat:*

Thus: No  $\lambda: \mathcal{P}R \rightarrow [0, \infty]$ :

$$\lambda(a, b) = b - a$$

(generalises length)

$$\lambda(r + A) = \lambda A$$

(translation invariant)

$$\lambda\left(\bigcup_{n=0}^{\infty} A_n\right) = \sum_{n=0}^{\infty} \lambda A_n$$

$\sigma$ -additive

Thm: no  $\lambda: \mathcal{P}R \rightarrow [0, \infty]$ :

$$\lambda(a, b) = b - a$$

(generalises length)

$$\lambda(r + A) = \lambda A$$

(translation invariant)

$$\lambda\left(\bigcup_{n=0}^{\infty} A_n\right) = \sum_{n=0}^{\infty} \lambda A_n$$

$\sigma$ -additive

Direct proof in standard analysis courses. Idea behind typical proof is:

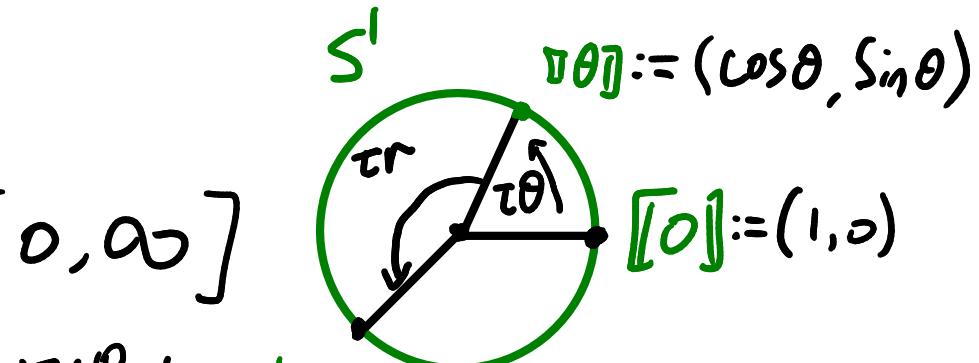
Thm: no  $\lambda: \mathcal{PS}' \rightarrow [0, \infty]$

s.t.

a) satisfy measure axioms for  $\mathcal{BS}' := \mathcal{PS}'$

b) invariant under rotations:  $E: \mathcal{BS}' \mapsto$

$$\lambda S' = \tau \quad (= 2\pi)$$



$$r: \mathbb{R} \mapsto \text{rotate}_r[\theta] := [\theta + \tau r]$$

$$\lambda \text{rotate}[E] = \lambda E$$

Reduce  $(S^i, \lambda^{S^i})$  to  $(R, \lambda^R)$  via restriction & push forward

$$\lambda^R_{|} := \lambda_{E \in P, i} \cdot \lambda_E : P[0,1] \rightarrow W$$

$$\lambda^{S^i} := \lambda_{E \in S^i} \cdot \lambda^R_{P[0,1]}(I - I^{-1}[E]) : PS^i \xrightarrow{I^{-1}} P[0,1] \xrightarrow{\lambda^R_{[0,1]}} W$$

noting

rotations in  $S^i \iff$  translations in  $R$

Since  $\exists \lambda^{S^i}$ , we have  $\exists \lambda^R$  either.

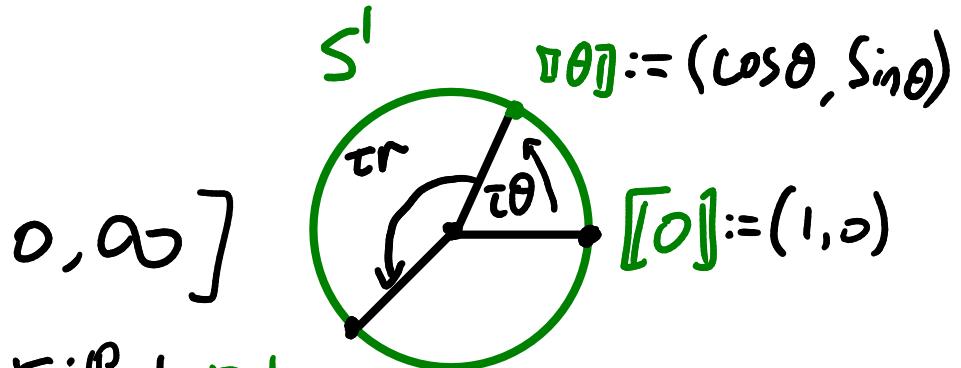
Thm: no  $\lambda: \mathcal{P}S' \rightarrow [0, \infty]$

st.

a) Satisfy measure axioms for  $BS := PS'$

b) invariant under rotations:  $E: BS' \vdash$

c)  $\lambda S' = \tau$  ( $\approx 2\pi$ )



$r: \mathbb{R} \vdash \text{rotate}_r [\theta] := [\theta + \tau r]$

$$\lambda \text{rotate}[E] = \lambda E$$

Proof:  $a+b \Rightarrow \neg c :$

1) Using axiom of choice (AoC):

$$S' = \bigcup_{i=0}^{\infty} E_i; \quad E_i = \text{rotate}_{r_i} [E_0]$$

$$2) \lambda S' = \sum_{i=0}^{\infty} \lambda E_i = \sum_i \lambda \text{rotate}_{r_i} E_0 = \sum_{i=0}^{\infty} \lambda E_0 = \begin{cases} \lambda E_0 = 0 : 0 \\ \lambda E_0 > 0 : \infty \end{cases} \neq \tau$$

Constructing  $E_i$ :

$$x, y : S' \vdash x \sim y := \exists q \in Q. \underset{q}{\text{rotate}} x = y \quad : \text{Prop}$$
$$\equiv \exists q \in [0,1] \cap Q. \text{rotate } x = y$$

$\sim$ -Equivalence classes:

$$x : S' \vdash [x]_{\sim} := \{ y \in S' \mid x \sim y \} \quad : \mathcal{P}S'$$

$$C := \{ [x]_{\sim} \in \mathcal{P}S' \mid x \in S' \}$$

$$\forall e \in C, e \neq \emptyset, \text{ so by AoC: } \exists \xi : C \rightarrow S'. \xi_e \in e.$$

---

NB:  $\xi$  injective

Take  $C_0 := \{\xi_e \in S' \mid e \in C\} \in \mathcal{PS}'$

Note:  $x \sim y, x, y \in C_0 \vdash x = y$ .

$q : Q \vdash C_q := \text{rotate}_q[C_0] \in \mathcal{PS}'$

Let  $(r_i)_{i=0}^{\infty}$  enumerate  $Q \cap [0, 1)$  st.  $r_0 = 0$

Take  $E_i := C_{r_i}$

By fiat:  $E_i = C_{r_i} = \text{rotate}_{r_i}[C_0] = \text{rotate}_{r_i}[E_0]$

RTP:  $S' = \bigcup_{i=0}^{\infty} E_i$

NB:  $x, y : S' \vdash$   
 $\text{any} : \text{Prop}$   
 $C = \sim\text{-equiv.}$   
 $\xi : C \rightarrow S'$   
 $e : C \vdash \xi_e \in \mathcal{E}$

$E_i \cap E_j = \emptyset, \quad i \neq j :$

---

$x \in E_1 \cap E_2 \Rightarrow \exists y_i \in \zeta. \quad x = \text{rotate}_{r_i} y_i$

$\Rightarrow y_1 \sim x \sim y_2 \Rightarrow y_1 = y_2 =: y$

$\Rightarrow \text{rotate}_{r_2 - r_1} y = y, \quad |r_2 - r_1| < 1$

$\Rightarrow r_1 = r_2$

$S = \bigcup_{i=0}^{\infty} E_i : x \in S'.$  letting  $e := \xi_{[x]_n} : \rho S'$

$\xi_e, x \in e \Rightarrow \xi_e \sim x$

$\Rightarrow \exists q \in (\mathbb{Q} \cap [0, 1]). \text{rotate}_q \xi_e = x.$

As  $\xi_e \in C_0 : x \in C_q.$  Find  $i$  s.t.  $r_i = q$

and  $x \in C_{r_i} = E_i.$



Takeaway: taking  $B/R := \mathcal{P}R$

Excludes measures such as:

length, area, volume

Workaround: only measure well-behaved subsets

Df: The Borel Subsets  $B_{\mathbb{R}} \subseteq \mathcal{P}(\mathbb{R})$ :

- Open intervals  $(a, b) \in B_{\mathbb{R}}$

Closure under  $\sigma$ -algebra operations:

$$\underline{\quad}$$

$$\emptyset \in B_{\mathbb{R}}$$

Empty set

$$\underline{A \in B_{\mathbb{R}}}$$

$$A^c := \mathbb{R} \setminus A \in B$$

↑  
complements

$$\overrightarrow{A} \in B_{\mathbb{R}}^N$$

$$\overline{\bigcup_{n=0}^{\infty} A_n \in B_{\mathbb{R}}}$$

countable unions

## Examples

discrete Countable:  $\{r\} = \bigcap_{\varepsilon \in \mathbb{Q}^+} (r-\varepsilon, r+\varepsilon) \in \mathcal{B}_{\mathbb{R}}$

$I$  countable  $\Rightarrow I = \bigcup_{r \in I} \{r\} \in \mathcal{B}_{\mathbb{R}}$

Closed intervals:  $[a,b] = (a,b) \cup \{a,b\}$

Non-examples?

More complicated: analytic, lebesgue

Df:

Measurable Space  $V = (V, \mathcal{B}_V)$

Set  
(Carrier)  
Family of  
Subsets  
 $\mathcal{B}_V \subseteq P(V)$

closed under  $\sigma$ -algebra operations:

—

$\emptyset \in \mathcal{B}_V$

Empty set

$A \in \mathcal{B}_V$

$A^c := V \setminus A \in \mathcal{B}_V$

↑  
complements

$\vec{A} \in \mathcal{B}_V^N$

$\overline{\bigcup_{n=0}^{\infty} A_n \in \mathcal{B}_V}$   
countable unions

Idea: Structure all spaces after the worst-case scenario

## Examples

- Discrete spaces

$$X^{\text{meas}} = (X, \mathcal{P}X)$$

- Euclidean spaces

$\mathbb{R}^n$  — replace intervals with  
charts  $\prod_{i=1}^n (a_i, b_i)$

Similarly

$$\{C \cap A \mid C \in \mathcal{B}_V\}$$

- Sub spaces:  $A \in \mathcal{P}V$      $A := (A, [\mathcal{B}_V] \cap A)$

- Products:  $A \times B := ([A] \times [B], \sigma([\mathcal{B}_A] \times [\mathcal{B}_B]))$

Def: Borel measurable functions  $f: V_1 \rightarrow V_2$

- functions  $f: V_1 \rightarrow V_2$
- inverse image preserves measurability:

$$f^{-1}[A] \in \mathcal{B}_{V_1} \iff A \in \mathcal{B}_{V_2}$$

### Examples

- $(+), (\cdot) : \mathbb{R}^2 \rightarrow \mathbb{R}$
- any continuous function  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$
- any function  $f: X^n \rightarrow V$
- $| - |, \sin: \mathbb{R} \rightarrow \mathbb{R}$

# Category Meas

Objects : Measurable spaces

Morphisms : Measurable functions

Identities:

$$id : V \rightarrow V$$

Composition:

$$f : V_2 \rightarrow V_3 \quad g : V_1 \rightarrow V_2$$

$$f \circ g : V_1 \rightarrow V_3$$

## Meas Category

Products, Co products / disjoint union, Subspaces

Categorical limits, colimits, but:

Thm [Arrow '61] No  $\sigma$ -algebras  $B_{B_R}, B_{R^R}$  for measurable

membership ←  $(\exists) : (B_R, B_{B_R}) \times R \rightarrow \text{Bool}$   
 $(U, r) \mapsto [r \in U]$

$\text{eval} : (\text{Meas}(R, \mathcal{V}R), B_{R^R}) \times R \rightarrow \mathcal{V}$   
 $(f, r) \mapsto f(r)$

Questions? Skip proof?

Proof (sketch) :

Borel hierarchy:

$$\Sigma^0_\omega \subset \Delta^0_1 \subset \Sigma^0_1 \subset \Delta^0_2 \subset \dots \subset \Delta^0_\omega \subset \dots \subset \Delta^0_{\omega+1}$$
$$\Pi^0_0 \subset \Pi^0_1 \subset \dots \subset \Pi^0_\omega \subset \dots$$

Stabilises at  $\Delta^0_{\omega_1} = B(\Sigma^0_\omega) = \Delta^0_{\omega_1 + 1}$

$$\text{rank } A := \min \{ \alpha < \omega_1 \mid A \in \Delta^0_\alpha \}$$

$$(\exists) : (\mathcal{B}_{\mathbb{R}}, \mathcal{B}_{\mathbb{R}}) \times \mathbb{R} \rightarrow \mathbb{R}$$

(U, r)  $\mapsto$  [r  $\in$  U]

for  $\mathcal{B}_{\mathbb{R}} = P(\mathcal{B}_{\mathbb{R}})$

If measurable:

$$\mathcal{B}_{V \times U} = \mathcal{B}([\mathcal{B}_V] \times [\mathcal{B}_U])$$

$$\alpha := \text{rank}((\exists)^{-1}[\text{true}]) < \omega,$$

Take  $A \in \mathcal{B}_{\mathbb{R}}$ ,  $\text{rank } A > \alpha$

But:

$$\alpha < \text{rank } A = \text{rank}((A, -)^{-1}[(\exists)^{-1}[\text{true}]]) \leq \text{rank}((\exists)^{-1}[\text{true}]) \leq \alpha$$

\*

More details in Ex. B

Sequential Higher-order Structure:

I Countable :  $V^{\mathbb{I}} = \prod_{i \in \mathbb{I}} V$

$\Rightarrow$  Some higher-order structure in Meas:

Cauchy  $\in B_{[-\infty, \infty]^N}$

$$\text{Cauchy} := \bigcap_{\epsilon \in \mathbb{Q}^+} \bigcup_{k \in \mathbb{N}} \bigcap_{\substack{m, n \in \mathbb{N} \\ m, n \geq k}} \{ \vec{y} \in [-\infty, \infty]^N \mid |y_m - y_n| < \epsilon \}$$

$$\limsup : [-\infty, \infty]^N \rightarrow [-\infty, \infty]$$

$$\lim : \text{Cauchy} \rightarrow \mathbb{R}$$

Compose higher-order building blocks:

lim IS measurable!  
}

$$\text{VanishingSeq}(\mathbb{R}) := \left\{ \vec{r} \in \mathbb{R}^N \mid \lim_{n \rightarrow \infty} r_n = 0 \right\} \in \mathcal{B}_{\mathbb{R}^N}$$

$$\text{approx}_- : \text{VanishingSeq}(\mathbb{R}^+) \times \mathbb{R} \rightarrow \mathbb{Q}^N$$

s.t.:  $|(\text{approx}_{\Delta} r)_n - r| < \Delta_n$

Slogan: Measurable by Type !

Not all operations of interest fit:

$$\limsup : ([-\infty, \infty]^{\mathbb{R}})^N \rightarrow [-\infty, \infty]^{\mathbb{R}}$$

$$\limsup := \lambda f. \lambda n. \limsup_{n \rightarrow \infty} f_n x$$

Intrinsically  
higher-order !

Want

Slogan: measurability by type!

But

For higher-order building blocks

defer measurability proofs until

we resume 1<sup>st</sup> order fragment  $\Rightarrow$  <sup>non</sup>composition

Plan:

- 1) type-driven Probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed) ✓
- 3) Quasi Borel spaces (Wed) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)

Please ask questions!

