

Higher-order building blocks for statistical modelling

Ohad Kammar
University of Edinburgh

PPS-PIHOC-DIAPASoN Workshop
17 February 2021



THE UNIVERSITY of EDINBURGH

informatics **lfcs**

Laboratory for Foundations
of Computer Science



supported by:



THE ROYAL
SOCIETY

The
Alan Turing
Institute

Facebook Research

Theorem (Aumann)

$S = 2, \mathbb{N}, \mathbb{R}, \dots$

No σ -algebra on $\mathbf{Meas}(\mathbb{R}, S)$ makes evaluation measurable:

$$\text{eval} : \mathbf{Meas}(\mathbb{R}, S) \times \mathbb{R} \rightarrow S$$

\implies bad fit for higher-order programming semantics:

\mathbf{Meas} is not Cartesian closed

higher-order measure theory

Bacci-et-al'18
Dahlqvist-Kozen'20
Ehrhard-Pagani-Tasson'18

Recent breakthroughs:

\times, \rightarrow
alongside
 Pr, \int

dependent
(co)products

This talk:

\sqcup, Π
Type, Prop, Dec

higher-order measure theory
with **quasi-Borel spaces**

Bacci-et-al'18
Dahlqvist-Kozen'20
Ehrhard-Pagani-Tasson'18

Recent breakthroughs:

\times, \rightarrow
alongside
 Pr, \int

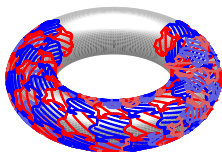
dependent
(co)products

This talk:

\sqcup, Π
Type, Prop, Dec

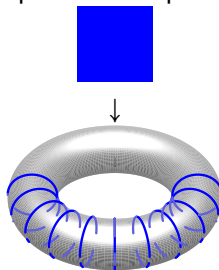
Intuition

measurable space



axiomatise
measurable events

quasi Borel space

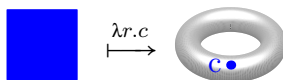


axiomatise
random elements

quasi-Borel space $X = (\iota X, \mathcal{R}_X)$

random element $\alpha \in \mathcal{R}_X \subseteq \iota X^{\mathbb{R}}$ axioms:

determinism. elements are random elements:

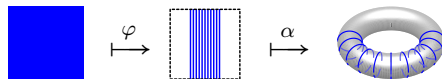
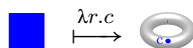


quasi-Borel space $X = (\iota X, \mathcal{R}_X)$

random element $\alpha \in \mathcal{R}_X \subseteq \iota X^{\mathbb{R}}$ axioms:

determinism.

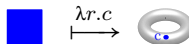
precomposition. $(\varphi \in \mathbf{Meas}(\mathbb{R}, \mathbb{R}))$



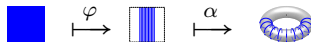
quasi-Borel space $X = (\iota X, \mathcal{R}_X)$

random element $\alpha \in \mathcal{R}_X \subseteq \iota X^{\mathbb{R}}$ axioms:

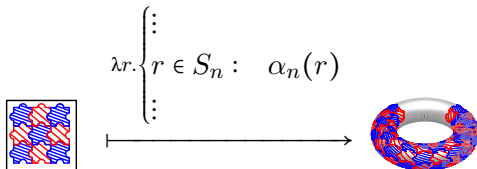
determinism.



precomposition. $(\varphi \in \mathbf{Meas}(\mathbb{R}, \mathbb{R}))$



recombination $(\mathbb{R} = \biguplus_{i=0}^{\infty} S_n)$



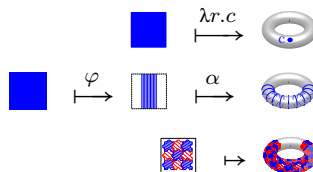
quasi-Borel space $X = (\llbracket X \rrbracket, \mathcal{R}_X)$

random element $\alpha \in \mathcal{R}_X \subseteq \llbracket X \rrbracket^{\mathbb{R}}$ axioms:

determinism.

precomposition. $(\varphi \in \mathbf{Meas}(\mathbb{R}, \mathbb{R}))$

recombination $(\mathbb{R} = \biguplus_{i=0}^{\infty} S_n)$



Morphisms $f : X \rightarrow Y$

functions $\llbracket f \rrbracket : \llbracket X \rrbracket \rightarrow \llbracket Y \rrbracket$:

$$\alpha \in \mathcal{R}_X \implies f \circ \alpha \in \mathcal{R}_Y$$

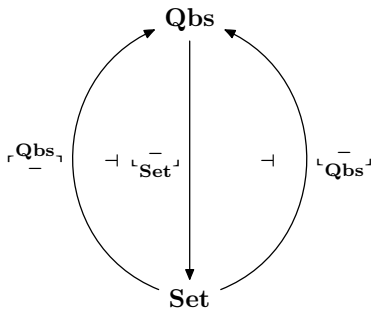
Free and cofree qbs

set A :

