

# Foundations for Type-Driven Probabilistic Modelling

Ohad Kammar  
University of Edinburgh

Laboratory for Software Science  
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Chair of Programming Languages  
Institute of Computer Science  
University of Tartu



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# Computational golden era

logic-rich & type-rich computation

statistical computation

# Computational golden era

## logic-rich & type-rich computation

- ▶ Expressive type systems: Haskell, OCaml, Rust, Agda, Idris
- ▶ Mechanised mathematics: Agda, Rocq, Isabelle/HOL, Lean
- ▶ Verification: SMT-powered real-world systems

## statistical computation

Generative modelling with efficient inference: Monte-Carlo simulation or gradient-based optimisation

# This course

Typed interface to probability/statistics

Every concept has:

- ▶ a type
- ▶ associated operations
- ▶ properties in terms of these operations.



Two implementations/models

course page

**discrete model**



familiar maths  
introductory

**full model**

supports discrete  
and  
continuous distributions  
same language

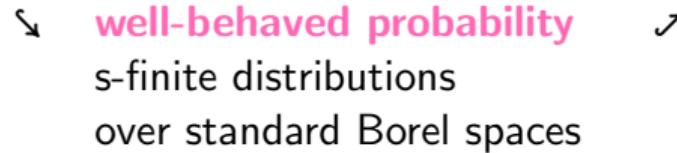
# Motivation: why foundations?

## discrete probability

countably supported distributions  
good type-structure  
**(this course)**

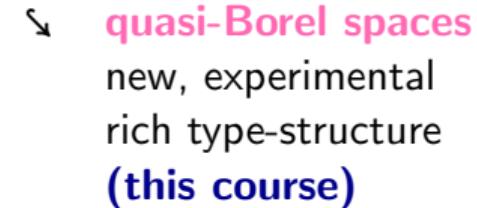
## measure theory

standard, established  
poor type-structure



## continuous probability

Lebesgue measure over  $\mathbb{R}^n$



## Takeaway

Use types to abstract away from the model

# Motivation: why types?

- ▶ **spotlights** meaningful operations

$$\int : (\text{Distribution} X) \times (\text{RandomVariable} X) \rightarrow [0, \infty]$$

- ▶ document **intent**:  
probability (**Distribution**  $X$ ) vs. density ( $X \rightarrow [0, \infty]$ ) vs. random variable
- ▶ succinctness: omit and elaborate details
- ▶ especially **formal** types, allow using theory correctly without fully understanding it

# Lecture plan

## Lecture 1: discrete model (today)

- ▶ Language of probability and distribution
- ▶ Discrete model
- ▶ Simply-typed probability
- ▶ Dependently-typed probability



course page

## Lecture 2: the full model

- ▶ Borel sets and measurable spaces
- ▶ Quasi-Borel spaces
- ▶ Type structure & standard Borel spaces
- ▶ Integration & random variables



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# Language of probability & distribution

$X$  type (=space) of **values/outcomes**

$\text{DX}$  type of **distributions/measures** over  $X$

$\text{PX} \subseteq \text{DX}$  sub-type of **probability distributions** over  $X$

$\mathcal{B}_X \subseteq \mathcal{P}X$  type of **events**: subsets we wish to measure

$\mathbb{W}$  type of **weights**: values in  $[0, \infty]$

$\int, \mathbb{E}$  Lebesgue integration and the expectation operation

Type judgements describe well-formed values/outcomes of a given type, e.g.:

$$\mu : \text{DX}, E : \mathcal{B}_X \vdash \text{Ce}_{\mu}[E] : \mathbb{W}$$

(measures weight  $\text{Ce}_{\mu}[E]$  of event  $E$  according to distribution  $\mu$ )

Propositions describe properties of well-formed values/outcomes of a given type, e.g.:

$$y_1, y_2 : Y \vdash y_1 \stackrel{Y}{=} y_2 : \text{Prop} \quad \mu : \text{PX}, E : \mathcal{B}_X \vdash \Pr_{\mu}[E] = \text{Ce}_{\mu}[E]$$

(probability of event according to probability distribution is its measure)

# Axioms for events and distributions

Empty event

$$\emptyset : \mathcal{B}_{\textcolor{red}{X}}$$

Empty events weight zero

$$\mu : \textcolor{blue}{D}X \vdash \textcolor{teal}{Ce}_{\mu}[\emptyset] = 0$$

# Axioms for events and distributions

Boolean Sub-algebra of Events

$$E : \mathcal{B}_X \vdash E^C : \mathcal{B}_X \quad E, F : \mathcal{B}_X \vdash E \cap F : \mathcal{B}_X \quad \text{so also: } E, F : \mathcal{B}_X \vdash X, E \cup F : \mathcal{B}_X$$

Disjoint additivity

$$w, v : \mathbb{W} \vdash w + v : \mathbb{W} \quad E, C : \mathcal{B}_X, \mu : \mathsf{DX} \vdash \underset{\mu}{\mathsf{Ce}}[E] = \underset{\mu}{\mathsf{Ce}}[E \cap C] + \underset{\mu}{\mathsf{Ce}}[E \cap C^C]$$

# Axioms for events and distributions

Boolean Sub-algebra of Events

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Disjoint additivity

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Exercise

Derive ‘axiomatically’ that:

- ▶ measurement is **monotone**:

$$\mu : \mathsf{DX}, E \subseteq F \vdash \underset{\mu}{\mathsf{Ce}}[E] \leq \underset{\mu}{\mathsf{Ce}}[F]$$

- ▶ the **inclusion-exclusion** principle:

$$\mu : \mathsf{DX}, E, F : \mathcal{B}_X \vdash \underset{\mu}{\mathsf{Ce}}[E \cup F] + \underset{\mu}{\mathsf{Ce}}[E \cap F] = \underset{\mu}{\mathsf{Ce}}[E] + \underset{\mu}{\mathsf{Ce}}[F]$$

# Axioms for events and distributions

Consider posets:

$$\omega := (\mathbb{N}, \leq) \quad (\mathcal{B}_X, \subseteq) \quad (\mathbb{W}, \leq)$$

**$\omega$ -chains** in a poset  $P = (\underline{P}, \leq)$ :

$$P^\omega := \{ p_\cdot \in \underline{P}^{\mathbb{N}} \mid p_0 \leq p_1 \leq \dots \}$$

Chain-closure of events and weights

$$E_\cdot : (\mathcal{B}_X, \subseteq)^\omega \vdash \bigcup_n E_n : \mathcal{B}_X \quad w_\cdot : (\mathbb{W}, \leq)^\omega \vdash \sup_n w_n : \mathbb{W}$$

Scott-continuity of measurement

$$E_\cdot : (\mathcal{B}_X, \subseteq)^\omega, \mu : \text{DX} \vdash \text{Ce}_\mu [\bigcup_n E_n] = \sup_n \text{Ce}_\mu [E_n]$$

# Axiom for probability

Probability distributions have total mass one

$$\text{PX} := \{\mu \in \text{DX} \mid \text{Ce}_\mu[X] = 1\} \quad \mu : \text{PX} \vdash \text{cast } \mu : \text{DX}$$

i.e., if we define:

$$\mathbb{I} := [0,1] \quad \mu : \text{PX}, E : \mathcal{B}_X \vdash \Pr_\mu[E] := \text{Ce}_{\text{cast } \mu}[E] : \mathbb{I}$$

then:

$$\mu : \text{PX} \vdash \Pr_\mu[X] = 1$$

# Integration

Lebesgue integration w.r.t. a distribution

$$\mu : \text{D}X, f : \mathbb{W}^X \vdash \int \mu(dx) f(x) : \mathbb{W}$$

(NB: We succinctly write  $\mathbb{W}^X$  for the type of functions  $X \rightarrow \mathbb{W}$ .)

Expectation w.r.t. a probability distribution

$$\mu : \text{P}X, f : \mathbb{W}^X \vdash \mathbb{E}_{x \sim \mu} [f(x)] := \int (\text{cast } \mu)(dx) f(x) : \mathbb{W}$$

We'll use variations on this notation, e.g.:

$$\int d\mu f, \int f d\mu, \int f(x) \mu(dx), \mathbb{E}_\mu [f]$$

# Summary

Have: Language and (some) axioms

Want: Model

Today: **discrete** model

Next week: **full** model

# Lecture plan

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# Discrete model

$X$ : types denote **sets**

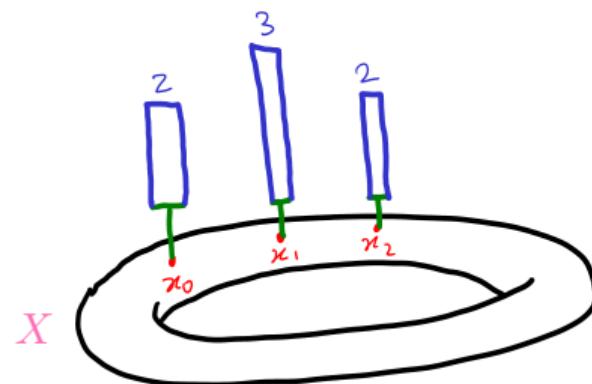
$\mathbf{D}X$ : set of **histograms**:

# Discrete model

$X$ : types denote **sets**

$\mathbb{D}X$ : set of **histograms**:

$\mathbb{D}X := \{\mu : X \rightarrow \mathbb{W} \mid \mu \text{ is } \mathbf{countably \ supported} \text{ (next slide)}\}$



$$\mu x_0 = 2 \quad \mu x_1 = 3 \quad \mu x_2 = 2$$

# Countably supported distributions

## Support

A subset  $S$  **supports** a weight function  $\mu : X \rightarrow \mathbb{W}$  when  $\mu$  is 0 outside  $S$ :

$$\mu : \mathbb{W}^X, S : \mathcal{P}X \vdash S \text{ supports } \mu := (\forall x : X. (\mu x > 0) \implies x \in S) : \text{Prop}$$

The subsets supporting a weight function  $\mu$  are closed under intersections.

$\implies$  There is a smallest supporting subset, called the **support** of  $\mu$ :

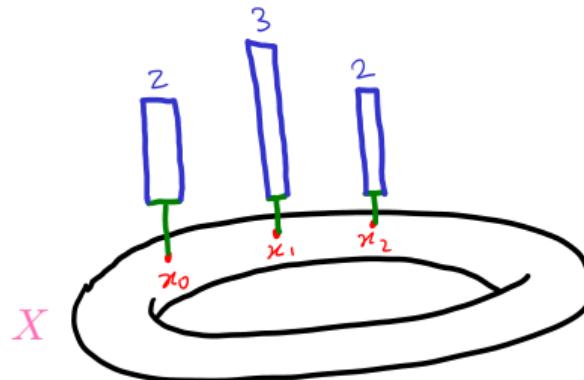
$$\mu : \mathbb{W}^X \vdash \text{supp } \mu := \{x \in X | \mu x > 0\}$$

# Discrete model

$X$ : types denote **sets**

$\mathbf{DX}$ : set of **histograms**:

$$\begin{aligned}\mathbf{DX} &:= \{\mu : X \rightarrow \mathbb{W} \mid \mu \text{ is } \mathbf{countably\ supported}\} \\ &:= \{\mu : X \rightarrow \mathbb{W} \mid \exists S \in \mathcal{P}X. S \text{ is countable}\} \\ &:= \{\mu : X \rightarrow \mathbb{W} \mid \text{supp } \mu \text{ is countable}\}\end{aligned}$$



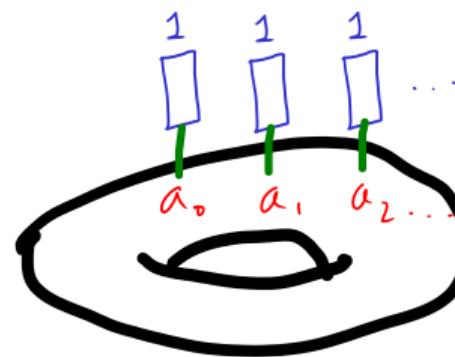
$$\mu x_0 = 2 \quad \mu x_1 = 3 \quad \mu x_2 = 2$$

# Example distributions

## Counting distribution

Counts the outcomes in a countable subset:

$$S : \mathcal{P}_{\text{ctbl}} X \vdash \#_S := \left( \lambda x. \begin{cases} x \in S & 1 \\ x \notin S & 0 \end{cases} \right) : \mathsf{D} X$$

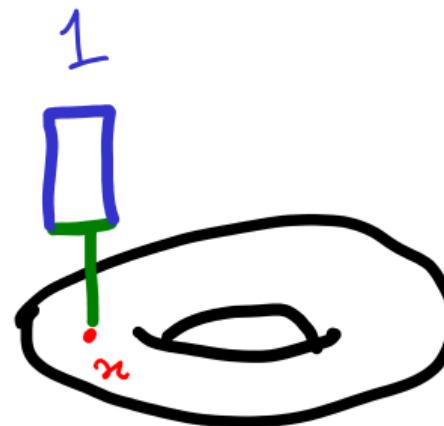


# Example distributions

Dirac

A point mass:

$$x : X \vdash \delta_x := \left( \lambda x'. \begin{cases} x' = x & 1 \\ x' \neq x & 0 \end{cases} \right) : \mathbf{D}X$$



(NB:  $x : X \vdash \delta_x = \#_{\{x\}}.$ )

# Example distributions

## Zero

No mass anywhere:

$$\vdash \mathbf{0} := \underline{0} := (\lambda x.0) : \mathbf{D}X$$

(NB:  $\vdash \mathbf{0} = \#_\emptyset.$ )

# Discrete model

$X$ : types denote **sets**

$\mathbf{D}X$ : set of **histograms**:

$$\mathbf{D}X := \{\mu : X \rightarrow \mathbb{W} \mid \mu \text{ is } \mathbf{countably\ supported}\}$$

$\mathcal{B}_X$ : **every subset** can be measured:

$$\mathcal{B}_X := \mathcal{P}X$$

Measurement: weighted sum of all (supported) outcomes:

$$\begin{aligned}\mu : \mathbf{D}X, E : \mathcal{B}_X \vdash \mathbf{Ce}_{\mu}[E] &:= \sum_{x \in E} \mu x \\ &:= \sum_{x \in E \cap \text{supp } \mu} \mu x\end{aligned}$$

NB:  $\mu : \mathbf{D}X, E : \mathcal{B}_X, S : \mathcal{P}_{\text{ctbl}}X, S \text{ supports } \mu \vdash \mathbf{Ce}_{\mu}[E] = \sum_{x \in E \cap S} \mu x.$

## Example measurements

(NB:  $\mu : \mathbf{D}X, E : \mathcal{B}_X, S : \mathcal{P}_{\text{ctbl}}X, S$  supports  $\mu \vdash \text{Ce}_\mu [E] = \sum_{x \in E \cap S} \mu x.$ )

### Counting distribution

counts supported outcomes

$$S : \mathcal{P}_{\text{ctbl}}X, E : \mathcal{B}_X \vdash \begin{array}{l} \text{Ce}_{\#_S}[E] = |E \cap S| := \begin{cases} E \cap S \text{ has } n \in \mathbb{N} \text{ elements:} & n \\ E \cap S \text{ is infinite:} & \infty \end{cases} \end{array}$$

# Example measurements

(NB:  $\mu : \text{DX}, E : \mathcal{B}_X, S : \mathcal{P}_{\text{ctbl}} X, S$  supports  $\mu \vdash \text{Ce}_\mu [E] = \sum_{x \in E \cap S} \mu x.$ )

## Counting distribution

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## Dirac

detects given outcome:

$$x : X, E : \mathcal{B}_X \vdash \text{Ce}_{\delta_x} [E] = \begin{cases} x \in E : & 1 \\ x \notin E : & 0 \end{cases}$$

# Example measurements

(NB:  $\mu : \text{DX}, E : \mathcal{B}_X, S : \mathcal{P}_{\text{ctbl}} X, S$  supports  $\mu \vdash \text{Ce}_\mu [E] = \sum_{x \in E \cap S} \mu x.$ )

## Counting distribution

counts supported outcomes

$$S : \mathcal{P}_{\text{ctbl}} X, E : \mathcal{B}_X \vdash \text{Ce}_{\#_S} [E] = |E \cap S| := \begin{cases} E \cap S \text{ has } n \in \mathbb{N} \text{ elements:} & n \\ E \cap S \text{ is infinite:} & \infty \end{cases}$$

## Dirac

detects given outcome:

$$x : X, E : \mathcal{B}_X \vdash \text{Ce}_{\delta_x} [E] = \begin{cases} x \in E : & 1 \\ x \notin E : & 0 \end{cases}$$

## Zero

measures every event as zero:

$$E : \mathcal{B}_X \vdash \text{Ce}_0 [E] = 0$$

# The discrete model validates the axioms

## Exercise

$$\mu : \mathbf{D} \quad \vdash_{\mu} \text{Ce}[\emptyset] = 0$$

$$E, C : \mathcal{B}_X, \mu : \mathbf{D} \quad \vdash_{\mu} \text{Ce}[E] = \text{Ce}_{\mu}[E \cap C] + \text{Ce}_{\mu}[E \cap C^c]$$

$$E_+ : (\mathcal{B}_X, \subseteq)^\omega, \mu : \mathbf{D}x \vdash_{\mu} \text{Ce}\left[\bigcup_n E_n\right] = \sup_n \text{Ce}_{\mu}[E_n]$$

# Parameterised distributions

## Kernel

$k : X \rightsquigarrow Y$  from  $X$  to  $Y$ : function  $k : X \rightarrow \mathbf{D}Y$ .

Kernels are open/parameterised distributions.

## Examples

Dirac and the counting distribution form kernels:

$$\delta_- : X \rightsquigarrow DX \quad \#_- : \mathcal{P}_{\text{ctbl}} X \rightsquigarrow DX$$

NB: This definition is **internal**: when we consider the full model, we will define kernels as those functions internal to the model rather than the set-theoretic functions.

# Action of kernels on distributions

Kock integral

$$\mu : \mathbf{D}X, k : (\mathbf{D}Y)^X \vdash \oint d\mu k : \mathbf{D}Y$$

This **distribution-valued** integral is implicit in many probability texts. It corresponds to integrating against an arbitrary weight function or random variable.

Discrete model interpretation

$$\begin{aligned}\oint d\mu k &:= \lambda y. \sum_{x \in X} \mu x \cdot k(x; y) \\ &:= \lambda y. \sum_{x \in \text{supp } \mu} \mu x \cdot k(x; y)\end{aligned}$$

NB1: we write  $k(x; y) := k(x)(y)$  for the uncurried function.

NB2:  $\mu : \mathbf{D}X, k : (\mathbf{D}Y)^X, S : \mathcal{P}_{\text{ctbl}} X, S \text{ supports } \mu \vdash \oint d\mu k = \lambda y. \sum_{x \in S} \mu x \cdot k(x; y)$

## Example

### Weak Disintegration Problem (non-standard terminology)

Input: distributions  $\mu : D\Theta$ ,  $\nu : DX$

Output: kernel  $k : \Theta \rightsquigarrow DX$  such that:  $\nu = \oint d\mu k$ .

Such a **weak disintegration** of  $\nu$  w.r.t.  $\mu$  provides an ‘explanation’ of an observed distribution  $\nu \in DX$  in terms of a given distribution on parameters  $\mu \in D\Theta$ . I use the term ‘explanation’ because it explains how the parameters transform into observations.

# Example

## Weak Disintegration Problem (non-standard terminology)

Input: distributions  $\mu : D\Theta$ ,  $\nu : DX$

Output: kernel  $k : \Theta \rightsquigarrow DX$  such that:  $\nu = \oint d\mu k$ .

### Example disintegration

For  $n \in \mathbb{N}$ , write  $\text{Fin } n := \{0, \dots, n - 1\}$ . For countable  $X$ , write  $\# := \#_X : DX$ .

Here is a disintegration of  $\# \in D((\text{Fin } 2)^{\text{Fin}(n+1)})$  w.r.t.  $\# \in D(\text{Fin } 2)$ :

$$k(x; f) := \begin{cases} fn = x : & 1 \\ \text{otherwise:} & 0 \end{cases} \quad \text{Indeed: } \left( \oint d\# k \right) f = \sum_{b \in \text{Fin } 2} \overbrace{\# b}^1 \cdot k(b; f) = k(0; f) + k(1; f)$$

$f : \text{Fin}(n+1) \rightarrow \text{Fin } 2$  function

so can take only one value: 0 or 1

$$\downarrow \\ = 1 = \# f$$

# Sub-type of probability distributions

## Sub-types

Given type  $X$  and  $x : X \vdash \varphi : \text{Prop}$ , take the **sub-type** and the **coercion** as follows:

$$\{x : X | \varphi\} \subseteq X \quad y : \{x : X | \varphi\} \vdash \text{cast } y := y : X$$

we **lift** values in  $X$  that satisfy  $\varphi$  to the sub-type:

$$\frac{\Gamma \vdash M : X \quad \Gamma \vdash \varphi [x \mapsto M]}{\Gamma \vdash \text{lift } M : \{x : X | \varphi\}} \quad \frac{\Gamma \vdash M : X \quad \Gamma \vdash \{\varphi\} x \mapsto M}{\Gamma \vdash \text{cast}(\text{lift } M) = M}$$

The axiom implies that **lift**  $M$  lifts  $M$  along **cast**. Moreover:

$$y : \{x \in X | \varphi\} \vdash \text{lift}(\text{cast } y) = y \quad y : \{x \in X | \varphi\} \vdash \varphi [x \mapsto \text{cast } y]$$

i.e., the lifting is unique and elements in the sub-type satisfy  $\varphi$ .