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# Plan

Def:  $V \in \text{Meas}$  is Standard Borel when

$$V \cong A \quad \text{for some } A \in \mathcal{B}_R$$

the "good part" of  $\text{Meas}$  – the subcategory

$$\text{Sbs} \hookrightarrow \text{Meas}$$

Sbs including

- Discrete  $\mathbb{I}$ ,  $\mathbb{I}$  countable
- Countable products of Sbs:

$$\mathbb{R}^n, \mathbb{R}^\mathbb{N}, \mathbb{Z}^n, \mathbb{N}^\mathbb{N}$$

- Borel subspaces of Sbs:

$$\mathbb{I} := [0, 1]$$

$$\mathbb{R}^+ := (0, \infty) \quad \mathbb{R}_{\geq 0} := [0, \infty]$$

- Countable coproducts of Sbs:

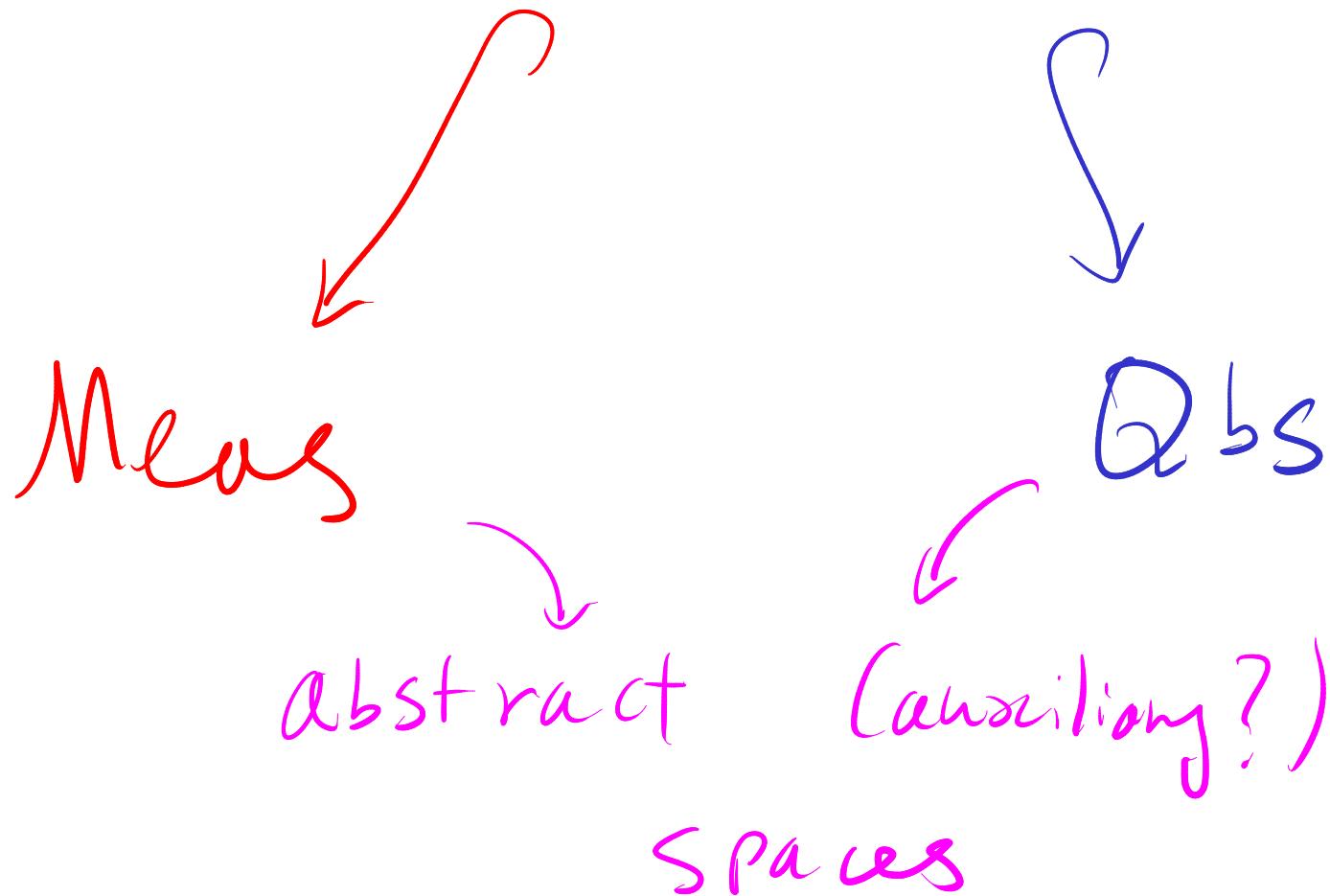
$$\mathbb{W} := [0, \infty]$$

$$\overline{\mathbb{R}} := [-\infty, \infty]$$

Conservative extensions:

Concrete spaces  
we "observe"

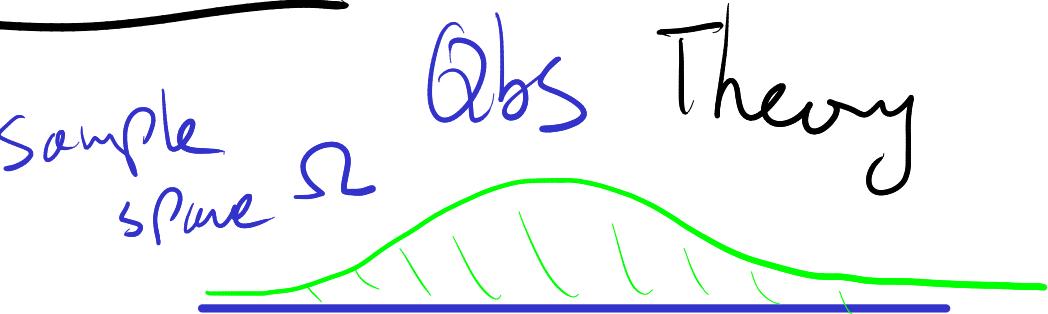
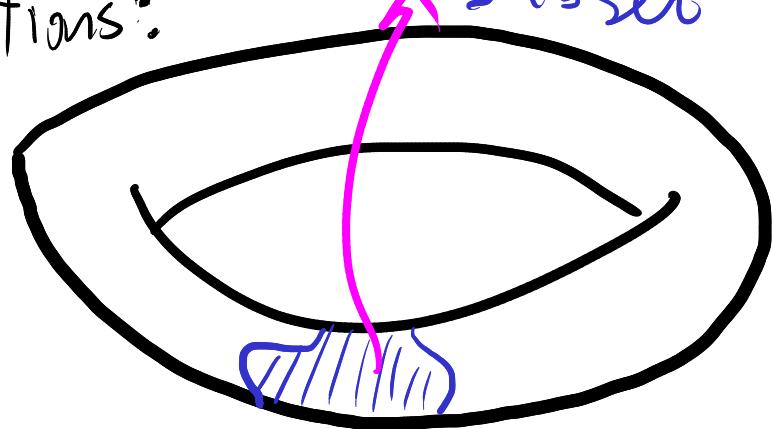
Standard Borel spaces



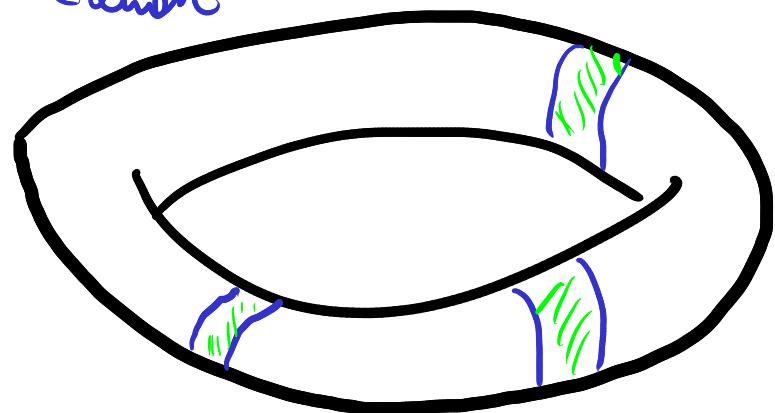
# Core idea

Measure Theory

Primitive notions:



random element  $\downarrow \alpha$



Derived

notions:

random

elements

$\alpha: \Omega \rightarrow \text{Space}$

measure

measurable  
subsets

Def: Quasi-Borel space  $X = (X, \mathcal{R}_X)$

$$\mathcal{R}_X \subseteq L^{\mathbb{R}_X}$$

Closed under:

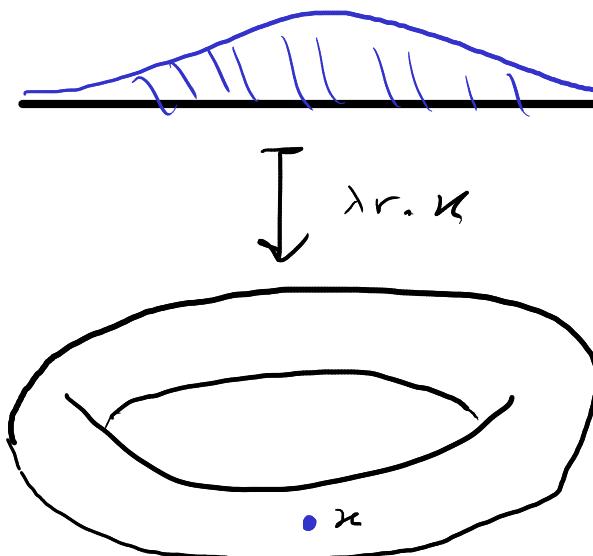
Set ↗  
"carrier"  
Set of  
functions  $\alpha: \mathbb{R} \rightarrow X$   
"random elements"

- Constants:

$$\begin{array}{c} x \in X \\ \hline (\lambda r. x) \in \mathcal{R}_X \end{array}$$

- precomposition:

- recombination



Def: Quasi-Borel space  $X = (LX, R_X)$

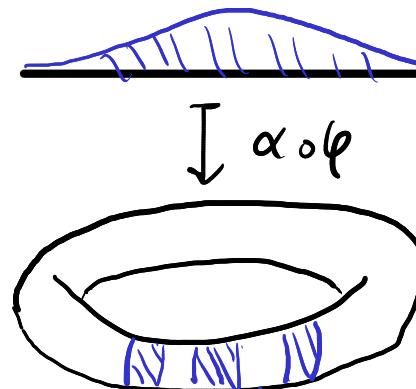
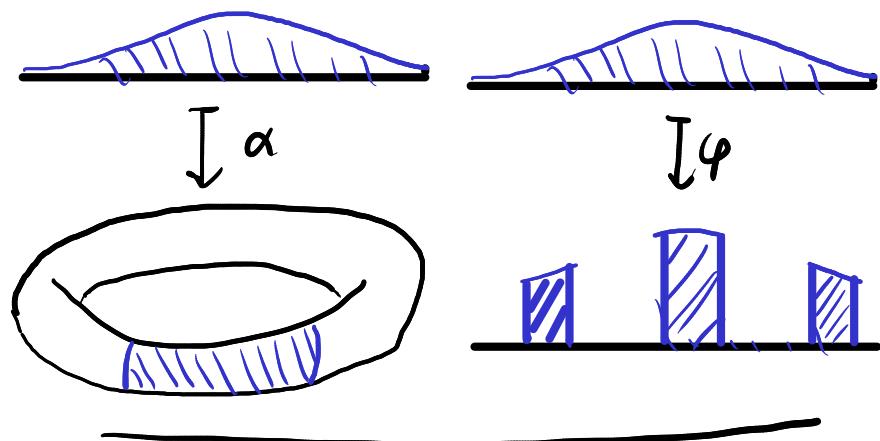
$$R_X \subseteq L^{R_J} \quad \text{closed under:}$$

- precomposition:

$$\alpha \in R_X \quad \varphi: \mathbb{R} \rightarrow \mathbb{R} \text{ in } Sbs$$

$$(\varphi \circ \alpha): \mathbb{R} \xrightarrow{\varphi} \mathbb{R} \xrightarrow{\alpha} LX \in R_X$$

Set  $\curvearrowleft$  Set of  
"carrier"  
"random elements"



Def: Quasi-Borel space  $X = (LX, RX)$

$$RX \subseteq LX^{\mathbb{N}}$$

Closed under:

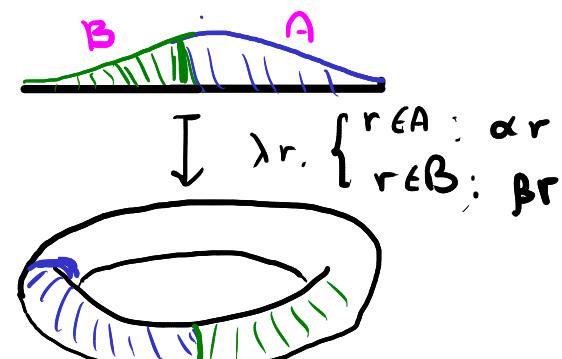
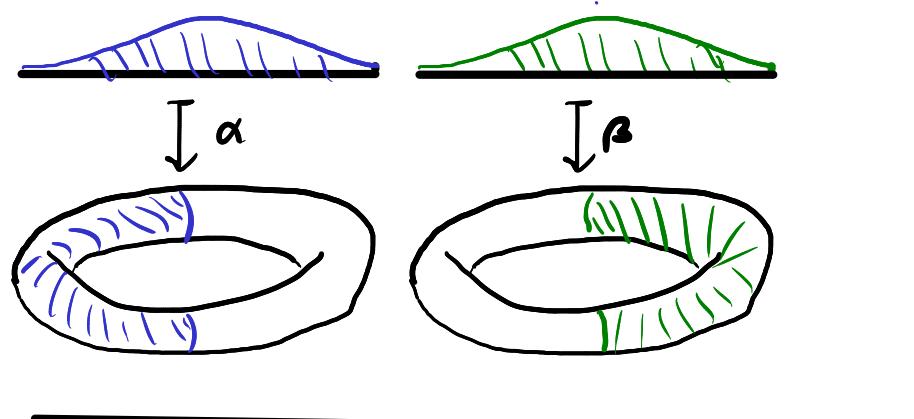
- recombination

$$\vec{\alpha} \in RX^{\mathbb{N}}$$
$$R = \bigcup_{n=0}^{\infty} A_n$$

$EB_R$

$$\lambda r. \left\{ \begin{array}{l} : \\ r \in A_n : \alpha_n r \\ : \end{array} \right.$$

Set ↗  
"carrier"  
Set of  
functions  $\alpha: \mathbb{N} \rightarrow X$   
"random elements"



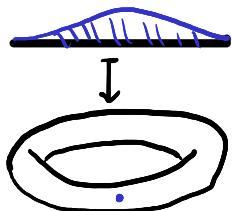
Ref: Quasi-Borel space  $X = (X_1, \mathcal{R}_X)$

$$\mathcal{R}_X \subseteq L^1(X_1, \mathbb{R})$$

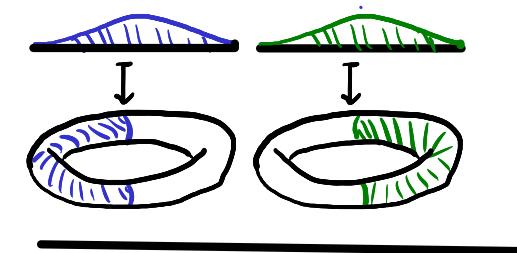
Closed under:

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Functions  $\alpha: \mathbb{R} \rightarrow X_1$   
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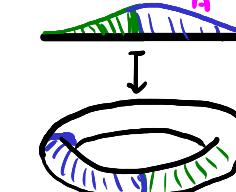
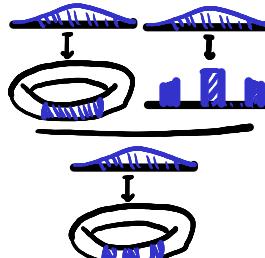
- Constants:



- recombination



- precomposition:



## Examples

recombination of  
constants

$$- \mathbb{R} = (\mathbb{R}, \text{Meas}(\mathbb{R}, \mathbb{R}))$$

qbs underlying  $\mathbb{R}$

$$- X \in \text{Set}, \quad \mathcal{X}^{\text{Qbs}} := (X, \sigma\text{-simple}(\mathbb{R}, X))$$

$\lambda r.$  {  
  : rEA<sub>n</sub>:  $x_n$   
  :  
  :}

discrete qbs on  $X$

$$- " \quad \mathcal{X}_{\text{Qbs}} := (X, X^{\mathbb{R}})$$

all functions

Indiscrete qbs on  $X$

Qbs morphism  $f: X \rightarrow Y$

- function  $f: X \rightarrow Y$

- $$\alpha \begin{matrix} \downarrow^R \\ \downarrow_X \\ \downarrow^L \end{matrix} \in R_X$$

---

$$\alpha \begin{matrix} \downarrow^R \\ \downarrow_X \\ \downarrow^L \\ f \downarrow \\ \downarrow^L_Y \end{matrix} \in R_Y$$

Example

- Constant functions

one qbs  
morphism

- σ - simple functions  
are qbs morphisms

Category Qbs



- identity, composition

## Full model

$$\text{type} : \text{Qbs} \quad \mathbb{W} := [0, \infty] \quad \mathcal{B}x := (\text{Thur})$$

$$DX := (\text{Fri})$$

$$PX := \left\{ \mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1 \right\} \quad (\text{Thu})$$

$$\underset{\mu}{\text{Ce}}[E] := (\text{Fri}) \quad S_x := (\text{Fri})$$

$$\phi \mu k := (\text{Fri})$$

Plan:

- 1) Type-driven Probability: discrete case (Mon + Tue)
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# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
4–16 December, 2023  
Canberra, ACT, Australia



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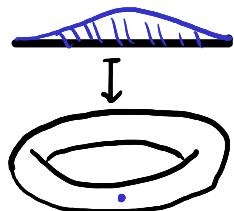
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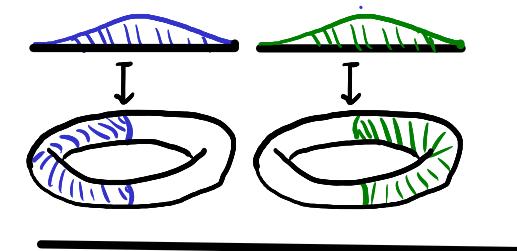
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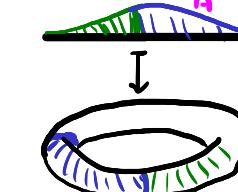
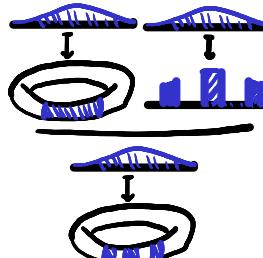
- Constants:



- recombination



- precomposition:



## Examples

recombination of  
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$$- \mathbb{R} = (\mathbb{R}, \text{Meas}(\mathbb{R}, \mathbb{R}))$$

qbs underlying  $\mathbb{R}$

$$- X \in \text{set}, \quad \mathcal{X}^{\text{Qbs}} := (X, \sigma\text{-simple}(\mathbb{R}, X))$$

$\lambda r.$  {  
  : rEA<sub>n</sub>:  $x_n$   
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  :}

discrete qbs on  $X$

$$- " \quad \mathcal{X}_{\text{Qbs}} := (X, X^{\mathbb{R}})$$

all functions

Indiscrete qbs on  $X$

Validate gbs axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(R, \mathbb{W}))$

- Constants:

$$E : B_{\mathbb{W}}, \kappa : \mathbb{W} \vdash$$

$$(Ar : R. x)^{-1}[E] = \begin{cases} x \in E : & R \\ x \notin E : & \emptyset \end{cases} \in B_R$$

✓

Validate gbs axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(R, \mathbb{W}))$

- Precomposition:

$\alpha: \text{Meas}(R, \mathbb{W}), \varphi: \text{Meas}(R, R) \vdash$

$$R \xrightarrow{\varphi} R \xrightarrow{\alpha} \mathbb{W} \quad \begin{matrix} \in \text{Meas}(R, \mathbb{W}) \\ \beta \end{matrix}$$

$\text{Meas}$  is a cat.