- ▶ free: $\mathcal{R}_{\ulcorner A \urcorner}^{\text{Qbs}} = \sigma\text{-simple functions:}$

$$\lambda r. \begin{cases} \vdots \\ r \in S_n : a_n \\ \vdots \end{cases}$$

- ▶ cofree: $\mathcal{R}_{\ulcorner \text{Qbs} \urcorner}^A = \text{all functions}$



qbses for:

$$2, \mathbb{N}, \mathbb{Q}$$

set-theoretic escape hatch

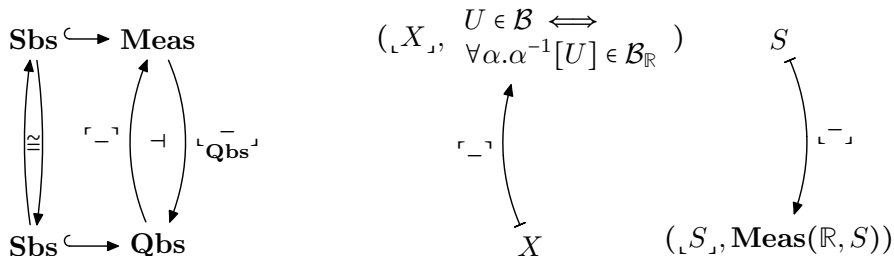
$\lrcorner \ulcorner \text{Set} \urcorner : \text{Qbs} \rightarrow \text{Set}$ generates limits and colimits

preserves
lifts
reflects



Using measure theory

Measurable space are carried by qbses:



Recover qbses for:

$$\mathbb{R}, \mathbb{W} := [0, \infty], \mathbb{I} := [0, 1]$$

Conservative extension for standard Borel spaces

Measure-theoretic escape hatch

Simple types

Simple products

Correlated random elements:

$$\mathcal{R}_{X \times Y} \xleftarrow[\cong]{(-,-)} \mathcal{R}_X \times \mathcal{R}_Y$$

Simple types

Simple products

Correlated random elements:

$$\mathcal{R}_{X \times Y} \xleftarrow[\cong]{(-,-)} \mathcal{R}_X \times \mathcal{R}_Y$$

Simple coproducts

Recombinations:

$$\alpha \in \mathcal{R}_{\coprod_{i \in \mathcal{I}} X_i} \iff \alpha = \lambda r. \begin{cases} \vdots \\ r \in S_n : (i_n, \alpha_n r) \\ \vdots \end{cases} \quad (\mathbb{R} = \biguplus_{n=0}^{\infty}, \alpha_n \in X_{i_n})$$

Simple types

Simple products

Correlated random elements:

$$\mathcal{R}_{X \times Y} \xleftarrow[\cong]{(-,-)} \mathcal{R}_X \times \mathcal{R}_Y$$

Simple coproducts

Recombinations:

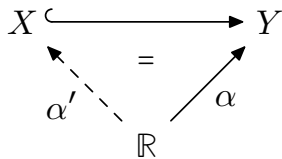
$$\alpha \in \mathcal{R}_{\coprod_{i \in \mathcal{I}} X_i} \iff \alpha = \lambda r. \begin{cases} \vdots \\ r \in S_n : (i_n, \alpha_n r) \\ \vdots \end{cases} \quad \left(\mathbb{R} = \biguplus_{n=0}^{\infty}, \alpha_n \in X_{i_n} \right)$$

Simple function spaces

$$\llbracket Y^X \rrbracket = \mathbf{Qbs}(X, Y) \quad \mathcal{R}_{Y^X} \xrightarrow[\cong]{\text{uncurry}} \mathbf{Qbs}(\mathbb{R} \times X, Y)$$

