A domain theory for quasi-Borel spaces

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Probabilistic programming

$$\llbracket - \rrbracket : programs \rightarrow distributions$$

▶ Continuous types: \mathbb{R} , $[0, \infty]$ scale distribution by rProbabilistic effects: $r:[0,\infty]$ normally $\mathbf{sample}: \mathbb{R}$ distributed $\mathbf{score}(r): \mathbf{1}$ sample evaluation order doesn't matter: s-finite distributions [Staton'17] Commutativity/Fubini quasi-Borel spaces [Heunen et al.'17] Traditional programming: modular implementation of Higher-order functions Bayesian inference algorithms [Ścibior et al.'18] Inductive types and bounded iteration domain theory Type and term recursion [this work]

Iso-recursive types: FPC

$$\Delta = \{\alpha_1, \dots, \alpha_n\}$$

$$\frac{\Delta,\alpha \vdash_{\mathbf{k}} \tau : \mathsf{type}}{\Delta \vdash_{\mathbf{k}} \mu \alpha. \tau : \mathsf{type}}$$

$$\frac{\Gamma \vdash t : \sigma[\alpha \mapsto \tau]}{\Gamma \vdash \tau.\mathbf{roll}\,(t) : \tau} (\tau = \mu \alpha.\sigma)$$

type recursion

$$\Gamma \vdash t : \mu \alpha . \sigma$$

$$\Gamma \vdash t : \mu \alpha. \sigma$$
 $\Gamma, x : \sigma[\alpha \mapsto \mu \alpha. \sigma] \vdash s : \tau$

 $\Gamma \vdash \mathbf{match} \ t \ \mathbf{with} \ \mathbf{roll} \ x \Rightarrow s : \tau$

 ω Cpo-enriched category

$$\llbracket \Delta \vdash_{\mathrm{k}} \tau : \mathsf{type}
rbracket : (\stackrel{\longleftarrow}{\mathcal{C}}^{\mathrm{op}})^n \times \mathcal{C}^n \xrightarrow{\longrightarrow} \stackrel{\longleftarrow}{\mathcal{C}}$$

Recursive types denote minimal invariants $[\![\Delta \vdash_{k} \mu \alpha. \tau : \mathsf{type}]\!]$

[Pitts'96]

locally continuous

functor

Challenge

- probabilistic powerdomain
- ightharpoonup commutativity/Fubini \leftarrow
- ▶ domain theory
- ▶ higher-order functions

traditional approach:

 $\mathsf{domain} \mapsto \mathsf{open} \ \mathsf{subsets} \mapsto \mathsf{Borel} \ \mathsf{subsets} \mapsto \mathsf{distributions}$

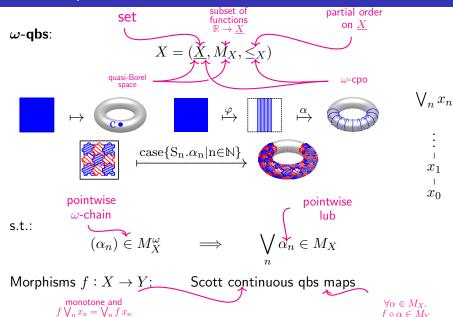
continuous domains [Jones-Plotkin'89]

our approach:

 $(domain, quasi-Borel \ space) \mapsto distributions$ separatebut compatible

open problem [Jung-Tix'98]

quasi-Borel pre-domains



Characterising $\omega \mathbf{Qbs}$

canonical map countable $F\mathbb{R} \to (F1)^{\mathbb{R}}$ product preserving injective $\mathbf{Sbs} \mapsto [\mathbf{Sbs^{op}}, \ \mathbf{Set} \]_{\mathbf{cpp}} \mapsto \ \mathsf{SepSh} \leftarrow$ qbs: Thm: $\mathbf{Qbs} \simeq \mathsf{SepSh}$ $\mathbf{Sbs} \mapsto [\mathbf{Sbs}^{\mathrm{op}}, \omega \mathbf{Cpo}]_{\mathrm{cpp}} \mapsto \omega \mathsf{SepSh}$ ω -qbs: $\omega \mathbf{Qbs} \simeq \omega \mathsf{SepSh}$ canonical map Thm: $F\mathbb{R} \to (F1)^{\mathbb{R}}$ full mono (order reflecting) \triangleright ω -cpos internal to Qbs as a quasi-topos strong monos: algebras for an essentially algebraic theory: $X \stackrel{f}{\rightarrowtail} Y$ $(f \circ)^{-1}[M_V] = M_X$ $\omega cpo + qbs + compatibility axiom$

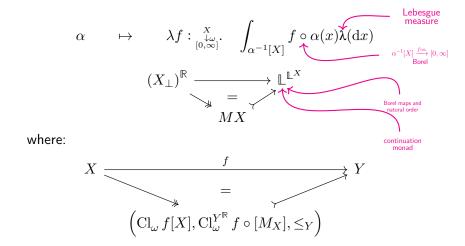
Axiomatic domain theory

Borel-open map
$$m: X \xrightarrow{\text{strong mono}} Y$$
:
$$\forall \beta \in M_Y. \qquad \beta^{-1}[m[X]] \in \mathcal{B}(\mathbb{R})$$

Borel-Scott open maps form a domain structure on $\omega \mathbf{Qbs}$

- → model axiomatic domain theory
- ⇒ solve recursive domain equations

A probabilistic powerdomain



Fact: densely strong epis closed under: product, lifting, $(-)^{\mathbb{R}}$ \implies monad-structure factorisation [McDermott-Kammar'18]

Summary

- $ightharpoonup \omega \mathbf{Qbs}$: separate, compatible domain and qbs structures
- Cartesian-closed, locally presentable
- Borel-Scott open maps model axiomatic domain theory
- Commutative probabilistic powerdomain models synthetic measure theory
- Adequate semantics for Probabilistic FPC