Explicitly:

$$(a \circ \varphi)^{-1}[E] \in \mathcal{B}R \xleftarrow{\varphi^{-1}} \alpha^{-1}[E] \in \mathcal{B}R \xleftarrow{\alpha^{-1}} E \in \mathcal{B}\mathbb{W} \quad \checkmark$$

Validate qbs axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(R, \mathbb{W}))$

- RL Combination

$$\begin{aligned} I^{\text{ctbl}}, \alpha: \text{Meas}(IR, \mathbb{W})^I, E_i: B_{IR}, R = \bigcup_{i \in I} E_i, F: B_W \vdash \\ \left( \exists r. \left\{ \begin{array}{l} : \\ r \in E_i : \alpha; r \\ \vdots \end{array} \right\}^{-1} [F] \right) \\ \beta := \bigcup_{i \in I} \alpha_i^{-1}[F] \wedge E_i \in B_R \end{aligned}$$

In fact:

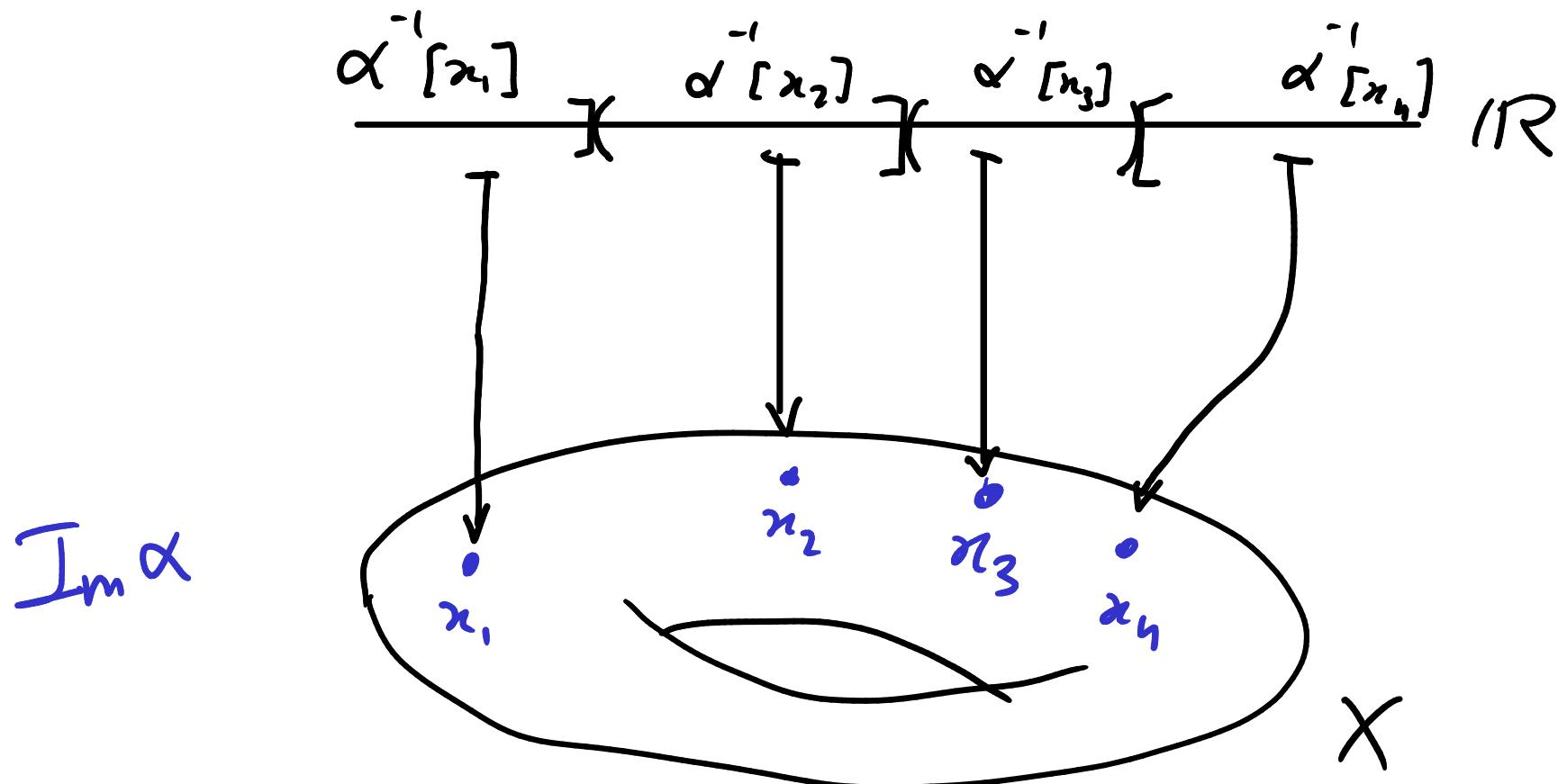
$$r \in \text{LHS} \Leftrightarrow \beta \vdash F \Leftrightarrow \exists i \in I. r \in E_i \wedge \alpha_i; r \vdash F \Leftrightarrow r \in \text{RHS}$$



## $\sigma$ -Simple function

$\alpha: R \rightarrow X$  s.t.  $Im \alpha := \alpha[R]$  is ctbl 1

$$\forall x \in Im \alpha . \alpha^{-1}[x] \in \mathcal{B}_R$$



Validate QBS axioms for:  $\Gamma^{\text{QBS}, \dagger} := (X, \sigma\text{-Simple}(X))$

- Constants

$$\text{Im}(\lambda r.x) = \{x\} \text{ ctbl} \quad \checkmark$$

NB:  $f \sigma\text{-Simple} : \text{Im } f \text{ ctbl} \wedge \tilde{f}[x] \in B_R$

$$y:X \vdash (\lambda r.x)^{-1}[y] = \begin{cases} x=y & : R \\ x \neq y & : \emptyset \end{cases} \in B_R \quad \checkmark$$

Validate q<sup>bs</sup>'s axioms for:  $X^{\text{qbs}, \dagger} := (X, \sigma\text{-simple}(X))$

• Precomposition:

$\alpha : \sigma\text{-simple}(X), \varphi : \text{Meas}(R, R) \vdash$

$$\text{Im}(\alpha \circ \varphi) \subseteq \text{Im} \alpha \text{ ctbl} \quad \checkmark$$

NB:  $f \text{ } \sigma\text{-simple} :$   
 $\text{Im } f \text{ ctbl}$  &  
 $\tilde{f}[x] \in \mathcal{B}_R$

$x : X \vdash$

$$(\alpha \circ \varphi)^{-1}[x] = \varphi^{-1}[\alpha^{-1}(x)] \in \mathcal{B}_R \quad \checkmark$$

$$\alpha^{-1}(x) \in \mathcal{B}_R$$

$\varphi : R \rightarrow R$  measurable

Validate qbs axioms for:  $X^{\text{qbs}} := (X, \sigma\text{-Simple}(X))$

• recombination:

$$\alpha := (\sigma\text{-Simple}(X))^I, E := B_R^I, R = \bigoplus_{i \in I} E_i$$

NB:  $f \sigma\text{-Simple} : \text{Im } f \subset \bigcup_{i \in I} f(E_i) \subseteq B_R$

$$\text{Im}[E_i \cdot \alpha_i]_{i \in I} \subseteq \bigcup_{i \in I} \text{Im } \alpha_i \quad \checkmark$$

$x : X \vdash$

$$[E_i \cdot \alpha_i]_{i \in I}^{-1}(x) = \bigcup_{i \in I} \alpha_i^{-1}[x] \cap E_i \in B_R \quad \checkmark$$

Prop:  $X : \text{Set}, A : \text{Qbs} \vdash$

•  $\forall f : X \rightarrow {}_L A_S . f : {}^r_X^{\text{Qbs}} \rightarrow A$

•  $\forall f : {}_L A_S \rightarrow X . f : A \rightarrow {}^X_L \text{Qbs}^l$

Prop:  $X : \text{Set}, A : \text{Qbs} \vdash$

- $\forall f : X \rightarrow {}_L A_S . f : {}^{r_{\text{Qbs}_S}}_X \rightarrow A$

Prf:  $\alpha : R_{{}^{r_{\text{Qbs}_S}}_X} \vdash \alpha \text{ } \sigma\text{-simple} \Rightarrow$

$$\alpha = [\bar{\alpha}[x]. \lambda r. x]_{x \in \text{Im } \alpha} \Rightarrow$$

$$(f \circ \alpha) = [\bar{\alpha}[x]. \lambda r. f x]_{x \in \text{Im } \alpha}$$

recombination

$\in R_A$

Borel

constant  $\in B_A$

ctbl

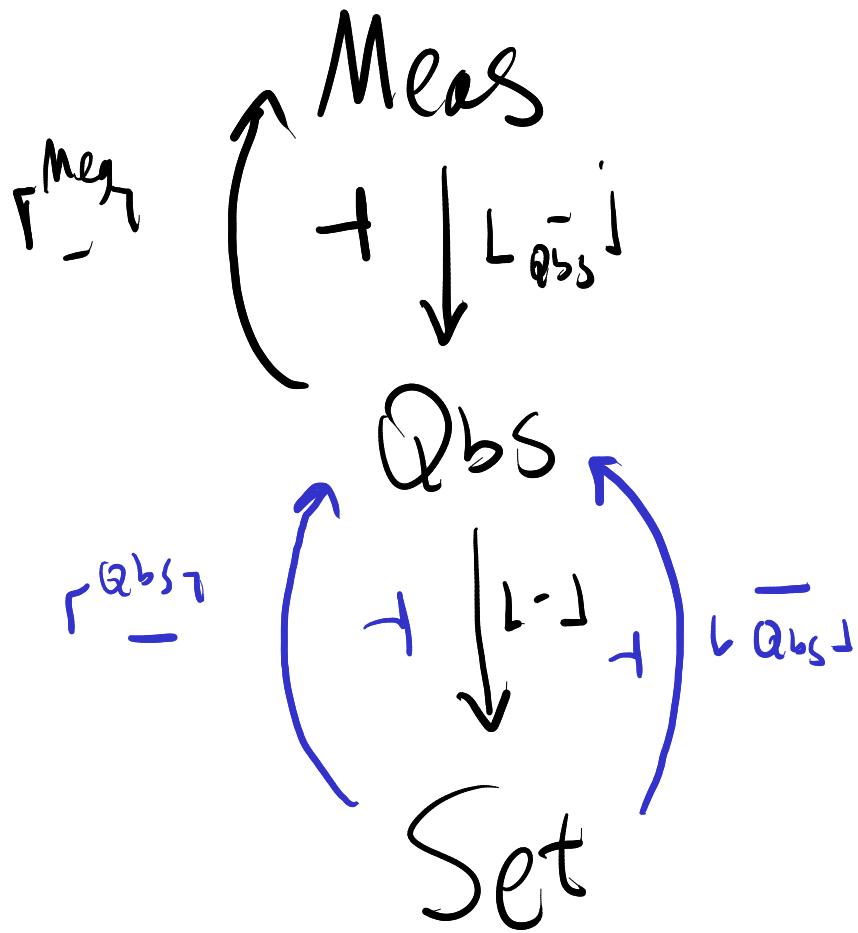
Prop:  $X : \text{Set}, A : \text{Qbs} \vdash$

- $\forall f : X \rightarrow {}_L A_S . f : {}^r_{X^{\text{Qbs}}} \rightarrow A$
- $\forall f : {}_L A_S \rightarrow X . f : A \rightarrow {}^r_{L^X_{\text{Qbs}}}$

Prf:  $\alpha : R_A \vdash (f \circ \alpha : R \rightarrow X) \in R_{{}^r_{X^{\text{Qbs}}}}$  always. ✓



# Useful adjunctions:



$$\begin{aligned} \underline{V}_{\text{Qbs}} &:= (\underline{V}_1, \text{Meas}(R, V)) \\ &\quad (V \in \text{Meas}) \\ \Gamma_X^{\text{Meas}} &:= \left\{ A \subseteq \underline{X}_1 \mid \forall \alpha \in R_X, \alpha^{-1}[A] \in B_R \right\} \end{aligned}$$

- limits (products, subspaces)  
and colimits (co-products, quotients)
- as in Set
- Slogan: every measurable space is carried by a qbs

## Example

Product  $(X \times Y, \pi_1, \pi_2)$ :

necessarily!

$$- L[X \times Y] = L[X_1 \times_1 Y]$$

$$- R_{X \times Y} = \{ \lambda r. (\alpha r, \beta r) \mid \alpha \in R_X, \beta \in R_Y \}$$

corresponding  
raw comb  
elements

rest of structure as in Set.

# Function Spaces

Straightforward |  
•

-  $\lfloor Y^X \rfloor := \text{Qbs}(X, Y)$

-  $R_{Y^X} := \text{Uncurry}[\text{Qbs}(R^{XX}, Y)]$

$$= \left\{ \alpha : R \rightarrow \lfloor Y^X \rfloor \mid \lambda(r, x). \alpha \circ x : R \times X \rightarrow Y \right\}$$

- eval :  $Y^X \times X \rightarrow Y$

$$\text{eval}(f, x) := fx$$

# Meas vs Obs

By generalities:

$\sigma$ -algebra  
on  $\text{Meas}(\mathbb{R}, \mathbb{R})$

meas -

$$\text{Meas} \xrightarrow{+} L_{\bar{\theta}_S^i} j$$

$\Gamma \vdash_{\text{IR}} M \in \Sigma$

Mos R

$$\mathbb{R} \times \mathbb{R}$$

$$\mathbb{R} \times \mathbb{R}$$

$$\cancel{\text{---} \rightarrow} \quad r[R] = [A]$$

No factorisation  
by  
Aumann's  
Theorem.

$\Gamma$  Meas  $\gamma$   
 $\beta V(a)$

$$R^R \times R^R \neq [R^R]^R \times [R^R]$$

## Simple Type Structure

"Simple" because:

- Simply-typed  $\lambda$ -calculus
- types are simple:  $A, B : \text{Type} \vdash B^A : \text{Type}$ 
  - no polymorphism
  - no term dependency
- contexts for terms:  $\Gamma \vdash t : A$ 
  - are simple:  $\Gamma = x_1 : A_1, \dots, x_n : A_n$ 
    - i.e.  $\text{List}(\text{Type})$

## Simple Type Structure

"Simple" because:

- interpretation is simple :

$$\llbracket x_1 : A_1, \dots, x_n : A_n \rrbracket := \prod_{i=1}^n A_i;$$

$$\llbracket \Gamma \vdash t : A \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow A$$

in QBS

## Simple Type Structure

Curry-Howard-Lambek

$$\frac{\Gamma \vdash t : A \quad \Gamma \vdash s : B}{\Gamma \vdash \langle t, s \rangle : A \times B}$$

$$\rightsquigarrow \llbracket \Gamma \rrbracket \xrightarrow{\lambda r. \langle tr, sr \rangle} A \times B$$

is measurable

$$\frac{\Gamma \vdash t : A \times B \quad \Gamma, x:A, y:B \vdash s : C}{\Gamma \vdash \text{let } (x,y)=t \text{ in } s : C}$$

$$\rightsquigarrow$$

measurability  
by  
type!

$$\lambda x. \text{let } (a,b)=tr \text{ in } s \vdash [x \mapsto a, y \mapsto b]$$

$$\llbracket \Gamma \rrbracket \longrightarrow C$$

is measurable. etc.

## Random Element Space

$R_X := X^R$  since  $\lfloor X^R \rfloor = R_X$  as sets.

Why?

( $\subseteq$ )  $\alpha \in \lfloor X \rfloor^R \Rightarrow \alpha: \mathbb{R} \rightarrow X$  in Qbs.

$\text{id}_{\mathbb{R}}: \mathbb{R} \rightarrow \mathbb{R}$  measurable  $\Rightarrow \text{id} \in R_{\mathbb{R}}$

$\Rightarrow \alpha = \alpha \circ \text{id} \in R_X$

( $\supseteq$ )  $\alpha \in R_X \Rightarrow \forall \varphi \in R_{\mathbb{R}} = \text{Meas}(\mathbb{R}, \mathbb{R})$ .  $\alpha \circ \varphi \in R_X \Rightarrow \alpha: \mathbb{R} \rightarrow X$   
 $\Rightarrow \alpha \in \lfloor X \rfloor^R$

## Subspaces

For  $X \in \mathbb{Q}bs$ ,  $A \subseteq X$ , set:

$$R_A := \left\{ \alpha: \mathbb{R} \rightarrow A \mid \alpha \in R_X \right\}$$

Then  $A = (A, R_A)$  is the *Subspace qbs*

We write  $A \hookrightarrow X$

## Borel Subspaces Ensemble

The  $\sigma$ -algebra  $B_X := \{ A \subseteq X \mid \forall \alpha \in R_X . \alpha^*[A] \in B_R \}$

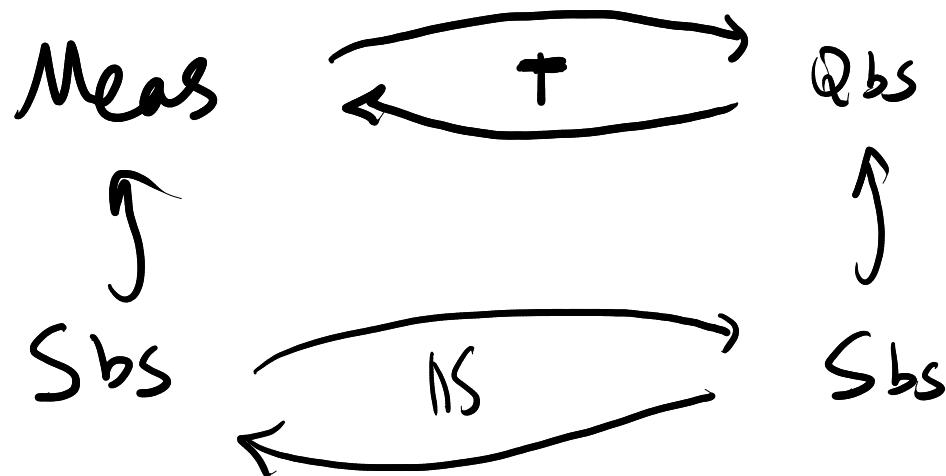
internalises as  $B_X = 2^X$ , the qbs of  
Borel Subsets.