# Sub-type of probability distributions

## Magnitude and probability distributions

$$\mu : \mathsf{D}X \vdash \|\mu\| := \mathsf{Ce}_{\mu}[X] : \mathbb{W} \quad \mathsf{P}X := \{\mu \in \mathsf{D}X \mid \|\mu\| = 1\} \quad \mathbb{I} := [0,1] := \{w \in \mathbb{W} \mid w \leq 1\}$$

## Event probability

$$\mu : \mathsf{P}X, E : \mathcal{B}_X \vdash \Pr_{\mu}[E] := \mathsf{lift} \left( \mathsf{Ce}_{\mathsf{cast} \mu}[E] \right) : \mathbb{I}$$

## Stochastic kernel

$k : X \rightsquigarrow Y$  from  $X$  to  $Y$ : function  $X \rightarrow \mathsf{P}Y$ .

NB: in the **discrete model** these distinctions and rules amount to pure pedantry. This pedantry will pay off in the **full model**.

# Lifting Dirac and Kock

## Lemma

Dirac kernels  $\delta_- : X \rightarrow DX$  lift along `cast`:

$$x : X \vdash \|\delta_x\| = \underset{\delta_x}{\text{Ce}}[X] = 1$$

so we can overload:

$$\begin{array}{ccc} \delta_- & \nearrow & \text{PX} \\ X & =: & \downarrow \text{cast} \\ \delta_- & \searrow & DX \end{array}$$

Kock integrals of stochastic kernels by probability distributions lift along `cast`:

$$\mu : \text{PX}, k : (\text{PY})^X \vdash \text{Ce}_{\oint(\text{cast } \mu)(dx) \text{cast}(k x)}[Y] = 1$$

so we can overload:

$$\begin{array}{ccc} (\text{PX}) \times (\text{PY})^X & \dashrightarrow^{\oint} & \text{PY} \\ \text{cast} \times (\text{cast} \circ) \downarrow & =: & \downarrow \text{cast} \\ (\text{DX}) \times (\text{DY})^X & \xrightarrow{\oint} & \text{DY} \end{array}$$

### Proposition

The triple  $(D, \delta_-, \oint)$  forms a monad over **Set**:

$$x : X, k : (DY)^X$$

$$\vdash \oint d\delta_x k = k x$$

$$\mu : DX$$

$$\vdash \oint \mu(dx) \delta_x = \mu$$

$$\mu : DX, k : (DY)^X, \ell : (DZ)^Y$$

$$\vdash \oint (\oint \mu(dx) k x) (dy) \ell y = \oint \mu(dx) \oint k(x; dy) \ell y$$

### Corollary

The triple  $(P, \delta_-, \oint)$  forms a monad over **Set**.

# Weighted average

## Lebesgue integral

Integration is the raison d'être for distributions:

$$\mu : \mathbf{D}X, f : \mathbb{W}^X \vdash \int d\mu f : \mathbb{W}$$

In the **discrete model**:

$$\int d\mu f := \sum_{x \in X} (\mu x) \cdot (f x) := \sum_{x \in \text{supp } \mu} (\mu x) \cdot (f x)$$

As usual, replace  $\text{supp } \mu$  by any countable supporting set:

$$\mu : \mathbf{D}X, f : \mathbb{W}^X, S : \mathcal{P}X, S \text{ supports } \mu \vdash \int d\mu f = \sum_{x \in S} (\mu x) \cdot (f x)$$

# Weighted average

## Expectation

To emphasise that some  $\mu$  is a probability distribution, we will use the notation:

$$\mu : \mathsf{P}X, f : \mathbb{W}^X \vdash \mathbb{E}_\mu [f] := \int d(\mathsf{cast} \mu) f : \mathbb{W}$$

When calculating, however, we will usually use  $\int$  and implicitly  $\mathsf{cast}$  any probability distribution to its corresponding distribution.

# Booleans

## Boolean type

The simplest kind of distinguishing outcomes:

$$\mathbb{B} := \{\text{True}, \text{False}\} \quad \frac{\Gamma \vdash M : \mathbb{B} \quad \Gamma \vdash N_1 : X \quad \Gamma \vdash N_2 : X}{\Gamma \vdash \text{if } M \text{ then } N_1 \text{ else } N_2 : X}$$

## Iverson bracket

Lets us replace Boolean propositions with arithmetic expressions:

$$b : \mathbb{B} \vdash [b] := (\text{if } b \text{ then } 1 \text{ else } 0) : \mathbb{W}$$

For example:

$$b : \mathbb{B}, w, v : \mathbb{W} \vdash \text{if } b \text{ then } w \text{ else } v = [b] \cdot w + (1 - [b]) \cdot v$$

# Simplest probabilistic model

## Bernoulli kernel

Single trial succeeding with the given probability:

$$\mathbf{B} : \mathbb{I} \rightsquigarrow \mathbb{B} \quad \mathbf{B}p := \lambda b. \begin{cases} b = \mathbf{True} : & p \\ b = \mathbf{False} : & 1 - p \end{cases}$$

For example, for a payoff of 10 units if the trial succeeds then the expected payoff is:

$$\mathbb{E}_{b \sim \mathbf{B} \frac{1}{4}} [[b] \cdot 10] = \frac{1}{4} \cdot 10 + (1 - \frac{1}{4}) \cdot 0 = \frac{10}{4} + 0 = \frac{5}{2}$$

# Events as functions

## Proposition

Membership testing induces an isomorphism between events and Boolean propositions:

$$(\in) : \mathcal{B}_X \xrightarrow{\cong} \mathbb{B}^X$$

Its inverse sends each Boolean property to the set of outcomes satisfying it:

$$\frac{x : X \vdash M : \mathbb{B}}{\{x \in X | M\} : \mathcal{B}_X} \quad \{x \in X | \varphi x\} := \{x \in X | \varphi x = \text{True}\}$$

## Characteristic function

represents an event as weight functions:  $E : \mathcal{B}_X \vdash [- \in E] : \mathbb{W}^X$

By the above proposition, every (internal)  $\{0, 1\}$ -valued weight function is the characteristic function of some event, namely, the inverse image of 1.

# Measurement through integration

## Lemma

We can replace event measurement by integration of characteristic functions:

$$\mu : \mathbf{D}X, E : \mathcal{B}_X \vdash \mathbf{Ce}_{\mu}[E] = \int \mu(dx) [x \in E]$$

We can deduce properties for  $\mathbf{Ce}[-]$  and  $\mathbf{Pr}[-]$  from those of the Lebesgue integral.

Notation:

$$\frac{\Gamma \vdash \mu : \mathbf{D}X \quad \Gamma, x : X \vdash M : \mathbb{B}}{\Gamma \vdash \mathbf{Ce}_{x \sim \mu}[M] := \mathbf{Ce}_{\mu}[\{x \in X | M\}] : \mathbb{W}}$$

and similarly for  $\mathbf{Pr}_{x \sim \mu}[M]$ .

# Language of probability & distribution (recap)

$X$  type of **values/outcomes**

$\text{DX}$  type of **distributions/measures** over  $X$

$\text{PX} \subseteq \text{DX}$  sub-type of **probability distributions** over  $X$

$\mathcal{B}_X \subseteq \mathcal{P}X$  type of **events**: subsets we wish to measure

$\mathbb{W}$  type of **weights**: values in  $[0, \infty]$

$\int, \mathbb{E}$  Lebesgue integration and the expectation operation

Type judgements describe well-formed values/outcomes of a given type, e.g.:

$$\mu : \text{DX}, E : \mathcal{B}_X \vdash \text{Ce}_{\mu}[E] : \mathbb{W}$$

(measures weight  $\text{Ce}_{\mu}[E]$  of event  $E$  according to distribution  $\mu$ )

Propositions describe properties of well-formed values/outcomes of a given type, e.g.:

$$y_1, y_2 : Y \vdash y_1 \stackrel{Y}{=} y_2 : \text{Prop} \quad \mu : \text{PX}, E : \mathcal{B}_X \vdash \text{cast}_{\mu} \text{Pr}[E] = \text{Ce}_{\mu}[E]$$

(probability of event according to probability distribution is its measure)

# Lecture plan

## Lecture 1: discrete model (today)

- ▶ Language of probability and distribution
- ▶ Discrete model
- ▶ Simply-typed probability
- ▶ Dependently-typed probability



course page

## Lecture 2: the full model

- ▶ Borel sets and measurable spaces
- ▶ Quasi-Borel spaces
- ▶ Type structure & standard Borel spaces
- ▶ Integration & random variables



ask questions on the  
Scottish PL Institute  
Zulip stream #qbs

# Simply-typed foundations for probabilistic modelling

## Compositional building blocks for modelling

- ▶ Affine combinations of distributions
- ▶ Product measures ( $\otimes$ ) :  $\mathbf{D}X \times \mathbf{D}Y \rightarrow \mathbf{D}(X \times Y)$
- ▶ Random elements and their laws (push-forward measure):  
 $(\lambda(\mu, \alpha) . \mu_\alpha) : \mathbf{D}\Omega \times X^\Omega \rightarrow \mathbf{D}X$

NB:

- ▶ Dirac kernel  $\delta_- : X \rightarrow \mathbf{D}X$
- ▶ Kock integration  
 $\oint : \mathbf{D}X \times (\mathbf{D}Y)^{\mathbf{D}X} \rightarrow \mathbf{D}Y$

## Standard vocabulary

- ▶ Joint and marginal distributions
- ▶ Independence
- ▶ Distribution/probability preservation and invariance
- ▶ Density and absolute continuity
- ▶ Almost certain/sure properties

# Simply-typed foundations for probabilistic modelling

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# Affine combinations of distributions: scaling

## Scaling distributions

$$w : \mathbb{W}, \mu : \mathbf{D}X \vdash w \cdot \mu : \mathbf{D}X$$

In the discrete model:

$$w \cdot \mu := \lambda x. w \cdot \mu x \quad \text{supp}(w \cdot \mu) \subseteq \text{supp } \mu$$

The function  $(\cdot) : \mathbb{W} \times \mathbf{D}X \rightarrow \mathbf{D}X$  is a **monoid action** for the monoid  $(\mathbb{W}, (\cdot), 1)$ :

$$\mu : \mathbf{D}X \vdash 1 \cdot \mu = \mu \quad w, v : \mathbb{W}, \mu : \mathbf{D}X \vdash w \cdot (v \cdot \mu) = (w \cdot v) \cdot \mu$$

Integration and measurement are homogeneous w.r.t. scaling:

$$w : \mathbb{W}, \mu : \mathbf{D}X, k : (\mathbf{D}Y)^X \vdash \oint d(w \cdot \mu)k = w \cdot \oint d\mu k$$

$$w : \mathbb{W}, \mu : \mathbf{D}X, f : \mathbb{W}^X \vdash \int d(w \cdot \mu)f = w \cdot \int d\mu f$$

$$w : \mathbb{W}, \mu : \mathbf{D}X, E : \mathcal{B}_X \vdash \underset{w \cdot \mu}{\text{Ce}}[f] = w \cdot \underset{\mu}{\text{Ce}}[f]$$

# Affine combinations of distributions: scaling

## Normalisation

$$\mu : \text{DX}, \|\mu\| \neq 0, \infty \vdash \frac{\mu}{\|\mu\|} := \text{lift} \left( \frac{1}{\|\mu\|} \cdot \mu \right) : \text{PX}$$

measurement is homogeneous

$$\text{Indeed: } \left\| \frac{\mu}{\|\mu\|} \right\| = \left\| \frac{1}{\|\mu\|} \cdot \mu \right\| = \frac{1}{\|\mu\|} \cdot \|\mu\| = 1$$

## Discrete uniform / categorical distribution

Random unbiased choice between finitely many options/categories:

$$S : \mathcal{P}_{\text{fin}}(X), S \neq \emptyset \vdash \mathbf{U}_S := \frac{\text{lift}\#_S}{\|\text{lift}\#_S\|} : \text{PX}$$

In the discrete model:

$$\mathbf{U}_S = \lambda x. \begin{cases} x \in S : & \frac{1}{|S|} \\ x \notin S : & 0 \end{cases}$$

so:  $x : X \vdash \mathbf{U}_{\{x\}} = \delta_x$ .

# Weights as distributions

Unit type

$$\textcolor{blue}{\mathbb{1}} := \{()\}$$

Proposition

The following two functions are mutually inverse:

$$\begin{array}{ccc} & \parallel - \parallel & \\ \textcolor{blue}{D}\textcolor{blue}{\mathbb{1}} & \xrightarrow{\hspace{2cm}} & \textcolor{blue}{\mathbb{W}} \\ & \xleftarrow{\hspace{2cm}} & \\ & (\cdot \delta_0) & \end{array}$$

Proof

Calculate:  $\mu : \textcolor{blue}{D}\textcolor{blue}{\mathbb{1}} \vdash \mu \mapsto \mu() \mapsto \lambda().\mu() = \mu$  and  $w : \textcolor{blue}{\mathbb{W}} \vdash w \mapsto \lambda().w \mapsto w$ . ■

# Internalising Lebesgue integration

## Proposition

We can recover Lebesgue integration from Kock integration:

$$\begin{array}{ccc} DX \times \mathbb{W}^X & \xrightarrow{\text{id} \times (\cong \circ)} & DX \times (D\mathbb{1})^X \\ \downarrow \int & = & \downarrow \oint \\ \mathbb{W} & \xleftarrow{\cong} & D\mathbb{1} \end{array}$$

Since measurement also reduced to Lebesgue integration, it usually suffices to prove properties of Kock integration and derive them for Lebesgue integration and for measurement.

# Affine combinations of distributions: addition

## Summation

$$\mu_- : (\text{DX})^I, I \text{ countable} \vdash \sum_{i \in I} \mu_i : \text{DX}$$

In the discrete model:

$$\sum_{i \in I} \mu_i := \lambda x. \sum_{i \in I} \mu_i x \quad \text{supp } \sum_{i \in I} \mu_i = \bigcup_{i \in I} \text{supp } \mu_i$$

## Affine and convex combinations

An **affine** combination is a countable sequence of weights  $w_- : \mathbb{W}^I$ .

It is **convex** when  $\sum_{i \in I} w_i = 1$ .

## Bernoulli revisited

We can express the Bernoulli distribution as follows:

$$p : \mathbb{I} \vdash \mathbf{B} p = \text{lift}(p \cdot \delta_{\text{True}} + (1 - p) \cdot \delta_{\text{False}}) : \text{PB}$$

# Affinity of integration and convexity of expectation

## Theorem (Multi-linearity)

The Kock and Lebesgue integrals and measurement are affine in each argument:

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, k : X \rightsquigarrow Y \vdash \oint d(\sum_{i \in I} w_i \cdot \mu_i)k = \sum_{i \in I} w_i \cdot \oint d\mu_i k$$

$$\mu : \mathbf{D}X, w_- : \mathbb{W}^I, k_- : (X \rightsquigarrow B)^I \vdash \oint d\mu(\sum_{i \in I} w_i \cdot k_i) = \sum_{i \in I} w_i \cdot \oint d\mu k_i$$

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, \varphi : \mathbb{W}^X \vdash \int d(\sum_{i \in I} w_i \cdot \mu_i)\varphi = \sum_{i \in I} w_i \cdot \int d\mu_i \varphi$$

$$\mu : \mathbf{D}X, w_- : \mathbb{W}^I, \varphi_- : (\mathbb{W}^X)^I \vdash \int d\mu(\sum_{i \in I} w_i \cdot \varphi_i) = \sum_{i \in I} w_i \cdot \int d\mu \varphi_i$$

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, E : \mathcal{B}_X \vdash \sum_{i \in I} \text{Ce}_{w_i \cdot \mu_i}[E] = \sum_{i \in I} w_i \cdot \text{Ce}_{\mu_i}[E]$$

# Weight arithmetic

This theorem, a working horse in probability, has several important consequences:

## Proposition

The isomorphism  $\mathbf{D}\mathbb{1} \cong \mathbb{W}$  is a  $\sigma$ -semiring isomorphism:

$$(\mathbf{D}\mathbb{1}, \sum, (\cdot)) \cong (\mathbb{W}, \sum, (\cdot))$$

and  $(\cdot) : \mathbb{W} \times \mathbf{D}\mathcal{X} \rightarrow \mathbf{D}\mathcal{X}$  makes each  $\mathbf{D}\mathcal{X}$  into a  $\mathbb{W}$ -module:

$$\left( \sum_{i \in I} w_i \right) \cdot \mu = \sum_{i \in I} (w_i \cdot \mu) \quad w \cdot \sum_{i \in I} \mu_i = \sum_{i \in I} w \cdot \mu_i$$

# Convex combinations of probability distributions

Lemma

**Convex** combination lifts to probability distributions:

$$w_- : \mathbb{W}^I, \mu_- : (\mathsf{P}X)^I, I \text{ countable}, \sum_{i \in I} w_i = 1 \vdash$$

$$\sum_{i \in I} w_i \cdot \mu_i := \text{lift} \sum_{i \in I} w_i \cdot (\text{cast } \mu_i) : \mathsf{P}X$$

**Proof**

Calculate:  $\left\| \sum_{i \in I} w_i \cdot (\text{cast } \mu_i) \right\| = \sum_{i \in I} w_i \cdot \|\text{cast } \mu_i\| = \sum_{i \in I} w_i \cdot 1 = 1$

■

# Convex combinations of probability distributions

## Corollary (Multi-convexity)

Stochastic Kock integration, expectation and measurement are convex:

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, k : X \rightsquigarrow Y, \sum_{i \in I} w_i = 1 \vdash \oint d(\sum_{i \in I} w_i \cdot \mu_i)k = \sum_{i \in I} w_i \cdot \oint d\mu_i k$$

$$\mu : \mathbf{D}X, w_- : \mathbb{W}^I, k_- : (X \rightsquigarrow B)^I, \sum_{i \in I} w_i = 1 \vdash \oint d\mu(\sum_{i \in I} w_i \cdot k_i) = \sum_{i \in I} w_i \cdot \oint d\mu k_i$$

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, \varphi : \mathbb{W}^X, \sum_{i \in I} w_i = 1 \vdash \mathbb{E}_{\sum_{i \in I} w_i \cdot \mu_i} [\varphi] = \sum_{i \in I} w_i \cdot \mathbb{E}_{\mu_i} [\varphi]$$

$$\mu : \mathbf{D}X, w_- : \mathbb{W}^I, \varphi_- : (\mathbb{W}^X)^I, \sum_{i \in I} w_i = 1 \vdash \mathbb{E}_\mu \left[ \sum_{i \in I} w_i \cdot \varphi_i \right] = \sum_{i \in I} w_i \cdot \mathbb{E}_\mu [\varphi_i]$$

$$\mu_- : (\mathbf{D}X)^I, w_- : \mathbb{W}^I, E : \mathcal{B}_X, \sum_{i \in I} w_i = 1 \vdash \Pr_{\sum_{i \in I} w_i \cdot \mu_i} [E] = \sum_{i \in I} w_i \cdot \Pr_{\mu_i} [E]$$

# Products

## Product distribution

$$\mu : \mathsf{D}X, \nu : \mathsf{D}Y \vdash \mu \otimes \nu := \oint \mu(dx) \oint \nu(dy) \delta_{(x,y)} : \mathsf{D}(X \times Y)$$

In the discrete model:

$$\mu \otimes \nu = \lambda(x, y) \cdot (\mu x) \cdot (\nu y) \quad \text{supp } (\mu \otimes \nu) = (\text{supp } \mu) \times (\text{supp } \nu)$$

Example: counting distribution on product space

$$S : \mathcal{P}_{\text{fin}}(X), T : \mathcal{P}_{\text{fin}}(Y) \vdash \#_{S \times T} \stackrel{\mathsf{D}(X \times Y)}{=} \#_S \otimes \#_T$$

Indeed:  $\text{supp } (\#_S \otimes \#_T) = S \times T = \text{supp } \#_{S \times T}$  and for  $(x, y) \in S \times T$ :

$$(\#_S \otimes \#_T)(x, y) = 1 \cdot 1 = 1 = \#_{S \times T}(x, y)$$

# Products

Notation:

$$\frac{\Gamma \vdash M : \mathbf{D}(X \times Y) \quad \Gamma, x : X, y : Y \vdash K : \mathbf{D}Z}{\Gamma \vdash \iint M(dx, dy)K := \oint dM(\lambda(x, y).K) : \mathbf{D}Z}$$

## Theorem (Fubini-Tonelli)

We can integrate products in any order:

$$\mu : \mathbf{D}X, \nu : \mathbf{D}Y, k : (\mathbf{D}Z)^{X \times Y} \vdash$$

$$\oint \mu(dx) \oint \nu(dy) k(x, y) = \iint (\mu \otimes \nu)(dx, dy) k(x, y) = \oint \nu(dy) \oint \mu(dx) k(x, y)$$

$$\mu : \mathbf{D}X, \nu : \mathbf{D}Y, \varphi : \mathbb{W}^{X \times Y} \vdash$$

$$\int \mu(dx) \int \nu(dy) \varphi(x, y) = \iint (\mu \otimes \nu)(dx, dy) \varphi(x, y) = \int \nu(dy) \int \mu(dx) \varphi(x, y)$$

# Applying Fubini-Tonelli

## Theorem (Rule of Product)

We can factor out products:

$$\begin{array}{ll} \mu : \mathsf{D}X, f : \mathbb{W}^X, \nu : \mathsf{D}Y, g : \mathbb{W}^Y \vdash & \iint (\mu \otimes \nu)(dx, dy) fx \cdot gy = \left( \int d\mu f \right) \cdot \left( \int d\nu g \right) \\ \mu : \mathsf{D}X, E : \mathcal{B}_X, \nu : \mathsf{D}Y, F : \mathcal{B}_Y \vdash & \underset{\mu \otimes \nu}{\mathsf{Ce}} [E \times F] = \underset{\mu}{\mathsf{Ce}} [E] \cdot \underset{\nu}{\mathsf{Ce}} [F] \end{array}$$

## Theorem

The product lifts to probability distributions:

$$\mu : \mathsf{P}X, \nu : \mathsf{P}Y \vdash (\mu \otimes \nu) := \mathsf{lift}(\mathsf{cast} \mu \otimes \mathsf{cast} \nu) : \mathsf{P}(X \times Y)$$

# Products

## Binomial distribution

the number of successful outcomes of  $n$  independent Bernoulli trials:

$$\mathbf{B}_n : \mathbb{I} \rightsquigarrow \mathsf{P}(\mathbf{Fin}(1+n)) \quad \mathbf{B}_0 p := \delta_0 : \mathsf{P}(\mathbf{Fin} 1)$$

$$\mathbf{B}_{1+n} p := \iint (\mathbf{B}_n p \otimes \mathbf{B} p)(dc, db) (\text{if } b \text{ then } \delta_{1+c} \text{ else } \delta_c) : \mathsf{P}(\mathbf{Fin}(2+n))$$

We can prove by induction on  $n$ , using Fubini-Tonelli and the Iverson bracket that:

$$p : \mathbb{I}, k : \mathbf{Fin}(1+n) \vdash \Pr_{c \sim \mathbf{B}_n p} [c = k] = \binom{n}{k}$$

# Push-forward distributions

Random element

in  $X$  any (internal) function:

$$\mu : D\Omega \vdash \alpha : \Omega \rightarrow X$$

Law

of a random element is the distribution:

$$\mu : D\Omega, \alpha : X^\Omega \vdash \mu_\alpha := \int \mu(d\omega) \delta_{\alpha\omega} : DX$$

Example

Represent outcomes of die roll by  $D6 := \{1, 2, \dots, 6\}$ , and two rolls by  $D6 \times D6$ .