Random element space: $\mathcal{R}_X := X^{\mathbb{R}}$

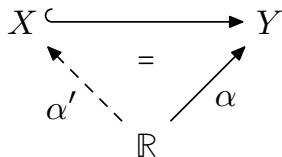
Subspaces



m injective and $\mathcal{R}_X = (m \circ)^{-1}[\mathcal{R}_Y]$

$\mathcal{Q} := \mathcal{Q}_{\text{bs}}^2$ subspace classifier

Subspaces



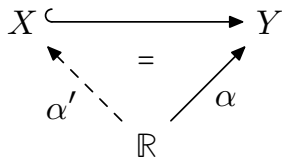
m injective and $\mathcal{R}_X = (m \circ)^{-1}[\mathcal{R}_Y]$

$\mathbb{Q} := \mathcal{L}_{\mathbf{Qbs}}^2$ subspace classifier

Example

- Use $\llbracket \mathbf{Prop} \rrbracket := \mathbb{Q}$ for reasoning/axiomatics [Sato et al.'19].

Subspaces



m injective and $\mathcal{R}_X = (m \circ)^{-1}[\mathcal{R}_Y]$

$\mathcal{Q} := \mathcal{Q}_{\text{bs}}^2$ subspace classifier

Example

- ▶ Use $\llbracket \mathbf{Prop} \rrbracket := \mathcal{Q}$ for reasoning/axiomatics [Sato et al.'19].
- ▶ Differentiation:

$$D_1\mathbb{R} := \{f : \mathbb{R} \rightarrow \mathbb{R} \mid f \text{ differentiable everywhere}\} \hookrightarrow \mathbb{R}^{\mathbb{R}}$$

$$\frac{d}{d} : D_1 \rightarrow \mathbb{R}^{\mathbb{R}}$$

$m : X \hookrightarrow Y$ when $m : X \hookrightarrow Y$ and $[- \in X] \in \Omega^Y$
factors through $2^Y \rightarrow \Omega^Y$

$m : X \hookrightarrow Y$ when $m : X \hookrightarrow Y$ and $[- \in X] \in \Omega^Y$
factors through $2^Y \rightarrow \Omega^Y$

Example

- higher-order Qbs-internal σ -algebra:

$$-^C : \mathcal{B}_X \rightarrow \mathcal{B}_X \quad \bigcap_{n=0}^{\infty} : \mathcal{B}_X^{\mathbb{N}} \rightarrow \mathcal{B}_X \quad -^{-1}[-] : Y^X \times \mathcal{B}_Y \rightarrow \mathcal{B}_X$$

$m : X \hookrightarrow Y$ when $m : X \hookrightarrow Y$ and $[- \in X] \in \Omega^Y$
factors through $2^Y \rightarrow \Omega^Y$

Example

- ▶ higher-order **Qbs**-internal σ -algebra:

$$-^C : \mathcal{B}_X \rightarrow \mathcal{B}_X \quad \bigcap_{n=0}^{\infty} : \mathcal{B}_X^{\mathbb{N}} \rightarrow \mathcal{B}_X \quad -^{-1}[-] : Y^X \times \mathcal{B}_Y \rightarrow \mathcal{B}_X$$

- ▶ Non-**Qbs**-morphisms:

$$\exists : \mathcal{B}_{X \times Y} \rightarrow \mathcal{B}_X \quad [- = \emptyset] : \mathcal{B}_X \rightarrow 2 \quad [- \sqsubseteq -] : \mathcal{B}_X^2 \rightarrow 2$$

Borel subspaces

$m : X \hookrightarrow Y$ when $m : X \hookrightarrow Y$ and $[- \in X] \in \Omega^Y$
factors through $2^Y \rightarrow \Omega^Y$

Example

- ▶ higher-order **Qbs**-internal σ -algebra:

$$-^C : \mathcal{B}_X \rightarrow \mathcal{B}_X \quad \bigcap_{n=0}^{\infty} : \mathcal{B}_X^{\mathbb{N}} \rightarrow \mathcal{B}_X \quad -^{-1}[-] : Y^X \times \mathcal{B}_Y \rightarrow \mathcal{B}_X$$

- ▶ Non-**Qbs**-morphisms:

$$\exists : \mathcal{B}_{X \times Y} \rightarrow \mathcal{B}_X \quad [- = \emptyset] : \mathcal{B}_X \rightarrow 2 \quad [- \sqsubseteq -] : \mathcal{B}_X^2 \rightarrow 2$$

- ▶ $\mathcal{B}_{\mathcal{B}_X}$: Borel-on-Borel sets [Sabok-Staton-Stein-Wolman'21]