$L(B_{(B_R)})$  are the Borel-on-Borel sets from  
descriptive set theory.  
(cf. [Sabou et al. '21])

# Standard Borel spaces

Def: A qbs  $S$  is **standard Borel** when

$$S \cong A \text{ for some } A \in \mathcal{B}_{\mathbb{R}}$$



**Slogan:** Qbs **Conservative extension** of Sbs

Example  $C_0 := \{f: \mathbb{R} \rightarrow \mathbb{R} \mid f \text{ continuous}\} \hookrightarrow \mathbb{R}^{\mathbb{R}}$

$C_0$  is sbs. (Well-known!)

Proof:

$$C_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$$

↑ sbs!

$$C'_0 := \left\{ g \in \mathbb{R}^{\mathbb{Q}} \mid \forall a, b \in \mathbb{Q}, \forall \varepsilon \in \mathbb{Q}^+ \exists \delta \in \mathbb{Q}^+ \forall p, q \in \mathbb{Q} \text{ s.t. } p, q \in [a, b], |p - q| < \delta \Rightarrow |g(p) - g(q)| < \varepsilon \right\}$$

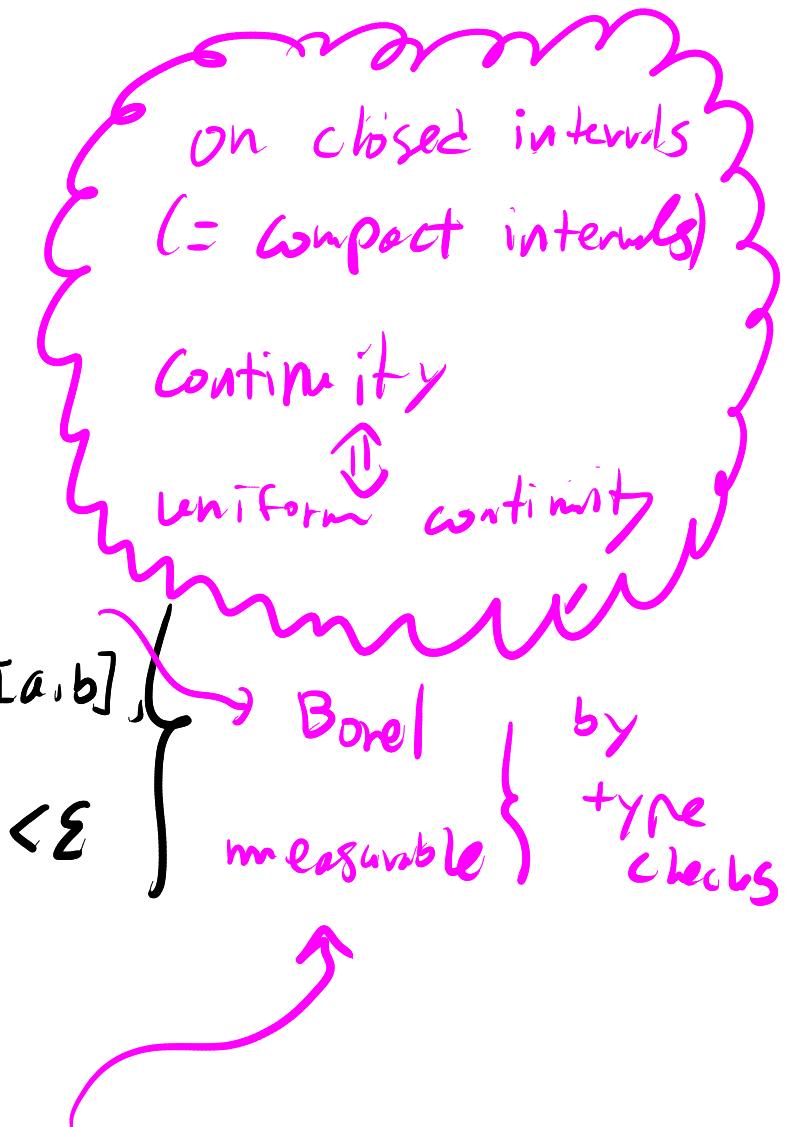
Then  $C_0 \cong C'_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$ :

$$C_0 \rightarrow C'_0$$

$$\varphi \mapsto \varphi|_{\mathbb{Q}}$$

$$C'_0 \rightarrow C_0$$

$$\varphi \mapsto \lambda r. \lim_{n \rightarrow \infty} g(\text{approx}_{\frac{1}{n}} \text{approx}_{\frac{1}{m}})_{n,m \in \mathbb{N}}$$



## Example (ctd)

$C_0$  is sbs, and  $\text{eval}: C_0 \times \mathbb{R} \rightarrow \mathbb{R}$   
is measurable.

Avoids:

- constructing complete separable metrics
- proving that evaluation is measurable w.r.t. metric  $\sigma$ -algebra.

Non-examples ~ [Sabok et al.'21]

$$-\left\{ A \in \mathcal{B}_{\mathbb{R}} \mid A \neq \emptyset \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$

$$-\left\{ (A_1, A_2) \in \mathcal{B}_{\mathbb{R}}^2 \mid A \subseteq B \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}^2$$

$$-\left\{ A \in \mathcal{B}_{\mathbb{R}} \mid A \text{ open} \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$

## Plan:

- 1) Type-driven Probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed) ✓
- 3) Quasi Borel spaces (Wed) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu) ✓
- 5) Integration & random variables (Fri)

Please ask questions!

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# Dependent Type Structure

+ types can contain terms : a type referring to a term

$$X:\text{Type}, E:B_X \vdash \{x \in X \mid x \in E\} : \text{Type}$$

a type, just like  
STLC

a term!

# Dependent Type Structure

+ types can contain terms :

$$X : \text{Type}, E : B_X \vdash \{x \in X \mid n \in E\} : \text{Type}$$

a type, just like  
STLC

a term!

a type referring  
to a term

Content formation:

$$\frac{\Gamma \vdash A : \text{Type}}{\Gamma, x : A \vdash}$$

# Dependent Type Structure

types denote spaces-in-Content

$$\begin{array}{c} \boxed{\Gamma \vdash A} \\ \downarrow \text{dep} \\ \boxed{\Gamma \vdash} \end{array}$$

Dependent types denote spaces-in-Content

$\Gamma \vdash \text{Content}$

$\Gamma \vdash A$

type in content

E.g.:

A

↓

1

simple types

$[\Gamma \vdash A]$

dep

$[\Gamma]$

Space in Content

Content Space

assigns

environment

$[E : B_A + \{x \in A \mid x \in E\}]$

$\{ (E, a) \in B_A^{X_A} \mid a \in E \}$

↓  
 $\pi_1$

$B_A$

decoher

## Content extension

$$\frac{\Gamma \vdash A}{\Gamma, a:A \vdash}$$

$$\frac{\llbracket \Gamma \vdash A \rrbracket}{\llbracket \Gamma \rrbracket \quad \llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket}$$

$\downarrow \text{dep}$

## Substitution

E.g. weakening

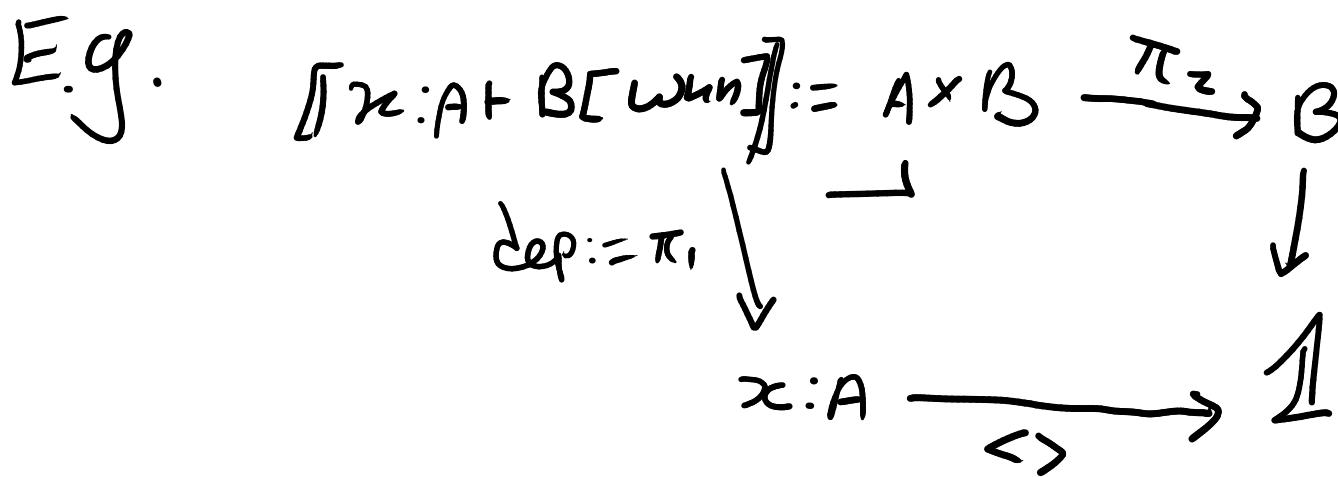
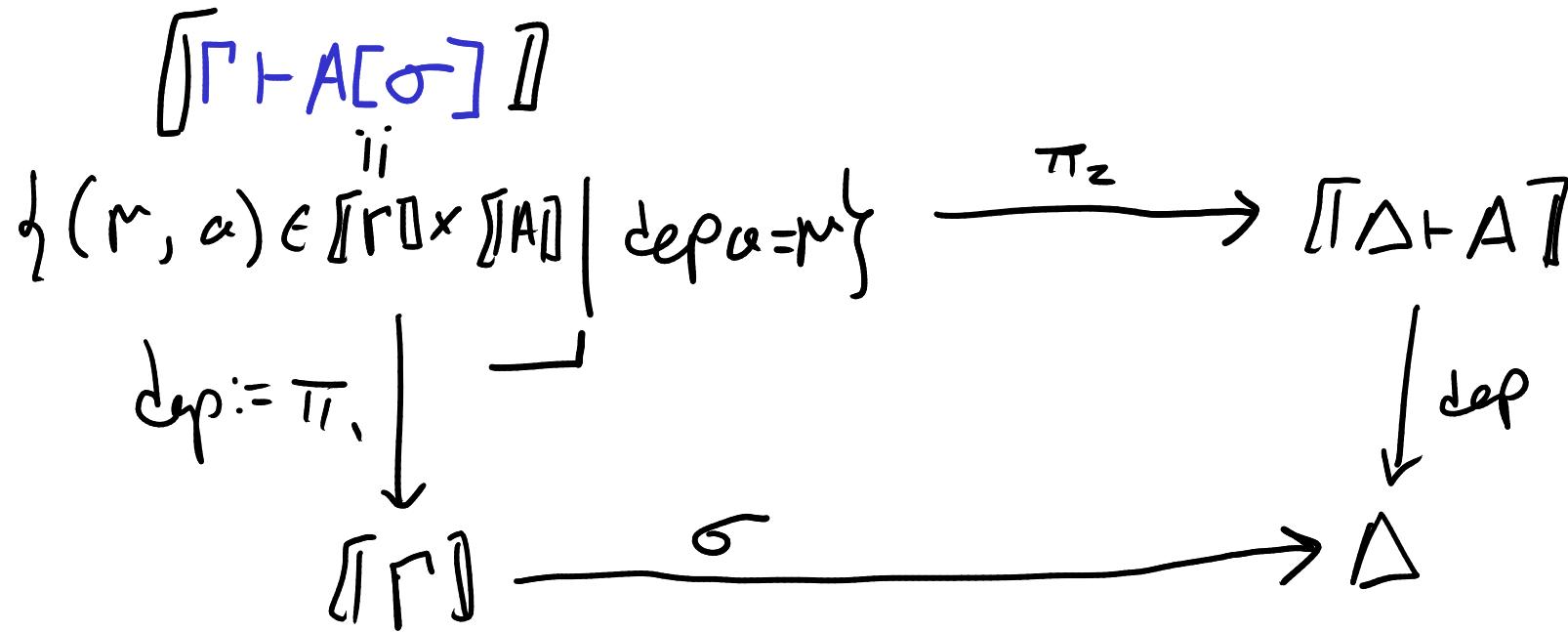
$$\Gamma \vdash \sigma : \Delta$$

$$\llbracket \sigma \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket \Delta \rrbracket$$

$$\Gamma, a:A \vdash \text{wkn} : \Gamma$$

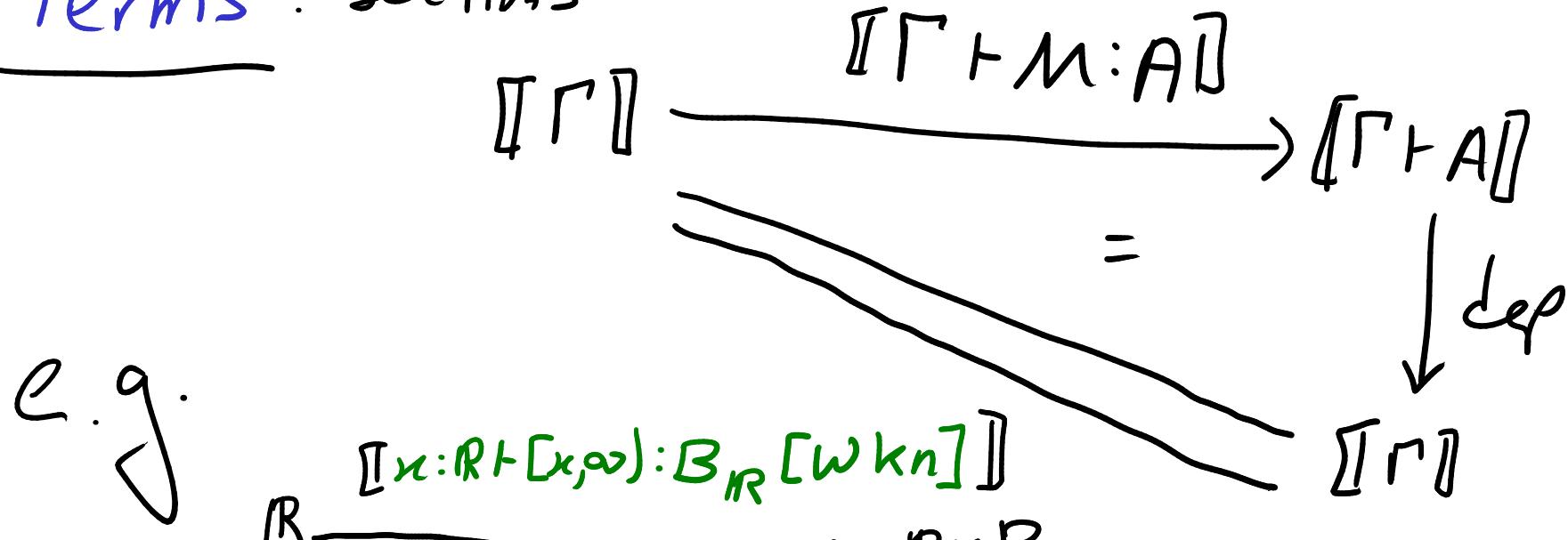
$$\llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket \xrightarrow[\text{dep}]{\text{wkn}} \llbracket \Gamma \rrbracket$$

# Action of Substitution on types



Simple type

## Terms : sections



E.g. Variables:  $\boxed{\Gamma, \alpha : A \vdash \alpha : A}$



Exercise:

action of substitution

$M[\sigma]$

## Dependent Pairs

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \prod_{a:A} B}$$

$$[\Gamma \prod_{a:A} A] := [\Gamma, a:A \vdash B]$$

$$\begin{aligned} &:= \downarrow \text{dep}_B \\ &\quad \begin{aligned} &[\Gamma, a:A] \\ &[\Gamma \vdash A] \\ &\downarrow \\ &[\Gamma] \end{aligned} \\ &\text{dep}_{\prod} \quad \swarrow \end{aligned}$$

## Dependent Products

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \prod_{a:A} B}$$

$$\prod_{a:A} B$$

$$[\Gamma \vdash \prod_{a:A} B] :=$$

$$\left\{ (m_0, f : \{ a \in [A] \mid \text{dep } a = m_0 \} \rightarrow [\Gamma, a:A \vdash B]) \middle| \right. \\ \left. \forall a \in [\Gamma, a:A]. \text{dep } a = m_0 \Rightarrow \text{dep}(fa) = a \right\}$$