The sum of the rolls is a random element:

$$(+ : D6 \times D6 \rightarrow \mathbb{N})$$

The law of the distribution  $\# \otimes \#$  counts the number of configurations in which the two rolls sum to a given number, e.g.:  $(\# \otimes \#)(+) : 1 \mapsto 0, 2 \mapsto 1$ .

# Push-forward distributions

Theorem (Law of the Unconscious Statistician)

Formulae for reparameterising integration and measurement:

$$\mu : \Omega, \alpha : X^\Omega, k : X \rightsquigarrow Y \vdash \oint d\mu_\alpha k = \oint d\mu(k \circ \alpha)$$

$$\mu : \Omega, \alpha : X^\Omega, f : \mathbb{W}^X \vdash \int d\mu_\alpha f = \int d\mu(f \circ \alpha)$$

$$\mu : \Omega, \alpha : X^\Omega, E : \mathcal{B}_X \vdash \underset{\mu_\alpha}{\text{Ce}}[E] = \underset{\mu}{\text{Ce}}[\alpha^{-1}[E]] = \underset{\omega \sim \mu}{\text{Ce}}[\alpha \omega \in E]$$

# Simply-typed foundations for probabilistic modelling

## Compositional building blocks for modelling

- ▶ Affine combinations of distributions
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NB:

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## Standard vocabulary

- ▶ Joint and marginal distributions
- ▶ Independence
- ▶ Distribution/probability preservation and invariance
- ▶ Density and absolute continuity
- ▶ Almost certain/sure properties

# Standard vocabulary: concepts concerning products

Let  $\pi_i : \prod_{i \in I} X_i \rightarrow X_i$  be the  $i$ -th projection.

**Joint distribution:**  $\mu : D(X \times Y)$ ,  $\mu : D(\prod_{i \in I} X_i)$

**Marginal distribution:** the law of a projection:

$$\mu : D\left(\prod_{i \in I} X_i\right) \vdash \mu_{\pi_i} : D X_i$$

Sometimes refers to any law of a r.e..

**Marginalisation:** the action of calculating a marginal distribution by integrating all other components.

## Exercise

$$\mu : P X, \nu : D X \vdash (\mu \otimes \nu)_{\pi_2} = \nu$$

# Independence

## Pairing random elements

$$\alpha : X^\Omega, \beta : Y^\Omega \vdash \lambda \omega. (\alpha \omega, \beta \omega) : (X \times Y)^\Omega$$

## Independent random elements

The joint law is the product of the marginals:

$$\mu : \mathsf{D}\Omega, \alpha : X^\Omega, \beta : Y^\Omega \vdash \alpha \perp_{\mu} \beta := \left( \mu_{(\alpha, \beta)} \stackrel{\mathsf{D}(X \times Y)}{=} \mu_\alpha \otimes \mu_\beta \right)$$

More generally, for finite  $I$ :

$$\mu : \mathsf{D}\Omega, \alpha_i : (X^\Omega)^I \vdash \perp_{\mu} \alpha_i := \left( \mu_{(\alpha_i)_i} \stackrel{\mathsf{D}(\prod_i X_i)}{=} \bigotimes_{i \in I} \mu_{\alpha_i} \right)$$

# Independence

## Example [Durett]

Model 3 independent coin tosses:

$$\text{Toss} := \{\text{Head}, \text{Tail}\} \quad \Omega := \text{Toss}^3 \quad \mu := \mathbf{U}_{\text{Toss}} \otimes \mathbf{U}_{\text{Toss}} \otimes \mathbf{U}_{\text{Toss}} : P\Omega$$

The outcome of the  $i^{\text{th}}$  coin toss is the random element  $\pi_i : \Omega \rightarrow \text{Toss}$ .

Consider the Boolean proposition in which the  $i^{\text{th}}$  and  $j^{\text{th}}$  tosses ( $i \neq j$ ) agree:

$$\text{Same}_{ij} := \lambda \omega. \pi_i \omega = \pi_j \omega : \Omega \rightarrow \mathbb{B}$$

Calculate:

LOTUS

$$\begin{aligned} \Pr_{\mu} [\text{Same}_{12}] &= \Pr_{(x,y) \sim \mu(\pi_1, \pi_2)} [x = y] = \Pr_{(x,y) \sim \mathbf{U} \otimes \mathbf{U}} [x = y] = \int \mathbf{U}(dx) \Pr_{y \sim \mathbf{U}} [x = y] \\ &= \frac{1}{2} \cdot \Pr_{y \sim \mathbf{U}} [\text{Head} = y] + \frac{1}{2} \cdot \Pr_{y \sim \mathbf{U}} [\text{Tail} = y] = \frac{1}{4} + \frac{1}{4} = \frac{1}{2} \end{aligned}$$

marginalisation

Fubini

# Independence

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Consider the Boolean proposition in which the  $i^{\text{th}}$  and  $j^{\text{th}}$  tosses ( $i \neq j$ ) agree:

$$\text{Same}_{ij} := \lambda \omega. \pi_i \omega = \pi_j \omega : \Omega \rightarrow \mathbb{B}$$

Therefore  $\mu_{\text{Same}_{12}} = \mathbf{U}_{\mathbb{B}}$  and similarly  $\mu_{\text{Same}_{ij}} = \mathbf{U}_{\mathbb{B}}$  for  $i \neq j$ .

# Independence

$\pi_1$ ,  $\text{Same}_{12}$ , and  $\text{Same}_{13}$  determine  $\pi_2, \pi_3$ , so:

$$\Pr_{\omega \sim \mu} [\text{Same}_{12}\omega = \text{True}, \text{Same}_{13}\omega = \text{True}]$$

Fubini-Tonelli

$$\begin{aligned} & \downarrow \\ &= \int \mathbf{U}_{\text{Toss}}(db_1) \Pr_{(b_2, b_3) \sim (\mathbf{U} \otimes \mathbf{U})} [\text{Same}_{12}(b_1, b_2, b_3) = \text{True}, \text{Same}_{13}(b_1, b_2, b_3) = \text{True}] \\ &= \frac{1}{2} \Pr_{(b_2, b_3) \sim (\mathbf{U} \otimes \mathbf{U})} [\text{Same}_{12}(\text{Head}, b_2, b_3) = \text{True}, \text{Same}_{13}(\text{Head}, b_2, b_3) = \text{True}] \\ &\quad + \frac{1}{2} \Pr_{(b_2, b_3) \sim (\mathbf{U} \otimes \mathbf{U})} [\text{Same}_{12}(\text{Tail}, b_2, b_3) = \text{True}, \text{Same}_{13}(\text{Tail}, b_2, b_3) = \text{True}] \\ &= \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} = \frac{1}{4} \end{aligned}$$

and similarly we get  $\frac{1}{4}$  in all other cases.

# Independence

## Example [Durett]

Model 3 independent coin tosses:

$$\text{Toss} := \{\text{Head}, \text{Tail}\} \quad \Omega := \text{Toss}^3 \quad \mu := \mathbf{U}_{\text{Toss}} \otimes \mathbf{U}_{\text{Toss}} \otimes \mathbf{U}_{\text{Toss}} : P\Omega$$

The outcome of the  $i^{\text{th}}$  coin toss is the random element  $\pi_i : \Omega \rightarrow \text{Toss}$ .

Consider the Boolean proposition in which the  $i^{\text{th}}$  and  $j^{\text{th}}$  tosses ( $i \neq j$ ) agree:

$$\text{Same}_{ij} := \lambda \omega. \pi_i \omega = \pi_j \omega : \Omega \rightarrow \mathbb{B}$$

Therefore  $\mu_{\text{Same}_{12}} = \mathbf{U}_{\mathbb{B}}$  and similarly  $\mu_{\text{Same}_{ij}} = \mathbf{U}_{\mathbb{B}}$  for  $i \neq j$ . So:

$$\mu_{(\text{Same}_{12}, \text{Same}_{13})} = \mathbf{U}_{\mathbb{B} \times \mathbb{B}} = \mathbf{U}_{\mathbb{B}} \otimes \mathbf{U}_{\mathbb{B}} = \mu_{\text{Same}_{12}} \otimes \mu_{\text{Same}_{13}}$$

So  $\text{Same}_{12} \perp \text{Same}_{13}$  even though their values depend on the outcome of the first toss.  
 $\mu$

# Distribution preservation

Distribution space  $(\Omega, \mu)$

A type  $\Omega$  equipped with a distribution  $\mu : D\Omega$ . Define **probability space** analogously.

Distribution preserving function

$f : (\Omega_1, \mu_1) \rightarrow (\Omega_2, \mu_2)$  is a function whose is the co domain distribution:

$$f : \Omega_1 \rightarrow \Omega_2 \quad (\mu_1)_f = \mu_2$$

$\mu : DX$  is **invariant** under  $f : X \rightarrow X$  when  $f : (X, \mu) \rightarrow (X, \mu)$  is dist. preserving.

Example

Consider the swapping function:  $\text{swap} := (\lambda(x, y). (y, x)) : X \times Y \rightarrow Y \times X$ . Then, for each  $\mu : DX$ ,  $\nu : DY$ , swapping is distribution preserving function:

$$\text{swap} : (X \times Y, \mu \otimes \nu) \rightarrow (Y \times X, \nu \otimes \mu)$$

$\text{swap}$  is invariant in the case  $X = Y$  and  $\mu = \nu$ .

# Density and scaling

## Density

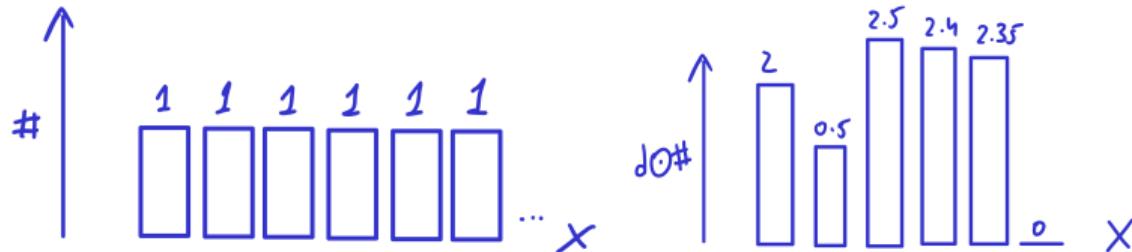
over  $X$  is any weight function  $f : X \rightarrow \mathbb{W}$ .

## Density scaling

We can scale a distribution by a density:

$$f : \mathbb{W}^X, \mu : \mathbf{D}X \vdash f \odot \mu := \int \mu(dx)(f, x) \cdot \delta_x : \mathbf{D}X$$

Scaling does not lift to probability distributions:  $\|f \odot \mu\| \neq 1$  even if  $\|\mu\| = 1$ .



# Density and scaling

## Density

over  $X$  is any weight function  $f : X \rightarrow \mathbb{W}$ .

## Density scaling

We can scale a distribution by a density:

$$f : \mathbb{W}^X, \mu : \mathsf{D}X \vdash f \odot \mu := \oint \mu(dx)(f, x) \cdot \delta_x : \mathsf{D}X$$

Scaling does not lift to probability distributions:  $\|f \odot \mu\| \neq 1$  even if  $\|\mu\| = 1$ .

## Warning!

The types of distributions and densities over  $X$  in the **discrete** model are close, but still **different**. They coincide on **countable** types, so people often confuse them. Types help us keep them separate.

# Density and absolute continuity

## Having density

This concept has several names in the literature:

$$\mu, \nu : \mathbf{D}X, f : \mathbb{W}^X \vdash \left( f = \frac{d\mu}{d\nu} \right) := (\mu = f \odot \nu) : \mathbf{Prop}$$

- ▶  $f$  is the **density** of  $\mu$  w.r.t.  $\nu$
- ▶  $f$  is a **Radon-Nikodym derivative** of  $\mu$  w.r.t.  $\nu$ .

## Absolute continuity

$\mu$  is **absolutely continuous** w.r.t.  $\nu$  when  $\mu$  has a density w.r.t.  $\nu$ :

$$\mu, \nu : \mathbf{D}X \vdash (\mu \ll \nu) := \exists f : \mathbb{W}^X. f = \frac{d\mu}{d\nu} : \mathbf{Prop}$$

# Density and absolute continuity

## Example

The **uniform distribution** is absolutely continuous w.r.t. the **counting measure** over the same support. Indeed, it has these two densities:

$$S : \mathcal{P}_{\text{fin}}(X) \vdash \left( \lambda x. \frac{1}{|S|} \right), \left( \lambda x. \begin{cases} x \in S : & \frac{1}{|S|} \\ x \notin S : & 0 \end{cases} \right) = \frac{d\mathbf{U}_S}{d\#_S}$$

These two densities are different, but they agree on the support, motivating the following concept.

# Almost certain/sure properties

## Almost certain event

is one we can assert without changing the distribution:

$$\frac{\Gamma \vdash \mu : \mathbf{D}X \quad \Gamma, x : X \vdash M : \mathbb{B}}{\Gamma \vdash \mu(\mathrm{d}x) \text{ almost certainly } M := [M] \odot \mu = \mu : \mathbf{Prop}}$$

For probabilities we define:

$$\frac{\Gamma \vdash \mu : \mathbf{P}X \quad \Gamma, x : X \vdash M : \mathbb{B}}{\Gamma \vdash \mu(\mathrm{d}x) \text{ almost surely } M := (\mathbf{cast}\ \mu)(\mathrm{d}x) \text{ almost certainly } M : \mathbf{Prop}}$$

# Existence and almost-sure uniqueness of densities

Theorem (Radon-Nikodym)

For **probability** distributions, we characterise absolute continuity as follows:

$$\mu, \nu : \mathbf{P} X \vdash (\mu \ll \nu) \iff \forall E : \mathcal{B}_X. \Pr_{\nu}[E] = 0 \implies \Pr_{\mu}[E] = 0$$

In that case, if  $f, g = \frac{d\mu}{d\nu}$  then  $\nu(dx)$  almost surely  $f x = g x$ .

In the **discrete model**, this characterisation amounts to  $\text{supp } \mu \subseteq \text{supp } \nu$ .

Example

For all countable  $X$ , we have:

$$\forall \mu : \mathbf{D} X. \mu \ll \#_X$$

Indeed, apply the Radon-Nikodym theorem, since  $\text{supp } \# = X$ .

Constructively, direct calculation shows:  $(\lambda x. \mu x) = \frac{d\mu}{d\#}$ .

# Simply-typed foundations for probabilistic modelling

## Compositional building blocks for modelling

- ▶ Affine combinations of distributions
- ▶ Product measures ( $\otimes$ ) :  $\mathbf{D}X \times \mathbf{D}Y \rightarrow \mathbf{D}(X \times Y)$
- ▶ Random elements and their laws (push-forward measure):  
 $(\lambda(\mu, \alpha) . \mu_\alpha) : \mathbf{D}\Omega \times X^\Omega \rightarrow \mathbf{D}X$

NB:

- ▶ Dirac kernel  $\delta_- : X \rightarrow \mathbf{D}X$
- ▶ Kock integration  
 $\oint : \mathbf{D}X \times (\mathbf{D}Y)^{\mathbf{D}X} \rightarrow \mathbf{D}Y$

## Standard vocabulary

- ▶ Joint and marginal distributions
- ▶ Independence
- ▶ Distribution/probability preservation and invariance
- ▶ Density and absolute continuity
- ▶ Almost certain/sure properties

# Lecture plan

## Lecture 1: discrete model (today)

- ▶ Language of probability and distribution
- ▶ Discrete model
- ▶ Simply-typed probability
- ▶ Dependently-typed probability



course page

## Lecture 2: the full model

- ▶ Borel sets and measurable spaces
- ▶ Quasi-Borel spaces
- ▶ Type structure & standard Borel spaces
- ▶ Integration & random variables



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# Type dependencies

## Example: Binomial kernels

We've defined, for every  $n \in \mathbb{N}$ , the binomial kernel:

$$\vdash \mathbf{B}_n : \mathbb{I} \rightsquigarrow \mathbf{Fin}(1 + n)$$

We will now look at **dependent-type** structure which allows us to view these as one kernel internally:

$$n : \mathbb{N} \vdash \mathbf{B}_n : \mathbb{I} \rightsquigarrow \mathbf{Fin}(1 + n)$$

# Family model

Family over an indexing set  $I$

consists of a sequence  $X_ = (X_i)_{i \in I}$  of sets.

We call each set  $X_i$  the **fibre over  $i$** .

Family  $F$

a pair  $F = (I, X_)$  consisting of (indexing) set  $I$  and a family  $X_$  over it.

Notation:  $F = I \vdash X_$

$= i : I \vdash X_i$ .

Example

The family  $n : \mathbb{N} \vdash \mathbf{Fin} n$  has  $\mathbb{N}$  as the indexing set. The fibre over  $n \in \mathbb{N}$  is:

$$\mathbf{Fin} n := \{0, 1, \dots, n - 1\}$$

# Family model

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Notation:  $F = I \vdash X_$

$$= i : I \vdash X_i.$$

Family map

$(\theta, f_ ) : (I \vdash X_ ) \rightarrow (J \vdash Y_ )$  is a pair of a function between the indexing sets and a sequence of functions between the corresponding fibres:

$$\theta : I \rightarrow J \quad (f_i : X_i \rightarrow Y_{\theta i})_{i \in I}$$

Notation:  $\theta \vdash f_$ . We won't use these maps explicitly, but they are the foundation.

# Terms in context

Dependent elements  $i : I \vdash M : X_i$

in family  $i : I \vdash X_i$  are  $I$ -indexed sequences of elements from the corresponding fibres:

$$(M \in X_i)_{i \in I}$$

## Example

We have the elements:

$$n : \mathbb{N} \vdash 0, \dots, n - 1 : \mathbf{Fin}\,n$$

## Subsumption

Every simple type becomes a family by ignoring the dependency through the constant family, e.g.,  $i : I \vdash \mathbb{N}$  and  $i : I \vdash 42 : \mathbb{N}$ .

# Simple functions

Fibred exponential

of two families over the same indexing set  $i : I \vdash X_i, Y_i$  is the family:

Family of distributions

$$i : I \vdash X_i \rightarrow Y_i$$

over a family  $i : I \vdash X_i$  is the family:

$$i : I \vdash \mathbf{D}X_i$$

Its sub-family of fibred **probability** distributions:

$$i : I \vdash \mathbf{P}X_i$$

Both have a **Dirac** distribution:

$$i : I \vdash \delta_- : X_i \rightarrow \mathbf{D}X_i \quad i : I \vdash \delta_- : X_i \rightarrow \mathbf{P}X_i$$

# Extension and dependent pairs

## Extension

of indexing set  $I$  by a **variable** of the family  $i : I \vdash X_i$  is the (indexing) set:

$$\coprod_{i \in I} X_i \coloneqq \bigcup_{i \in I} \{i\} \times X_i = \left\{ (i, x) \in I \times \bigcup_{i \in I} X_i \mid x \in X_i \right\}$$

Notation:  $(i : I, x : X_i) \coloneqq \coprod_{i \in I} X_i$  and we'll often write  $i, x$  instead of  $(i, x)$ .