- Unrestricted Giry:

$$\mathcal{G}_1 X := \left\{ \mu : \mathcal{B}_X^{\text{Qbs}} \rightarrow \mathbb{W} \mid \mu \text{ is a measure on } \mathcal{B}_X^{\text{Qbs}} \right\}$$

$$\mathcal{R}_{G_1 X} := \left\{ k : \mathbb{R} \times \mathcal{B}_X^{\text{Qbs}} \rightarrow \mathbb{W} \mid k \text{ is a kernel} \right\}$$

Analogous to the measure-theoretic Giry

But: careful to evaluate only σ -simple random Borel sets

- Unrestricted Giry:

$$\llbracket G_1 X \rrbracket := \left\{ \mu : \llbracket \mathcal{B}_X^{\text{Qbs}} \rrbracket \rightarrow \mathbb{W} \mid \text{measure on } \llbracket X^{\text{Meas}} \rrbracket \right\}$$

$$\mathcal{R}_{G_1 X} := \left\{ k : \mathbb{R} \times \llbracket \mathcal{B}_X^{\text{Qbs}} \rrbracket \rightarrow \mathbb{W} \mid k \text{ is a kernel} \right\}$$

Analogous to the measure-theoretic Giry

But: careful to evaluate only σ -simple random Borel sets

- Following can evaluate any random Borel set:

$$\llbracket G_2 X \rrbracket := \left\{ \mu : \mathcal{B}_X \rightarrow \mathbb{W} \mid \text{measure on } \llbracket X^{\text{Meas}} \rrbracket \right\}$$

$$\mathcal{R}_{G_2 X} := \{ k : \mathbb{R} \times \mathcal{B}_X \rightarrow \mathbb{W} \mid k \text{ is a kernel} \}$$

$$\mathcal{M}X := \{v_*\lambda_\Omega | v : \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \text{ a } \sigma\text{-finite standard measure space}\}$$

$$\mathcal{R}_{MX} := \{v_* \circ k | k : \mathbb{R} \times \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \dots \text{ ditto } \dots\}$$

$\mathcal{M}X := \{v_* \lambda_\Omega \mid v : \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \text{ a } \sigma\text{-finite standard measure space}\}$

$\mathcal{R}_{MX} := \{v_* \circ k \mid k : \mathbb{R} \times \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \dots \text{ ditto } \dots\}$

- ▶ For standard Borel spaces S, T :
 - ▶ $\mathcal{M}S$ are the s-finite measures
 - ▶ $(\mathcal{M}S)^T$ are the s-finite kernels

Fully-definable semantic domain for first-order Monte Carlo models [Staton'17]

$\mathbf{M}X := \{v_* \lambda_\Omega \mid v : \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \text{ a } \sigma\text{-finite standard measure space}\}$

$\mathcal{R}_{MX} := \{v_* \circ k \mid k : \mathbb{R} \times \Omega \rightarrow X, \Omega \hookrightarrow \mathbb{R} \dots \text{ ditto } \dots\}$

- ▶ For standard Borel spaces S, T :
 - ▶ $\mathbf{M}S$ are the s-finite measures
 - ▶ $(\mathbf{M}S)^T$ are the s-finite kernels

Fully-definable semantic domain for first-order Monte Carlo models [Staton'17]

- ▶ Integration is commutative
- ▶ Models synthetic measure theory [Kock'12]
- ▶ Probabilistic fragment $\mathbf{P}X := \{\mu \mid \mu X = 1\} \hookrightarrow \mathbf{M}X$ with de Finetti's theorem [Heunen et al.'17]
- ▶ $\mathbf{P}X$ also models name generation [Sabok et al.'21]

Syntactic spaces and recursive domains

\mathbf{Qbs} is locally presentable

\implies initial algebra semantics for inductive types

Syntactic spaces and recursive domains

\mathbf{Qbs} is locally presentable

\implies initial algebra semantics for inductive types

Example

- ▶ Syntactic spaces for operational semantics
- ▶ Meta-programming data structures for Monte Carlo inference [Ścibior et al.'18, Lew et al.'20]
- ▶ Extends to **recursive** types with domain theory [Vákár-Kammar-Staton'19]
- ▶ Opportunity: abstract syntax with binding [Fiore-Plotkin-Turi'99]