Exercise: find the random elements.

aha:  $(a:A) \rightarrow B$

## Full model

$$\text{type : Obs} \quad W := [0, \infty] \quad \mathcal{B}_X \cong \mathcal{B}^X$$

$$DX := (\text{Fr}_i)$$

$$PX := \left\{ \mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1 \right\}$$

$$\underset{\mu}{\text{Ce}}[E] := (\text{Fr}_i) \quad S_x := (\text{Fr}_i)$$

$$\phi \mu k := (\text{Fr}_i)$$

Plan:

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# Foundations for type-driven probabilistic modelling

Ohad Kammar  
University of Edinburgh

Logic Summer School  
Australian National University  
4–16 December, 2023  
Canberra, ACT, Australia



THE UNIVERSITY OF EDINBURGH

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Institute

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# Partiality cf. [Väkär et al., '19]

A Borel embedding  $e: X \rightarrow Y$

- injective function  $e: [X] \rightarrow [Y]$
- its image is Borel:  $e[[x]] \in \mathcal{B}_Y$
- $e$  is Strong:  $\alpha \in R_X \Leftrightarrow e \circ \alpha \in R_Y$

## Examples

- $\mathbb{N} \rightarrow \mathbb{Z}$
- $S$  is sbs  $\Leftrightarrow \exists S \subseteq \mathbb{R}$

Def: A Partial map  $f: X \rightarrow Y$  is a morphism

$$f: X \rightarrow Y \amalg \{\perp\}$$

Its domain of definition

$$f: (Y \amalg \{\perp\})^X \vdash \text{Dom } f := \{x \in X \mid f_x \neq \perp\} : \text{Type}$$

Depent-type  
interpretation

$$\begin{array}{ccc} \llbracket \text{Dom } f \rrbracket & \longrightarrow & \{g \in Y \mid g \in E\} \\ \downarrow \text{dep} & & \downarrow \text{dep} \\ \llbracket f : (Y \amalg \{\perp\})^X \rrbracket \llbracket \underset{E \mapsto \{x \mid f_x \neq \perp\}}{\overrightarrow{x}} \rrbracket & & \llbracket E : \mathcal{B}_Y \rrbracket \end{array}$$

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## Full model

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$$\underset{\mu}{\text{Ce}}[E] := (\text{Fr}_i) \quad S_x := (\text{Fr}_i)$$

$$\phi \mu k := (\text{Fr}_i)$$

Def: A measure  $\mu$  over  $\mathbb{R}$  is a function

$$\mu : \mathcal{B}_{\mathbb{R}} \rightarrow \mathbb{W} := [0, \infty]$$

Satisfying the measure axioms:

$$E : \mathcal{B}^\omega \rightarrow$$

$$\mu \phi = 0, \quad \mu E = \mu(E \cap F) + \mu(E \cap F^c), \quad \mu(\bigvee_n E_n) = \sup_n \mu E_n$$

For measurable spaces, replace  $\mathbb{R}$  with  $V$

We write  $[GV]$  for the set of measures on  $V$

For abs  $X$ , take  $[G^{\tau_{\text{meas}}} X]$

Thm (Lebesgue measure):

There is a unique measure  $\lambda \in \mathcal{L}G(\mathbb{R})$ , s.t.:

$$\lambda(a, b) = b - a$$

Thm (Lebesgue measure):

There is a unique measure  $\lambda \in \mathcal{L}G(\mathbb{R})$ , s.t.:

$$\lambda(a, b) = b - a$$

Proof Sketch (standard analysis textbook):

- 1) restrict attention to  $(0, 1]$  & extend via  $\sigma$ -additivity
- 2) Take  $\Sigma_0 \subseteq \mathcal{B}_{(0, 1]}$   $E \in \Sigma_0 \Leftrightarrow E = \bigcup_{i=1}^n (a_i, b_i)$
- 3) Defining  $\lambda: \Sigma_0 \rightarrow \mathbb{W}$ ,  $\lambda \bigcup_{i=1}^n (a_i, b_i) := \sum_{i=1}^n (b_i - a_i)$  independent of
- 4)  $\lambda \phi = 0$ ,  $\lambda E = \lambda(E \cap F) + \lambda(E \cap F^c)$  (straightforward)

Up

5) Technical gadget:  $\forall (E_n \supseteq E_{n+1})$  in  $\Sigma_0$ ,

$$\inf \lambda_{E_n} > 0 \Rightarrow \bigcap E_n \neq \emptyset.$$

6)  $\lambda$  is continuous on  $\Sigma_0$ : If  $(E_n \subseteq E_{n+1})_n$  in  $\Sigma_0$

$$\text{and } \bigcup_n E_n \in \Sigma_0 \text{ then } \lambda \bigcup E_n = \sup_n \lambda_{E_n}$$

7) Noting that:  $\Sigma_0$  is a Boolean algebra

$$\leftarrow \sigma(\Sigma_0) = \mathcal{B}_{\{0,1\}}$$

We use Caratheodory's extension theorem:

$\lambda$  extends uniquely to  $\lambda : \mathcal{B}_{\{0,1\}} \rightarrow W$ .

# The Unrestricted Giry Spaces

Equip  $\lfloor GV \rfloor$  with two qbs structures:

$$X \quad R_{GV} := \left\{ \alpha: R \rightarrow GV \mid \forall A \in B_V, \exists r, \alpha(r, A): R \rightarrow W \right\}$$

✓  $GV \hookrightarrow W^{B_X}$

- $\alpha$  is a kernel.
- Fewer random elements
- $R_{GV} \subseteq R_{G'V}$
- Lebesgue integral measurable in both arguments.

## Farewell Meas

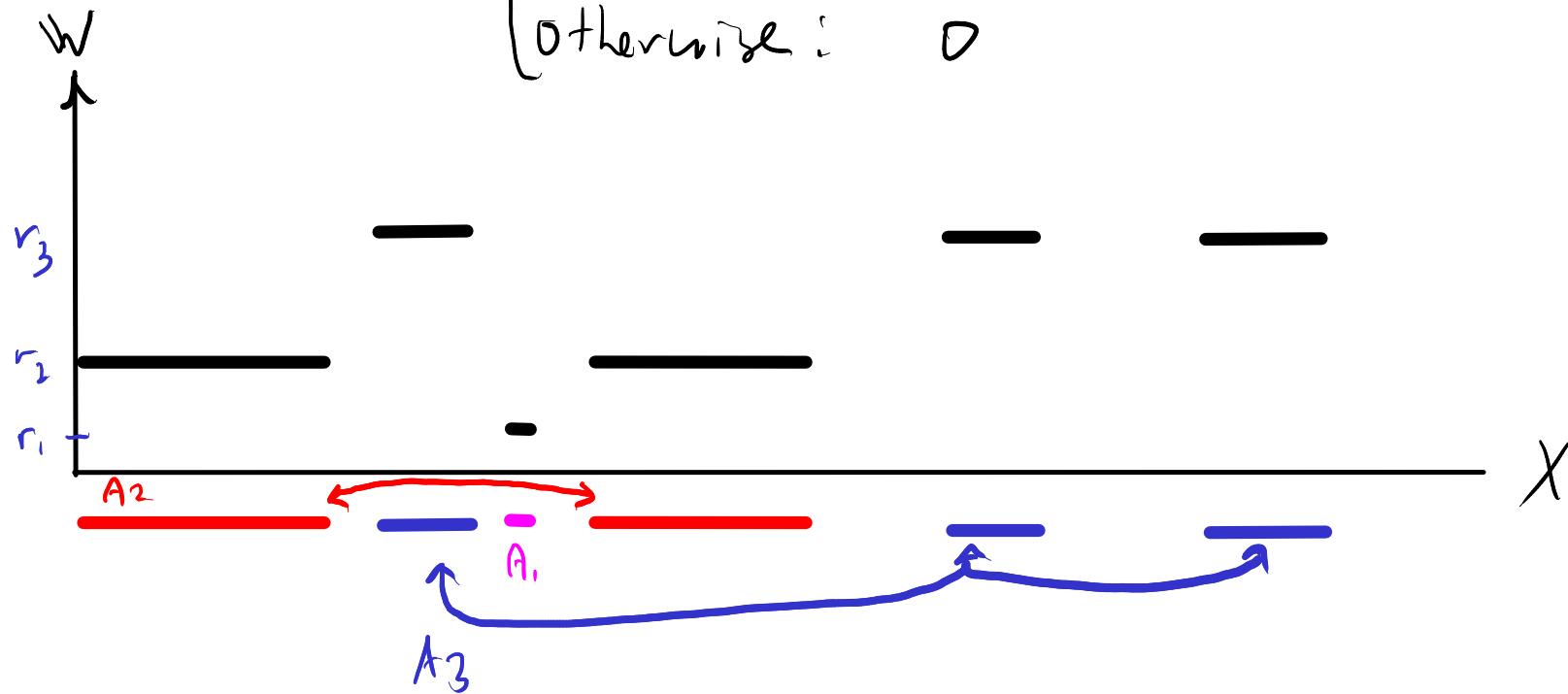
Now on:

1. All spaces are quasi-Borel (upcoming)
2. "measurable function" means qbs morphism!

Def: Simple function  $\varphi: X \rightarrow W$  when

$\exists n \in \mathbb{N}$ ,  $\vec{A} \in \mathcal{B}_X^n$ ,  $A_i \cap A_j = \emptyset$ ,  $r^i \in W$  s.t.  
 $(i \neq j)$

$$\varphi(x) = \begin{cases} \vdots & \vdots \\ x \in A_i & r_i \\ \vdots & \vdots \\ \text{otherwise: } & 0 \end{cases}$$



Encoder into a space:

$$\text{SimpleCode} := \coprod_{n \in \mathbb{N}} \mathcal{B}_X^n \times \mathcal{W}^n$$

$$\text{Simple} := \{ f \in \mathcal{W}^X \mid f \text{ simple} \} \hookrightarrow \mathcal{W}^X$$

and define an interpretation:

$$[\![\cdot]\!]: \text{SimpleCode} \longrightarrow \text{Simple}$$

$$[\![(\vec{n}, \vec{A}, \vec{r})]\!] := \sum_{i=1}^n r_i \cdot [\![\cdot \in A_i]\!]$$

↳ characteristic function  
for  $A_i$

Lemma:  $f: X \rightarrow W$  is measurable → remember!  
qbs  
morphism!

iff  $f = \lim_{n \rightarrow \infty} f_n$  for some monotone sequence

$f_n \in \text{Simple}$ .

Moreover, we have measurable such choice.

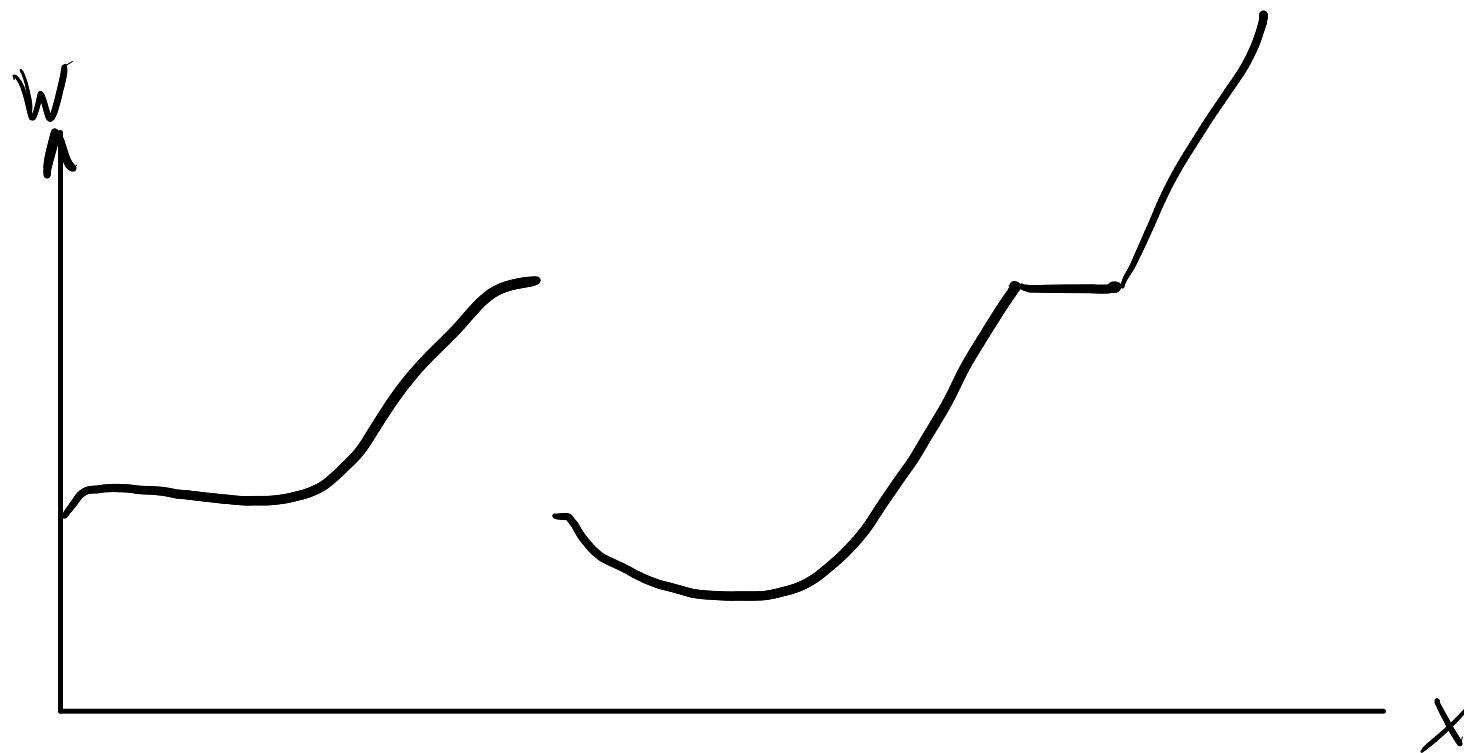
Simple Approx:

$$\left\{ \vec{\alpha} \in \mathbb{R}^+ \mid \Delta_n \rightarrow 0 \right\} \times \left\{ \vec{\alpha}' \in W^{IN} \mid \begin{array}{l} \vec{\alpha} \text{ monotone} \\ a_n \rightarrow \infty \end{array} \right\} \times W^X \rightarrow \text{SimpleCode}$$