## Dependent pairs

$$\frac{i : I \vdash X_i \quad i : I, x : X_i \vdash Y_{i,x}}{i : I \vdash (x : X_i) \times (Y_{i,x}) \coloneqq \coprod_{x \in X_i} Y_{i,x}}$$

# Functions and kernels

## Dependent functions

we identify a function  $f$  with a tuple  $(fx)_x$  as usual:

$$\frac{i : I \vdash X_i \quad i : I, x : X_i \vdash Y_{i,x}}{i : I \vdash ((x : X) \rightarrow Y_{i,x}) \coloneqq \prod_{x \in X} Y_{i,x}}$$

Dependent kernels  $i : I \vdash k : (x : X_i) \rightsquigarrow Y_{i,x}$

are dependent elements:

$$i : I \vdash k : (x : X_i) \rightarrow \mathsf{D}Y_{i,x}$$

Dependent **stochastic** kernels  $i : I \vdash k : (x : X_i) \rightsquigarrow Y_{i,x}$  are similarly:

$$i : I \vdash k : (x : X_i) \rightarrow \mathsf{P}Y_{i,x}$$

# Integration

## Dependent Kock integral

$$i : I, \mu : \mathbf{D}X_i, k : (x : X_i) \rightsquigarrow Y_{i,x} \vdash \oint d\mu k : \mathbf{D}Y_{i,x}$$

and in the **discrete model** we define it for  $i, \mu, k$  as in the simply-typed case:

$$(\oint d\mu k)y \coloneqq \sum_{x \in X_i} \mu x \cdot k(x; y) : \mathbb{W}$$

Through the identification  $\mathbb{W} \cong \mathbf{D}\mathbb{1}$  and characteristic functions, we reduce dependent Lebesgue integration and measurement to dependent Kock integration:

$$i : I, \mu : \mathbf{D}X_i, f : (x : X_i) \rightarrow \mathbb{W} \vdash \int d\mu f : \mathbb{W} \quad i : I, \mu : \mathbf{D}X_i, E : \mathcal{B}_{X_i} \vdash \text{Ce}_\mu [E] : \mathbb{W}$$
$$\int d\mu f = \sum_{x \in X} \mu x \cdot f x \quad \text{Ce}_\mu [E] = \sum_{x \in E} \mu x$$

# Random variables

Let  $\overline{\mathbb{R}} := [-\infty, \infty]$  be the extended real line.

## Signed and unsigned random variable

in a probability space  $(\Omega, \mu)$  are random elements  $\alpha : \Omega \rightarrow \overline{\mathbb{R}}$  and  $\alpha : \Omega \rightarrow \mathbb{W}$ .

The **positive** and **negative parts** are unsigned random variables  $\alpha^\pm : \overline{\mathbb{R}}^\Omega \rightarrow \mathbb{W}^\Omega$ :

$$\alpha^+ := \lambda \omega. \max(\alpha \omega, 0) = [\alpha \geq 0] \cdot |\alpha| \quad \alpha^- := \lambda \omega. -\min(\alpha \omega, 0) = [\alpha \leq 0] \cdot |\alpha|$$

An unsigned r.v.  $\alpha$  is **Lebesgue integrable** when its Lebesgue integral is finite:

$$\int d\mu \alpha < \infty.$$

For a (signed) r.v.  $\alpha$ , when either  $\alpha^+$  or  $\alpha^-$  is Lebesgue integrable, we define:

$$\mu : \mathbf{DX}, \alpha : \overline{\mathbb{R}}^X, \int d\mu \alpha^+, \int d\mu \alpha^- < \infty \vdash \int d\mu \alpha := \int d\mu \alpha^+ - \int d\mu \alpha^-$$

A signed variable is **Lebesgue integrable** when both its parts are Lebesgue integrable.

# Random variable spaces

Lebesgue integrability is a Boolean property:

$$\mu : \text{DX}, \alpha : X \rightarrow \bar{\mathbb{R}} \vdash \alpha \text{ integrable} := \int d\mu \alpha^+ < \infty \wedge \int d\mu \alpha^- < \infty : \mathbb{B}$$

Lebesgue spaces ensemble

is the family:

$$i : I, p : [1, \infty), \mu : \mathsf{P}X_i \vdash \mathcal{L}_p(X_i, \mu) := \{\alpha : X_i \rightarrow \bar{\mathbb{R}} \mid \alpha^p \text{ integrable}\}$$

Every fibre has a vector space structure and a norm (almost a Banach space!):

$$i : I, p : [1, \infty), \mu : \mathsf{P}X_i, \alpha : \mathcal{L}_p(X_i, \mu) \vdash \|\alpha\|_p := \sqrt[p]{\mathbb{E}_\mu [\|\alpha\|^p]} : \mathbb{W}$$

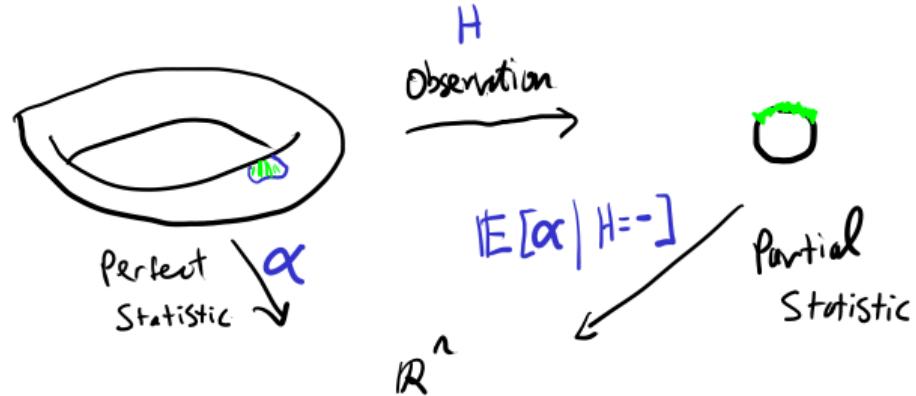
and the fibre 2 has an inner product (almost a Hilbert space!):

$$i : I, \mu : \mathsf{P}X_i, \alpha, \beta : \mathcal{L}_2(X_i, \mu) \vdash (\alpha, \beta) := \sqrt{\mathbb{E}_\mu [\alpha \cdot \beta]} : \mathbb{W}$$

# Conditioning

Situation:

- ▶ Statistical model  $\mu : D\Omega$   
(voters in the next election)
- ▶ Perfect statistic  $\alpha : \Omega \rightarrow \mathbb{R}$   
(expected winning candidate)
- ▶ Observation  $H : D \rightarrow X$   
(poll voting intention)



Conditional expectation of  $\alpha$  along  $H$  w.r.t  $\mu$

Statistic  $\beta : X \rightarrow \mathbb{R}$  that 'best' approximates  $H \circ \alpha$  statistically. Halmos and Doob's definition: any measurement we make of  $\beta$  agrees with measurement of  $\alpha$ :

$$\mu : D\Omega, H : \Omega \rightarrow X, \alpha : \mathcal{L}_1(\Omega, \mu), \beta : \mathcal{L}_1(X, \mu_H) \vdash$$

$$\left( \beta = \underset{\mu}{\mathbb{E}} [\alpha | H = -] \right) \doteq \left( \forall \varphi : \mathcal{L}_1 X, \mu_H. \int d\mu_H \beta \cdot \varphi = \int d\mu \alpha (\varphi \circ H) \right) \quad : \text{Prop}$$

# Conditioning

Theorem (Kolmogorov)

Every random variable has a conditional expectation:

$$\mu : D\Omega, H : \Omega \rightarrow X, \alpha : \mathcal{L}_1(\Omega, \mu) \vdash \exists \beta : \mathcal{L}_1(X, \mu_H). \beta = \underset{\mu}{\mathbb{E}} [\alpha | H = -]$$

Therefore:

Corollary (Internal conditional expectation)

In the **discrete model** we have a dependent function:

$$\mathbb{E}_- [- | - = -] :$$

$$(\mu : D\Omega) \rightarrow (H : \Omega \rightarrow X) \rightarrow (\alpha : \mathcal{L}_1(\Omega, \mu)) \rightarrow \left\{ \beta : \mathcal{L}_1(X, \mu_H) \middle| \beta = \underset{\mu}{\mathbb{E}} [\alpha | H = -] \right\}$$

# Conditioning

Conditional probability

of event is a conditional expectation of its characteristic function:

$$\mu : \mathsf{P}\Omega, H : \Omega \rightarrow X, E : \mathcal{B}_\Omega, \alpha : \mathcal{L}_1(\Omega, \mu) \vdash \\ \left( \alpha = \Pr_{\mu} [E|H = -] \right) \coloneqq \left( \alpha = \mathbb{E}_{\omega \sim \mu} [\omega \in E|H = -] \right) : \mathsf{Prop}$$

Regular conditional probability

a kernel that agrees with the conditional expectation of the characteristic functions:

$$\mu : \mathsf{P}\Omega, H : \Omega \rightarrow X, k : X \rightsquigarrow \Omega \vdash \\ \left( k = \Pr_{\mu} [-|H = -] \right) \coloneqq \left( \forall E \in \mathcal{B}_\Omega. k(-; E) = \mathbb{E}_{\omega \sim \mu} [\omega \in E|H = -] \right) : \mathsf{Prop}$$

# Conditioning

Kolmogorov's theorem does **not** ensure the existence of a regular conditional probability, although the constructive, discrete, definition does.

# Conditioning

Disintegration of  $H : \Omega \rightarrow X$  w.r.t.  $\mu : D\Omega$

is a kernel  $k : X \rightsquigarrow \Omega$  that weakly disintegrates the law of  $H$  and, moreover,  $d\mathbf{x}$ -almost surely each probability distribution  $k(\mathbf{x})$  is in the fibre:

$$\begin{aligned} \mu : D\Omega, H : \Omega \rightarrow X, k : X \rightsquigarrow \Omega \vdash k \text{ disintegrates } H \text{ w.r.t. } \mu \\ := \mu = \oint d\mu_H k, \\ \mu_H(d\mathbf{x}) \text{ almost surely } k(\mathbf{x}; d\omega) \text{ almost surely } H\omega = \mathbf{x} \quad : \text{Prop} \end{aligned}$$

In the **discrete model** we have an internal disintegration:

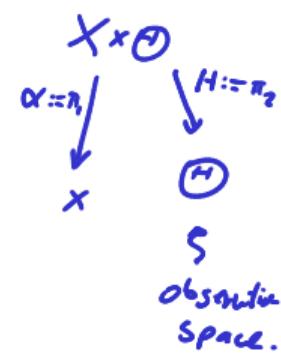
$$\begin{aligned} -^{\dagger-} : (H : \Omega \rightarrow X) \rightarrow (\mu : P\Omega) \rightarrow \{k : X \rightsquigarrow \Omega | k \text{ disintegrates } H \text{ w.r.t. } \mu\} \\ H^{\dagger\mu} := \lambda x. \begin{cases} \mu x > 0 : \frac{[H=x]}{\mu_H x} \odot \mu \\ \mu x = 0 : \mu \end{cases} \end{aligned}$$

## Bayes's Thm (discrete version, adapted from Williams):

Let  $\lambda : P(X \times \Theta)$  joint probability distribution.

Assume  $\mu : D_X, V : D_{\Theta}$  s.t.  $\lambda \ll \mu \otimes V$ .

with  $d_{X,\Theta} = \frac{d\lambda}{d(\mu \otimes V)}$ .



Obs 1:  $d_X : W^X$       then  $d_X = \frac{d\lambda}{d\mu}$   
 $d_X := \lambda x. \int V(d\theta) d_{X,\Theta}(x, \theta)$

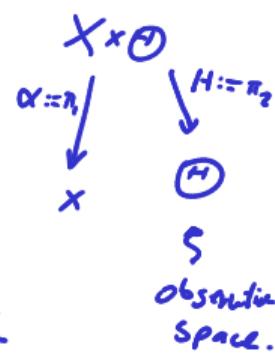
& similarly  $(d_{\Theta} : W^\Theta) := \lambda \theta. \int \mu(dx) d_{X,\Theta}(x, \theta) = \frac{d\lambda_\Theta}{d\mu}$

## Bayes's Thm (discrete version, adapted from Williams):

Let  $\lambda: P(X \times \Theta)$  joint probability distribution.

Assume  $\mu: D_X, V: D_{\Theta}$  s.t.  $\lambda \ll \mu \otimes V$ .

with  $d_{X,H} = \frac{d\lambda}{d(\mu \otimes V)}$  .  $d_X = \frac{d\lambda}{d\mu}$   $d_{\Theta} = \frac{d\lambda_H}{dV}$



Let  $d_{X|H}^{(-|-)}: X \times \Theta \rightarrow W$

$$d_{X|H}^{(-|\theta)}(x|\theta) := \begin{cases} d_H \theta \neq 0: & \frac{d_{X,H}(x,\theta)}{d_H \theta} \\ \text{o.w.:} & 0 \end{cases}$$

$$\lambda_{X|H=-}: \Theta \rightarrow P_X$$

$$\lambda_{X|H=\theta} := d_{X|H}^{(-|\theta)} \otimes \mu$$

Bayes's formula:

$$P_r[-|H=-] = \lambda_{X|H=-}$$

# Lecture plan

## Lecture 1: discrete model

- ▶ Motivation
- ▶ Language of probability and distribution
- ▶ Discrete model
- ▶ Simply-typed probability
- ▶ Dependently-typed probability



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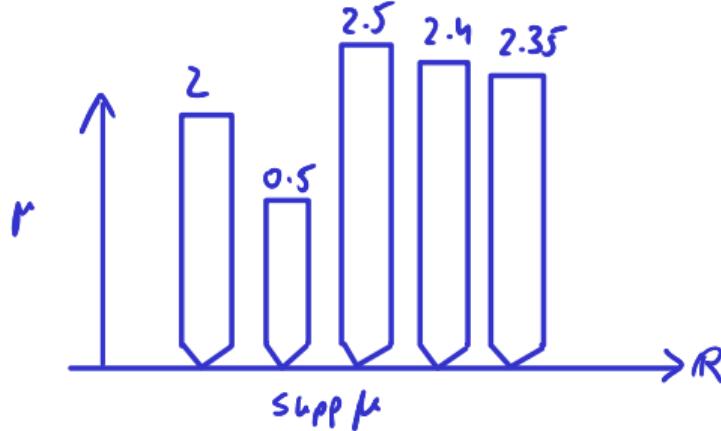
## Lecture 2: the full model

- ▶ Borel sets and measurable spaces
- ▶ Quasi-Borel spaces
- ▶ Type structure & standard Borel spaces
- ▶ Integration & random variables



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discrete model measure Only histograms:



Want :

- lengths
- areas
- volumes .

Continuous Caveat:

Then: No  $\lambda: \mathcal{P}R \rightarrow [0, \infty]$ :

$$\lambda(a, b) = b - a \quad (\text{generalises length})$$

$$\lambda(r + A) = \lambda A \quad (\text{translation invariant})$$

$$\lambda\left(\bigcup_{n=0}^{\infty} A_n\right) = \sum_{n=0}^{\infty} \lambda A_n \quad (\sigma\text{-additive})$$

Takeaway: Taking  $\mathcal{B}/R := \mathcal{P}R$

Excludes measures such as:

length, area, volume

Workaround: only measure well-behaved subsets

Df: The Borel Subsets  $B_{\mathbb{R}} \subseteq \mathcal{P}(\mathbb{R})$ :

- Open intervals  $(a, b) \in B_{\mathbb{R}}$

Closure under  $\sigma$ -algebra operations:

$$\frac{}{\emptyset \in B_{\mathbb{R}}} \quad \frac{A \in B_{\mathbb{R}}}{A^c := \mathbb{R} \setminus A \in B} \quad \frac{}{\overline{A} \in B_{\mathbb{R}}^N}$$

Empty set      complements

$$\frac{\bigcup_{n=0}^{\infty} A_n \in B_{\mathbb{R}}}{\text{countable unions}}$$

## Examples

discrete Countable:  $\{r\} = \bigcap_{\varepsilon \in \mathbb{Q}^+} (r-\varepsilon, r+\varepsilon) \in \mathcal{B}_{\mathbb{R}}$

$I$  countable  $\Rightarrow I = \bigcup_{r \in I} \{r\} \in \mathcal{B}_{\mathbb{R}}$

Closed intervals:  $[a,b] = (a,b) \cup \{a,b\}$

Non-examples?

More complicated: analytic, lebesgue

Df: Measurable Space  $V = (\mathbb{V}, \mathcal{B}_V)$

Set  $\hookrightarrow$   
(Carrier)  $\quad \quad \quad$  Family of  
Subsets  
 $\mathcal{B}_V \subseteq \mathcal{P}(\mathbb{V})$

Closed under  $\sigma$ -algebra operations:

$$\underline{\quad}$$

$\emptyset \in \mathcal{B}_V$   
Empty set

$$A \in \mathcal{B}_V$$

$A^c := \mathbb{V} \setminus A \in \mathcal{B}_V$   
Complements

$$\vec{A} \in \mathcal{B}_V^N$$

$\overline{\bigcup_{n=0}^{\infty} A_n \in \mathcal{B}_V}$   
Countable unions

Idea: Structure all spaces after the worst-case scenario

## Examples

- Discrete spaces  $\overset{\text{meas.}}{X} = (X, \mathcal{P}X)$
- Euclidean spaces  $\mathbb{R}^n$  — replace intervals with  
charts  $\prod_{i=1}^n (a_i, b_i)$   
 $\mathbb{R}^{\mathbb{N}}$  similarly
- Sub spaces:  $A \in \mathcal{P}V$   $A := (A, [\mathcal{B}_V] \cap A)$
- Products:  $A \times B := (\_A \times \_B, \sigma([\mathcal{B}_A] \times [\mathcal{B}_B]))$

$$\{C \cap A \mid C \in \mathcal{B}_V\}$$

/

Def: Borel measurable functions  $f: V_1 \rightarrow V_2$

- functions  $f: V_{1,} \rightarrow V_{2,}$
- inverse image preserves measurability:

$$f^{-1}[A] \in \mathcal{B}_{V_1} \iff A \in \mathcal{B}_{V_2}$$

### Examples

- $(+), (\cdot): \mathbb{R}^2 \rightarrow \mathbb{R}$
- $| - |, \sin: \mathbb{R} \rightarrow \mathbb{R}$
- any continuous function  $f: \mathbb{R}^n \rightarrow \mathbb{R}^m$
- any function  $f: X \rightarrow V$

# Category Meas

Objects: Measurable spaces

Morphisms: Measurable functions

Identities:

$$id : V \rightarrow V$$

Composition:

$$\begin{array}{ccc} f : V_2 \rightarrow V_3 & & g : V_1 \rightarrow V_2 \\ \searrow & & \swarrow \\ & f \circ g : V_1 \rightarrow V_3 & \end{array}$$

## Meas Category

Products, Coproducts / disjoint union, Subspaces

Categorical limits, colimits, but:

Theorem [Aumann '61] No  $\sigma$ -algebras  $B_{\mathbb{B}_{\mathbb{R}}}$ ,  $B_{\mathbb{R}^{\mathbb{R}}}$  for measurable

membership ←  
predicate  $\rightarrow$  :  $(B_{\mathbb{R}}, B_{\mathbb{B}_{\mathbb{R}}}) \times \mathbb{R} \rightarrow \text{Bool}$   
 $(U, r) \mapsto [r \in U]$

eval :  $(\text{Meas}(\mathbb{R}, \mathbb{R}), B_{\mathbb{R}^{\mathbb{R}}}) \times \mathbb{R} \rightarrow \mathbb{R}$   
 $(f, r) \mapsto f(r)$

Sequential Higher-order structure:

$$\text{I Countable : } V^{\mathbb{I}} = \prod_{i \in \mathbb{I}} V$$

$\Rightarrow$  Some higher-order structure in Meas:

$$\text{Cauchy} \in \mathcal{B}_{[-\infty, \infty]^{\mathbb{N}}}$$

$$\text{Cauchy} := \bigcap_{\varepsilon \in \mathbb{Q}^+} \bigcup_{k \in \mathbb{N}} \bigcap_{\substack{m, n \in \mathbb{N} \\ m, n \geq k}} \{ \vec{y} \in [-\infty, \infty]^{\mathbb{N}} \mid |y_m - y_n| < \varepsilon \}$$

$$\lim \text{Sup} : [-\infty, \infty]^{\mathbb{N}} \rightarrow [-\infty, \infty] \quad \lim : \text{Cauchy} \rightarrow \mathbb{R}$$

Compose higher-order building blocks:

lim is measurable!

$$\text{VanishingSeq}(\mathbb{R}) := \left\{ \vec{r} \in \mathbb{R}^N \mid \lim_{n \rightarrow \infty} r_n = 0 \right\} \subseteq \mathcal{B}_{\mathbb{R}^N}$$

$$\text{approx\_} : \text{VanishingSeq}(\mathbb{R}^+) \times \mathbb{R} \longrightarrow \mathbb{Q}^N$$

$$\text{s.t.: } |(\text{approx}_{\Delta} \vec{r})_n - r| < \Delta_n$$

Slogan: Measurable by Type !