Monadic operational semantics

[Dal Lago et al.'17, Gavazzo'19, Vákár et al.'19]

$$\frac{k_1(t) w_1 \quad k_2(t, w_1) w_2 \quad \dots \quad k_n(t, w_1, \dots, w_n) v}{l(t) f(t, w_1, \dots, w_n, v)}$$

means

$$\begin{aligned} l(t) &:= k_1(t) && \gg= \lambda w_1. \\ &k_2(t, w_1) && \gg= \lambda w_2. \dots \\ &k_n(t, w_1, \dots, w_{n-1}) && \gg= \lambda v. \\ &\delta_{f(t, w_1, \dots, w_n, v)} \end{aligned}$$

and $l(t) := k_1(t) \quad l(t) := k_2(t)$ means $l(t) := k_1(t) + k_2(t)$

Monadic operational semantics

[Dal Lago et al.'17, Gavazzo'19, Vákár et al.'19]

$$\frac{k_1(t) w_1 \quad k_2(t, w_1) w_2 \quad \dots \quad k_n(t, w_1, \dots, w_n) v}{l(t) f(t, w_1, \dots, w_n, v)}$$

means

$$\begin{aligned} l(t) &:= k_1(t) && \gg= \lambda w_1. \\ & \quad k_2(t, w_1) && \gg= \lambda w_2. \dots \\ & \quad k_n(t, w_1, \dots, w_{n-1}) && \gg= \lambda v. \\ & \quad \delta_{f(t, w_1, \dots, w_n, v)} \end{aligned}$$

and $l(t) := k_1(t) \quad l(t) := k_2(t)$ means $l(t) := k_1(t) + k_2(t)$

Example

$$\frac{t \Downarrow_n \underline{0} \quad s_1 \Downarrow_n v}{\text{match } t \text{ with } \{\underline{0} \rightarrow s_1 \mid _ \rightarrow s_2\} \Downarrow_n v}$$

$$\frac{t \Downarrow_n \underline{r} \quad s_2 \Downarrow_n v}{\text{match } t \text{ with } \{\underline{0} \rightarrow s_1 \mid _ \rightarrow s_2\} \Downarrow_n v} (r \neq 0)$$



Refinement types:

$$\frac{B : X \rightarrow \Omega^Y}{\prod_{x:X} Bx, \coprod_{x:X} Bx} \quad \prod_{x:X} Bx := \{f : X \rightarrow Y \mid \forall x \in X. fx \in Bx\} \hookrightarrow Y^X$$

$$\coprod_{x:X} Bx := \{(x, y) \in X \times Y \mid y \in Bx\} \hookrightarrow X \times Y$$

Example

Lebesgue spaces and modes of convergence:

$$\mathcal{L}_-^p \in \prod_{\substack{\lambda \in \mathbf{P}\Omega \\ p \in \mathbb{R}}} \left\{ f : \Omega \rightarrow [-\infty, \infty] \mid \int d\omega |f \omega|^p < \infty \right\}$$

- ▶ Observation: **S**bs closed under dependent pairs
- ▶ To get dependent types, want “good” universe **Type**

Monte Carlo inference

$$\text{model} = \rho \odot \lambda := \lambda \varphi. \int \lambda(d\omega) \rho(\omega) \cdot \varphi(\omega)$$

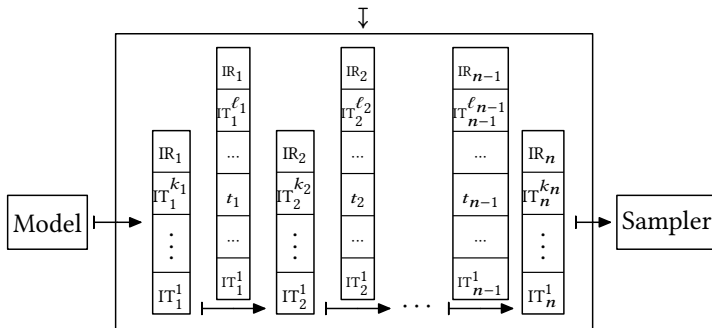
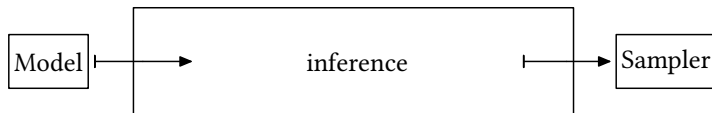
Diagram illustrating the components of the model definition:

- likelihood** (pink arrow pointing to ρ)
- prior** (pink arrow pointing to λ)

Defined programmatically:

sample : $\mathbf{M}\mathbb{R}$ #Uniform distribution on $[0, 1]$
score : $\mathbb{R} \rightarrow \mathbf{M}\mathbb{1}$

Modular inference [Ścibior et al.'18a+b, Lew et al.'20]



cf. inference with handlers [Bingham et al.'19] and ongoing



Thank you!

Exact conditioning