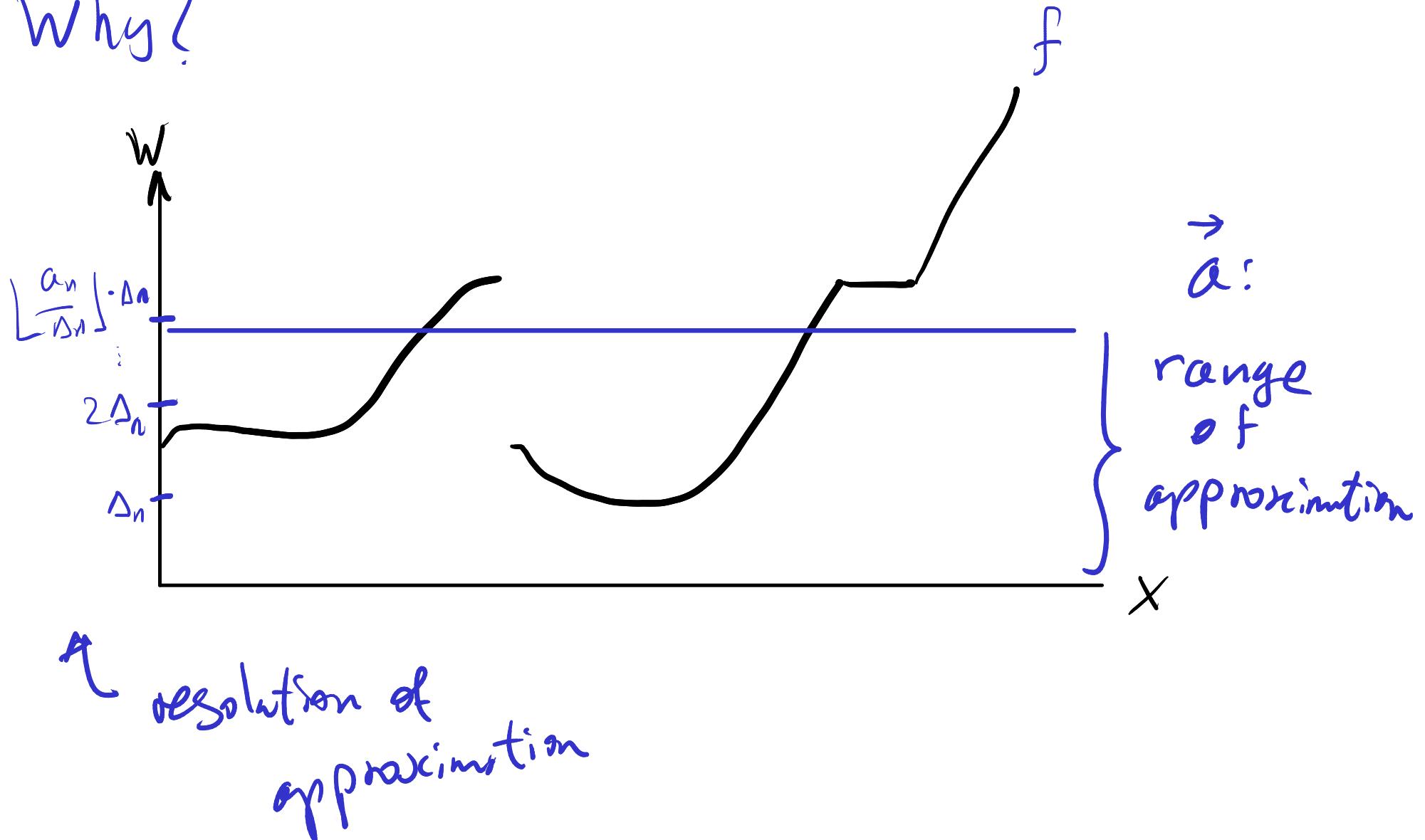
$\uparrow$   
rate of convergence

$\uparrow$   
range of approximation

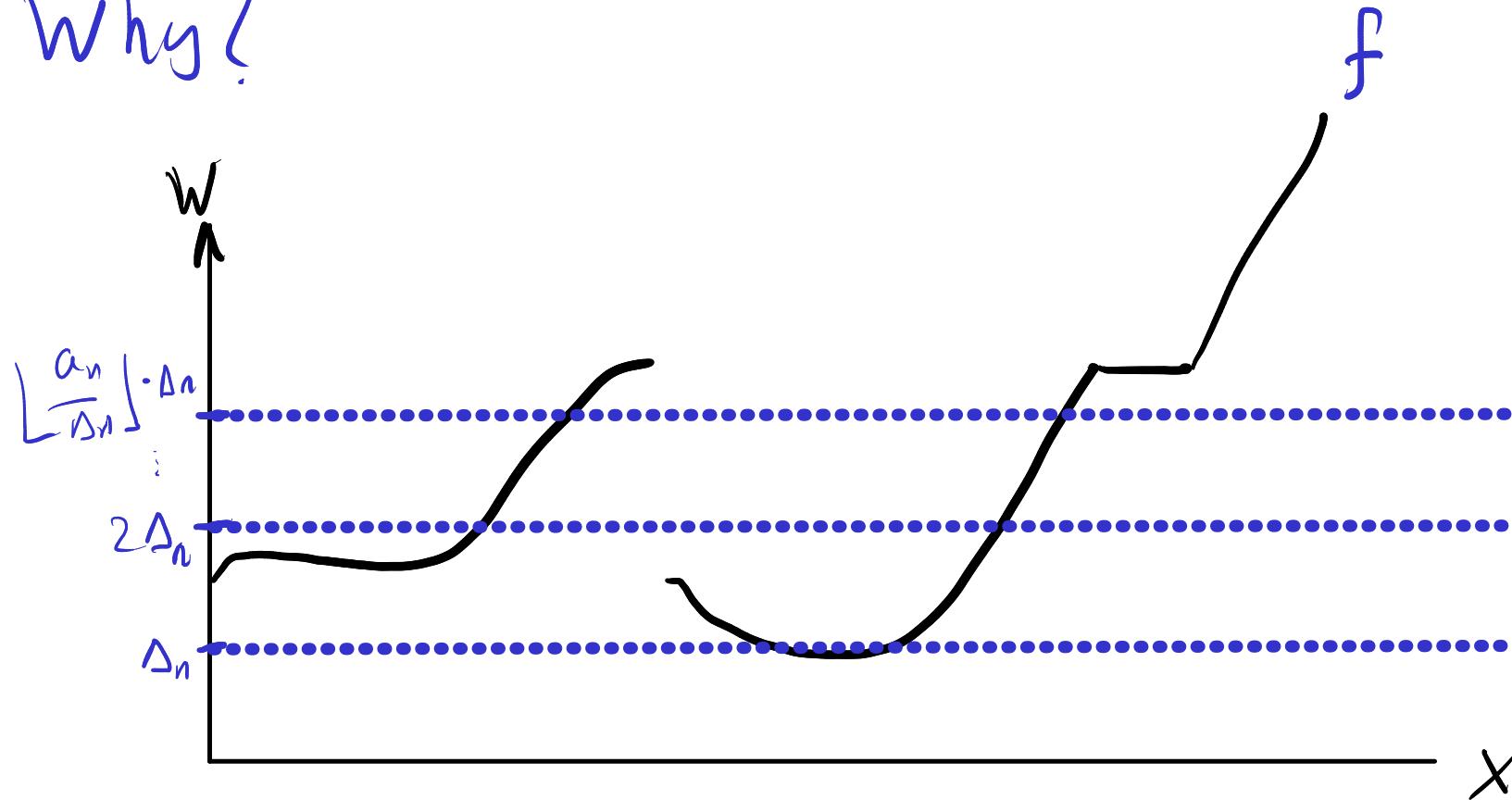
Why?



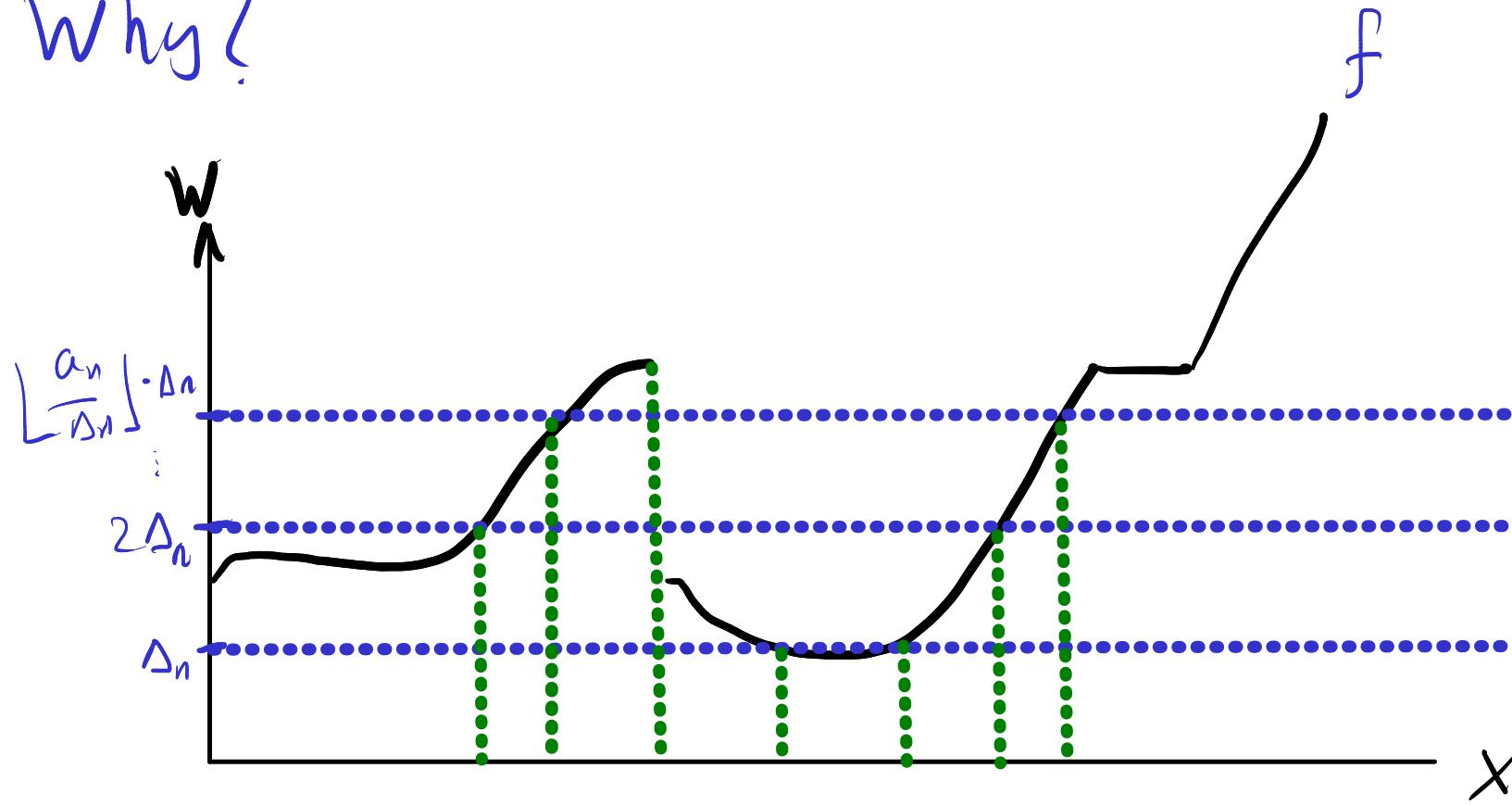
Why?



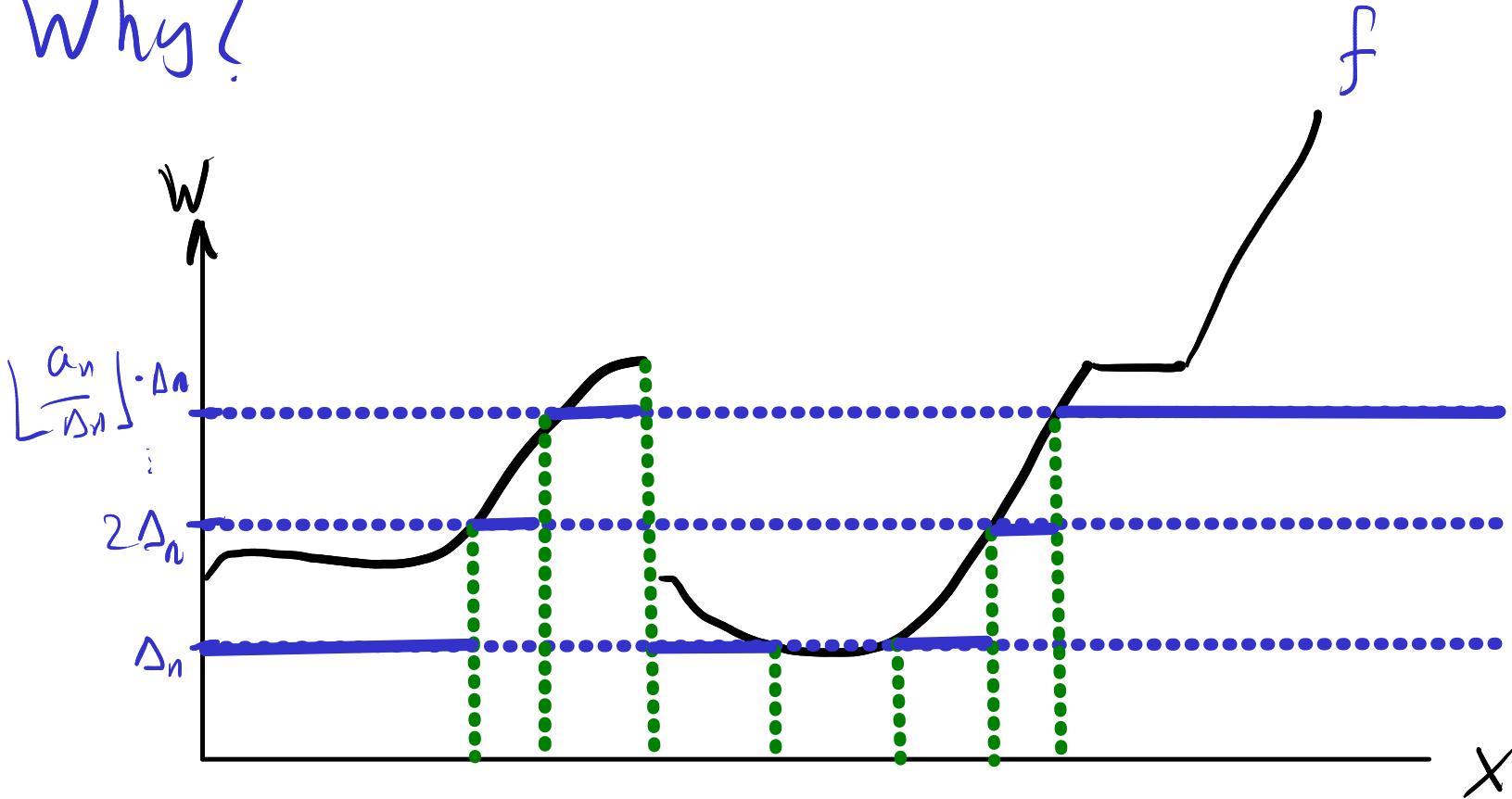
Why?



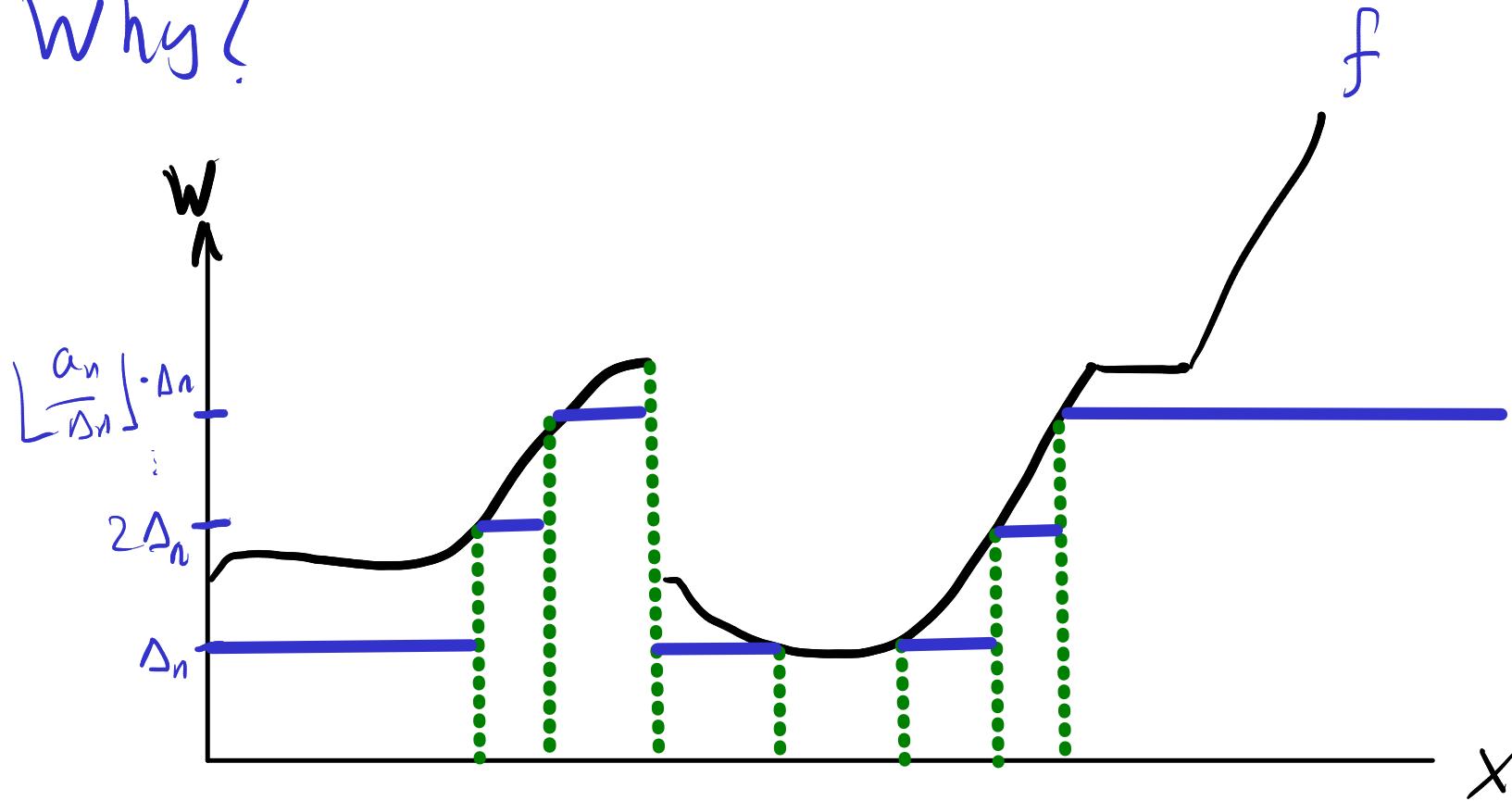
Why?