Not all operations of interest fit:

$$\limsup : ([-\infty, \infty]^{\mathbb{R}})^{\mathbb{N}} \rightarrow [-\infty, \infty]^{\mathbb{R}}$$

$$\limsup := \lambda \vec{f}. \lambda x. \limsup_{n \rightarrow \infty} f_n x$$

Intrinsically  
higher-order !

Want

Slogan: measurability by type!

But

For higher-order building blocks

defer measurability proofs until

we resume 1<sup>st</sup> order fragment  $\Rightarrow$  <sup>non</sup>composition

# Plan

Def:  $V \in \text{Meas}$  is Standard Borel when

$$V \cong A \quad \text{for some } A \in \mathcal{B}_R$$

the "good part" of  $\text{Meas}$  - the subcategory

$$\text{Sbs} \hookrightarrow \text{Meas}$$

Sbs including

- Discrete  $\mathbb{I}$ ,  $\mathbb{I}$  countable
- Countable products of Sbs:

$$\mathbb{R}^n, \mathbb{R}^\mathbb{N}, \mathbb{Z}^n, \mathbb{N}^\mathbb{N}$$

- ~ Borel subspaces of Sbs:

$$\mathbb{I} := [0, 1]$$

$$\mathbb{R}^+ := (0, \infty) \quad \mathbb{R}_{\geq 0} := [0, \infty]$$

- Countable coproducts of Sbs:

$$\mathbb{W} := [0, \infty]$$

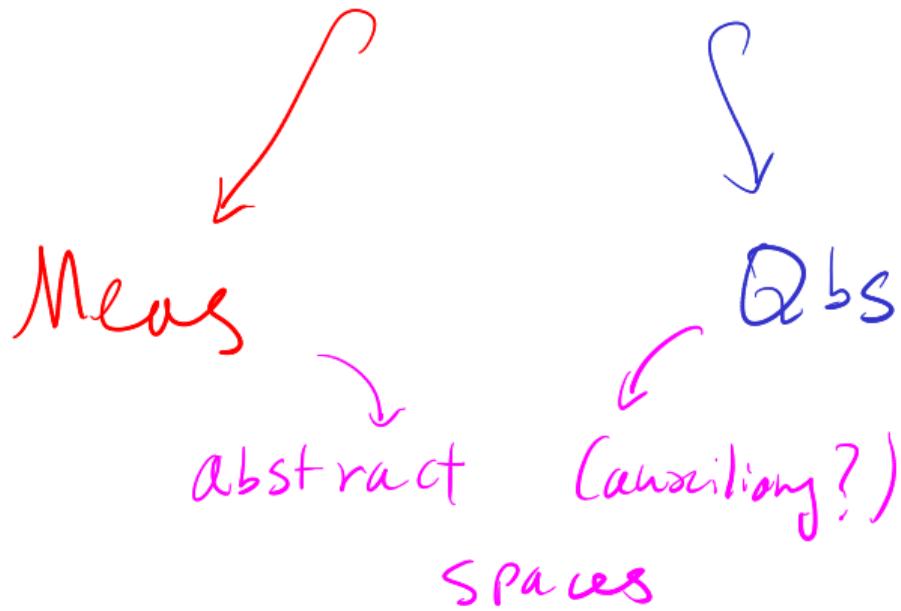
$$\overline{\mathbb{R}} := [-\infty, \infty]$$

Conservative extensions:

Concrete spaces

we "observe"

Standard Borel spaces



Plan:

- 1) Type-driven probability: discrete case ✓
  - 2) Borel sets & measurable spaces ✓
  - 3) Quasi Borel spaces
  - 4) Type structure & standard Borel spaces
  - 5) Integration & random variables
- Lecture 1
- Lecture 2

Please ask questions!

smile

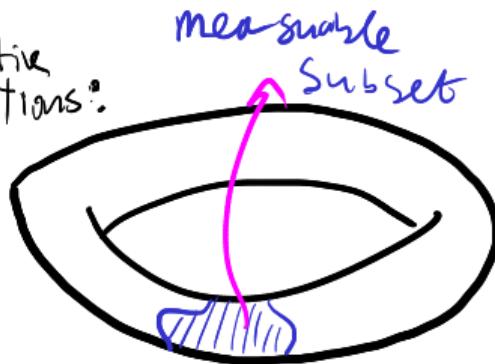


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# Cone ikeu

Measure Theory

Primitive  
notions:

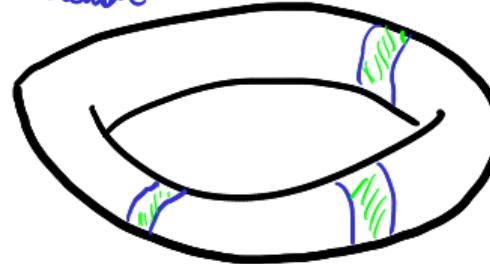


Abs Theory

Sample  
space  $\Omega$

random  
element

$\downarrow \alpha$



Derived  
notions:

random  
elements  
 $\alpha: \Omega \rightarrow \text{Space}$

measure

Events

$E \in \mathcal{B}_X$

Def: Quasi-Borel space  $X = (X_1, \mathcal{R}_X)$

$$\mathcal{R}_X \subseteq {}^{\text{L}(I\!\!R)}X_1$$

Closed under:

Set ↗  
"carrier"

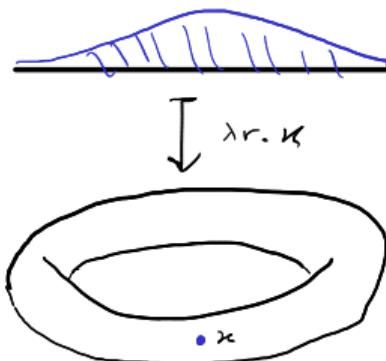
Set of  
functions  $\alpha: I\!\!R \rightarrow X_1$   
"random elements"

- Constants:

$$\frac{x \in X_1}{(\lambda r. x) \in \mathcal{R}_X}$$

- precomposition:

- recombination



Def: Quasi-Borel space  $X = (X_1, \mathcal{R}_X)$

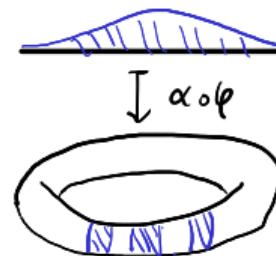
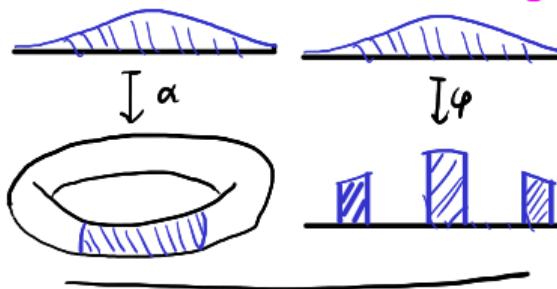
$$\mathcal{R}_X \subseteq {}^L X_1 \quad \text{closed under:}$$

- precomposition:

$$\alpha \in \mathcal{R}_X \quad \varphi: \mathbb{R} \rightarrow \mathbb{R} \text{ in } \mathcal{S}_{\mathbb{R}}$$

$$(\varphi \circ \alpha): \mathbb{R} \xrightarrow{\varphi} \mathbb{R} \xrightarrow{\alpha} X_1 \in \mathcal{R}_X$$

Set ↗  
"carrier"  
Set of  
functions  $\alpha: \mathbb{R} \rightarrow X_1$   
"random elements"



Def: Quasi-Borel space  $X = (X_1, \mathcal{R}_X)$

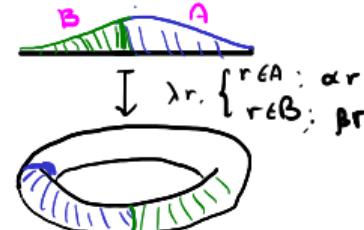
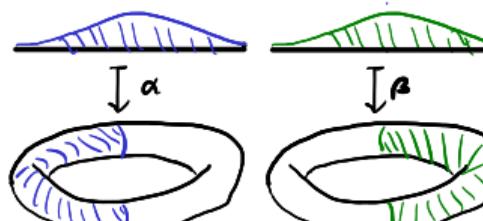
$$\mathcal{R}_X \subseteq {}^L X_1 \quad \text{closed under:}$$

- recombination

$$\vec{\alpha} \in R_X^N \quad \mathcal{R} = \bigcup_{n=0}^{\infty} A_n \quad | \quad \mathcal{E} \mathcal{B}_R$$

$$x_r. \left\{ \begin{array}{l} r \in A_n : \alpha_n r \\ \vdots \\ r \in A_0 : \alpha_0 r \end{array} \right.$$

Set ↗  
"carrier"  
Set of  
functions  $\alpha: \mathbb{R} \rightarrow X_1$   
"random elements"



Ref: Quasi-Borel space  $X = (LX, R_X)$

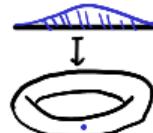
$$R_X \subseteq L^{(R_X)}$$

Closed under:

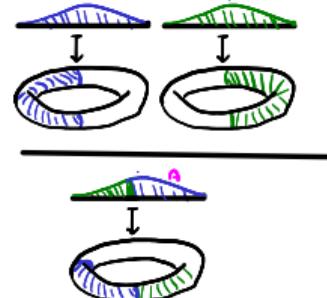
Set  
"carrier"

Set of  
functions  $\alpha: \mathbb{R} \rightarrow X$   
"random elements"

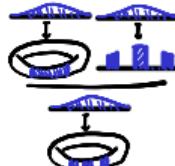
- Constants:



- recombination



- precomposition:



## Examples

-  $\mathbb{R} = (\mathbb{R}, \text{Meas}(\mathbb{R}, \mathbb{R}))$

qbs underlying  $\mathbb{R}$

-  $X \in \text{set}, \quad \lceil X \rceil := (X, \sigma\text{-simple}(\mathbb{R}, X))$

discrete qbs on  $X$

- "  $\lceil X \rceil_{\text{Qbs}} := (X, X^{\mathbb{R}})$

all functions

Indiscrete qbs on  $X$

recombination of  
constants

$\lambda_r, \begin{cases} : \\ \text{reA}n: x_n \\ : \end{cases}$

Qbs morphism  $f: X \rightarrow Y$

- function  $f: X_i \rightarrow Y_j$

- $\alpha^R_{f\downarrow} \in R_X$

$$\begin{array}{c} R \\ \alpha \downarrow \\ x \\ f \downarrow \\ y \end{array} \in R_Y$$

Example

- Constant functions

one qbs  
morphism

- $\sigma$ -simple functions

are qbs morphisms

Category Qbs  $\Leftarrow$  - identity, composition

## Full model

$\text{type} : \mathbb{Q}_{\text{bs}}$      $\mathbb{W} := [0, \infty]$      $\mathcal{B}X := \boxed{\quad}$

$\mathcal{D}X := \boxed{\quad}$

$\mathcal{P}X := \left\{ \mu \in \mathcal{D}X \mid \underset{\mu}{\text{Ce}}[X] = 1 \right\}$

$\underset{\mu}{\text{Ce}}[E] := \boxed{\quad}$      $\delta_x := \boxed{\quad}$

$\phi \mu k := \boxed{\quad}$

Plan:

- 1) Type-driven probability: discrete case ✓
- 2) Borel sets & measurable spaces ✓
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Lecture 1

Lecture 2

Please ask questions!

smile



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Lecture 1

Lecture 2

Please ask questions!

smile



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## Full model

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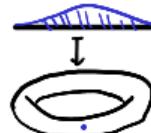
Closed under:

Set  
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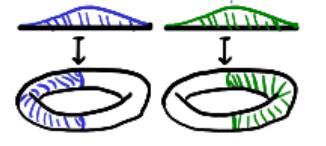
Set of

functions  $\alpha: \mathbb{R} \rightarrow X$   
"random elements"

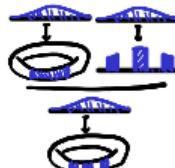
- Constants:



- recombination



- precomposition:



## Examples

-  $\mathbb{R} = (\mathbb{R}, \text{Meas}(\mathbb{R}, \mathbb{R}))$

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all functions

Indiscrete qbs on  $X$

recombination of  
constants

$\lambda_r, \begin{cases} : \\ \text{reA}n: x_n \\ : \end{cases}$

Validate gbs axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(R, \mathbb{W}))$

- Constants:

$E : \mathcal{B}_{\mathbb{W}}, n : \mathbb{W} \vdash$

$$(\lambda r : R. x)^{-1}[E] = \begin{cases} x \in E : & R \\ n \notin E : & \emptyset \end{cases} \in \mathcal{B}_R$$

✓

Validate q's axioms for:  $\mathbb{W} := ([0, \infty], \text{Meas}(R, \mathbb{W}))$

- Precomposition:

$\alpha: \text{Meas}(R, \mathbb{W}), \varphi: \text{Meas}(R, R) \vdash$

$$R \xrightarrow{\varphi} R \xrightarrow{\alpha} \mathbb{W} \quad \in \text{Meas}(R, \mathbb{W})$$

$\Downarrow$   
Meas is a cat.

Explicitly:

$$(\alpha \circ \varphi)[E] \in \beta R \xleftarrow{\varphi^*} \varphi[E] \in \beta R \xleftarrow{\alpha^*} E \in \beta \mathbb{W} \quad \checkmark$$

Validate qbs axioms for:  $\mathcal{W} := ([0, \infty], \text{Meas}(R, \mathcal{W}))$

- RC Combination

$I \text{ ctbl}, \alpha : \text{Meas}(R, \mathcal{W})^I, E_i : B_{IR}^I, R = \bigcup_{i \in I} E_i, F : \mathcal{B}_{\mathcal{W}} \vdash$

$$\left( \exists r. \left\{ \begin{array}{l} r \in E_i : \alpha_i[r] \\ \vdots \end{array} \right\} \right)^{\neg\neg} [F]$$

$$\beta := \bigcup_{i \in I} \alpha_i^{-1}[F] \cap E_i \in B_{IR}$$

In fact:

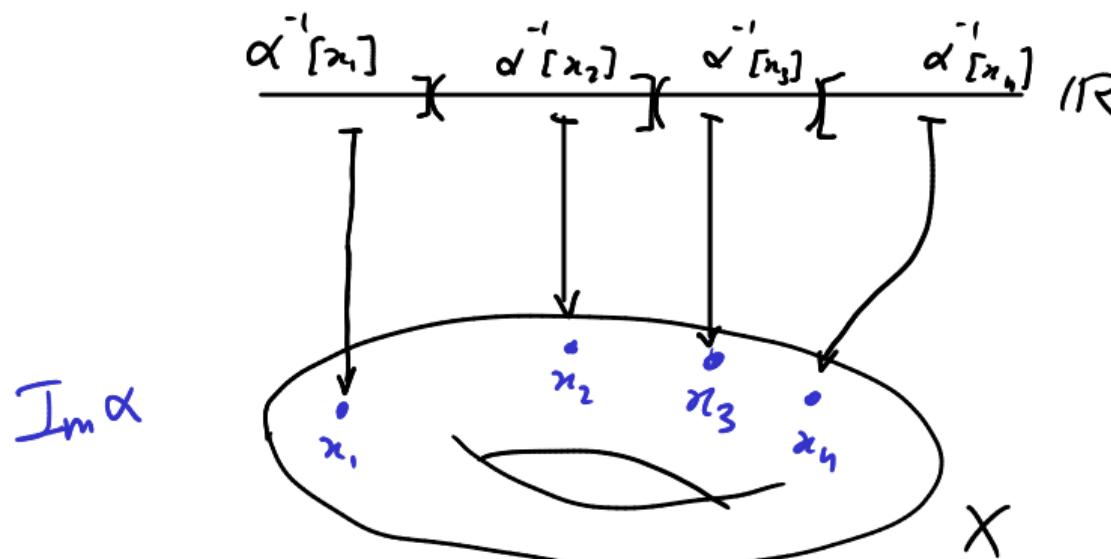
$$r \in \text{LHS} \Leftrightarrow \beta[r \in F] \Leftrightarrow \exists i \in I. r \in E_i \wedge \alpha_i[r \in F] \Leftrightarrow r \in \text{RHS}$$



## σ-Simple function

$\alpha: \mathbb{R} \rightarrow X$  s.t.  $\text{Im } \alpha := \alpha[\mathbb{R}]$  is ctbl 1

$\forall x \in \text{Im } \alpha. \alpha^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$



Validate g's axioms for:  $\stackrel{\text{Qbs}}{\vdash} X := (X, \sigma\text{-simple}(X))$

- Constants

$$\text{Im}(\lambda r. n) = \{n\} \text{ ctbl } \checkmark$$

NB:  $f$   $\sigma$ -Simple:  
 $\text{Im } f$  ctbl 1  
 $\tilde{f}[x] \in B_R$

$$g: X \vdash (\lambda r. n)^{-1}[y] = \begin{cases} x=y: R \\ x \neq y: \emptyset \end{cases} \in B_R \checkmark$$

Validate g's axioms for:  $\overset{\text{Qbs}}{X} := (X, \sigma\text{-simple}(X))$

- Precomposition:

$\alpha : \sigma\text{-simple}(X), \varphi : \text{Meas}(\mathbb{R}, \mathbb{R}) \vdash$

$$\text{Im}(\alpha \circ \varphi) \subseteq \text{Im} \alpha \text{ ctbl} \quad \checkmark$$

NB:  $f \text{ } \sigma\text{-Simple:}$   
 $\text{Im } f \text{ ctbl} \wedge$   
 $f^{-1}[x] \in \mathcal{B}_{\mathbb{R}}$

$x : X \vdash$

$$(\alpha \circ \varphi)^{-1}[x] = \varphi^{-1}[\alpha^{-1}(x)] \in \mathcal{B}_{\mathbb{R}} \quad \checkmark$$

$$\alpha^{-1}(x) \in \mathcal{B}_{\mathbb{R}}$$

$\varphi : \mathbb{R} \rightarrow \mathbb{R}$  measurable

Validate g's axioms for:  $\stackrel{\text{Qbs}}{[X]} := (X, \sigma\text{-simple}(X))$

• recombination:

$$\alpha_- : (\sigma\text{-simple}(X))^I, E : \mathcal{B}_{IR}^I, R = \bigcup_{i \in I} E_i \vdash$$

NB:  $f \sigma\text{-Simple} : \text{Im } f \subset \bigcup_{i \in I} E_i$   
 $f^{-1}[x] \in \mathcal{B}_{IR}$

$$\text{Im}[E_i, \alpha_i]_{i \in I} \subseteq \bigcup_{i \in I} \text{Im } \alpha_i \quad \text{ctbl} \quad \checkmark$$

$$x : X \vdash [E_i, \alpha_i]_{i \in I}^{-1}(x) = \bigcup_{i \in I} \alpha_i^{-1}[x] \cap E_i \in \mathcal{B}_{IR} \quad \checkmark$$

Prop:  $X: \text{Set}, A: \text{Qbs} \vdash$

•  $\forall f: X \rightarrow {}_L A_s . \quad f: {}^{\text{Qbs}} X \longrightarrow A$

•  $\forall f: {}_L A_s \longrightarrow X . \quad f: A \longrightarrow {}^X_{\text{Qbs}}$

Prop:  $X: \text{Set}, A: \text{Qbs} \vdash$

- $\forall f: X \rightarrow {}_L A_S . \quad f: {}^{r^{\text{Qbs}}_A} X \longrightarrow A$

Prf:  $\alpha: R_{r^{\text{Qbs}}_A} \vdash \alpha \text{ } \sigma\text{-simple} \Rightarrow$

$$\alpha = [\alpha^{-1}[x].\lambda r.x]_{x \in \text{Im } \alpha} \Rightarrow$$

$$(f \circ \alpha) = [\alpha^{-1}[x].\lambda r.fx]_{x \in \text{Im } \alpha} \stackrel{\text{recombination}}{\in} R_A$$

✓

$\uparrow$        $\uparrow$        $\uparrow$

Borel      constant  $\in B_A$       ctbl

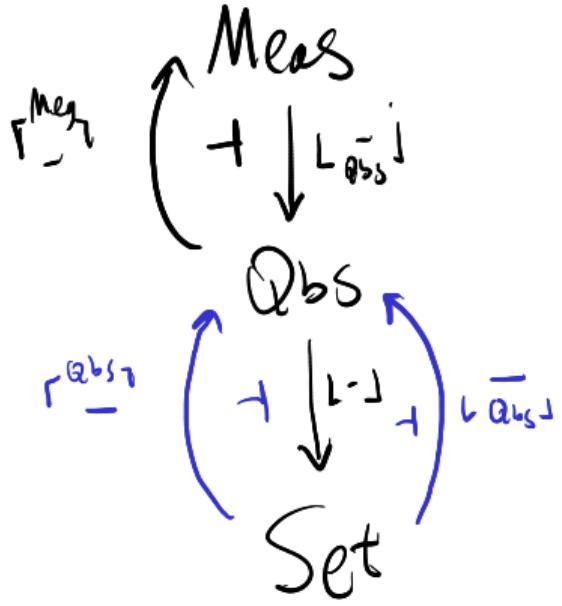
Prop:  $X: \text{Set}, A: \text{Qbs} \vdash$

- $\forall f: X \rightarrow {}_L A_1 . f: {}^{X^{\text{Qbs}}}_L \rightarrow A$
- $\forall f: {}_L A_1 \rightarrow X . f: A \rightarrow {}_{\text{Qbs}}^X$

Prf:  $\alpha: R_A \vdash (f \circ \alpha: R \rightarrow X) \in R_{{}^{X^{\text{Qbs}}}_L}$  always. ✓



# Useful adjunctions:



$$\mathcal{L}_{\text{Qbs}} := (\mathcal{L}_{\text{V}}, \text{Meas}(R, V))$$

$(V \in \text{Meas})$

$$\mathcal{L}_{\text{X}}^{\text{Meas}} := \left\{ A \subseteq X_j \mid \forall a \in R_x, a^{-1}[A] \in \mathcal{B}_R \right\}$$

- limits (products, subspaces)  
and colimits (coproducts, quotients)
- as in Set

- Slogan: every measurable space is carried by a qbs

## Example

Product  $(X \times Y, \pi_1, \pi_2)$ :

-  $L_{X \times Y} = L_{X_1 \times_1 Y_1}$  necessarily!