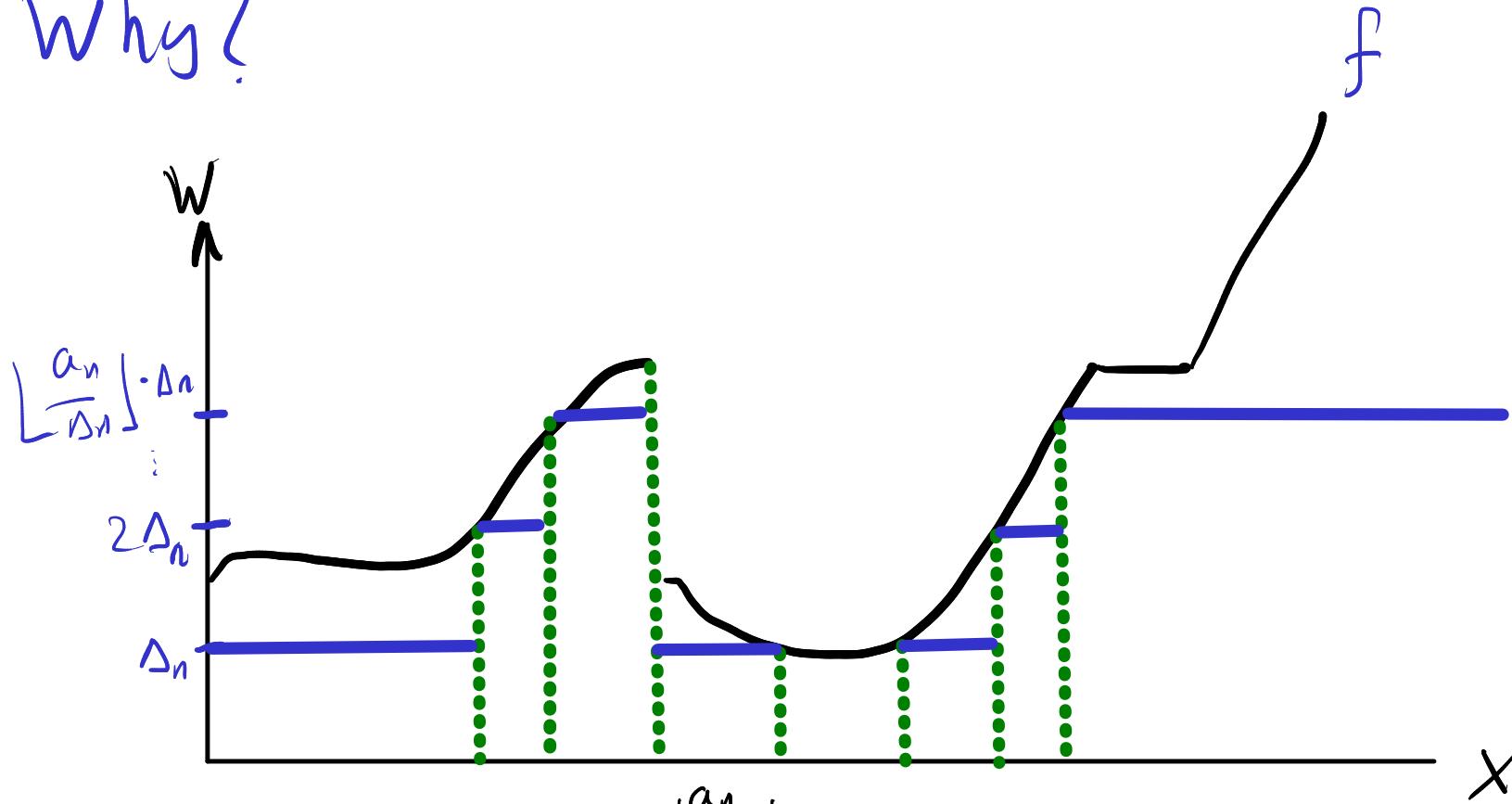
Why?



Why?

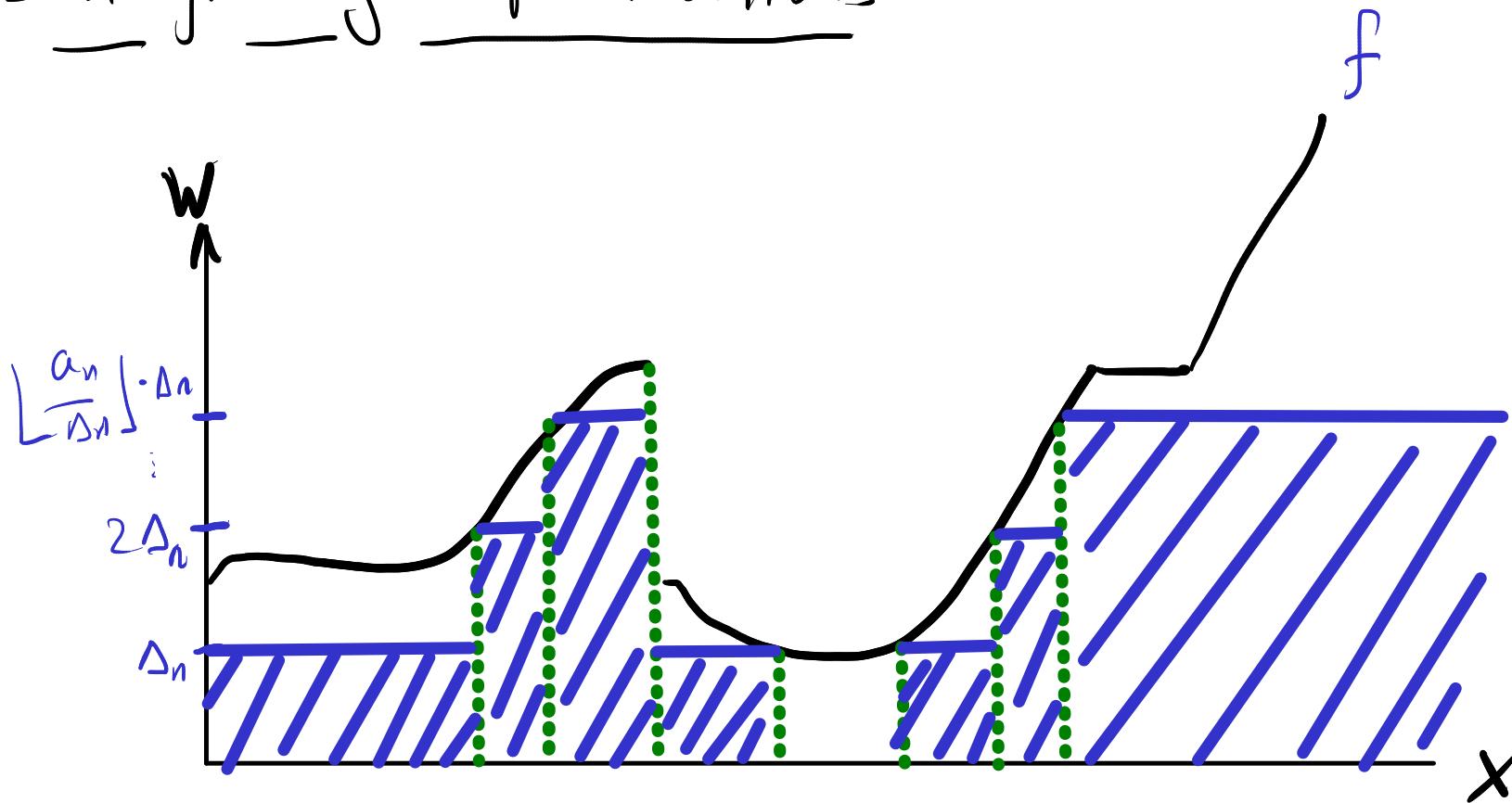


Why?



$$\left\| \text{Simple Approx} \xrightarrow{\Delta, \alpha} f \right\| := \sum_{i=1}^{\lfloor \frac{a_n}{\Delta_n} \rfloor} i \cdot \Delta_n [i \cdot \Delta_n \leq f < (i+1) \Delta_n] + \lfloor \frac{a_n}{\Delta_n} \rfloor \Delta_n \cdot [f \geq \lfloor \frac{a_n}{\Delta_n} \rfloor \cdot \Delta_n] \in \text{Simple}$$

# Integrating Simple Functions



$\int : G X \times \text{Simple Code} \rightarrow \mathbb{W}$

$$\int \mu(n, \vec{A}, \vec{r}) := \sum_{I \subseteq \{1, \dots, n\}} \left( \sum_{i \in I} r_i \right) \cdot \mu \left( \bigcap_{i \in I} A_i \setminus \bigcup_{i \notin I} A_i \right)$$

# Integration

Proper higher-order operation

$$\int : Gx \times W^X \rightarrow W$$

$$\int^\mu f := \sup \left\{ \int^\mu \varphi \mid \varphi \in \text{Simple}, \quad \varphi \leq f \right\}$$

we also  
write

$$= \lim_{n \rightarrow \infty} \int^\mu (\text{Simple Approx}_{\vec{\Delta}, \vec{a}} f)_n \sim \text{measurable by type}$$

$$\int^\mu (\Delta n) t$$

$$\text{for } \int^\mu (\lambda x, t)$$

for  $\frac{a_n}{\Delta n} \rightarrow 0$ , e.g.  $\Delta n = \frac{1}{2^n}$   $a_n = n$ .

resolution

The unrestricted Giry      Strong Monad

Dirac:

$$\delta: X \rightarrow Gx$$

$$x \mapsto \lambda A. \begin{cases} x \in A : 1 \\ x \notin A : 0 \end{cases}$$

Unlike the unrestricted  
Giry on Meas.

but: non-commutative

Kleisli extension/Kock integral:

$$\oint: Gx \times Gp^X \rightarrow Gp$$

$$\oint \mu f := \lambda A. \int \mu(dx) f(x; A)$$

(Fubini Rule,  
just like in  
Meas)

Fubini-Tonelli; fails

$$\# \in G/R$$

$$\# E := \begin{cases} E \text{ finite:} & |E| \\ \text{o.w.:} & \infty \end{cases}$$

$$\lambda \in G/R$$

lebesgue

$$k: R \times R \rightarrow W \cong G/1$$

$$\int \#(\lambda r) \underbrace{\int \lambda(x) k(x,y)}_{y: R + \{<\} \mapsto \lambda(y) \cdot 1 + \lambda(y)^c \cdot 0 = 0} = \int \# \underline{0} = \underline{0} \stackrel{?}{=} 0$$

$k(x,y) := [x=y]$

#

$$\int \lambda(dx) \underbrace{\int \#(dr) k(x,y)}_{x: R + \{<\} \mapsto \# \{x\} \cdot 1 + 0 = 1} = \int \lambda(x) \delta_x \stackrel{?}{=} \infty$$

## Randomisable measures monad

$D \rightarrow G$

$$\lambda A. \int_{\text{Dom } \alpha} \lambda (\text{Dom } \alpha)$$

$$LDX := \left\{ \lambda \alpha \mid \alpha: \mathbb{R} \rightarrow X \right\}$$

$$R_{Dx} := \left\{ \lambda x. \lambda \alpha_x \mid \alpha: \mathbb{R} \times \mathbb{R} \rightarrow X \right\}$$

$$\delta: x \rightarrow Dx \quad \oint: D^{\Gamma \times} (DX) \rightarrow Dx \quad \text{lift along } D \rightarrow G.$$

$D$  validates our measure axioms including Fubini-Tonelli:  
 $\mu \in DX, \nu \in DY$

$$\oint \mu(dx) \oint \nu(dy) \delta_{(x,y)} = \oint \nu(dy) \oint \mu(dx) \delta_{(x,y)} =: \mu \otimes \nu$$

Thm: For  $S$ ,  $\text{PS}, D_{\leq 1} S, D_{<\infty} S \in \text{Sbs}$   
and agree with their Counterparts on  $\text{Meas}$ .

$$DS_S = \{ \mu \mid \mu \text{ } S\text{-finite} \} \quad \text{See [Staton'16]}$$

$$R_{DS} = \{ K: R \rightarrow G0 \mid K \text{ } S\text{-finite kernel} \}$$

Open: Is there a counterpart to  $D$  in  $\text{Meas}$ ?

More modestly, is  $DS \in \text{Sbs}$ ?

(Hypothesis: **No**)

## Distribution Submonads

A measure space

$$\Omega = (\Omega, \mu)$$

is a gbs  $\Omega$  with  
 $\mu \in D_X$ .

Similarly:-  
finite measure space  
- (sub) probability space.

$$P_X := \left\{ \mu \in D_X \mid \mu X = 1 \right\}$$

$$P_{\leq 1} X := \left\{ \mu \in D_X \mid \mu X \leq 1 \right\}$$

$$P_{<\infty} X := \left\{ \mu \in D_X \mid \mu X < \infty \right\}$$

$$D_X$$

## Full model

$$\begin{aligned} \text{type : Obs} & \quad W := [0, \infty] \quad \mathcal{B}^X \cong \mathcal{B}^X \\ DX := & \left( \{\lambda_\alpha \mid \alpha : R \rightarrow X\}, \{\lambda_r, \lambda_{\alpha(r,-)} \mid \alpha : R \times R \rightarrow X\} \right) \\ P_X := & \left\{ \mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1 \right\} \\ \underset{\mu}{\text{Ce}}[E] := & \mu E \quad \delta_x := E \mapsto \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases} \\ \oint \mu k := & \lambda E. \int \mu(\lambda x) k(x; E) \end{aligned}$$

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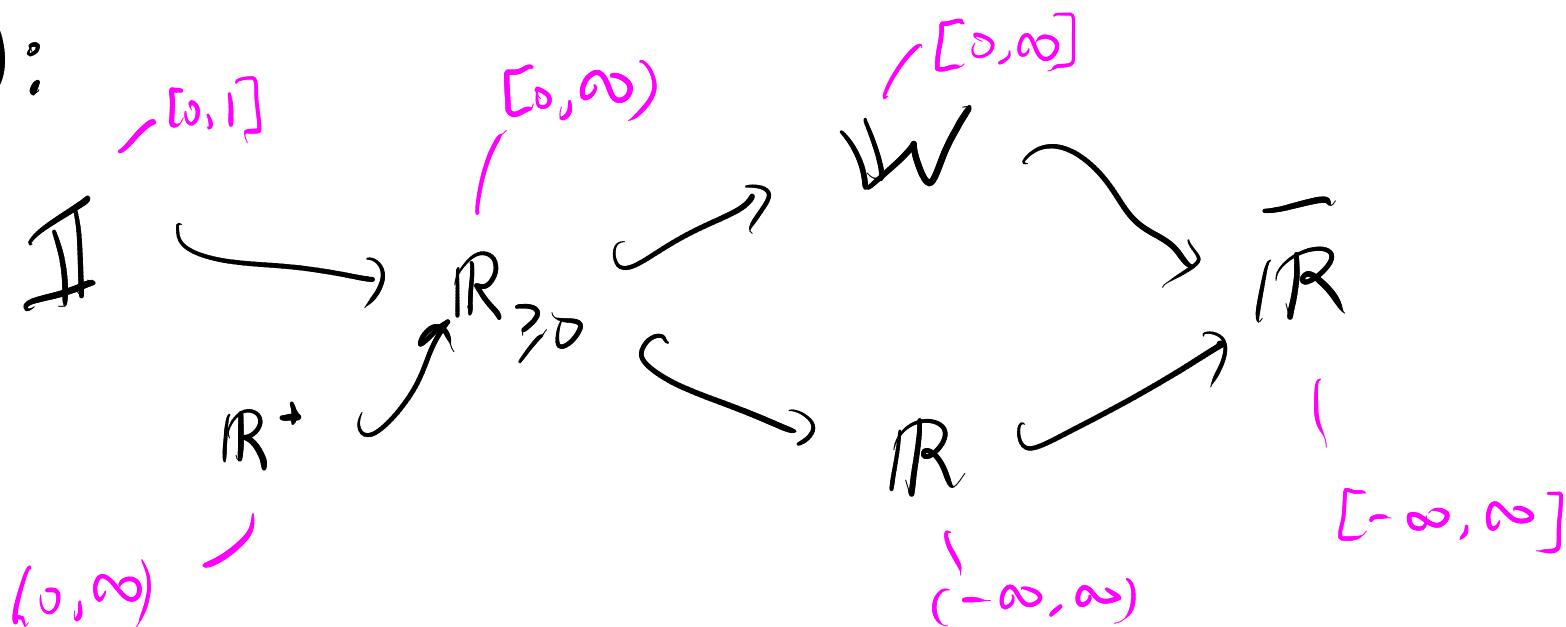
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Random variable:  $\xi : \Omega \rightarrow \mathbb{H} \subseteq \bar{\mathbb{R}}$

$\mathbb{H}:$



-  $\Omega$  is a space

-  $\mathbb{R}^\Omega$  measurable vector space:

$$\alpha \xi + \zeta := \lambda \omega \cdot \alpha \cdot \xi \omega + \zeta \omega$$

-  $W^\Omega$  measurable  $\sigma$ -Semi-module  
for  $W$ :

$$\sum_{n=0}^{\infty} \alpha_n \xi_n := \lambda \omega \cdot \sum_{n=0}^{\infty} \alpha_n \cdot \xi_n$$

$$\Pr_\lambda : P_{\Omega} \times B_{\Omega} \rightarrow \mathbb{W}$$

$$\Pr_\lambda A := \text{eval}(\lambda, A) = \lambda A$$

Probability Space  $\mathcal{R} = (\Omega, \lambda_\Omega)$

$P : P_{\Omega} \vdash$  "  $P_\lambda$  holds  $\lambda(\omega)$ -almost surely"  
for some  $Q \subseteq \Omega$ ,  $P \models Q$ ,  $[- \in Q] \cdot \lambda = \lambda$

Example  $(\xi, \zeta \in \Theta^\Omega)$

$\xi = \zeta$  a.s. when  $\Pr_{w \sim \lambda} [\xi_w \neq \zeta_w] = 0$

# Integrating Random Variables (as discretely)

$(-)_{+}, (-)_{-} : \bar{\mathbb{R}}^n \rightarrow \mathbb{W}^n$  in Qbs!

$$\xi_{+} := \max(\xi, 0) \quad \xi_{-} := \max(-\xi, 0)$$

$$\text{So: } \xi = \xi_{+} - \xi_{-}$$

$$\int : P\mathcal{R} \times \mathbb{W}^n \longrightarrow \mathbb{W} \quad \begin{cases} \text{respects} \\ \text{a.s. equality:} \end{cases}$$

$$\int \lambda \xi := \int \lambda \xi_{+} - \int \lambda \xi_{-} \quad \xi_{+} = \xi \text{ (a.s.)} \\ \Rightarrow \int \lambda \xi = \int \xi.$$

## Example

$$\lambda: P\Omega \vdash ASConverg(\bar{\mathbb{R}})^{\omega} : B(\bar{\mathbb{R}}^{N \times \omega})$$
$$:= \left\{ \vec{\zeta} \in \bar{\mathbb{R}}^{N \times \omega} \mid \Pr_{w \sim \lambda} [\lim \vec{\zeta}_n w \neq \perp] \right\}$$

So;

$$\lim^{\text{as}}: \bar{\mathbb{R}}^{N \times \omega} \rightarrow \bar{\mathbb{R}}^\Omega$$
$$\text{Dom } \lim^{\text{as}} := ASConverg(\bar{\mathbb{R}})^\omega$$

$$\lim^{\text{as}} \vec{\zeta} := \text{a.s. limsup}_{n \rightarrow \infty} f_n w$$



$\lim^{\text{as}}$  respects a.s. equality.