-  $R_{X \times Y} = \{ \lambda r. (\alpha r, \beta r) \mid \alpha \in R_X, \beta \in R_Y \}$

correlated

random  
elements

rest of structure as in Set.

# Function Spaces

Straightforward!

$$- \mathbb{Y}^X := \text{Qbs}(X, \mathbb{P})$$

$$- \mathbb{R}_{Y^X} := \text{uncurry}[\text{Qbs}(\mathbb{R} \times X, Y)]$$

$$= \left\{ \alpha: \mathbb{R} \rightarrow \mathbb{Y}^X \mid \lambda(r, x). \alpha r x: \mathbb{R} \times X \rightarrow Y \right\}$$

$$- \text{eval}: Y^X \times X \rightarrow Y$$
$$\text{eval}(f, x) := fx$$

# Meas vs Qbs

By generalities:

$\sigma$ -algebra  
on  $\text{Meas}(\mathbb{R}, \mathbb{R})$

$$\Gamma^{\text{Meas}}_{\mathbb{R}}$$

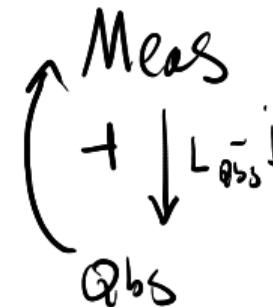
$$\mathbb{R} \times \mathbb{R}$$

$$\Gamma^{\text{Meas}}_{\mathbb{R}}$$

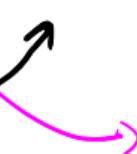
$$\mathbb{R} \times \mathbb{R}$$



$$\Gamma^{\text{Meas}}$$



$$\Gamma_{\mathbb{R}} = \mathbb{R}$$



$$\Gamma^{\text{Meas}}_{\text{Eval}}$$

$$\left( \frac{s}{\Gamma_{\mathbb{R}} \times \mathbb{R}} \neq \Gamma_{\mathbb{R}^2} \times \Gamma_{\mathbb{R}} \right)$$

No factorisation  
by  
Aumann's  
Theorem.

## Simple Type Structure

"Simple" because:

- Simply-typed  $\lambda$ -calculus
- types are simple:  $A, B : \text{Type} \vdash B^A : \text{Type}$ 
  - no polymorphism
  - no term dependency
- contexts for terms:  $\Gamma \vdash t : A$ 
  - are simple:  $\Gamma = x_1 : A_1, \dots, x_n : A_n$
  - i.e.  $\text{List}(\text{Type})$

## Simple Type Structure

"Simple" because:

- interpretation is simple:

$$\llbracket x_1 : A_1, \dots, x_n : A_n \rrbracket := \prod_{i=1}^n A_i$$

$$\llbracket \Gamma \vdash t : A \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow A$$

in QBS

# Simple Type Structure Curry-Howard-Lambek

$$\frac{\Gamma \vdash t : A \quad \Gamma \vdash s : B}{\Gamma \vdash \langle t, s \rangle : A \times B} \rightsquigarrow \boxed{\Gamma} \xrightarrow{\lambda r. \langle tr, sr \rangle} A \times B$$

is measurable

---

$$\frac{\Gamma \vdash t : A \times B \quad \Gamma, x:A, y:B \vdash s : C}{\Gamma \vdash \text{let } (x,y) = t \text{ in } s : C} \rightsquigarrow$$

measurability  
by  
type!

---

$$\boxed{\Gamma} \xrightarrow{\lambda r. \text{let } (a,b) = tr \text{ in } sr[x \mapsto a, y \mapsto b]} C$$

is measurable. etc.

## Random elent space

$R_X := X^R$  since  $\lfloor X^R \rfloor = R_X$  as sets.

Why?

( $\subseteq$ )  $\alpha \in \lfloor X \rfloor^R \Rightarrow \alpha: \mathbb{R} \rightarrow X$  in Qbs.

$\text{id}_{\mathbb{R}}: \mathbb{R} \rightarrow \mathbb{R}$  measurable  $\Rightarrow \text{id} \in R_{\mathbb{R}}$

$\Rightarrow \alpha = \alpha \circ \text{id} \in R_X$

Pre composition

( $\supseteq$ )  $\alpha \in R_X \Rightarrow \exists \psi \in R_{\mathbb{R}} = \text{Meas}(\mathbb{R}, \mathbb{R})$ .  $\alpha \circ \psi \in R_X \stackrel{\checkmark}{=} \alpha: \mathbb{R} \rightarrow X$   
 $\Rightarrow \alpha \in \lfloor X \rfloor^R$

## Subspaces

For  $X \in \text{Qbs}$ ,  $A \subseteq X$ , set:

$$R_A := \left\{ \alpha: \mathbb{R} \rightarrow A \mid \alpha \in R_X \right\}$$

Then  $A = (A, R_A)$  is the *subspace qbs*

We write  $A \hookrightarrow X$

## Borel Subspaces Ensemble

The  $\sigma$ -algebra  $B_X := \left\{ A \subseteq X \mid \forall \alpha \in R_X . \alpha^*[A] \in B_{R_X} \right\}$

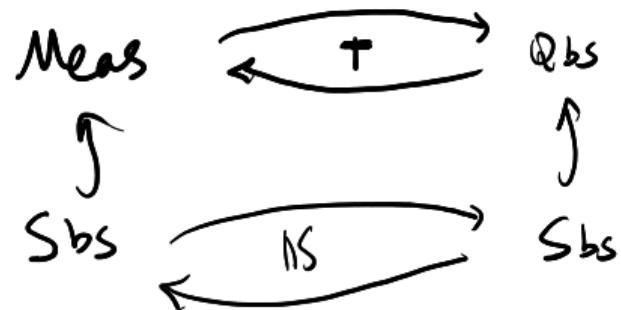
internalises as  $B_X = 2^X$ , the qbs of  
Borel subsets.

$L^{(B_{R_X})}$  are the Borel-on-Borel sets from  
descriptive set theory.  
(Cf.. [Sabou et al. '21])

## Standard Borel Spaces

Def: A qbs  $S$  is standard Borel when

$$S \cong A \text{ for some } A \in \mathcal{B}_{\mathbb{R}}$$



Slogan:  $Qbs$  conservative extension of  $Sbs$

Example  $C_0 := \{f: \mathbb{R} \rightarrow \mathbb{R} \mid f \text{ continuous}\} \hookrightarrow \mathbb{R}^{\mathbb{R}}$

$C_0$  is sbs. (Well-known!)

Proof:

$$C'_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$$

$$C'_0 := \left\{ g \in \mathbb{R}^{\mathbb{Q}} \mid \begin{array}{l} \forall a, b \in \mathbb{Q}, \varepsilon \in \mathbb{Q}^+ \\ \exists \delta \in \mathbb{Q}^+ \forall p, q \in \mathbb{Q} \cap [a, b], |p - q| < \delta \Rightarrow |g(p) - g(q)| < \varepsilon \end{array} \right\}$$

on closed intervals  
(= compact intervals)  
Continuity  
Uniform continuity  
Borel measurable  
by type checks

then  $C_0 \cong C'_0 \in B_{\mathbb{R}^{\mathbb{Q}}}$ :

$$c_0 \mapsto c'_0$$

$$c'_0 \mapsto c_0$$

$$\varphi \mapsto \varphi|_{\mathbb{Q}}$$

$$\varphi \mapsto \lambda r. \lim_{n \rightarrow \infty} g(\text{approx}_{\frac{1}{m}} \varphi)_{\frac{n}{m}}$$

## Example (ctd)

$C_0$  is sbs, and eval:  $C_0 \times \mathbb{R} \rightarrow \mathbb{R}$   
is measurable.

Avoids;

- constructing complete separable metrics
- proving that evalution is measurable  
w.r.t. metric  $\sigma$ -algebra.

Non-examples ~ [Sabok et al.'21]

$$- \left\{ A \in \mathcal{B}_{\mathbb{R}} \mid A \neq \emptyset \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$

$$- \left\{ (A_1, A_2) \in \mathcal{B}_{\mathbb{R}}^2 \mid A \subseteq B \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}^2$$

$$- \left\{ A \in \mathcal{B}_{\mathbb{R}} \mid A \text{ open} \right\} \hookrightarrow \mathcal{B}_{\mathbb{R}}$$

Plan:

- 1) Type-driven probability: discrete case ✓
  - 2) Borel sets & measurable spaces ✓
  - 3) Quasi Borel spaces ✓
  - 4) Type structure & standard Borel spaces
  - 5) Integration & random variables
- Lecture 1
- Lecture 2



Please ask questions!  
smile

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## Dependent Type Structure

Types can contain terms :

$X:\text{Type}, E:B_X \vdash \{x \in X \mid x \in E\} : \text{Type}$

a type referring  
to a term

a type, just like  
STLC

a term!

## Dependent Type Structure

types can contain terms :

$$X:\text{Type}, E:B_X \vdash \left\{ x \in X \mid x \in E \right\} : \text{Type}$$

a type, just like  
STLC

a term!

↑      ↗  
a type referring  
to a term

Content formation:

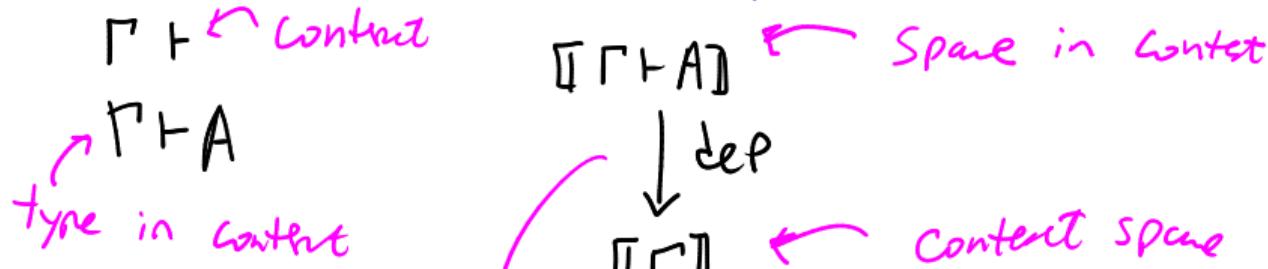
$$\frac{\Gamma \vdash A : \text{Type}}{\Gamma, x:A \vdash}$$

## Dependent Type Structure

types denote spaces-in-context

$$\begin{array}{c} \llbracket \Gamma \vdash A \rrbracket \\ \downarrow \text{def} \\ \llbracket \Gamma \rrbracket \end{array}$$

Dependent types denote spaces-in-Content



E.g.:

A



1

simple types

assigns environment

$$\llbracket E : B_A + \{x \in A \mid x \in E\} \rrbracket$$
$$\{ (E, a) \in B_A^{x \in A} \mid a \in E \}$$
$$\downarrow \pi_1$$
$$B_A$$

decoder

## Content extension

$$\frac{\Gamma \vdash A}{\Gamma, a:A \vdash}$$

$$\frac{}{\llbracket \Gamma \vdash A \rrbracket} \text{ dep} \downarrow \llbracket \Gamma \rrbracket \quad \llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket$$

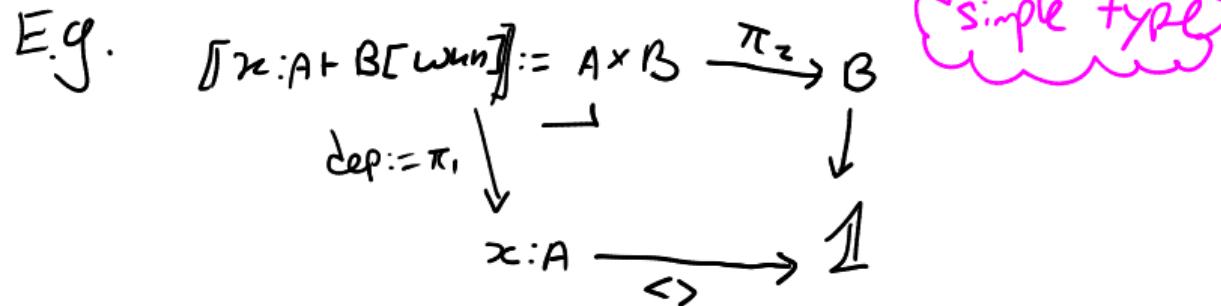
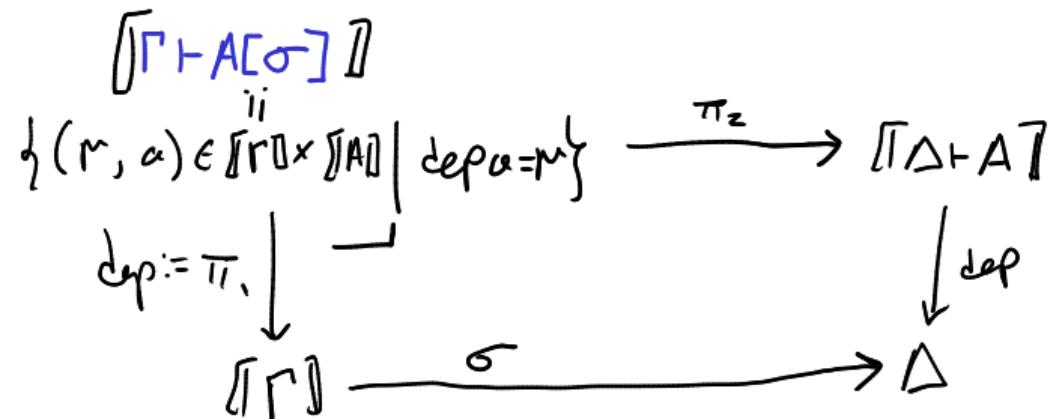
## Substitution

$$\frac{\Gamma \vdash \sigma : \Delta}{\llbracket \sigma \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket \Delta \rrbracket}$$

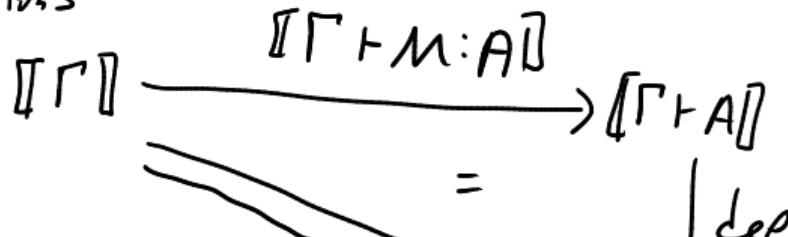
E.g. Weakening  $\Gamma, a:A \vdash \text{wkn} : \Gamma$

$$\llbracket \Gamma, a:A \rrbracket := \llbracket \Gamma \vdash A \rrbracket \xrightarrow[\text{dep}]{\text{wkn}} \llbracket \Gamma \rrbracket$$

## Action of Substitution on types



Terms : sections



e.g.

$$R \xrightarrow{\boxed{x : R \vdash [x, \alpha] : B_R[wkn]}} R \times B_R$$

=

$$R \xrightarrow{\pi_1} R$$

E.g. Variables:  $\boxed{\Gamma, \alpha : A \vdash \alpha : A}$

$$\boxed{\Gamma, \alpha : A} \xrightarrow{< \text{it}, \text{dep} >_{\Gamma \vdash A}} \boxed{\Gamma, \alpha : A \vdash A[wkn]}$$

=

$\downarrow \text{dep}$

Exercise:

action of substitution  
 $M[\sigma]$

## Dependent Pairs

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \prod_{a:A} B}$$

$$\llbracket \prod_{a:A} A \rrbracket := \llbracket \Gamma, a:A \vdash B \rrbracket$$

:=

$$\begin{array}{c} \downarrow \text{dep}_B \\ \llbracket \Gamma, a:A \rrbracket \\ \llbracket \Gamma \vdash A \rrbracket \\ \downarrow \\ \llbracket \Gamma \rrbracket \end{array}$$

$\text{dep}_{\prod}$

## Dependent Products

$$\frac{\Gamma, a:A \vdash B}{\Gamma \vdash \underset{a \in A}{\pi} B}$$

aha:  $(a:A) \rightarrow B$

$$[\Gamma \vdash \underset{a \in A}{\pi} B] :=$$

$$\left\{ (m_0, f : \{ a \in [\Gamma] \mid \text{dep } a = m_0 \} \rightarrow [\Gamma, a:A \vdash B]) \middle| \right. \\ \left. \forall a \in [\Gamma, a:A]. \text{dep } a = m_0 \Rightarrow \text{dep}(fa) = a \right\}$$

Exercise: find the random elements.

## Full model

type : Obs     $\mathbb{W} := [0, \infty]$      $\mathcal{B}^X \cong \mathcal{B}^X$

$\Omega^X := (\text{Fri})$

$P^X := \{\mu \in \Omega^X \mid \underset{\mu}{\text{Ce}}[X] = 1\}$

$\underset{\mu}{\text{Ce}}[E] := (\text{Fri})$      $\delta_x := (\text{Fri})$

$\phi \mu k := (\text{Fri})$

Plan:

- 1) Type-driven probability: discrete case ✓
  - 2) Borel sets & measurable spaces ✓
  - 3) Quasi Borel spaces ✓
  - 4) Type structure & standard Borel spaces ✓
  - 5) Integration & random variables
- Lecture 1
- Lecture 2

Please ask questions!

smile



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## Partiality cf. [Väkär et al., '19]

A Borel embedding  $e: X \hookrightarrow Y$

- injective function  $e: [X] \rightarrow [Y]$
- its image is Borel:  $e[X] \in \mathcal{B}_Y$
- $e$  is Strong:  $\alpha \in R_X \Leftrightarrow e \circ \alpha \in R_Y$

## Examples

- $\mathbb{1} \hookrightarrow \mathbb{2}$
- $S$  is sbs  $\Leftrightarrow \exists S \hookrightarrow \mathbb{R}$

Def: A Partial map  $f: X \rightarrow Y$  is a morphism

$$f: X \rightarrow Y \amalg \{\perp\}$$

Its domain of definition

$$f: (Y \amalg \{\perp\})^X \rightarrow \text{Dom } f := \left\{ x \in X \mid f_x \neq \perp \right\} : \text{Type}$$

Depent-type  
interpretation

$$\begin{array}{ccc} \llbracket \text{Dom } f \rrbracket & \xrightarrow{\quad} & \left\{ g \in Y \mid g \in E \right\} \\ \downarrow \text{dep} & & \downarrow \text{dep} \\ \left[ f : (Y \amalg \{\perp\})^X \right] \left[ \frac{E \vdash x : \text{Dom } f}{f_x \neq \perp} \right] & \xrightarrow{\quad} & \llbracket E : \mathcal{B}_Y \rrbracket \end{array}$$

Plan:

- 1) Type-driven probability: discrete case ✓
  - 2) Borel sets & measurable spaces ✓
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Please ask questions!  
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## Full model

type : Obs       $\mathbb{W} := [0, \infty]$        $\mathcal{B}^X := \mathcal{B}^X$

$\mathcal{D}^X := \square$

$P^X := \{\mu \in \mathcal{D}^X \mid \underset{\mu}{\text{Ce}}[X] = 1\}$

$\underset{\mu}{\text{Ce}}[E] := \square$        $S_n := \square$

$\phi \mu k := \square$

Def: A measure  $\mu$  over  $\mathbb{R}$  is a function

$$\mu : \mathcal{B}_{\mathbb{R}} \rightarrow \mathbb{W} := [0, \infty]$$

Satisfying the measure axioms:

$$E : \mathcal{B}^\omega \vdash$$

$$\mu \emptyset = 0, \quad \mu E = \mu(E \cap F) + \mu(E \cap F^c), \quad \mu(\bigvee_n E_n) = \sup_n \mu E_n$$

For measurable spaces, replace  $\mathbb{R}$  with  $V$

We write  $[GV]$  for the set of measures on  $V$

For qbs  $X$ , take  $[G^{r_{\text{meas}}} X]$

Thm (Lebesgue measure):

There is a unique measure  $\lambda \in \mathcal{L}(R)$ , s.t.:

$$\lambda(a, b) = b - a$$

## The unrestricted Giry spaces

Equip  $\llbracket GV \rrbracket$  with two gbs structures:

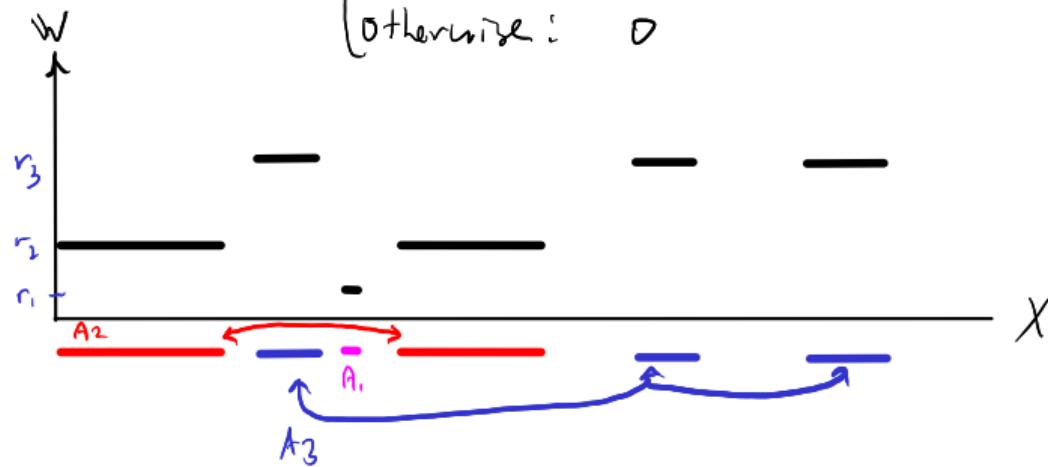
X  $R_{GV} := \{ \alpha: R \rightarrow GV \mid \forall A \in \mathcal{B}_V, \lambda r. \alpha(r, A): R \rightarrow W \}$

- ✓  $GV \hookrightarrow W^{B_X}$
- $\hookrightarrow \alpha$  is a kernel.
  - Fewer random elents
  - $R_{GV} \subseteq R_{G'V}$
  - Lebesgue integral measurable in both arguments.  
(upcoming)

Def: Simple function  $\varphi: X \rightarrow W$  when

$\exists n \in \mathbb{N}, \vec{A} \in \mathcal{B}_X^n, A_i \cap A_j = \emptyset, \vec{r} \in W$  s.t.

$$\varphi_x = \begin{cases} \vdots & \\ x \in A_i : & r_i \\ \vdots & \\ \text{otherwise:} & 0 \end{cases}$$



Encode into a space:

$$\text{SimpleCode} := \bigcup_{n \in \mathbb{N}} \mathcal{B}_X^n \times W^n$$

$$\text{Simple} := \{ f \in W^X \mid f \text{ simple} \} \hookrightarrow W^X$$

and define an interpretation:

$$[\![ - ]\!]: \text{SimpleCode} \longrightarrow \text{Simple}$$

$$[\![ (n, \vec{A}, \vec{r}) ]\!] := \sum_{i=1}^n r_i \cdot [\!- \in A_i]\!]$$

↳ characteristic function  
for  $A_i$

Lemma:  $f: X \rightarrow W$  is measurable → remember!  
qbs  
morphism!

iff  $f = \lim_{n \rightarrow \infty} f_n$  for some monotone sequence

$f_n \in \text{Simple}$ .

Moreover, we have measurable such choice:

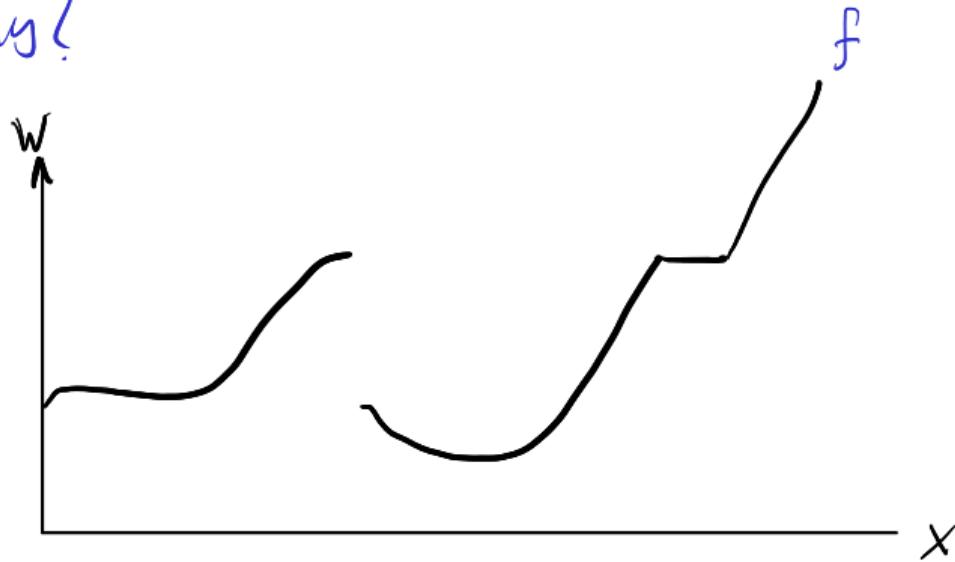
Simple Approx:

$$\left\{ \vec{\alpha} \in \mathbb{R}^+ \mid \Delta_n \rightarrow 0 \right\} \times \left\{ \vec{\alpha}' \in W^N \mid \begin{array}{l} \vec{\alpha}' \text{ monotone} \\ a_n \rightarrow \infty \end{array} \right\} \times W \xrightarrow{X} \text{SimpleCode}$$

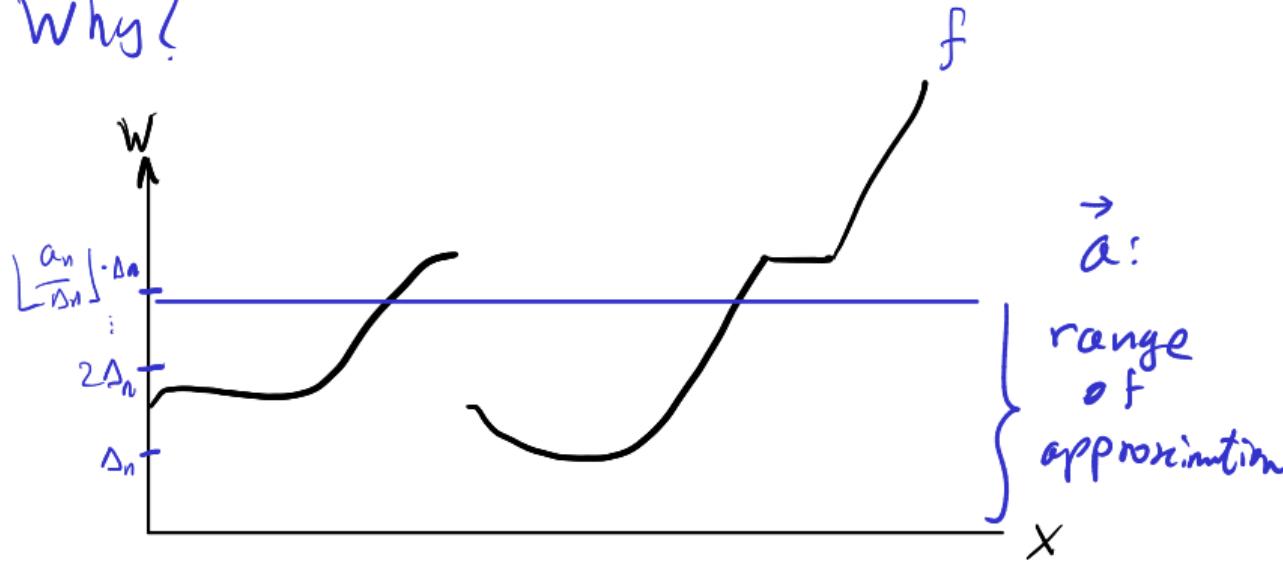
↑  
rate of convergence

↑  
Range of approximation

Why?

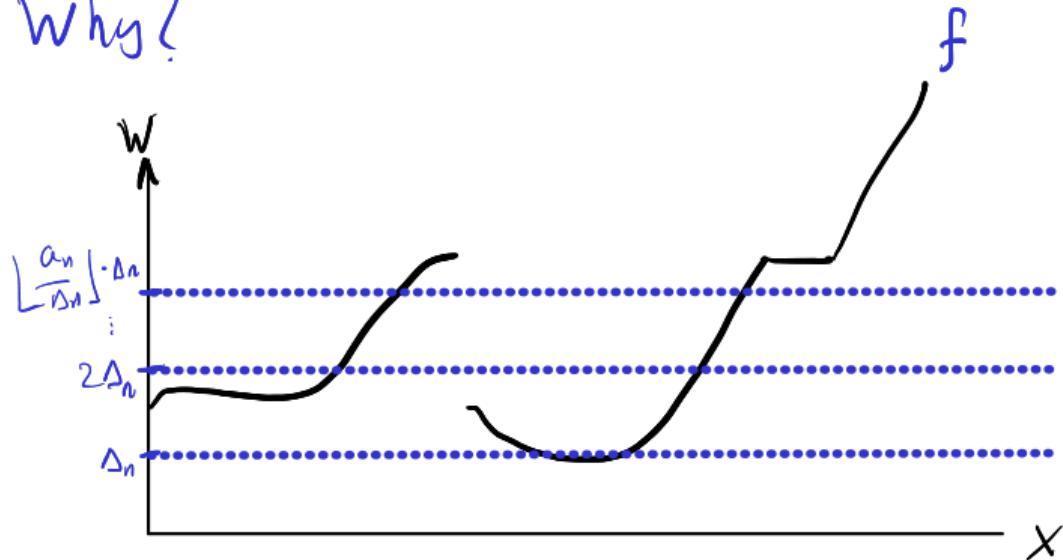


Why?

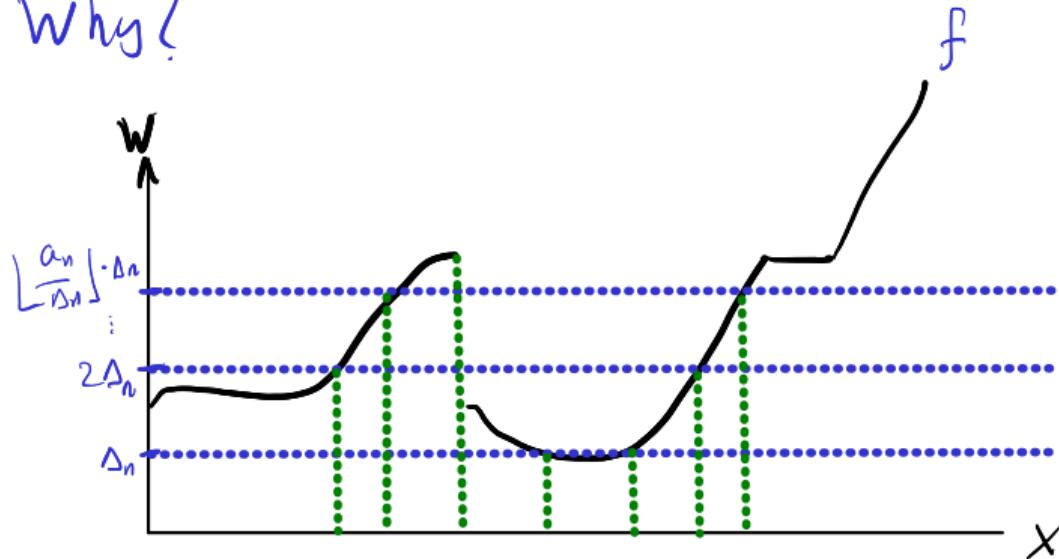


↗ resolution of  
approximation

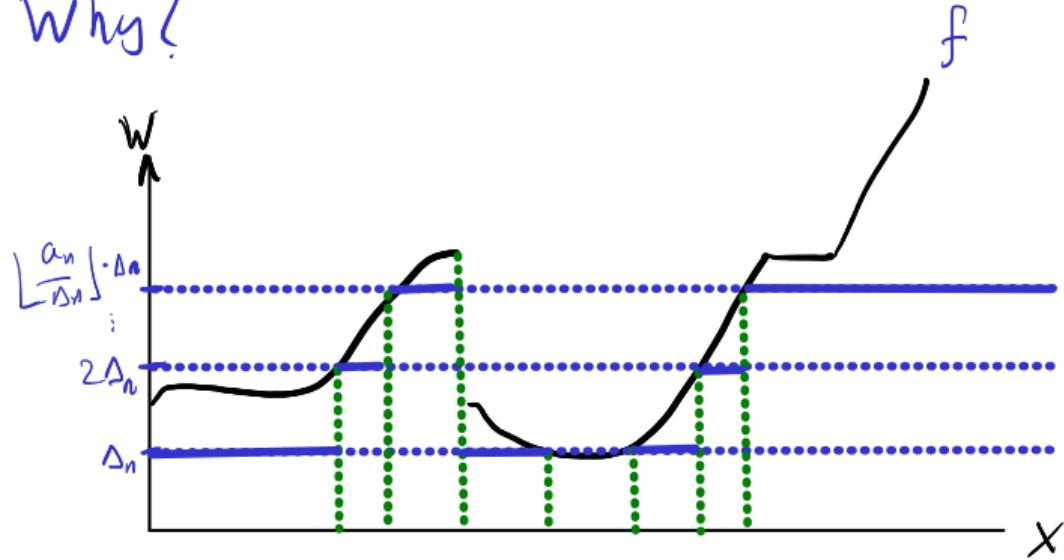
Why?



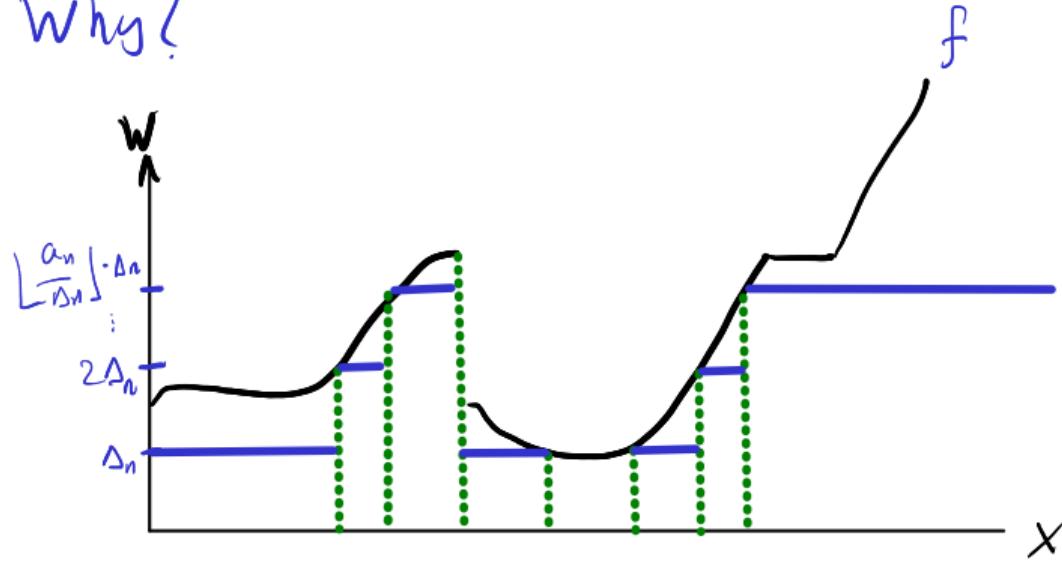
Why?



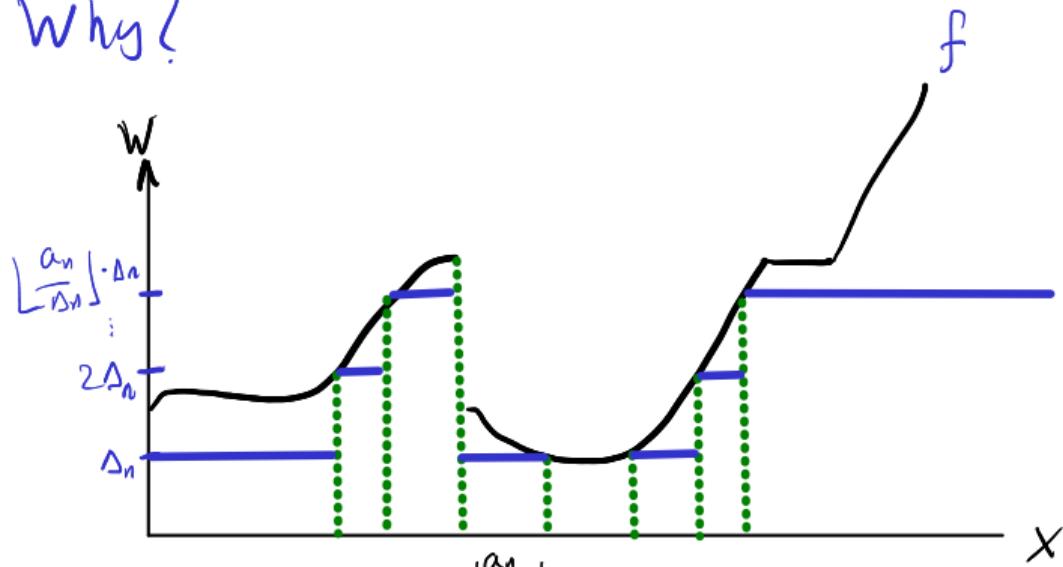
Why?



Why?

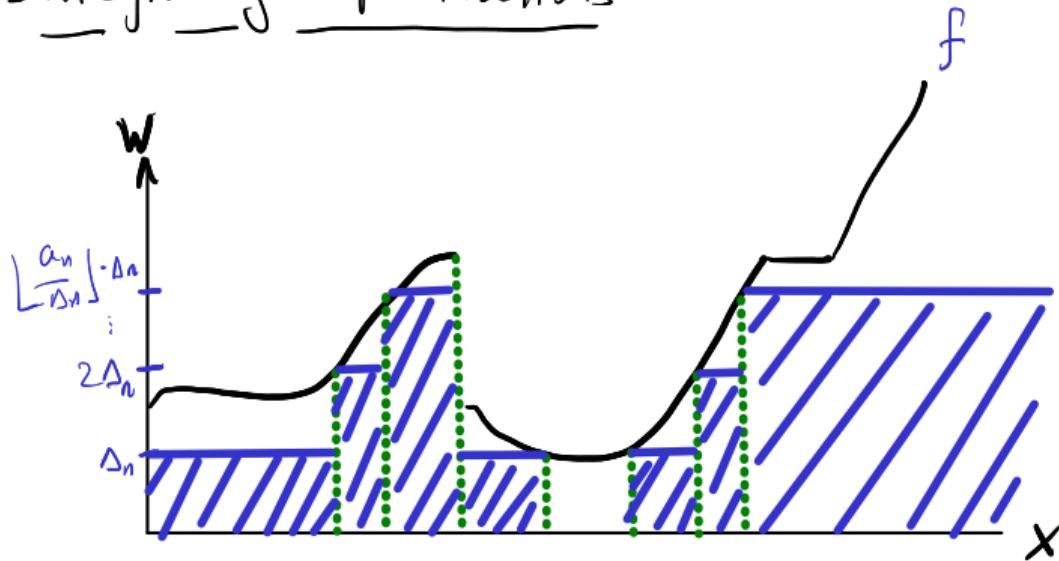


Why?



$$\| \text{Simple Approx} \overrightarrow{\Delta, \alpha} f \| := \sum_{i=1}^{\lfloor a_n / \Delta_n \rfloor} i \cdot \Delta_n [i \cdot \Delta_n \leq f < (i+1) \Delta_n] + \lfloor a_n / \Delta_n \rfloor \Delta_n \cdot [f \geq \lfloor a_n / \Delta_n \rfloor \cdot \Delta_n] \in \text{Simple}$$

## Integrating Simple Functions



$$\left\{ \begin{array}{l} : Gx \times \text{Simple Code} \rightarrow W \end{array} \right.$$

$$\int \mu(n, \vec{A}, \vec{r}) := \sum_{I \subseteq \{1, \dots, n\}} \left( \sum_{i \in I} r_i \right) \cdot \mu\left(\bigcap_{i \in I} A_i \setminus \bigcup_{i \notin I} A_i\right)$$

# Integration

proper higher-order  
operation

$$\int : Gx \times W^X \rightarrow W$$

$$\int \mu f := \sup \left\{ \int \mu \varphi \mid \varphi \in \text{Simple}, \varphi \leq f \right\}$$

we also  
write

$$= \lim_{n \rightarrow \infty} \int \mu (\text{Simple Approx}_{\Delta, \vec{a}} f)_n \sim \begin{matrix} \text{measurable} \\ \text{by} \\ \text{type} \end{matrix}$$

$$\int \mu(\delta x) t$$

$$\text{for } \int \mu(\lambda x, t)$$

$$\text{for } \frac{a_n}{\Delta_n} \rightarrow 0, \text{ e.g. } \Delta_n = \frac{1}{2^n}, a_n = n.$$

resolution

The unrestricted Giry      Strong Monad

Dirac:

$$\delta: X \rightarrow Gx$$

$$x \mapsto \lambda A. \begin{cases} x \in A : 1 \\ x \notin A : 0 \end{cases}$$

Unlike the unrestricted  
Giry on  $\text{Meas}$ .

but: non-commutative

Kleisli extension/ Kock integral:

$$\oint: Gx \times Gr^X \rightarrow Gr$$

$$\oint \mu f := \lambda A. \int \mu(dx) f(x; A)$$

(Fubini fails,  
just like in  
 $\text{Meas}$ )

Fubini - Tonelli; fails

$$\# \in G/R \quad \# E := \begin{cases} E \text{ finite:} & |E| \\ \text{o.w.:} & \infty \end{cases}$$

$\lambda \in G/R$  lebesgue  $k: \mathbb{R} \times \mathbb{R} \rightarrow W \cong G\mathbb{1}$

$$\int \#(dx) \underbrace{\int \lambda(dx) k(x,y)}_{y \in \mathbb{R} + \{\leftrightarrow\}} = \int \# \underline{0} = \underline{0} \stackrel{?}{=} 0$$
$$k(x,y) := [x=y]$$
$$y \in \mathbb{R} + \{\leftrightarrow\} \mapsto \lambda\{y\} \cdot 1 + \lambda\{y\} \cdot 0 = 0 \quad \text{if}$$

$$\int \lambda(dx) \underbrace{\#(dr)}_{x \in \mathbb{R} + \{\leftrightarrow\}} k(x,y) = \int \lambda(\{x\}) \delta_y \stackrel{?}{=} \infty$$
$$x \in \mathbb{R} + \{\leftrightarrow\} \mapsto \{x\} \cdot 1 + 0 = 1$$

## Randomisable measures monad

$$D \rightarrow G$$

$\lambda A. \int_{\text{Dom } \alpha} \lambda (\text{Dom } \alpha)$

$$L D X := \left\{ \lambda \alpha \mid \alpha : \mathbb{R} \rightarrow X \right\}$$

Lebesgue measure

$$R_{Dx} := \left\{ \lambda x. \lambda_{\alpha x} \mid \alpha : \mathbb{R} \times \mathbb{R} \rightarrow X \right\}$$

$$\delta : x \rightarrow Dx \quad \oint : D^{\Gamma \times (Dx)} \rightarrow Dx \quad \text{lift along } D \gg 6.$$

D validates our measure axioms including Fubini-Tonelli:

$$\mu \in DX, \nu \in DY \vdash$$

$$\oint \mu(dx) \oint \nu(dy) \delta_{(x,y)} = \oint \nu(dy) \oint \mu(dx) \delta_{(x,y)} =: \mu \otimes \nu$$

Thm: For  $sbs\ S$ ,  $PS, D_{\leq 1}S, D_{<\infty}S \in Sbs$

and agree with their counterparts on  $\text{Meas}$ .

$$DS_S = \{\mu \mid \mu \text{ s-finite}\} \quad \text{see [Staton'16]}$$

$$R_{DS} = \left\{ K: \mathbb{R} \rightarrow G(D) \mid K \text{ s-finite kernel} \right\}$$

Open: Is there a counterpart to  $D$  in  $\text{Meas}$ ?