Thm (monotone convergence):

Let  $\vec{\xi} \in \mathbb{W}^{N \times n}$   $\lambda$ -a.s. monotone.

$$\xi = \lim_{n \rightarrow \infty} \xi_n \quad (\text{a.s.})$$



$$\int \lambda \xi = \lim_{n \rightarrow \infty} \int \lambda \xi_n$$

Lebesgue Space  $\left( \Omega \text{ Prob. Space}, P \in [1, \infty) \right)$

$P: [1, \infty), \lambda: P\Omega \vdash L_{(\Omega, \lambda)}^P: B(\mathbb{R}^\Omega)$

$$:= \left\{ \xi \in \mathbb{R}^\Omega \mid \int \lambda |\xi|^P < \infty \right\} \hookrightarrow \mathbb{R}^\Omega$$

Ensemble  $L_\Omega := \prod_{\lambda \in P\Omega} L_{(\Omega, \lambda)}^P$

$$L \quad P \leq q \Rightarrow L_\Omega^P \supseteq L_\Omega^q$$

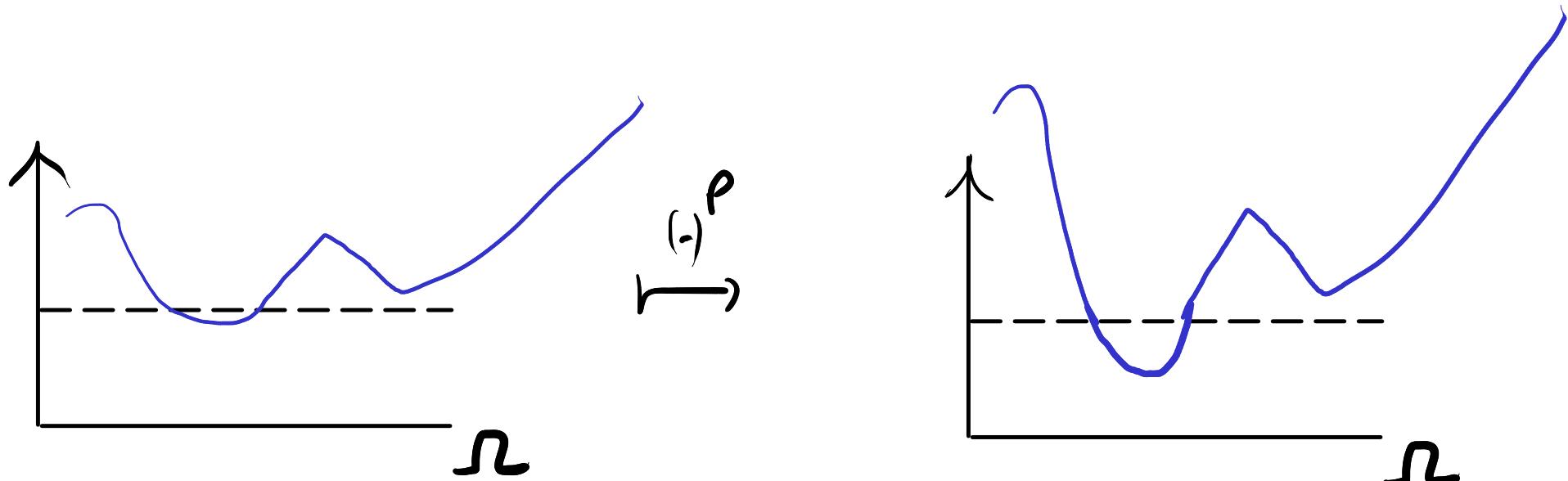
$L^p$  semi norms

$$\| - \| : \prod_{P,\lambda} L_{(2,\lambda)}^p \rightarrow \mathbb{R}_{\geq 0} \quad \|\xi\|_p := \sqrt[p]{\int \lambda |\xi|^p}$$

$L^2$  inner product

$$\langle \cdot, \cdot \rangle : \prod_{P,\lambda} L_{(2,\lambda)}^p \times L_{(2,\lambda)}^p \rightarrow \mathbb{R}$$

$$\langle \xi, \eta \rangle := \int \lambda \xi \eta$$



## Statistics

### Expectation

$$\mathbb{E} : \prod_{\lambda} \mathcal{L}^1 \rightarrow \mathbb{R}$$

$$\mathbb{E}_{\lambda} \xi := \int_{\lambda} \xi$$

### Covariance and Correlation

$$\text{Cov}, \text{Corr} : \prod_{\lambda} \mathcal{L}^2 \rightarrow \mathbb{R}$$

$$\text{Cov}(\xi, \zeta) := \langle \xi - \mathbb{E} \xi, \zeta - \mathbb{E} \zeta \rangle$$

$$\text{Corr}(\xi, \zeta) := \frac{\langle \xi, \zeta \rangle}{\|\xi\|_2 \cdot \|\zeta\|_2} = \cos(\text{angle}(\xi, \zeta))$$

## Sequential limits

$P: [1, \infty)$ ,  $\lambda: P X \vdash$  Cauchy  $L_{(R,\lambda)}^P: B(L_{(R,\lambda)}^P)^{IN}$

$$:= \left\{ \vec{\Sigma} \mid \forall \varepsilon \in \mathbb{Q}^+ \exists \kappa \in \mathbb{N} \quad \forall m, n \geq \kappa, \quad \| \Sigma_{n+m} - \Sigma_{n+m} \|_P < \varepsilon \right\}$$

Thm:  $L_{(R,\lambda)}^P$  is Cauchy-complete

$\lim: \text{Cauchy } L_{(R,\lambda)}^P \rightarrow L^P$  (convergence in mean)

Why?

1. Every Cauchy sequence has an a.s. converging subseq.
2. We can find it measurable

## Example

Theorem (dominated convergence)

For  $\tilde{z}_n, z \in L^1$  s.t.  $\tilde{z}_n \leq z$  a.s.:

1.  $\lim^{\text{as}} \tilde{z} \in L^1$

2.  $\lim^1 \tilde{z} = \lim^{\text{as}} \tilde{z}$

3.  $\lim_{n \rightarrow \infty} \int \tilde{z}_n = \int \lim_{n \rightarrow \infty} \tilde{z}_n$

## Separability

Def:  $L^P$  separable: has countable dense subset

Fact: Separability is property of  $\lambda_2$ :

TFAE:

- $\exists p \geq 1$ .  $L^p$  separable
- $\forall p \geq 1$ .  $L^p$  separable

Measurable separability in  $I \hookrightarrow P\Omega \times [1, \infty)$

$$\vec{\beta} : \prod_{(\lambda, p) \in I} L^p_{(\Omega, \lambda)} \xrightarrow{IN} \text{S.t.}$$

$$\left\{ \vec{\beta}_n^{(p)} \mid n \in \mathbb{N} \right\} \text{ dense in } L^p_{(\Omega, \lambda)}$$

Prop. - Every SBS  $S$  measurable separable in

$$PS \times [1, \infty)$$

-  $I \hookrightarrow P\Omega \times \{2\}$  measurably separable

$$\Rightarrow \exists \vec{\beta} \in \prod_{\lambda \in I} L^2_{(\Omega, \lambda)} \text{ Orthonormal System}$$

$$\begin{aligned} \langle \beta_n, \beta_m \rangle &= 0 \\ \|\beta_n\|_2 &= 1 \\ (\beta_n) &\text{ dense} \end{aligned}$$

Example

Let  $S \subset L^2$  closed Vector Subspace.

Orthogonal decomposition linear in fact.

$$\langle P, P^\perp \rangle : L^2 \rightarrow S \times S^\perp$$

When  $S$  is separable with orthonormal system  $\beta$

We have a measurable version of

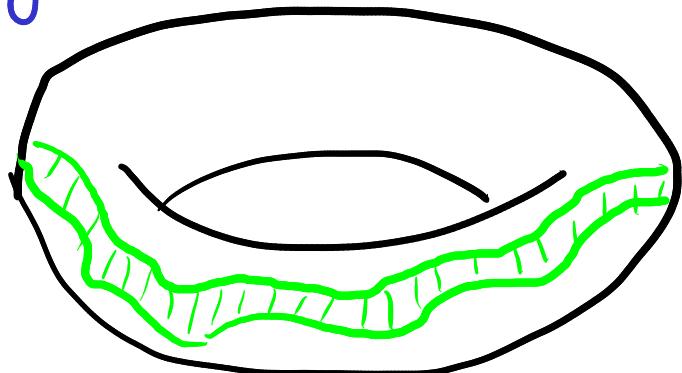
$$\langle P, P^\perp \rangle : L^2 \rightarrow S \times S^\perp$$

$$P\xi := \sum_{n=0}^{\infty} \langle \xi, \beta_n \rangle \beta_n$$

$$P^\perp := I_d - P$$

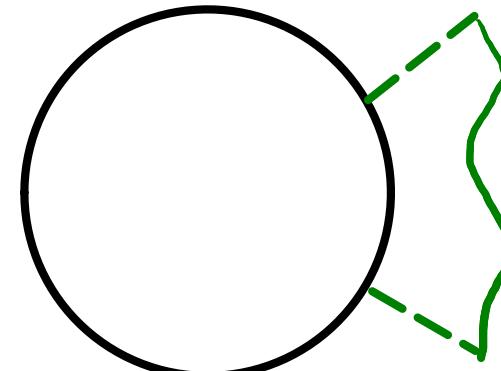
# Kolmogorov's Conditional Expectation

↳ ground truth space



$H$   
observation

( $H$ ) Sample space



↳ conditional expectation

$$\mathbb{E}[\xi | H = -]$$

Observed  
statistic

$\xi$   
Statistic  
of interest

R

# Kolmogorov's Conditional Expectation

A Conditional expectation

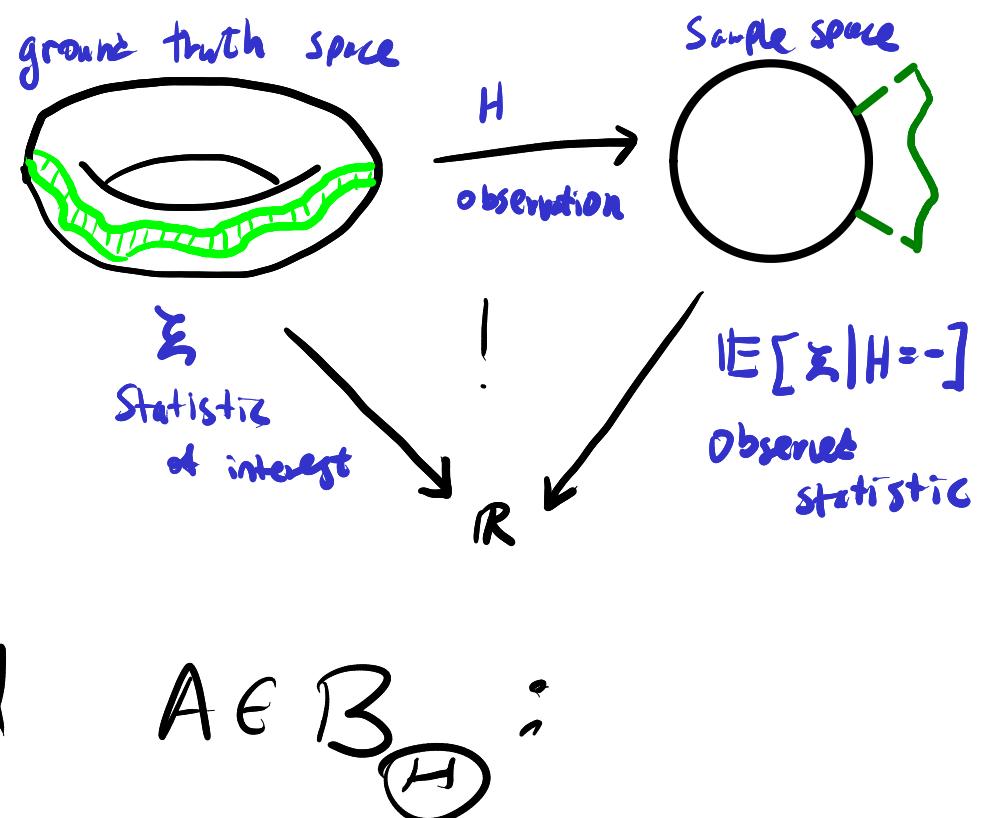
of  $\xi \in \mathcal{L}_\Omega$  wrt

$H: \Omega \rightarrow \mathbb{H}$  is

$\xi \in \mathcal{L}_{(H)}$  s.t. for all  $A \in \mathcal{B}_{(H)}$ :

$$\int_A \mu \xi = \int_{H^{-1}[A]} \lambda \xi$$

where  $\mu := \lambda_H$

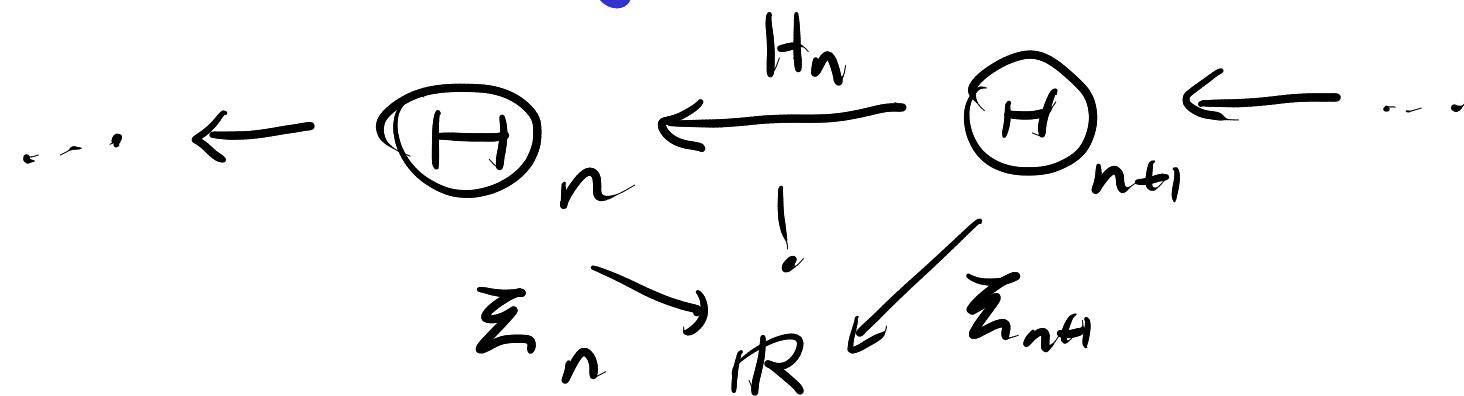


# Conditional expectations

1. unique a.s.

2. fundamental to Modern Probability, e.g.:

a Martingale



$$\text{S.t. } \xi_n = \mathbb{E}[\xi_{n+1} | H_n = -]$$

Theorem (Existence)

- $\exists \mathbb{E}[-|H=-] : \int'_{L(\Omega, \lambda)} \rightarrow \int'_{L(\mathbb{D}, \mu)}$
- When  $(\Omega, \lambda)$  is Separable  
 $\mathbb{E}[-|H=-] : \int'_{L(\Omega, \lambda)} \rightarrow \int'_{L(\mathbb{D}, \mu)}$
- When  $H$  is  $\mathcal{I}$ -measurably separable  
 $\mathbb{E}[-|-\cdot-\cdot] : \prod_{\substack{H \in \mathbb{D} \\ \lambda \in H^{\perp}[\mathcal{I}]}} \int'_{L(\Omega, \lambda)} \rightarrow \int'_{L(\mathbb{D}, \mu)}$

## Plan:

- 1) Type-driven Probability: discrete case (Mon + Tue)
- 2) Borel sets & measurable spaces (Wed)
- 3) Quasi Borel spaces (Wed) Simple type structure (Thu)
- 4) Dependent type structure & standard Borel spaces (Thu)
- 5) Integration & random variables (Fri)



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## Discrete model

$$\text{type} : \text{set} \quad \mathbb{W} := [0, \infty] \quad \mathcal{B}X := \mathcal{P}X$$

$$DX := \{\mu : X \rightarrow \mathbb{W} \mid \text{Supp } \mu \text{ countable}\}$$

$$PX := \{\mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1\}$$

$$\underset{\mu}{\text{Ce}}[E] := \sum_{x \in E} \mu_x \quad \delta_x := \lambda x'. \begin{cases} x = x': 0 \\ x \neq x': 1 \end{cases}$$

$$\phi \mu k := \lambda x. \sum_{m \in \Gamma} \mu^m \cdot k(m; x)$$

## Full model

$$\begin{aligned} \text{type : Qbs} \quad w := [0, \infty] \quad \mathcal{B}^X &\cong \mathcal{B}^X \\ DX := (\{\lambda_\alpha \mid \alpha : R \rightarrow X\}, \{\lambda_r, \lambda_{\alpha(r,-)} \mid \alpha : R \times R \rightarrow X\}) \\ P_X := \{\mu \in DX \mid \underset{\mu}{C_e}[X] = 1\} \\ C_e[E] := \mu E \quad \delta_x := E \mapsto \begin{cases} x \in E : 1 \\ x \notin E : 0 \end{cases} \\ \oint \mu k := \lambda E. \int \mu(\lambda) k(x; E) \end{aligned}$$