More modestly, is  $DS \in Sbs$ ?

(Hypothesis: **No**)

## Distribution Submonads

A measure space

$$\Omega = (\Omega, \mu)$$

is a qbs  $\Omega$  with  
 $\mu \in D\chi$ .

Similarly:- finite measure space  
- (Sub) probability space.

$$P\chi := \left\{ \mu \in D\chi \mid \mu \chi = 1 \right\}$$

$$P_{\leq 1} \chi := \left\{ \mu \in D\chi \mid \mu \chi \leq 1 \right\}$$

$$P_{<\infty} \chi := \left\{ \mu \in D\chi \mid \mu \chi < \infty \right\}$$

$$D\chi$$

## Full model

$$\text{type: Obs} \quad W := [0, \infty] \quad \mathcal{B}^X \cong \mathcal{B}^X$$

$$DX := (\{\lambda_\alpha | \alpha: R \rightarrow X\}, \{\lambda_{r,\alpha} | \alpha: R \times R \rightarrow X\})$$

$$PX := \{\mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1\}$$

$$\underset{\mu}{\text{Ce}}[E] := \mu E \quad S_n := E \mapsto \begin{cases} n \in E : 1 \\ n \notin E : 0 \end{cases}$$

$$\oint \mu k := \lambda E. \int \mu(\lambda x) k(x; E)$$

Plan:

- 1) Type-driven probability: discrete case ✓
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- Lecture 1
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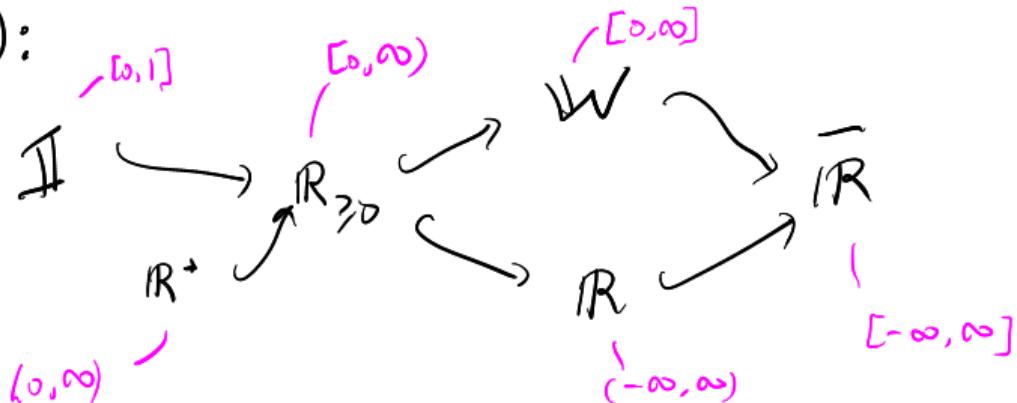


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Random variable:  $\xi : \Omega \rightarrow \mathbb{H} \hookrightarrow \bar{\mathbb{R}}$

$\mathbb{H}:$



-  $\Omega$  is a Space

-  $\mathbb{R}^n$  measurable vector space:

$$\alpha \xi + \zeta := \lambda w. \alpha \cdot \xi w + \zeta w$$

-  $W^n$  measurable  $\sigma$ -Semi-module  
for  $W$ :

$$\sum_{n=0}^{\infty} \alpha_n \xi_m :=$$

$$\lambda w. \sum_{n=0}^{\infty} \alpha_n \cdot \xi_{m,n}$$

$$\Pr_r : P_{\Omega} \times \mathcal{B}_n \rightarrow \mathbb{W}$$

$$\Pr_{\lambda} A := \text{eval}(\lambda, A) = \lambda A$$

Probability Space  $\mathcal{R} = (\Omega, \lambda_{\Omega})$

$P : P_{\Omega} \vdash$  "  $P_{\Omega}$  holds  $\lambda(\omega)$ -almost surely"

for some  $Q \subseteq \Omega$ ,  $P \models Q$ ,  $[-\infty Q] \cdot \lambda = \lambda$

Example  $(\xi, \zeta \in \mathbb{H}^{\Omega})$

$\xi = \zeta$  a.s. when  $\Pr_{\omega \sim \lambda} [\xi_{\omega} \neq \zeta_{\omega}] = 0$

Integrating Random Variables (as discretely) in Qbs!

$$(-)_+, (-)_- : \mathbb{R}^n \rightarrow \mathbb{W}^n$$

$$\xi_+ := \max(\xi, 0) \quad \xi_- := \max(-\xi, 0)$$

$$\text{So: } \xi = \xi_+ - \xi_-$$

$$\int : P\mathcal{R} \times \mathbb{W}^n \longrightarrow \mathbb{W} \quad \left. \begin{array}{l} \text{respects} \\ \text{a.s. equality:} \end{array} \right.$$

$$\int \lambda \xi := \int \lambda \xi_+ - \int \lambda \xi_- \quad \begin{aligned} \xi &= \zeta \text{ (a.s.)} \\ \Rightarrow \int \lambda \xi &= \int \lambda \zeta. \end{aligned}$$

Example

$$\lambda : P\Omega + \text{ASConv}(I\bar{\mathbb{R}})^{\Omega} : B(\bar{\mathbb{R}}^{N \times \Omega}) \\ := \left\{ \vec{\zeta} \in \bar{\mathbb{R}}^{N \times \Omega} \mid \Pr_{w \sim \lambda} [\lim \vec{\zeta}_n w \neq \perp] \right\}$$

So:

$$\lim^{\text{as}} : \bar{\mathbb{R}}^{N \times \Omega} \rightarrow \bar{\mathbb{R}}^\Omega \quad \text{Dom } \lim^{\text{as}} := \text{ASConv}(I\bar{\mathbb{R}})^\Omega$$

$$\lim^{\text{as}} \vec{\zeta} := \lambda_w. \limsup_{n \rightarrow \infty} f_n w$$

L  $\lim^{\text{as}}$  respects a.s. equality.

Thm (monotone convergence):

let  $\vec{\xi} \in \mathbb{W}^{N \times 2}$   $\lambda$ -a.s. monotone.

$$\vec{\xi}_l = \lim_{n \rightarrow \infty} \vec{\xi}_n \quad (\text{a.s.})$$



$$\int \lambda \vec{\xi}_l = \lim_{n \rightarrow \infty} \int \lambda \vec{\xi}_n$$

Lebesgue Space  $\left( \Omega \text{ Prob. Space}, P \in [1, \infty) \right)$

$P: [1, \infty), \lambda: P\Omega \vdash L_{(\Omega, \lambda)}^P: \mathcal{B}(\mathbb{R}^\Omega)$

$$:= \left\{ \xi \in \mathbb{R}^\Omega \mid \int |\xi|^P < \infty \right\} \hookrightarrow \mathbb{R}^\Omega$$

Ensemble  $L_\Omega := \prod_{\lambda \in P\Omega} L_{(\Omega, \lambda)}^P$   
 $P \in [1, \infty)$

$$L \quad P \leq q \Rightarrow L_\Omega^P \supseteq L_\Omega^q$$

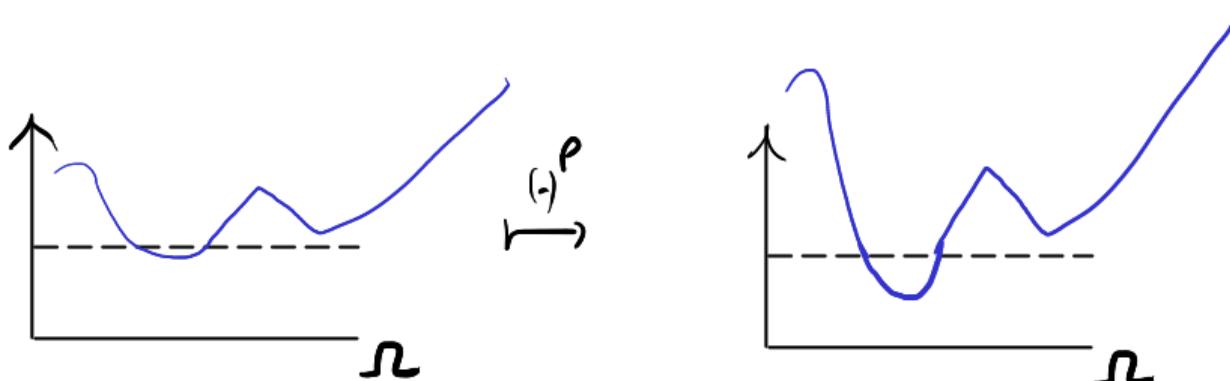
$L^p$  semi norms

$$\| \cdot \| : \bigcup_{p,\lambda} L_{(2,\lambda)}^p \rightarrow \mathbb{R}_{\geq 0} \quad \| \xi \|_p := \sqrt[p]{\int \lambda |\xi|^p}$$

$L^2$  inner product

$$\langle \cdot, \cdot \rangle : \bigcup_{p,\lambda} L_{(2,\lambda)}^p \times L_{(2,\lambda)}^p \rightarrow \mathbb{R}$$

$$\langle \xi, \eta \rangle_{p,\lambda} := \int \lambda \xi \eta$$



## Statistics

Expectation

$$\mathbb{E}: \bigcup_{\lambda} \mathcal{L}^1 \rightarrow \mathbb{R}$$

$$\mathbb{E}_{\lambda} \xi := \int_{\lambda} \xi$$

Covariance and Correlation

$$\text{Cov}, \text{Corr}: \bigcup_{\lambda} \mathcal{L}^2 \rightarrow \mathbb{R}$$

$$\text{Cov}(\xi, \zeta) := \langle \xi - \mathbb{E}\xi, \zeta - \mathbb{E}\zeta \rangle$$

$$\text{Corr}(\xi, \zeta) := \frac{\langle \xi, \zeta \rangle}{\|\xi\|_2 \|\zeta\|_2} = \cos(\text{angle}(\xi, \zeta))$$

## Sequential limits

$$\rho: [1, \infty), \lambda: P X \rightarrow \text{Cauchy } L_{\mathbb{Q}}^{\rho} : \mathcal{B}\left(\int_{P(X)}^{\rho}\right)^N$$
$$:= \left\{ \vec{\Sigma} \mid \forall \varepsilon \in \mathbb{Q}^+ \exists n \in \mathbb{N} \quad \forall m, n \geq n \quad \|\Sigma_{n+m} - \Sigma_n\|_{\rho} < \varepsilon \right\}$$

Thm:  $L_{\mathbb{Q}}^{\rho}$  is Cauchy-complete

$\lim$ : Cauchy  $L_{\mathbb{Q}}^{\rho} \rightarrow L^{\rho}$  (convergence in mean)

Why?

1. Every Cauchy sequence has an a.s. converging subseq.
2. We can find it measurable

## Example

Theorem (dominated convergence)

For  $\vec{z}_n, \vec{z} \in \mathbb{F}^1$  s.t.  $\vec{z}_n \leq \vec{z}$  a.s.:

$$1. \lim_{n \rightarrow \infty} \vec{z}_n \in \mathbb{F}^1$$

$$2. \lim_{n \rightarrow \infty} \vec{z}_n = \lim_{n \rightarrow \infty} \vec{z}$$

$$3. \lim_{n \rightarrow \infty} \int \lambda \vec{z}_n = \int \lambda \lim_{n \rightarrow \infty} \vec{z}_n$$

## Separability

Def:  $L^P$  separable: has countable dense subset

Fact: Separability is property of  $\lambda_2$ :

TFAE:

- $\exists p \geq 1$ .  $L^p$  separable
- $\forall p \geq 1$ .  $L^p$  separable

Measurable separability in  $I \hookrightarrow P\Omega \times [1, \infty)$

$$\vec{\beta} : \prod_{(\lambda, p) \in I} L^p_{(\Omega, \lambda)}^{(N)} \quad \text{s.t.}$$

$$\left\{ \vec{\beta}_n^{(p)} \mid n \in \mathbb{N} \right\} \text{ dense in } L^p_{(\Omega, \lambda)}$$

Prop. - Every SBS  $S$  measurable separable in

$$PS \times [1, \infty)$$

-  $I \hookrightarrow P\Omega \times \{2\}$  measurable separable

$$\langle \beta_n, \beta_m \rangle = 0$$

$$\|\beta_n\|_2 = 1$$

$\Rightarrow \exists \vec{\beta} \in \prod_{\lambda \in I} L^2_{(\Omega, \lambda)}$  orthonormal system  $(\beta_n)$  dense

## Escape

Let  $S \hookrightarrow L^2$  closed Vector Subspace.

Orthogonal decomposition linear in fact.

$$\langle P, P^\perp \rangle : L^2 \xrightarrow{P} S \times S^\perp$$

When  $S$  is separable with orthonormal system  $\beta$

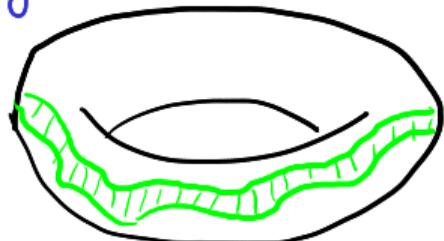
We have a measurable version of

$$\langle P, P^\perp \rangle : L^2 \rightarrow S \times S^\perp$$

$$P\xi := \sum_{n=0}^{\infty} \langle \xi, \beta_n \rangle \beta_n \quad P^\perp := \text{Id} - P.$$

# Kolmogorov's Conditional Expectation

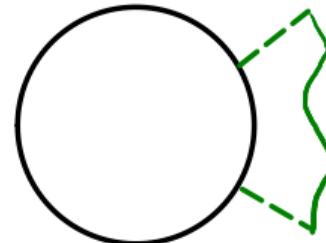
ground truth space



(H)

Sample space

H  
observation



$\Sigma$   
Statistic  
of interest

Conditional expectation

$$\mathbb{E}[\Sigma | H = -]$$

Observed  
statistic

R

# Kolmogorov's Conditional Expectation

A Conditional expectation

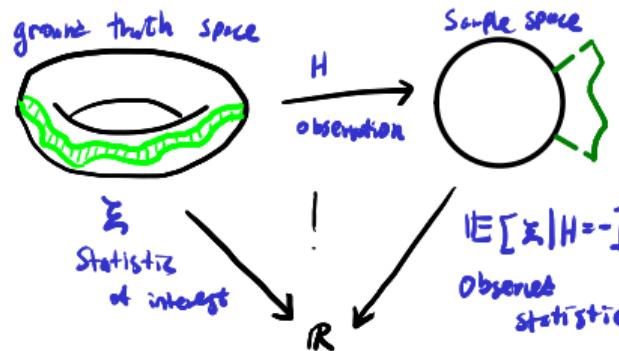
of  $Z \in \mathcal{L}_n^1$  wrt

$H: \Omega \rightarrow \mathbb{H}$  is

$Z \in \mathcal{L}_{\mathbb{H}}^1$  s.t. for all  $A \in \mathcal{B}_{\mathbb{H}}$ :

$$\int_A \mu Z = \int_{H^{-1}[A]} \lambda Z$$

where  $\mu := \lambda_H$

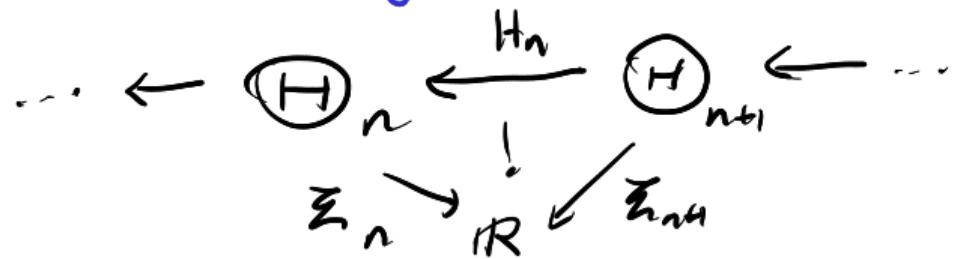


# Conditional expectations

1. unique a.s.

2. fundamental to modern probability, e.g.:

a Martingale



$$\text{St. } \bar{Z}_n = \mathbb{E}[\bar{Z}_{n+1} | H_n = -]$$

Theorem (Existence)

- $\exists \mathbb{E}[-|H=-]: \mathcal{L}_{(\Omega, \lambda)} \rightarrow \mathcal{F}_{(\Theta, \mu)}^I$

- When  $(\Omega, \lambda)$  is Separable

$$\mathbb{E}[-|H=-]: \mathcal{L}_{(\Omega, \lambda)} \rightarrow \mathcal{F}_{(\Theta, \mu)}^I$$

- When  $\Theta$  is  $I^I$ -measurably separable

$$\mathbb{E}[-|H=-]: \prod_{\substack{H \in \Theta \\ \lambda \in H_I^{-1}[I]}} \mathcal{L}_{(\Omega, \lambda)} \rightarrow \mathcal{F}_{(\Theta, \mu)}^I$$

Plan:

- 1) Type-driven probability: discrete case ✓
- 2) Borel sets & measurable spaces ✓
- 3) Quasi Borel spaces ✓
- 4) Type structure & standard Borel spaces ✓
- 5) Integration & random variables ✓

Lecture 1

Lecture 2

Please ask questions!

smile



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## Discrete model

type : set       $\mathbb{W} := [0, \infty]$        $\mathcal{B}X := \mathcal{P}X$

$\mathcal{D}X := \{\mu : X \rightarrow \mathbb{W} \mid \text{Supp } \mu \text{ countable}\}$

$\mathcal{P}X := \{\mu \in \mathcal{D}X \mid \underset{x}{\text{Ce}}[x] = 1\}$

$\underset{\mu}{\text{Ce}}[E] := \sum_{x \in E} \mu_x \quad \delta_x := \lambda x'. \begin{cases} x = x' : 0 \\ x \neq x' : 1 \end{cases}$

$\oint \mu k := \lambda x. \sum_{m \in \Gamma} \mu^m \cdot k(m; x)$

## Full model

$$\text{type: Obs} \quad W := [0, \infty] \quad \mathcal{B}^X \cong \mathcal{B}^X$$

$$DX := (\{\lambda_\alpha | \alpha: R \rightarrow X\}, \{\lambda_{r,\alpha} | \alpha: R \times R \rightarrow X\})$$

$$PX := \{\mu \in DX \mid \underset{\mu}{\text{Ce}}[X] = 1\}$$

$$\underset{\mu}{\text{Ce}}[E] := \mu E \quad S_n := E \mapsto \begin{cases} n \in E : 1 \\ n \notin E : 0 \end{cases}$$

$$\oint \mu k := \lambda E. \int \mu(\lambda x) k(x; E)$$

Enough!

Lunch.